



University of Minho  
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Project: Report

# Marketing Digital Outdoor with gesture interaction — Analysis

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# List of Abbreviations

<b>Notation</b>	<b>Description</b>	<b>First used on page nr.</b>
AC	Alternating Current	12
ACF	Aggregated channel feature	79
ANN	Artificial Neural Networks	83
API	Application Programming Interface	4
BN	Billions	2
BSD	Berkeley Software Distribution	73
BSP	Board Support Package	13
CAGR	Compound Annual Growth Rate	3
CLI	Command Line Interface	13
CNN	Convolutional Neural Network	77
COTS	Commercial off-the-shelf	5
CPS	Cyber–Physical Systems	1
CPU	Central Processing Unit	80
CV	Computer Vision	15
DB	Database	12
DBMS	Database Management System	85
DC	Direct Current	12
DDL	Data Definition Language	94
DML	Data Manipulation Language	95
DOOH	Digital Out-Of-Home	2
DPM	Deformable part models	79
ER	Entity-Relationship	85, 87
ERD	Entities-Relationships diagram	92

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**List of Abbreviations**

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<b>Notation</b>	<b>Description</b>	<b>First used on page nr.</b>
GIF	Graphics Interchange Format	2
GPU	Graphics Processing Unit	80
GUI	Graphical User Interface	99
HD	High-Definition	78
HOG	Histogram of Oriented Gradients	79
HW	Hardware	5
I/O	Input/Output	65
IOT	Internet of Things	78
IP	Internet Protocol	15
IPC	Inter-Process Communication	59
KNN	K-nearest neighbor	83
LAN	Local Area Network	62
MAC		
	Media Access Control	62
MDO	Marketing Digital Outdoor	2
MDO-L	MDO Local System	9
MDO-RC	MDO Remote Client	9
MDO-RS	MDO Remote Server	9
NB	Naive Bayes	83
OCR	Optical Character Recognition	77
OMT	Object-Modeling Technique	58
OOSE	Object Oriented Software Engineering	58
OS	Operating System	13
OSI	Open Systems Interconnection	13
PCB	Printed Circuit Board	5
PGID	Process Group ID	68
PID	Process ID	68
POSIX	Portable Operating System Interface	58

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**List of Abbreviations**

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<b>Notation</b>	<b>Description</b>	<b>First used on page nr.</b>
R&D	Research and Development	3
RDBMS	Relational Database Management System	14
ROI	Region Of Interest	83
SID	Session ID	68
SoC	System-on-a-Chip	12
SQL	Structured Query Language	85
SVM	Support Vector Machine	79
SW	Software	5
TCP	Transmission Control Protocol	63
TCP/IP	Transmission Control Protocol/Internet Protocol	9
UDP	User Datagram Protocol	63
UI	User Interface	6
UML	Unified Modeling Language	57
USD	United States Dollar	78
WAL	Write-Ahead Log	88
WLAN	Wireless local Area Network	62

# **1. Introduction**

The present work, within the scope of the Embedded Systems course, consists in the project of a Cyber–Physical Systems (CPS), i.e., a system that provides seamless integration between the cyber and physical worlds [1]. The Waterfall methodology is used for the project development, providing a systematic approach to problem solving and paving the way for project's success.

In this chapter are presented the project's context and motivation, the problem statement — clearly defining the problem, the market research — defining the product's market share and opportunities, the project goals, the project planning and the document outline.

## **1.1. Context and motivation**

COVID pandemics presented a landmark on human interaction, greatly reducing the contact between people and surfaces. Thus, it is an imperative to provide people with contactless interfaces for everyday tasks. People redefined their purchasing behaviors, leading to a massive growth of the online shopping. However, some business sectors, like clothing or perfumes, cannot provide the same user experience when moving online. Therefore, one proposes to close that gap by providing a marketing digital outdoor for brands to advertise and gather customers with contactless interaction.

Scenting marketing is a great approach to draw people into stores. Olfactory sense is the fastest way to the brain, thus, providing an exceptional opportunity for marketing [2] — “75% of the emotions we generate on a daily basis are affected by smell. Next to sight, it is the most important sense we have” [3].

Combining that with additional stimuli, like sight and sound, can significantly boost the marketing outcome. Brands can buy advertisement space and time, selecting the videoclips to be displayed and the fragrance to be used at specific times, drawing the customers into their stores.

Marketing also leverages from better user experience, thus, user interaction is a must-have, providing the opportunity to interact with the customer. In this sense, when users approach the outdoor a gesture-based interface will be provided for a brand immersive experience, where the user can take pictures or create GIFs with brand specific image filters and share them through their social media, with the opportunity to gain

several benefits.

## 1.2. Problem statement

The first step of the project is to clearly define the problem, taking into consideration the problem's context and motivation and exploiting the market opportunities.

The project consists of a Marketing Digital Outdoor (MDO) with sound and video display, and fragrance emission selected by the brands, providing a gesture-based interface for user interaction to create pictures and GIFs, brand-specific, and share them on social media. It is comprised of several modes:

- normal mode (advertisement mode): the MDO will provide sound, video and fragrance outputs.
- interaction mode: When a user approaches the device, the MDO will go into interaction mode, turning on and displaying the camera feed and waiting for recognizable gestures to provide additional functionalities, such as brand-specific image filters.
- multimedia mode: in this mode the facial detection is applied, enabling the user to select and apply different brand-specific image filters and take pictures or create a GIF.
- sharing mode: after a user take a picture or create a GIF, it can share it across social media.

Brands can buy advertisement space and time, selecting the videoclips to be displayed and the fragrance to be used at specific times, drawing the customers into their stores. Customers can be captivated by the combination of sensorial stimuli, the gesture-based interaction, the immersive user experience provided by the brands – feeling they belong in a TV advertisement, and the opportunity to gain several benefits, e.g., discount coupons.

## 1.3. Market research

A Digital Outdoor is essentially a traditional outdoor advertising powered up by technology. The pros of a digital outdoor compared to a traditional one is mostly the way that it captivates the attention of consumers in a more dynamic way. It can also change its advertisement according to certain conditions, such as weather and/or time. Some researches tell that the British public sees over 1.1 Billions (BN) digital outdoor advertisements over a week [4], which can tell how much digital marketing is valued nowadays.

When talking about numbers, “At the end of 2020, despite the COVID wipe-out, the Digital Out-Of-Home (DOOH) market was estimated to be worth \$41.06 BN, but by 2026, nearly two out of three (65%) advertising executives predict this will rise to between \$50 BN and \$55 BN. A further 16% expect it to be worth between \$55 BN and \$60 BN, and 14% estimate it will be even bigger” [5].

#### 1.4. Project goals

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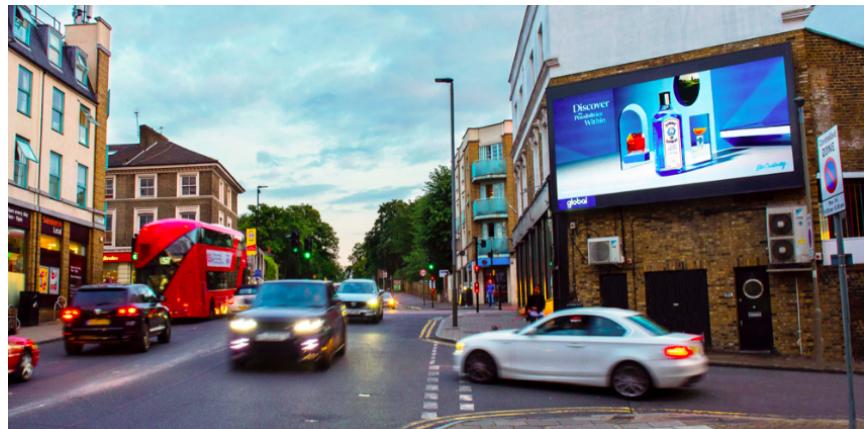


Figure 1.1.: Example of a Digital Outdoor, withdrawn from [4]

Scent market is the art of taking a company's brand identity, marketing messages, target audience and creating a scent that amplifies these values. That's because "a scent has the ability to influence behavior and trigger memories almost instantaneously. When smell is combined with other marketing cues, it can amplify a brand experience and establish a long lasting connection with consumers" [6].

Ambient scent uses fragrance to enhance the experience of consumers with different purposes, whereas scents in scent branding are unique to each company's identity. According to a Samsung study: "when consumers were exposed to a company scent, shopping time was increased by 26% and they visited three times more product categories" [7]. Also, "the digital scent technology market is expected to grow from \$1.0 BN in 2021 to \$1.5 BN by 2026, at a Compound Annual Growth Rate (CAGR) of 9.2%." [8].

The market growth can be attributed to several factors, such as expanding application and advancements in e-nose technologies, increasing use of e-nose devices for disease diagnostic applications, emerging Research and Development (R&D) activities to invent e-nose to sniff out COVID-19, and rising use of e-nose in food industry for quality assurance in production, storage, and display.

## 1.4. Project goals

The project aims to develop a CPS for multi-sensory marketing with contactless user interaction. The key goals identified and the respective path to attain them are:

1. devise a device with audio and video outputs, as well as fragrance diffusion: understand audio and video streaming and study fragrance nebulizer technologies.
2. create a contactless user interface based on gestures through computer vision: identify user gestures through computer vision and match them to interface callbacks; a virtual keyboard may be required

## 1.5. Project planning



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Figure 1.2.: Attractive Opportunities in the Digital Scent Technology Market, withdrawn from [8]

for user input.

3. devise a distributed architecture to convey brand advertisement information to the local device: understand distributed architectures and apply them for optimal data flow; create a remote client-server model to convey information from the brands to the device through remote cloud database services; devise adequate data frames to convey information to the local device; create a local server to respond to the remote server requests.
4. apply facial detection to the camera feed and subsequently apply image filters specific to each brand: understand facial detection algorithms and apply them to the camera feed; apply image filters on top of the identified faces through a specialized Application Programming Interface (API).
5. enable image and GIFs sharing to social media for increased brand awareness: understand how to use social media APIs for media sharing.

## 1.5. Project planning

In Appendix A is illustrated the Gantt chart for the project (Fig. A.1), containing the tasks' descriptions. It should be noted that the project follows the Waterfall project methodology, which is meant to be iterative.

The tasks are described as follows:

- Project planning: in the project planning, a brainstorming about conceivable devices takes place, whose viability is then assessed, resulting in the problem statement (Milestone 0). A market research

## 1.6. Report Outline

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is performed to assess the product's market space and opportunities. Finally, an initial version of the project planning is conceived to define a feasible timeline for the suggested tasks.

- Analysis: in this phase an overview of the system is conceived, presenting a global picture of the problem and a viable solution. The requirements and constraints are elicited, defining the required features and environmental restrictions on the solution. The system architecture is then derived and subsequently decomposed into subsystems to ease the development, consisting of the events, use cases, dynamic operation of the system and the flow of events throughout the system. Finally, the theoretical foundations for the project development are presented.
- Design: at this stage the analysis specification is reviewed, and the HW and SW and the respective interfaces are fully specified. The HW specification yields the respective document, enabling the component selection, preferably Commercial off-the-shelf (COTS), and shipping. The SW specification is separately performed in the subsystems identified, yielding the SW specifications documentation (milestone).
- Implementation: product implementation which is done by modular integration. The HW is tested and the SW is implemented in the target platforms, yielding the SW source code as a deliverable (milestone). The designed HW circuits are then tested in breadboards for verification and the corresponding Printed Circuit Board (PCB) is designed, manufactured and assembled. After designing the PCB, the enclosure is designed to accommodate all HW components, manufactured and assembled. Lastly, the system configuration is performed, yielding prototype alpha of the product.
- Tests: modular tests and integrated tests are performed regarding the HW and SW components and a functional testing is conducted.
- Functional Verification/Validation: System verification is conducted to validate overall function. Regarding validation, it is conducted by an external agent, where a user should try to interact with the designed prototype.
- Documentation: throughout the project the several phases will be documented, comprising several milestones, namely: problem statement; analysis; design; implementation; and final.

## 1.6. Report Outline

This report is organized as follows:

- In Chapter 1 is presented the project's context and motivation, the problem statement, the market research, the project goals, and project planning.
- In Chapter 2, the product requirements are derived — defining the client expectations for the product

## **1.6. Report Outline**

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— as well as the project constraints — what the environments limits about the product. Based on the set of requirements and constraints, a system overview is produced, capturing the main features and interactions with the system, as well as its key components. Then, the system architecture is devised, comprising both hardware and software domains. Next, the system is decomposed into subsystems, presenting a deeper analysis over it, comprising its user mock-ups, events, use cases diagram, dynamic operation and flow of events. A budget estimation is conducted to evaluate the project's costs for both the scale-model and real-scale prototypes.

- Chapter 3 lays out the theoretical foundations for project development, namely the project development methodologies and associated tools, concurrency, the communications technologies and client–server architecture, daemons and device drivers, fragrance diffusion technologies, computer vision for facial detection and gesture recognition, database management, motion detection, camera recording, image filtering APIs, GIF generation, social media APIs, the User Interface (UI) framework, and the file transfer protocols.
- Lastly, the appendices (see Section 3.17) contain detailed information about project planning and development.

## **2. Analysis**

In the analysis phase, the product requirements are derived — defining the client expectations for the product — as well as the project constraints — what the environment limits about the product. Based on the set of requirements and constraints, a system overview is produced, capturing the main features and interactions with the system, as well as its key components.

Then, the system architecture is devised, comprising both hardware and software domains. Next, the system is decomposed into subsystems, presenting a deeper analysis over it, comprising its user mock-ups, events, use cases diagram, dynamic operation and flow of events.

### **2.1. Requirements and Constraints**

The development requirements are divided into functional and non-functional if they pertain to main functionality or secondary one, respectively. Additionally, the constraints of the project are classified as technical or non-technical.

#### **2.1.1. Functional requirements**

- Advertising through a screen and speakers;
- Have fragrance diffusion;
- Take pictures and GIFs;
- Detect a user in range of the device;
- Contactless user interaction through gesture recognition;
- Camera feed and facial detection;
- Apply brand-specific image filters;
- Enable sharing multimedia across social media;
- Provide a remote user interface for brands to purchase and configure the advertisements;
- Provide a remote user interface for company staff to monitor and control the MDO local system.

### **2.1.2. Non-functional requirements**

- Low power consumption;
- Provide user-friendly interfaces;
- Have low latency between local system and remote server;
- Use wireless communication between the local and remote systems.

### **2.1.3. Technical constraints**

- Use device drivers;
- Use Makefiles;
- Use C/C++;
- Use Raspberry Pi as the development board;
- Use compatible HW with the development board;
- Use buildroot;
- Social media APIs for sharing multimedia
- Image filtering through specialized APIs.

### **2.1.4. Non-technical constraints**

- Project duration: one semester (circa 20 weeks);
- Pair work flow;
- Limited budget;
- Scale model prototype.

## **2.2. System overview**

The system overview presents a global view of the system, considering its main features, components and interactions. It is not intended to be complete, but rather provide a basis for the outline of the system architecture. Fig. 2.1 presents the MDO system overview.

Considering the system interactions, three main actors were identified:

1. Brand: represents the brands contracting the advertisement services;
2. Administrator: the development company staff, which can monitor and control the outdoor (administrative privileges).

## 2.2. System overview

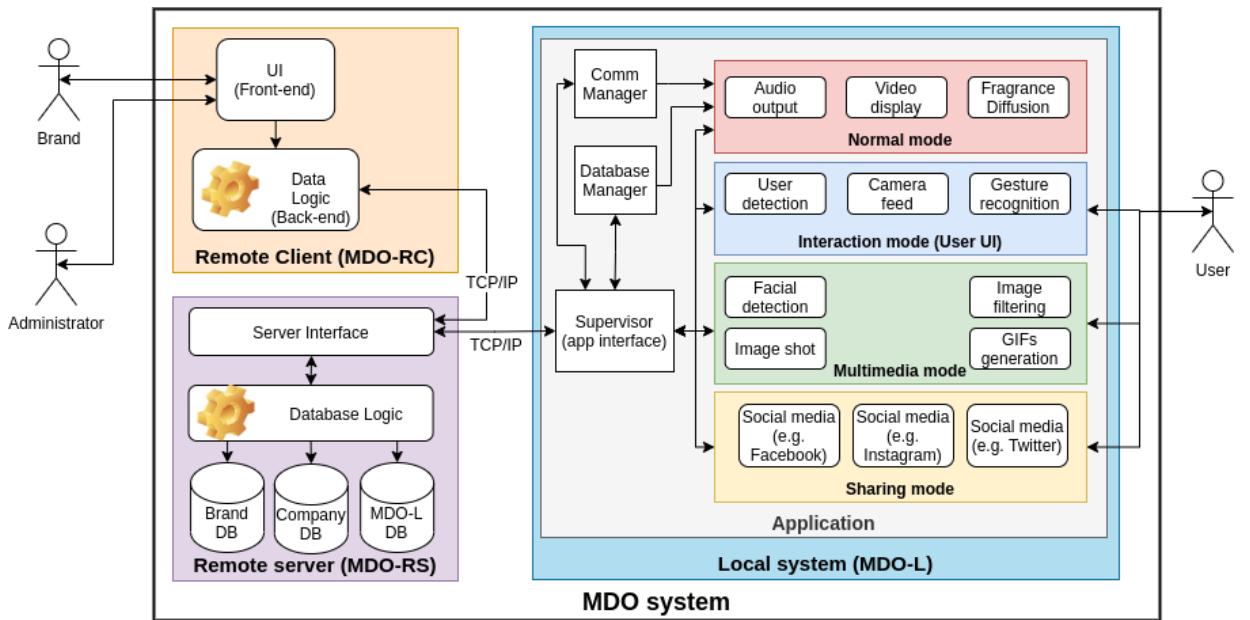


Figure 2.1.: MDO system overview

3. User: the user (the target audience of the advertisement) interacting with the system.

Considering the data flow across the **MDO system**, three main subsystems were identified: **MDO Remote Client (MDO-RC)**, **MDO Remote Server (MDO-RS)**, and **MDO Local System (MDO-L)**.

The rational behind this initial decomposition is explained next.

### 2.2.1. MDO Remote Client

The Brand and Administrator members require a remote UI (front-end) to interact with the system: the former to configure the advertisements being displayed at the MDO and purchase them; the latter to remotely monitor and control the operation of the MDO. Thus, it is clear that an authentication mechanism must be provided for the remote UI.

The data is then dispatched to the back-end, where it is processed and feed back to the UI user and/or sent to the remote server, via Transmission Control Protocol/Internet Protocol (TCP/IP) comprising the data logic component of the UI.

### 2.2.2. MDO Remote Server

Although the MDO-RC could communicate directly with the MDO-L, this is not desirable or a good architecture mainly due to: communications failure could result in data loss, compromising the system's integrity; the

## 2.2. System overview

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remote client and the local system become tightly coupled, meaning the remote client must be aware of all the available local systems; if the data storage in the local system fails, the remote client would have to provide the backup information.

Thus, a remote server component is included, providing the access and management of the system databases, pertaining to the Brand, Company, and MDO Local system. The first two provide the historical information of the **Brand** and **Administrator** entities, and the last one the information related to all of the **MDO-L** systems in operation.

The main functions of the **MDO-RS** are:

- UI requests responses: when a UI user requests/modifies some information from the database, the server must provide/update it.
- MDO-L monitoring and control: provide command dispatch and feedback to the **Administrator** staff for remote monitoring and control of the device.
- MDO-L update: periodically check for start times of each MDO-L device and transfer the relevant data to it.

The server interface is the responsible for managing the requests and respective responses from the remote client and for periodically send the update data to all MDO-L devices.

### 2.2.3. MDO Local system

The MDO local system (MDO-L) is the marketing device, interacting with the user to display multi-sensory advertisements. As aforementioned in Section 1.2, it is comprised of four modes:

- normal mode: the MDO provides sound, video and fragrance outputs. It is the default mode.
- interaction mode: When a user approaches the device, the MDO will go into interaction mode, turning on and displaying the camera feed and waiting for recognizable gestures to provide additional functionalities, such as brand-specific image filters. This is the **User UI**.
- multimedia mode: in this mode the facial detection is applied, enabling the user to select and apply different brand-specific image filters and take pictures or create a GIF.
- sharing mode: after a user take a picture or create a GIF, it can share it across social media.

The user interaction is considered to be a higher priority activity than the advertisements, so when a User interacts with the system, the **normal mode** is overriden by the **Interaction mode**, thus, halting the advertisements.

The MDO-L application communicates with the remote server (**MDO-RS**) through the **Supervisor** via TCP/IP to handle requests from **Administrator** members to monitor and control the device through the **Supervisor** or to update the advertisements. Additionally, the **Supervisor** oversees the application mode

### 2.3. System architecture

and the communication (Comm Manager) and database (Database manager) managers to handle system events.

## 2.3. System architecture

In this section, the system architecture is devised in the HW and SW components, using the system overview as a starting point.

### 2.3.1. Hardware architecture

Fig. 2.2 illustrates an initial hardware big picture that fulfils the system's goals, meeting its requirements and constraints. As it can bee seen, the diagram is divided in four distinct parts: **External Environment**, **Local System**, **Remote Server** and **Remote Client**.

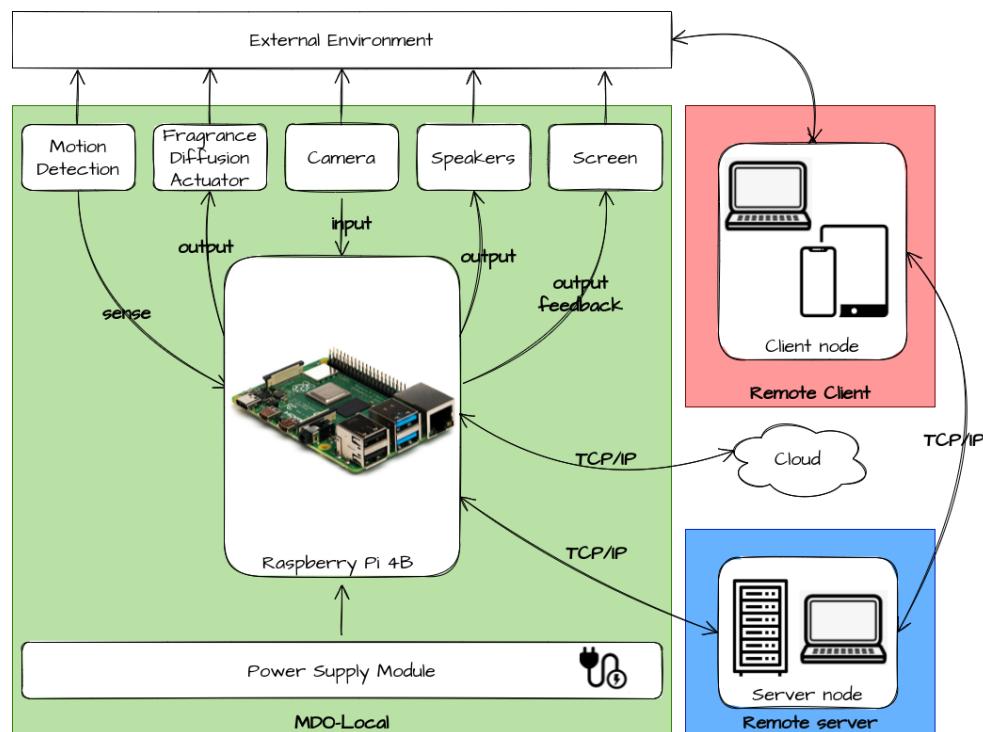


Figure 2.2.: HW architecture diagram

Firstly, the **External Environment** represents all the environment that interacts with the system. In this case, these are all its users — normal users, brands and company staff (Administrator role).

### 2.3. System architecture

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Secondly, the **Local System** is composed of the main controller, which is the Raspberry Pi 4B. This System-on-a-Chip (SoC) is responsible to control all the **Local System** and to establish a connection with the **Remote Server** through its included WiFi module. Additionally, the **Local system** also communicates with the Cloud to share contents on social media, and, potentially, to access image filtering APIs. The **Local System** is powered through a Alternating Current (AC)-Direct Current (DC) power converter, and, potentially, a step-down converter – **Power Supply Module**. The main board has several blocks connected to it:

- Motion Detection: used to detect the users and switch from normal mode to interaction mode;
- Fragrance Diffusion Actuator: used to diffuse the fragrance into the air;
- Camera: used to capture image that is then processed;
- Speakers: used to reproduce advertisements sounds;
- Screen: used to display video clips of advertisements.

In third place, the **Remote Server** has a server node running in another machine that can be one computer or a main frame. The remote server stores all databases which the **Remote Client** and **Local System** may need to access and serves as a proxy server to enable the **Admin** users to control and monitor the **Local System**.

Lastly, the **Remote Client** runs the MDO management application, which can be deploy to a computer (like the Raspberry Pi), a tablet or a smartphone.

#### 2.3.2. Software architecture

In this section the SW architecture for MDO-RC, MDO-RS, and MDO-L subsystems is presented, defining its SW stack.

##### **MDO remote client**

Fig. 2.3 illustrates the SW architecture for the remote client, representing its SW stack. It is comprised of the following layers:

- Application: contains the remote client application. The **Brand** and **Admin** members interact with the **UI**, which is the visual part of the interface. The **UI engine** is notified and handles all UI events – internal or external – providing the **UI** with feedback for its users. The relevant commands are then parsed – **Parser component** – and responded. The commands are then translated to the appropriate DB queries and responded through the **DB Manager**. The **Comm Manager** is responsible for encapsulating the DB queries into the respective TCP/IP frames to be sent to the **Remote Server** as well as unwrap the incoming server responses.

### 2.3. System architecture

- Middleware: contains the TCP/IP framework supporting these communication protocols as part of OSI model for internet applications. It manages the incoming/outgoing TCP/IP frames by providing the adequate protocol handshaking and queueing and timing aspects of the bytes to send/receive.
- OS & BSP – Operating System (OS) & Board Support Package (BSP): it contains the low-level and communication drivers required to handle input (keyboard/touch), output (screen) and communication to the Remote Server.

It should be noted that for desktop and mobile applications, the **Middleware** and **OS & BSP** layers are usually abstracted by the OS, thus, the relevant APIs should be used.

