

**Supplementary Document**  
for  
The  $k$ -RF Measures for Labeled Trees

by E. Khayatian, G. Valiente, L. Zhang

**Contents**

<b>1</b>	<b>Proof of Propositions 1–3 in Section 4</b>	<b>2</b>
<b>2</b>	<b>Supplementary figures for the frequency distribution of pair-wise <math>k</math>-RF scores</b>	<b>7</b>
<b>3</b>	<b>Proof of Propositions 4 and 5</b>	<b>9</b>

# 1 Proof of Propositions 1–3 in Section 4

**Proposition 1.** *Let  $T$  and  $S$  be two 1-labeled trees over  $L(T)$  and  $L(S)$ , respectively.*

(a) *Let  $|L(S) \cap L(T)| \leq 2$ ,  $k \geq 1$ , and  $|E(T)| \geq 2$ . Then,*

$$d_{k\text{-RF}}(S, T) = |E(S)| + |E(T)|.$$

(b) *Suppose that  $L(S) \neq L(T)$  and  $k < \min\{\text{diam}(T), \text{diam}(S)\}$ . Then,*

$$k + 1 \leq d_{k\text{-RF}}(S, T) \leq |E(S)| + |E(T)|.$$

*In addition, if the trees have the same size, then we have*

$$2(k + 1) \leq d_{k\text{-RF}}(S, T) \leq |E(S)| + |E(T)|.$$

(c) *For  $k = 0$ ,  $d_{k\text{-RF}}(S, T) = |E(S) \triangle E(T)|$ .*

(d) *If  $k \geq \max\{\text{diam}(T), \text{diam}(S)\} - 1$ , then  $d_{k\text{-RF}}(S, T) = d_{RF}(S, T)$ .*

*Proof.* Note that if  $k \geq 1$  and  $|E(T)| \geq 2$ , then each  $P_T(e, k)$  involves at least three labels. Thus, it is obvious that if  $|L(S) \cap L(T)| \leq 2$ , then for every  $e \in E(T)$ ,  $\acute{e} \in E(S)$ , we have  $P_T(e, k) \neq P_S(\acute{e}, k)$ . Thus, we have  $\mathcal{P}_k(S) \cap \mathcal{P}_k(T) = \emptyset$ , which implies that

$$d_{k\text{-RF}}(S, T) = |(\mathcal{P}_k(S) \cup \mathcal{P}_k(T)) \setminus (\mathcal{P}_k(S) \cap \mathcal{P}_k(T))| = |E(S)| + |E(T)|.$$

Thus, (a) is true.

For part (b), without loss of generality, we can assume that  $e = (u, v) \in E(T)$  such that  $\ell(v) \not\subseteq L(S)$ . Obviously,  $P_T(e, k)$  does not belong to  $\mathcal{P}_k(S)$ . Thus,  $d_{k\text{-RF}}(S, T) \geq 1$ . Now, if  $\text{diam}(T) > k \geq 1$ , we have either  $\deg(u) \geq 2$  or  $\deg(v) \geq 2$ . Thus, we can find another edge, namely  $e_1 = \{u_1, u\}$  such that  $P_T(e_1, k)$  does not belong to  $\mathcal{P}_k(S)$ . If  $k = 1$ , we have  $k + 1 \leq d_{k\text{-RF}}(S, T)$ . Otherwise, since  $\text{diam}(T) > k \geq 2$ , we have either  $\deg(v) \geq 2$  or  $\deg(u_1) \geq 2$  or  $\deg(u) \geq 3$  and degree of the third node adjacent to  $u$  is at least 2. Hence, we can find another edge, namely  $e_2$  such that  $P_T(e_2, k)$  does not belong to  $\mathcal{P}_k(S)$ . Continuing this process, we can find  $e_1, \dots, e_k$  such that  $P_T(e_1, k), \dots, P_T(e_k, k)$  do not belong to  $\mathcal{P}_k(S)$ . Thus, we have  $k + 1 \leq d_{k\text{-RF}}(S, T)$ . In addition, if the trees have the same size, we can find  $k$  distinct edge-induced partitions in  $\mathcal{P}_k(S) \setminus \mathcal{P}_k(T)$ , which implies that  $2(k + 1) \leq d_{k\text{-RF}}(S, T)$ . Finally, since  $d_{k\text{-RF}}(S, T) \leq |E(S)| + |E(T)|$ , the statement holds.

