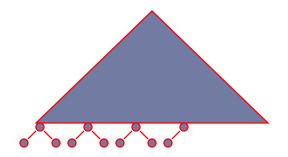
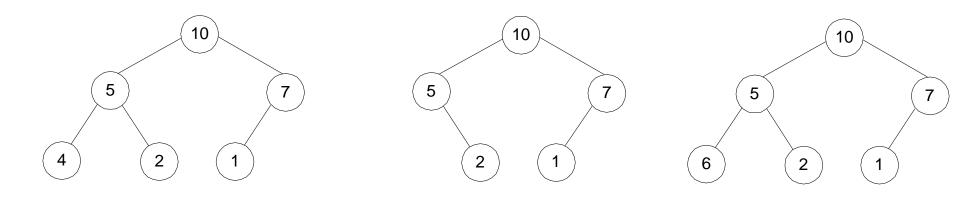
# Transform and Conquer

# Heaps and Heapsort

- A *heap* is a binary tree with a key at each nodes such that it is essentially complete
- A *complete* tree has all its levels are full except possibly the last level, where only some rightmost keys may be missing
- ▶ The key at each node is  $\geq$  keys at its children

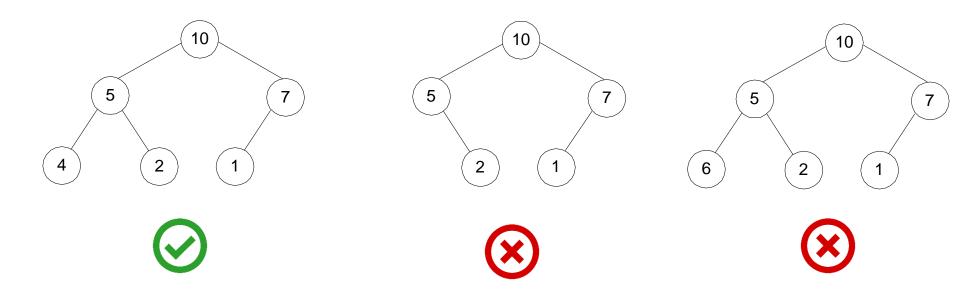


## Illustration of the heap's definition





### Illustration of the heap's definition



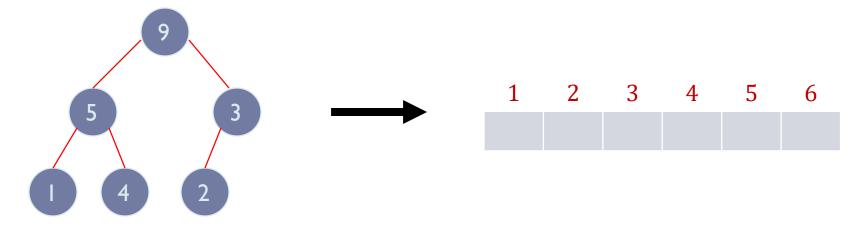
- Heap's elements are ordered top down (along any path down from its root)
- But they are not ordered left to right

### Some Important Properties of a Heap

- Given n, there exists a unique binary tree with n nodes that is essentially complete, with  $h = \lfloor \log_2 n \rfloor$
- ▶ The root contains the largest key
- The sub-tree rooted at any node of a heap is also a heap
- A heap can be represented as an array

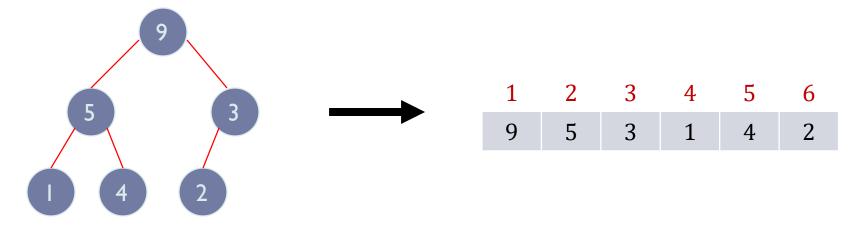
## Heap's Array Representation

- Store heap's elements in an array
- $\triangleright$  Elements indexed, for convenience, 1 to n
- In top-down left-to-right order



