

Upgrades as a Service

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http://github.com/Eeko/mediawiki_uuas/

Abstract

This report examines several ways to perform upgrades for service applications without usage downtime by leveraging the possibilities of using on-demand computing resources provided by Infrastructure as a Service -providers. The study has produced a proof-of-concept of a tool capable of providing a limited live-upgrade from MediaWiki 1.4 to 1.5, originally requiring over 22 hours of write-locking to the system when it was applied to Wikipedia in 2005.

1 Introduction

This document is the end report for an independent research project performed for EURECOM¹ semester project done between July 2011 - January 2012. The purpose of the project was to research and demonstrate possibilities to leverage flexible cloud-infrastructure to provide online service updates with very little or no downtime for the end-user. The project was directed by Dr. Tudor Dumitras from Symantec Research Labs and supervised by Prof. Marc Dacier from EURECOM.

2 Problem

The high level issue under research is the possibility of performing system-updates in scale without affecting the availability of service to the end-user. With system updates changing the functionality of software, the usual case requires some unavailability period while modifications to the system are made. For this study, we reviewed the Wikipedia upgrade 1.5 from 2005, which required a 22 hour write lock due to a significant database-schema change requiring a re-write of the entire article database.²

¹<http://www.eurecom.fr/>

²http://meta.wikimedia.org/wiki/MediaWiki_1.5_upgrade

There are two conventional methods to avoid availability breaks in a distributed system. One is to perform the upgrade as a switch-over, where the system is split in two halves. First one part of the system gets updated upgraded while the other part serves the clients. When the update is completed, the updated system is switched to be the client-serving end and the other part applies the update in turn. The second way is to perform the upgrade as a “rolling wave”. Here the upgrade is applied to individual nodes of the distributed system in successive order. This allows for a greater accuracy in failure localization and reduces risks of failures as the entire system (or significant parts of it) do not get compromised for upgrade-errors.

However, neither of these approaches allow a downtimeless upgrade if the upgrade causes backwards incompatibilities. Any updates into the non-updated systems should reflect into the updated system as well. In this study, such incompatibility appears with the significant database-schema change of MediaWiki 1.5.

2.1 Leveraging elastic computing resources for updates

The examined method of avoiding the incompatibility issues with downtimeless upgrade is to use external computing resources to flexibly clone the existing service into a “parallel universe”, where the upgrade can be applied without touching the existing system providing service to the clients. When the upgrade is successfully applied to a machine cloned from a corresponding existing resource, the system performs some kind of catch-up with the changes inserted into the original client-serving machine and starts routing the client requests into itself. [1]

Modern cloud computing infrastructures provide us with a flexible platform for creating and utilizing external resources as needed. Applications running within Infrastructure as a Service (IaaS) providers such as Amazon EC2, Rackspace Cloud Servers or OpenNebula are by default running on virtualized hardware and are thus very easily replicable without large and permanent investments in big hardware.

2.2 MediaWiki 1.4 to 1.5 upgrade

The June 2005 update to MediaWiki 1.5 was primarily introduced to perform the schema change examined in this project. Few of the new features, such as logging the page-rename history and revised permalinking are dependant on the new schema. But those features do not interfere with online-upgrading, since the relevant tables can be re-generated separately from the existing database if needed.

In the 1.4 version of the schema, individual articles contained entries in two tables. The “cur”-table contained the most recent revisions of the articles in their entirety. The “old”-table contained the corresponding article history, usually listing up several related (old wikipedia revisions) entries for a single article in cur-table. The entries in the tables are connected by the cur-table unique id of “cur_namespace + cur_title”. A completely new article appears as a new