For part (c), note that we may represent each node of a 1-labeled tree with its unique label. As a result,  $P_T(e, 0) = e$  and  $P_S(\bar{e}, 0) = e$  for  $e \in E(T)$  and  $\bar{e} \in E(S)$ . Thus, (c) follows. Finally, Part (d) follows from the definition of the  $k$ -RF measure.  $\square$

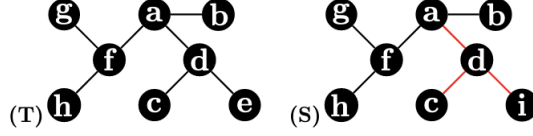


Figure 1: Two 1-labeled trees with the same size. The pairs induced by red edges in  $S$  are the only pairs different from all edge-induced pairs in  $T$ . Thus, the 1-RF score of  $S$  and  $T$  is 6.

**Proposition 2.** *Let  $k \geq 0$  be an integer. The  $k$ -RF dissimilarity measure is a pseudometric on the space of all 1-labeled trees. In other words, it satisfies the non-negativity, symmetry and triangle inequality conditions.*

*Proof.* Let  $T_1$ ,  $T_2$ , and  $T_3$  be three 1-labeled trees. Then, we show that the following inequalities hold.

$$\begin{aligned} d_{k\text{-RF}}(T_1, T_2) &\leq d_{k\text{-RF}}(T_1, T_3) + d_{k\text{-RF}}(T_3, T_2), \\ d_{k\text{-RF}}(T_1, T_2) &\geq 0, \\ d_{k\text{-RF}}(T_1, T_2) &= d_{k\text{-RF}}(T_2, T_1). \end{aligned}$$

Let  $\mathcal{P}_k(T_1)$ ,  $\mathcal{P}_k(T_2)$  and  $\mathcal{P}_k(T_3)$  be the three sets corresponding to  $T_1, T_2, T_3$ , respectively. The first inequality follows from the following relation:

$$\mathcal{P}_k(T_1) \triangle \mathcal{P}_k(T_2) \subseteq (\mathcal{P}_k(T_1) \triangle \mathcal{P}_k(T_3)) \cup (\mathcal{P}_k(T_3) \triangle \mathcal{P}_k(T_2)).$$

Also, the second and third inequalities follow from the definition of the  $k$ -RF measure.  $\square$

**Remark 1.** *Propositions 1 and 2, by the same proof, also hold for 1-labeled rooted trees.*

**Lemma 1.** *Let  $k \geq 0$  and let  $T$  be a 1-labeled rooted tree with  $n$  nodes. For all  $w$  and all  $i \leq k$ ,  $D_i(w) = \{w\} \cup \{x \in D_T(w) | d(w, x) \leq i\}$  and  $L(D_i(w))$  can be computed in  $O(kn)$  set operations, where  $D_T(w)$  consists of all the descendants of  $w$ .*

*Proof.* There are  $n$  labels in  $T$ . By ordering the  $n$  labels, we represent each label subset as a  $n$ -bit string, in which the  $i$ -th bit is 1 if and only if the  $i$ -th label is in the subset.

We prove the statement by induction. In the case  $k = 0$ ,  $D_0(w) = \{w\}$  and  $L(D_0(w)) = \{\ell(w)\}$ . Clearly, all the  $D_0(w)$  ( $w \in V(T)$ ) can be computed in  $O(n)$  set operations.

