The Library of GPGPU-Powered Sparse Boolean Linear Algebra Operations

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Abstract—Sparse matrices are widely applicable in data analysis while the theory of matrix processing is well-established. There are a wide range of algorithms for basic operations such as matrix-matrix and matrix-vector multiplication, factorization, etc. To facilitate data analysis, GraphBLAS API provides a set of building blocks and allows for reducing algorithms to sparse linear algebra operations. While GPGPU utilization for highperformance linear algebra is common, the high complexity of GPGPU programming makes the implementation of GraphBLAS API on GPGPU challenging. In this work, we present a GPGPU library of sparse operations for an important case — Boolean algebra. The library is based on modern algorithms for sparse matrix processing. We provide a Python wrapper for the library to simplify its use in applied solutions. Our evaluation shows that operations specialized for Boolean matrices can be up to 2 times faster and consume up to YYY times less memory than generic operations from modern libraries. We hope that our results help to move the development of a GPGPU version of GraphBLAS API forward.

Index Terms—sparse linear algebra, GPGPU, boolean semiring, sparse boolean matrix

I. INTRODUCTION

One technique to efficiently solve a data analysis problem is to formulate it in terms of operations over vectors and matrices (in terms of linear algebra). This way it is possible to employ a set of reliable mathematical tools and solutions. Another advantage of this approach is the ability to evaluate the problem with zero-cost by high-performance linear algebra libraries, which utilize modern hardware, provide various optimization techniques, and allow one to prototype a solution in code with predefined building blocks quickly and safely. GraphBLAS API¹ [1] is one of the standards that introduce such building blocks. GraphBLAS takes into account the sparsity of data by using sparse formats of matrices and vectors, and generalizes the building blocks by operating with arbitrary monoids and semirings. While initially GraphBLAS was focused on graph analysis, it was shown that the proposed approach can be successfully used for data analysis in other areas, such as computational biology [2] and machine learning [3].

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¹GraphBLAS project web page: https://graphblas.github.io/. Access date: 19.01.2021.

GPGPU utilization for data analysis and for linear algebra operations is a promising way to high-performance data analysis because GPGPU is much more powerful in parallel data processing. Unfortunately, GPGPU programming is very challenging. It introduces heterogeneous device model into the system, memory traffic, and data operations limitations, as well as requires taking into account vendor-specific capabilities. Best to our knowledge, there is no complete implementation of GraphBLAS API on GPGPU, except for the GraphBLAST project² [4], which is currently in active development.

The sparsity of data introduces issues with load balancing, irregular data access, thus sparsity complicates the implementation of high-performance algorithms for sparse linear algebra on GPGPU even more. As a result, there is a huge number of different formats for sparse matrices and vectors representation, such as CSR, COO, Quad-tree, and a huge number of algorithms for operations over these formats. See [5] for a significant survey of sparse matrix-matrix multiplication algorithms. Algorithms for different operations, such as matrix-matrix multiplication and matrix-vector multiplication are developed independently. Thus, there are no sparse linear algebra libraries based on the state-of-the-art algorithms. Moreover, existing libraries, such as cuSparse³, clSparse⁴ [6], or more modern CUSP⁵ or bhSparse⁶ [7], are focused on numerical computations over floats or doubles, not on generic data processing over arbitrary semirings which is required for GraphBLAS API implementation.

An important partial case of linear algebra is the sparse Boolean linear algebra. Boolean algebra is suitable for problems over a finite set of values, such as transitive closure of a relation or a graph, regular and context-free path queries for graphs [8], as well as parsing for different classes of

²GraphBLAST project: https://github.com/gunrock/graphblast. Access date: 19.01.2021.

³NVIDIA sparse matrix library (in Cuda) https://docs.nvidia.com/cuda/cusparse/. Access date: 19.01.2021.

⁴Sparse linear library functions in OpenCL: http://clmathlibraries.github.io/ clSPARSE/. Access date: 19.01.2021.

⁵CUSP sparse linear algebra library: https://cusplibrary.github.io/modules.html. Access date: 19.01.2021.

⁶bhSparse sparse matrix multiplication library: https://github.com/weifengliu-ssslab/bhSPARSE. Access date: 19.01.2021.

languages, such as Context-Free [9], Boolean and Conjunctive [10], Multiple Context-Free(MCFL) [11]. Moreover, some operations over the Boolean semiring can be used as building blocks for algorithms over other semirings. For example, to compute the shape of the result of the operation. Thus, sparse Boolean linear algebra is an important partial case both as a way to solve applied problems and as a building block for other algorithms. However, sparse Boolean linear algebra on GPGPU is still not presented, because of its high specificity.