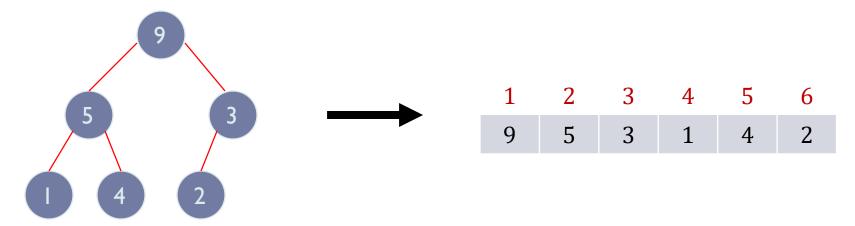
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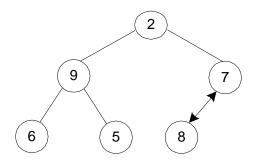
- Left child of node j is at 2j
- Right child of node j is at 2j + 1
- Parent of node j is at  $\lfloor \frac{j}{2} \rfloor$
- Parental nodes are represented in the first  $\lfloor \frac{n}{2} \rfloor$  locations

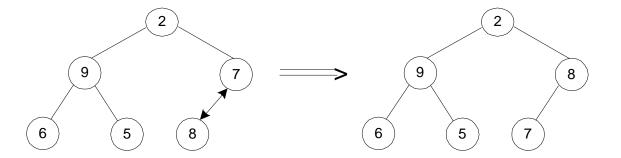
## Heap Construction (bottom-up)

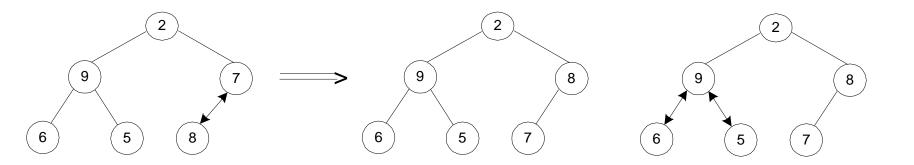
Step 0: Initialize the structure with keys in the order given

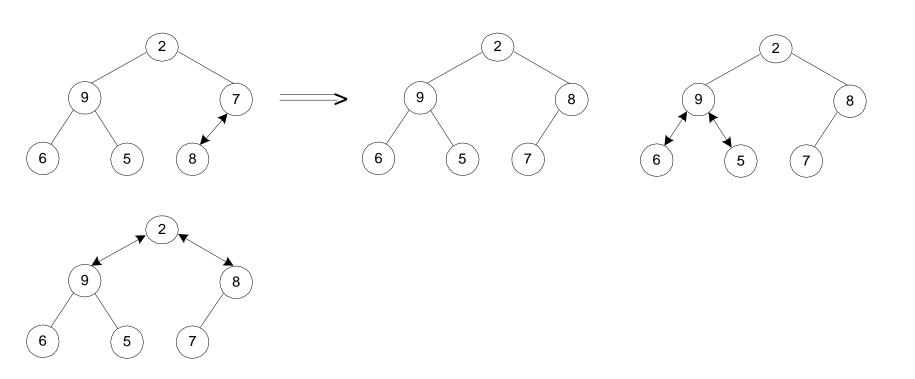
Step 1: Starting with the last (rightmost) parental node, fix the heap rooted at it, if it doesn't satisfy the heap condition: keep exchanging it with its largest child until the heap condition holds

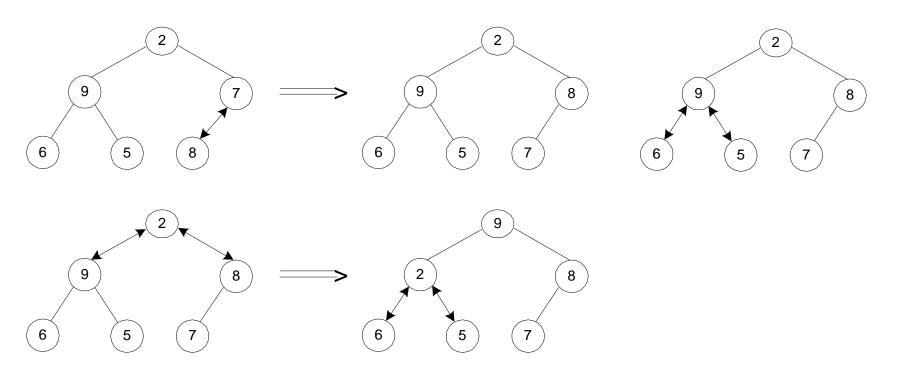
Step 2: Repeat Step 1 for the preceding parental node