Assume that all the  $D_{k-1}(w)$  ( $w \in V(T)$ ) can be computed in at most  $2kn$  set operations. Assume  $w$  has  $d_w$  children  $u_1, u_2, \dots, u_{d(w)}$ . Then,

$$D_k(w) = \{w\} \cup \left( \bigcup_{i=1}^{d_w} D_{k-1}(u_i) \right)$$

and so

$$L(D_k(w)) = \{\ell(w)\} \cup \left( \bigcup_{i=1}^{d_w} L(D_{k-1}(u_i)) \right).$$

This implies that all  $L(D_k(w))$  can be computed from all  $\ell(w)$  and  $L(D_{k-1}(w))$  ( $w \in V(T)$ ) using  $\sum_{w \in V(T)} (1 + d_w) = 2n - 1$  set operations. In total, we can compute all  $L(D_k(w))$  in at most  $2n - 1 + 2kn = 2(k+1)n$  set operations. Note that this proof implies a dynamic programming algorithm for computing the label subsets in  $2(k+1)n = O(kn)$  set operations.  $\square$

**Lemma 2.** *Let  $k \geq 0$  and  $T$  be a 1-labeled rooted tree with  $n$  nodes. Using  $L(D_i(w))$  ( $w \in V(T), 0 \leq i \leq k$ ), we can compute  $L(B_k(w))$  for all  $w$  in  $O(kn)$  set operations, where*

$$B_k(w) = \{x \in V(T) \mid \exists y \in A_T(w) \cup \{w\} : d(y, w) + d(y, x) \leq k\}.$$

*Proof.* Let  $T$  be a 1-labeled tree and  $r$  be its root. For any node  $w \in V(T)$ , we assume the unique path from  $r$  to  $w$  be:

$$w_0 = r, w_1, \dots, w_t = w.$$

Then, we have that

$$B_k(w_t) = \bigcup_{i=0}^{\min(k, t)} D_{k-i}(w_{t-i}).$$

Given all  $L(D_i(w))$ ,  $L(B_k(w))$  is a union of at most  $k$  known label subsets and so can be computed in at most  $k$  set operations for each node  $w$ . In total, we can compute all  $L(B_k(w))$  in  $O(kn)$  set operations.  $\square$

**Proposition 3.** *Let  $S$  and  $T$  be two 1-labeled trees each with  $n$  nodes and  $k \geq 0$ . Then,  $d_{k\text{-RF}}(S, T)$  can be computed in  $O(kn)$  time.*

*Proof.* We first consider the rooted tree case. Let  $T$  and  $S$  be two 1-labeled rooted trees with  $n$  nodes. Without loss of generality, we may assume that

$S$  and  $T$  are labeled with the same set  $L$ ,  $|L| = n$ . By Lemma 1 and Lemma 2, We can compute all  $P_X(e, k)$  ( $e \in E(X)$ ) in  $O(kn)$  set operations for  $X = S, T$ . Since each edge-induced is an ordered pair of label subsets and we represent each label subset using a  $n$ -bit string, we consider  $P_X(e, k)$  as a  $2n$ -bit string. In this way, we sort all the edge-induced partitions for each tree in  $O(n)$  time by radix sort (that is, indexing) and then compute the symmetric difference of the two set of edge-induced partitions in  $O(n)$ . This concludes the proof.

In the unrooted case, we first root the trees at a leaf. In this way, we can compute all the edge-induced partitions in the derived rooted trees in  $O(kn)$  time. Since the edge-induced partitions are unordered pairs of label subsets in the original trees, we rearrange the two label subsets of each obtained partition in such a way that the smallest label in the first subset is smaller than every label in the second one. After the rearrangement, we can use radix-sort the partitions and compute the  $k$ -RF score in linear time.  $\square$

**Remark 2.** Let  $k \geq 0$  be an integer. If  $T$  is a 1-labeled rooted tree over  $L(T)$ , then we have  $P_T(e, k) = P_T(\acute{e}, k)$  if and only if  $e = \acute{e}$ . Moreover, each  $P_T(e, k)$  consists of two disjoint sets.

**Theorem 1.** The 0-RF dissimilarity measure is a metric on the space of all 1-labeled rooted trees. In other words, besides satisfying non-negativity, symmetry and triangle inequality conditions, it is zero for two trees if and only if the trees are isomorphic.

