

CS2102 Database Systems
2014/2015 Semester I
Tutorial #5 Normal Forms

1. Consider the relation R (A, B, C, D, E) with the set of FDs
 $F = \{ \{A\} \rightarrow \{B\}, \{BC\} \rightarrow \{E\}, \{ED\} \rightarrow \{A\} \}$.
 - a. List all the keys of R. **CDE, ACD, BCD**
 - b. Is R in 3NF? **Yes, because B, E, A are all parts of keys.**
 - c. Is R in BCNF? **No, because none of A, BC, ED contain a key**

2. Consider the relation R(A, B, C, D, E) with the set of FDs $F = \{A \rightarrow B, C \rightarrow D\}$.
 - a. If R in 3NF? If not, decompose it into a collection of 3NF relations.
Key is ACE. R is not in 3NF. Decompose to AB, CD, ACE

 - b. If R in BCNF? If not, decompose it into a collection of BCNF relations.
No. BCNF decomposition: AB, CD, ACE

3. Suppose we have the following instance of a relation S with three attributes ABC.

A	B	C
1	2	3
4	2	3
5	3	3

Which of the following dependencies can you infer does not hold over relation S?

- a. $A \rightarrow B$
- b. **$BC \rightarrow A$ does not hold. Look at tuples $\langle 1, 2, 3 \rangle$ and $\langle 4, 2, 3 \rangle$**
- c. $B \rightarrow C$

Can you identify any dependencies that hold over relation S?

No, given just an instance of S, we can say that certain dependencies (e.g., $A \rightarrow B$ and $B \rightarrow C$) are not violated by this instance, but we cannot say that these dependencies hold with respect to S. To say that an FD holds w.r.t a relation is to make a statement about ALL allowable instances of that relation.