In this paper, we present the implementation of sparse boolean linear algebra operations as a stand-alone selfsufficient programming library for the two most popular GPGPU platforms: NVIDIA Cuda⁷ and OpenCL⁸. Cuda is a GPGPU technology for NVIDIA devices which employs some platform-specific facilities, such as unified memory mechanism, and make architectural assumptions which gives more optimizations space at the cost of portability. OpenCL is a platform-agnostic API standard, which allows for running computations on different platforms, such as multi-threaded CPUs, GPUs, and FPGAs. Our implementation relies on modern techniques of sparse matrices processing and exploits some optimizations, related to the boolean data processing. Moreover, we provide a Python API to simplify utilization of our library. Preliminary evaluation shows that such operation as matrix-matrix multiplication specialized for Boolean matrices can be up to 2 times faster and consume up to YYY times less memory in comparison with generic operations from such libraries as CUSP or CuSparse.

II. ABOUT LIBRARIES

Implemented sparse boolean linear algebra libraries for OpenCL and NVIDIA Cuda platforms are called *clBool*⁹ and *cuBool*¹⁰ respectively. Projects are hosted at GitHub platform. The source code is licensed under MIT license. The build process is straightforward: it is configured with CMake tool and requires extra setup only of platform-specific development kits.

Libraries architecture is briefly depicted at figure 1. The core of the libraries is written in the C++ programming languages, which is well-suited for performance and resource critical computational tasks. Actual GPU related logic is presented in platform specific backends: Cuda and OpenCL, which use respective technologies for resources and GPU executable code management. cuBool library exposes C compatible API, what gives expressiveness and allows to embed that API into other execution environments by interoperability mechanisms. Pycubool module encapsulates such functionality and provides it for high-level Python runtime.

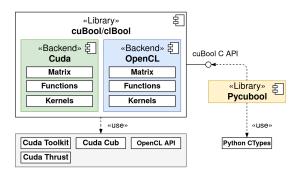


Fig. 1. Conceptual sparse boolean linear algebra library architecture

It is worth mention, that it is convenient to create the single library with common interface and several backends for different execution targets. At this time clBool and cuBool are distinct libraries, but they can be integrated into single library. This integration is something to be done in near future. This process requires careful selection of the interface to allow the end user properly configure the library for specific tasks, as well as provide the option to automatically select a specific implementation depending on the capabilities of the target device.

Libraries operate on boolean semiring with values set $\{true, false\}$ with false as a neutral element, '+' operation defined as logical or and '*' defined as logical and. Values are also denoted as $\{1, 0\}$ respectively, and the abbreviation nnz(M) gives the number of non-zero cells of the matrix M.

Main primitive is sparse matrix of boolean values, stored in one of the sparse formats. Sparse vector primitive is not presented, since its utilization is relatively rare presented in practical computational tasks. But its support is something to be added in far future. Primary available operations and functions are following.

- Create sparse matrix M of size $m \times n$.
- \bullet Delete sparse matrix M and release all its internal resources.
- Fill the matrix M with values $L = \{(i,j)_k\}_k$. The result of this operation is $M_{i,j} = 1$ for each $(i,j) \in L$, and $M_{i,j} = 0$ for the rest of matrix values.
- Read matrix M values $L = \{(i, j) \mid M_{i,j} = 1\}.$
- Matrix-matrix multiply-add operation $C += M \times N$.
- Matrix-matrix add operation M += N.
- Matrix-matrix Kronecker product $K = M \otimes N$.

III. IMPLEMENTATION DETAILS

In this section we discuss the particular implementation details of the proposed libraries. Although general and architectural specifics are similar, the actual internal storage formats and algorithms are different. With this development strategy we address the potential problem of processing the sparse data with different values distribution, as well as the problem of proper balancing between time of the execution and memory consumption.

⁷CUDA is a platform and programming model for NVIDIA devices. Home page: https://developer.nvidia.com/CUDA-zone. Access date: 19.01.2021.

⁸OpenCL is an open standard for parallel programming of heterogeneous systems. Home page: https://www.khronos.org/opencl/. Access date: 19.01.2021.

⁹clBool project: https://replace/me/with/actual/url. Access date: 03.02.2021.
¹⁰cuBool project: https://github.com/JetBrains-Research/cuBool. Access date: 03.02.2021.