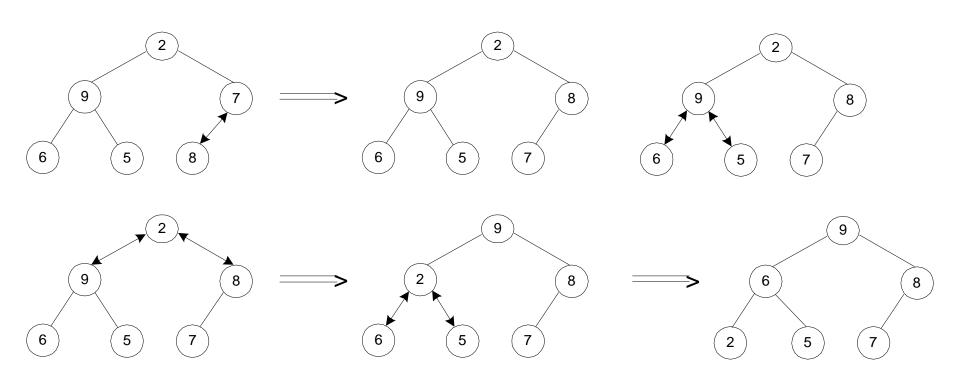












## Bottom-up heap construction

```
Algorithm HeapBottomUp(H[1..n])
//Constructs a heap from the elements of a given array
// by the bottom-up algorithm
//Input: An array H[1..n] of orderable items
//Output: A heap H[1..n]
for i \leftarrow \lfloor n/2 \rfloor downto 1 do
    k \leftarrow i; \quad v \leftarrow H[k]
    heap \leftarrow \mathbf{false}
    while not heap and 2*k \le n do
           i \leftarrow 2 * k
            if j < n //there are two children
               if H[j] < H[j+1] j \leftarrow j+1
            if v \geq H[j]
                  heap \leftarrow true
            else H[k] \leftarrow H[j]; \quad k \leftarrow j
     H[k] \leftarrow v
```

### Heapsort

Stage 1: Construct a heap for a given list of *n* keys

Stage 2: Repeat operation of root removal n-1 times:

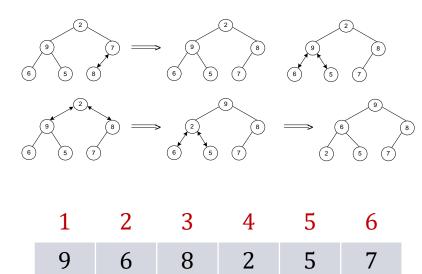
- Exchange keys in the root and in the last (rightmost) leaf
- Decrease heap size by 1
- If necessary, swap new root with larger child until the heap condition holds

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)

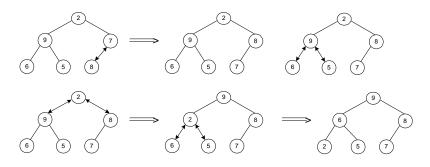
Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)



Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)



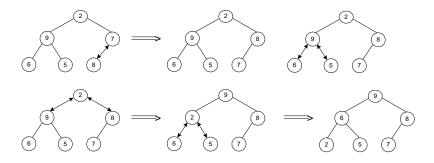
 1
 2
 3
 4
 5
 6

 9
 6
 8
 2
 5
 7

1	2	3	4	5	6
9	6	8	2	5	7

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)

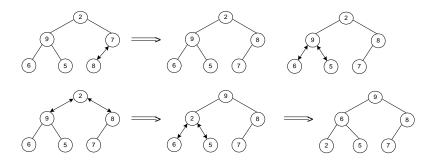


1	2	3	4	5	6
9	6	8	2	5	7

1	2	3	4	5	6
9	6	8	2	5	7
7	6	8	2	5	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

