## Notebook UNosnovatos

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# $1 \quad C++$

# 1.1 C++ plantilla

```
#include <bits/stdc++.h>
using namespace std;
#define sz(arr) ((int) arr.size())
typedef long long 11;
typedef pair<int, int> ii;
typedef vector<ii> vii;
typedef vector<int> vi;
typedef vector<long long> v1;
const int INF = 1e9;
const ll INFL = 1e18;
const int MOD = 1e9+7;
const double EPS = 1e-9;
int dirx[4] = \{0, -1, 1, 0\};
int diry[4] = \{-1,0,0,1\};
int dr[] = \{1, 1, 0, -1, -1, -1, 0, 1\};
int dc[] = \{0, 1, 1, 1, 0, -1, -1, -1\};
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0);
    // freopen("file.in", "r", stdin);
    // freopen("file.out", "w", stdout);
    return 0;
```

# 2 Estructuras de Datos

# 2.1 Disjoint Set Union

```
struct dsu{
    vi p, size;
    int num_sets;
    int maxSize;

    dsu(int n) {
        p.assign(n, 0);
            size.assign(n, 1);
            num_sets = n;
            for (int i = 0; i<n; i++) p[i] = i;
    }

int find_set(int i) {return (p[i] == i) ? i : (p[i] = find_set(p[i]));}

bool is_same_set(int i, int j) {return find_set(i) == find_set(j);}

void unionSet(int i, int j){</pre>
```

```
if (!is_same_set(i, j)){
                 int a = find_set(i), b = find_set(j);
                 if (size[a] < size[b])</pre>
                      swap(a, b);
                 p[b] = \bar{a};
                 size[a] += size[b];
                 maxSize = max(size[a], maxSize);
                 num_sets--;
};
```

## 2.2 Fenwick Tree

```
#define LSOne(S) ((S) & -(S))
struct fenwick_tree{
    vl ft; int n;
    fenwick_tree(int n): n (n) {ft.assign(n+1, 0);}
    ll rsq(int j) {
        11 \text{ sum } = 0;
        for(; j; j -= LSOne(j)) sum += ft[j];
        return sum;
    ll rsq(int i, int j) {return rsq(j) - (i == 1 ? 0 :
       rsq(i-1));}
    void upd(int i, ll v) {
        for (; i <= n; i += LSOne(i)) ft[i] += v;</pre>
};
```

# 2.3 Segment Tree

```
int nullValue = 0;
struct nodeST{
    nodeST *left, *right;
    int 1, r; 11 value, lazy, lazy1;
    nodeST(vi \&v, int l, int r) : l(l), r(r) {
        int m = (1+r) >> 1;
        lazv = 0;
        lazv1 = 0;
        if (1!=r) {
            left = new nodeST(v, 1, m);
            right = new nodeST(v, m+1, r);
            value = opt(left->value, right->value);
        else
            value = v[1];
```

```
ll opt(ll leftValue, ll rightValue) {
    return leftValue + rightValue;
void propagate(){
    if(lazv1){
        value = lazv1 * (r-l+1);
        if (1 != r) {
            left->lazy1 = lazy1, right->lazy1 = lazy1
            left->lazy = 0, right->lazy = 0;
        lazv1 = 0;
        lazv = 0:
    else{
        value += lazy * (r-l+1);
        if (1 != r) {
            if(left->lazy1) left->lazy1 += lazy;
            else left->lazy += lazy;
            if(right->lazy1) right->lazy1 += lazy;
            else right->lazv += lazv;
        lazv = 0;
ll get(int i, int j){
    propagate();
    if (l>=i && r<=j) return value;</pre>
    if (l>j || r<i) return nullValue;</pre>
    return opt(left->get(i, j), right->get(i, j));
void upd(int i, int j, int nv) {
    propagate();
    if (1>j || r<i) return;</pre>
    if (1>=i && r<=j) {
        lazy += nv;
        propagate();
        // value = nv;
        return;
    left->upd(i, j, nv);
    right->upd(i, j, nv);
    value = opt(left->value, right->value);
void upd(int k, int nv) {
    if (1>k || r<k) return;</pre>
    if (1>=k && r<=k) {
        value = nv;
        return;
```

```
3 GRAFOS
```

```
left->upd(k, nv);
    right->upd(k, nv);

    value = opt(left->value, right->value);
}

void upd1(int i, int j, int nv) {
    propagate();
    if (l>j || r<i) return;
    if (l>=i && r<=j) {
        lazy = 0;
        lazy1 = nv;
        propagate();
        return;
    }

    left->upd1(i, j, nv);
    value = opt(left->value, right->value);
}

};
```