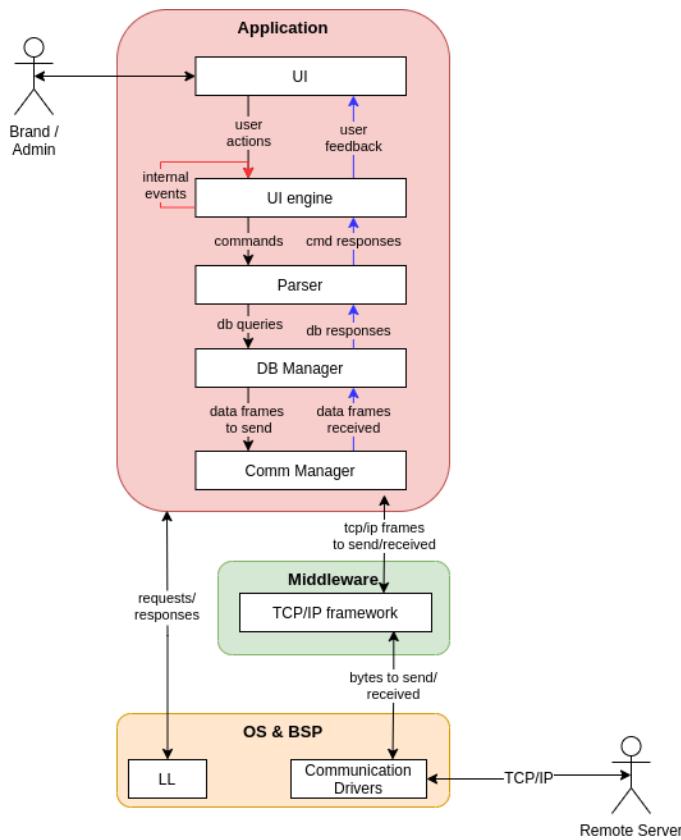


Figure 2.3.: SW architecture diagram: remote client

### MDO remote server

Fig. 2.4 illustrates the SW stack for architecture for the remote server. It is comprised of the following layers:

- Application: contains the remote server application. It provides a Command Line Interface (CLI) to handle **Remote client** requests. The CLI engine is notified and handles all UI events — internal or external — providing the appropriate feedback. The relevant commands are then parsed — **Parser**

### 2.3. System architecture

component — and responded: DB queries are handled by the **Relational Database Management System (RDBMS)** issuing DB transactions; other commands received from the **Remote Client** are handled internally and translated, being dispatched to the **Local System** by the **Comm Manager** (via **Communication drivers**). Internal events can also trigger the **RDBMS** to issue database transactions for the **Remote Client** or **Local System**. The **Comm Manager** is responsible for wrapping/unwrapping the data frames received by or sent to the **Remote Client** or **Local System**.

- **Middleware:** contains the RDBMS framework supporting the management of the relational databases using database transactions.
- **OS & BSP** – OS & BSP: it contains the **Communication drivers** to handle requests from the **Remote Client**, and the **File I/O** drivers to manipulate DB transactions from/to storage.

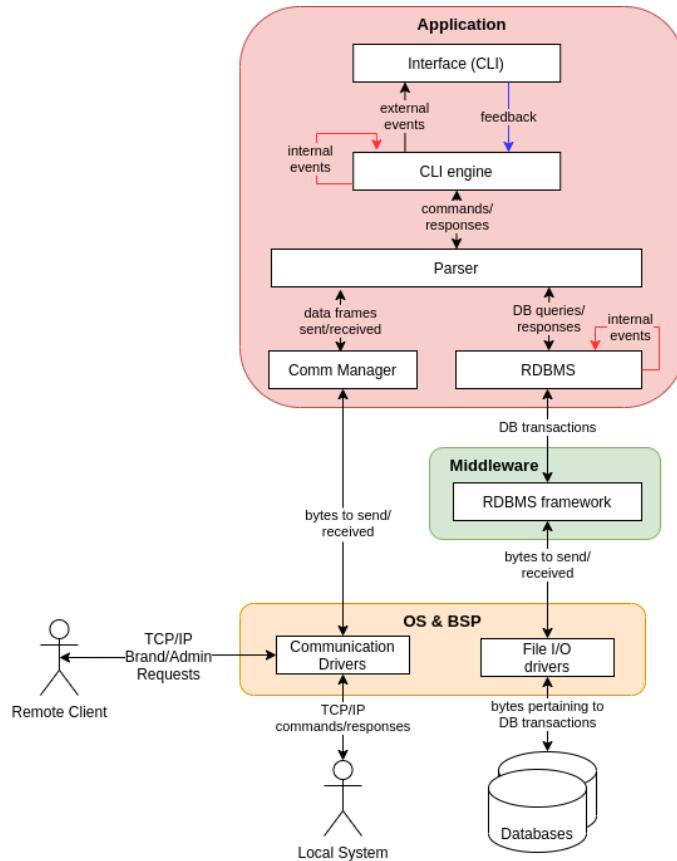


Figure 2.4.: SW architecture diagram: remote server

It should be noted that the **Remote Server** main functions are:

- provide relational databases for easier management of all entities and respective data in the system;
- decompose the relationship many-to-many, between the remote clients and local systems — many remote clients may want to connect to different local systems;

#### 2.4. Subsystem decomposition

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- decouple the architecture as the Remote Client should not know the Internet Protocol (IP) address of every local system it may potentially try to access, acting as a proxy server.

#### MDO local system

Fig. 2.5 illustrates the SW stack for architecture for the Local System. It is comprised of the following layers:

- Application: contains the local system application. It provides a UI to handle User interaction. The Interface engine is notified and handles all UI events — internal or external — through gesture recognition, providing the appropriate feedback. The relevant commands are then parsed — Supervisor component — and responded: DB queries are handled by the Database manager issuing DB transactions for internal databases; commands received from the Remote Server to monitor or control the system are handled internally and responded back by the Comm Manager (via Communication drivers); mode management is performed. Internal events can also trigger the Database manager to issue database transactions to update the Local System. The Comm Manager is responsible for wrapping/unwrapping the data frames received by or sent to the Remote Server.
- Middleware: contains: the DB framework supporting the management of the internal databases using database transactions; the Computer Vision (CV) framework that handles gesture and facial detection; image filtering and GIF frameworks for multimedia; social media framework.
- OS & BSP – OS & BSP: contains: the Communication drivers to handle requests from the Remote Server and for social media sharing, and, potentially the API calls to cloud-based image filtering frameworks, depending on the application profiling; the File I/O drivers to manipulate internal DB transactions from/to storage; audio, video and fragrance diffuser actuator drivers for normal mode; the camera driver for camera feed; the detection sensor driver to signal a User is in range, triggering the switch from normal mode to interaction mode.

The Local system is a soft real-time system, as no mandatory deadlines must be met.

## 2.4. Subsystem decomposition

In this section the system is decomposed into subsystems and, for each subsystem, a more detailed analysis is performed yielding its user mock-ups, events, use case diagram, dynamic operation and the flow of events throughout the subsystem.

As aforementioned, the subsystems identified are: Remote Client (MDO-RC), Remote Server (MDO-RS), and Local System (MDO-L).

## 2.4. Subsystem decomposition

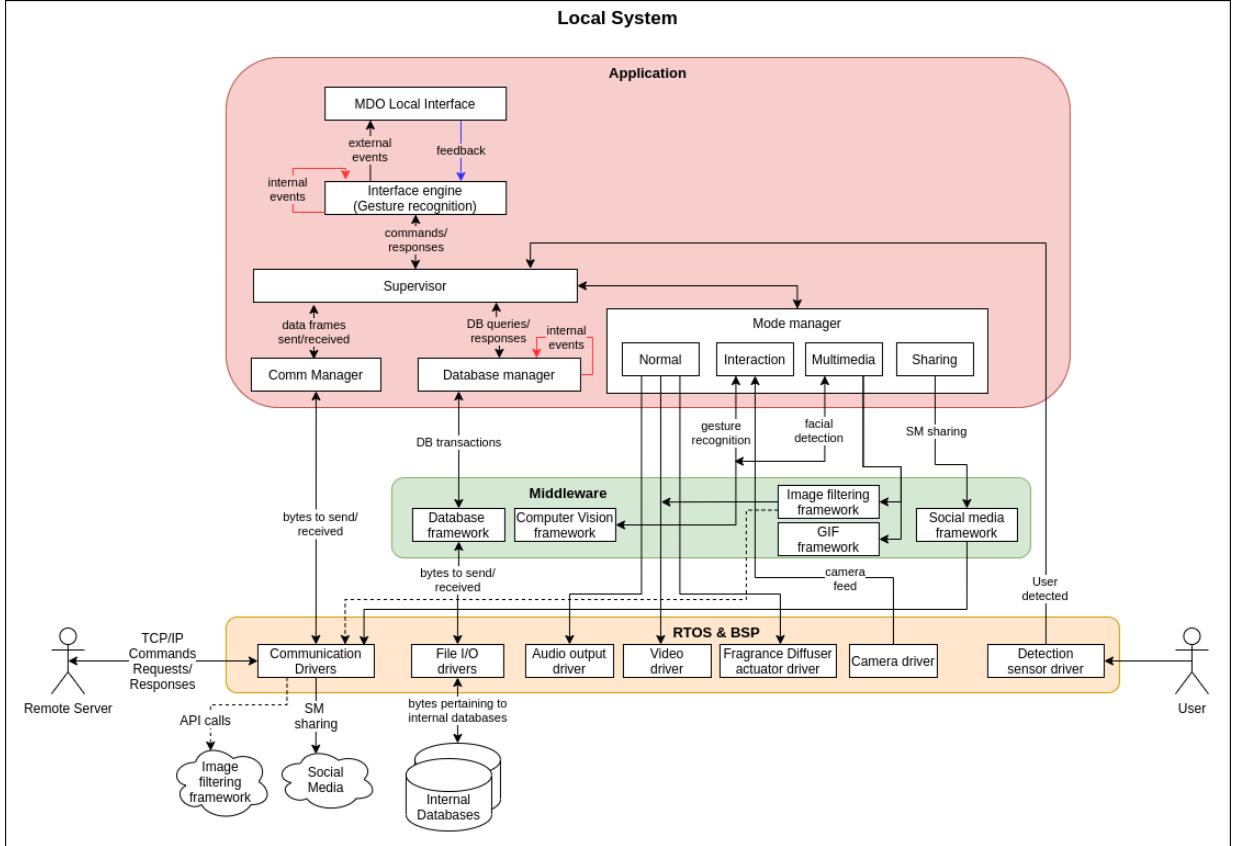


Figure 2.5.: SW architecture diagram: local system

### 2.4.1. Remote Client

In this section the remote client is analyzed, considering its events, use cases, dynamic operation and the flow of events.

#### User mock-ups

In Fig. 2.6 is illustrated the user mock-ups for the remote client. It intends to clarify how does the UI works for the two actors: Brands and Company (staff).

The initial state of the MDO-RC's UI is depicted in thick border outline: the 'Sign In' window. If the User makes a mistake in its username and/or password, it will be shown an error message. Also, the 'Sign In' window has an option to recover the password, triggering the dispatch of an e-mail. If the User still remembers its credentials, the app flows through one out of two possibilities: if the user is an admin, goes to the admin main menu, otherwise if the user is a brand, it will appear the brand main menu.

Firstly, the Admin workflow:

## 2.4. Subsystem decomposition

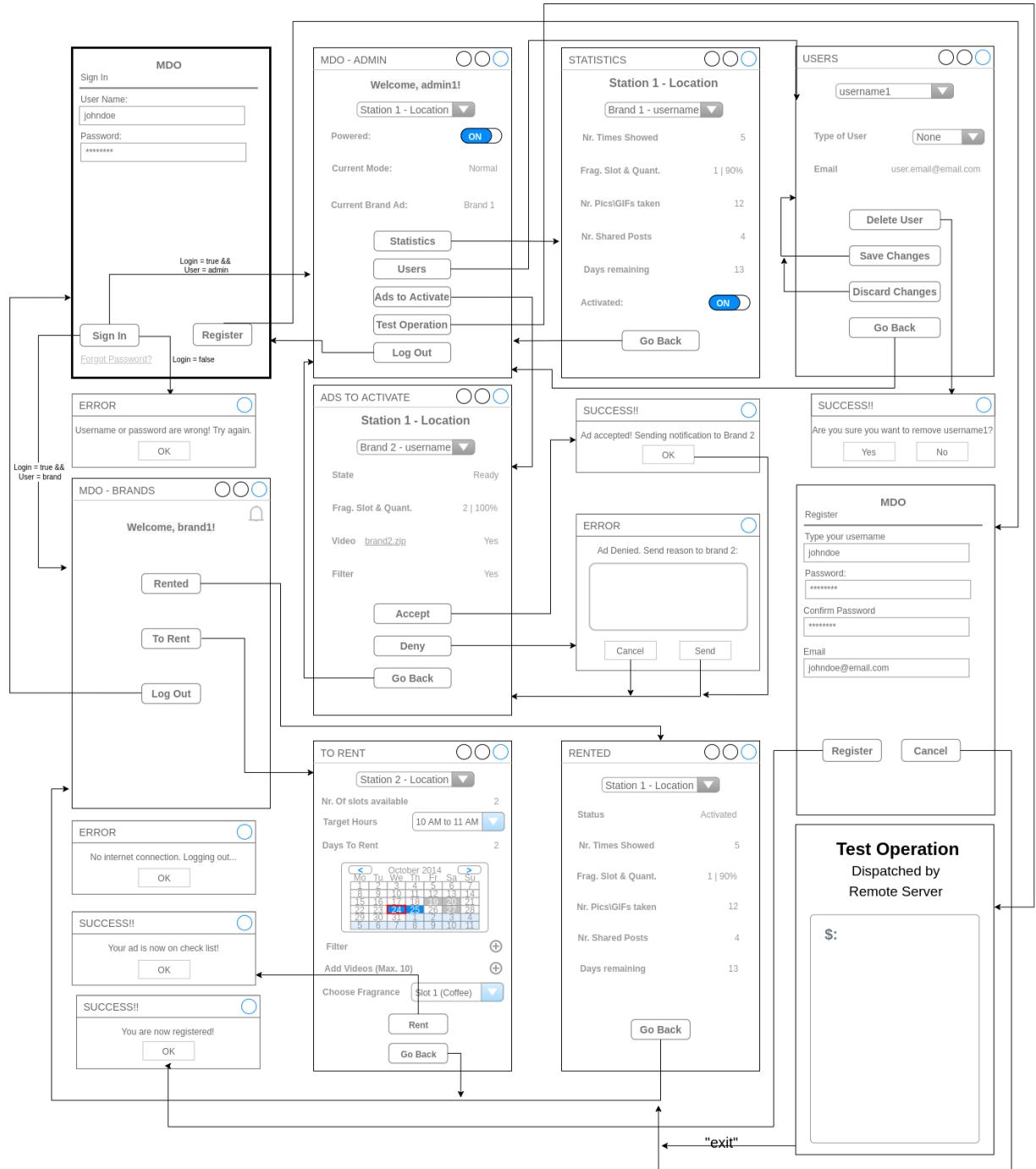


Figure 2.6.: User mockups: remote client

- The Admin main menu contains a drop down button with all the available stations. Choosing one of them, the Admin can turn it On/Off, see it's current mode and the current brand ad being displayed. Also, the Admin can log out and choose between two different paths:

## 2.4. Subsystem decomposition

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- Statistics: It is possible to see various statistics of all different brands that are currently playing on the station: the number of times that the ad was shown, the number of pictures/GIFs and shared posts, the fragrance slot and quantity (percentage) and the days remaining for the rent to end. It is also possible to deactivate the advertisement if something wrong occurs and go back to the previous menu.
- Users: In this window, the **Admin** can manage all users and see their information, changing their type or deleting them from the database.
- Ads to Activate: In this window, the **Admin** can handle all the ads that the brands are intending to rent. For that, the **Admin** needs to validate the ads': verify video's content (checking if all videos are appropriate), if it has a filter, a fragrance and a time slot. After that, the **Admin** can either accept or deny the ad. If it accepts the ad, it is shown a success message and the ad is added to the station with its preferences. Otherwise, the **Admin** indicates the denial reason, which is subsequently sent to the **Brand**'s email.

Secondly, the **Brand** workflow:

- The **Brand** main menu contains a welcome message, a notification bell to see if another ad was accepted or denied and three buttons - Rented, To Rent and Log Out. The 'Log Out' button logs the **Brand** out of its account, the other two buttons switch to different widgets:
  - Rented: The **Brand** can see all statistics of all its rented ads on different stations that it rented. The statistics are: status, number of times the ad was shown, the fragrance slot and quantity (percentage), the number of pictures/GIFs taken, the number of shared posts and the number of days remaining to end its rent.
  - To Rent: The **Brand** can rent ads in the same station or in other stations. For that, the **Brand** selects the target hours and then a calendar displays the available dates. Then after choosing the days, the **Brand** needs to upload a filter and a compressed multimedia archive with a maximum of ten videos. Finally, the **Brand** needs to select the fragrance and select 'Rent'. After that, a success message will be shown and the ad will enter in a waiting queue for an **Admin** to validate.

It is also possible to register a new user through the 'Register' button. This opens a window to type a username, a password, confirm the password and the e-mail. If everything is in order, the user is created with the default user type of Brand.

Finally, at any time, it can occur the loss of internet connection, which triggers an error message informing the automatic log out of the account.

#### 2.4. Subsystem decomposition

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Table 2.1.: Events: remote client

<b>Event</b>	<b>System response</b>	<b>Source</b>	<b>Type</b>
Login	The system verifies if the user credentials are correct and what type of user is and asks for data from databases	User	Asynchronous
Verify internet connection	Periodically verify internet connection	Remote Client	Synchronous
Statistics	Request to the Remote Server all the information to show statistics from all stations and brands	User (Admin)	Asynchronous
Accept/Deny ad	Send information to the Remote Server if the ad is either accepted or denied and if so, why	User (Admin)	Asynchronous
Power On/Off Station	Send command to Remote Server to Power On/Off a certain station	User (Admin)	Asynchronous
Rented	Request to the Remote Server all the information to show statistics from all stations the brand rented	User (Brand)	Asynchronous
Rent	Send to the Remote Server all the information of rent from the brand, all the videos and the filter	User (Brand)	Asynchronous
Test Operation	The System dispatches the command kine provided by the Remote Server	User (Admin)	Asynchronous
Forgot Password	Send e-mail to the user that has forgotten his password	User	Asynchronous

### Events

Table 2.1 presents the most relevant events for the Local system, categorizing them by their source and synchrony and linking it to the system's intended response.

### Use cases

Fig. 2.7 depicts the use cases diagram for the Remote Client, describing how the system should respond under various conditions to a request from one of the stakeholders to deliver a specific goal.

The Admin and the Brand interact with the Remote Client and this last interacts with the Remote Server to process commands, such as query databases or power on/off machines.

The Admin can Manage the Station, which includes Power On/Off Station, Manage Ads to Activate, Enable/Disable an Ad and test its operation. It can also manage users, removing or modifying them. All these use cases are processed from the Remote Client and are requested to the Remote Server.

## 2.4. Subsystem decomposition

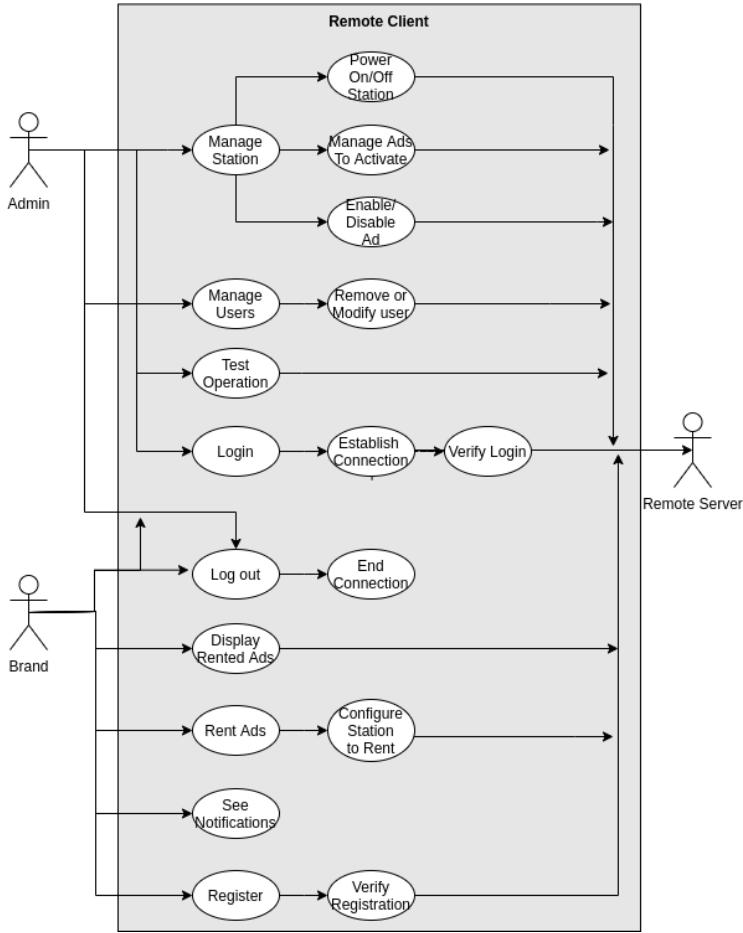


Figure 2.7.: Use cases: remote client

The Brand can see Rented Ads, Rent Ads, See notifications and register. All these cases are also processed from the Remote Client and are requested to the Remote Server.

There are some use cases that are common to the Admin and to the Brand: Login and Logout.

### Dynamic operation

Fig. 2.8 depicts the state machine diagram for the Remote Client, illustrating its dynamic behavior.

There are two main states:

- **Initialization:** the application is initialized. The settings are loaded and if invalid they are restored. The WiFi communication is setup, signaling the communication status and if valid, an IP address is returned.
- **Execution:** after the initialization is successful, the system goes into the **Execution** macro composite state with several concurrent activities, modeled as composite states too. However, it should be noted

## 2.4. Subsystem decomposition

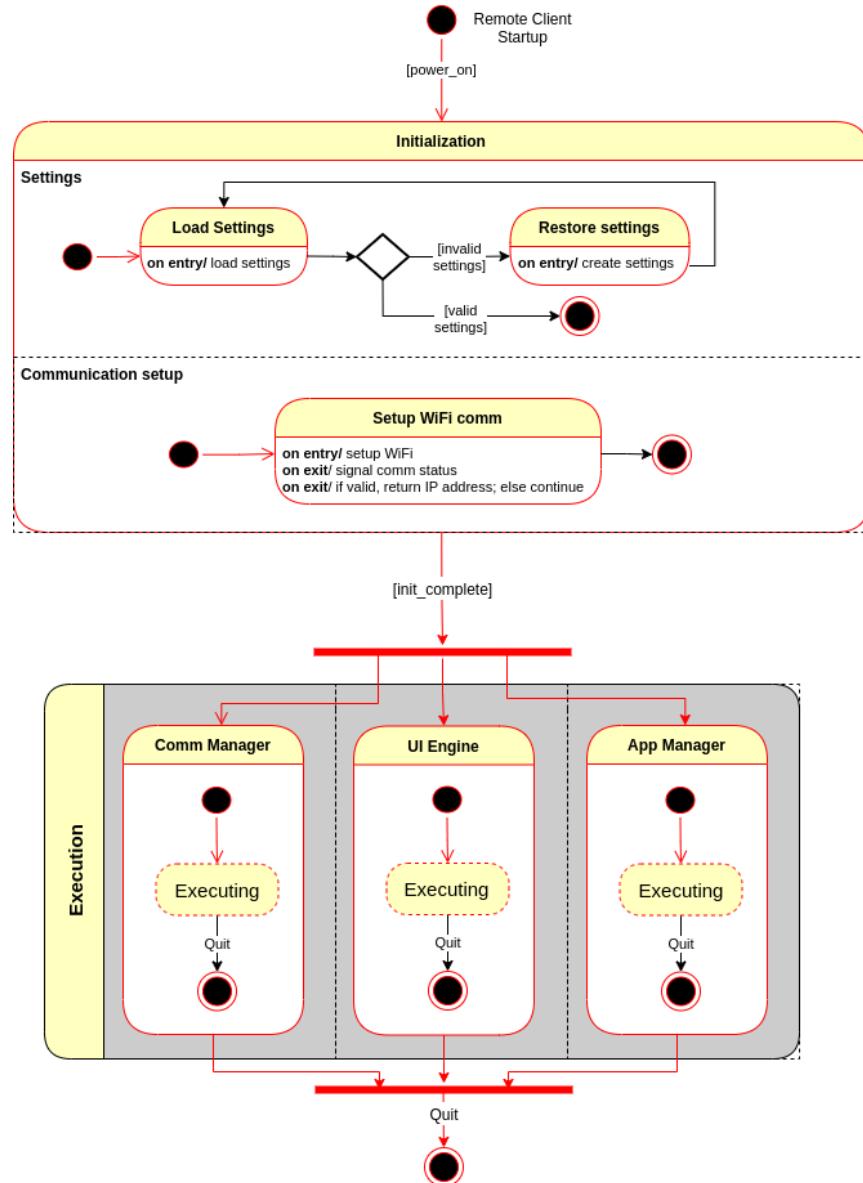


Figure 2.8.: State Machine Diagram: remote client

that there is only one actual state for the device, although at the perceivable time scale they appear to happen simultaneously. These activities are communication management (Comm Manager), interface management (UI Engine) and application manager (App Manager), and are executed forever until system's power off. They are detailed next.

## 2.4. Subsystem decomposition

### Communication Manager

Fig. 2.34 depicts the state machine diagram for the **Comm Manager** component. Upon successful initialization the **Comm Manager** goes to **Idle**, listening for incoming connections. When a remote node tries to connect, it makes a connection request which can be accepted or denied. If the connection is accepted and the node authenticates successfully the **Comm Manager** is ready for bidirectional communication. When a message is received from the remote node, it is written to **TX msg queue** and the **Supervisor** is notified. When a message must be sent to the remote, it is read from the **TX msg queue** and sent to the recipient. If the connection goes down, it is restarted, going into **Idle** state again.

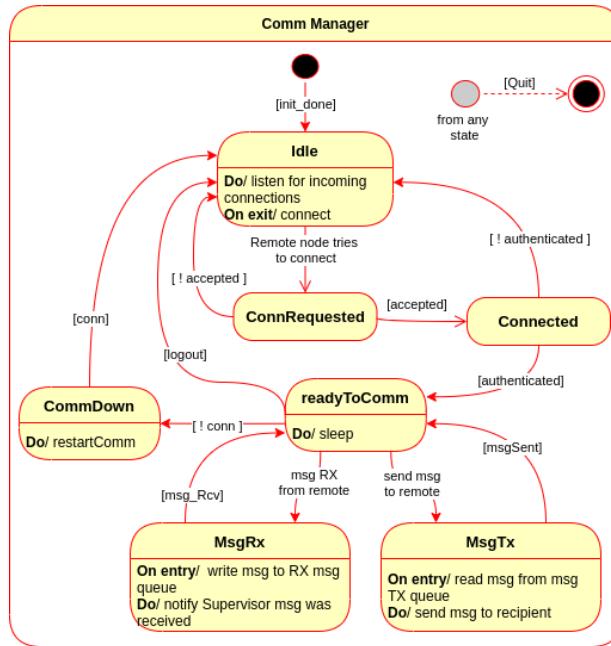


Figure 2.9.: State Machine Diagram: remote client – Communication Manager

### App Manager

Fig. 2.10 depicts the state machine diagram for the **App Manager** component. Upon successful initialization the **App Manager** goes to **Login**, waiting for some action.