In the 1.5 version, the tables are split into three tables. The new “revision” table is formed by combining the previous old- and cur tables and is intended to represent the relevant metadata for all article insertions and updates of the system. Such as the user who edited the page, the relevant timestamps and comments of the new update. A revision-entry is connected to a relevant entry in “text”-table, which contain the articles themselves. An entry in revisions also connect to an entry in the “page”-table, which represents an individual article with a revision-history in the revision-table. The most current article is stored in the “page latest” pointer towards a singular entry in revision-table.

```

cur_id                page_id
cur_namespace         page_namespace
cur_title             page_title
cur_text              page_restrictions
cur_comment           page_counter
cur_user              page_is_redirect
cur_user_text         page_is_new
cur_timestamp         page_random
cur_restrictions      page_touched
cur_counter           page_latest
cur_is_redirect
cur_minor_edit
cur_is_new
cur_random
cur_touched
inverse_timestamp

d:
old_id
old_namespace
old_title
old_text
old_comment
old_user
old_user_text
old_timestamp
old_minor_edit
old_flags
inverse_timestamp

revision:
rev_id
rev_page
rev_comment
rev_user
rev_user_text
rev_timestamp
inverse_timestamp
rev_minor_edit

text:
old_id
old_text
old_flags

```

Figure 1: Database Schema Changes in MediaWiki 1.5 [3]

3 Solution

We built a small prototype of a software-stack capable of reading the MySQL query logs in real time from the system providing service to the end-user. Whenever it detects an update to the article-tables under update, it would create a translation of those queries compatible with the new schema. After the standard, non-modified 1.4 to 1.5 update is applied to the parallel universe clone of the system, we use an external program hooking into the new database. This program uses the recorder modifications to mimic inserting article updates to the updated database. When the system under update has reached a synchronous state with the live-system, we can shut down the old system and the upgrade programs and route all traffic to the updated system.

3.1 Implementation details

The original test-system under study is a small-sized³ Amazon EC2-instance running a software stack⁴ capable of running MediaWiki 1.4 with a custom test-database for a set of test articles.

3.1.1 System replication tools

To automate the system replication procedure, we developed a series of bash-scripts leveraging the Amazon EC2-tools and knowledge of the details of the system under upgrade. Mainly we require the Amazon instance running details (instance number, hostname) and the application information (database name, host, username and password) for running the replication stack. The system is designed in a way, that we can use an external node with ssh- and EC2-tools access to the Amazon Instances to download the necessary programs from repository and start performing the upgrade process centrally, without touching the running instance providing service for the clients.

The main scripts to initiate the upgrading process are as follows:

- `configs.conf` – A sourceable configuration file to set the required environmental variables in the bash-scripts. Requires manual modifications to point to the EC2-node to be replicated and for the necessary database knowledge.
- `prepare_for_cloning.sh`⁵ – Intended for installing a necessary stack of software to the node to be replicated. Such as Python 2.7 and mysql-python required by the updater software. Should also ensure that the required program-versions are available for the updating scripts at the locations specified in them.

³<http://aws.amazon.com/ec2/instance-types/>

⁴Amazon Linux 2011.2 with PHP 5.2, Apache 2 and MySQL 5.1

⁵Not implemented yet as of February 2, 2012.

- `create_aws_replica.sh` – Initiates the cloning process by copying the targeted node disk-image into a Amazon AMI (Requiring a brief shutdown of the said instance.) and starting a new identical instance with the said image. Creates a modified. `configs.replica` -file to include the necessary instance details of the replicated instance needed by the rest of the scripts.
- `setup_replica.sh` – Copies the necessary scripts and programs into the new instance.
- `start_update.sh` – Makes some necessary database-access modifications into the query-translating programs and launches an SSH-pipe into the original MediaWiki 1.4 node to stream the query log into a file to be readable by the local transaction-catchup programs. This is to be run in the new, replicated instance.
- `std_update_mwiki14-15.sh` – This script contains tools to download the newer MediaWiki 1.5 version and for running the standard upgrade-procedure to create a new copy of MediaWiki 1.5 running on top of the restructured database. This is to be run in the new, replicated instance and requires superuser access for the necessary Apache configuration and reboots.

3.1.2 Query parser, translator and mapper

The rest of the software is a series of python-programs used by `update_mediawiki.py`. `Update_mediawiki.py` requires the path to the file where the original systems query-log has been streamed as an argument and eventually writes all updates detected for existing articles into the new database schema within the parallel universe.

The program components are as follows:

- `update_mediawiki.py` – The entry-point of the program. Contains a `main()` method executing the translator-program from `translator.py`.
- `translator.py` – Contains the logic needed to use the query-log parser program (`parser.py`), how to interpret its returns and to translate them into SQL-queries writable by the database-hookup component. (`mysql_connect.py`)
- `parser.py` – Contains the logic required to detect and parse relevant INSERT and UPDATE queries from MySQL query-logs. The lines we wish to detect are the ones making article-updating modifications into an original MediaWiki 1.4 database.
- `mysql_connect.py` – Helper methods used to interact with a MySQL database.