*Proof.* Let  $S$  and  $T$  be two 1-labeled rooted trees. By the proof of Proposition 2, it is enough to show that if  $d_{0\text{-RF}}(S, T) = 0$  (equivalently,  $\mathcal{OP}_0(T) = \mathcal{OP}_0(S)$ ), then  $S \cong T$ . Clearly, the equality  $\mathcal{OP}_0(T) = \mathcal{OP}_0(S)$  implies that for each  $e = (u, v) \in E(T)$ , there is one and only one edge  $\bar{e} = (\bar{u}, \bar{v}) \in E(S)$  such that  $\ell(u) = \ell(\bar{u})$  and  $\ell(v) = \ell(\bar{v})$  (note that  $P_T(e, 0) = (\ell(v), \ell(u))$  and  $P_S(\bar{e}, 0) = (\ell(\bar{v}), \ell(\bar{u}))$ ). Now, we define  $f : V(T) \rightarrow V(S)$ , which maps  $v$  and  $u$  to  $\bar{v}$  and  $\bar{u}$ , respectively. Obviously,  $f$  defines an isomorphism between  $S$  and  $T$ .  $\square$

**Theorem 2.** Let  $k \geq 1$  be an integer. The  $k$ -RF dissimilarity measure is a metric on the space of all 1-labeled rooted trees. In other words, besides satisfying non-negativity, symmetry and triangle inequality conditions, it is zero for two trees if and only if the trees are isomorphic.

*Proof.* Let  $S$  and  $T$  be two 1-labeled rooted trees. By the proof of Proposition 2, it is enough to show that if  $d_{k\text{-RF}}(S, T) = 0$  (equivalently,  $\mathcal{OP}_k(T) = \mathcal{OP}_k(S)$ ), then  $S \cong T$ . We have  $V(T) = \cup_{l=\text{depth}(T)-1}^0 V_l(T) \cup \text{Leaf}(T)$  and

$V(S) = \cup_{l=\text{depth}(T)-1}^0 V_l(S) \cup \text{Leaf}(S)$ , where

$$V_l(T) = \{v \in V(T) \mid v \text{ is in level } l\}.$$

( $V_l(S)$  is defined similarly.) Now, we define  $f : V(T) \rightarrow V(S)$  as follows.

Clearly,  $v \in \text{Leaf}(T)$  (respectively,  $w \in \text{Leaf}(S)$ ) if and only if  $P_T(e, k) = (\ell(v), A)$  (respectively,  $P_S(\bar{e}, k) = (\ell(w), B)$ ), where  $e$  (resp.  $\bar{e}$ ) is the edge incident to  $v$  (respectively,  $w$ ),  $A \subseteq L(T)$  (resp.  $B \subseteq L(S)$ ), and  $\ell(v)$  (respectively,  $\ell(w)$ ) is a set with only one element. Now, we define  $f(v) = u$  where  $u \in \text{Leaf}(S)$  and  $\ell(v) = \ell(u)$ . Thus,  $f$  is a one-to-one correspondence between  $\text{Leaf}(T)$  and  $\text{Leaf}(S)$ , where corresponding nodes have the same labels.

Suppose we have defined

$$f : \cup_{l=\text{depth}(T)-1}^n V_l(T) \cup \text{Leaf}(T) \rightarrow \cup_{l=\text{depth}(S)-1}^n V_l(S) \cup \text{Leaf}(S)$$