#### Stage 1 (heap construction)

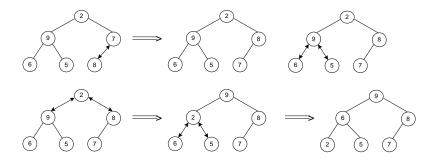


1	2	3	4	5	6
9	6	8	2	5	7

				5	
					7
7	6	8	2	5	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

#### Stage 1 (heap construction)

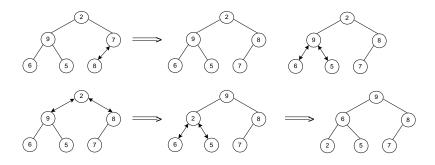


1	2	3	4	5	6
9	6	8	2	5	7

1	2	3	4	5	6
9	6	8	2	5	7
7	6	8	2	5	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

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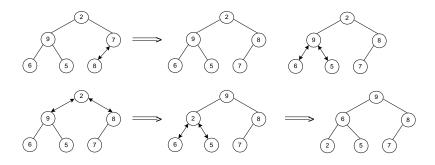


1	2	3	4	5	6
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		3			
9	6	8	2	5	7
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Stage 1 (heap construction)

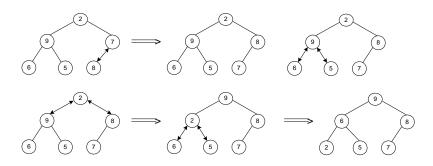


1	2	3	4	5	6
9	6	8	2	5	7

1	2	3	4	5	6
		8			
8	6	7	2	5	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)

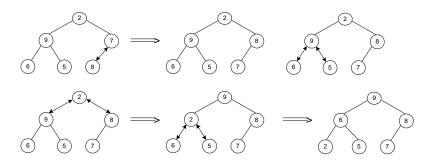


1	2	3	4	5	6
9	6	8	2	5	7

		3			
9	6	8	2	5	7
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Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)

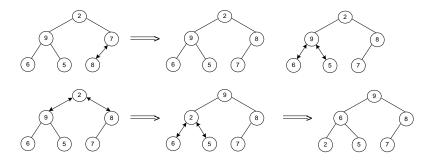


1	2	3	4	5	6
9	6	8	2	5	7

	2				
9	6	8	2	5	7
8	6	7	2	5	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

#### Stage 1 (heap construction)

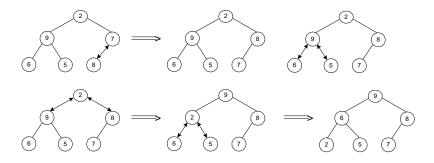


1	2	3	4	5	6
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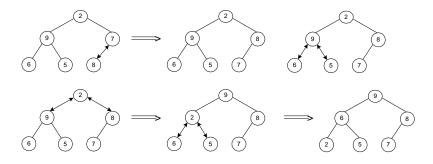


1	2	3	4	5	6
9	6	8	2	5	7

1	2	3	4	5	6	
				5		
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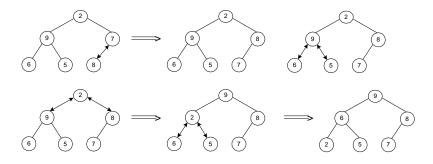


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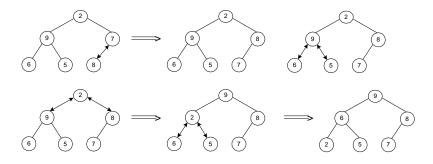


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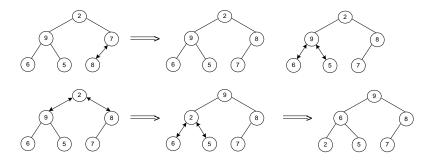


1	2	3	4	5	6
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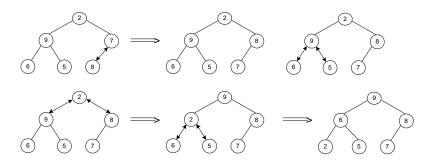


1	2	3	4	5	6
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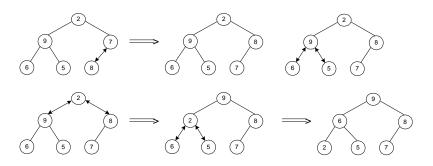


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	2				
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Stage 1 (heap construction)

