

# Mathematical Logic HW3

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1. (b)

$$(A \leftrightarrow (\neg A \wedge B)) \rightarrow (B \wedge \neg B)$$

$A$	$B$	$\neg A \wedge B$	$A \leftrightarrow (\neg A \wedge B)$	$(A \leftrightarrow (\neg A \wedge B)) \rightarrow (B \wedge \neg B)$
$T$	$T$	$F$	$F$	$T$
$T$	$F$	$F$	$F$	$T$
$F$	$T$	$T$	$F$	$T$
$F$	$F$	$F$	$T$	$F$

i. DNF:

$$(A \wedge B) \vee (A \wedge \neg B) \vee (\neg A \wedge B)$$

ii. CNF:

$$A \vee B$$

(d)

$$(x \rightarrow y) \rightarrow (z \wedge (y \downarrow x))$$

$x$	$y$	$z$	$x \rightarrow y$	$y \downarrow x$	$z \wedge (y \downarrow x)$	$(x \rightarrow y) \rightarrow (z \wedge (y \downarrow x))$
$T$	$T$	$T$	$T$	$T$	$T$	$T$
$T$	$T$	$F$	$T$	$T$	$F$	$F$
$T$	$F$	$T$	$F$	$T$	$T$	$T$
$T$	$F$	$F$	$F$	$T$	$F$	$T$
$F$	$T$	$T$	$T$	$T$	$T$	$T$
$F$	$T$	$F$	$T$	$T$	$F$	$F$
$F$	$F$	$T$	$T$	$F$	$F$	$F$
$F$	$F$	$F$	$T$	$F$	$F$	$F$

i. DNF:

$$(x \wedge y \wedge z) \vee (x \wedge \neg y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \vee (\neg x \wedge y \wedge z)$$

ii. CNF:

$$(\neg x \vee \neg y \vee z) \wedge (x \vee \neg y \vee z) \wedge (x \vee y \vee \neg z) \wedge (x \vee y \vee z)$$

2. (b)

$$((p \downarrow q) \rightarrow r) \uparrow p$$

$p$	$q$	$r$	$p \downarrow q$	$(p \downarrow q) \rightarrow r$	$((p \downarrow q) \rightarrow r) \uparrow p$
$T$	$T$	$T$	$F$	$T$	$F$
$T$	$T$	$F$	$F$	$T$	$F$
$T$	$F$	$T$	$F$	$T$	$F$
$T$	$F$	$F$	$F$	$T$	$F$
$F$	$T$	$T$	$F$	$T$	$T$
$F$	$T$	$F$	$F$	$T$	$T$
$F$	$F$	$T$	$T$	$T$	$T$
$F$	$F$	$F$	$T$	$F$	$T$

As shown in the truth table,  $((p \downarrow q) \rightarrow r) \uparrow p$  is neither a tautology nor a contradiction.

3. (b)

$$\begin{aligned}
p \oplus q &= (p \wedge \neg q) \vee (\neg p \wedge q) \\
&= (\neg \neg p \wedge \neg q) \vee (\neg p \wedge \neg \neg q) \\
&= (\neg(p \downarrow p) \wedge \neg q) \vee (\neg p \wedge \neg(q \downarrow q)) \\
&= ((p \downarrow p) \downarrow q) \vee (p \downarrow (q \downarrow q)) \\
&= \neg(((p \downarrow p) \downarrow q) \downarrow (p \downarrow (q \downarrow q))) \\
&= (((p \downarrow p) \downarrow q) \downarrow (p \downarrow (q \downarrow q))) \downarrow (((p \downarrow p) \downarrow q) \downarrow (p \downarrow (q \downarrow q)))
\end{aligned}$$

4. (b)

$$f(x, y, z) = \neg x \wedge (\neg y \rightarrow \neg z)$$

$$\begin{aligned}
f(a, a, a) &= \neg a \wedge (\neg a \rightarrow \neg a) \\
&= \neg a \wedge T \\
&= \neg a
\end{aligned}$$

$$\begin{aligned}
f(a, a, f(a, a, a)) &= f(a, a, \neg a) \\
&= \neg a \wedge (\neg a \rightarrow a) \\
&= \neg a \wedge \neg(\neg a \wedge \neg a) \\
&= \neg a \wedge \neg \neg a \\
&= \neg a \wedge a \\
&= F
\end{aligned}$$

$$\begin{aligned}
f(f(a, a, f(a, a, a)), b, b) &= f(F, b, b) \\
&= T \wedge (\neg b \rightarrow \neg b) \\
&= T \wedge T \\
&= T
\end{aligned}$$

5. (b) We must show that  $\{\oplus, \wedge, \leftrightarrow\}$  are complete. One way to do that, is to show that with these operators we are able to make a  $\downarrow$ . We are able to make a  $\downarrow$  as follows:

$a$	$b$	$a \oplus b$	$a \wedge b$	$a \leftrightarrow b$	$(a \oplus b) \leftrightarrow (a \wedge b)$
$T$	$T$	$F$	$T$	$T$	$F$
$T$	$F$	$T$	$F$	$F$	$F$
$F$	$T$	$T$	$F$	$F$	$F$
$F$	$F$	$F$	$F$	$T$	$T$

As the table shows,  $(a \oplus b) \leftrightarrow (a \wedge b) = a \downarrow b$ , and we know that  $\downarrow$  on its own is itself complete, thus the set of operators  $\{\oplus, \wedge, \leftrightarrow\}$  is also complete.

(d) We must show that the functions  $f$  and  $g$  defined below are complete.

$$\begin{aligned}
f(x, y, z) &= y \rightarrow ((x \wedge \neg z) \vee (z \wedge \neg x)) \\
g(x, y) &= \neg x \vee y
\end{aligned}$$

$x$	$y$	$z$	$f(x, y, z)$	$g(x, y)$
$T$	$T$	$T$	$F$	$T$
$T$	$T$	$F$	$T$	$T$
$T$	$F$	$T$	$T$	$F$
$T$	$F$	$F$	$T$	$F$
$F$	$T$	$T$	$T$	$T$
$F$	$T$	$F$	$F$	$T$
$F$	$F$	$T$	$T$	$T$
$F$	$F$	$F$	$T$	$T$

$$\begin{aligned} g(a, a) &= \neg a \vee a \\ &= T \end{aligned}$$

$$\begin{aligned} f(g(a, a), a, b) &= f(T, a, b) \\ &= a \rightarrow ((T \wedge \neg b) \vee (b \wedge F)) \\ &= a \rightarrow (\neg b \vee F) \\ &= a \rightarrow \neg b \\ &= \neg(a \wedge \neg \neg b) \\ &= \neg(a \wedge b) \\ &= a \uparrow b \end{aligned}$$

We already know that  $\uparrow$  is complete, and we can make a  $\uparrow$  using the functions  $f$  and  $g$ , so it follows that  $f$  and  $g$  are complete.