CS 300 HOMEWORK 5

Q1)

Step 1

	Weight Knowi	
Α	infinity	False
В	infinity	False
С	2	False
D	infinity	False
Е	infinity	False
F	infinity	False
G	0	True
Н	2	False

Step 2

	Weight	Known
Α	Infinity	False
В	6	False
С	2	True
D	infinity	False
Е	infinity	False
F	8	False
G	0	True
Н	2	False

Step 3

	Weight	Known
Α	6 False	
В	6	False
С	2	True
D	infinity	False
Е	infinity	False
F	8	False
G	0	True
Н	2	True



Step 6

	Weight	Known
Α	6	True
В	6	True
С	2	True
D	10	False
Е	7	True
F	8	False
G	0	True
Н	2	True

Step 5

	Weight Knowr	
Α	6	True
В	6	True
С	2	True
D	infinity	False
Е	7	False
F	8	False
G	0	True
Н	2	True

Step 4

	Weight	Known
Α	6	True
В	6	False
С	2	True
D	infinity	False
Е	15	False
F	8	False
G	0	True
Н	2	True



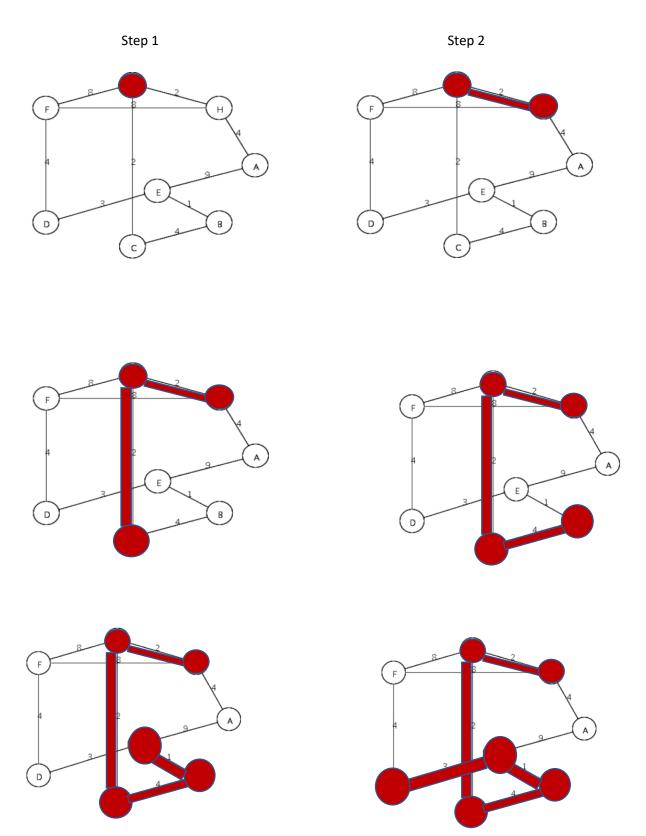
Step 7

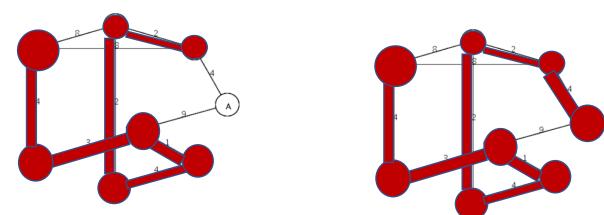
	Weight	Known
Α	6	True
В	6	True
С	2	True
D	10	False
Е	7	True
F	8	True
G	0	True
Н	2	True

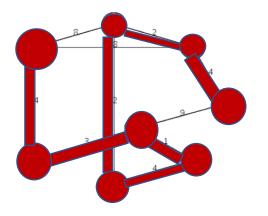
Step 8

	Weight Know	
Α	6	True
В	6	True
С	2	True
D	10	True
Е	7	True
F	8	True
G	0	True
Н	2	True

Selected vertices and edges are shown in red for every step.

