

COMPUTER NETWORKS FINAL EXAMINATION-2020

- · Eisha Tir Raazia
- · 17k-3730
- · Sec C

QUESTION # 01

17K-3730

L = (0 mod 10) +1 = 1 Mbits.

P1 = 4+1 # = 5 Mbps.

 $R_2 = 7+1 = 8 \text{ Mbps}$ $R_3 = 4+1 = 5 \text{ Mbps}$

 $d_1 = 1 + 1 = 2 \times km$ $d_2 = 7 + 1 = 8 \times km$ $d_3 = 3 + 1 = 4 \times km$

dproc = 0+1 = 1 Msec.

dqueving (pkt1) = 1 + 1=2 MSec = TQ,

dqueving (pkt2) = 7 + 1=8 MSec = TQ2

dqueving (pkt3) = 3 + 1=4 MSec = TQ3

Parameters (pkt4) = 7 + 1-8 MSec = TQ3

3 = 2.5 × 108 m/s.

Transmission delay & $T_{A-X} = L = 0.2$ or $\frac{1}{5}$ sec R_1 $T_{X-Y} = L = 0.125$ or $\frac{1}{8}$ sec R_2

Ty-B = L = 0.2 or 1/5 sec.

To = Transmission Delay(A-B) = 0.2 +0.2 +0.125 = 0.525 sec

Propagation delay &

Tprop =
$$\frac{d}{3} = \frac{d}{1} + \frac{d}{2} + \frac{d}{3} = \frac{(2+8+4)\times10^3}{8.5\times10^8}$$

 $T_{prop} = 5.6 \times 10^{-5} \text{s}$

$$D_{i} = PDelay (pk+1) = dprec + TO_{i} + 700 + 700 + 700 + 700 + 10000 + 1000 + 1000 + 1000 + 1000 + 1000 + 1000 + 1000 + 1000 + 1000$$

D. = Ppelay (Pk+2) = dptoc + TD2 + TD + Tprop.

Send to end) = 10x10⁻⁶ + 8x10⁻⁶ + 0.525 * + 5.6x10⁻⁵

= 0.5250 sec

Dy= Poetry (PKt4) = aproc + TO3 + To + Tprop

endtoend = 1×10-6 + 10.8 ×10-6 + 0.525 + 5.6 ×10-5

= 0.525 sec

Nodal delay = D, + D2 + D3 + D4 = 2.1 sec.

(A)

- (i) HTTP is used-
- (ii) Bachan's client use request message.
- (iii°) Messages are:

 1st Action required (put)

 and Action required (delete)

 3rd Action required (put)
- (iv) Put www.bachan.com/image1.jpg.

 delete www.bachan.com/image1.jpg.

 put www.bachan.com/image2.jpg.
- (V) HTTP version 1.1 as put and delete isn't available in HTTP version 1.0.

(B)

- 1) First, client start communication with server
 - · Sewer waits to be connected.
 - · After the connection is established, the program starts running using inler process communication.
 - 2) It's client server model.
 - 3) IP addless and port no identification req.
 - 4) Sockets used for communication.

	ar	Typest hop
255-255-255-128-	19 168-65-128	13.62-13.5
September System		Who P
	124 13/13/1	MXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXXX
Mark State of the		16×600-25-3
9/0	17.500	1926168-64-194

(4)

(A)

Network	Next	Interface.	
10.100.16.0/22	-	ml	
192-168-65-128/25	1 -1	mo	
10-100-22-3 /24		m3	
192-168-69-194/26	A STATE	m2	
ofher	-	m2	
	192-168-65-128/25 192-168-65-128/25 192-168-65-194/26	10.100.16.0/22 - 192.168.65.128/25 - 10.100.22.3/24 - 192.168.65.194/26 -	

- (B) It will route to interface "m3"
- (C) It will route to interface m3"
- (D) It will route to interface "m2"

- (A) Minimum time required is more than los, los for router 2 to start and few muser for each neighbours. It's adjacent to computing the routing table.
- (B) R1 and R2, path b/c it's common for all the router traffic, this will utilize even R3 16 R3, R4, R4 If we send pkd to R2.
- (C) For link R, and R2 to stop utilize the other routers, R, can cost 7 so that rest of routers will not use this path in this it's own routers, it will take cost of 7 because it will increase cost for all the routers.

		(1	A)					
N	D(A),P(A)	B	C	D	E	F	9	
2,	2,2	5,2	4,2	00	00	∞	00	
ZA		A,P	00	9,4	∞	14,A	∞	
2AB			5,B	8, B	7,B	00	∞0	
ZABC				00	9,0	∞	00	
ZABCE				11, E		00	(B,E	
ZABCED						∞	16,0	
ZABCEDG	THICK AND ADDRESS.			The state of the s			16,6	

Z-A least cost is 2.
Z-G least cost is 16

80

Dijsktra computed least cost path by considering the only cost edges which has minimum weight -

Date:

QUESTION # 09

(A)

There are 3 techniques for it &

- a) Parity checking

 In it, single bit parity detection and

 Do bit parity is also used
 - b) Sheck-sum method

 Sused at fromsport layer and detect
 enor by adding all 16bit nos, 2 taking
 the complement.
- c) Cyclic redundancy check

 -> used at the data link layer and

 it's more powerful ever detection

 coding-
 - I with check summing.

171 000-	SLIN
17k-3730	JAN JAN



Date:-

(B)

 $G = 2^3 + 2^2 + 1 = 1101$ G - 1 = 9 = 7 + 1r= 3

databits = 1911 19th Nov.

= 01110 1110 1110 ==

011101000

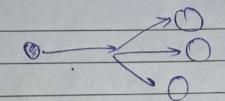
101001000

Databits = 011101111 with G=101



(A)

(i) When address is not in switch's forwarding table, the frame is treated like broadcast frame and sent to all hosts on network also known as "flooding"



A
1
149

(III) No, because witch at only maintains

9 table that tells which Mac addlers

of host is reachable via which

port so, port doesn't need seperate

address.

(B)

Host A

IP: 176.35.196.2

mac: ee:2c:bd:8f:1b:a5

A seds ARP request to B when it knows

That B's IP is 176.36.100.2, the MAC
address that it get in return will be

routers MAC address 70:00:d6:cb:cq:e8

As, the router will forward the pakets to

remote host, knowing the destination IP

of host B (176.36.100.2)

A creates link layer frame with RIs MAC

address as destination, frame contain A to

B IP datagram