# 3 Grafos

#### 3.1 DFS

```
//O(V+E)
int vertices, aristas;
vector<int> dfs_num(vertices+1, -1); //Vector del estado de cada vertice (visitado o no visitado)

const int NO_VISITADO = -1;
const int VISITADO = 1;
vector<vector<int>> adj(vertices + 1); //Lista adjunta del grafo

// Complejidad O(V + E)
void dfs(int v){
    dfs_num[v] = VISITADO;
    //Se recorren los vecinos
    for (int i = 0; i < (int) adj[v].size(); i++){
        if (dfs_num[adj[v][i]] == NO_VISITADO) {
            dfs(adj[v][i]);
        }
    }
}</pre>
```

## 3.2 BFS

## 3.3 Puntos de articulación y puentes

```
//Puntos de articulacion: son vertices que desconectan el
    grafo
//Puentes: son aristas que desconectan el grafo
//Usar para grafos dirigidos
vi dfs_num, dfs_low, dfs_parent, articulation_vertex;
int dfsNumberCounter, dfsRoot, rootChildren;
vector<vii> adj;
void articulationPointAndBridge(int u) {
    dfs_num[u] = dfsNumberCounter++;
    dfs_low[u] = dfs_num[u]; // dfs_low[u] <= dfs_num[u]
    for (auto &[v, w] : adj[u]) {
        if (dfs_num[v] == -1) { // una arista de arbol
            dfs parent[v] = u;
            if (u == dfsRoot) ++rootChildren; // vaso
               especial, raiz
            articulationPointAndBridge(v);
            if (dfs low[v] >= dfs num[u]) // para puntos
               de articulacion
                articulation vertex[u] = 1;
            if (dfs_low[v] > dfs_num[u]) // para puentes
                printf(" (%d, %d) is a bridge\n", u, v);
            dfs_low[u] = min(dfs_low[u], dfs_low[v]); //
        else if (v != dfs parent[u]) // si es ciclo no
           trivial
            dfs low[u] = min(dfs low[u], dfs num[v]); //
               entonces actualizar
int main(){
    dfs_num.assign(V, -1); dfs_low.assign(V, 0);
```

# 3.4 Orden Topologico

```
//Orden de un grafo estilo malla curricular de
    prerrequisitos
vector<vi> adj;
vi dfs_num;
vi ts;

void dfs(int v) {
    dfs_num[v] = 1;
    for (int i = 0; i < (int) adj[v].size(); i++) {
        if (dfs_num[adj[v][i]] != 1) {
            dfs(adj[v][i]);
        }
    }
    ts.push_back(v);
}
//Imprimir el vector ts al reves: reverse(ts.begin(), ts.end());</pre>
```

# 3.5 Algoritmo de Khan

```
//ALgoritmo de orden topologico
//DAG: Grafo aciclico dirigido
int n, m;
vector<vi> adj;
vi grado;
vi orden;

void khan() {
    queue<int> q;
```

```
for (int i = 1; i<=n; i++) {</pre>
        if (!grado[i]) q.push(i);
    int nodo;
    while(!q.empty()){
        nodo = q.front(); q.pop();
        orden.push_back(nodo);
        for (int v : adj[nodo]) {
            grado[v]--;
            if (qrado[v] == 0) q.push(v);
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0);
    cin >> n >> m;
    adj.resize(n+1);
    grado.resize(n+1);
    for (int i = 0; i<m; i++) {
        int x, y; cin >> x >> y;
        adj[x].push_back(y);
        grado[y]++;
    khan();
    if (orden.size() == n) {
        for (int i : orden) cout << i;</pre>
    else{
        cout << "No DAG"; //No es un grafo aciclico</pre>
            dirigido (tiene un ciclo)
```