such that for each  $n \leq l \leq \text{depth}(T)-1$ ,  $f|_{V_l(T)} : V_l(T) \rightarrow V_l(S)$  and  $f|_{\text{Leaf}(T)} : \text{Leaf}(T) \rightarrow \text{Leaf}(S)$  is a one-to-one correspondence and corresponding nodes have the same labels. In addition, suppose that  $(v_1, v_2) \in E(T)$  if and only if  $(f(v_1), f(v_2)) \in E(S)$  for  $v_1, v_2 \in \cup_{l=\text{depth}(T)-1}^n V_l(T) \cup \text{Leaf}(T)$ . Now, let  $v \in V_{n-1}(T) \setminus \text{Leaf}(T)$ , and consider the edge connecting  $v$  to its parent. The pair induced by this edge is in the form  $(\ell(v) \cup (L(D_k(v)) \cap L(V_n(T))) \cup \dots \cup (L(D_k(v)) \cap L(\text{Leaf}(T))), A)$ , where  $A \subseteq L(T)$ . By our assumption, there is  $(u_1, u) \in E(S)$  inducing the same pair. On the other side, by properties of  $f$ , there are  $B_l \subseteq V_l(S), B \subseteq \text{Leaf}(S)$  with  $L(B_l) = L(D_k(v)) \cap L(V_l(T))$  and  $L(B) = L(D_k(v)) \cap L(\text{Leaf}(T))$  such that the first term in the pair is in the form  $\ell(u) \cup L(B_n) \cup \dots \cup L(B_{\text{depth}(T)-1}) \cup L(B)$ . Thus,  $\ell(u) = \ell(v)$ , and by definition of pairs assigned to a rooted tree, each node in  $B_n$  must be adjacent to  $u$  and  $u \in V_{n-1}(S)$ . Now, we define  $f(v) = u$ . Clearly,  $f : \cup_{l=\text{depth}(T)-1}^{n-1} V_l(T) \cup \text{Leaf}(T) \rightarrow \cup_{l=\text{depth}(S)-1}^{n-1} V_l(S) \cup \text{Leaf}(S)$  such that  $f|_{V_{n-1}(T)} : V_{n-1}(T) \rightarrow V_{n-1}(S)$  is a one-to-one correspondence and corresponding nodes have the same labels. Moreover, we have  $(v_1, v_2) \in E(T)$  if and only if  $(f(v_1), f(v_2)) \in E(S)$  for  $v_1, v_2 \in \cup_{l=\text{depth}(T)-1}^{n-1} V_l(T) \cup \text{Leaf}(T)$ .

Continuing the process, we infer that  $f : \cup_{l=\text{depth}(T)-1}^0 V_l(T) \cup \text{Leaf}(T) \rightarrow \cup_{l=\text{depth}(S)-1}^0 V_l(S) \cup \text{Leaf}(S)$  is an isomorphism which implies that  $S \cong T$ . (Note that  $\text{root}(T) \in V_0(T)$  is mapped to  $\text{root}(S) \in V_0(S)$  by  $f$ .)  $\square$

**Corollary 1.** *Let  $k \geq 0$  be an integer. The  $k$ -RF dissimilarity measure is a metric on the space of all 1-labeled trees.*

*Proof.* If  $k = 0$ , the statement follows from the same proof as for Proposition 1. Now, let  $S$  and  $T$  be two 1-labeled trees and  $k \geq 1$ . By Proposition 2,

it is enough to show that if  $d_{k\text{-RF}}(S, T) = 0$  (equivalently,  $\mathcal{P}_k(T) = \mathcal{P}_k(S)$ ), then  $S \cong T$ . Using the same strategy stated in the proof of Proposition 2, we can infer that there is a one-to-one correspondence between  $\text{Leaf}(T)$  and  $\text{Leaf}(S)$ , where corresponding nodes have the same labels. Let  $u \in \text{Leaf}(T)$  and  $v \in \text{Leaf}(S)$  be such that  $\ell(u) = \ell(v)$ . Then, we root  $T$  and  $S$  at  $u$  and  $v$  respectively. We call the induced rooted trees  $\dot{T}$  and  $\dot{S}$ .

Now, we show that  $\mathcal{OP}_k(\dot{T}) = \mathcal{OP}_k(\dot{S})$ . Let  $e_u \in E(T)$  and  $e_v \in E(S)$  are the edges leaving  $u$  and  $v$ , respectively. Since  $\ell(u) = \ell(v)$  and  $\mathcal{P}_k(T) = \mathcal{P}_k(S)$ , we have  $P_{\dot{T}}(e_u, k) = P_{\dot{S}}(e_v, k)$ . Now, suppose that the end points of  $e_u$  and  $e_v$  are  $u_1$  and  $v_1$ , respectively. Let  $e_{u_1}$  leave  $u_1$ . Then, we can find one and only one edge, namely  $e_{v_1}$  which leaves  $v_1$  such that  $P_{\dot{T}}(e_{u_1}, k) = P_{\dot{S}}(e_{v_1}, k)$  and  $\ell(u_1) = \ell(v_1)$ . Continuing this process, we have  $\mathcal{OP}_k(\dot{T}) = \mathcal{OP}_k(\dot{S})$ . Finally, the statement follows from Proposition 2.  $\square$

