## Mathematical Logic HW2

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Throught this document I will use T to represent a tautology, and F to represent a contradiction.

1. (b)  $x \to x$  is a tautology, since  $\to$  is only false when the left operand is true but the right operand is false. Since in  $x \to x$  both operands are the same, it cannot occur that the left operand be true and the right one be false. Therefore  $(x \to x) \to x = T \to x$ . However,  $T \to x$  is only true when x is true, so  $(x \to x) \to x = x$ . Therefore:

$$(((x \to x) \to x) \to x) \to x$$

$$= (x \to x) \to x$$

$$= x$$

which is neither a tautology or a contradiction.

(d) As the truth table below shows, when A is false and B and C are true, the proposition is false, otherwise it is true. Therefore this proposition is neither a tautology nor a contradiction.

A	B	C	$((A \vee \neg B) \to (A \vee (B \wedge C))) \to (C \to A)$
F	F	F	T
F	F	T	T
F	Τ	F	T
F	Τ	Т	F
Т	F	F	T
Τ	F	Т	T
Т	Τ	F	T
Т	Τ	Т	T

- 2. (a) "Yesterday, the sun rose in the west" is simply an atomic proposition. It is not a contradiction, even though we know that factually it is incorrect.
  - (b) "It's rainy only if it's cloudy." is neither a tautology nor a contradiction, since although we know that in the world we live in we need clouds for there to be rain, regarding mathematical logic, we can in theory assign a truth value to "It's rainy" and a false value to "it's cloudy", thus giving our proposition a value of falsehood in that case.
  - (c) "If my name is Ron, then my name is not Ron." is neither a tautology nor a contradiction. Let R be the atomic proposition "my name is Ron". The former proposition is equivalent to  $R \to \neg R$ . If R takes the value 'True', then  $\neg R$  takes the value 'False', in which case  $R \to \neg R = T \to F = F$ , proving that the proposition in question is not a tautology. However, if R takes the value 'False', then  $R \to \neg R = F \to T$ , which evaluates to 'True', so the proposition is also not a contradiction.
- 3. (b)

$$(A \to B) \land (C \to B) \Leftrightarrow (A \lor C) \to B$$

Proof:

$$\begin{split} (A \to B) \wedge (C \to B) \\ = (\neg (A \wedge \neg B)) \wedge (\neg (C \wedge \neg B)) \\ = (\neg A \vee B) \wedge (\neg C \vee B) \\ = (\neg A \wedge \neg C) \vee B \\ = \neg (A \vee C) \vee B \\ = \neg ((A \vee C) \wedge \neg B) \\ = (A \vee C) \to B \end{split}$$

4. (b) Let  $\alpha =$  "a is larger than b",  $\beta =$  "a is larger than 0",  $\gamma =$  "b is larger than 0",  $\delta =$  "b is equal to 0". We want to determine if

$$(\alpha \wedge (\gamma \vee \delta)) \to \beta \iff (\alpha \wedge \neg \beta) \to (\neg \gamma \wedge \neg \delta)$$

$$(\alpha \wedge (\gamma \vee \delta)) \to \beta$$

$$= \neg (\alpha \wedge (\gamma \vee \delta) \wedge \neg \beta)$$

$$= \neg \alpha \vee \neg (\gamma \vee \delta) \vee \beta$$

$$= \neg \alpha \vee \beta \vee (\neg \gamma \wedge \neg \delta)$$

$$= \neg (\alpha \wedge \neg \beta) \vee (\neg \gamma \wedge \neg \delta)$$

$$= \neg ((\alpha \wedge \neg \beta) \wedge \neg (\neg \gamma \wedge \neg \delta))$$

$$= (\alpha \wedge \neg \beta) \to (\neg \gamma \wedge \neg \delta)$$

5. (b)

$$\begin{aligned} & (X \vee Y) \wedge (X \vee \neg Y) \\ = & X \vee (Y \wedge \neg Y) \\ = & X \vee F \\ = & X \end{aligned}$$

(d)

$$\begin{split} &(X \wedge Y \wedge Z) \vee (X \wedge Y \wedge \neg Z) \vee (X \wedge \neg Y) \\ = &X \wedge ((Y \wedge Z) \vee (Y \wedge \neg Z) \vee \neg Y) \\ = &X \wedge ((Y \wedge (Z \vee \neg Z)) \vee \neg Y) \\ = &X \wedge ((Y \wedge T) \vee \neg Y) \\ = &X \wedge (Y \vee \neg Y) \\ = &X \wedge T \\ = &X \end{split}$$

- 6. (b) Let P be false, and R and S be true.  $P \vee R$  is true because R is true.  $P \to S$  is true because P is false.  $R \to S$  is true because both R and S are true. This is an example where all three given propositions hold, yet P is false. Therefore the given implications do not hold.
  - (d) Let R be false, and P and S be true.  $P \vee R$  is true because P is true.  $R \to S$  is true because R is false.  $P \to S$  is true because both P and S are true. This is an example where all three given propositions hold, yet  $R \vee \neg S$  is false. Therefore the given implications do not hold.

