Announcements

Upcoming schedule ...

Start thinking about final project; form teams of 2 people (1-3 ok)

Final 2 labs will be devoted to working on your project (work on proposal as team; check-in)

Late days apply through Assign7, not final project

Single and double buffering review

font_get_char demonstration

Interrupts



Blocking I/O

```
while (1) {
   read_char_from_keyboard();
   update_screen();
}
```

Read fn loops until char is received - blocking

char arrives

update_screen read ·····→ time

Blocking I/O

```
while (1) {
   read_char_from_keyboard();
   update_screen();
}
```

How long does it take to send a scan code? - 11kHz, 11 bits/scan code

How long does it take to update the screen?

What could go wrong?

code/glkeyboard

Blocking I/O

```
while (1) {
   read_char_from_keyboard();
   update_screen();
}
```

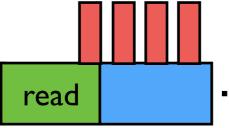
char arrives

Blocking I/O

```
while (1) {
   read_char_from_keyboard();
   update_screen();
}
```

chars arrive

time



The Problem

Need long-running computations (graphics, computations, applications, etc.).

Need to respond to external events quickly.

How could we change this code?

```
while (1) {
   read_char_from_keyboard();
   update_screen();
}
```

Interrupts

```
when a scan code arrives {
  add scan code to buffer();
while (1) {
  // Doesn't block
  while (read chars from keyboard()) {}
  update screen();
}
```

proc

update_screen

Interrupts to the Rescue

Cause processor to pause what it's doing and immediately execute interrupt code, returning to original code when done.

- **External events (reset, timer, GPIO)**
- Internal events (bad memory access, software trigger)

Critical for responsive systems.

Using interrupts exercises everything you've learned so far.

■ Architecture, assembly, linking, memory, C, peripherals

They'll complete your interactive graphics console.

code/blink

blink.c

Timer causes an interrupt once per second

while() loop in main() is interrupted

Interrupt handler increments counter

while() checks for counter change

Toggles ACT_LED and prints counter

Why is counter declared volatile?

How is interrupt_handler called?

Interrupt Vectors

8 Kinds of Interrupts

```
vectors:
  ldr pc, reset asm
  ldr pc, _undefined_instruction_asm
  ldr pc, software interrupt asm
  ldr pc, prefetch abort asm
  ldr pc, _data_abort_asm
  ldr pc, reset asm
  ldr pc, _interrupt_asm
  ldr pc, fast asm
                              .word impossible asm
reset asm:
_undefined_instruction_asm:
                              .word impossible asm
software interrupt asm:
                              .word impossible asm
prefetch abort asm:
                              .word impossible asm
_data_abort_asm:
                              .word impossible asm
                              .word interrupt asm
interrupt asm:
                              .word impossible asm
_fast_asm:
```

_vectors_end

Interrupt Vectors

```
vectors:
  ldr pc, reset asm
  ldr pc, _undefined_instruction_asm
  ldr pc, _software_interrupt_asm
  ldr pc, _prefetch_abort_asm
  ldr pc, _data_abort_asm
  ldr pc, _reserved_asm
                               branch to interrupt asm
  ldr pc, interrupt asm:
  ldr pc, _fast_asm
                              .word impossible asm
_reset_asm:
undefined instruction asm:
                              .word impossible asm
software interrupt asm:
                              .word impossible asm
_prefetch_abort_asm:
                              .word impossible asm
_data_abort_asm:
                              .word impossible_asm
_interrupt_asm:
                              .word interrupt asm
fast asm:
                              .word impossible asm
_vectors_end
```

How is interrupt_asm is called?
What is the difference between interrupt_asm
and _interrupt_asm?

Interrupt Vectors

```
vectors:
  ldr pc, reset asm
  ldr pc, _undefined_instruction_asm
  ldr pc, _software_interrupt_asm
  ldr pc, _prefetch_abort_asm
  ldr pc, _data_abort_asm
 ldr pc, _reset_asm
                               branch to interrupt asm
  ldr pc, interrupt asm:
  ldr pc, _fast_asm
                              .word impossible asm
reset asm:
undefined instruction asm:
                              .word impossible asm
software interrupt asm:
                              .word impossible asm
_prefetch_abort_asm:
                              .word impossible asm
_data_abort_asm:
                              .word impossible_asm
_interrupt_asm:
                              .word interrupt asm
                              .word impossible asm
fast asm:
_vectors_end
```