3.2 Problems with the implementation

Our query-detector and translator approach mainly requires us to understand and re-implement large portions of the application logic within an external

framework. This requires a significant amount of manual labor and would intuitively be more suitable to be integrated directly into the standard upgrading mechanisms instead of providing an external framework. More notably, individual upgrade-instructions are not recyclable for other upgrades; neither can we leverage the existing SQL-upgrade instructions to automate the logic-programming.

Another issue is the limits of the data extractable from the query logs. Much of the details ending up to the database can be programmed and computed to be performed by the database itself without necessarily revealing them in the query logs. E.g. generating entries via database-triggers and the auto-increments of id's are usually done within the database and can't be read from default query-logs. In the lower levels the database may perform optimizations or transaction aborts not necessarily visible to the logs. Or they can be insufficiently hard to predict and react for in large scale parsing. For adequate understanding of the workings of the database, the visible plain-text query-logs are likely insufficient. The approach would be more suited to be done by using the existing database replication infrastructures and binary-logs, which reliably reveal the internal database-actions in detail.

Though due to the elasticity of computing resources and low cost of upgrade-failures, failed upgrades can be tried again as often as needed. This opens us for possibilities of performing a less-refined probabilistic upgrade, where we only need a chance for individual upgrade to succeed and a number of computers performing the upgrade. After a unit passes tests for schema-equivalence, one can use standard database-replication suites to duplicate the relevant infrastructure to match the existing system.

The implementation looks into the updates as individual transaction one at a time, which is necessary as the program simulates a working application performing similar actions in a live use scenario. This is hardly efficient for larger data sets and is somewhat error prone; should there be unexpected modifications (such as manual inserts) to the database not detectable by the developed application. Another way would be to use the query-logs to create records of data requiring action after a stage of upgrade has been completed. For example, two updates to the same article could be marked as a single entry to a table of "touched" article-id's. Then we make an external query to the original database to stream the necessary changes into the parallel universe. This does not free us from implementing some application logic, as actions such as deleting rows or modifying their unique id's would have to be represented in the tracking logic of tainted-entries. Neither is it granted that the query-logs available present us with enough data to identify the tainted items. For example, an INSERT-query might enter their unique id as NULL and auto-increment it in the database or application-logic. Such incrementation based on the MAX(ID)-value of the new flattened text-table of MediaWiki 1.5 was required in our implementation.

It is necessary to note that the current implementation does not manage fully without downtime, since creating an identical real-time duplicate in Amazon EC2 requires a downtime to make a copy of the image. However, it should be possible to make relevant replications in virtualized production environments, since equivalent copies of the database can be made without shutting down, by using the standard redundancy replication procedures provided by every major

RDBMS.

4 Alternate approaches

During the course of the study, several other methods of performing the online-upgrade were speculated of and experimented with.

4.1 Using existing database-replication

One approach we experimented with was trying to leverage an existing database-reflection infrastructure. Initially, the GORDA database replication toolkit⁶ showed much promise. With a flexible middleware providing accurate reflections and access to the inner workings of the database, we would be assured to have all the necessary data needed for keeping track of the changes during the upgrade. Ideally we could use similar architecture used for redundancy replications of a database and modify their outputs in a way which would accommodate for the changes in database schema. GORDA promised to offer a flexible API to those reflections given that we'd use a supported database.[2]

However, there proved to be a number of issues with this approach. First of all, the tried toolkit for GORDA somewhat compatible with MediaWiki was one hooking up into PostgreSQL database. For MediaWiki 1.4, the PostgreSQL -support was considered to be “experimental” and it was recommended to use MySQL for production database. After the schema upgrade in 1.5, the PostgreSQL-support was officially discarded, though the unmodified components were still within the source code. Making the software run on PostgreSQL (or Apache Derby, which is “default” database for GORDA) requires a number of extra modifications.

