

A dark blue vertical bar is positioned on the left side of the slide, spanning the height of the first rectangular box.

Transform and Conquer

A light blue vertical bar is positioned on the left side of the slide, spanning the height of the second rectangular box.

Transform and Conquer

This group of techniques solves a problem by a *transformation*

- ▶ to a simpler/more convenient instance of the same problem (*instance simplification*)
- ▶ to a different representation of the same instance (*representation change*)
- ▶ to a different problem for which an algorithm is already available (*problem reduction*)

Instance simplification - Presorting

Solve a problem's instance by transforming it into another simpler/easier instance of the same problem

Presorting

- ▶ Many problems involving lists are easier when the list is sorted.
 - ▶ searching
 - ▶ computing the median (selection problem)
 - ▶ checking if all elements are distinct (element uniqueness)

Also:

- ▶ Topological sorting helps solving some problems for dags.
- ▶ Presorting is used in many geometric algorithms.

How fast can we sort ?

Efficiency of algorithms involving sorting depends on efficiency of sorting.

$\lceil \log_2 n! \rceil \approx n \log_2 n$ comparisons are necessary in the worst case to sort a list of size n by any comparison-based algorithm.

About $n \log_2 n$ comparisons are also sufficient to sort array of size n (by mergesort).

Searching with presorting

Problem: Search for a given K in $A[0..n-1]$

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Presorting-based algorithm:

Stage 1 Sort the array by an efficient sorting algorithm

Stage 2 Apply binary search

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- ▶ Good or bad?
- ▶ Why do we have our dictionaries, telephone directories, etc. sorted?

Element Uniqueness with Presorting

- ▶ **Brute force algorithm**
 - Compare all pairs of elements

Element Uniqueness with Presorting

- ▶ **Brute force algorithm**

Compare all pairs of elements

Efficiency: $O(n^2)$

Element Uniqueness with Presorting

- ▶ Presorting-based algorithm

Element Uniqueness with Presorting

- ▶ **Presorting-based algorithm**

- Stage 1: sort by efficient sorting algorithm (e.g. mergesort)

- Stage 2: scan array to check pairs of adjacent elements

Element Uniqueness with Presorting

► Presorting-based algorithm

Stage 1: sort by efficient sorting algorithm (e.g. mergesort)

Stage 2: scan array to check pairs of adjacent elements

Efficiency: $\Theta(n \log n) + O(n) = \Theta(n \log n)$

Gaussian Elimination

Given: A system of n linear equations in n unknowns with an arbitrary coefficient matrix.

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2$$

...

$$a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n$$

Gaussian Elimination

Given: A system of n linear equations in n unknowns with an arbitrary coefficient matrix.

Transform to: An equivalent system of n linear equations in n unknowns with an upper triangular coefficient matrix.

Solve the latter by substitutions starting with the last equation and moving up to the first one.

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2$$

$$a'_{11}x_1 + a'_{12}x_2 + \dots + a'_{1n}x_n = b'_1$$

$$a'_{22}x_2 + \dots + a'_{2n}x_n = b'_2$$



$$a_{n1}x_1 + a_{n2}x_2 + \dots + a_{nn}x_n = b_n$$

$$a'_{nn}x_n = b'_n$$

Gaussian Elimination (cont.)

The transformation is accomplished by a sequence of elementary operations on the system's coefficient matrix (which don't change the system's solution):

for $i \leftarrow 1$ to $n-1$ do
 replace each of the subsequent rows (i.e., rows $i+1, \dots, n$) by
 a difference between that row and an appropriate multiple
 of the i -th row to make the new coefficient in the i -th column
 of that row 0

Gaussian Elimination (cont.)

ALGORITHM *GaussElimination*($A[1..n, 1..n]$, $b[1..n]$)

//Applies Gaussian elimination to matrix A of a system's coefficients,

//augmented with vector b of the system's right-hand side values

//Input: Matrix $A[1..n, 1..n]$ and column-vector $b[1..n]$

//Output: An equivalent upper-triangular matrix in place of A with the

//corresponding right-hand side values in the $(n + 1)$ st column

for $i \leftarrow 1$ **to** n **do** $A[i, n + 1] \leftarrow b[i]$ //augments the matrix

for $i \leftarrow 1$ **to** $n - 1$ **do**

for $j \leftarrow i + 1$ **to** n **do**

for $k \leftarrow i$ **to** $n + 1$ **do**

$A[j, k] \leftarrow A[j, k] - A[i, k] * A[j, i] / A[i, i]$

Example of Gaussian Elimination

$$\begin{array}{rclcl} 2x_1 & - & 4x_2 & + & x_3 & = & 6 \\ 3x_1 & - & x_2 & + & x_3 & = & 11 \\ x_1 & + & x_2 & - & x_3 & = & -3 \end{array}$$