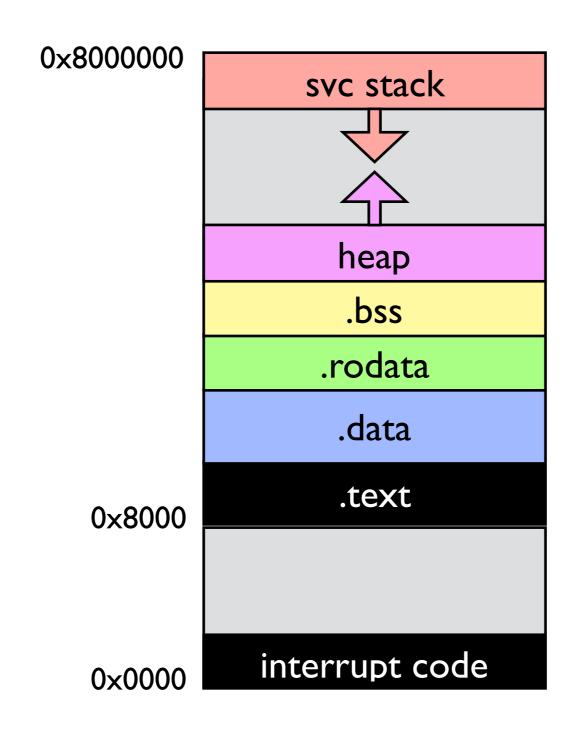
How does the system know what to call when an interrupt occurs?

cstart.c

```
static int * const RPI_INTERRUPT_VECTOR_BASE = 0x0;
int* vectorsdst = RPI_INTERRUPT_VECTOR_BASE;
int* vectors = &_vectors;
int* vectors_end = &_vectors_end;
while (vectors < vectors_end) {
    *vectorsdst++ = *vectors++;
}</pre>
```

Where do we put the interrupt vectors?
Why do we need to copy them?
Where are _vectors and _vectors_end defined?

CPU Address Space



```
vectors:
  ldr pc, _reset_asm
  ldr pc, _undefined_instruction_asm
  ldr pc, _software_interrupt_asm
  ldr pc, _prefetch_abort_asm
  ldr pc, _data_abort_asm
  ldr pc, _reset_asm
  ldr pc, _interrupt_asm:
                                branch to interrupt_asm
  ldr pc, _fast_asm
                              .word impossible asm
reset asm:
_undefined_instruction_asm:
                              .word impossible_asm
_software_interrupt_asm:
                              .word impossible asm
prefetch_abort_asm:
                              .word impossible_asm
                              .word impossible asm
_data_abort_asm:
_interrupt_asm:
                              .word interrupt asm
_fast_asm:
                              .word impossible_asm
vectors end
```

Why does this code work if it is copied to address 0?

```
Position independent code
% cd ../vector
% make main.list
% cat main.list
                            pc-relative addressing
00008040 <_vectors>:
   8040:
           ldr pc, [pc, #24]
                              ; 8060 <_impossible_asm>
   8044:
           ldr pc, [pc, #20]
                              ; 8060 <_impossible_asm>
                             ; 8060 < impossible_asm>
   8048: ldr pc, [pc, #16]
   804c: ldr pc, [pc, #12]
                             ; 8060 <_impossible_asm>
   8050: ldr pc, [pc, #8]
                             ; 8060 <_impossible_asm>
   8054: ldr pc, [pc, #4] ; 8060 <_impossible_asm>
   8058: ldr pc, [pc, #4] ; 8064 <_interrupt_asm>
           ldr pc, [pc, #-4]
                             ; 8060 <_impossible_asm>
   805c:
00008060 <_impossible_asm>:
   8060: .word
                  0x00008084
00008064 <_interrupt_asm>:
                  0x00008068
   8064:
           .word
00008084 <impossible_asm>:
00008068 <interrupt_asm>:
```