## 2 Supplementary figures for the frequency distribution of pairwise $k$ -RF scores

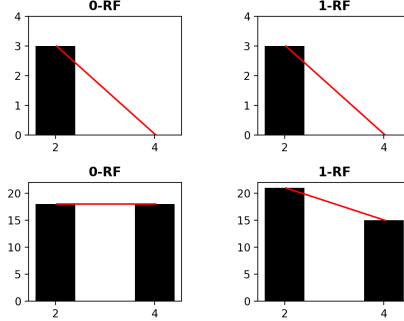


Figure 2: The frequency distribution of  $k$ -RF scores for the 1-labeled 3-node unrooted (top row) and rooted (bottom row) trees.

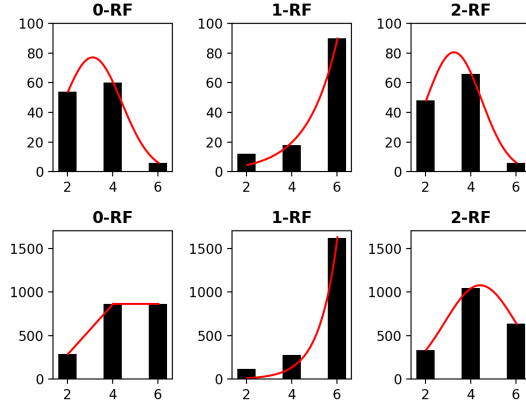


Figure 3: The frequency distribution of  $k$ -RF scores for the 1-labeled 4-node unrooted (top row) and rooted (bottom row) trees.

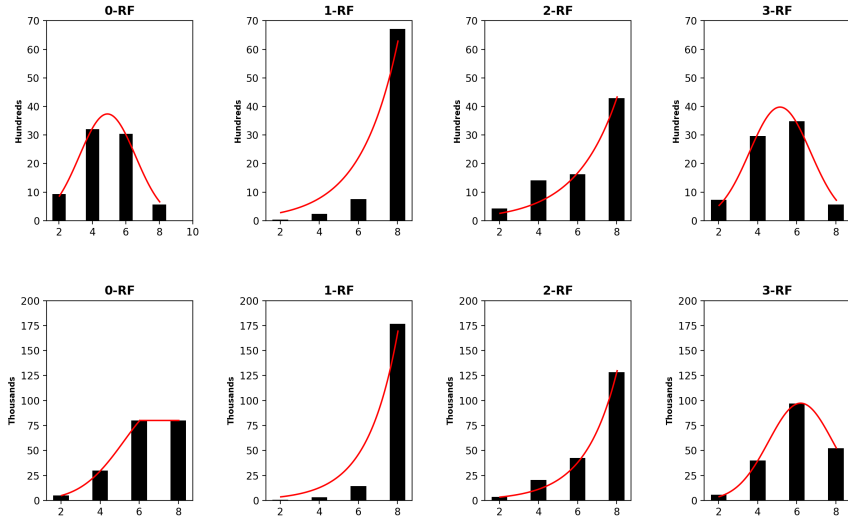


Figure 4: The frequency distribution of  $k$ -RF scores for the 1-labeled 5-node unrooted (top row) and rooted (bottom row) trees.