Secondly, GORDA itself have more or less been abandoned recently. The documentation is in inconsistent state and in some cases become completely unavailable. For example, the documentation wiki of Escada Replication Server (which is required by later versions of GORDA) has disappeared since. Utilizing GORDA proved to require a tremendous amount of reverse engineering, which became even more troublesome as the database-hookup components lacked a number of important features, such as namespaces for database tables.

Should GORDA have worked and if we'd had enough time to figure out how to port MediaWiki 1.5 to PostgreSQL or another more supported database, it would definitely solve many problems imminent with our other approaches. Such as the consistency problems with application-call or incremental standard upgrading for renewing database subsets.

4.2 Using similarity in application calls

Another idea to provide upgrades as a service for an ongoing database upgrade would be to move the upgrade-synchronization entirely away from the

⁶<http://gorda.di.uminho.pt/>

database-layer. If an upgrade touches only the underlying database layer and the application interface connecting to it, one could cache and re-route identical application calls to a parallel-universe backend replicating similar functionality within different schema. This kind of upgrade would naturally suit a typical 3-tiered web-application, where the user-transactions provided by web-server are separated from a dynamic content engine and data storages.

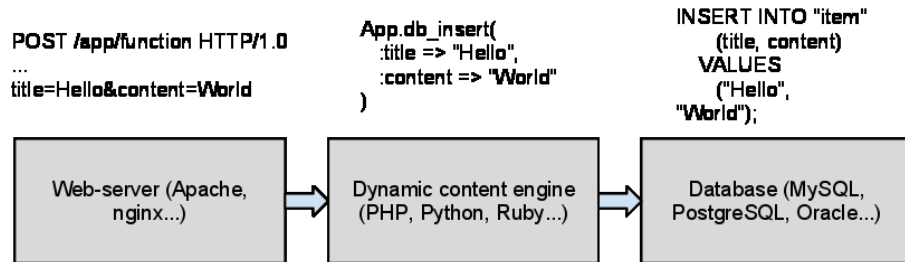


Figure 2: A typical 3-tiered web application

Should the interface to the frontend-engine stay identical, this kind of approach would provide a fairly easy and elegant solution for making a seamless transition to new software version without requiring any adaptations or query mappings between the systems. In the case of MediaWiki 1.4 to 1.5 upgrade, the main focus of the upgrade was this database-schema change. The upgrade could have been split into parts only affecting the application & database layers whilst keeping the user facing web server interface the same. Should an upgrade provide any UI modifications, those can be made in an another upgrade-package keeping the database untouched.

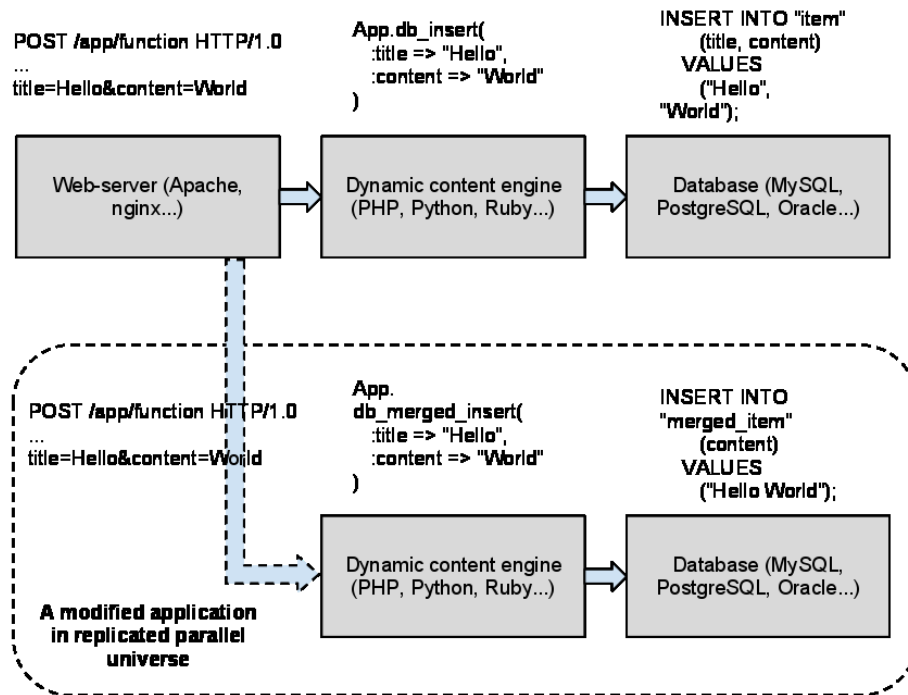


Figure 3: A 3-tiered web application routing interface-queries to original and upgraded backends

