# **Fundamentals of Data Structures**

# **Laboratory Projects 2**

**Tree Traversals** 

### **Chapter 1: Introduction**

#### 1.1 The content of the question

Given the partial results of a binary tree's traversals in in-order, pre-order, post-order. In this task, we are supposed to build a tree. After having solved the problem, we need to output the complete traversals of in-order, pre-order, post-order, and level order.

In the input, the first line of the input is the number of nodes you need to deal with in the tree. A represents the current number is missing, you are supposed to find it using your program. Finally, output your answer.

#### 1.2 The input and output sample(given by the question)

#### Sample Input 1

```
9
3 - 2 1 7 9 - 4 6
9 - 5 3 2 1 - 6 4
3 1 - - 7 - 6 8 -
```

#### **Sample Output 1**

```
3 5 2 1 7 9 8 4 6  /*in-order*/
9 7 5 3 2 1 8 6 4  /*pre-order*/
3 1 2 5 7 4 6 8 9  /*post-order*/
9 7 8 5 6 3 2 4 1  /*lever-order*/
```

#### Sample Input 2

```
3
----
-1-
1--
```

#### **Sample Output2**

```
Impossible
```

### **Chapter 2: Algorithm Specification**

#### 2.1 Searching Algorithm(searching for the node)

```
int Search(int L1,int R1,int L2,int R2,int L3,int R3)
  int root=pre[L2] | | post[R2];
 /*The last node of post-order or the first node of pre-order is the root number.*/
 for(int i=L1;i<=R1;i++)</pre>
    if(in[i]==root) p=i; /*find the root*/
 while(cannot find the root)
    if(in[i]==0) p=i; /*traverse all the i to find the position that have chance to be
the root*/
   for(int j=L2+1; j<=L2+length; j++)</pre>
      for(int k=L3;k<=L3+length;k++)</pre>
        judge(root);/*judge whether it's the root*/
       if(root) break;
      }
   i++;
    /*using enumeratio to judge the root in in-order.*/
  lc[root]=Search(newL1,newR1,newR2,newR2,newR3);//searching in the left subtree
  rc[root]=Search(newwL1,newwR1,newwL2,newwR2,newwR3);//searching in the right
subtree
  return root;
}
```

#### 2.2 BFS Algorithm

**Pseudo Code** 

```
void bfs(int root)
{
  int level[105]; //It's used to store the traversal of level order.
  int head=-1,tail=0;
  level[0]=root;
  while(head!=tail)
  {
    head++;
    if(lc[level[head]]) level[++tail]=lc[level[head]]//have a left child
    if(rc[level[head]]) level[++tail]=rc[level[head]]//have a right child
  }
}
```

#### 2.3 Main Function

```
int main()
{
    scanf("%d",&n);
    input(in[]);
    input(pre[]);
    input(post[]);
    int root = Search(1,n,1,n,1,n);
    bfs(root);
    output(in[]);
    output(pre[]);
    output(post[]);
    output(level[]);
}
```

## **Chapter 3: Test cases**

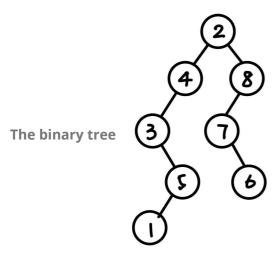
#### 3.1 Case 1: The Binary tree is a Queue

**Input Sample** 

```
9
3 5 9 - 4 2 7 - 6
2 4 - 3 9 1 8 - 6
3 - 9 5 4 7 - 8 2
```

#### **Output Sample**

```
3 5 9 1 4 2 7 8 6
2 4 5 3 9 1 8 7 6
3 1 9 5 4 7 6 8 2
2 4 8 5 7 6 3 9 1
```



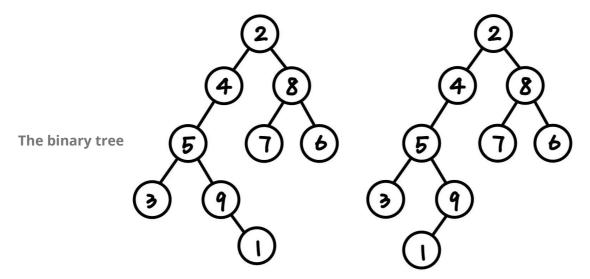
#### 3.2 Case 2: The tree have 2 possibilities

**Input Sample** 

```
9
- 5 9 - - - 8 - 6
2 4 5 - 9 - 8 - 7
3 1 - - 4 - 6 - 2
```

#### **Output Sample**

```
3 5 9 1 4 2 8 7 6
2 4 5 3 9 1 8 6 7
3 1 9 5 4 7 6 8 2
2 4 8 5 6 3 9 7 1
```



3.3 Case 3: The input is incorrect(we cannot build such a tree)

#### **Input Sample**

```
9
2 - 9 5 4 - - 8 1
1 - 4 3 - 9 - 7 6
- - 3 4 1 - - 8 2
```

#### **Output Sample**

```
Impossible
```

#### 3.5 Case 5: The tree has the minumum n.

**Input Sample** 

```
1
1
1
```

#### **Output Sample**

```
1
1
1
```

## **Chapter 4: Analysis and Comments**

#### 4.1 Complexity of Brute-Force Algorithm

4.1.1 The main idea of BFA

In the in-order traversal, BFA will enumerate all the possible root.

4.1.2 Time Complexity of the Algorithm

Time Complexity = O(N!)

The first root have n cases, the second n-1 cases......

4.1.3 Space Complexity of the Algorithm

Space Complexity = O(N)

We use an array to store the searching result.

#### 4.2 Complexity of BFS Algorithm

#### 4.2.1Time Complexity of the Algorithm

Time Complexity =O(N)

The worst case is that this binary tree is a queue, therefore we need O(N) to traverse all the node.