Example of Gaussian Elimination

$$\begin{aligned} 2x_1 - 4x_2 + x_3 &= 6 \\ 3x_1 - x_2 + x_3 &= 11 \\ x_1 + x_2 - x_3 &= -3 \end{aligned}$$

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right|$$

Example of Gaussian Elimination

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Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right|$$

Eliminate the first column
After the first row

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right| \begin{array}{l} \\ -\left(\frac{3}{2}\right) \times \text{row } 1 \\ -\left(\frac{1}{2}\right) \times \text{row } 1 \end{array}$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right| \quad \begin{array}{l} \\ -\left(\frac{3}{2}\right) \times [2 \quad -4 \quad 1 \quad 6] \\ -\left(\frac{1}{2}\right) \times [2 \quad -4 \quad 1 \quad 6] \end{array}$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right| \quad \begin{array}{l} \\ -[3 \quad -6 \quad \frac{3}{2} \quad 9] \\ -[1 \quad -2 \quad \frac{1}{2} \quad 3] \end{array}$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \right| \begin{array}{l} \\ -[3 \quad -6 \quad \frac{3}{2} \quad 9] \\ -[1 \quad -2 \quad \frac{1}{2} \quad 3] \end{array} \Rightarrow \left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ & & & \\ & & & \end{array} \right|$$

Example of Gaussian Elimination

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Example of Gaussian Elimination

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Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right|$$

Eliminate the second column
After the second row

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right| \quad -\left(\frac{3}{5}\right) \times \text{row } 2$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right| \quad -\left(\frac{3}{5}\right) \times [0 \ 5 \ -\frac{1}{2} \ 2]$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right| \quad -\left[0 \quad 3 \quad -\frac{3}{10} \quad \frac{6}{5} \right]$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right| \begin{array}{c} \\ \\ -[0 \ 3 \ -\frac{3}{10} \ \frac{6}{5}] \end{array} \Rightarrow \left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 3 & -\frac{3}{2} & -6 \end{array} \right|$$

Example of Gaussian Elimination

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Example of Gaussian Elimination

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$$-\frac{6}{5}x_3 = -\frac{36}{5}$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 0 & -\frac{6}{5} & -\frac{36}{5} \end{array} \right|$$

$$5x_2 - \frac{1}{2}x_3 = 2$$

$$-\frac{6}{5}x_3 = -\frac{36}{5}$$

Example of Gaussian Elimination

$$\left| \begin{array}{cccc} 2 & -4 & 1 & 6 \\ 0 & 5 & -\frac{1}{2} & 2 \\ 0 & 0 & -\frac{6}{5} & -\frac{36}{5} \end{array} \right| \quad \begin{array}{l} 2x_1 - 4x_2 + x_3 = 6 \\ 5x_2 - \frac{1}{2}x_3 = 2 \\ -\frac{6}{5}x_3 = -\frac{36}{5} \end{array}$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6$$

$$5x_2 - \frac{1}{2}x_3 = 2$$

$$-\frac{6}{5}x_3 = -\frac{36}{5}$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6$$

$$5x_2 - \frac{1}{2}x_3 = 2$$

$$-\frac{6}{5}x_3 = -\frac{36}{5} \quad \rightarrow x_3 = 6$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6$$

$$5x_2 - \frac{1}{2}x_3 = 2 \quad \rightarrow 5x_2 - \frac{1}{2} \times 6 = 2$$

$$-\frac{6}{5}x_3 = -\frac{36}{5} \quad \rightarrow x_3 = 6$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6$$

$$5x_2 - \frac{1}{2}x_3 = 2 \quad \rightarrow 5x_2 - \frac{1}{2} \times 6 = 2 \quad \rightarrow x_2 = 1$$

$$-\frac{6}{5}x_3 = -\frac{36}{5} \quad \rightarrow x_3 = 6$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6 \quad \rightarrow 2x_1 - 4 \times 1 + 6 = 6$$