Not every kind of software update would work with this approach. Several upgrades implement new functionality with a need for wide modifications to every layer of the software. Though often such upgrades can be split into incremental parts where the heavy database-modifying and writing operations can be done before the corresponding modifications are introduced to the frontend. As this requires additional engineering-consideration for upgradeability, it might not be feasible to provide upgrading as an external service for application-upgrades modifying the entire stack.

Other issue would be the lack of guarantees for consistency in the entire stack. Should there be failures in the middle- or database tier of the live-system, a dumb frontend-replicator would not detect them nor present adequate information to account the inconsistency in the parallel universe. Such issues could be coped with most standard redundancy & reliability techniques utilized in the system. The redundancy technology can be made to monitor and require confirmation of commits from the backend or from redundancy replication interfaces. Or with cheap replicable virtual hardware, we could settle with eventual consistency where we just restart an upgrade-procedure until the backend passes sufficient tests of equivalence.

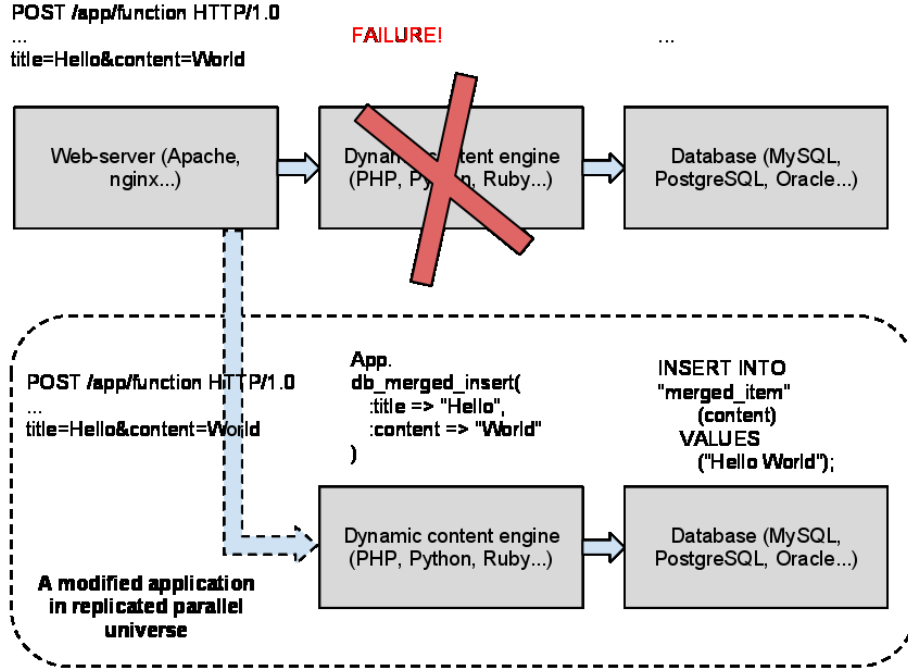


Figure 4: A faulty parallel-run leaving the two systems in inequivalent states

4.3 Using the existing upgrade tools to decreasing database subsets

One of the more intriguing approaches would be to create a framework to be able to read the existing schema-upgrade scripts available and deduct the upgrade logic and resulting table from those. Since the target tables are expanding leading to increasing time-requirements of re-applying the updates to complete tables, we need to be able to divide the work into smaller subsets as new items and updates get inserted during the online-upgrade.

An intuitive way for such division would be to split the dataset under update by their timestamps, so that we only re-run the standard upgrade script for new items inserted after the last known item in the databases under upgrade was received. However, under some upgrades this will provide an incompatible and possibly broken result due to the unpredictability of the live-system updates.

