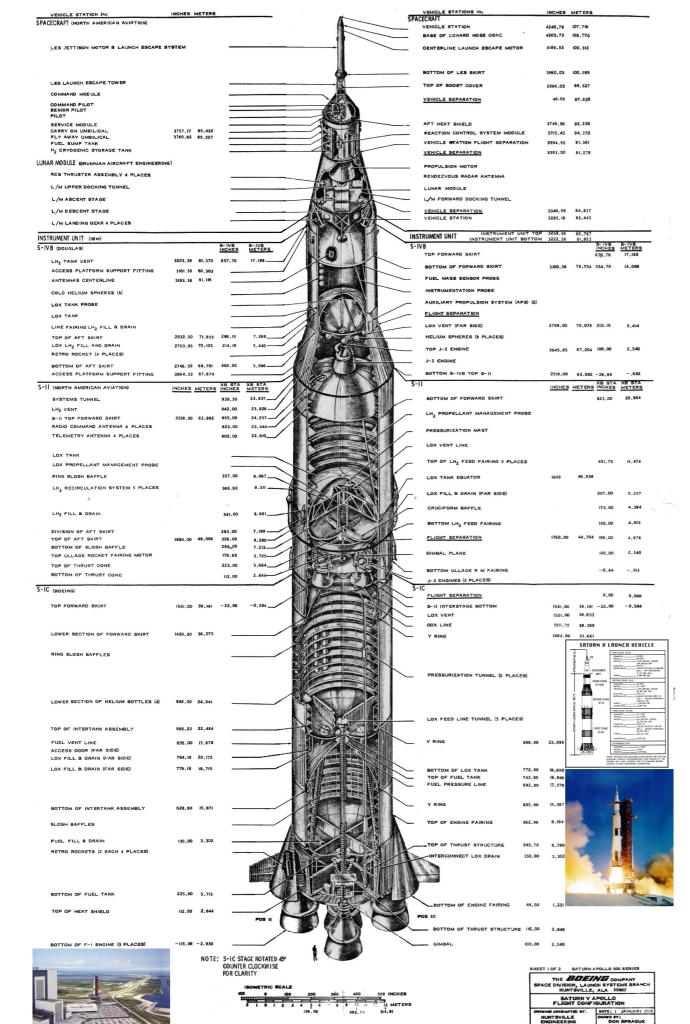
Performance





Command Module 64,000 lbs

Saturn V 6,200,000 lbs

Payload 1.5% of total weight





Demo: optimize gl_clear

Optimization Strategies

- Compiler flags leverage automatic optimizations
- Inlining Reduce function call overhead
- Avoid volatile enable more compiler optimizations
- Aggregate loads/stores less instructions per memory operation
- Loop optimizations code hoisting, combination, unrolling
- Manual assembly Be smarter than the compiler

$$S_{latency}(s) = \frac{1}{(1-p) + \frac{p}{s}}$$

- $S_{latency}$ the theoretical speedup of the execution of the entire program
- s the speedup of the part of the program you're optimizing
- p the proportion of execution time that the part the program you're optimizing originally occupied



Suppose your original program took t cycles to execute



The program is divided into two distinct portions:

- Part A takes 75% of the time
- Part B takes 25% of the time



If we optimize part B to make it 5 times faster, this only reduces the overall computation time slightly



If we optimize part A to make it just twice as fast, we get a greater overall speedup

Part B is 25% of the overall program (p = .25) and we speed it up by a factor of 5 (s = 5)

$$S_{latency} = \frac{1}{1 - .25 + \frac{.25}{5}} = 1.25$$

Overall program speedup is 1.25

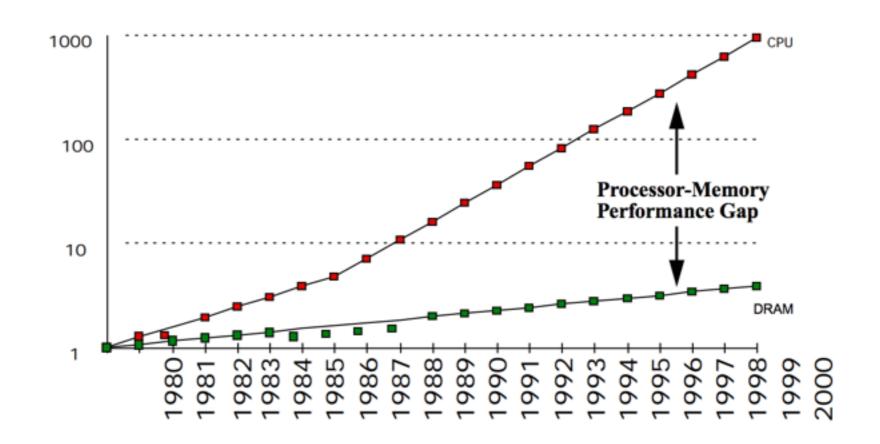
Part A is 75% of the overall program (p = .75) and we speed it up by a factor of 2 (s = 2)

$$S_{latency} = \frac{1}{1 - .75 + \frac{.75}{2}} = 1.60$$

Overall program speedup is 1.60

Beyond Software Optimization

Processor-Memory Performance Gap



[Hennessy, J.L.; Patterson, D.A. Computer Organization and Design, 2nd ed.]