Interrupt Handler

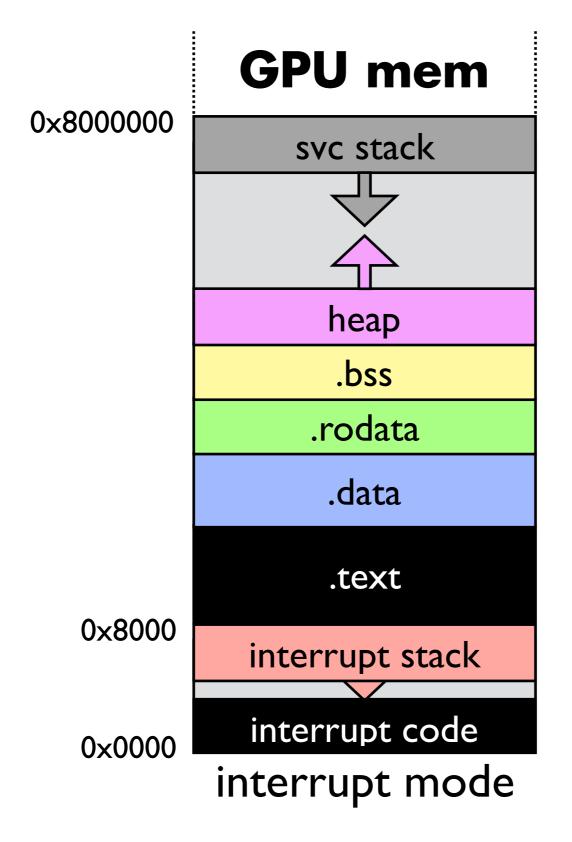
Interrupt Handler

```
interrupt_asm:
   mov sp, #0x8000
    sub 1r, 1r, #4
    push {r0-r12,lr}
   mov r0, lr
    bl interrupt handler
    pop {r0-r12, lr}
    movs pc, lr
```

```
interrupt asm:
    mov sp, #0x8000
    sub 1r, 1r, #4
    push {r0-r12,1r}
    mov r0, lr
    bl interrupt handler
    pop {r0-r12, lr}
    movs pc, lr
```

Why do we save all of the registers? Where do we save the registers?

Interrupt Stack



Interrupt occurs right before instruction

```
Disassembly of section .text:
00008000 <_start>:
   8000: e3a0d902
                                   sp, #32768 ; 0x8000
                            mov
   8004: eb000001
                            bl
                                   8010 < cstart>
00008008 <hang>:
   8008: eb000039
                                   80f4 <led_on>
                            bl
   800c: eaffffe
                                   800c <hang+0x4>
                            b
00008010 <_cstart>:
                                   {fp, lr} Interrupt!
         e92d4800
                            push
```

What is the pc when the interrupt occurs? Where can we store that information?

8010:

```
interrupt_asm:
    mov sp, #0x8000
    sub 1r, 1r, #4
    push {r0-r12,lr}
    mov r0, lr
    bl interrupt handler
    pop {r0-r12, lr}
    movs pc, lr
```

Why do we subtract 4 from 1r? What is the value passed to interrupt_handler?

Processor Modes

User - unprivileged mode

IRQ - interrupt mode

FIQ - fast interrupt mode

Supervisor - privileged mode, entered on reset

Abort - memory access violation

Undefined - undefined instruction

System - privileged mode that shares user regs

Shared / Unshared Registers

User32	FIQ32	Supervisor32	Abort32	IRQ32	Undefined32
R0	R0	R0	R0	R0	R0
R1	R1	R1	R1	R1	R1
R2	R2	R2	R2	R2	R2
R3	R3	R3	R3	R3	R3
R4	R4	R4	R4	R4	R4
R5	R5	R5	R5	R5	R5
R6	R6	R6	R6	R6	R6
R7	R7	R7	R7	R7	R7
R8	R8_fig	R8	R8	R8	R8
R9	R9_fig	R9	R9	R9	R9
R10	R10_fig	R10	R10	R10	R10
R11	R11_fig	R11	R11	R11	R11
R12	R12_fig	R12	R12	R12	R12
R13	R13_fig	R13_svc	R13_abt	R13_irq	R13_und
R14	R14_fig	R14_svc	R14_abt	R14_irq	R14_und
R15 (PC)	R15 (PC)	R15 (PC)	R15 (PC)	R15 (PC)	R15 (PC)
		Program Stat	us Registers		
CPSR	CPSR	CPSR	CPSR	CPSR	CPSR
	SPSR_fig	SPSR_svc	SPSR_abt	SPSR_irq	SPSR_und

Return from Interrupt

```
interrupt_asm:
    mov sp, #0x8000
    sub 1r, 1r, #4
    push {r0-r12,1r}
    mov r0, lr
    bl interrupt handler
    pop {r0-r12, lr}
    movs pc, lr
```

Restore the registers movs causes a return from interrupt

Can this code be written in C?

Summary

Interrupts allow external events to trigger code to run with very little delay: responsiveness despite long-running functions

■ They bring together everything you've learned so far

Running code at arbitrary points is dangerous!

■ Copies of Ir and sp, use separate stack

Interrupt vectors are at 0x0-0x1c

- Have to copy them there at boot time
- Generating safe assembly requires explicitly embedding addresses

Next time: using and writing interrupts (the return of GPIO)