## 3.6 Floodfill

```
//Relleno por difusion-etiquetado/coloreado de
    componentes conexos
//Recorrer matrices como grafos implicitos
//Pueden usar los vectores dirx y diry en lugar de dr y
    dc si se requiere
vector<string> grid;
int R, C, ans;
int floodfill(int r, int c, char c1, char c2) {
    //Devuelve tamano de CC
    if (r < 0 || r >= R || c< 0 || c >= C) return 0;
        //fuera de la rejilla
    if (grid[r][c] != c1) return 0;
        //No tiene color c1
```

# 3.7 Algoritmo Kosajaru

int ans = 1;

return ans;

cin.tie(0);

cin >> R; cin >> C;

int main() {

qrid[r][c] = c2;

vertice (r, c) tiene color c1

ans += floodfill(r + dr[d], c + dc[d], c1, c2);

c) a c2 para evitar ciclos

cout << floodfill(0, 0, 'W', '.');

for (int d = 0; d < 8; d++) {

ios::sync\_with\_stdio(false);

```
//Encontrar las componentes fuertemente conexas en un
   grafo dirigido
//Componente fuertemente conexa: es un grupo de nodos en
   el que hav
//un camino dirigido desde cualquier nodo hasta cualquier
    otro nodo dentro del grupo.
void Kosaraju(int u, int pass) {
    dfs num[u] = 1;
    vii &neighbor = (pass == 1) ? AL[u] : AL_T[u];
    for (auto &[v, w] : neighbor)
        if (dfs num[v] == UNVISITED)
            Kosaraju(v, pass);
    S.push back(u);
int main(){
    S.clear();
    dfs num.assign(N, UNVISITED);
    for (int u = 0; u < N; ++u)
        if (dfs num[u] == UNVISITED)
            Kosaraju(u, 1);
    numSCC = 0;
    dfs num.assign(N, UNVISITED);
    for (int i = N-1; i >= 0; --i)
        if (dfs num[S[i]] == UNVISITED)
            ++numSCC, Kosaraju(S[i], 2);
    printf("There are %d SCCs\n", numSCC);
```

//suma 1 a ans porque el

//Colorea el vertice (r.

# 3.8 Dijkstra

```
//Camino mas cortos
```

```
//NO USAR CON PESOS NEGATIVOS, usar Bellman Ford o SPFA(
   mas rapido)
// O ((V+E) *log V)
vi dijkstra(vector<vii> &adj, int s, int V) {
    vi dist(V+1, INT_MAX); dist[s] = 0;
    priority queue<ii, vii, greater<ii>> pq; pq.push(ii
        (0, s);
    while(!pq.empty()){
        ii front = pq.top(); pq.pop();
        int d = front.first, u = front.second;
        if (d > dist[u]) continue;
        for (int j = 0; j < (int)adj[u].size(); j++) {</pre>
            ii v = adi[u][i];
            if (dist[u] + v.second < dist[v.first]){</pre>
                dist[v.first] = dist[u] + v.second;
                pq.push(ii(dist[v.first], v.first));
    return dist;
```

# 3.9 Bellman Ford

```
//Lo mismo que dijkstra pero con pesos negativos
//O(E*V)
void bellman_ford() {
    vi dist(V, INF); dist[s] = 0;
    for (int i = 0; i < V-1; ++i)
        for (int u = 0; u < V; ++u)
            if (dist[u] != INF)
            for (auto &[v, w] : adj[u])
                dist[v] = min(dist[v], dist[u]+w);
```

# 3.10 Floyd Warshall

```
//Camino minimo entre todos los pares de vertices
int main() {
    ios::sync_with_stdio(false);
    cin.tie(0);
    int V; cin >> V;
    vector<vi> adjMat(V+1, vi(V+1));
    //Condicion previa: adjMat[i][j] contiene peso de la
       arista (i, j)
    //o INF si no existe esa arista
    for (int k = 0; k < V; k++)
        for (int i = 0; i<V; i++)
            for (int j = 0; j<V; j++)
                adjMat[i][j] = min(adjMat[i][j], adjMat[i
                   [k] + adjMat[k][j]);
```