In an example, we have a database of two tables representing a list of current states and saved histories of those states:

state_id	state_content
1	1_Fourth_state
2	2_Second_state

Table 1: State-table

history_id	history_content	history_link_to_state
1	1_First_state	1
2	1_Second_state	1
3	2_First_state	2
4	1_Third_state	1

Table 2: History-table

Suppose, that a schema upgrade would flatten these said tables into one table containing both the current state of the items and the given history of said items. An upgrade would be done with the following SQL-code:

— Note, that the primary id's of the table are auto-incremented.
— This is similar to how Mediawiki 1.4 to 1.5 upgrade handles the
— database-flatten operation.

```
INSERT INTO "history" (history_content , history_link_to_state)
SELECT (state_content , state_id)
FROM "state";
```

After the update, the new table would look like this:

history_id	history_content	history_link_to_state
1	1_First_state	1
2	1_Second_state	1
3	2_First_state	2
4	1_Third_state	1
5	1_Fourth_state	1
6	2_Second_state	2

Table 3: Merged history-table

However, suppose that we receive a third state to the original-system during the time taken by the upgrade of the system and it receives several updates for it. We have sufficient translation logic in place to only apply the INSERT-queries for items entered to the database after we begun merging our previous entries. The code would work something like this:

— We first apply the modifications from history-table

```
INSERT INTO "history" (history_content , history_link_to_state)
SELECT (history_content , state_id)
FROM "olddb.history"
WHERE timestamp > last_update_time;
```

— Then we flatten the items from the table of current states

```
INSERT INTO "history" (history_content , history_link_to_state)
SELECT (state_content , history_link_to_state)
FROM "olddb.state"
WHERE timestamp > last_update_time;
```

Then when we would receive the following rows into the database:

history_id	history_content	history_link_to_state
5	3_First_state	3

Table 4: New row in history-table inserted during the update

state_id	state_content
3	3_Second_state

Table 5: New row in state-table inserted during the update

We would end up with a merged table looking like the following:

history_id	history_content	history_link_to_state
1	1_First_state	1
2	1_Second_state	1
3	2_First_state	2
4	1_Third_state	1
5	1_Fourth_state	1
6	2_Second_state	2
7	3_First_state	3
8	3_Second_state	3

Table 6: Merged history-table after incremental upgrade

If the upgrade would have been done with write-locks and the two tables would be merged after all 8 commits were received in the same sequential order as in our online-example, the resulting table would look like this:

history_id	history_content	history_link_to_state
1	1_First_state	1
2	1_Second_state	1
3	2_First_state	2
4	1_Third_state	1
5	3_First_state	3
6	1_Fourth_state	1
7	2_Second_state	2
8	3_Second_state	3

Table 7: Merged-history table without online-upgrading

The id's for items in the different upgrade-approaches are not equivalent. If it is used as an external-id in somewhere else in database-logic without the necessary modifications, we will encounter in faulty behaviour. To fix this, we need to implement some amount of application- or database-logic to upgrade corresponding tables with history_id references.

5 Conclusions

The implementation provided fails in its initial goal of performing an online upgrade in a scale comparable to a real-life use scenario. (I.e. the entire Wikipedia and updates performed to it during the upgrade.) Though it manages to illustrate the difficulties in the approach. Namely the requirement to model much of the application logic within the mapper and the unreliability of the used data-source. (Human readable query-logs from MySQL.)

In addition, we have examined a number of other approaches capable of providing a non-downtime upgrade-procedure in a replicated system and shown via contradiction, why some of them would be inadequate for selected, common database-modifying procedures. Though used as an initial approach for the upgrade, we were unsuccessful in utilizing the most promising method of leveraging database-reflection interfaces.

5.1 Future work

Although the approach with GORDA proved to be unfeasible in this scope, different methods to leverage the reflections from database-internals to provide upgrades as a service are still to be examined. Another possible ways to get sufficient data could be to hook up into existing database replication protocols or into the binary-logs used by the said replication protocols.

And even if the other introduced methods to perform upgrades externally in replicated environments are not generalizable, we have still to examine whether they might be sufficient for individual cases such as the MediaWiki upgrade presented. Especially given the simplicity of rerouting high-level application calls for replicated cloud-servers, designing upgrades to support it could prove to be a decent engineering practice for online-services.

References

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