CS2102 Database Systems

SCHEMA REFINEMENT: NORMAL FORMS

Normal Forms

- A normal form restricts the set of data dependencies that are allowed to hold on a schema to avoid certain undesirable redundancy and update problems in the database.

Normal Forms

- There are several normal forms, each providing guidance on good schema designs
- We focus on two normal forms that are based on FDs:
  - Boyce-Codd Normal Form (BCNF)
  - Third Normal Form (3NF)
- Definitions of BCNF and 3NF assume that each FD is of the form $X \rightarrow A$ where A is a single attribute.

Boyce-Codd Normal Form (BCNF)

- A relation schema R (with FDs F) is in Boyce-Codd normal form if for every non-trivial FD $X \rightarrow A$ in F, X is a superkey.

- A non-trivial FD $X \rightarrow A$ that holds on R is said to violate BCNF if X is not a superkey of R

Example

- Consider the MovieList schema with FDs
  F = { title $\rightarrow$ director, address $\rightarrow$ phone, {address, time} $\rightarrow$ title }
- Recall that the only key is {address, time}
- FDs in F that violate BCNF are
  - title $\rightarrow$ director
  - address $\rightarrow$ phone
- Thus, MovieList is not in BCNF

Decomposition into BCNF

- Let $X \rightarrow A$ be an FD in F that causes violation of BCNF
- Decompose R into

$$R_1 = XA$$

$$R_2 = R - A$$

- If R_1 or R_2 is not in BCNF, then decompose them further as described.

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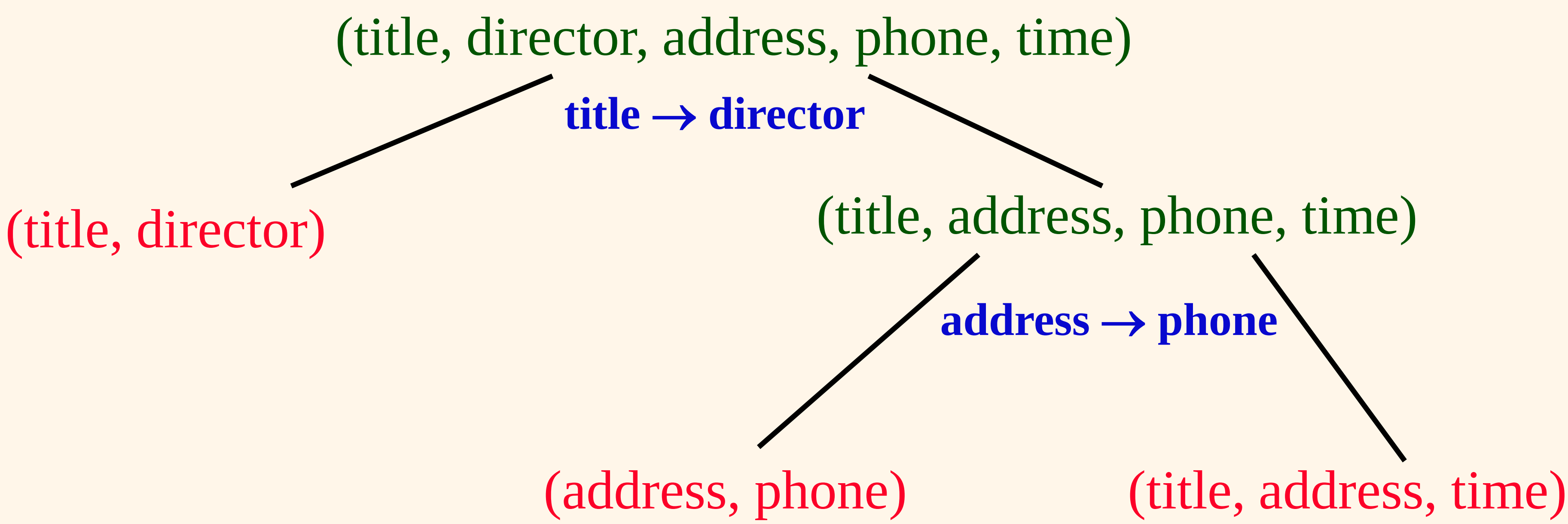


Decomposition into BCNF

- Decomposition $\{R_1, R_2, \ldots, R_n\}$ is in BCNF if each R_i is in BCNF (w.r.t. F_{Ri})
- BCNF decompositions are lossless join decomposition
- But, not all schema has a dependency-preserving BCNF decomposition