#### 3.11 MST Kruskal

```
//Arbol de minima expansion
//O(E*log V)
int main() {
    int n, m;
    cin >> n >> m;
    vector<pair<int, ii>> adj; //Los pares son: {peso, {
       vertice, vecino}}
    for (int i = 0; i<m; i++) {</pre>
        int x, y, w; cin >> x >> y >> w;
        adj.push_back(make_pair(w, ii(x, y)));
    sort(adj.begin(), adj.end());
    int mst costo = 0, tomados = 0;
    dsu UF(n);
    for (int i = 0; i<m && tomados < n-1; i++) {</pre>
        pair<int, ii> front = adj[i];
        if (!UF.is same set(front.second.first, front.
            second.second)){
            tomados++;
            mst costo += front.first;
            UF.unionSet(front.second.first, front.second.
                second);
    cout << mst_costo;</pre>
```

# 3.12 Shortest Path Faster Algorithm

```
//Algoritmo mas rapido de ruta minima
//O(V*E) peor caso, O(E) en promedio.
ll spfa(vector<vii>& adj, ll s, ll n) {
  vl d(n+1, INFL);
  vector<bool> inqueue(n, false);
  queue<ll> q;

  d[s] = 0;
  q.push(s);
  inqueue[s] = true;
  while (!q.empty()) {
    ll v = q.front();
    q.pop();
    inqueue[v] = false;

  for (auto edge : adj[v]) {
    ll to = edge.first;
```

```
ll len = edge.second;
if (d[v] + len < d[to]) {
        d[to] = d[v] + len;
        if (!inqueue[to]) {
            q.push(to);
            inqueue[to] = true;
        }
    }
}
return d[n];
}</pre>
```

## 4 Matematicas

#### 4.1 Criba de Eratostenes

```
// O(N log log N)
ll sieve size;
bitset<10000010> bs;
                        //10^7 es el limite aprox
                        //Lista compacta de primos
void sieve(ll upperbound) {
                                     //Rango = [0..limite]
    sieve size = upperbound+1;
                                     //Para incluir al
       limite
    bs.set();
                                     //Todo unos
    bs[0] = bs[1] = 0;
                                     //0 y 1 (no son
       primos)
    for (ll i = 2; i < _sieve_size; ++i) if (bs[i]) {</pre>
        for (ll j = i*i; j < _sieve_size; j += i) bs[j] =
        p.push_back(i);
                                    //Anadir primo i a la
            lista
```

# 4.2 Descomposicion en primos (y mas cosas)

```
// O( sqrt(N) / log(sqrt(N)) )
vl primeFactors(ll N) {
    vl factors;
    for (int i = 0; (i < (int)p.size()) && (p[i]*p[i] <=</pre>
        while (N%p[i] == 0) {
                                     //Hallado un primo
           para N
            N /= p[i];
                                     //Eliminarlo de N
            factors.push_back(p[i]);
    if (N != 1) factors.push back(N); //El N restante es
       primo
    return factors;
int main(){
    sieve(10000000);
//Variantes del algoritmo
//Contar el numero de divisores de N
int numDiv(ll N) {
    int ans = 1;
                 //Empezar con ans = 1
    for (int i = 0; (i < (int)p.size()) && (p[i]*p[i] <=</pre>
       N); ++i) {
        int power = 0; //Contar la potencia
        while (N%p[i] == 0) \{ N /= p[i]; ++power; \}
        ans *= power+1; //Sequir la formula
    return (N != 1) ? 2*ans : ans; //Ultimo factor = N^1
//Suma de los divisores de N
//N = a^i * b^i * ... * c^k => N = (a^i+1) - 1) / (a-1)
   + ...
ll sumDiv(ll N) {
    ll ans = 1;
                        // empezar con ans = 1
    for (int i = 0; (i < (int)p.size()) && (p[i]*p[i] <=</pre>
       N); ++i) {
        ll multiplier = p[i], total = 1;
        while (N%p[i] == 0) {
            N /= p[i];
            total += multiplier;
            multiplier *= p[i];
                                             // total para
        ans *= total;
                                             // este
           factor primo
    if (N != 1) ans \star= (N+1); // N^2-1/N-1 = N+1
    return ans;
```