A user can register itself by pressing the 'Register' button which leads to **Register** state: if succeeds, it returns to **Login** state. If the 'Login' button is pressed, the system goes to **Validation** state, determining its type:

- **Admin** – **Admin Mode**: the Admin has several can view statistics (**Statistics**), manage all users (**Users**), manage all ads to activate (**Ads To Activate**) and test operations on the machines (**Test Operation**).

## 2.4. Subsystem decomposition

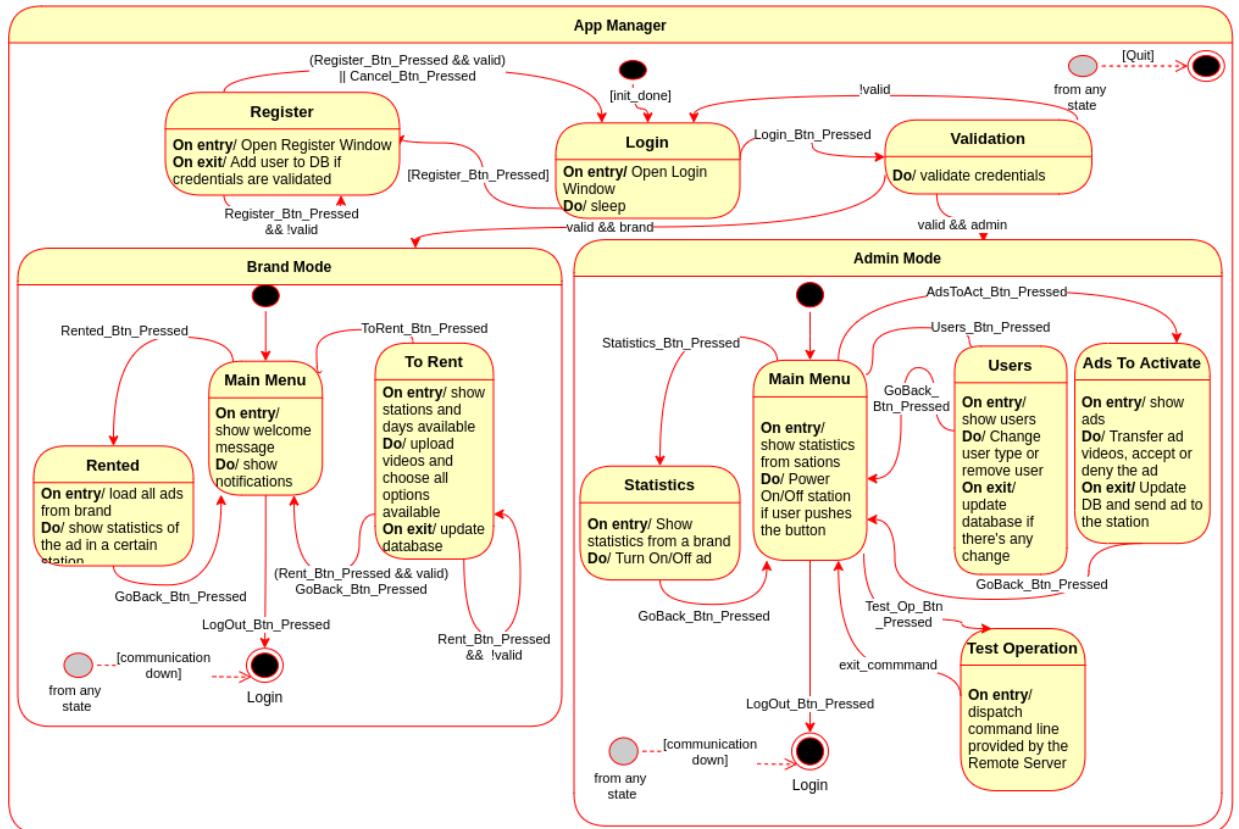


Figure 2.10.: State Machine Diagram: remote client – App Manager

- **Brand – Brand mode:** the Brand can see all its ads (**Rented**), see notifications and messages (**Main Menu**) and rent new ads (**To Rent**).

These two states are terminated by pressing the 'Log Out' button, which redirects to **Login** state.

If, in any state, a critical error occurs, that can cause an unexpected quit of the **App Manager**, leading to the application abnormal shutdown.

### Flow of events

The flow of events throughout the system is described using a sequence diagram, comprising the interactions between the most relevant system's entities. It is usually pictured as the visual representation of an use case. The main sequence diagrams are illustrated next (Fig. 2.11 through Fig. 2.16).

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As it can be seen in Fig. 2.11 to Fig. 2.17, the user interacts with the **UI**, then this last interacts with the **UI Engine**, interacting with the **Remote Client Back-End** in order to process and execute all the information and commands needed. There's an alternate way to go to the user, that's because on the authentication

## 2.4. Subsystem decomposition

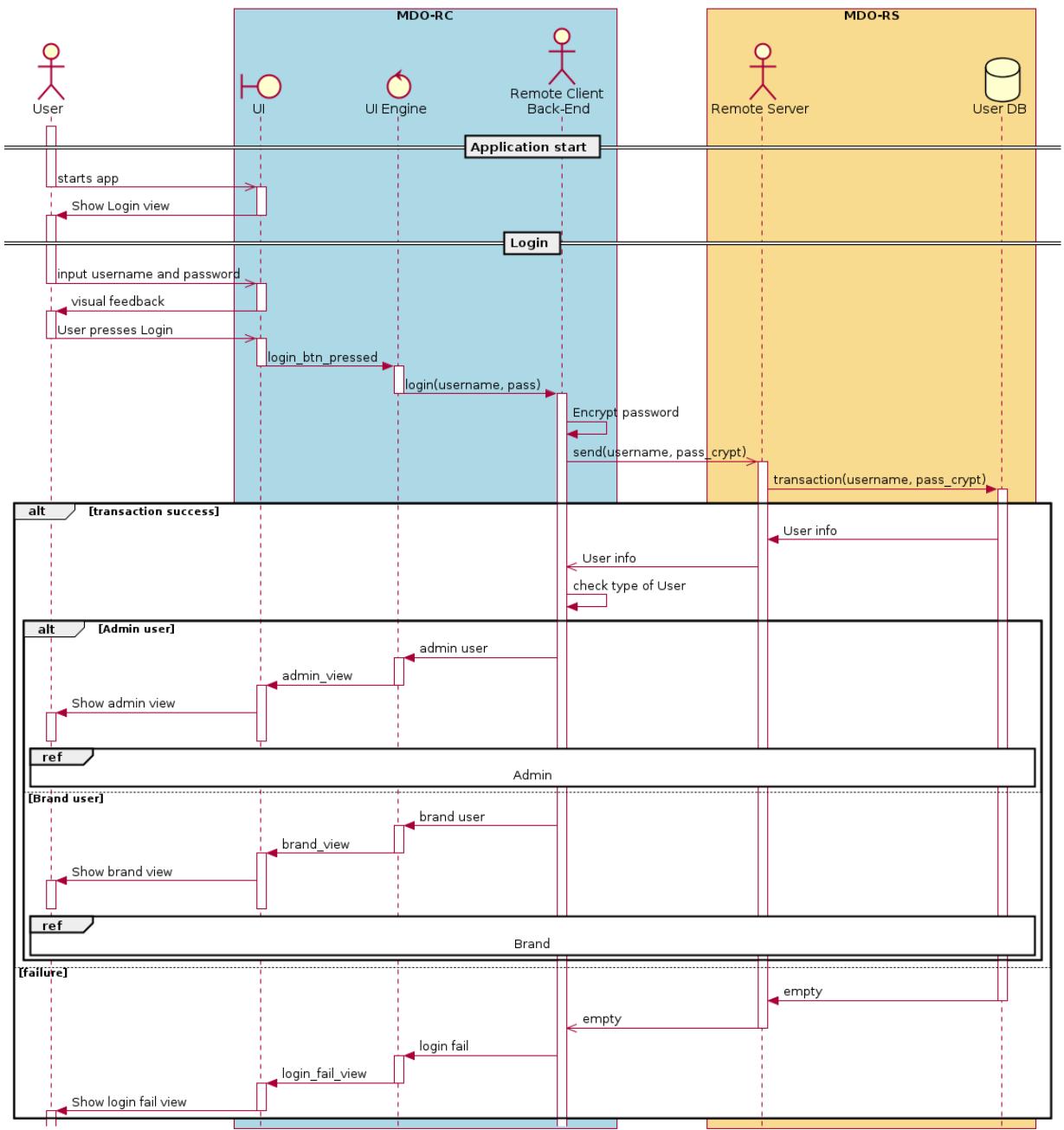


Figure 2.11.: Sequence Diagram: remote client – Login

the **Remote Client Back-End** will discover if the user is an **Admin** or a **Brand**. On both cases, it shows its main menu and it can end the sequence through the 'Logout'. In each one of the cases there's alternative sequences to occur, depending in what the **User** decides to do. Also, in each alternative choice, the **Remote Client** can interact with different **Databases**, either to update them or to ask some info.

## 2.4. Subsystem decomposition

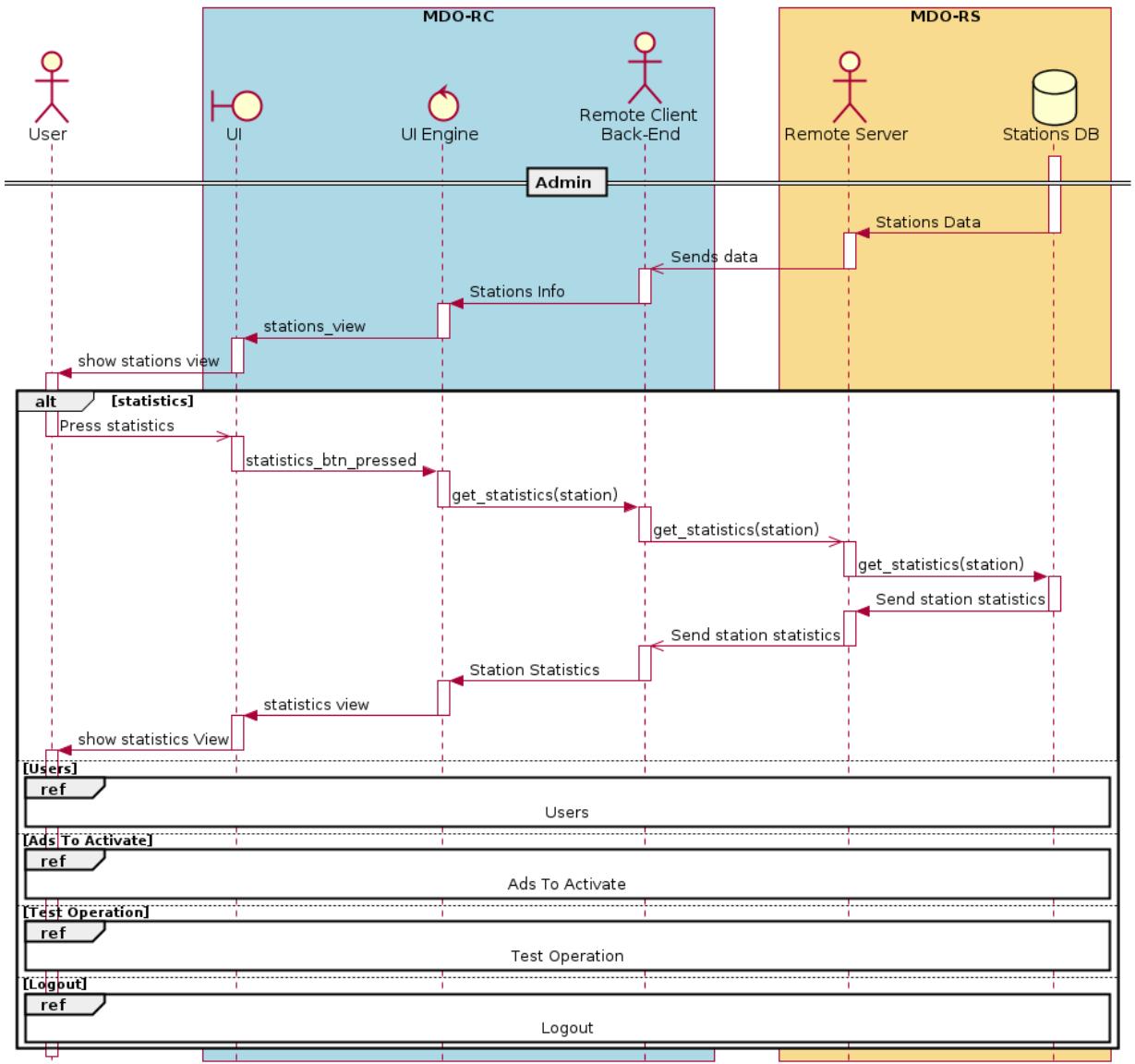


Figure 2.12.: Sequence Diagram: remote client – admin statistics

===== As it can be seen, the user interacts with the UI, whose events are tracked by the UI Engine triggering the appropriate callback and dispatching data to the Remote Client Back-End for adequate processing.

There are two flow paths, pertaining to type of User – Admin or Brand – as a result of the User authentication. On both cases, it shows its main menu and it can end the sequence through the 'Logout'.

In each one of the cases there's alternative sequences to occur, depending of what the User decides to do. Also, in each alternative choice, the Remote Client can interact with different Databases, either to

## 2.4. Subsystem decomposition

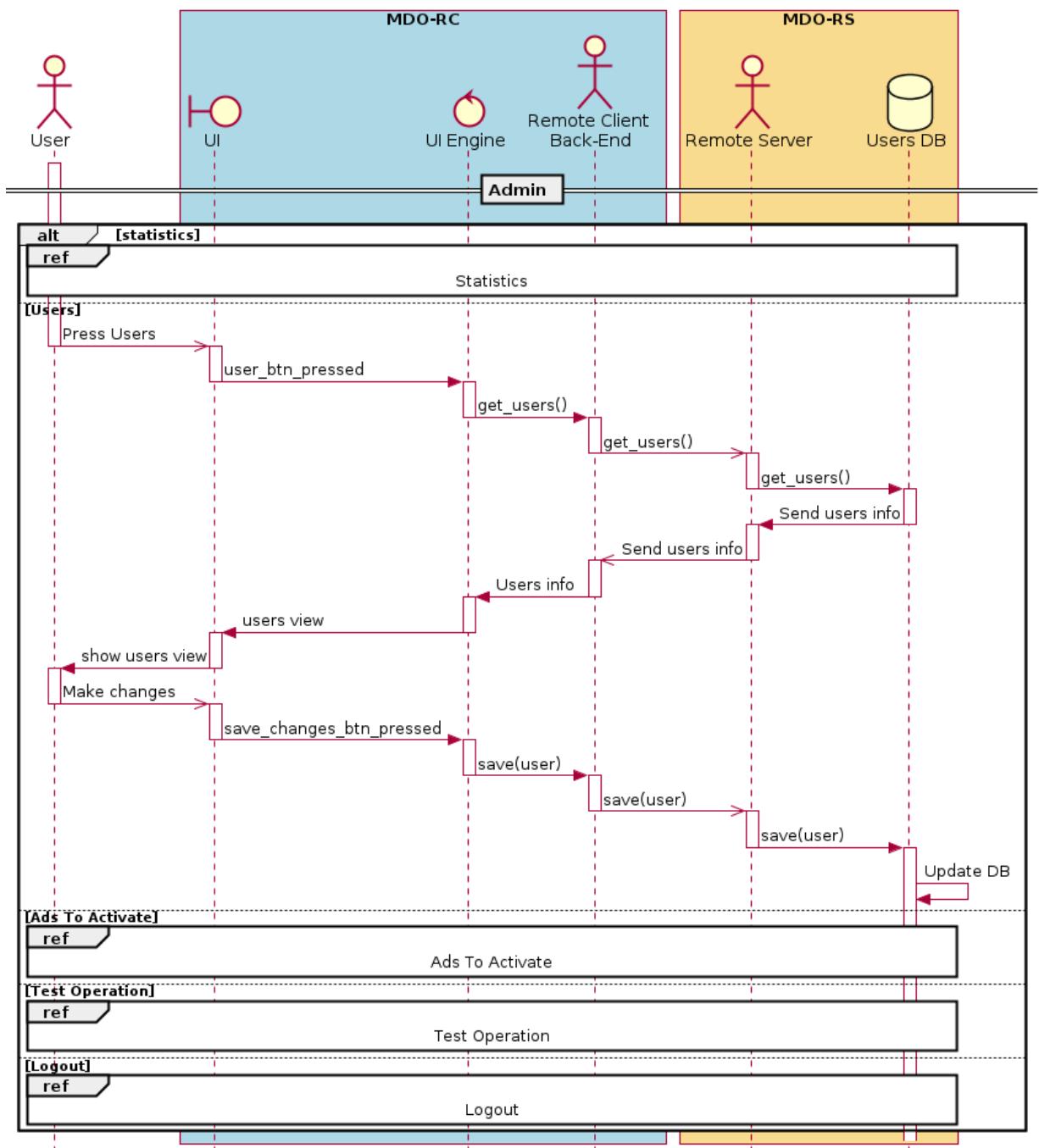


Figure 2.13.: Sequence Diagram: remote client – admin users

query or update them. »»»> 1b64a1fa2e419d79801eaaa806f4d9675f06902e

## 2.4. Subsystem decomposition

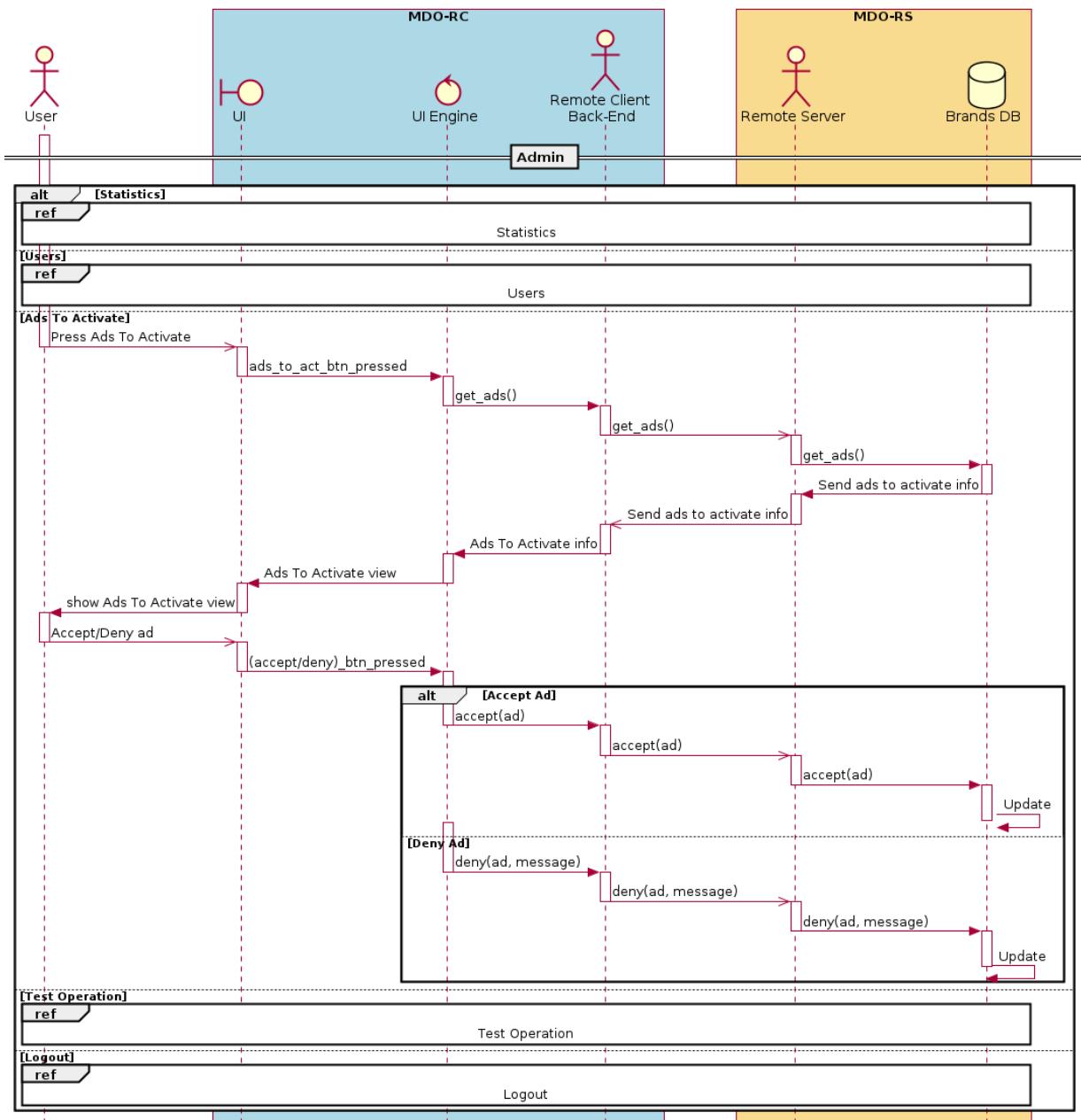


Figure 2.14.: Sequence Diagram: remote client — admin ads to activate

### 2.4.2. Remote server

In this section the remote server is analyzed, considering its events, use cases, dynamic operation and the flow of events.

## 2.4. Subsystem decomposition

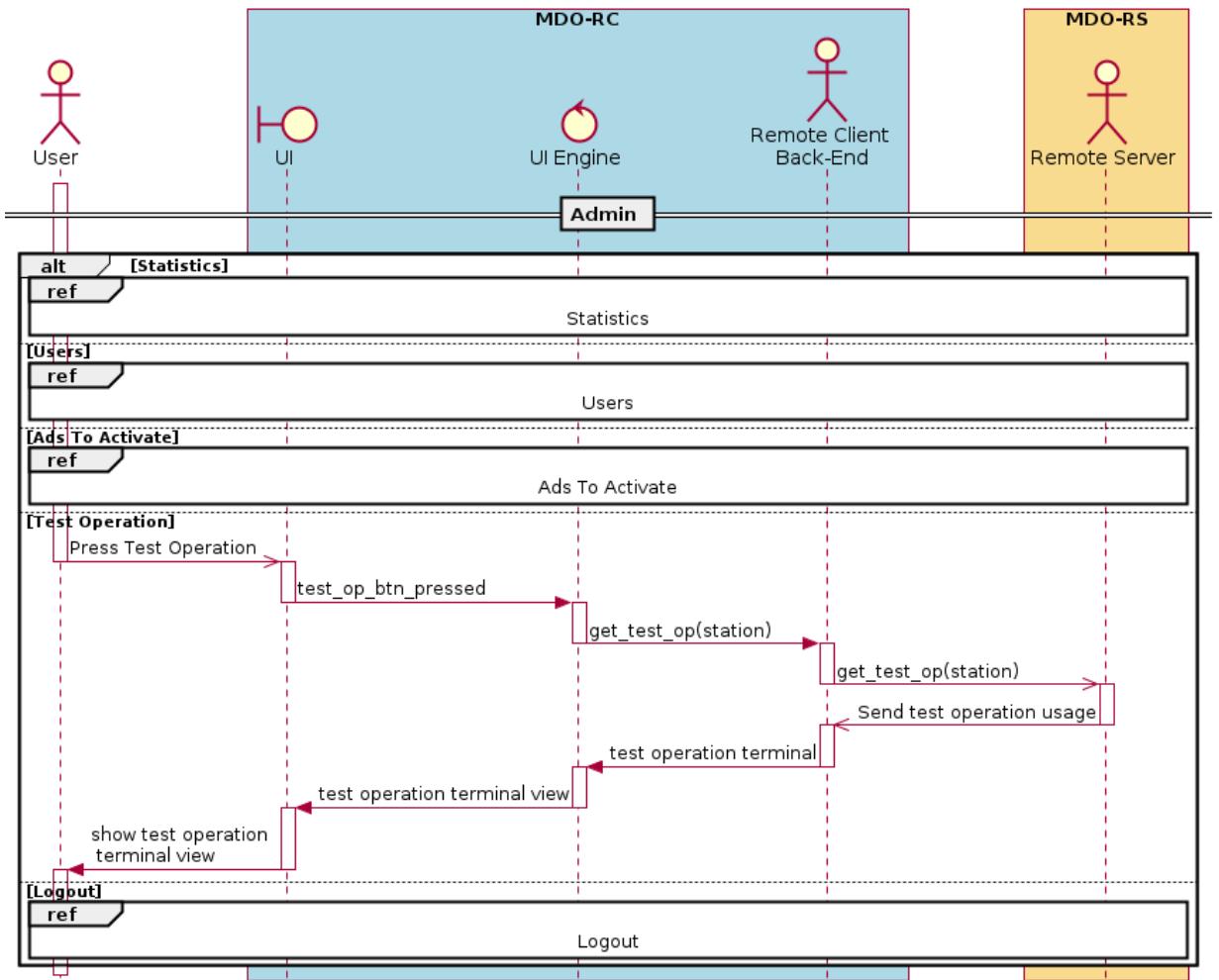


Figure 2.15.: Sequence Diagram: remote client – admin test operation

### User mock-ups

Fig. 2.18 illustrates the user mock-ups for the Remote Server. It intends to mimic the user interaction with the Remote server interface, clarifying the user actions and the respective responses, as well as the workflow, defining the Remote Server interface.

It consists of a CLI providing basic commands to authenticate an user, perform operations over a DB and test the operation of a designated Local System (only available to administrator users).

