

# CS2102 Database Systems

Semester 1 2019/2020

## Tutorial 07 (*Selected Answers*)

### Quiz

Questions 1-6 uses the following schema  $R(A, B, C, D)$  with set of FDs  $F = \{AB \rightarrow C, B \rightarrow A, C \rightarrow D\}$

- Which of the following FDs are *logically implied* by  $F$ ?
  - $A \rightarrow D$
  - $AB \rightarrow D$
  - $B \rightarrow C$
  - $B \rightarrow D$
  - $B \rightarrow CD$
  - $C \rightarrow A$
- Which of the following set of FDs are equivalent to  $F$ ?
  - $G_1 = \{B \rightarrow A, B \rightarrow D, C \rightarrow D\}$
  - $G_2 = \{B \rightarrow A, C \rightarrow D, B \rightarrow D, B \rightarrow C\}$
  - $G_3 = \{B \rightarrow AC, C \rightarrow D\}$
- Compute the attribute closure of  $B$  w.r.t.  $F$
- What is/are the key(s) of  $R$  w.r.t.  $F$ ?
- What are the prime attributes of  $R$  w.r.t.  $F$ ?
- Compute one minimal cover of  $F$

### Tutorial Questions

[Discussion: 7(a), 7(b), 8(a), 8(b), 9(ab), 10, 11]

- The more *extended* set of Armstrong's axioms has union and decomposition. Proof them only using Armstrong's axioms.
  - Union if  $a \rightarrow b$  and  $a \rightarrow c$ , then  $a \rightarrow bc$
  - Decomposition if  $a \rightarrow bc$ , then  $a \rightarrow b$

**Solution:** [the proof-assisted version]

- Proof  
 $A \rightarrow B$  [Given]  
 $A \rightarrow C$  [Given]  
 $A \rightarrow AB$  [Augmentation (1) with A]  
 $AB \rightarrow BC$  [Augmentation (2) with B]  
 $A \rightarrow BC$  [Transitivity (3) and (4)]
- Proof  
 $A \rightarrow BC$  [Given]  
 $BC \rightarrow B$  [Reflexivity]  
 $A \rightarrow B$  [Transitivity (1) and (2)]

- The more *extended* set of Armstrong's axioms has two more rules in addition to union and decomposition. For each of the rules below, proof them using only Armstrong's axiom.
  - Pseudo-transitivity if  $a \rightarrow b$  and  $bc \rightarrow d$ , then  $ac \rightarrow d$
  - Composition if  $a \rightarrow b$  and  $c \rightarrow d$ , then  $ac \rightarrow bd$

**Solution:** [the proof-assisted version]

- Proof  
 $A \rightarrow B$  [Given]  
 $BC \rightarrow D$  [Given]

$AC \rightarrow BC$  [Augmentation (1) with C]

$AC \rightarrow D$  [Transitivity (3) and (2)]

b) Proof

$A \rightarrow B$  [Given]

$C \rightarrow D$  [Given]

$AC \rightarrow BC$  [Augmentation (1) with C]

$BC \rightarrow BD$  [Augmentation (2) with B]

$AC \rightarrow BD$  [Transitivity (3) and (4)]

9. Consider  $R(A, B, C, D, E, G)$  with FDs  $F = \{ABC \rightarrow E, BD \rightarrow A, CG \rightarrow B\}$ .

a) Use extended Armstrong's axioms to show that  $F \models CDG \rightarrow E$

b) Compute  $CDG^+$

c) Find all the keys of  $R$

**Solution:** [the proof-assisted version is given in the appendix]

a) Proof

$ABC \rightarrow E$  [Given]

$BD \rightarrow A$  [Given]

$CG \rightarrow B$  [Given]

$CDG \rightarrow BD$  [Augmentation (3) with D]

$CDG \rightarrow A$  [Transitivity (4) and (2)]

$CDG \rightarrow CG$  [Reflexivity]

$CDG \rightarrow B$  [Transitivity (6) and (3)]

$CDG \rightarrow BC$  [Augmentation (7) with C]

$CDG \rightarrow ABC$  [Union (5,8)]

$CDG \rightarrow E$  [Transitivity (9) and (1)]

b)  $CDG^+ = ABCDEG = R$

c) From (b), we know that  $CDG$  is a superkey of  $R$ . Since  $DG^+ = DG$ ;  $CG^+ = CGB$ ; and  $CD^+ = CD$ , we can conclude that  $CDG$  is a key of  $R$ .

Since each of  $C$ ,  $D$ , or  $G$  does not appear on the RHS of any FD in  $F$ , it follows that

$C \notin (R - C)^+$ ;  $D \notin (R - D)^+$ ; and  $G \notin (R - G)^+$ .

Therefore, any key of  $R$  must be a superset of  $CDG$ . Since  $CDG$  is already a key of  $R$ ,  $CDG$  is the only key of  $R$ .