#### 4.2.2 Space Complexity of the Algorithm

```
Space Complexity =O(N)
```

That's because we use an array to store the searching result.

## **Chapter 5: Declaration**

I hereby declare that all the work done in this project is of my independent effort.

## Appendix: Source Code in C

```
#include<stdio.h>
#include<stdlib.h>
#include<stdbool.h>
int pre[105], in[105], post[105],level[105];
int lc[105], rc[105], root, n;
int f[105];//0-not visited,1-visited
void quit();
void bfs(int root);
int solve(int L1, int R1, int L2, int R2, int L3, int R3);
int main()
{
 int i;
   scanf("%d",&n);
   for (i = 1; i <= n; i++)
      scanf("%d",&in[i]);
   f[in[i]] = 1;
    for (i = 1; i \le n; i++)
      scanf("%d",&pre[i]);
    f[pre[i]] = 1;
    for (i = 1; i \le n; i++)
      scanf("%d",&post[i]);
    f[post[i]] = 1;
  }
    root = solve(1, n, 1, n, 1, n);
   bfs(root);
```

```
//print in-order
    printf("%d",in[1]);
    for(i=2;i<=n;i++)
      printf(" %d",in[i]);
   printf("\n");
    //pre-order
   printf("%d",pre[1]);
    for(i=2;i<=n;i++)
      printf(" %d",pre[i]);
   printf("\n");
    //post-order
    printf("%d",post[1]);
    for(i=2;i<=n;i++)
      printf(" %d",post[i]);
   printf("\n");
    //lever-order
 printf("%d",level[0]);
    for(i=1;i<n;i++)
      printf(" %d",level[i]);
   return 0;
}
void quit()
{
   printf("Impossible");
   exit(0);
}
void bfs(int root)
{
    int head = -1, tail = 0;
 //head represent the number of nodes we have asserted, tail represent the present one.
    level[0]=root;
  //the first node when we traverse in level-order is the root.
   while(head!=tail)
    {
        head++;
        //has a left child
    if (lc[level[head]])
    {
            tail++;
      level[tail]=lc[level[head]];
    }
    //has a right child
   if (rc[level[head]])
            tail++;
      level[tail]=rc[level[head]];
    }
    }
```

```
int solve(int L1, int R1, int L2, int R2, int L3, int R3)
 int i,j,k,tmp,size,root,p;
 int cnt=0, cnt1 = 0, cnt2 = 0;
  int fl = 0, flag = 0;
   if (L1 > R1) return 0;
   //have traversaled in-order
   if (L1 == R1)
  {
        tmp = in[L1] | pre[L2] | post[L3];
        if (!tmp)
    {
            for (i = n; i >= 1; i--)
                if (!f[i])
        {
                    tmp = i;
                    f[i] = 1;
                    break;
                }
        in[L1] = pre[L2] = post[L3] = tmp;
        return in[L1];
    }
    //impossible case: the last node of post-order doesn't equals to the first node of
pre-order
   if (pre[L2] && post[R3] && pre[L2]!=post[R3])
    //not both of the first node of pre-order and the last node of post-order are 0
   if (!pre[L2] | !post[R3])
        pre[L2] = post[R3] = pre[L2] + post[R3];
        //let pre[L2] and post[R3] contains the same node
    //we can't find the node
   if (!pre[L2])
   quit();
   root = pre[L2];
  p = L1;
  cnt = 0;
   for (i = L1; i <= R1; i++)
     if(in[i] == 0)
        cnt++;
      //count the nodes missing in in-order[L1,R1].
   if (root == in[i])
    {
            flag = 1;
        p=i;
        }
    }
```

```
//cant find the root and only one node missing, then it's the root.
    if(!flag && cnt==1)
      for(i=L1;i<=R1;i++)</pre>
        if(!in[i])
        in[i]=root;
        p = i;
        flag=1;
    if(flag) size = p - L1;
    else
  {
        for (p = L1 + 1; p \le R1; p++)
          if (in[p] == 0)
      {
               size = p - L1;
               //we need to judge which one is the root.
                //pre[L2 + 1 \sim L2 + size] root, L2+1, L2+2^{\circ} \neq R2
                //post[L3 \sim L3 + size - 1] L3,L3+1,L3+2° \neq R3
                int cnt1 = 0, cnt2 = 0;
                for (i = L2 + 1; i \le L2 + size; i++)
                     if (!pre[i]) cnt1++;//the number of 0 in pre-order
                 for (i = L3; i \le L3 + size - 1; i++)
                    if (!post[i]) cnt2++;//the number of 0 in post-order
                 for (i = L2 + 1; i \le L2 + size; i++)
                     if (pre[i])
          {
                         fl = 0;
                         for (j = L3; j \le L3 + size - 1; j++)
                             if (pre[i] == post[j]) fl = 1;
                         if (!fl) cnt2--;
                 for (i = L3; i \le L3 + size - 1; i++)
                     if (post[i])
          {
                         fl = 0;
                         for (j = L2 + 1; j \le L2 + size; j++)
                             if (post[i] == pre[j]) fl = 1;
                         if (!fl) cnt1--;
                if (cnt1 \ge 0\&\&cnt2 \ge 0) break;
                 //the number is all missing in pre-order and post-order -> it's the
root.
          }
    pre[L2] = root;
    post[R3] = root;
    in[p] = pre[L2];
    size = p - L1;
```

```
if(p > R1) quit();
    lc[root] = solve(L1, p - 1, L2 + 1, L2 + size, L3, L3 + size - 1);//searching the
left tree
    rc[root] = solve(p + 1, R1, L2 + size + 1, R2, L3 + size, R3 - 1);//searching the
right tree
    return root;
}
```