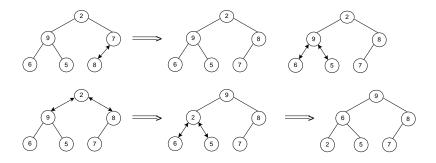


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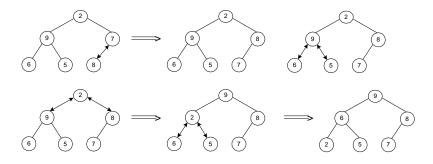


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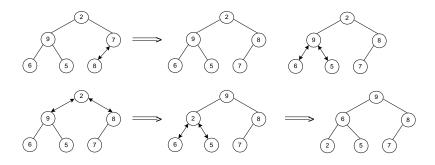


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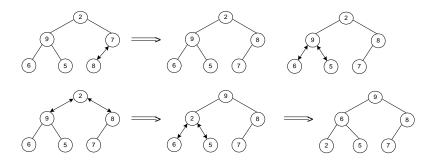


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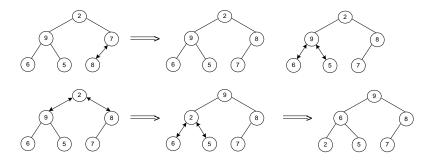


1	2	3	4	5	6
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1	2	3	4	5	6
				5	
6	2	5	7	8	9

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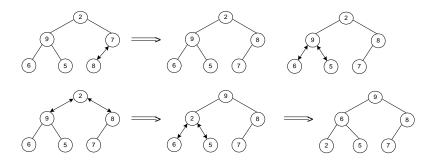


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9	6	8	2	5	7

1	2	3	4	5	6
	6				
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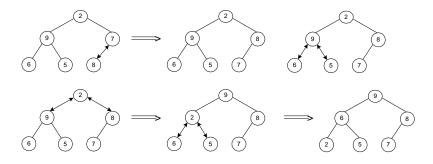


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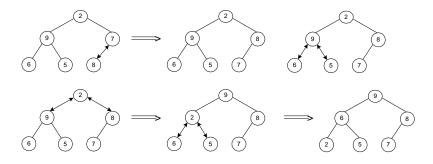


1	2	3	4	5	6
9	6	8	2	5	7

				5	
					7
5	2	6	7	8	9

Sort the list 2, 9, 7, 6, 5, 8 by heapsort

Stage 1 (heap construction)

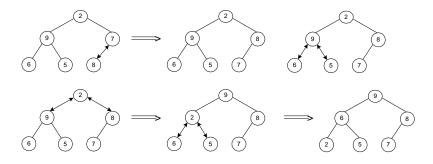


1	2	3	4	5	6
9	6	8	2	5	7

	2				
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Sort the list 2, 9, 7, 6, 5, 8 by heapsort

#### Stage 1 (heap construction)

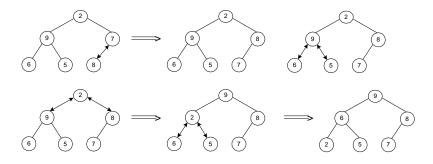


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	2				
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Sort the list 2, 9, 7, 6, 5, 8 by heapsort

#### Stage 1 (heap construction)



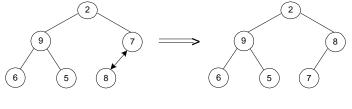
1	2	3	4	5	6
9	6	8	2	5	7

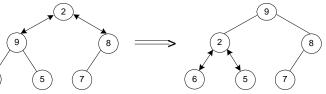
	2				
9	6	8	2	5	7
2	5	6	7	8	9

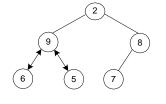
Sort the list 2, 9, 7, 6, 5, 8 by heapsort

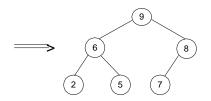
#### Stage I (heap construction)

- 2 9 7 6 5 8
- 2 9 8 6 5 7
- 2 9 8 6 5 7
- 9 2 8 6 5 7
- 9 6 8 2 5 7









- 9 6 8 2 5 7
- 7 6 8 2 5 | 9
- 8 6 7 2 5 | 9
- 5 6 7 2 | 8 9
- 7 6 5 2 8 9
- 2 6 5 | 7 8 9
- 6 2 5 | 7 8 9
- 5 2 | 6 7 8 9
- <u>5</u> 2 | 6 7 8 9
- 2 | 5 6 7 8 9