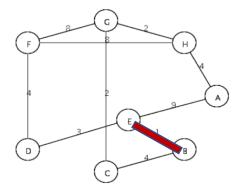




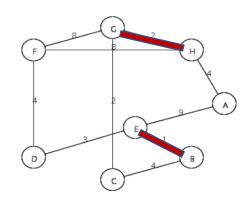


The selected edges are shown in red at every step of the program.

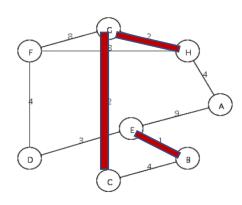
Step 1



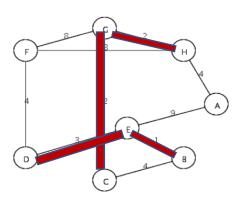
Step 2



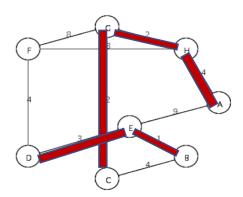
Step 3



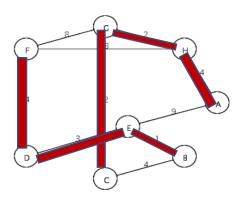
Step 4



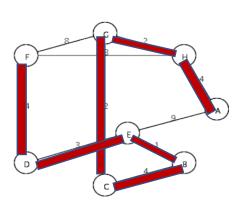
Step 6



Step 7







Q4)

Step 1: Queue: S

Step 2: Queue: S - B - A - D

Step 3: Queue: S - B - A - D (We can not add new vertex to the queue since there is no adjacent vertex that is unknown)

Step 4: Queue: S - B - A - D - C

Step 5: Queue: S - B - A - D - C - E - T - F

Step 6: Queue: S - B - A - D - C - E - T - F

Step 7: Queue: S - B - A - D - C - E - T - F - G

Step 8: Queue: S - B - A - D - C - E - T - F - G

Step 9: Queue: S - B - A - D - C - E - I - F - G

Step 10: Queue: S - B - A - D - C - E - T - F - G

Our breadth first search terminates until the queue is empty.

If I print the vertices after I dequeue it from the queue, the output will be below.

Output: S - B - A - D - C - E - T - F - G

a) Step 1: Stack: S

Step 2: Stack:

В

S

Step 3: Stack

D

В

S

Step 4: Stack:

Ε

D

В

S

Step 5: Stack:

G

Ε

D

В

S

Step 6: Stack:

Т

G

E D

В

S

Step 7: Stack:	F
	Т
	G
	E
	D
	В
	S

Step 8: Stack:	Т
	G
	Ε
	D
	В
	S

Step 13: Stack:

Step 14: Stack:

Step 15: Stack:

Step 16: Stack:

S

Step 17: Stack:

Our DFS terminate when there is no element in the stack.

Output: S - B - D - E - G - T - F - A - C (If I print the vertices that I visit that are unknown).

b) Post Orders:

S: 9

B: 6

A: 8

D: 5

C: 7

E: 4

G: 3

T: 2

F: 1

c) Pre Orders:

S: 1

B: 2

A: 8

D: 3

C: 9

E: 4

G: 5

T: 6

F: 7

d)

Tree arcs: (S, B), (B, D), (D, E), (E, G), (G, T), (T, F), (S, A), (A, C)

Forward arcs: (S, D), (D, T), (D, F),

Backward arcs: (B, S), (G, E)

Cross arcs: (C, E), (E, T), (C, D)

Q6) Since the vertex A and G are degree 0, I enqueue them in to the queue(I picked the first one arbitrarily), after that, when I dequeue the vertex from the queue, I looked the adjacent vertices of that vertex and if the vertex's indegree is zero when I decrement the indegree of that vertex by one I add it to the queue.

$$A-G-B-D-C-E-F$$

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