To test the operation of a Local System, an Admin can:

- Normal mode: add, delete, play or stop video, audio and fragrance;
- Interactive & Multimedia modes — camera: turn on/off the camera, apply facial detection, use an image filter, take a picture or create a GIF;

## 2.4. Subsystem decomposition

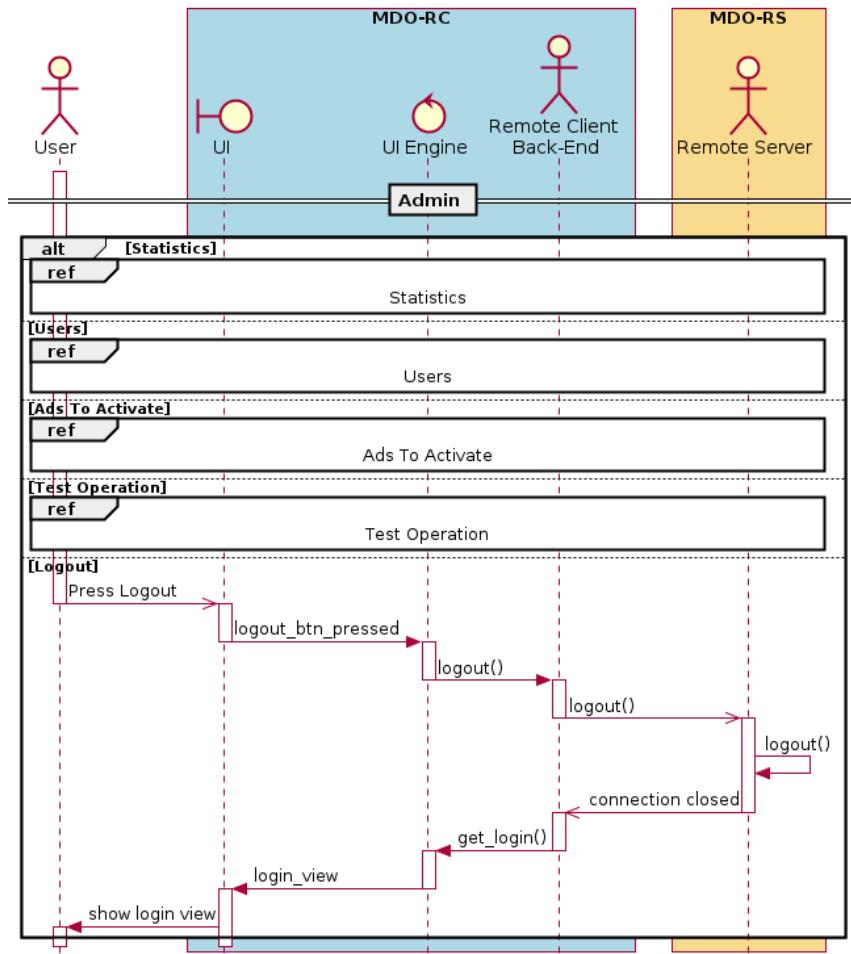


Figure 2.16.: Sequence Diagram: remote client – admin logout

- Sharing mode: share to a designated social media network a post, containing a message and attachment (picture or GIF).

## Events

Table 2.2 presents the most relevant events for the Remote Server, categorizing them by their source and synchrony and linking it to the system's intended response.

## Use cases

Fig. 2.19 depicts the use cases diagram for the Remote Server, describing how the system should respond under various conditions to a request from one of the stakeholders to deliver a specific goal.

## 2.4. Subsystem decomposition

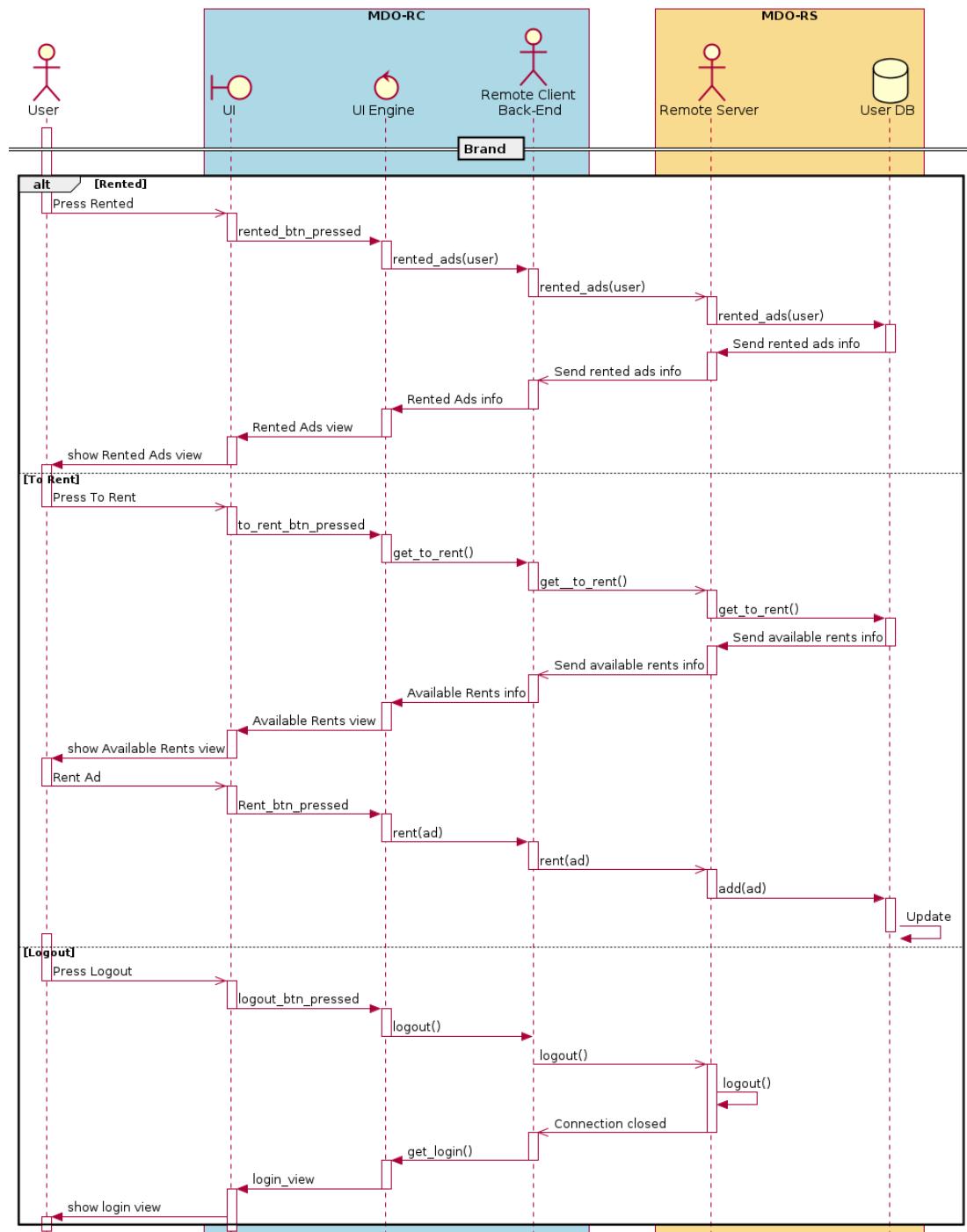


Figure 2.17.: Sequence Diagram: remote client - brand

As it can be seen, the Remote Client can interact through various modes: Help, Authenticate User, Interact with databases, Test Operation and Disconnect.

## 2.4. Subsystem decomposition

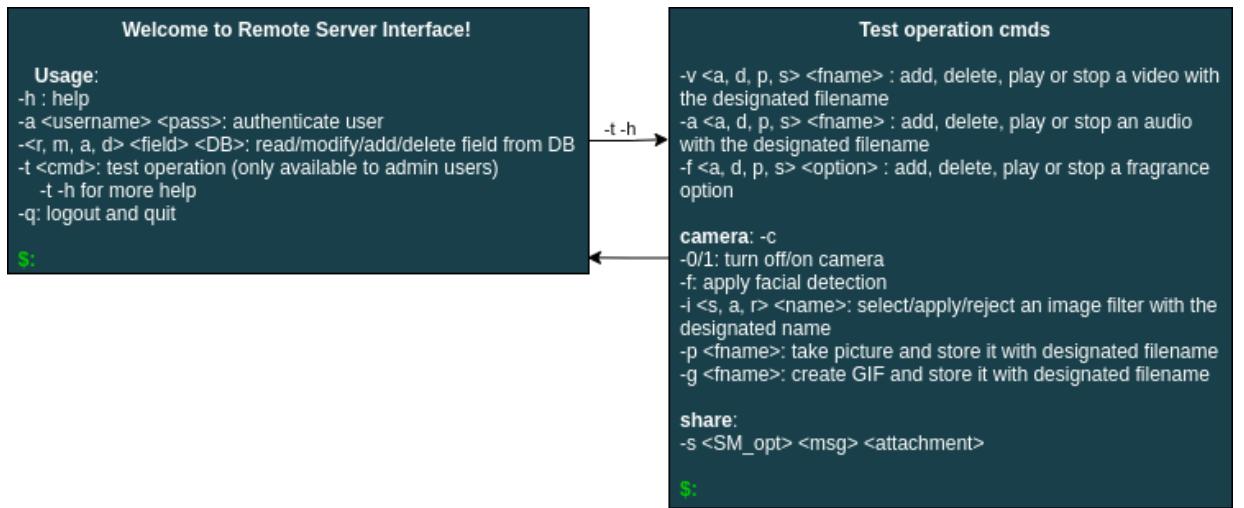


Figure 2.18.: User mock-ups: Remote Server

Table 2.2.: Events: Remote Server

Event	System response	Source	Type
Power on	Initialize RDBMS and go to Idle mode	System maintainer	Asynchronous
Connection Requested	Accept/refuse connection	Remote Client	Asynchronous
Connection Accepted	Start listening for commands	Remote Client	Asynchronous
Authenticate	Query User DB to validate user credentials. If valid, login user.	Remote Client	Asynchronous
Help	Send help information	Remote Client	Asynchronous
Logout	Logout user, close connection and go to Idle mode	Remote Client	Asynchronous
Check WiFi connection	Periodically check WiFi connection	Remote Client	Synchronous
Connection timeout	Logout user, close connection and go to Idle mode	Remote Server	Synchronous
DB management	Read/modify/add/delete data from DB	Remote Client	Asynchronous
Update stations	Update all ready-to-run stations with ads data	Remote Server	Synchronous
Command invalid	Inform RC that command is invalid	Remote Server	Synchronous
Station notification	Store station notification into DB	Local System	Asynchronous
Test Operation RC	Parse command originated from RC and, if valid, dispatch it to designated station	Remote Client (Admin)	Asynchronous
Test Operation Callback	Provide command dispatch to original RC	Local System	Asynchronous

When interacting with the databases, it is possible to read, modify, add or delete some field from a Database. The operation of the Local System can be tested — Test Operation — if the User is an Admin, namely: manage audio, video, fragrance and camera and test the share option.

## 2.4. Subsystem decomposition

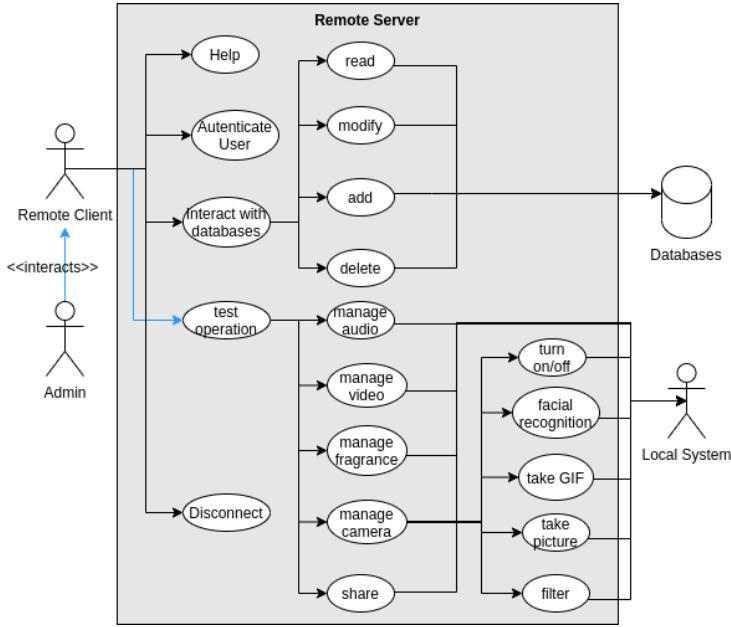


Figure 2.19.: Use cases: remote server

The Manage Camera use case is subdivided into: Turn On/Off, facial detection, take GIF, take picture and filter.

### Dynamic operation

Fig. 2.20 depicts the state machine diagram for the Local System, illustrating its dynamic behavior. There are two main states:

- **Initialization:** the Remote Server is initialized. The settings are loaded and if invalid they are restored. The WiFi communication is setup, signaling the communication status and if valid, an IP address is returned. Lastly, the RDBMS is configured and started: if any error occurs the device goes into the **Critical Error** state, dumping the error to a log file and waiting for reset; otherwise, the initialization is complete.
- **Execution:** after the initialization is successful, the system goes into the **Execution** macro composite state with several concurrent activities, modeled as composite states too. However, it should be noted that there is only one actual state for the device, although at the perceivable time scale they appear to happen simultaneously. These activities are communication management (**Comm Manager**), DB management (**DB manager**), and request handling (**Request Handler**), and are executed forever until the system's power off. They are detailed next.

## 2.4. Subsystem decomposition

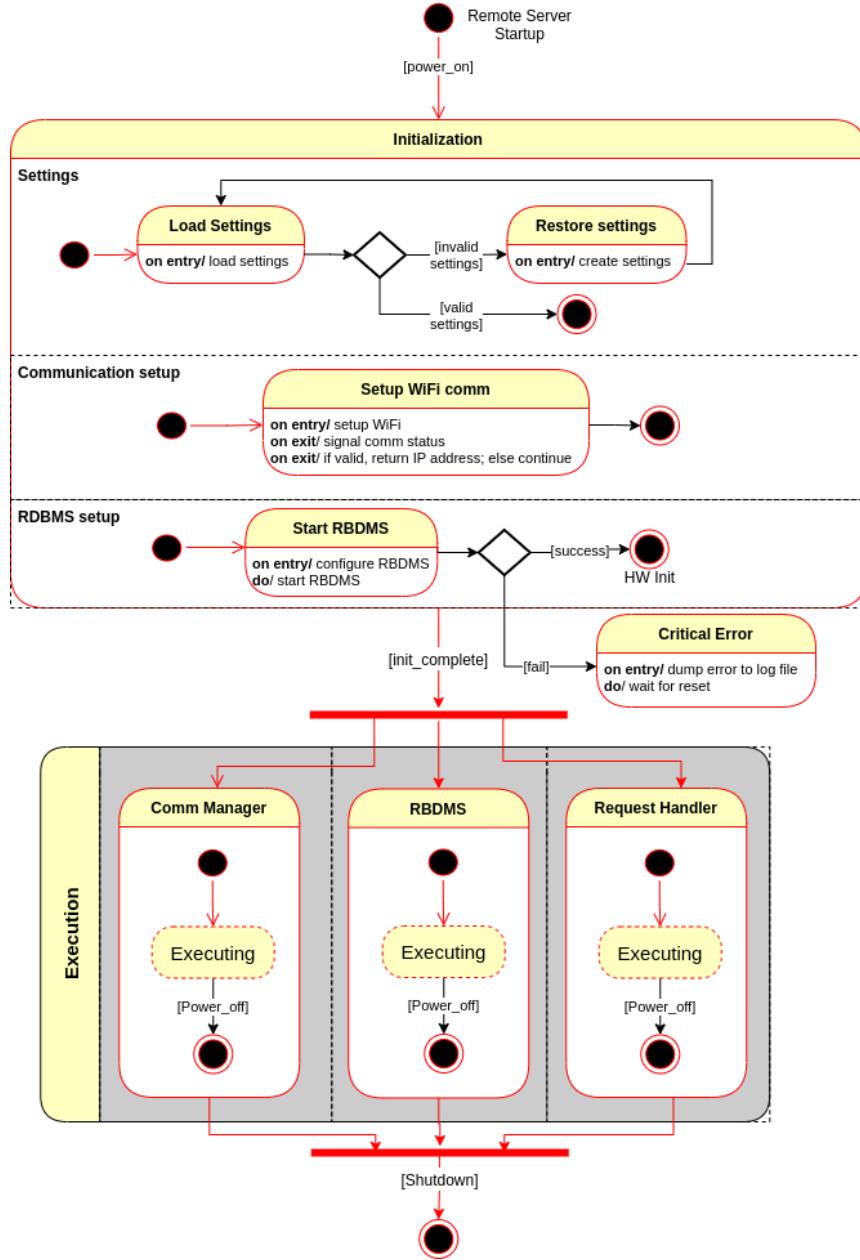


Figure 2.20.: State machine diagram: Remote server

### Communication Manager

Fig. 2.21 depicts the state machine diagram for the **Comm Manager** component. Upon successful initialization the **Comm Manager** goes to **Idle**, listening for incoming connections. When a remote node tries to connects, it makes a connection request which can be accepted or denied. If the connection is accepted and the node authenticates successfully the **Comm Manager** is ready for bidirectional communication. When a message is received from the remote node, it is written to TX msg queue and the **Request Handler**

## 2.4. Subsystem decomposition

is notified. When a message must be sent to the remote, it is read from the TX msg queue and sent to the recipient. If the connection goes down, it is restarted, going into Idle state again.

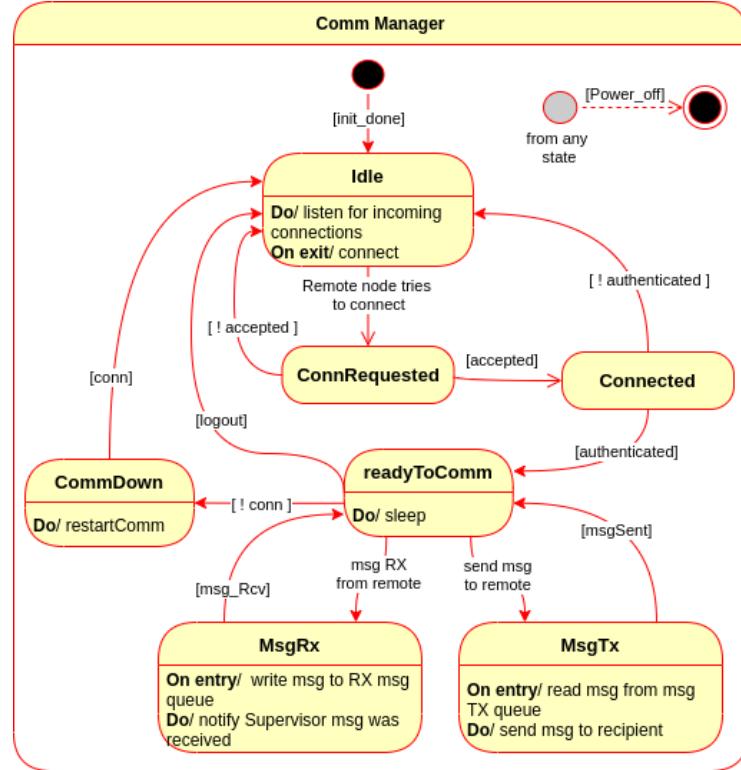


Figure 2.21.: State machine diagram: Remote Server – Comm Manager

### Database Manager

Fig. 2.22 depicts the state machine diagram for the DB Manager component. Upon successful initialization the DB Manager goes to Idle, waiting for incoming DB requests.

When a request arrives, it is parsed, checking its validity. If the request is a DB query, a transaction is read from the respective DB to the RX transaction queue and the Supervisor is notified that there is a transaction to read. Otherwise, if the request is a DB update the transaction is written from the TX transaction queue to the DB and the Supervisor is notified that the DB was updated.

Alternatively, the RDBMS can be triggered to update a station (update\_station event), retrieving the station and operation data to update. A data frame is composed and a server request is created, signaling it to the Request Handler to process it.

## 2.4. Subsystem decomposition

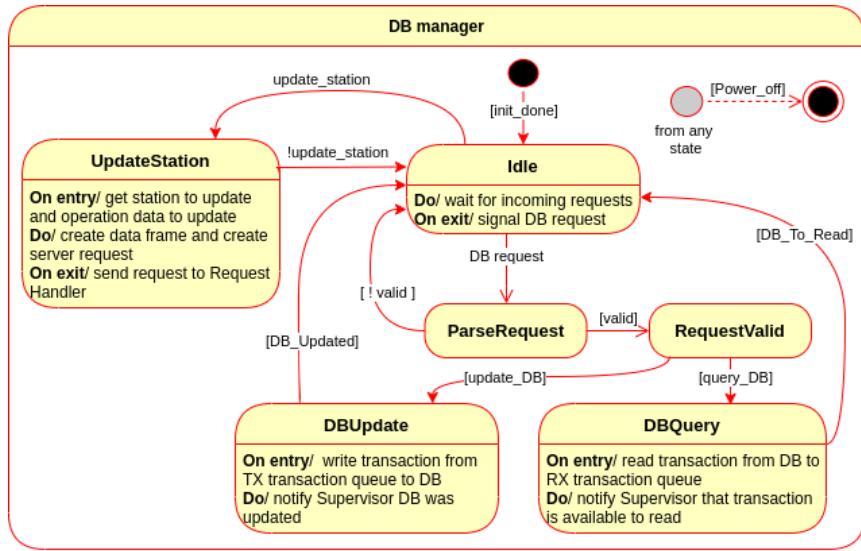


Figure 2.22.: State machine diagram: Remote Server – DB Manager

### Request Handler

Fig. 2.23 depicts the state machine diagram for the Request Handler component, which handles incoming requests from the Remote Client, Local System, or internally (to update stations). When a request arrives, it is parsed, and, if valid, the appropriate callback is triggered, processing the request and returning its output.

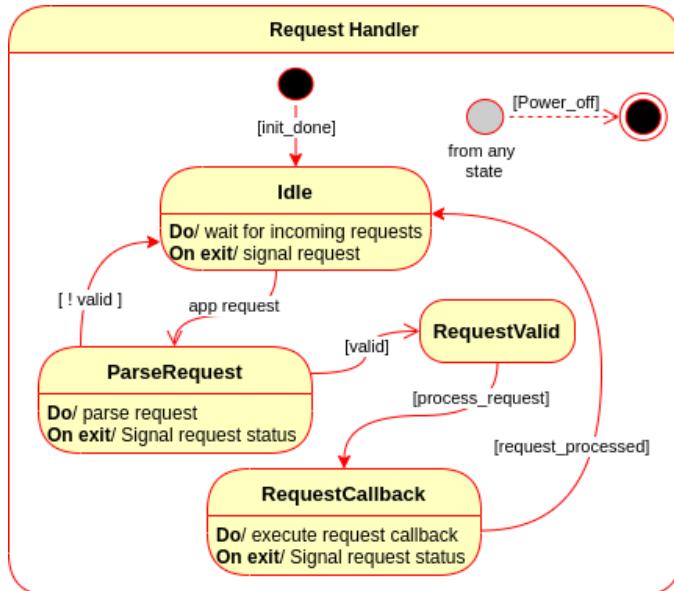


Figure 2.23.: State machine diagram: Remote Server – Request Handler

## Flow of events

The flow of events throughout the system is described using a sequence diagram, comprising the interactions between the most relevant system's entities. It is usually pictured as the visual representation of an use case. The main sequence diagrams are illustrated next (Fig. 2.24 through Fig. 2.30).

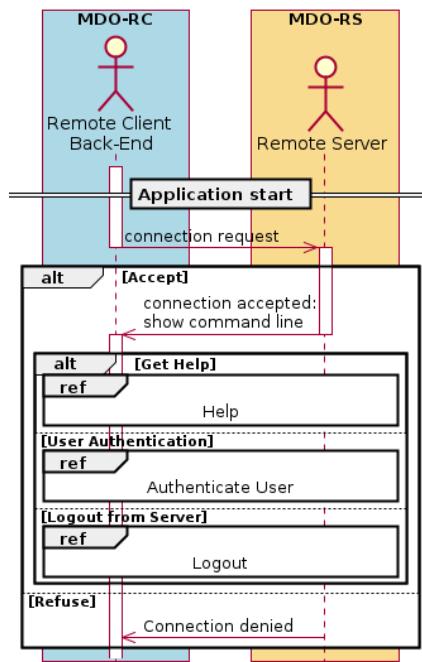


Figure 2.24.: Sequence diagram: Remote Server

The main interaction in all diagrams occurs between the **Remote Client** and the **Remote System**.

In first place, Fig. 2.24 shows the **Application Start**, starting with a connection request that it can be either accepted or denied.

In second place, Fig. 2.25 shows the **Authentication** process, where the **Remote Client** sends the user-name and the password to the **Remote Server**. The latter searches for the username in the **User DB**, if it is not found, the authentication fails, otherwise, the **Remote Server** compares the passwords and sends the authentication status to the **Remote Client**. It is only possible to access the **Manage DBs** and the **Test Operation** if the authentication succeeds.

In Fig. 2.26 the **Remote Client** requests to the **Remote Server** to manage a Database's field. The **Remote Server** makes the request to the specific **Database**, receives the response and returns it to the **Remote Client**.

In Fig. 2.27, Fig. 2.28 and Fig. 2.29 it can be seen the **Test Operation**. As said previously, only the **Admin** can access this part, where it can test the functionality of any **Local System**. The **Admin** sends

## 2.4. Subsystem decomposition

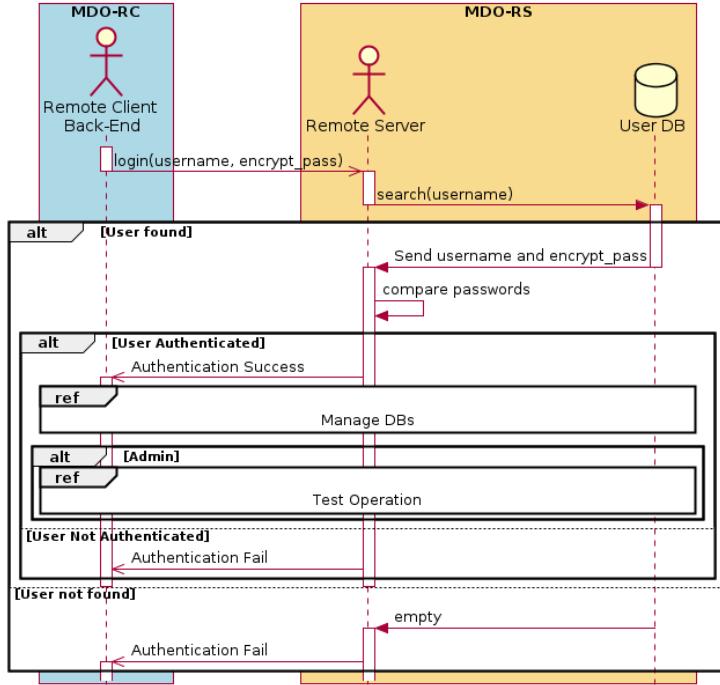


Figure 2.25.: Sequence diagram: Remote Server – Authentication

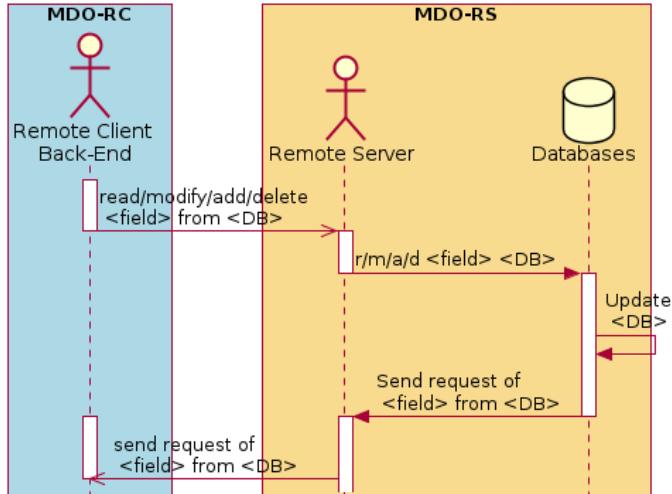


Figure 2.26.: Sequence diagram: Remote Server – Manage Databases

the request to the Remote Client that then sends it to the Remote Server with the specific station. Then, the Remote Server interacts with the specific station, making the specific Local System act according to the command that was sent. For every test operation there's always a command feedback to indicate to the User its execution status. The references Apply Filter, Take Picture and Create GIF in Fig. 2.28 and Sharing Mode in Fig. 2.29 are the responsibility of the Local System part and will be explained there.

