Data Structures 1 - HW4 - Abraham Murciano
1) Given an array of size n, the way to implement 2 stacks would be to have I stack starting at position 0, going upwards, using one variable to store the index of the 'top' element of that stack. Then the second stack would begin at position n-1 and elements would be pushed onto the previous elements of the array. Again there would be a priable to store the index of the top element of that stack. (see illustration telow)
array 4: 42 A3 B4 88 82 81  a=2 b=6
Stack At > Stack B
2) class stack {    queue Q;    int size;    void push (int x) { // complexity of push is $\Theta(size)$ Q. enqueue( $x$ );    for (int i=0; i < size; i++)-{       Q. enqueue (Q. dequeue());    }
S)ze++;
int pop() { //complexity of pop 15 B(1)  if (is Empty()) if (is Empty); cannot pop";  throw "Stack is empty; cannot pop";  }  size;  return Q. dequeue; }
bool is Empty ) { //complexity of is Empty is O(1) return size == 0; }
3) a) roid to (axwell 01, queue 02) { int Q1 size = Q1. size();  for (int i=0; i < Q1 size; i++) {  int temp = Q1. dequeve();  if (second) (Q2, temp)) {  Q1. enqueve(temp);  }
bool search (queue Q, int x) {     for (int i=0; i < Q.size(); i++) {         if (x == Q.dequeue()) {             return true;         }     }
return false;
b) complexity is O(nxm) because complexity of search is O(m), and the bop in too rans n times, each time calling seach, so O(nxm)