Memory System Performance

| Processor | Alpha 21164 | |
|-------------------------|--------------------|----------------|
| Machine | AlphaServer 8200 | |
| Clock Rate | 300 MHz | |
| Memory Performance | Latency | Bandwidth |
| I Cache (8KB on chip) | 6.7 ns (2 clocks) | 4800 MB/sec |
| D Cache (8KB on chip) | 6.7 ns (2 clocks) | 4800 MB/sec |
| L2 Cache (96KB on chip) | 20 ns (6 clocks) | 4800 MB/sec |
| L3 Cache (4MB off chip) | 26 ns (8 clocks) | 960 MB/sec |
| Main Memory Subsystem | 253 ns (76 clocks) | 1200 MB/sec |
| Single DRAM component | ≈60ns (18 clocks) | ≈30–100 MB/sec |
| | | |

[Patterson, David, et al. "A case for intelligent RAM."]

Moving data between the cpu and memory is the bottleneck

strcpy

```
for (int i = 0; i <= strlen(src); i++) {
    dst[i] = src[i]
}</pre>
```

All we're doing is loading data from memory into the CPU and storing it back into memory

Avoiding the Memory Bottleneck

- Raspberry Pi has a DMA Controller that allows us to read and write memory without having to go through the processor (avoiding the load/store latency)
- Section 4 of BCM2835-ARM-Peripherals.pdf

Demo: dma_strcpy.c

Measuring Performance

Don't optimize blind

Profiling

- Analyze your program at runtime to measure characteristics of interest
 - Space/time complexity
 - Frequency of certain instructions
 - Frequency and Duration of Function Calls
- Most often used for guiding optimization

Taking Measurements

- Hardware interrupts gprof.c
- Code Instrumentation timer
- Also: instruction set simulation, OS hooks, performance counters
- Many techniques rely on sampling (statistical profilers) to trade off accuracy for speed

Visualizing Measurements

Chrome Developer Tools

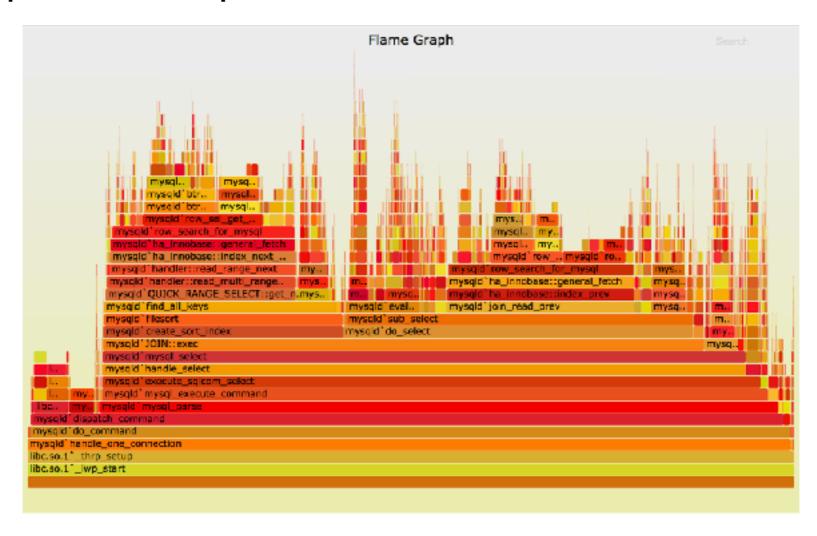


- Output profiling information in standard format (linux perf)
- Use thlorenz/cpuprofilify to convert into .cpuprofile format

