

Give regular expressions for each of the following languages over the alphabet $\{0, 1\}$.

1 All strings containing the substring 000 .

Solution: $(0 + 1)^*000(0 + 1)^*$

2 All strings *not* containing the substring 000 .

Solution: $(1 + 01 + 001)^*(\varepsilon + 0 + 00)$

Solution: $(\varepsilon + 0 + 00)(1(\varepsilon + 0 + 00))^*$

3 All strings in which every run of 0 s has length at least 3.

Solution: $(1 + 0000^*)^*$

Solution: $(\varepsilon + 1)((\varepsilon + 0000^*)1)^*(\varepsilon + 0000^*)$

4 All strings in which 1 does not appear after a substring 000 .

Solution: $(1 + 01 + 001)^*0^*$

5 All strings containing at least three 0 s.

Solution: $(0 + 1)^*0(0 + 1)^*0(0 + 1)^*0(0 + 1)^*$

Solution: $1^*01^*01^*0(0 + 1)^*$ or $(0 + 1)^*01^*01^*01^*$

6 Every string except 000 . (**Hint:** Don't try to be clever.)

Solution: Every string $w \neq 000$ satisfies one of three conditions: Either $|w| < 3$, or $|w| = 3$ and $w \neq 000$, or $|w| > 3$. The first two cases include only a finite number of strings, so we just list them explicitly. The last case includes *all* strings of length at least 4.

$$\begin{aligned} & \varepsilon + 0 + 1 + 00 + 01 + 10 + 11 \\ & + 001 + 010 + 011 + 100 + 101 + 110 + 111 \\ & + (1 + 0)(1 + 0)(1 + 0)(1 + 0)(1 + 0)^* \end{aligned}$$

Solution: $\varepsilon + 0 + 00 + (1 + 01 + 001 + 000(1 + 0))(1 + 0)^*$

7 All strings w such that *in every prefix of w* , the number of 0 s and 1 s differ by at most 1.

Solution: Equivalently, strings that alternate between 0 s and 1 s: $(01 + 10)^*(\varepsilon + 0 + 1)$

8 (Hard.) All strings containing at least two 0 s and at least one 1 .

Solution: There are three possibilities for how such a string can begin:

- Start with 00 , then any number of 0 s, then 1 , then anything.
- Start with 01 , then any number of 1 s, then 0 , then anything.
- Start with 1 , then a substring with exactly two 0 s, then anything.

All together: $000^*1(0 + 1)^* + 011^*0(0 + 1)^* + 11^*01^*0(0 + 1)^*$

Or equivalently: $(000^*1 + 011^*0 + 11^*01^*0)(0 + 1)^*$

Solution:

There are three possibilities for how the three required symbols are ordered:

- Contains a **1** before two **0**s: $(0 + 1)^* 1 (0 + 1)^* 0 (0 + 1)^* 0 (0 + 1)^*$
- Contains a **1** between two **0**s: $(0 + 1)^* 0 (0 + 1)^* 1 (0 + 1)^* 0 (0 + 1)^*$
- Contains a **1** after two **0**s: $(0 + 1)^* 0 (0 + 1)^* 0 (0 + 1)^* 1 (0 + 1)^*$

So putting these cases together, we get the following:

$$\begin{aligned} & (0 + 1)^* 1 (0 + 1)^* 0 (0 + 1)^* 0 (0 + 1)^* \\ & + (0 + 1)^* 0 (0 + 1)^* 1 (0 + 1)^* 0 (0 + 1)^* \\ & + (0 + 1)^* 0 (0 + 1)^* 0 (0 + 1)^* 1 (0 + 1)^* \end{aligned}$$

Solution: $(0 + 1)^* (101^* 0 + 010\ 011^* 0 + 01^* 01) (0 + 1)^*$

9 (Hard.) All strings w such that *in every prefix of w* , the number of **0**s and **1**s differ by at most 2.

Solution: $(0(01)^* 1 + 1(10)^* 0)^* \cdot (\varepsilon + 0(01)^*(0 + \varepsilon) + 1(10)^*(1 + \varepsilon))$

10 (Really hard.) All strings in which the substring **000** appears an even number of times.
(For example, **0001000** and **0000** are in this language, but **00000** is not.)

Solution: Every string in $\{0, 1\}^*$ alternates between (possibly empty) blocks of **0**s and individual **1**s; that is, $\{0, 1\}^* = (0^* 1)^* 0^*$. Trivially, every **000** substring is contained in some block of **0**s. Our strategy is to consider which blocks of **0**s contain an even or odd number of **000** substrings.

Let X denote the set of all strings in **0** * with an even number of **000** substrings. We easily observe that $X = \{0^n \mid n = 1 \text{ or } n \text{ is even}\} = 0 + (00)^*$.

Let Y denote the set of all strings in **0** * with an *odd* number of **000** substrings. We easily observe that $Y = \{0^n \mid n > 1 \text{ and } n \text{ is odd}\} = 000(00)^*$.

We immediately have $0^* = X + Y$ and therefore $\{0, 1\}^* = ((X + Y)1)^*(X + Y)$.

Finally, let L denote the set of all strings in $\{0, 1\}^*$ with an even number of **000** substrings. A string $w \in \{0, 1\}^*$ is in L if and only if an odd number of blocks of **0**s in w are in Y ; the remaining blocks of **0**s are all in X .

$$L = ((X1)^* Y1 \cdot (X1)^* Y1)^* (X1)^* X$$

Plugging in the expressions for X and Y gives us the following regular expression for L :

$$\left(((0 + (00)^*)1)^* \cdot 000(00)^* 1 \cdot ((0 + (00)^*)1)^* \cdot 000(00)^* 1 \right)^* \cdot ((0 + (00)^*)1)^* \cdot (0 + (00)^*)$$

Whew!