## 2.4. Subsystem decomposition

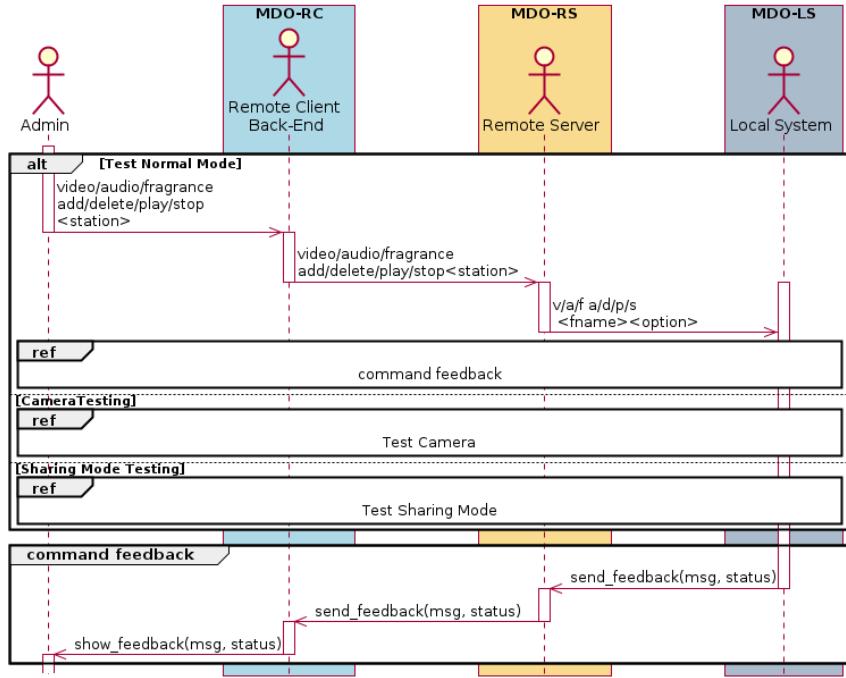


Figure 2.27.: Sequence diagram: Remote Server – Test Operation

### 2.4.3. Local system

In this section the local system is analyzed, considering its events, use cases, dynamic operation and the flow of events.

#### User mock-ups

Fig. 2.31 illustrates the user mock-ups for the local system. It intends to mimic the user interaction with the local system, clarifying the user actions (gestures) and the respective responses, as well as the workflow, comprising its four modes.

The initial state of the MDO-L's UI is depicted in thick border outline, after a User has been detected – Interaction mode. On the left it is the camera feed and on the right the commands ribbon, containing the hints to use the system and the available options. As it can be seen, the User can choose an option by hovering with pointing finger over the desired option for a designated amount of time (e.g., 3 seconds).

The workflow can be as follows:

- If the User selects the Image filter option, the Image filtering view is shown, presenting the options to select filters (which can be scrolled through palm raising/lowering), to cancel or accept the image filter. If a filter is selected `filter1_pressed`, it is applied, and if accepted it will return to Interaction mode.

## 2.4. Subsystem decomposition

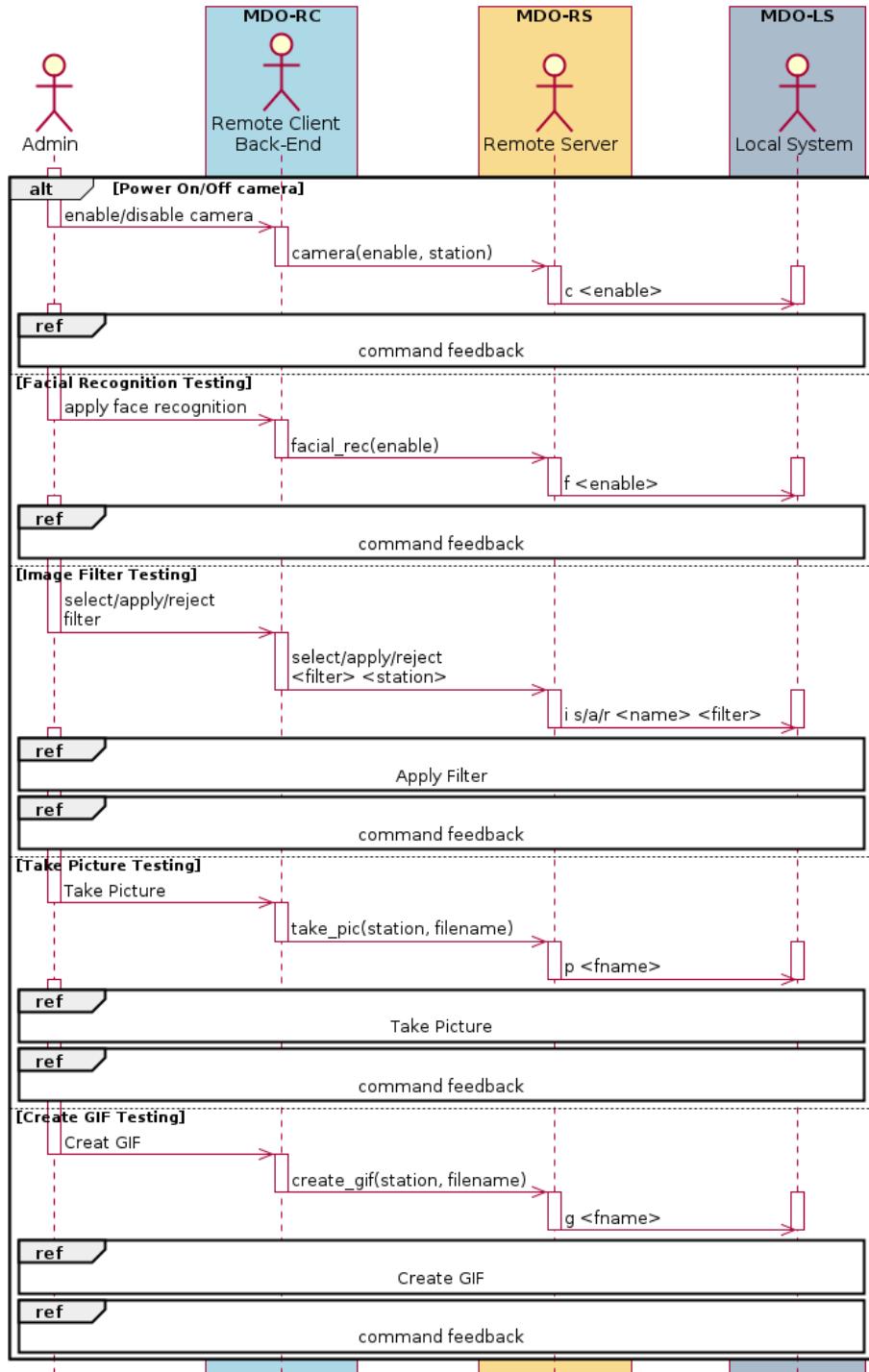


Figure 2.28.: Sequence diagram: Remote Server — Test Operation Camera

mode, keeping the filter on.

- If the User selects the Take Pic option, Picture mode is started with a timer to allow the User to get

## 2.4. Subsystem decomposition

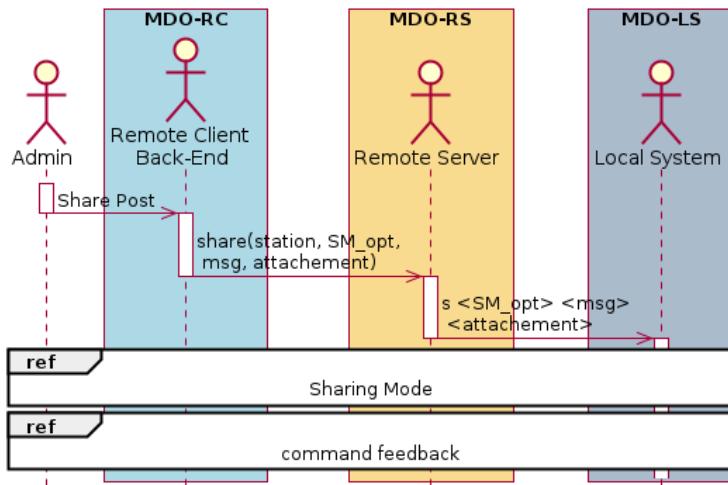


Figure 2.29.: Sequence diagram: Remote Server – Test Operation Share

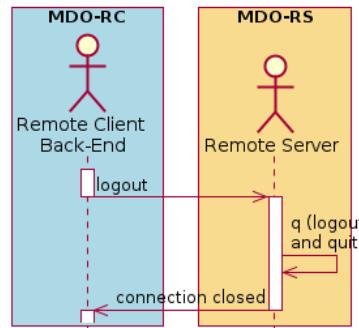


Figure 2.30.: Sequence diagram: Remote Server – logout

ready before actually taking the picture. The User can Cancel – returning to main menu – or Share – starting Sharing mode.

- If the User selects the Create GIF option, GIF mode (setup) is started with a timer to allow the User to get ready before actually creating the GIF. After the **setup\_timer** is elapsed, the GIF mode (operation) starts, displaying a dial with the GIF duration until being complete. When the **gif\_timer** elapses, the GIF is created, enabling the User to Cancel – returning to main menu – or to Share – starting Sharing mode.
- Lastly, in the Sharing mode, the User can Cancel – returning to main menu – or select the social media network. After selecting the social media, the User can edit the post by entering its customized message and, if Share is pressed, a message box will appear displaying the status of the post sharing – Success or Error.

## 2.4. Subsystem decomposition

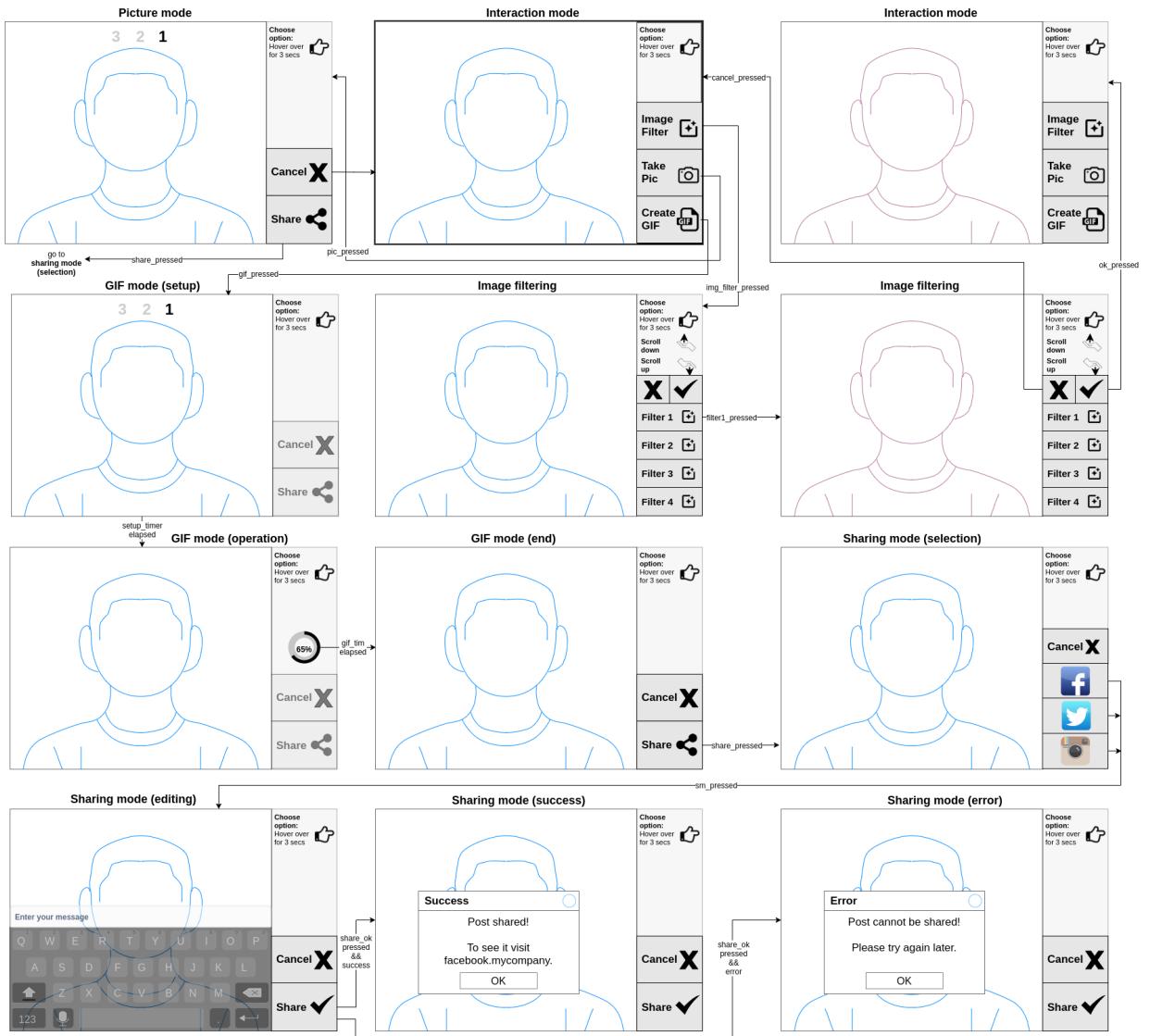


Figure 2.31.: User mock-ups: local system

## Events

Table 2.3 presents the most relevant events for the Local system, categorizing them by their source and synchrony and linking it to the system's intended response. A further division is done separating UI events from the remaining ones.

## 2.4. Subsystem decomposition

---

Table 2.3.: Events: local system

Event	System response	Source	Type
Power on	Initialize sensors and go to Normal mode	System maintainer	Asynchronous
User detected	Turn on camera feed and go to Interaction mode	User	Asynchronous
Command received	Parse it and respond	Remote Server	Asynchronous
Database update	Request update of internal databases to Remote Server	Database manager	Asynchronous
Enable fragrance diffuser	Enable fragrance diffusion for a predefined period of time	Local System	Synchronous
Video ended	Playback the next video on the queue	Local System	Synchronous
Check WiFi connection	Periodically check WiFi connection	Local System	Synchronous
<b>UI events</b>			
Option selected	Track the option selected and inform the UI engine	User	Asynchronous
Image filter pressed	Go to Image filter view	User	Asynchronous
Filter selected	Detect User's face and apply filter	User	Asynchronous
Pic pressed	Go to Picture mode	User	Asynchronous
Pic setup elapsed	Take picture	Local System	Synchronous
GIF pressed	Go to GIF mode	User	Asynchronous
GIF setup elapsed	Go to GIF operation	Local System	Synchronous
GIF operation elapsed	Finish GIF	Local System	Synchronous
Share mode pressed	Go to Sharing mode (selection)	User	Asynchronous
Keyboard pressed	Give feedback to user	User	Asynchronous
Share post pressed	Upload post to designated social media	User	Asynchronous
Share post status	Inform user about shared post status	Cloud	Asynchronous

## Use cases

Fig. 2.32 depicts the use cases diagram for the Local System, describing how the system should respond under various conditions to a request from one of the stakeholders to deliver a specific goal.

The Admin interacts with the Remote Client (through its UI) requesting the Remote Server to process commands, getting the state of the device, adding a video or selecting the fragrance. Additionally, the Admin may test the operation of the device: play video, test audio, nebulize fragrance or test the camera. This last one tests the main functionalities the User also utilizes, namely: select image filter, apply/reject image filter, take picture, create GIF or share multimedia on the social media.

## 2.4. Subsystem decomposition

A precondition for the interaction of the Admin with the Local System is the establishment of a remote connection between the Remote Server and the Local system, verifying its credentials. However, there is another important use case for this remote connection: the update of the Local System's internal databases from the Remote server which will sent appropriate commands for this purpose.

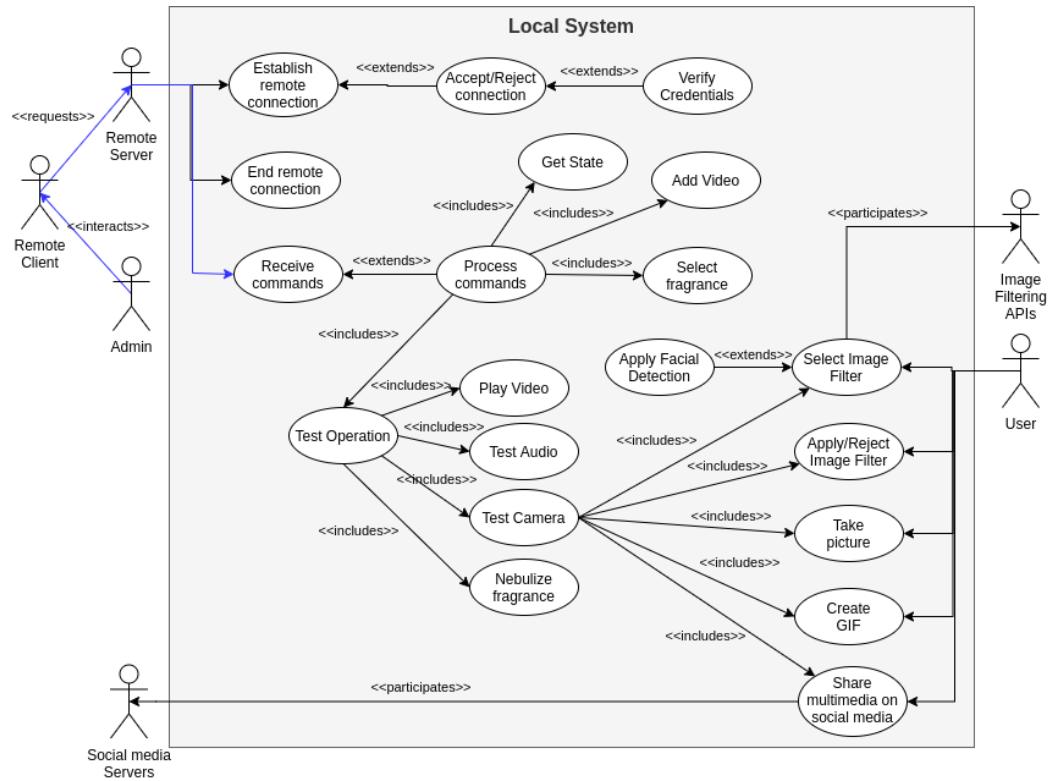


Figure 2.32.: Use cases diagram: local system

## Dynamic operation

Fig. 2.33 depicts the state machine diagram for the Local System, illustrating its dynamic behavior.

There are two main states:

- **Initialization:** the device is initialized. The settings and DBs are loaded and if invalid they are restored. The WiFi communication is setup, signaling the communication status and if valid, an IP address is returned. Lastly, the HW is initialized, checking its presence, configuring it and testing the configuration: if any error occurs the device goes into the **Critical Error** state, dumping the error to a log file and waiting for reset; otherwise, the initialization is complete.
- **Execution:** after the initialization is successful, the system goes into the **Execution** macro composite state with several concurrent activities, modeled as composite states too. However, it should be noted

## 2.4. Subsystem decomposition

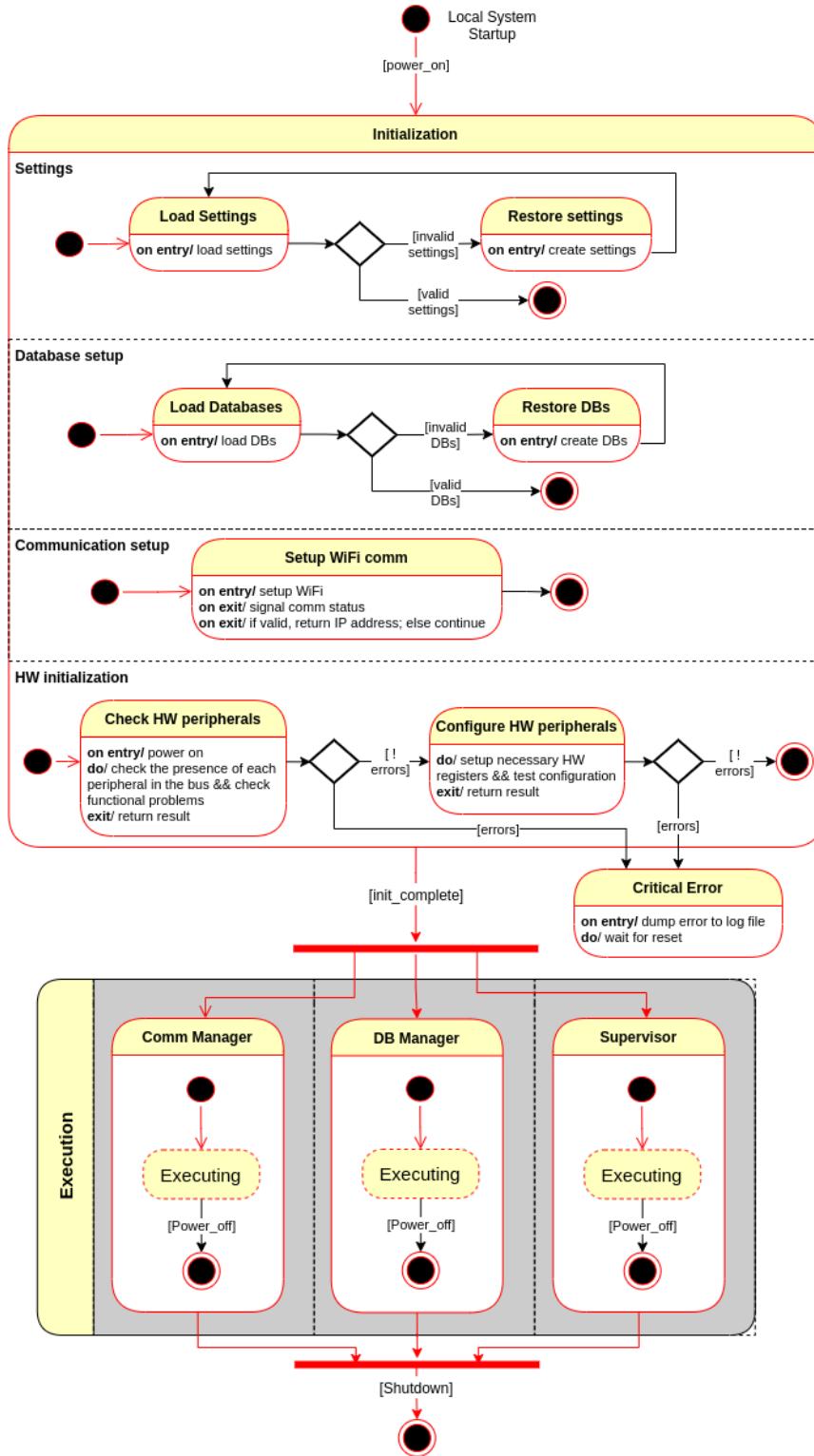


Figure 2.33.: State machine diagram: local system

## 2.4. Subsystem decomposition

that there is only one actual state for the device, although at the perceivable time scale they appear to happen simultaneously. These activities are communication management (Comm Manager), DB management (DB manager), and application supervision (Supervisor), and are executed forever until system's power off. They are detailed next.

### Communication Manager

Fig. 2.34 depicts the state machine diagram for the Comm Manager component. Upon successful initialization the Comm Manager goes to Idle, listening for incoming connections. When a remote node tries to connects, it makes a connection request which can be accepted or denied. If the connection is accepted and the node authenticates successfully the Comm Manager is ready for bidirectional communication. When a message is received from the remote node, it is written to TX msg queue and the Supervisor is notified. When a message must be sent to the remote, it is read from the TX msg queue and sent to the recipient. If the connection goes down, it is restarted, going into Idle state again.

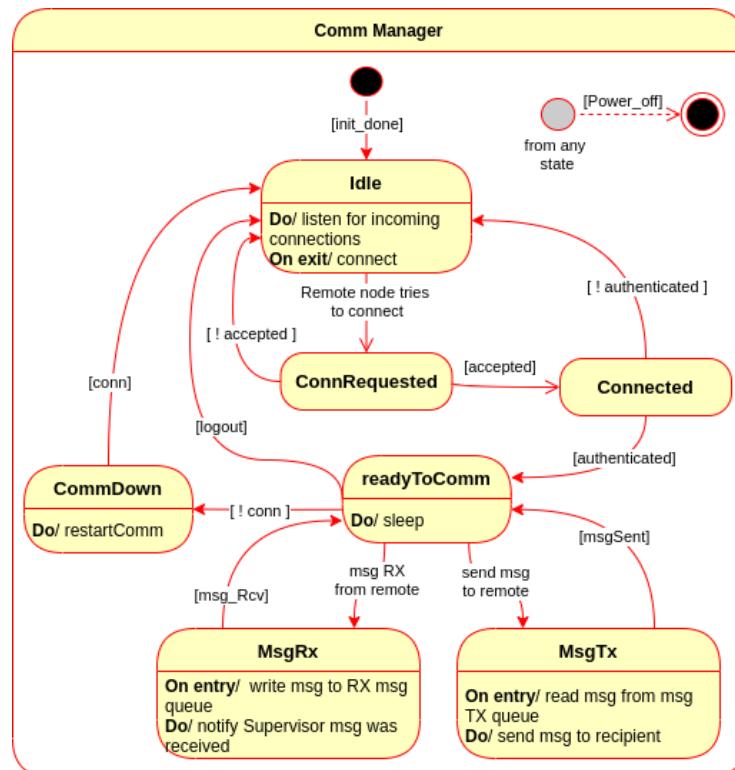


Figure 2.34.: State machine diagram: local system – Comm Manager

## 2.4. Subsystem decomposition

### Database Manager

Fig. 2.35 depicts the state machine diagram for the DB Manager component. Upon successful initialization the DB Manager goes to **Idle**, waiting for incoming DB requests. When a request arrives, it is parsed, checking its validity. If the request is a DB query, a transaction is read from the respective DB to the RX transaction queue and the Supervisor is notified that there is a transaction to read. Otherwise, if the request is a DB update the transaction is written from the TX transaction queue to the DB and the Supervisor is notified that the DB was updated.

