MATH50003 Numerical Analysis

II.2 Reals

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Part II

Representing Numbers

- 1. Integers via modular arithmetic
- 2. Reals via floating point
- 3. Floating point arithmetic and bounding errors
- 4. Interval arithmetic for rigorous computations

Ariane 5 rocket explosion

Learn floating point, or else...



II.2.1 Real numbers in binary

We can represent any real number using binary digits

Definition 6 (real binary format). For $b_1, b_2, \ldots \in \{0, 1\}$, Denote a non-negative real number in binary format by:

$$(B_p \dots B_0.b_1b_2b_3\dots)_2 := (B_p \dots B_0)_2 + \sum_{k=1}^{\infty} \frac{b_k}{2^k}.$$

Example 11 (rational in binary)

II.2.2 Floating-point numbers

How do we represent an uncountable set with only p-bits?

Bit Format: $s q_{Q-1} \dots q_0 b_1 \dots b_s$

Definition 7 (floating-point numbers). Given integers σ (the exponential shift), Q (the number of exponent bits) and S (the precision), define the set of Floating-point numbers by dividing into normal, sub-normal, and special number subsets:

$$F_{\sigma,Q,S} := F_{\sigma,Q,S}^{\text{normal}} \cup F_{\sigma,Q,S}^{\text{sub}} \cup F^{\text{special}}.$$

How do bits dictate whether its normal/sub/special?

Look at exponent. 3 examples:

0 10000 1010000000

1 00000 1100000000

1 11111 0000000000

II.2.3 IEEE float-point numbers

What exponent shift/number of bits/precision is used in practice?

$$F_{16} := F_{15,5,10}$$

$$F_{32} := F_{127,8,23}$$

$$F_{64} := F_{1023,11,52}$$

$$F_{\sigma,Q,S}^{\text{normal}} := \{ \pm 2^{q-\sigma} \times (1.b_1b_2b_3 \dots b_S)_2 : 1 \le q < 2^Q - 1 \}.$$

Example 12 (interpreting 16-bits as a float). Consider the number with bits

0 10000 1010000000

Half-precision $F \cdot - F$

$$F_{16} := F_{15,5,10}$$

Example 13 (rational to 16-bits). How is the number 1/3 stored in F_{16} ?

II.2.4 Sub-normal and special numbers

Sub-normal have exponent bits all 0, special have all 1

$$F_{\sigma,Q,S}^{\text{sub}} := \{ \pm 2^{1-\sigma} \times (0.b_1b_2b_3 \dots b_S)_2 \}.$$

Example 14 (subnormal in 16-bits). Consider the number with bits

1 00000 1100000000

$$F^{\mathrm{special}} := \{\infty, -\infty, \mathrm{NaN}\}$$

Example 15 (special in 16-bits). The number with bits

1 11111 0000000000

On the other hand, the number with bits

1 11111 000000001

Time for code.