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Hidden parallelepiped in Hawk

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Abstract

Some abstract here

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Thank you to some people

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Chapter 1

Introduction

1.1 Context and motivation

Digital signatures are an important part of secure communication today. The most used cryptographic scheme used for digital signatures today is DSA (Digital Signature Algorithm) or RSA (Rivest, Shamir, and Adleman) signatures (source). However, in 1994, Peter Shor developed Shor's algorithm, which, given a large enough quantum computer, is able to solve the hard problems DSA and RSA is based upon, namely the Discrete Logarithm Problem and Prime Factorization(source). Whether big enough quantum computers will emerge any time soon is debatable. However, measures against this potential looming threat has already begun. In 2016, NIST (National Institute of Standards and Technology) announced a standardization process for new standard schemes for KEMs (Key Encapsulation Methods) and digital signatures that have strong security against quantum computers (source). Many of the submissions to this process, including KRYSTALS-Dilithium which is to be standardized, are based on lattice problems that are believed to be hard to solve for both classical and quantum computers (source).

Cryptographic schemes based on lattice problems are not an entirely new phenomenon, however. NTRU-Sign, published in 2003(source), is a digital signature scheme based on the hardness of the Closest Vector Problem (source). The original scheme was broken due to Phong. Q. Nguyen & Oded Regev in 2006 [3], who showed that by observing enough signatures generated with one secret key, one can retrieve the secret key. A newer

digital signature scheme, Hawk (source), submitted to NIST’s standardization process, is a scheme similar that of NTRU and GGH. The goal of this thesis is to try and adapt the Hidden Parallelepiped Problem attack to Hawk [2]. Hawk spec here [1]

1.2 Objectives

The objective for this thesis consists of two main parts:

- **Implementation of Hawk in Rust.** As the first part of the thesis I implement the Hawk digital signature scheme in the Rust programming language. Implementing a scheme on ones own is a good way to actually learn how it works. I chose to implement it in Rust for the sake of learning the programming language. Moreover, having ones own version makes it easier to experiment, adjust and modify to ones need. It would also be a challenge to understand and work with complicated source code someone else has written.
- **Cryptanalysis and experimentation.** As part two of the thesis I want to do cryptanalysis of Hawk. The goal is to use the "Learning a parallelepiped" attack [3] and adjusting it to try and break Hawk. This requires both theoretical and practical work, and experiments will be implemented in Rust.

1.3 Thesis outline

Chapter 2 will introduce important notions and mathematical background used in this thesis. Chapter 3 will introduce Hawk and the *Learning a Parallelepiped* attack and implementations of these. In Chapter 4, a version HPP attack aimed at Hawk will be presented. Chapter 5 will show results, and Chapter 6 will discuss future work.

Chapter 2

Background

2.1 Asymmetric cryptography

2.1.1 Digital signatures

2.1.2 Hash-Then-Sign

2.1.3 GGH

2.2 Linear algebra & lattices

2.2.1 Lattices

2.3 Probability theory

2.3.1 Distributions

Chapter 3

Hawk and *Learning a parallelepiped*

3.1 Hawk

In the following Hawk, the digital signature scheme, will be presented, as in [1]

3.1.1 Simple Hawk

Present simple sketch of keygen, sign and ver, as well as an example in dimension 2

3.1.2 Key generation

3.1.3 Signature generation

3.1.4 Signature verification

3.2 Learning a parallelepiped

The paper *Learning a Parallelepiped: Cryptanalysis of GGH and NTRU Signatures* by Phong Q. Nguyen and Oded Regev from 2006 introduced a method for breaking digital signature schemes based on the GGH scheme [3]. Essentially, by observing enough *message, signature* pairs generated by a secret key one can deduce this secret key. The attack broke the NTRU-Sign scheme by observing as little as 400 signatures.

3.2.1 Solving the Hidden Parallelepiped Problem

First we define an idealized version of both the problem to solve and the solution as proposed in [3]. Then we discuss the problem and solution in a practical context.

Hidden Parallelepiped Problem (HPP): Let $V = [\mathbf{v}_1, \dots, \mathbf{v}_n]$ be a secret $n \times n$ matrix and $V \in \mathcal{GL}_n(\mathbb{R})$. Define by $\mathcal{P}(V) = \{\sum_{i=1}^n x_i \mathbf{v}_i : x_i \in [-1, 1]\}$ a secret n -dimensional parallelepiped defined by V . Given a polynomial (in n) number of vector samples uniformly distributed over $\mathcal{P}(V)$, recover the rows \mathbf{v}_i of V .

We solve this problem in two main steps:

1. Transforming the hidden parallelepiped into a hidden hypercube:
2. Learning the hypercube:

Parallelepiped \rightarrow Hypercube The first step of the attack is to convert our hidden parallelepiped $\mathcal{P}(V)$ into a hidden hypercube $\mathcal{P}(C)$ with orthogonal basis vectors. Essentially, one moves each sampled point in accordance to a transformation matrix L computed the following way:

- Approximate $G \approx V^T V$ using our samples $\mathcal{X} = \{\mathbf{x}_1, \dots, \mathbf{x}_u\}$
- Compute L such that $LL^T = G^{-1}$
- Then $C = VL$
- By multiplying our samples \mathcal{X} to the left by L , our samples are now uniformly distributed over the hidden hypercube $\mathcal{P}(C)$

Learning a Hypercube ABCDEFG

3.2.2 Covariance matrix

3.2.3 Hidden parallelepiped to hidden hypercube

3.2.4 Gradient descent

3.2.5 Example in dimension 2

3.2.6 Hawk resistance against HPP

Chapter 4

Implementation

Introduction to the implementation part of the thesis

4.1 Implementation of Hawk

Something something about the implementation of Hawk in Rust. Mentions of sampling, integer/float types, speed and comparison to Hawk team's C code?

4.2 Implementation of HPP

Something something about the implementation of HPP in Rust. Something about speedup/parallelization of gradient descent?

Chapter 5

Adapting HPP to Hawk

In this chapter we investigate the steps needed to possibly apply the Hidden Parallelepiped Problem to the Hawk digital signature scheme.

5.1 Covariance matrix secret key in Hawk

Nothing yet I'm afraid

5.2 Secret Key Recovery

Since \mathbf{x} follows some distribution close to some normal distribution, we hope that enough vectors \mathbf{w} will disclose some information about \mathbf{B}^{-1} . If we know \mathbf{B}^{-1} we know \mathbf{B} . This is the goal.

Bibliography

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URL: <https://hawk-sign.info/>.
- [2] Léo Ducas, Eamonn W. Postlethwaite, Ludo N. Pulles, and Wessel van Woerden. Hawk: Module LIP makes lattice signatures fast, compact and simple. Cryptology ePrint Archive, Paper 2022/1155, 2022.
URL: <https://eprint.iacr.org/2022/1155>.
- [3] Phong Q. Nguyen and Oded Regev. Learning a parallelepiped: Cryptanalysis of ggh and ntru signatures, 2009.

Appendix A

Generated code

Listing A.1: Source code of something

```
1 println!("Goodbye World");
```