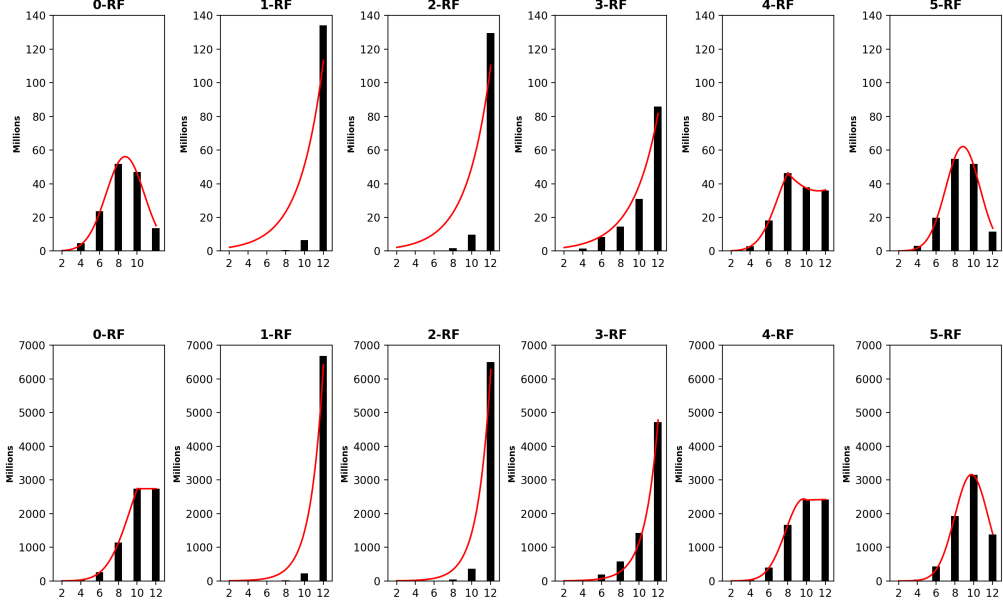


Figure 5: The frequency distribution of  $k$ -RF scores for the 1-labeled 7-node unrooted (top row) and rooted (bottom row) trees.

### 3 Proof of Propositions 4 and 5

**Proposition 4.** *Let  $k \geq 1$  be an integer. The  $k$ -RF dissimilarity measure is a pseudometric on the space of all trees whose nodes are labeled by multi-subsets of their corresponding label sets. In other words, it satisfies the non-negativity, symmetry and triangle inequality conditions.*

*Proof.* Let  $T_1$ ,  $T_2$ , and  $T_3$  be three labeled trees. Then, we have

$$\begin{aligned} d_{k\text{-RF}}(T_1, T_2) &\leq d_{k\text{-RF}}(T_1, T_3) + d_{k\text{-RF}}(T_3, T_2), \\ d_{k\text{-RF}}(T_1, T_2) &\geq 0, \\ d_{k\text{-RF}}(T_1, T_2) &= d_{k\text{-RF}}(T_2, T_1). \end{aligned}$$

Let  $\mathcal{P}_k(T_1)$ ,  $\mathcal{P}_k(T_2)$  and  $\mathcal{P}_k(T_3)$  be the three multisets corresponding to  $T_1$ ,  $T_2$ ,  $T_3$ , respectively. We show that the first inequality holds.

If  $x^{m(x)} \in \mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_2)$ , we have either  $x^{m(x)} \in \mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_2)$  or  $x^{m(x)} \in \mathcal{P}_k(T_2) \setminus_m \mathcal{P}_k(T_1)$ . Assume  $x^{m(x)} \in \mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_2)$ . Then, we have  $m_{\mathcal{P}_k(T_1)}(x) > m_{\mathcal{P}_k(T_2)}(x)$ . Now, if  $x \notin \text{Supp}(\mathcal{P}_k(T_3) \setminus_m \mathcal{P}_k(T_2))$ , we have  $m_{\mathcal{P}_k(T_1)}(x) > m_{\mathcal{P}_k(T_2)}(x) \geq m_{\mathcal{P}_k(T_3)}(x)$ , which implies that  $x \in$

$\text{Supp}(\mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_3))$  and

$$m_{\mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_3)}(x) = m_{\mathcal{P}_k(T_1)}(x) - m_{\mathcal{P}_k(T_3)}(x) \geq m_{\mathcal{P}_k(T_1)}(x) - m_{\mathcal{P}_k(T_2)}(x) = m(x).$$

Thus, we have

$$m(x) \leq m_{\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_3)}(x) + m_{\mathcal{P}_k(T_3) \triangle_m \mathcal{P}_k(T_2)}(x).$$