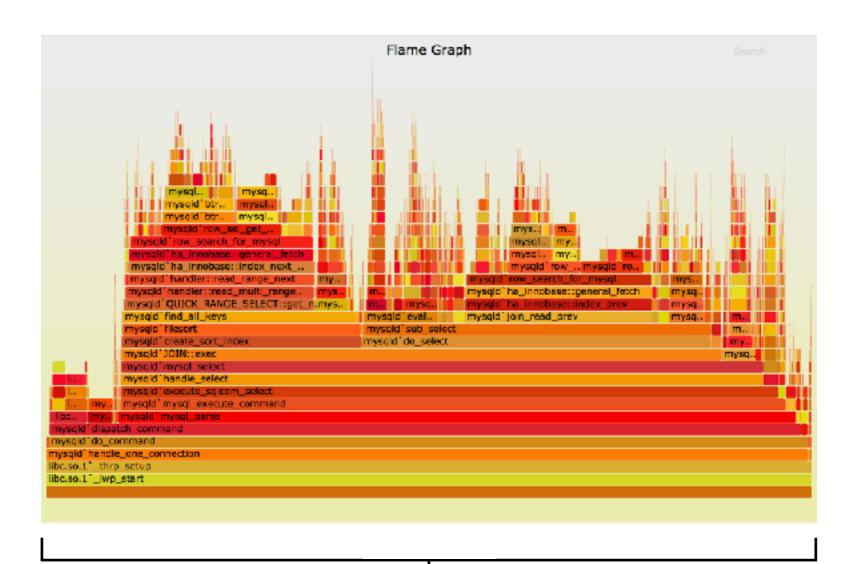
Demo: stackprof.c

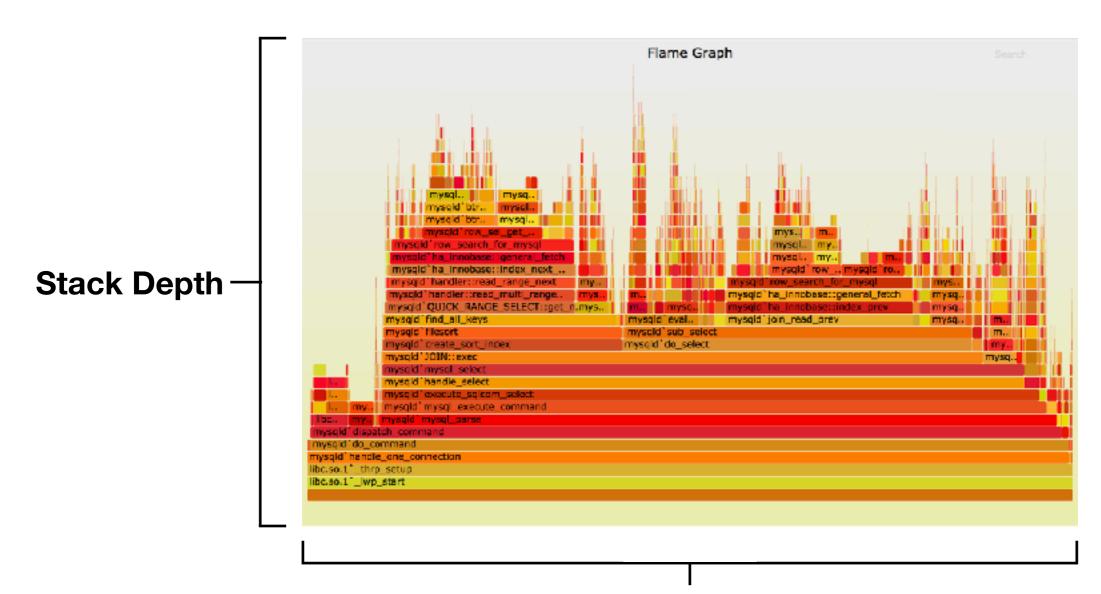
Visualizing Measurements

Flame Graph - designed to quickly and accurately find the most frequent code paths

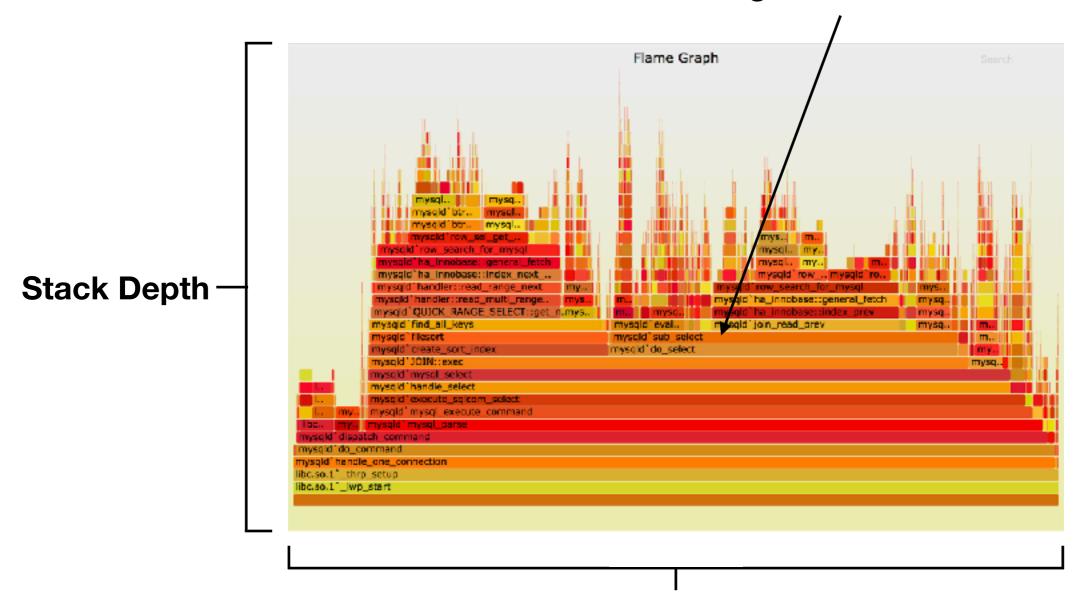


http://www.brendangregg.com/flamegraphs.html





Rectangle is a stack frame



Wider the frame, the more often is was on present in the stack