Example

- Consider R (course, prof, time) with FDs
  F = { course $\rightarrow$ prof, {prof, time} $\rightarrow$ course }
- Keys are {course, time} and {prof, time}
- R is not in BCNF because course is not a superkey
- Decomposition into R_1 (course, prof) and R_2 (course, time) is a lossless join but does not preserve the FD {prof, time} $\rightarrow$ course

Third Normal Form (3NF)

- 3NF is a less restrictive normal form that always guarantees a lossless join decomposition that preserves dependencies.
- A relation schema R (with FDs F) is in third normal form if for every non-trivial FD $X \rightarrow A$ in F (where A is a single attribute), either X is a superkey or A is a prime attribute.
- A non-trivial FD $X \rightarrow A$ that holds on R is said to violate 3NF if X is not a superkey of R and A is a nonprime attribute
- R in BCNF $\Rightarrow$ R in 3NF

Example

- Consider again R (course, prof, time) with FDs { course $\rightarrow$ prof, {prof, time} $\rightarrow$ course }
- Keys are {course, time} and {prof, time}
- R is in 3NF because both prof and course are prime attributes

Instance of R

prof	time	course
Codd	Tue 3pm	DB101
Codd	Thur 9am	DB101
Gray	Tue 4pm	CS323
Gray	Fri 10am	IT201

Decomposition into 3NF

- Synthesis Approach
- Input: Schema R with FDs F which is a minimal cover
- Output: A dependency preserving, lossless join 3NF decomposition of R

Decomposition into 3NF (cont'd)

- Initialize $D = \phi$
- Apply union rule to combine FDs with same LHS into a single FD.
  - Let $F = \{f_1, f_2, \ldots, f_n\}$ be the resultant set of FDs
- For each f_i of the form $X_i \rightarrow A_i$ do
  - Create a relation schema $R_i (X_i, A_i)$ for FD f_i
  - Insert the schema R_i into D
- Choose a key K of R and insert a relation schema $R_{n+1}(K)$ into D
- Remove redundant schema from D
  - Delete R_i from D if $R_i \subseteq R_j$ where $R_j \in D$
- Return D

set of relations in third normal form

Example

- Consider R(A, B, C, D, E) with FDs F = {ABCD → E, E → D, A → B, AC → D} key is AC
- A minimal cover of F is {AC → E, E → D, A → B}
- Only key is AC these 2 fds violate 3nf
- R is not in 3NF because A → B violates 3NF (A is not a superkey and B is not a prime attribute)
- 3NF decomposition of R
  - Create a schema for each FD: R_1 (A, C, E), R_2 (E, D), R_3 (A, B)
  - Create a schema for a key of R: R_4 (A, C)
  - Remove redundant schema: R_4 is redundant because $R_4 \subseteq R_1$
  - 3NF decomposition is R_1 (A, C, E), R_2 (E, D), R_3 (A, B)

Remarks on 3NF Decomposition

- A decomposition $\{R_1, R_2, \ldots, R_n\}$ is in 3NF if each R_i is in 3NF (w.r.t. F_{Ri})
- The 3NF decomposition produced by synthesis approach may not be unique
  - Choice of minimal cover
  - Choice of redundant relation schema being removed

BCNF vs. 3NF

- **BCNF** is lossless join (may not be dependency preserving
- **3NF** is lossless join and dependency preserving
- Recall R(course, prof, time) with FDs { course $\rightarrow$ prof, {prof, time} $\rightarrow$ course }
  - Keys are {course, time} and {prof, time}
  - R is in 3NF but not in BCNF
  - BCNF decomposition { R_1(course, prof), R_2(course, time) } is lossless but not dependency preserving

Another Example

- Consider schema Contract (contractid, supplierid, projectid, deptid, partid, qty, value)
- CSJDPQV for short
- Contract C is an agreement that supplier S will supply Q items of part P to project J associated with department D; value of this contract is V
  - Contract id C is a key: $C \rightarrow CSJDPQV$
  - A project purchase a part using a single contract: $JP \rightarrow C$
  - A department purchase at most one part from a supplier: $SD \rightarrow P$
  - Each project deals with a single supplier: $J \rightarrow S$