7.

$$\begin{split} \alpha &= A \wedge B \wedge ((A \wedge B) \rightarrow C) \\ \alpha &= A \wedge B \wedge \neg ((A \wedge B) \wedge \neg C) \\ \alpha &= A \wedge B \wedge (\neg (A \wedge B) \vee C) \\ \alpha &= A \wedge B \wedge (\neg A \vee \neg B \vee C) \\ \alpha &= (A \wedge B \wedge \neg A) \vee (A \wedge B \wedge \neg B) \vee (A \wedge B \wedge C) \\ \alpha &= F \vee F \vee (A \wedge B \wedge C) \\ \alpha &= A \wedge B \wedge C \end{split}$$

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\beta = A \land C \land (B \to (A \land C)) \land (C \to B)
\beta = A \land C \land \neg (B \land \neg (A \land C)) \land \neg (C \land \neg B)
\beta = A \wedge C \wedge (\neg B \vee (A \wedge C)) \wedge (\neg C \vee B)
\beta = A \wedge C \wedge (\neg B \vee A) \wedge (\neg B \vee C) \wedge (\neg C \vee B)
\beta = A \wedge (\neg B \vee A) \wedge (\neg B \vee C) \wedge ((C \wedge \neg C) \vee (C \wedge B))
\beta = A \wedge (\neg B \vee A) \wedge (\neg B \vee C) \wedge (F \vee (C \wedge B))
\beta = A \land B \land C \land (\neg B \lor A) \land (\neg B \lor C)
\beta = A \wedge C \wedge ((B \wedge \neg B) \vee (B \wedge A)) \wedge ((B \wedge \neg B) \vee (B \wedge C))
\beta = A \wedge C \wedge (F \vee (B \wedge A)) \wedge (F \vee (B \wedge C))
\beta = A \wedge C \wedge B \wedge A \wedge B \wedge C
\beta = A \wedge B \wedge C
\gamma = (C \land B) \land ((C \lor A) \land \neg(C \to \neg A))
\gamma = C \wedge B \wedge ((C \vee A) \wedge (C \wedge A))
\gamma = C \wedge B \wedge (C \vee A) \wedge C \wedge A
\gamma = A \wedge B \wedge C \wedge (C \vee A)
\gamma = A \wedge B \wedge C
\delta = C \land (B \lor (A \to C)) \land (C \to (A \lor B))
\delta = C \wedge (B \vee \neg (A \wedge \neg C)) \wedge \neg (C \wedge \neg (A \vee B))
\delta = C \wedge (B \vee \neg A \vee C) \wedge (\neg C \vee A \vee B)
\delta = C \wedge (\neg C \vee A \vee B)
\delta = (C \land \neg C) \lor (C \land \neg A) \lor (C \land \neg B)
\delta = (C \land \neg A) \lor (C \land \neg B)
\delta = C \wedge (\neg A \vee \neg B)
\delta = \neg (A \land B) \land C
\epsilon = (A \land B \land C \land D) \lor (A \land ((\neg B \land C) \to D) \land (A \to \neg (B \to \neg C)))
\epsilon = (A \land B \land C \land D) \lor (A \land \neg((\neg B \land C) \land \neg D) \land \neg(A \land \neg(B \land C)))
\epsilon = (A \land B \land C \land D) \lor (A \land (\neg(\neg B \land C) \lor D) \land (\neg A \lor (B \land C)))
\epsilon = (A \land B \land C \land D) \lor (A \land ((B \lor \neg C) \lor D) \land (\neg A \lor (B \land C)))
\epsilon = (A \land B \land C \land D) \lor (A \land (B \lor \neg C \lor D) \land (\neg A \lor (B \land C)))
\epsilon = (A \land B \land C \land D) \lor ((B \lor \neg C \lor D) \land ((A \land \neg A) \lor (A \land (B \land C))))
\epsilon = (A \land B \land C \land D) \lor ((B \lor \neg C \lor D) \land A \land B \land C)
\epsilon = (A \land B \land C \land D) \lor (A \land B \land C)
\epsilon = A \wedge B \wedge C
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 $\alpha$ ,  $\beta$ ,  $\gamma$  and  $\epsilon$  are all equivalent.  $\delta$  is not equivalent to any of them.