

1 Introduction

After printing many review materials double-sided, we find it disturbing to flip them upside down during the review. Assume that N pages of the materials are double-sided printed, each side containing one of N chapters in the textbook. We wonder if there is a way to flip some of them so that we can cover all the N chapters at once.

In this program, we need to find out one flipping solution, which can let us cover all the N chapters at once.

For example if the Input is following :

```
5
3 2
2 5
4 1
1 3
2 1
```

The output should be : 3 5 4 1 2 (ignore the echo of input)

That means if we just flip the second page (swap 2 5 , the third line of input), and then we will get the first column as the answer, which let us cover all the 5 chapters at once.

2 Algorithm Specification

The first thing we should do is read the input and do process of them. We use **trap x20** to read. Because the N may be larger than 10, we should use the special character like **Space** and **Enter** as a signal. After finishing the read of a number ,we transfer it into digit rather than ascii code and store it into memory.

```
1      num ← 0 \\ initialize
2      trap x20
3      trap x21
4      while R0 != \n
5          num ← 10 * num + R0 - '0'
6      for i = 0 to num-1
7          num ← 0 \\ initialize
8          trap x20
9          trap x21
10         while R0 != ' '
11             num ← 10 * num + R0 - '0'
12         page[2*i] ← num
13         num ← 0 \\ initialize
14         trap x20
15         trap x21
16         while R0 != ' '
17             num ← 10 * num + R0 - '0'
18         page[2*i + 1] ← num
19
```

Then, we use the recursion to determine the flip condition. To simplify the coding procedure, we use an array to store whether one page need flip. Once we determine the i th page need to flip, then we recursively call function itself to get the flip condition of $(i+1)$ th page to n th page. This procedure will form a complete binary tree. When the program comes to the leave of the binary tree, the program will call the check function to check

whether in this flip condition, we can cover all the N chapters at once.

```
1  function flip(int n)
2      if n == num
3          call check()
4          return
5      flip[n] ← 0
6      call flip(n-1)
7      flip[n] ← 1
8      call flip(n-1)
9      flip[n] ← 0 \\recover to be init
10
```

When the program comes to the leave of binary tree, the check function will be called. In the check function, we check whether the chosen chapter of each page can cover from 1 to num. Firstly, according to the flip condition, we take out the front chapter or back chapter of each page and store the result into memory. Then we form a label area, according to the result we get before to set the label area. That means if the result contain chapter 4, then label[3] will be set to 1. Finally, we go through the label area. If it's all 1 in label area, then we just output the result we obtained before and halt the program.

```
1  function check()
2      for i = 0 to num - 1
3          if flip[i] == 1
4              result[i] ← page[2*i + 1]
5          else
6              result[i] ← page[2*i]
7      // clear the label area
8      for i = 0 to num - 1
9          label[i] ← 0
10     for i = 0 to num - 1
11         label[result[i]-1] ← 1
12     for i = 0 to num - 1
13         if label[i] == 0
14             return
15     //output
16     for i = 0 to num - 1
17         R0 ← result[i]
18         trap x21
19     halt
20
```

3 Q and A

- Q: what is your recursive algorithm?

A: in the program, we do the recursive algorithm according to the flip condition. Once we set the ith page to be flipped or not, we then recursively set the (i+1)th to nth pages. In this way, the program will form a binary tree. When the program comes to the leave of binary tree, we call the check function to check the condition.