A. cuBool

cuBool is sparse boolean linear algebra implementation specifically for NVIDIA Cuda platform. Core of this library relies on Cuda C/C++ language and API, what with NVCC compiler allows intermix C++ with Cuda specifics. Also cuBool employs NVIDIA Thrust auxiliary library, which provides implementation for generic data containers and operations, such as *iterating*, *exclusive or inclusive scan*, *map* and etc., which are executed on Cuda device. That allows express algorithms in terms of high-level optimised primitives, what increases code readability and reduces time for development.

Sparse matrix is stored in the *compressed sparse row* (CSR) format with only two arrays: rowspt for row offset indices and cols for columns indices. Boolean matrices has no actual values, thus 1 values are encoded only as (i,j) pairs. It allows to store matrix M of size $m \times n$ in $(m+nnz(M)) \times sizeof(index_t)$ bytes of GPU memory, where $index_t$ is type of stored indices, for simplicity can be selected as $uint32_t$.

The algorithm Nspasrse [12] is used for matrix-matrix multiplication. This algorithm is a boolean values case adaptation of the state-of-the-art, efficient and memory saving sparse general matrix multiplication (SpGEMM) algorithm, proposed in Yusuke Nagasaka et al. research [13]. This algorithm was selected because it gives promising relatively small memory footprint for large matrices processing, as well as it competes with other major Cuda SpGEMM implementations, such as cuSPARSE or CUSP.

Matrix-matrix addition is based on GPU Merge Path algorithm [14] with dynamic work balancing and two pass processing. These optimizations give better workload dispatch among execution blocks and allow more precise memory allocations in order to keep memory footprint small respectively.

B. clBool

clBool is sparse boolean linear algebra implementation for OpenCL platform. This library is implemented in the C++ with OpenCL kernels, stored as separate source files, loaded on demand at runtime.

Sparse matrix primitive is stored in *coordinate format* (COO) with two arrays: rows and cols for row and column indices of the stored non-zero values. For the matrix M of size $m \times n$ memory consumption is $2 \times nnz(M) \times sizeof(index_t)$. This format was selected instead of CSR, because COO gives better memory footprint for very sparse matrices with a lot of empty rows.

Matrix-matrix multiplication implementation is based on the algorithm, proposed in Weifeng Liu et al. research [15]. It is multi-step algorithm with dynamic workload balancing, which operates on CSR matrices. Since clBool primary primitive is COO matrix, before actual matrix-matrix multiplication the input matrices are converted into *doubly compressed sparse row* (DCSR) format, described in A. Buluc et al. work [16]. This algorithm is suitable for OpenCL implementation, what is confirmed with its utilisation in clSPARSE library.

Matrix-matrix addition is based on GPU Merge Path algorithm as well. Since all COO matrix values are stored in the

TABLE I MATRIX DATA

| Matrix M | #rows | Nnz of M | Nnz of M ² | Nnz of $M + M^2$ |
|----------------|---------|------------|-----------------------|------------------|
| wing | 62032 | 243088 | 714200 | 917178 |
| luxembourg_osm | 114599 | 239332 | 393261 | 632185 |
| roadNet-PA | 1090920 | 3083796 | 7238920 | 9931528 |
| roadNet-TX | 1393383 | 3843320 | 8903897 | 12264987 |
| belgium_osm | 1441295 | 3099940 | 5323073 | 8408599 |
| roadNet-CA | 1971281 | 5533214 | 12908450 | 17743342 |

continuous manner, its merge can be completed at single time, compared to CSR matrix merge computed on a per row basis. This operation is implemented in a classic one pass fashion: it allocates single merge buffer of size nnz(A) + nnz(B)) before actual merge of matrices A and B, what can negatively affect memory consumption for large matrices with lots of duplicated non-zero values at the same positions.

IV. EVALUATION

We evaluate the applicability of the proposed libraries for analysis of some real-world matrix data. The experiments are designed as a computational tasks, that arise as stand-alone or intermediate steps in the solving of practical problems.

For evaluation, we used a PC with Ubuntu 20.04 installed. It has Intel core i7-6700 CPU, 3.4GHz, DDR4 64Gb RAM and Geforce 1070Ti GPU with 8Gb VRAM. We only measure the execution time of the operations themselves. The actual data is assumed to be loaded into the VRAM or RAM respectively in the appropriate format, required for the target tested framework. Time to load data from disc and prepare initial matrices state is excluded from the time measurements.

We use four sparse matrix libraries, CUSP, cuSPARSE, clSPARSE for GPU and SuiteSparse for CPU. CUSP provides template based implementation for operations, however it does not provide extra optimizations especially for boolean case values. cuSPARSE and clSPARSE both provide operations only for general types, such as float or double. But this limitation can be ignored, if we consider non-zero float values as *true* one. SuiteSparse is a GraphBLAS API reference implementation for CPU, which allows to use build-in boolean semiring.