$$5x_2 - \frac{1}{2}x_3 = 2 \quad \rightarrow 5x_2 - \frac{1}{2} \times 6 = 2 \quad \rightarrow x_2 = 1$$

$$-\frac{6}{5}x_3 = -\frac{36}{5} \quad \rightarrow x_3 = 6$$

Example of Gaussian Elimination

$$2x_1 - 4x_2 + x_3 = 6 \quad \rightarrow 2x_1 - 4 \times 1 + 6 = 6 \quad \rightarrow x_1 = 2$$

$$5x_2 - \frac{1}{2}x_3 = 2 \quad \rightarrow 5x_2 - \frac{1}{2} \times 6 = 2 \quad \rightarrow x_2 = 1$$

$$-\frac{6}{5}x_3 = -\frac{36}{5} \quad \rightarrow x_3 = 6$$

Example of Gaussian Elimination

Solve

$$\begin{array}{rrcr} 2x_1 & -4x_2 & +x_3 & = 6 \\ 3x_1 & -x_2 & +x_3 & = 11 \\ x_1 & +x_2 & -x_3 & = -3 \end{array}$$

Gaussian elimination

$$\begin{array}{rrrr} 2 & -4 & 1 & 6 \\ 3 & -1 & 1 & 11 \\ 1 & 1 & -1 & -3 \end{array} \begin{array}{l} \\ \text{row2} - (3/2)*\text{row1} \\ \text{row3} - (1/2)*\text{row1} \end{array} \rightarrow \begin{array}{rrrr} 2 & -4 & 1 & 6 \\ 0 & 5 & -1/2 & 2 \\ 0 & 3 & -3/2 & -6 \end{array} \begin{array}{l} \\ \\ \text{row3} - (3/5)*\text{row2} \end{array}$$

\downarrow

$$\begin{array}{rrrr} 2 & -4 & 1 & 6 \\ 0 & 5 & -1/2 & 2 \\ 0 & 0 & -6/5 & -36/5 \end{array}$$

Backward substitution

$$\begin{aligned} x_3 &= (-36/5) / (-6/5) = 6 \\ x_2 &= (2 + (1/2)*6) / 5 = 1 \\ x_1 &= (6 - 6 + 4*1) / 2 = 2 \end{aligned}$$

Pseudocode and Efficiency of Gaussian Elimination

Stage 1: Reduction to the upper-triangular matrix

```
for i ← 1 to n-1 do
    for j ← i+1 to n do
        for k ← i to n+1 do
             $A[j, k] \leftarrow A[j, k] - A[i, k] * A[j, i] / A[i, i]$ 
```

Stage 2: Backward substitution

```
for j ← n downto 1 do
    t ← 0
    for k ← j +1 to n do
         $t \leftarrow t + A[j, k] * x[k]$ 
     $x[j] \leftarrow (A[j, n+1] - t) / A[j, j]$ 
```

Pseudocode and Efficiency of Gaussian Elimination

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```

Efficiency: $\Theta(n^3) + \Theta(n^2) = \Theta(n^3)$

Discussion

Solve the following system by Gaussian elimination.

$$x_1 + x_2 + x_3 = 2$$

$$2x_1 + x_2 + x_3 = 3$$

$$x_1 - x_2 + 3x_3 = 8.$$

Discussion

$$\begin{bmatrix} 1 & 1 & 1 & 2 \\ 2 & 1 & 1 & 3 \\ 1 & -1 & 3 & 8 \end{bmatrix} \quad \begin{array}{l} \text{row 2} - \frac{2}{1}\text{row 1} \\ \text{row 3} - \frac{1}{1}\text{row 1} \end{array}$$

$$\begin{bmatrix} 1 & 1 & 1 & 2 \\ 0 & -1 & -1 & -1 \\ 0 & -2 & 2 & 6 \end{bmatrix} \quad \text{row 3} - \frac{-2}{-1}\text{row 2}$$

$$\begin{bmatrix} 1 & 1 & 1 & 2 \\ 0 & -1 & -1 & -1 \\ 0 & 0 & 4 & 8 \end{bmatrix}$$

Then, by backward substitutions, we obtain the solution as follows:

$$x_3 = 8/4 = 2, \quad x_2 = (-1 + x_3)/(-1) = -1, \quad \text{and} \quad x_1 = (2 - x_3 - x_2)/1 = 1.$$