# 4.3 Prueba de primalidad

```
ll sieve size;
bitset<10000010> bs;
vl p;
void sieve(ll upperbound) {
    _sieve_size = upperbound+1;
    bs.set();
    bs[0] = bs[1] = 0;
    for (ll i = 2; i < _sieve_size; ++i) if (bs[i]) {</pre>
        for (ll j = i*i; j < _sieve_size; j += i) bs[j] =</pre>
        p.push_back(i);
bool isPrime(ll N) {
    if (N < _sieve_size) return bs[N]; // O(1) primos</pre>
    for (int i = 0; i < (int)p.size() && p[i]*p[i] <= N;
       ++i)
        if (N%p[i] == 0)
            return false;
    return true;
                        // 0 ( sqrt(N) / log(sqrt(N)) )
       para N > 10^7
   //Nata: solo se garantiza para N <= (ultimo primo de
   p)^2 = 9.99 * 10^13
```

#### 4.4 Criba Modificada

```
//Criba modificada
Si hay que determinar el numero de factores primos para
   muchos (o un rango) de enteros.
La mejor solucion es el algoritmo de criba modificada O(N
    log log N)
int numDiffPFarr[MAX N+10] = \{0\}; // e.g., MAX N = 10^7
for (int i = 2; i <= MAX_N; ++i)</pre>
    if (numDiffPFarr[i] == 0) // i is a prime number
        for (int j = i; j <= MAX_N; j += i)
            ++numDiffPFarr[j]; // j is a multiple of i
//Similar para EulerPhi
int EulerPhi[MAX N+10];
for (int i = 1; i <= MAX N; ++i) EulerPhi[i] = i;</pre>
for (int i = 2; i <= MAX N; ++i)</pre>
    if (EulerPhi[i] == i) // i is a prime number
        for (int j = i; j <= MAX_N; j += i)
            EulerPhi[j] = (EulerPhi[j]/i) * (i-1);
```

#### 4.5 Funcion Totient de Euler

//EulerPhi(N): contar el numero de enteros positivos < N
 que son primos relativos a N.</pre>

```
//El vector p es el que genera la criba de eratostenes
//Phi(N) = N * productoria(1 - (1/pi))
ll EulerPhi(ll N) {
    11 ans = N; // Empezar con ans = N
    for (int i = 0; (i < (int)p.size()) && (p[i]*p[i] <=</pre>
       N): ++i) {
        if (N%p[i] == 0) ans -= ans/p[i]; //contar
            factores
        while (N%p[i] == 0) N /= p[i];
                                           //primos unicos
    if (N != 1) ans -= ans/N; // ultimo factor
    return ans:
```

# 4.6 Exponenciacion binaria

```
ll binpow(ll b, ll n, ll m) {
    b %= m;
    11 \text{ res} = 1;
    while (n > 0) {
        if (n & 1)
             res = res * b % m;
        b = b * b % m;
        n >>= 1;
    return res % m;
```

## 4.7 Fibonacci Matriz

```
def mult (matriz1, matriz2):
    fila 1 = [matriz1[0][0] * matriz2[0][0] + matriz1
       [0][1] * matriz2[1][0], matriz1[0][0] * matriz2
       [0][1] + matriz1[0][1] * matriz2[1][1]]
    fila 2 = [matriz1[1][0] * matriz2[0][0] + matriz1
       [1][1] * matriz2[1][0], matriz1[1][0] * matriz2
       [0][1] + matriz1[1][1] * matriz2[1][1]]
    return [fila 1, fila 2]
def mult_vector(matriz, vector):
    a = matriz[0][0] * vector[0] + matriz[0][1] * vector
    b = matriz[1][0] * vector[0] + matriz[1][1] * vector
       [1]
    return [a, b]
def modulos(matriz, n):
   matriz[0][0] %= n
   matriz[0][1] %= n
   matriz[1][0] %= n
   matriz[1][1] %= n
   return matriz
def exp_bin(b, n, m):
```

```
res = [[1, 0], [0, 1]]
    while n > 0:
        if n % 2 == 1:
            res = mult(modulos(res, m), modulos(b, m))
        b = mult (modulos (b, m), modulos (b, m))
        n //= 2
    return modulos(res, m)
matriz = [[1, 1], [1, 0]]
vector = [1, 0]
\# a = list(map(int, input().split()))
m = exp bin(matriz, int(input()), 1000000007)
v = mult_vector(m, vector)
print(v[1] % 1000000007)
```