Example – BCNF Decomposition

- FDs F = { C $\rightarrow$ CSJDPQV, JP $\rightarrow$ C, SD $\rightarrow$ P, J $\rightarrow$ S }
- From JP $\rightarrow$ C, C $\rightarrow$ CSJDPQV and transitivity, we have JP $\rightarrow$ CSJDPQV JP is key too
- SD $\rightarrow$ P violates BCNF since SD is not a key, decompose CSJDPQV into CSJDQV and SDP
- From J $\rightarrow$ S, decompose CSJDQV into JS and CJDQV
- Decomposition is lossless
- Decomposition does not preserve FD JP $\rightarrow$ C
  - Need to join the two relations to check that the FD is not violated.
  - Can add a relation CJP to the decomposition if CJP is in BCNF

Example – 3NF Synthesis

- FDs F = { $C \rightarrow CSJDPQV$, $JP \rightarrow C$, $SD \rightarrow P$, $J \rightarrow S$ }
- F is not a minimal cover.
  - Replace $C \rightarrow CSJDPQV$ with { $C \rightarrow S$, $C \rightarrow J$, $C \rightarrow D$, $C \rightarrow P$, $C \rightarrow Q$, $C \rightarrow V$ }
  - Remove $C \rightarrow P$ from F since it is implied by $C \rightarrow S$, $C \rightarrow D$ and $SD \rightarrow P$
  - Remove $C \rightarrow S$ from F since it is implied by $C \rightarrow J$ and $J \rightarrow S$
- Minimal cover F' = { $C \rightarrow J$, $C \rightarrow D$, $C \rightarrow Q$, $C \rightarrow V$,

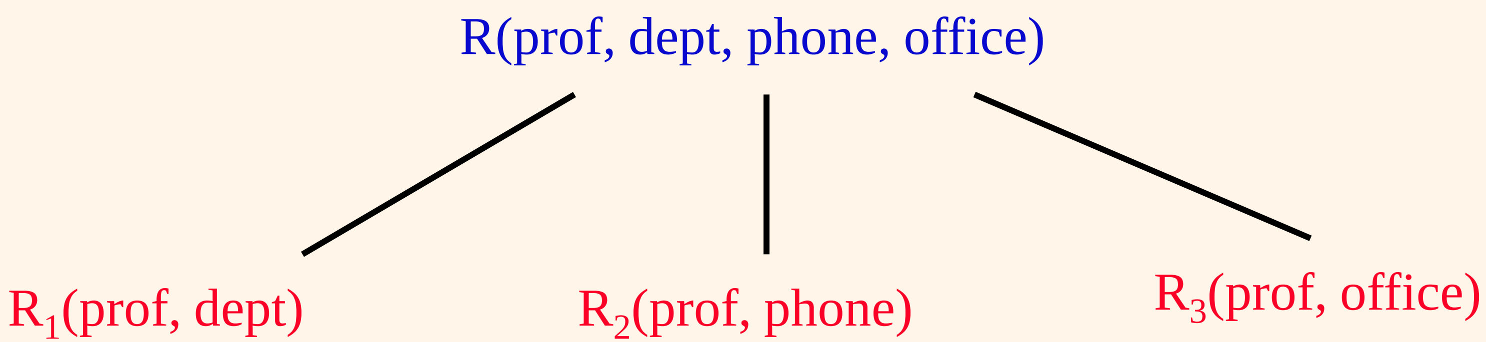
  $JP \rightarrow C$, $SD \rightarrow P$, $J \rightarrow S$ }

Example – 3NF Synthesis

- Minimal cover F′ = {$C \rightarrow J$, $C \rightarrow D$, $C \rightarrow Q$, $C \rightarrow V$,
  $JP \rightarrow C$, $SD \rightarrow P$, $J \rightarrow S$ }
- Combine FDs with same LHS F′ = {$C \rightarrow JDQV$,
  $JP \rightarrow C$, $SD \rightarrow P$, $J \rightarrow S$ }
- Create relations CJDQV, CJP, SDP, JS
- Remark: You can combine relations with C as key
  - e.g., CJDQV and CJP to CJDQVP

Remarks on Decomposition

- Decomposition is a last resort to solve problems of redundancy and anomalies
- Too much decomposition can be harmful
- Example: R(prof, dept, phone, office) with FD { prof $\rightarrow$ dept,phone,office }



- Consider de-normalization for performance reasons