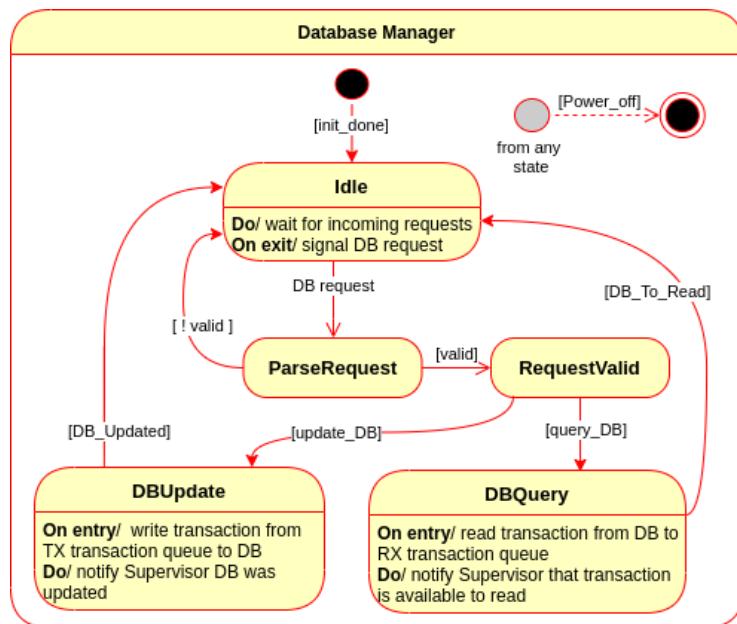


Figure 2.35.: State machine diagram: local system – DB Manager

### Supervisor

Fig. 2.36 depicts the state machine diagram for the Supervisor component, comprising two tasks running in ‘parallel’ (Fig. 2.36a):

- Request Handler (Fig. 2.36b): handles incoming requests from the Remote server. When a request arrives, it is parsed, and, if valid, the appropriate callback is triggered, processing the request and returning its output.
- Mode manager (Fig. 2.36c): Upon successful initialization the Mode Manager goes to **Idle**, and it is ‘awake’ if it is time to play the advertisements or if a user is detected. If the former is verified—Normal mode — the device retrieves video and fragrance data from the DB and plays video and nebulizes fragrance. If the latter is verified — Interaction mode the device turns on the camera and mirrors the feed on the display, waiting for a recognizable gesture.

## 2.4. Subsystem decomposition

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If the **User** choose to select an image filter, take a picture or create a GIF, the device goes into **Multimedia mode**, returning back to **Interaction mode** after its exit condition or after a timeout.

Lastly, if the **User** chooses to share the image or GIF created, it must select the social media network, edit the post to enter some message and confirm the sharing, returning to **Interaction mode**. If no user interaction happens for a while, the device returns back to **Idle mode**.

### Flow of events

The flow of events throughout the system is described using a sequence diagram, comprising the interactions between the most relevant system's entities. It is usually pictured as the visual representation of an use case. The main sequence diagrams are illustrated next.

#### Normal mode

Fig. 2.37 depicts the **Normal mode**'s sequence diagram. The blue area delimits the **MDO-L** system, comprising the **Local System Back-End**.

When its time to play the advertisements, the **Normal mode** is activated, retrieving video, audio and fragrance from the internal DB. Then two parallel activities are executed:

- Video playback: while its time to play the advertisements, a video is played from the video list. When it finishes, it moves the next video in the playback queue.
- Fragrance diffusion: while there are timestamps for fragrance diffusion, diffuse fragrance between start and stop times and sleep on the other occasions.

#### Interaction mode

Fig. 2.38 depicts the **Interaction mode**'s sequence diagram. The blue area delimits the **MDO-L** system, comprising the **Gesture Recognition Engine**, the **UI engine** and the **Local System Back-End**.

When the **User** is in range (asynchronous event), the camera is activated, and two parallel activities are executed:

- mirror camera feed: while the **User** is in range and active, the **UI engine** will grab frame from the camera and display it on the window providing visual feedback to the **User**.
- gesture recognition and processing: if a gesture is recognized by the **Gesture Recognition Engine** it is dispatched to the **Local System Back-End** which will process it according to the following cases: **Select Image filter**, **Take Pic**, and **Create GIF**, showing the respective view in the **UI** and triggering the associate sequence diagram (indicated by the **ref** keyword).

## 2.4. Subsystem decomposition

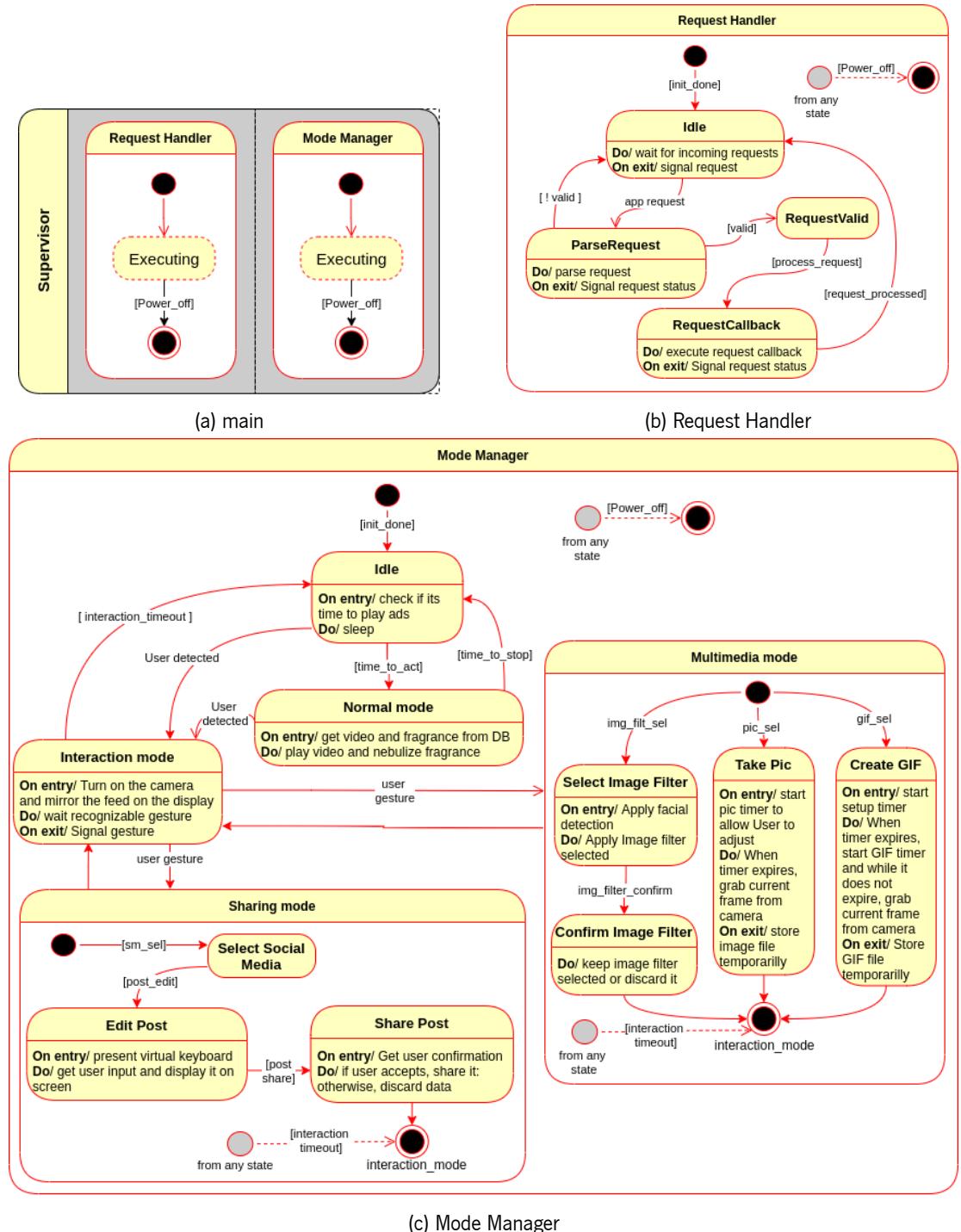


Figure 2.36.: State machine diagram: local system – Supervisor

### Multimedia mode

Fig. 2.39 through Fig. 2.41 depicts the Multimedia mode's sequence diagrams, namely:

## 2.4. Subsystem decomposition

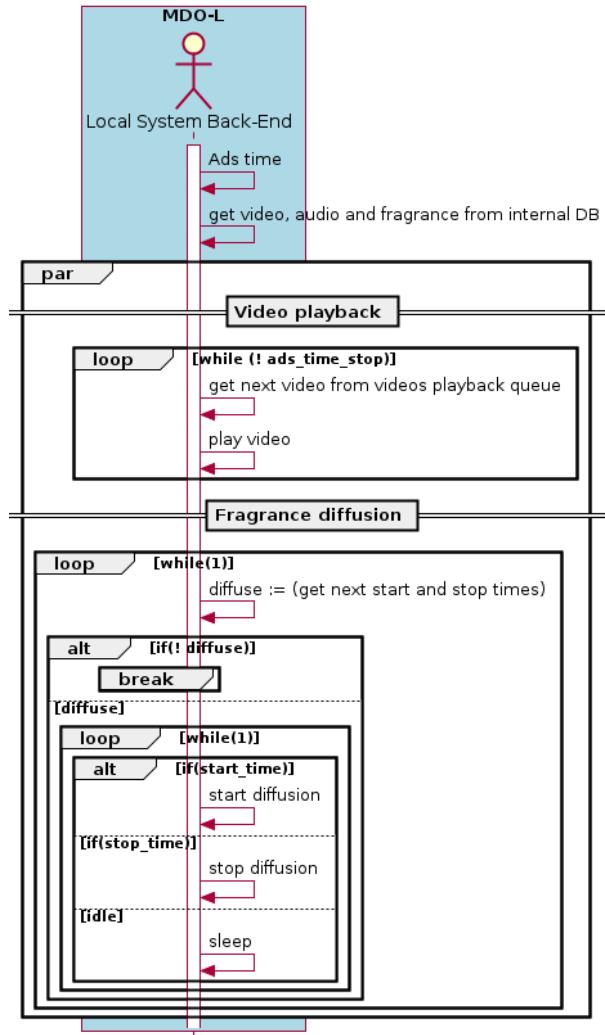


Figure 2.37.: Sequence diagram: local system – Normal mode

- **Select image filter** (Fig. 2.39): after the **Image filter view** is presented to the **User**, he/she can make a gesture to select the filter, which upon being recognized by the **Gesture Recognition Engine** it is dispatched by the **UI Engine** to the **Local System back-end**. Facial detection is then applied, and while the filter is active, a request is made to **Image Filtering APIs** to apply the designated filter, showing it to the **User** – **Apply filter** reference.  
If the **User** accepts the filter, it returns to **Interaction mode** with the filter simultaneously on (**Apply filter**). Otherwise, if the **User** cancels it, it simply returns to **Interaction mode**.
- **Take picture**: after **Picture mode** is initiated, the **Local System back-end** starts a timer to allow the **User** to get in position, and while the timer is running, the time remaining is presented to the **User**. When the timer elapses the picture is stored internally.

## 2.4. Subsystem decomposition

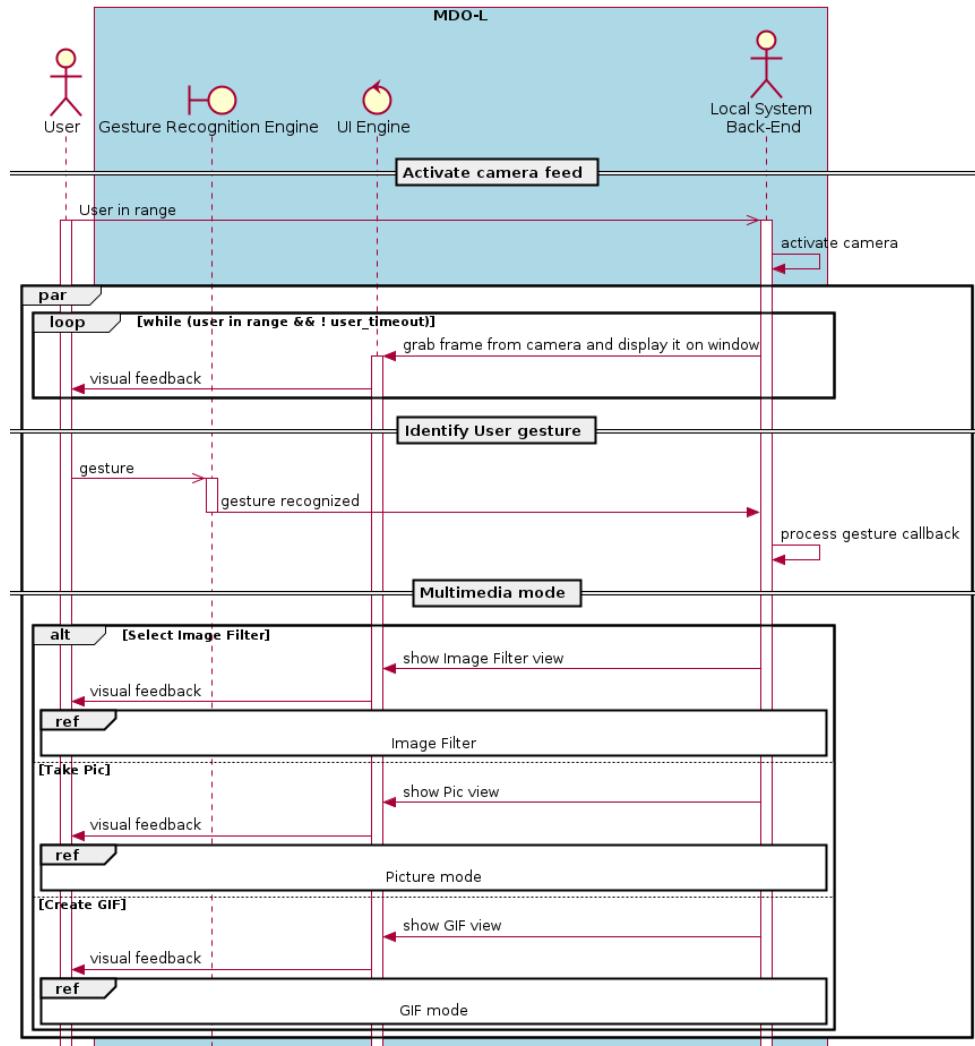


Figure 2.38.: Sequence diagram: local system – Interaction mode

- **Create GIF:** after GIF mode is initiated, the Local System back-end starts a timer to allow the User to get in position, and while the timer (`gif_setup_timer`) is running, the time remaining is presented to the User. When the timer elapses the GIF creation can start, with another timer (`gif_oper_timer`) being started, and while the timer is running the remaining time is shown to User but in a dial form. When this timer elapses the GIF is stored internally.

### Sharing mode

Fig. 2.42 depicts the Sharing mode's sequence diagram.

The User starts by selecting the social media platform (with a gesture), which upon being recognized is dispatched to the Local System back-end identifying the social media selected. Then, the social me-

## 2.4. Subsystem decomposition

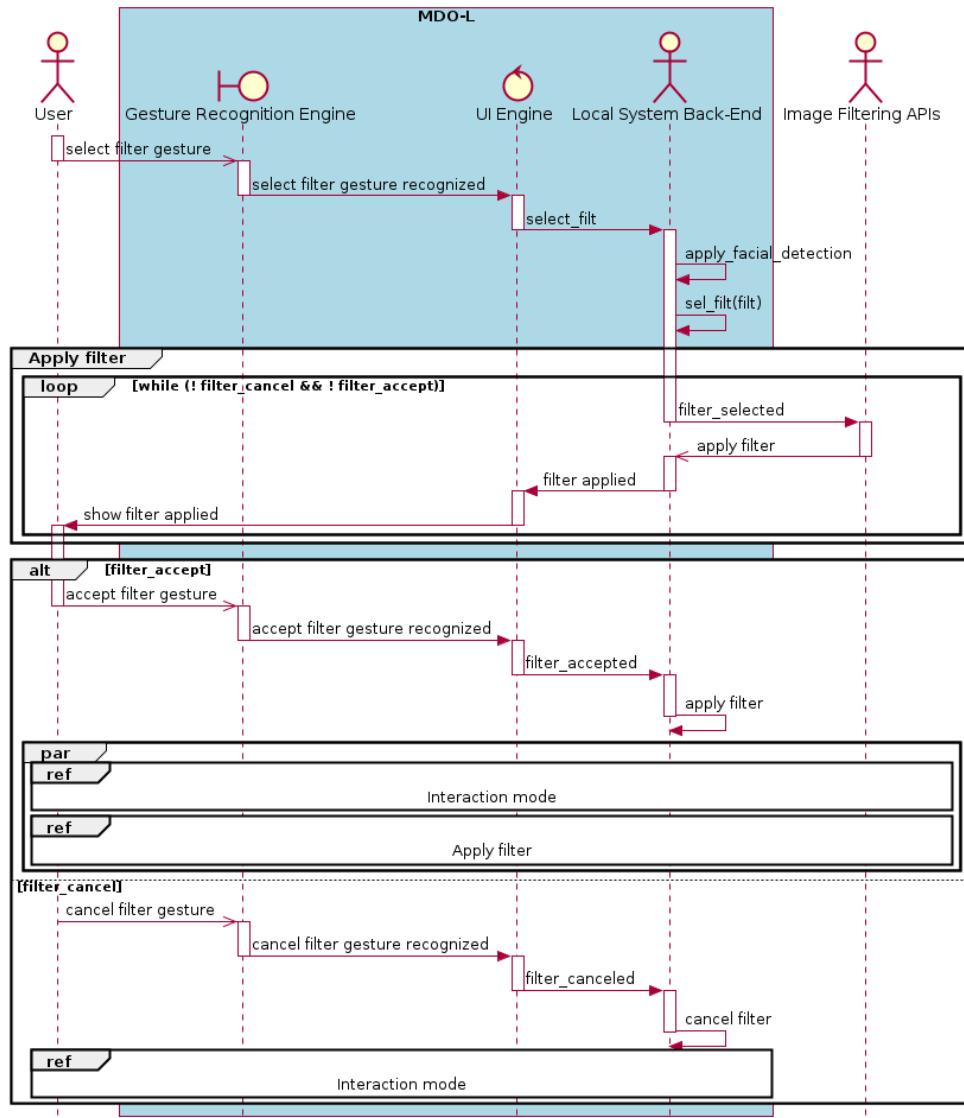


Figure 2.39.: Sequence diagram: local system – Multimedia mode (select image filter)

dia parameters are configured and the attachment is set to the last multimedia file. After social media configuration, the Post Edit view is shown to the User.

In the Post editing mode, while the User does not decide to share or cancel the post, the selected character from the virtual keyboard is visually feed back to the User.

When the User decides to share or cancel the post, it will trigger share\_post or cancel\_post callbacks, with the latter going into Interaction mode.

Upon triggering the share\_post callback, Local System back-end tries to perform the login in the required social media platform, requesting it to one of its servers. If the login succeeds, the post is sent to Social

## 2.5. Budget estimation

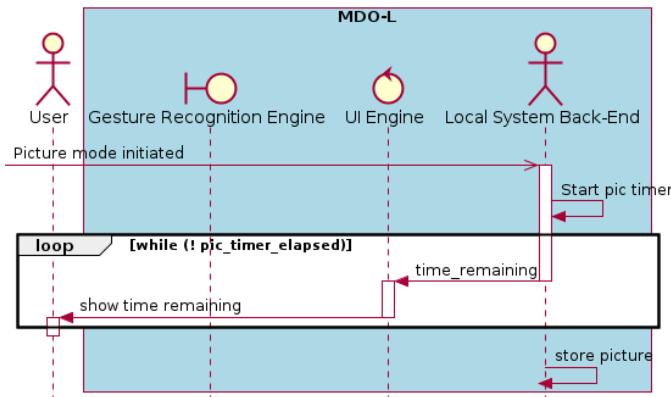


Figure 2.40.: Sequence diagram: local system – Multimedia mode (take picture)

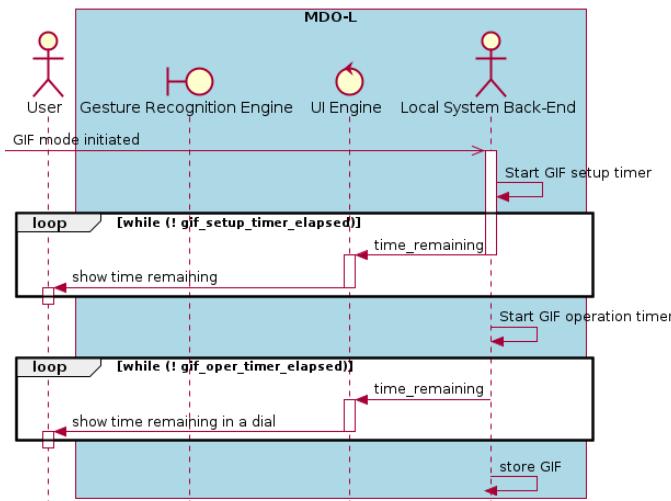


Figure 2.41.: Sequence diagram: local system – Multimedia mode (create GIF)

Media servers, which will return its status. Upon completion a dialog box will be presented to the User, informing it of share post status: success or failure (in case the login or the post transmission fails).

## 2.5. Budget estimation

Table 2.4 presents the budget estimation for the project.

Two classes of prototypes were considered, scale-model and real-scale, due to project feasibility concerns: the real-scale prototype envisions a large display (55 inch) whose cost is prohibitive; thus, a trade-off between functionality and cost needed to be performed.

Additionally, several types of costs were assessed, namely:

## 2.5. Budget estimation

Table 2.4.: Budget estimation

		Scale-model Prototype	Real-scale Prototype		
		Item	Cost (€) *	Item	Cost (€) *
<b>HW</b>	Raspberry Pi 4B	50	Raspberry Pi 4B	50	50
	User Detection sensor	3	User Detection sensor mesh (4)	12	12
	LCD display 10" (non-touch)	55	LCD display 55" Full HD (non-touch)	1500	1500
	Fragrance diffusion actuator	5	Fragrance diffusion actuator mesh (4)	20	20
	Camera 8 MP	32	Camera 8 MP	32	32
	Speakers	5	Speakers	30	30
	Power supply	10	Power supply	30	30
	PCB	8	PCB	16	16
<b>Mechanical Structure</b>	3D printed + screws	20	Built-in with display + HW packaging	100	100
			Full HW packaging	350	350
	<b>Physical Prototype cost</b>	<b>188</b>	<b>Physical Prototype cost **</b>	<b>2040</b>	
<b>SW development</b>	Remote Client: 500 h ***	5000	Remote Client: 500 h ***	5000	5000
	Remote Server: 300 h ***	3000	Remote Server: 300 h ***	3000	3000
	Local system: 1000 h ***	10000	Local system: 1000 h ***	10000	10000
	<b>SW development cost</b>	<b>18000</b>	<b>SW development cost</b>	<b>18000</b>	
<b>Operational costs</b>	Local System power consumption ****	26,28	Local System power consumption ****	197,1	197,1
	Server operation cost *****	420	Server operation cost *****	420	420
	<b>Yearly Operational cost</b>	<b>446.28</b>	<b>Yearly Operational cost</b>	<b>617.1</b>	
	<b>Total cost</b>	<b>18634.28</b>	<b>Total cost</b>	<b>20657.1</b>	

\* tax included

\*\* considering the most expensive option

\*\*\* 10 €/h

\*\*\*\* 24h/7d for 1 year

\*\*\*\*\* yearly cost

- physical prototype cost – comprises HW and mechanical structure of the device: all estimation costs were made as a mean value between all trustworthy suppliers. The physical prototype cost for the scale-model prototype is about 188 EUR, while for the real-scale prototype is about 2,040 EUR, with the main difference being due to the 55 inch display cost (1,500 EUR) and the full HW packaging (350 EUR).
- SW development cost: the development cost for all software components, namely **Remote Client**, **Remote Server**, and **Local System**, yielding 18 000 hours of development, which represents about 3 months of work for a two people team. This cost is the same for both prototypes as the scale factor is only associated to HW.
- operational costs – comprises power consumption and server operation costs on a yearly basis. The server operation is the same for both prototypes, but the power consumption is more than 7 times more.
- total cost: sum of physical prototype, SW development and operational costs. The total cost for the

## **2.5. Budget estimation**

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scale-model prototype is about 18,635 EUR, while for the real-scale one it is about 20,657 EUR, with the main difference being due to HW and power consumption costs.

Retail price and break-even analysis were not assessed at this point, as several business models may be used for that purpose.