For performance evaluations, we selected N various square matrices which are widely used for sparse matrices benchmarks from the Sparse Matrix Collection at University of Florida¹¹. The name and size of the matrix data are summarized in the table I.

The results of the evaluation are summarized in the tables below. Time is measured in seconds unless specified otherwise. The result for each experiment is average over 20 runs. The cell is left blank if the time limit is exceeded, or if there is not enough memory to allocate the data and internal auxiliary structures.

The first experiment is intended to measure performance of the matrix-matrix multiplication as $M \times M$. The results

¹¹T. Davis. The SuiteSparse Matrix Collection (the University of Florida Sparse Matrix Collection). Home page: https://sparse.tamu.edu/. Access date: 23.01.2021.

TABLE II

MATRIX-MATRIX MULTIPLICATION EVALUATION RESULTS, TIME IS

MEASURED IN SECONDS.

| Matrix M | CuBool | CUSP | CuSprs | ClBool | ClSprs | SuiteSprs |
|----------------|--------|-------|--------|--------|--------|-----------|
| wing | 0.003 | 0.007 | 0.021 | 0.007 | - | 0.007 |
| luxembourg_osm | 0.005 | 0.005 | 0.002 | 0.010 | - | 0.003 |
| roadNet-PA | 0.023 | 0.043 | 0.037 | 0.047 | - | 0.067 |
| roadNet-TX | 0.030 | 0.053 | 0.047 | 0.057 | - | 0.084 |
| belgium_osm | 0.027 | 0.034 | 0.030 | 0.054 | - | 0.061 |
| roadNet-CA | 0.039 | 0.076 | 0.071 | 0.078 | - | 0.121 |

TABLE III
ELEMENT-WISE MATRIX-MATRIX ADDITION, TIME IS MEASURED IN SECONDS.

| Matrix M | CuBool | CUSP | CuSprs | ClBool | ClSprs | SuiteSprs |
|----------------|--------|-------|--------|--------|--------|-----------|
| wing | 0.001 | 0.002 | 0.003 | 0.006 | - | 0.003 |
| luxembourg_osm | 0.002 | 0.002 | 0.001 | 0.006 | - | 0.002 |
| roadNet-PA | 0.015 | 0.011 | 0.013 | 0.073 | - | 0.035 |
| roadNet-TX | 0.019 | 0.014 | 0.015 | 0.076 | - | 0.045 |
| belgium_osm | 0.019 | 0.010 | 0.011 | 0.064 | - | 0.027 |
| roadNet-CA | 0.027 | 0.019 | 0.020 | 0.141 | - | 0.062 |

are presented in the table II. We can see that for a relatively small matrices, the results of all libraries are comparable. However, cuBool shows generally better performance among competitors. clBool gives good performance, constantly better than SuiteSparse, and comparable to cuSPRASE or CUSP in some cases. However, there is still space for optimizations, so it requires an in deep investigation in our further research.

The second experiment is intended to measure performance of the element-wise matrix-matrix addition as $M+M^2$, where evaluation of matrix M^2 is excluded from measurements. The results are presented in the table III. The numbers obtained in this experiment are less unambiguous than in the previous experiment. Although libraries are still comparable for small matrices, results vary greatly for large matrices. CUSP shows nearly best performance among almost all experiments. cuS-PRASE is also comparable with it in this aspect. Libraries cuBool and clBool lag behind insignificantly, and generally keep their results within acceptable limits. It is worth mention, that CUSP matrix-matrix addition implementation has significant memory consumption, what can negatively affect on processing of huge data. cuSPARSE performance can degraded in such case as well, since its implementation is based on hashing, what is very sensitive for out-of shared blocks memory access for large data processing.

V. CONCLUSION

In this work we present a library for sparse Boolean linear algebra which implements such basic operations, as matrix-matrix multiplication and element-wise matrix-matrix addition in both, Cuda C and OpenCL C.

The first direction of the future work is to integrate all parts (OpenCL and Cuda backends) to a single library and improve its documentation and prepare to publish. Also, it is necessary to publish a Python package.

Another important step is to evaluate the library on different algorithms and devices. Namely, algorithms for RPQ and CFPQ should be implemented and evaluated on related data sets. Also, it is necessary to evaluate OpenCL version on FPGA which may require additional technical effort and code changes.

Finally, we plan to discuss with GraphBLAS community possible ways to use our library as a backend for GraphBLAST or SuiteSparse in case of Boolean computations. Moreover, it may be possible to use implemented algorithms as a base for generalization to arbitrary semirings.

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