## 4.8 GCD y LCM

```
//O(\log 10 \ n) \ n == \max(a, b)
int qcd(int a, int b) { return b == 0 ? a : qcd(b, a%b);
int lcm(int a, int b) { return a / gcd(a, b) * b; }
//gcd(a, b, c) = gcd(a, gcd(b, c))
```

## Geometria

## 5.1 Puntos

```
// Punto entero
struct point{
                   11 x, y;
                    point(ll x, ll y): x(x), y(y) {}
};
// Punto flotante
struct point{
                    double x, y;
                    point (double _x, double _y): x(_x), y(_y) {}
                    bool operator == (point other) const{
                                       return (fabs(x-other.x) < EPS) && (fabs(y-other.y) <
                                                        EPS);
                    };
};
// Distancia entre dos puntos
double dist(point p1, point p2) {
                    return sqrt((p1.x-p2.x)*(p1.x-p2.x)+(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p1.y-p2.y)*(p
                                     -p2.v));
// Rotacion de un punto
point rotate(point p, double theta) {
                    // rotar por theta grados respecto al origen (0,0
                    double rad = theta*(M_PI/180);
```

#### 5.2 Lineas

```
// Linea de flotantes de la forma ax+by+c=0
struct line{double a,b,c;};
// Creacion de linea con dos puntos
// b=1 para lineas no verticales y b =0 para verticales
void pointsToLine(point p1, point p2, line& 1) {
    if (fabs(p1.x-p2.x)<EPS) {
        l.a=1.0; l.b=0.0; l.c=-p1.x;
    }else{
        l.a= -double(p1.y-p2.y)/(p1.x-p2.x);
        1.b = 1.0;
        1.c= -double(1.a*p1.x)-p1.y;
// Comprobacion de lineas paralelas
bool areParallel(line 11, line 12) {
    return (fabs(11.a-12.a) < EPS) && (fabs(11.b-12.b) < EPS)
// Comprobacion de lineas iquales
bool areSame(line 11, line 12) {
    return areParallel(11,12) && (fabs(11.c-12.c) <EPS);</pre>
// Disntacia de un punto a una linea
double distPointToLineaEq(line 1, point p) {
    return fabs(1.a*p.x + 1.b*p.y + 1.c)/sqrt(1.a*1.a+1.b
       *1.b);
bool areIntersect(line 11, line 12, point& p) {
    if (areParallel(11,12)) return false;
    // resolver sistema 2x2
    p.x = (12.b*11.c - 11.b*12.c)/(12.a*11.b - 11.a*12.b)
    // CS: comprobar linea vertical -> div por cero
    if (fabs(11.b) > EPS) p.y = -(11.a * p.x + 11.c);
    else p.v = -(12.a*p.x + 12.c);
    return true;
```

#### 5.3 Vectores

```
// Creacion de un vector
struct vec{
```

```
double x, v:
    vec(double x, double y) : x(x), y(y) {}
};
// Puntos a vector
vec toVec(point a, point b) {
    return vec(b.x-a.x , b.y-a.y);
// Escalar un vector
vec scale(vec v, double s){
    // s no negatico:
    // <1 mas corto
    // 1 iqual
    // >1 mas largo
    return vec(v.x*s,v.v*s);
// Trasladar p segun v
point traslate(point p, vec v) {
    return point(p.x+v.x , p.y+v.y);
// Producto Punto
double dot(vec a, vec b) {
    return (a.x*b.x + a.y*b.y);
// Cuadrado de la norma
double norm sq(vec v) {
    return v.x*v.x + v.y*v.y;
// Angulo formado por aob
double angle (point a, point o, point b) {
    vec oa = toVec(o, a);
    vec ob = toVec(o,b);
    return acos(dot(oa,ob)/sqrt(norm sq(oa)*norm sq(ob)))
// Producto cruz
double cross(vec a, vec b) {
    return (a.x*b.y) - (a.y*b.x);
// Lado respecto una linea pg
bool ccw(point p,point q,point r) {
    // Devuelve verdadero si el punto r esta a la
       izquierda de la linea pq
    return cross(toVec(p,q),toVec(p,r))>0;
// Colinear
bool collinear(point p, point q, point r) {
    return fabs(cross(toVec(p,q), toVec(p,r))) < EPS;
```