10. Consider the schema  $R(A, B, C, D, E)$  with FDs  $F = \{AB \rightarrow CDE, AC \rightarrow BDE, B \rightarrow C, C \rightarrow B, C \rightarrow D, B \rightarrow E\}$ . Find all the prime attributes of  $R$ . Show your working.

**Solution:**

Observe that attribute  $A$  does not appear in any RHS of any FDs in  $F$ . Therefore,  $A$  must appear in every key. By process of enumeration of all subset of  $R$  containing  $A$ , starting from the simplest:

- $A^+ = A$   $A$  is NOT a superkey
- $AB^+ = ABCDE$   $AB$  is a superkey
  - $B^+ = BCDE$   $B$  is NOT a superkey  $\Rightarrow AB$  is a key
- $AC^+ = ABCDE$   $AC$  is a superkey
  - $C^+ = BCDE$   $C$  is NOT a superkey  $\Rightarrow AC$  is a key

- $AD^+ = AD$   $AD$  is NOT a superkey
- $AE^+ = AE$   $AE$  is NOT a superkey
- $ADE^+ = ADE$   $ADE$  is NOT a superkey

The two keys of  $R$  are  $AB$  and  $AC$ . Hence, the prime attributes are  $A$ ,  $B$ , and  $C$ .

11. Consider the schema  $R(A, B, C, D, E)$  with FDs  $F = \{AB \rightarrow CDE, AC \rightarrow BDE, B \rightarrow C, C \rightarrow B, C \rightarrow D, B \rightarrow E\}$ . Find one minimal cover of  $F$ . Show your working.

**Solution:**

Start with  $G = F = \{AB \rightarrow CDE, AC \rightarrow BDE, B \rightarrow C, C \rightarrow B, C \rightarrow D, B \rightarrow E\}$

In this example, we don't start working from  $a \rightarrow A$  to simplify the notation.  $a \rightarrow A$  can always be obtained by decomposition rule of extended Armstrong's axioms.

Remove redundant attributes:

- Consider  $AB \rightarrow CDE$ 
  - $A^+ = A$  So  $B$  is not a redundant attribute in FD
  - $B^+ = BCDE$  So  $A$  is a redundant attribute in FD
    - Let  $G = \{B \rightarrow CDE, AC \rightarrow BDE, B \rightarrow C, C \rightarrow B, C \rightarrow D, B \rightarrow E\}$
- Consider  $AC \rightarrow BDE$ 
  - $A^+ = A$  So  $C$  is not a redundant attribute in FD
  - $C^+ = BCDE$  So  $A$  is a redundant attribute in FD
    - Let  $G = \{B \rightarrow CDE, C \rightarrow BDE, B \rightarrow C, C \rightarrow B, C \rightarrow D, B \rightarrow E\}$
- No other FD need to be considered since all LHS are a single attribute now

Remove redundant FDs

- Let  $G = \{B \rightarrow C, B \rightarrow D, B \rightarrow E, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$ 
  - after applying decomposition rule and removing duplicates
- Consider  $B \rightarrow C$ 
  - $B^+ = BDE$  w.r.t.  $\{B \rightarrow D, B \rightarrow E, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$ 
    - $B \rightarrow C$  is not redundant
  - Thus,  $G = \{B \rightarrow C, B \rightarrow D, B \rightarrow E, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$
- Consider  $B \rightarrow D$ 
  - $B^+ = BCDE$  w.r.t.  $\{B \rightarrow C, B \rightarrow E, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$ 
    - $B \rightarrow D$  is redundant
  - Thus,  $G = \{B \rightarrow C, B \rightarrow E, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$
- Consider  $B \rightarrow E$ 
  - $B^+ = BCDE$  w.r.t.  $\{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$ 
    - $B \rightarrow E$  is redundant
  - Thus,  $G = \{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$
- Consider  $C \rightarrow B$ 
  - $C^+ = CDE$  w.r.t.  $\{B \rightarrow C, C \rightarrow D, C \rightarrow E\}$ 
    - $C \rightarrow B$  is not redundant
  - Thus,  $G = \{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$
- Consider  $C \rightarrow D$ 
  - $C^+ = BCE$  w.r.t.  $\{B \rightarrow C, C \rightarrow B, C \rightarrow E\}$ 
    - $C \rightarrow D$  is not redundant
  - Thus,  $G = \{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$
- Consider  $C \rightarrow E$

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- $C^+ = BCD$  w.r.t.  $\{B \rightarrow C, C \rightarrow B, C \rightarrow D\}$ 
  - $C \rightarrow E$  is not redundant
- Thus,  $G = \{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$

Thus,  $G = \{B \rightarrow C, C \rightarrow B, C \rightarrow D, C \rightarrow E\}$  is a minimal cover of  $F$