## 2.5. Budget estimation

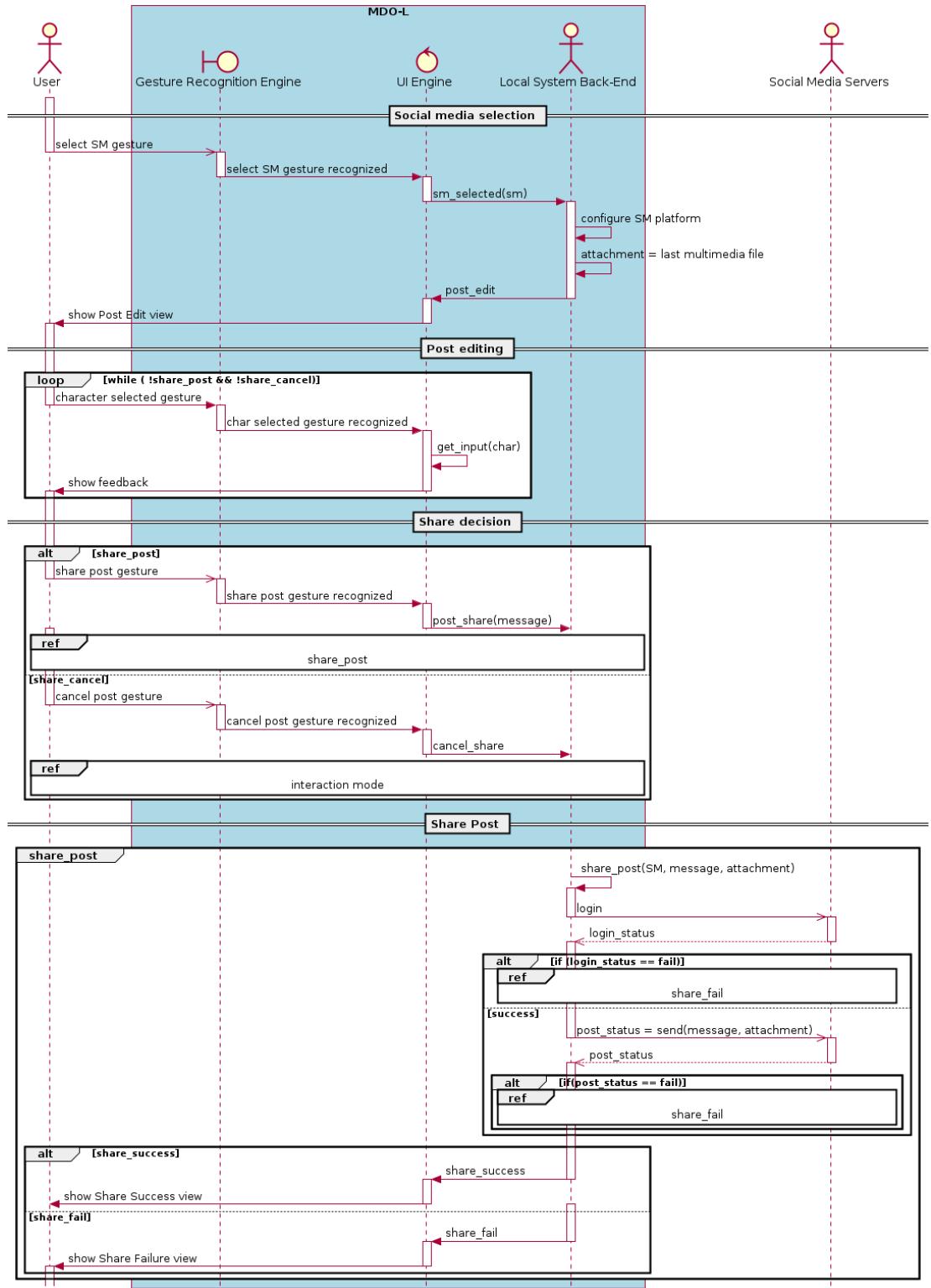


Figure 2.42.: Sequence diagram: local system – Sharing mode

# **3. Theoretical foundations**

In this chapter the theoretical foundations are outlined, providing the basic technical knowledge to undertake the project.

## **3.1. Project methodology**

In this section the project methodologies tools are outlined, easing the development process.

### **3.1.1. Waterfall model**

For the domain-specific design of software the waterfall methodology is used. The waterfall model (fig. 3.1) represents the first effort to conveniently tackle the increasing complexity in the software development process, being credited to Royce, in 1970, the first formal description of the model, even though he did not coin the term [9]. It envisions the optimal method as a linear sequence of phases, starting from requirement elicitation to system testing and product shipment [10] with the process flowing from the top to the bottom, like a cascading waterfall.

In general, the phase sequence is as follows: analysis, design, implementation, verification and maintenance.

1. Firstly, the project requirements are elicited, identifying the key requirements and constraints the system being developed must meet from the end-user perspective, captured in natural language in a product requirements document.
2. In the analysis phase, the developer should convert the application level knowledge, enlisted as requirements, to the solution domain knowledge resulting in analysis models, schema and business rules.
3. In the design phase, a thorough specification is written allowing the transition to the implementation phase, yielding the decomposition in subsystems and the software architecture of the system.

### 3.1. Project methodology

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4. In the implementation stage, the system is developed, following the specification, resulting in the source code.
5. Next, after system assembly and integration, a verification phase occurs and system tests are performed, with the systematic discovery and debugging of defects.
6. Lastly, the system becomes a product and, after deployment, the maintenance phase start, during the product life time.

While this cycle occurs, several transitions between multiple phases might happen, since an incomplete specification or new knowledge about the system, might result in the need to rethink the document.

The advantages of the waterfall model are: it is simple and easy to understand and use and the phases do not overlap; they are completed sequentially. However, it presents some drawbacks namely: difficulty to tackle change and high complexity and the high amounts of risk and uncertainty. However, in the present work, due to its simplicity, the waterfall model proves its usefulness and will be used along the project.

As a reference in the sequence of phases and the expected outcomes from each one, it will be used the chain of development activities and their products depicted in fig. 3.2 (withdrawn from [11]).

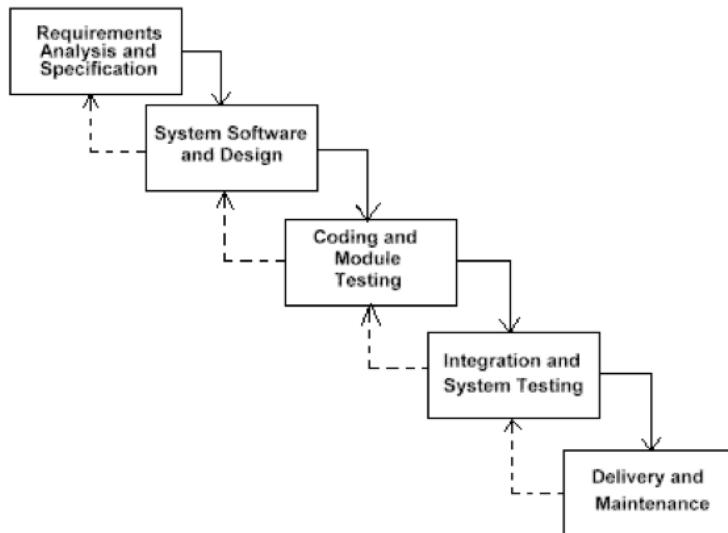


Figure 3.1.: Waterfall model diagram

### 3.1.2. Unified Modeling Language (UML)

To aid the software development process, a notation is required, to articulate complex ideas succinctly and precisely. The notation chosen was the Unified Modeling Language (UML), as it provides a spectrum of

notations for representing different aspects of a system and has been accepted as a standard notation in the software industry [11].

The goal of UML is to provide a standard notation that can be used by all object-oriented methods and to select and integrate the best elements of precursor software notations, namely Object-Modeling Technique (OMT), Booch, and Object Oriented Software Engineering (OOSE) [11]. It provides constructs for a broad range of systems and activities (e.g., distributed systems, analysis, system design, deployment). System development focuses on three different models of the system (fig. 3.2) [11]:

1. **The functional model:** represented in UML with use case diagrams, describes the functionality of the system from the user's point of view.
2. **The object model:** represented in UML with class diagrams, describes the structure of the system in terms of objects, attributes, associations, and operations.
3. **The dynamic model:** represented in UML with interaction diagrams, state-machine diagrams, and activity diagrams, describes the internal behaviour of the system.

## 3.2. Concurrency

Concurrency is used to refer to things that appear to happen at the same time, but which may occur serially [12], like the case of a multithreaded execution in a single processor system. Two concurrent tasks may start, execute and finish in overlapping instants of time, without the two being executed at the same time. As defined by the Portable Operating System Interface (POSIX) specification, a concurrent execution requires that a function that suspends the calling thread shall not suspend other threads, indefinitely.

This concept is different from parallelism. Parallelism refers to the simultaneous execution of tasks, like the one of a multithreaded program in a multiprocessor system. Two parallel tasks are executed at the same time and, as such, they require the execution in exclusivity in independent processors.

Every concurrent system provides three important facilities [12]:

- **Execution Context:** refers to the concurrent entity state. It allows the context switch and it must maintain the entities states, independently.
- **Scheduling:** in a concurrent system, the scheduling decides what context should execute at any given time.
- **Synchronization:** this allows the management of shared resources between the concurrent execution contexts.

### 3.3. Threads versus Processes

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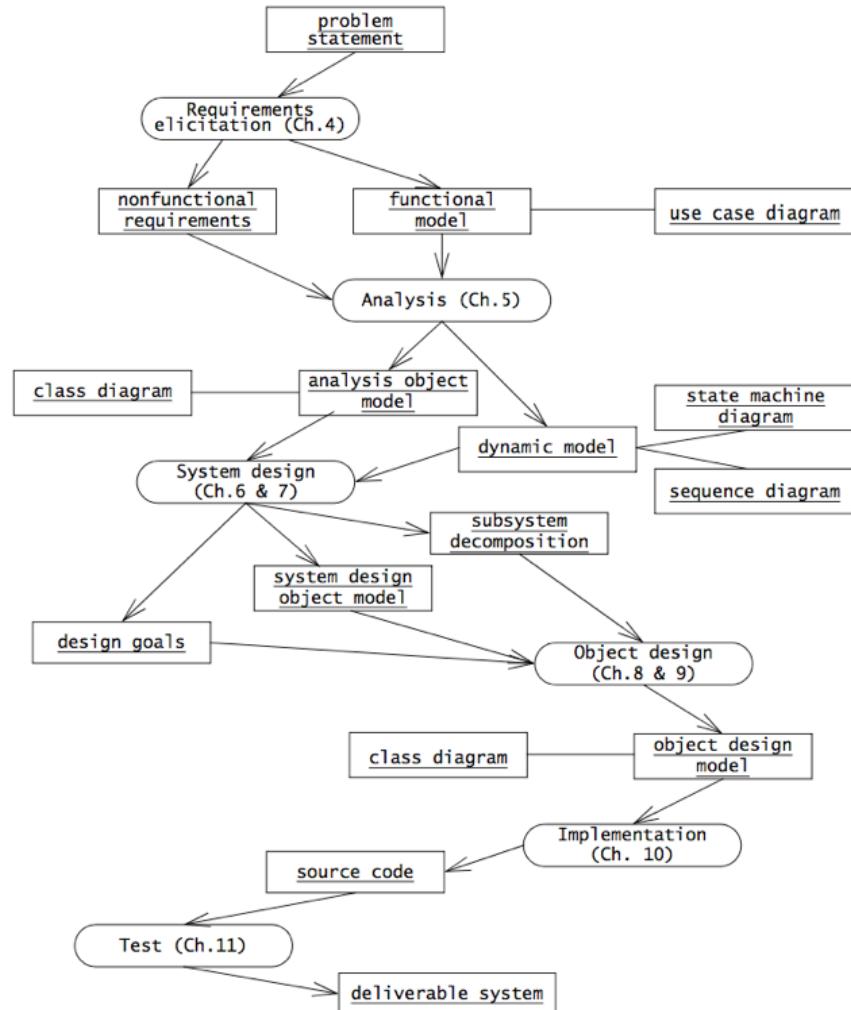


Figure 3.2.: An overview of the object-oriented software engineering development and their products. This diagram depicts only logical dependencies among work products (withdrawn from [11])

## 3.3. Threads versus Processes

Threads and processes are two mechanisms to design an application to perform multiple tasks concurrently. A single process can contain multiple threads. In this section, it is briefly presented some of the factors that might influence the choice of whether to implement an application as a group of threads or as a group of processes.

The advantages of a multithreaded approach are [13]:

- Data sharing: Sharing data between threads is easy, as all of them share the same data and heap address spaces. By contrast, sharing data between processes requires the usage of an Inter-Process Communication (IPC) mechanism.

### 3.3. Threads versus Processes

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- Context switching: Thread creation is faster than process creation; context-switch time may be lower for threads than for processes.

Using threads can have some disadvantages compared to using processes [13]:

- Thread safety: When programming with threads, one needs to ensure that the functions we call are thread-safe, i.e., can be invoked by multiple threads at the same time. Multiprocess applications don't need to be concerned with this.
- Isolation: A bug in one thread (e.g., modifying memory via an incorrect pointer) can damage all of the threads in the process, since they share the same address space and other attributes. By contrast, processes are more isolated from one another.
- Memory usage: Each thread is competing for use of the finite virtual address space of the host process. In particular, each thread's stack and thread-specific data (or thread-local storage) consumes a part of the process virtual address space, which is consequently unavailable for other threads. Although the available virtual address space is large, this factor may be a significant limitation for processes employing large numbers of threads or threads that require large amounts of memory. By contrast, separate processes can each employ the full range of available virtual memory (subject to the limitations of RAM and swap space).

Summarizing, the key factors to consider when designing a concurrent application (multithread or multiprocess) are [13]:

- In a multithreaded process, multiple threads are concurrently executing the same program. All of the threads share the same global and heap variables, but each thread has a private stack for local variables. The threads in a process also share a number of other attributes, including process ID, open file descriptors, signal dispositions, current working directory, and resource limits.
- The key difference between threads and processes is the easier sharing of information that threads provide, and this is the main reason that some application designs map better onto a multithread design than onto a multiprocess design.
- Threads can also provide better performance for some operations (e.g., thread creation is faster than process creation), but this factor is usually secondary in influencing the choice of threads versus processes.

#### 3.3.1. Pthreads API

In the late 1980s and early 1990s, several different threading APIs existed. Thus, a standard and portable implementation was required, leading to standardization of the POSIX threads API — Pthreads — by the POSIX.1c in 1995 [13].

Pthreads is a standardized model for dividing a program into subtasks whose execution can be interleaved or run in parallel [12].

### Thread creation

When a program is started, the resulting process consists of a single thread, called the initial or main thread. Additional threads can be created using the function `pthread_create()` [13]:

```
1 #include <pthread.h>
int pthread_create(pthread_t *thread, const pthread_attr_t *attr,
3                 void *(*start)(void *), void *arg);
```

This function takes as arguments a pointer to a buffer of type `pthread_t` – into which the unique identifier for this thread is copied (thread ID) before `pthread_create()` returns, the thread attributes, a function pointer containing the so called worker function, and the worker function arguments.

The new thread then starts execution by calling the function identified by `start` with the argument `arg` (i.e. `start(args)`). The thread that calls `pthread_create()` continues execution with the next statement that follows the call. The `arg` argument is declared as `void *`, allowing generic data to be passed to the worker function. This function returns 0 on success, or a positive number indicating the error occurred.

### Thread termination

The execution of a thread terminates in one of the following ways [13]:

- the thread's worker function performs a `return` specifying a return value for the thread;
- the thread calls `pthread_exit()`;
- the thread is canceled using `pthread_cancel()`;
- any of the thread calls `exit()`, or the main thread performs a `return`, causing all threads in the process to terminate immediately.

The `pthread_exit()` function terminates the calling thread and returns a value via `retval` that is available to another thread in the same process that calls `pthread_join()`.

```
1 #include <pthread.h>
void pthread_exit(void *retval);
```

### Joining with a terminated thread

The `pthread_join()` function waits for the thread identified by `thread` to terminate [13].

```
1 #include <pthread.h>
2 int pthread_join(pthread_t thread, void ** retval);
```

It is important to note that:

- if the thread has already terminated, `pthread_join()` returns immediately;
- calling `pthread_join()` for a thread ID that has been previously joined can lead to unpredictable behavior;
- if a thread is not detached, it must be joined with `pthread_join()`, otherwise it produces a ‘zombie’ thread, wasting system resources.

#### Detaching a thread

By default, a thread is joinable, meaning that when it terminates, another thread can obtain its return status using `pthread_join()`. Sometimes, the thread’s return status is irrelevant: one simply wants the system to automatically clean up and remove the thread when it terminates. In this case, the thread can be detached, by making a call to `pthread_detach()` specifying the thread’s identifier in `thread` [13].

```
1 #include <pthread.h>
2 int pthread_detach(pthread_t thread);
```

Once a thread has been detached, it is no longer possible to use `pthread_join()` to obtain its return status, and the thread can’t be made joinable again. Another important note is that `pthread_detach()` only controls what happens after a thread terminates, not how or when it terminates. If another thread calls `exit()` or the main thread returns, all threads in the process are immediately terminated, regardless of whether they are joinable or detached.

## 3.4. Communications

The communications technologies and the associated tools used for the project development are briefly described next.

### 3.4.1. IEEE 802.11 – Wi-Fi

IEEE 802.11, commonly known as Wi-Fi, is part of the IEEE 802 set of Local Area Network (LAN) protocols, and specifies the set of Media Access Control (MAC) and physical layer protocols for implementing Wireless local Area Network (WLAN) communication in a wide spectrum of frequencies, ranging from 2.4–60 GHz.

## TCP/IP

The most commonly used protocols for Internet communications, including Wi-Fi, are Transmission Control Protocol (TCP) and IP, usually associated together, being part of the OSI model (Fig. 3.3), which characterises and standardises the communication functions of a telecommunication or computing system, being agnostic to their underlying internal structure and technology.

A computer protocol is a standardised procedure for the exchange and transmission of data between devices, as requested for the application processes. The TCP provides services at the Transport layer, handling the reliable, unduplicated and sequenced delivery of data [14], while the UDP provides data transportation without guaranteed data delivery or acknowledgments. The TCP can be thought of a reliable version of User Datagram Protocol (UDP), generalizing. The IP part of the TCP/IP suite, providing services at the Network layer, is used to make origin and destination addresses available to route data across networks.

These protocols are applied in sequence to the user's data to create a frame that can be transmitted from the sending application to the receiving application. The receiver reverses the procedure to obtain the original user's data and pass them to the receiving application [14].

Another interesting fact, due to the technology agnostic aspect of the OSI Model, is that IP and the higher-level protocols may be implemented on several kinds of physical nets.

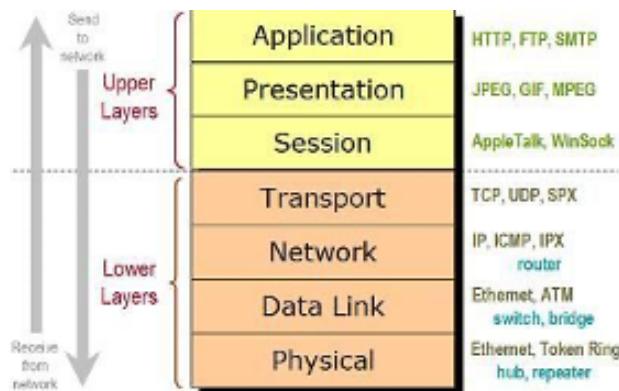


Figure 3.3.: OSI model

### 3.4.2. Network programming – sockets

Computer systems implement multiple processes which require an identifier. As such, the IP address is not enough to uniquely identify the origin/destination of data to be transmitted, and the port number is added. This combination of an IP address and port number is sometimes called a network socket [15], allowing data to be delivered to multiple processes in the same machine – same IP address. It is the socket pair (the

4-tuple consisting of the client IP address, client port number, server IP address, and server port number) that specifies the two end points that uniquely identifies each TCP connection in an internet [15].

In a broader sense, a socket can be described as a method of IPC that allows data to be exchanged between applications, either on the same host (computer) or on different hosts connected by a network [13], as a local interface to a system, created by the applications and controlled by the operating system, allowing an application process to simultaneously send and receive messages from other processes.

The Socket API was created in UNIX BSD 4.1 in 1981, with widespread implementation in UNIX BSD 4.2 [13]. It implements the Client-Server paradigm and implement several (standard) functions to access the operating system network resources, through system calls, in Linux [13].

There are two generic ways to use sockets: for outgoing connections – client socket – and for incoming connections – server socket. Fig. 3.4 illustrates the required steps to obtain a connected socket:

1. When a socket is initially created is mostly unuseful.
2. Binding the server socket associates it to an unique network tuple (address and port number), enabling it to be uniquely addressed.
3. When a socket server goes into listening mode, the remote devices can initiate the connection procedure, referring to its unique network tuple.
4. When the socket server accepts a connection, it spawns a new socket which is connected to the remote device, and the endpoints can effectively communicate. The server socket is ready to accept new incoming connections.

#### 3.4.3. Client/server model

The client/server model is the most common form of network architecture used in data communications today [17]. A client is a system or application that request the activity of a service provider system or application, called servers, to accomplish specific tasks. The client/server concept functionally divides the execution of a unit of work between activities initiated by the end user (client) and resource responses (services) to the activity request as a cooperative environment [17]. The client, typically handling user interactions and data exchange/modification in the user's behalf, makes a request for a service, and a server, often requiring some resource management (synchronization and access to the resource), performs that service, responding to the client requests with either data or status information [18].

An example of a simple client-server model using the Socket API, through system calls, is presented in Fig. 3.5. The operation of sockets can be explained as follows [13]:

- The `socket()` system call creates a new socket, establishing the protocols under which they should communicate. For both client and server to communicate, each of them must create a socket.

### 3.4. Communications

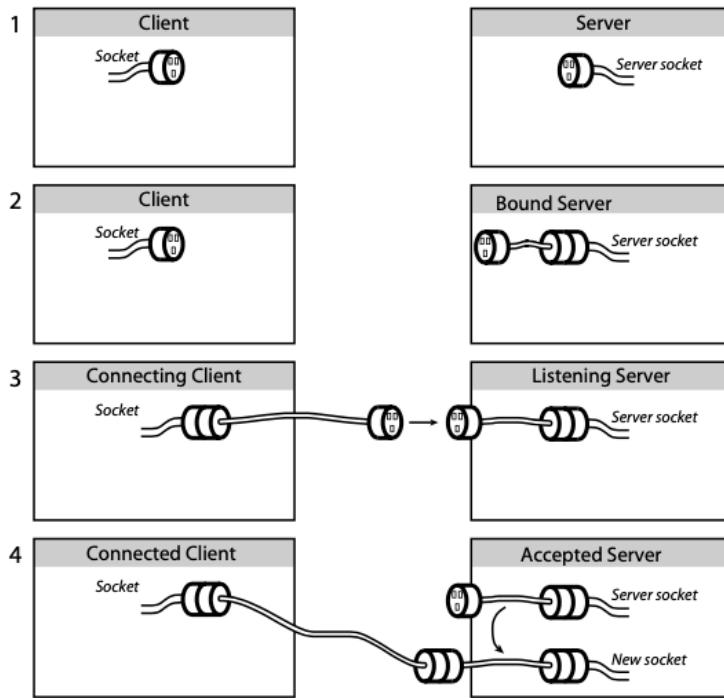


Figure 3.4.: Steps to obtain a connected socket (withdrawn from [16])

- Communication via a stream socket is analogous to a telephone call. One application must connect its socket to another application's socket before communication can take place. Two sockets are connected as follows:
  1. One application, assuming the role of server, calls `bind()` to bind the socket to a well-known address, and then calls `listen()` to notify the kernel it is ready to accept incoming connections.
  2. The other application, assuming the role of client, establishes the connection by calling `connect()`, specifying the address of the socket to which the connection is to be made.
  3. The server then accepts the connection using `accept()`. If the `accept()` is performed before the client application calls `connect()`, then the `accept()` blocks.
- Once a connection has been established, data can be transmitted in both directions between the applications (analogous to a bidirectional telephone conversation) until one of them closes the connection using `close()`.
- Communication is performed using the conventional `read()` and `write()` system calls or via a number of socket-specific system calls (such as `send()` and `recv()`) that provide additional functionality. By default, these system calls block if the Input/Output (I/O) operation can't be completed immediately. However, nonblocking I/O is also possible.

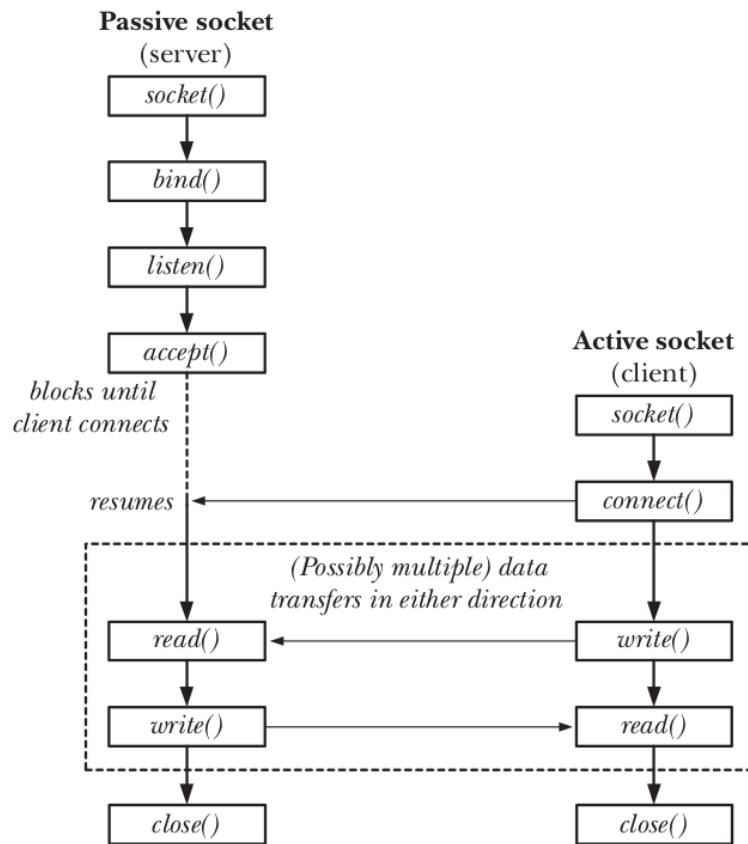


Figure 3.5.: Overview of UNIX system calls with sockets implementing a server/client paradigm (withdrawn from [13])

## 3.5. Daemons

In this section daemons are introduced, showing how to create them, handle possible errors and how to communicate with them.

### 3.5.1. What is a Daemon?

A daemon is a background process that runs without user input and usually provides some service, either for the system as a whole or for user programs. [19]

Normally, daemons are started when a system boots and, unless forcibly terminated, run until system shutdown. Because a daemon does not have a controlling terminal, any output, either to stderr or stdout, requires special handling. They often run with **superuser privilege** because they use privileged ports (1 - 1024) or because they have access to some sort of privileged resource. Generally, daemons are process

group leaders and session leaders, a daemon's parent is the init process, which has the PID of 1 (daemon is an orphan process inherited by init).

### 3.5.2. How to create a Daemon

The steps to create a daemon are:

1. Fork and exit in the parent process;
2. Create a new session in the child using the setsid call;
3. Make the root directory, "/", the child process's working directory;
4. Change the child process's umask to 0;
5. Close any unneeded file descriptor the child inherited.

#### Fork and exit in the parent process

A daemon is started from a shell script or the command line. Daemons are unlike application programs because they are not interactive, i.e., they run in the background and, as a result, do not have the controlling terminal.

The parent forks and exits as the first step toward getting rid of the controlling terminal (they only need a terminal interface long enough to get started).

```
pid_t pid;
// create a new process

4 pid = fork();
if (pid < 0) { // error trying to execute fork
6   ERROR("fork failure");
8   exit(EXIT_FAILURE);
}

10 if (pid > 0) // parent process (exit)
    exit(EXIT_SUCCESS);
12 // child process continues the execution
...
```

#### Create a new session in the child using the setsid call

Calling setsid accomplishes several things:

- It creates a new session if the calling process is not a process group leader, making the calling process the session leader of the new session;

- It makes the calling process the process group leader of the new process group;
- It sets the Process Group ID (PGID) and the Session ID (SID) to the Process ID (PID) of the calling process;
- It dissociates the new session from any controlling tty.

```
1 pid_t sid;
2 sid = setsid(); // create a new session
3 if (sid < 0) {
4     ERROR("setsid failure");
5     exit(EXIT_FAILURE);
6 }
```

- Each process is a member of a process group, which is a collection of one or more processes generally associated with each other for the purposes of job control (# cat ship-inventory.txt | grep booty | sort).
- When a new user first logs into a machine, the login process creates a new session that consists of a single process, the user's login shell. The login shell functions as the session leader.