On the other hand, if  $x \in \text{Supp}(\mathcal{P}_k(T_3) \setminus_m \mathcal{P}_k(T_2))$  and  $m_{\mathcal{P}_k(T_3)}(x) \geq m_{\mathcal{P}_k(T_1)}(x)$ , then

$$m_{\mathcal{P}_k(T_3) \setminus_m \mathcal{P}_k(T_2)}(x) = m_{\mathcal{P}_k(T_3)}(x) - m_{\mathcal{P}_k(T_2)}(x) \geq m_{\mathcal{P}_k(T_1)}(x) - m_{\mathcal{P}_k(T_2)}(x) = m(x).$$

If  $x \in \text{Supp}(\mathcal{P}_k(T_3) \setminus_m \mathcal{P}_k(T_2))$  and  $m_{\mathcal{P}_k(T_3)}(x) < m_{\mathcal{P}_k(T_1)}(x)$ , we have  $m_{\mathcal{P}_k(T_1)}(x) > m_{\mathcal{P}_k(T_3)}(x) > m_{\mathcal{P}_k(T_2)}(x)$ , which implies that  $x \in \text{Supp}(\mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_3))$ .

Thus, we have

$$m(x) = m_{\mathcal{P}_k(T_1) \setminus_m \mathcal{P}_k(T_3)}(x) + m_{\mathcal{P}_k(T_3) \setminus_m \mathcal{P}_k(T_2)}(x) \leq m_{\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_3)}(x) + m_{\mathcal{P}_k(T_3) \triangle_m \mathcal{P}_k(T_2)}(x).$$

Similarly, if  $x^{m(x)} \in \mathcal{P}_k(T_2) \setminus_m \mathcal{P}_k(T_1)$ , then we obtain the same result.

To summarize, we have

$$\text{Supp}(\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_2)) \subseteq \text{Supp}(\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_3)) \cup \text{Supp}(\mathcal{P}_k(T_3) \triangle_m \mathcal{P}_k(T_2)).$$

In addition, for each  $x \in \text{Supp}(\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_2))$ , we have

$$m_{\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_2)}(x) \leq m_{\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_3)}(x) + m_{\mathcal{P}_k(T_3) \triangle_m \mathcal{P}_k(T_2)}(x).$$

Therefore, we have

$$|\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_2)| \leq |\mathcal{P}_k(T_1) \triangle_m \mathcal{P}_k(T_3)| + |\mathcal{P}_k(T_3) \triangle_m \mathcal{P}_k(T_2)|.$$

Thus, the first inequality holds. Also, the second inequality and the third equality follow from the definition of the  $k$ -RF measures.  $\square$

**Proposition 5.** *Let  $k \geq 0$  and  $S, T$  be two (rooted) trees whose nodes are labeled by  $L(S)$  and  $L(T)$ , respectively. Then,  $d_{k\text{-RF}}(S, T)$  can be computed in  $O(n(k + |L(S)| + |L(T)|))$  if the total multiplicity of each label is upper bounded by a constant.*

*Proof.* An algorithm in the 1-labeled case can be modified as follows for computing  $k$ -RF multiset-labeled unrooted and rooted trees:

- Represent each label multiset as a  $(|L(S)| + |L(T)|)$ -digit number, in which the digit at position  $j$  is the multiplicity of the  $j$ -th label. Compute all edge-induced partitions in  $O(kn)$  set operations.

- Radix-sort all the edge-induced partitions in each tree in  $O(n(|L(S)| + |L(T)|))$  time, where we use the assumption that the total multiplicity of each label is bounded by a constant  $B$ .
- Compute the symmetric difference of the set of the edge-induced partitions in the two input trees in  $O(n)$  time.

□

**Theorem 3.** *Let  $k \geq 1$  be an integer. The  $k$ -RF dissimilarity measure is a pseudometric on the space of all multiset-labeled rooted trees. In other words, it satisfies the non-negativity, symmetry and triangle inequality conditions.*

*Proof.* It is exactly the same as proof of Proposition 4.

□