# 5.4 Poligonos

```
// Crear un poligono
// la idea es crearlo con algun orden ya sea horario o
   anti-horario
// v debe cerrarse
vector<point> Poligono;
// Perimetro de un poligono
double perimeter(const vector<point>& P) {
    double result =0.0;
    for (int i =0;i<(int)P.size()-1;i++)result+= dist(P[i</pre>
       ],P[i+1]);
    return result;
// Area de un poligono
double area(const vector<point>& P) {
    // la mitad del determinante
    double result = 0.0, x1, y1, x2, y2;
    for (int i =0; i<(int)P.size()-1; i++) {
        x1 = P[i].x;
        x2 = P[i+1].x;
        y1 = P[i].y;
        y2 = P[i+1].y;
        result += (x1*y2 - x2*y1);
    return fabs(result/2.0);
// Comprobacion de si es Convexto un poligono
bool isConvex(const vector<point>& P) {
    int sz = (int)P.size();
    if (sz<=3) return false;</pre>
    bool isLeft = ccw(P[0], P[1], P[2]);
    for (int i =1;i<sz-1;i++)</pre>
        if (ccw(P[i],P[i+1],P[(i+2)==sz ? 1:i+2])!=isLeft
            return false;
    return true;
// Comprobar si un punto esta dentro de un poligono
bool inPoligono(point pt, const vector<point>& P) {
    // P puede ser concavo/convexo
    if ((int)P.size()==0) return false;
    double sum =0;
    for (int i =0;i<(int)P.size()-1;i++){</pre>
        if (ccw(pt,P[i],P[i+1]))
            sum += angle(P[i],pt,P[i+1]); // izquierda/
                anti-horario
        else sum -= angle(P[i],pt,P[i+1]);// derecha/
           horario
    return fabs(fabs(sum)-2*M_PI)<EPS;</pre>
```

#### 5.5 Convex Hull

```
struct pt{
    double x, y;
    pt (double x, double y): x(x), y(y) {}
int orientation(pt a, pt b, pt c) {
    double v = a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y);
    if (v < 0) return -1; // horario</pre>
    if (v > 0) return +1; // anti-horario
    return 0;
bool cw(pt a, pt b, pt c, bool include_collinear) {
    int o = orientation(a, b, c);
    return o < 0 || (include collinear && o == 0);</pre>
bool collinear(pt a, pt b, pt c) { return orientation(a,
   b, c) == 0;
void convex hull(vector<pt>& a, bool include collinear =
    pt p0 = *min_element(a.begin(), a.end(), [](pt a, pt
       b) {
        return make pair(a.y, a.x) < make pair(b.y, b.x);
    });
    sort(a.begin(), a.end(), [&p0](const pt& a, const pt&
        int o = orientation(p0, a, b);
        if (0 == 0)
            return (p0.x-a.x) * (p0.x-a.x) + (p0.y-a.y) * (p0
                 < (p0.x-b.x) * (p0.x-b.x) + (p0.y-b.y) * (p0.y-b.y)
                    y-b.y);
        return o < 0;
    });
    if (include collinear)
        int i = (int)a.size()-1;
        while (i \ge 0 \&\& collinear(p0, a[i], a.back())) i
        reverse(a.begin()+i+1, a.end());
    vector<pt> st;
    for (int i = 0; i < (int)a.size(); i++) {</pre>
        while (st.size() > 1 && !cw(st[st.size()-2], st.
            back(), a[i], include_collinear))
            st.pop back();
        st.push back(a[i]);
    a = st;
int main() {
```

```
ios::sync_with_stdio(false);
cin.tie(0);

ll n; cin>>n;
vector<pt> Puntos;
for (int i =0;i<n;i++) {
    double x,y;cin>>x>>y;
    pt punto(x,y);
    Puntos.push_back(punto);
```

```
}
convex_hull(Puntos, true);
cout<<Puntos.size()<<ln;
for (pt punto:Puntos){
    cout<<(ll)punto.x<<" "<<(ll)punto.y<<ln;
}
}</pre>
```