# Analysis of Heapsort

Stage 1: Build heap for a given list of *n* keys worst-case

$$C(n) = \sum_{i=0}^{h-1} 2^{i} \times 2(h-i) = 2(n - \log(n+1)) \in \Theta(n)$$
# nodes at level i

Stage 2: Repeat operation of root removal *n*-1 times (fix heap) worst-case

$$C(n) = \sum_{i=1}^{n-1} 2\log i \in \Theta(n\log n)$$

Both worst-case and average-case efficiency:  $\theta(n \log n)$ 

In-place: yes

# Priority Queue

- A *priority queue* is a set of elements with numerical priorities:
  - find element with highest priority
  - delete element with highest priority
  - insert element with assigned priority (see below)
- Heap is a very efficient way for implementing priority queues
- Two ways to handle priority queue in which highest priority = smallest number highest priority = largest number

### Insertion of a New Element into a Heap

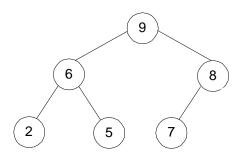
- Insert the new element at last position in heap.
- Compare it with its parent and, if it violates heap condition, exchange them
- Continue comparing the new element with nodes up the tree until the heap condition is satisfied



### Insertion of a New Element into a Heap

- Insert the new element at last position in heap.
- Compare it with its parent and, if it violates heap condition, exchange them
- Continue comparing the new element with nodes up the tree until the heap condition is satisfied

Example: Insert key 10

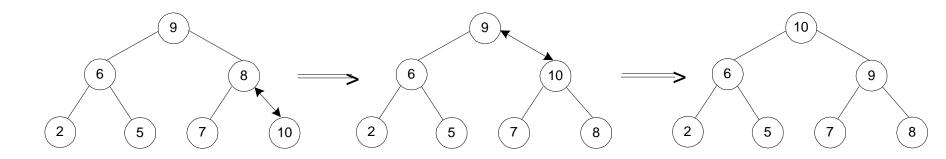




### Insertion of a New Element into a Heap

- Insert the new element at last position in heap.
- Compare it with its parent and, if it violates heap condition, exchange them
- Continue comparing the new element with nodes up the tree until the heap condition is satisfied

Example: Insert key 10



Efficiency:  $O(\log n)$ 



- Sort the following lists by heapsort by using the array representation of heaps.
  - a. 1, 2, 3, 4, 5 (in increasing order)
  - b. 5, 4, 3, 2, 1 (in increasing order)
  - c. S, O, R, T, I, N, G (in alphabetical order)

6. a. Sort 1, 2, 3, 4, 5 by heapsort

Heap Construction							
1	<b>2</b>	3	4	5			
1	5	3	4	2			
1	5	3	4	2			
5	4	3	1	2			

b. Sort 5, 4, 3, 2, 1 (in increasing order) by heapsort

Heap Construction 5 4 3 2 1 5 4 3 2 1

c. Sort S, O, R, T, I, N, G (in alphabetic order) by heapsort Heap Construction Maximum Deletions SS R O I Ο  $\mathbf{R}$   $\mathbf{T}$ O R T I N G R O I N G O R G I S TROING Т O N G  $\mathbf{R}$ G R Ν G Ν O

Spaghetti sort Imagine a handful of uncooked spaghetti, individual rods whose lengths represent numbers that need to be sorted.

- a. Outline a "spaghetti sort"—a sorting algorithm that takes advantage of this unorthodox representation.
- b. What does this example of computer science folklore (see [Dew93]) have to do with the topic of this chapter in general and heapsort in particular?

- 11. a. After the bunch of spaghetti rods is put in a vertical position on a tabletop, repeatedly take the tallest rod among the remaining ones out until no more rods are left. This will sort the rods in decreasing order of their lengths.
  - b. The method shares with heapsort its principal idea: represent the items to be sorted in a way that makes finding and deleting the largest item a simple task. From a more general perspective, the spaghetti sort is an example, albeit a rather exotic one, of a representation-change algorithm.

# Programming Exercise

- Implement the heapsort algorithm
- Develop test cases to test your code
- Compare Heapsort with other sorting algorithms