## Make the root directory, "/", the child process's working directory

This is necessary because any process whose current directory is on a mounted file system will prevent that file system from being unmounted.

Making "/" a daemon's working directory is a safe way to avoid this possibility.

```
// make '/' the root directory
2 if (chdir("/") < 0) {
3     ERROR("chdir failure");
4     exit(EXIT_FAILURE);
5 }

// continue executing the child process
8 ...
```

## Change the child process's umask to 0

This step is necessary to prevent the daemon's inherited umask from interfering with the creation of files and directories.

Consider the following scenario:

- A daemon inherits a umask of 055, which masks out read and execute permissions for group and other.
- Resetting the daemon's umask to 0 prevents such situation.

```
1 // resetting umask to 0
2 umask(0);

3 // continue executing the child process
4 ...
5 ...
```

### Close any unneeded file descriptor the child inherited

This is simply a common sense step. There is no reason for a child to keep open descriptors inherited from parent. The list of potential file descriptors to close includes at least stdin, stdout and stderr.

```
1 // close unneeded file descriptors
2 close(STDIN_FILENO);
3 close(STDOUT_FILENO);
4 close(STDERR_FILENO);

5 // continue executing the child process
6 ...
7 ...
```

### 3.5.3. How to handle errors

There's the **problem** that once a daemon calls setsid, it no longer has the controlling terminal an so it has nowhere to send output that would normally go to stdout or stderr (such as error messages).

The **solution** is that, fortunately, the standard utility for this purpose is the **syslog** service, provided by the system logging daemon, **syslogd**.

#### Handling Errors with syslog

syslogd is a daemon that allow to save log messages from other daemons or applications. The relevant interface is defined in <syslog.h> header file. The API is simple, **openlog** opens the log, **syslog** writes a message to it, and **closelog** close the log.

The function prototypes are listed here:

```
1 #include <syslog.h>

3 void openlog(char *ident, int option, int facility);
4 void closelog(void);
5 void syslog(int priority, char *format, ...);
```

### 3.5.4. Communicating with a Daemon

To communicate with a daemon, you send it signals that cause it to respond in a given way. For example, it is typically necessary to force a daemon to reread its configuration file. The most common way to do this is to send a **SIGHUP** signal to the daemon.

When you execute the command "kill PID" on command line, the signal **SIGINT** is sent to daemon to terminate the daemon execution.

## 3.6. Device drivers

A device driver is a small piece of software that tells the operating system and other software how to communicate with a piece of hardware [20].

When the kernel recognizes that a certain action were requested to a device, it calls an appropriate function from the driver, and transfers the process control from the user to the driver function. After the driver function ends its execution, it gives the control back to the user space process.

The device driver provides the following characteristics:

- i. A set of functions to communicate with the hardware device, and a standardized kernel interface;
- ii. It shoud be an autonomous component, able to be dynamically added to and removed from the OS;
- iii. Data flux control between the user space program and device;
- iv. A user defined section, so the device driver can be visible as a node in the /dev, for the OS.

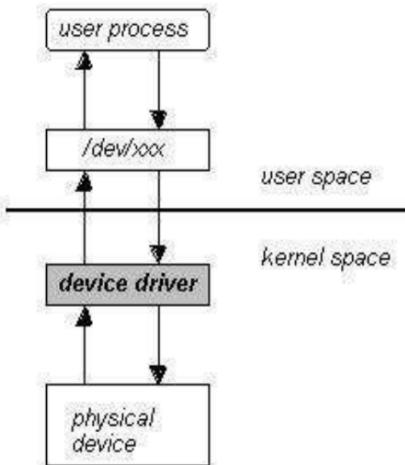


Figure 3.6.: Example of the usage of a device driver (withdrawn from [21])

Character device drivers and Block device drivers are the two of the most important device drivers. Block device drivers are accessed by the user space program by a system buffer that works like a data cache.

There's no need for management or allocation routines as the system transfers the data from/to the device. Character device drivers communicate directly with the user space program, so no buffer is required [21].

### 3.6.1. Kernel mode vs Application

#### How can an application access module services?

- i. When one of the kernel functions associate to the driver is called, the function `module_init()` calls the function `register_capability()` to register the driver services on kernel - on character drivers this is done by `chrdev_region()` and `alloc_chrdev_region()`.
- ii. The function `register_capability()` saves a pointer to an argument in the internal data structure `capabilities[]`.
- iii. The system defines a protocol of how the application will have access to `capabilities[]` by system calls.
- iv. When the driver is no more necessary, the module can be unregistered from kernel, freeing the entry on `capabilities[]` using the function `module_exit()`.

The modules run on kernel space, while the applications run on user space.

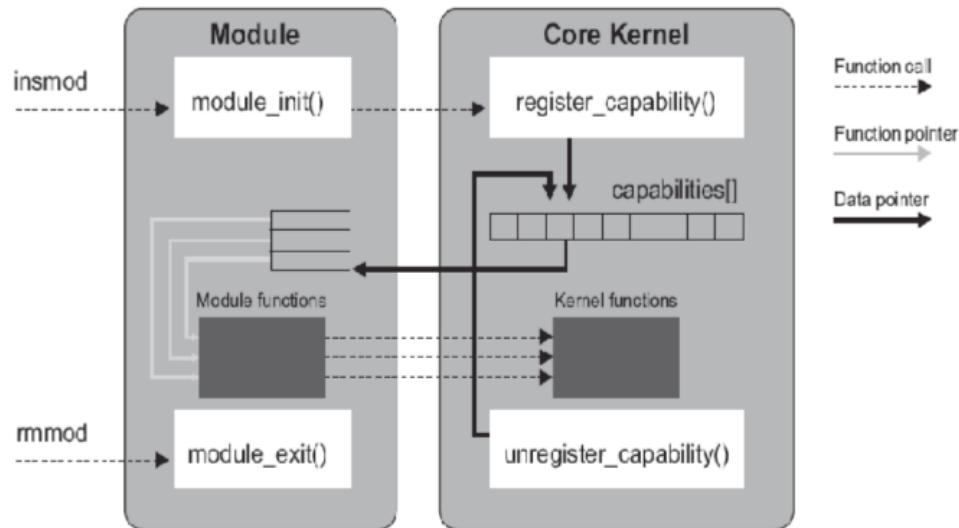


Figure 3.7.: Application accessing to module services (withdrawn from [21])

Linux kernel offers a set of functions to the user space for the interaction with hardware. Linux represents and manages a byte sequence using files. Almost all devices are represented by a sequence of bytes, so linux also represents the I/O devices as special files called nodes with entries in /dev directory.

### 3.7. Build system and Makefiles

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In the kernel space, Linux offers a set of functions that interact directly with the hardware and allows data transfer between kernel space and user space.

Normally, for each function in user space there is a matching function in the kernel space allowing data transfer between kernel and user space.

Events	User Functions	Kernel Functions
Load a module	insmod	module_init()
Open a device	fopen( )	file_operations: open
Read from a device	fread( )	file_operations: read
Write to a device	fwrite( )	file_operations: write
Close a device	fclose( )	file_operations: release
Remove a module	rmmod	module_exit()

Figure 3.8.: Examples of drivers event and corresponding interface functions between kernel space and user space (withdrawn from [21])

To associate the normal files to the kernel mode, the linux system uses the major number and the minor number:

- **Major number** is normally used by linux system to map I/O requests to driver code. By other words, identifies the driver associated to the device.
- **Minor number** is dedicated to internal use and it is used by kernel to determine exactly which device was referenced. Depending on the way the driver was coded, this number can have a direct meaning, as a device counter, or it can be a simple driver iterator on the local devices array on kernel.

To do the association, it should be created a file or node in /dev directory, (calling mknod as root user) that will be used to access the driver.

## 3.7. Build system and Makefiles

Make is essentially a dependency-tracking build utility, which is most commonly used as a build automation tool to generate executable programs and libraries from source code by reading files called **Makefiles**. These files specify how to derive the target program, defining the prerequisites (dependencies) and the rule for generating it [22]. Nonetheless, the **Make** utility is also used as automation build tool for compiling **LaTeX** documents or to manage the workflow of data analysis projects [23].

Make was originally created by Stuart Feldman in April 1976 at Bell Labs, after the experience of a coworker in futilely debugging a program of his where the executable was accidentally not being updated with changes [24]:

'Make originated with a visit from Steve Johnson (author of yacc, etc.), storming into my office, cursing the Fates that had caused him to waste a morning debugging a correct program (bug had been fixed, file hadn't been compiled, cc \*.o was therefore unaffected). As I had spent a part of the previous evening coping with the same disaster on a project I was working on, the idea of a tool to solve it came up. It began with an elaborate idea of a dependency analyzer, boiled down to something much simpler, and turned into Make that weekend. Use of tools that were still wet was part of the culture. Makefiles were text files, not magically encoded binaries, because that was the Unix ethos: printable, debuggable, understandable stuff.'

— Stuart Feldman, The Art of Unix Programming, Eric S. Raymond 2003

Before Make's introduction, the Unix build system most commonly consisted of OS dependent `make` and `install` shell scripts accompanying their program's source code. Make enabled the combination of commands for the different targets into a single file and to abstract out dependency tracking and archive handling, making it a cornerstone to pave the way for modern build environments.

Several versions of the Make utility were implemented, by porting them into different OSes or by rewriting from scratch, like the `BSD Make` (Berkeley Software Distribution (BSD) systems), `GNU make` (Linux and MacOS) or `nmake` (Windows). The `GNU make` is the standard implementation for Linux- and MacOS-based systems [25], and, thus, it will be the one being addressed here.

The GNU implementation of `make` — written by Richard Stallman and Roland McGrath [22] — tracks down the targets that need to be recompiled and executes the defined commands in the `makefile` containing the rules for those targets. If the flag `-f` is not provided to `make`, it will look for makefiles `GNUmakefile` (not recommended — only if the file is specific to `GNU make`), `makefile`, and `Makefile` in that order [22]. `make` will then update the target if it depends on prerequisite files that have been modified since the target was last modified, or if the target does not exist.

#### 3.7.1. Makefile syntax and example

A simple `makefile` consists of rules with the following shape [22]:

```
target ... : prerequisites ...
<TAB>recipe
<TAB>...
<TAB>...
```

The basic syntax can be explained as follows [22]:

- target: it is usually the name of a file that is generated by a program; examples of targets are executable or object files. A target can also be the name of an action to carry out, such as `clean`, being named as phony targets.
- prerequisite: file that is used as input to create the target. A target often depends on several files.
- recipe: it is an action that make carries out. A recipe may have more than one command, either on the same line or each on its own line. A `tab` character is required at the beginning of every recipe line.

Usually a recipe is in a rule with prerequisites and serves to create a target file if any of the prerequisites change. However, the rule that specifies a recipe for the target need not have prerequisites. For example, the rule containing the `delete` command associated with the target `clean` does not have prerequisites.
- rule: explains how and when to remake certain files which are the targets of the particular rule. `make` carries out the recipe on the prerequisites to create or update the target. A rule can also explain how and when to carry out an action.

Listing 3.1 presents an `Makefile` example, with a preamble documenting the syntax for easier comprehension (lines 1 through 51). Line 54 defines the project name by assigning it to a variable `PROJ`. Line 57 and 58 define the file paths for source code and binaries using macro expansion. Line 61 defines the compilation flags — `CFLAGS` — which is known makefile variable. Line 62 defines the command `RM` to remove files and folders recursively. Lines 65 and 66 define the `Doxygen` documentation paths. Lines 69 and 71 retrieve database and C-source files by using the wildcard expansion which matches the patterns (globs) defined after. Lines 74 through 79 exemplify the usage of conditionals to add different compilation flags. Line 81 comment shows a bad practice, which is to define the `CC` variable — `C Compiler` — to a specific compiler, thus, making it not portable. The alternative is to set the `CC` outside of the makefile. Line 83 finalizes the initialization part of the makefile by performing pattern substitution to obtain all object files from C-source files.

Line 86 starts the rules part of the makefile, defining the default rule and stating that to build `$(PROJ)` executable it is required the object files (`$(OBJ)`) and one can compile it using the `C compiler`, the dependencies and the required libraries. When `make` tries to evaluate the dependency list, if the object files haven't been produced yet, it will generate them using the rule from line 95 and then it can execute rule `$(PROJ)`.

Line 101 states that calling `make all` will invoke the default rule and echo ‘Compiling project’.

Line 105 states that calling `make install` will invoke the `all` and `clean` rules to compile the application and install the binaries at the appropriate location.

Line 111 contains the phony targets, i.e., targets whose names represent actions like `clean`, `mrproper`,

### 3.7. Build system and Makefiles

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clean-all, and doc: clean removes all objects files, mrproper and clean-all will remove also the executable file generated, and doc will generate the documentation using Doxygen.

```
1 ##### Makefile #####
# a Makefile works accordingly to the following syntax:
# target: dependencies
#   <TAB> command1
#   <TAB> command2
6 #
# - all: the name of the rules to be executed
# - teste.o: teste.c (can be interpreted as destination_file: origin_file)
# - clean: deletes the intermediary files.
# - mrproper: deletes everything that must be modified
11 #
# ----- Defining variables -----
# VAR_NAME=VAL
# and to use the variable we placed between $(): $(VAR_NAME)
# #Using wildcards to help define variables
16 # SRC=$(wildcard *.c)
# OBJ=$(SRC:.c=.o)
# -----
# #Silent Mode (no echo in the command line)
# add @ to the beginning of the command. Ex:
21 #      @$(CC) -o $@ $^
# ----- Internal Variables -----
# $@ Name of the rule
# $< Name of the 1st dependency
# $^ List of dependencies
26 # $? List of dependencies more recent than the rule
# $* Name of the file with suffix
# -----
# Rules without dependencies, e.g., actions like clean:, must be called
# explicitly, otherwise they wont be executed.
31 # - make clean will remove all objects files from folder, as described below
# ----- Interference rules -----
# Generic rules called by default
# - .c.o: make a .o file from a .c file
# - %.o: %.c: the same thing. The line teste.o:teste.c can be modified with this
36 # rule
# - .PHONY: This rule avoids conflicts.
# - If one has a clean file in the directory, nothing will happen when one
# executes make clean
# - So .PHONY says that clean and mrproper should be executed even if
41 # files with that name exist current path
```

### 3.7. Build system and Makefiles

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```
# -----
# ----- make install -----
# Automating the program's installation with the install: rule
#   - install: Places the binary or executable inside a given folder , e.g., /bin
46 #   or /usr/bin in Linux. It can be any other , using the command mv or cp to
#   move or copy
#   - Lets create 2 vars:
#     - prefix=path/to/proj
#     - bindir=$(prefix)/bin
51 #   - and we add the rule install:all

# PROJ name
PROJ=Electric -gym

56 # File Paths
SRC_DIR=.
BIN_DIR=${SRC_DIR}/bin

#LIBS=-lform -lncurses
61 CFLAGS=-Wall #-W -ansi -pedantic # options passed to the compiler
RM=rm -rf

# documentation
DOXYGEN=/Applications/Doxygen.app/Contents/Resources/doxygen
66 DOXYFILE=../Doxyfile

# Databases
DB = $(wildcard *.db)

71 SRC := $(wildcard *.c)

# Debug settings
DEBUG ?= 1
ifeq ($(DEBUG), 1)
76     CFLAGS+=-g
else
    CFLAGS+=-DNDEBUG
endif

81 # CC=gcc #defining the compiler name for C or C++ (here its gcc) // bad practice
# to define it (not portable)
OBJ := $(SRC:.c=.o)
```

### 3.7. Build system and Makefiles

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```
# Building executable (w/ linking)
86 $(PROJ): $(OBJ)
    # the compiler does the linking between the 2 objects
    # gcc -o teste teste.o main.o
    # $@ = teste:      #Rule name
    # $^ = teste.o main.o  #dependency list
91    @echo "Creating executable"
    $(CC) -o $@ $^ ${LIBS}

# Building object files from sources
%.o: ${SRC_DIR}/%.c # make a .o file from a .c file
96 # gcc -o teste.o -c teste.c -W -Wall -ansi -pedantic
    # $< Name of the 1st dependency
    @echo "Building object files"
    $(CC) -o $@ -c $^ $(CFLAGS)

101 all: $(PROJ)
    @echo "Compiling project"

# Install: run make and then make install
install: all clean
106    @echo "Installing binaries"
    @mkdir -p $(BIN_DIR)
    @mv $(PROJ) $(BIN_DIR)/
    @mv $(BIN_DIR) ../

111 .PHONY: clean mrproper clean-all doc
clean-all: clean mrproper
clean:
    @- $(RM) $(OBJ)
    # @- $(RM) $(DB)
116 mrproper: clean
    @$(RM) $(PROJ)
    # Documentation
doc:
    @echo "Generating documentation"
121    @$(DOXYGEN) $(DOXYFILE)
```

Listing 3.1: Makefile example

### 3.7.2. Other build systems for C/C++

Make is just part of the simplest building system for C/C++. But there are other ones which are noteworthy[26]:

- Autotools – Automake and Autoconf: In 1991, David J. MacKenzie got tired of customizing Makefile for the 20 platforms he had to deal with. Instead, he handcrafted a little shell script called configure to automatically adjust the Makefile. Compiling his package was now as simple as running `./configure && make` [27]. The Autotools are tools that will create a GNU Build System for your package. Autoconf mostly focuses on `configure` and `automake` on Makefiles.

Simply put, `automake` aims to allow the programmer to write a makefile in a higher-level language, rather than having to write the whole makefile manually. However, the system is very complex and has a steep learning curve.

- CMake: CMake reads the projects to build from `CMakeLists.txt` files written in a language of its own. From there it generates some Makefiles (or equivalent project files for Xcode or Visual Studio for example) that can be used to build the project.

The main criticism against CMake is its language, requiring another syntax to be learned and that it generates lot of intermediate files.

## 3.8. Source code documentation

### 3.8.1. Doxygen

## 3.9. Nebulizing technologies

aaaaaaaa

## 3.10. Computer vision

Computer vision is a vast field, but can broadly be defined as the transformation of data from a still image or video camera into either a decision or a new representation to achieve some particular goal [28].

The input data may include some contextual information such as ‘the camera is mounted in a car’ or ‘laser range finder indicates an object is 1 meter away.’ The decision might be ‘there is a person in this

scene' or 'there are 14 tumor cells on this slide.' A new representation might mean turning a color image into a grayscale image or removing camera motion from an image sequence [28].

In this section, computer vision frameworks/libraries are listed, with special focus on OpenCV, and computer vision algorithms for face detection and hand gesture recognition are analyzed.

### **3.10.1. Computer vision frameworks**

There are several noteworthy computer vision frameworks, namely [29]:

1. [Google Cloud's Vision API](#): it is an easy-to-use image recognition technology that lets developers understand the content of an image by applying powerful machine learning models. It enables key vision detection features within an application – face, and landmark detection, image labeling, Optical Character Recognition (OCR), and explicit content tagging – and image classification into millions of predefined categories.
2. [YOLOv3](#): YOLO (You Only Look Once) is a state-of-the-art, real-time object detection system among the most widely used deep learning-based object detection methods. It considers object detection as a regression problem, directly predicting the class probabilities and bounding box offsets from full images with a single feed-forward Convolutional Neural Network (CNN). YOLOv3 eliminates region proposal generation and feature resampling and encapsulates all stages in a single network to form a true end-to-end detection system.
3. [TensorFlow](#): it is a free, open-source framework for creating algorithms to develop a user-friendly Graphical Framework called TensorFlow Graphical Framework (TF-GraF) for object detection API, which is widely applied to solve complex tasks efficiently in agriculture, engineering, and medicine. The TF-GraF provides independent virtual environments for amateurs and beginners to design, train, and deploy machine intelligence models without coding or CLI in the client-side.
4. [libfacedetection](#): it is an open-source library for face detection in images. It uses a pre-trained CNN, enabling face detection on inputs with a size greater than  $10 \times 10$  pixels. The source code is not dependant on any other libraries. A C++ compiler is required to compile the source under various platforms, such as Windows, Linux, ARM, etc..
5. [Raster Vision](#): it is an open-source Python framework to build computer vision models on satellite, aerial, and other large sets of images (including oblique drone imagery), using deep learning or machine learning models. It has built-in support for chip classification, object detection, and semantic segmentation with backends using PyTorch and Tensorflow. The framework is also extensible to new data sources, tasks (e.g., object detection), backend (e.g., TF Object Detection API), and cloud providers.

6. SOD: it is an embedded, modern cross-platform computer vision and machine learning software library. It exposes a set of APIs for deep-learning, advanced media analysis, and processing, including real-time, multi-class object detection, and model training on embedded systems with limited computational resource and Internet of Things (IOT) devices. Designed for computational efficiency and with a strong focus on real-time applications, SOD includes a comprehensive set of both classic and state-of-the-art deep-neural networks with their pre-trained models. Although it is open source, the pre-trained models are charged (one time fee only – up to 30 United States Dollar (USD)).
7. Face\_recognition: it is a facial recognition API for Python and the command line, built with deep learning using dlib60's state-of-the-art face recognition. The model has an accuracy of 99.38% on the Labeled Faces in the Wild benchmark.
8. JeelizFaceFilter: it is a lightweight and robust face tracking library, designed for augmented reality face filters. This JavaScript library can detect and track the face in real-time from the webcam video feed captured, enabling the developers to solve computer-vision problems directly from the browser. The key features include face detection, face tracking, face rotation detection, mouth opening detection, multiple face detection, and tracking, video acquisition with High-Definition (HD) video ability, etc.
9. OpenCV: it is an open-source computer vision and machine learning software library, built to provide a common infrastructure for computer vision applications and accelerate the use of machine perception in commercial products. OpenCV was designed for computational efficiency and with a strong focus on real-time applications. It is written in optimized C++ and can take advantage of multicore processors, with wrappers written in Python, Java, and Matlab, and supporting Windows, Linux, Android and Mac OS. The library has more than 2500 optimized algorithms, including a comprehensive set of both classic and state-of-the-art computer vision and machine learning algorithms. These algorithms can be used to detect and recognize faces, identify objects, classify human actions in videos, track camera movements, track moving objects, extract 3D models of objects and produce 3D point clouds from stereo cameras. It can stitch images together to produce a high-resolution image of an entire scene, find similar images from an image database, remove red eyes from images taken using flash, follow eye movements, recognize scenery and establish markers to overlay it with augmented reality.

### **3.10.2. Face detection**

Face detection has been studied for decades in the computer vision literature. Modern face detection algorithms can be categorized into four categories [30]: cascade based methods, part based methods, channel feature based methods, and neural network based methods. A detailed survey can be found in

[31, 32].

The seminal work by Viola and Jones [33] introduces integral image to compute Haar-like features in constant time. These features are then used to learn AdaBoost classifier with cascade structure for face detection. Various later studies follow a similar pipeline, with SURF cascade [34] achieving competitive performance.

Although the Viola-Jones are extremely fast, they require severe tuning to prevent false-positives or complete misses. In 2005, Dalal and Triggs [35] demonstrated that the Histogram of Oriented Gradients (HOG) image descriptor and a Linear Support Vector Machine (SVM) could be used to train highly accurate object classifiers – or in their particular study, human detectors. However, due to the nature of HOG that are only suited for frontal face detection as it is not invariant to changes in rotation and viewing angle.

One of the well-known part based methods is Deformable part models (DPM) [36]. Deformable part models define face as a collection of parts and model the connections of parts through Latent SVM. The part based methods are more robust to occlusion compared with cascade-based methods.

Aggregated channel feature (ACF) is first proposed by Dollar et al. [37] to solve pedestrian detection. Later on, Yang et al. [38] applied this idea on face detection. In particular, features such as gradient histogram, integral histogram, and color channels are combined and used to learn boosting classifier with cascade structure.

Recent studies [39, 40] show that face detection can be further improved by using deep learning, leveraging the high capacity of deep convolutional networks.

Considering a more pragmatic approach, the following tips are useful when selecting a face detection method [41]:

1. OpenCV's Haar cascades: Use OpenCV's Haar cascades when speed is your primary concern (e.g., when considering an embedded device like the Raspberry Pi). Haar cascades aren't as accurate as their HOG + Linear SVM and deep learning-based counterparts, but they compensate it in raw speed. False-positive detections is highly likely and parameter tuning is required when calling `detectMultiScale`.
2. HOG + linear SVM detector: Use dlib's HOG + Linear SVM detector when Haar cascades are not accurate enough, but you cannot commit to the computational requirements of a deep learning-based face detector. The HOG + Linear SVM object detector is a classic algorithm in the computer vision literature that is still relevant today. However, running HOG + Linear SVM on a Central Processing Unit (CPU) will likely be too slow for an embedded device.
3. CNN face detection: Use dlib's CNN face detection when requiring extremely accurate face detections. However, there is a tradeoff – with higher accuracy comes slower run-time. This method cannot run in real-time on a laptop/desktop CPU, and even with Graphics Processing Unit (GPU) acceleration, you'll struggle to hit real-time performance.

4. DNN face detector: Use OpenCV's DNN face detector as a good balance. As a deep learning-based face detector, this method is accurate — and since it's a shallow network with a single-shot detector backbone, it's easily capable of running in real-time on a CPU. Furthermore, since you can use the model with OpenCV's `cv2.dnn` module, that also implies that (1) you can increase speed further by using a GPU or (2) utilizing the Movidius NCS on your embedded device.

## Haar cascade classifier

Object detection using Haar feature-based cascade classifiers is an effective method proposed by Viola and Jones [42]. It is a machine learning based approach where a cascade function is trained from a lot of positive and negative images. It is then used to detect objects in other images [43].

The algorithm can be explained in four stages [44]:

1. Calculating Haar Features: the first step is to collect the Haar features. A Haar feature is essentially calculations that are performed on adjacent rectangular regions at a specific location in a detection window. The calculation involves summing the pixel intensities in each region and calculating the differences between the sums. Some examples of Haar features are illustrated in Fig. 3.9.
2. Creating Integral Images: instead of computing at every pixel, sub-rectangles and its array references are created, which are then used to compute the Haar features (see Fig. 3.10).
3. Using Adaboost: to determine the best features from the hundreds of thousands of Haar features, one requires a machine learning model. Adaboost essentially chooses the best features and trains the classifiers to use them. It uses a combination of ‘weak classifiers’ to create a ‘strong classifier’ that the algorithm can use to detect objects. Weak learners are created by moving a window over the input image, and computing Haar features for each subsection of the image. This difference is compared to a learned threshold that separates non-objects from objects. Because these are ‘weak classifiers’, a large number of Haar features is needed for accuracy to form a strong classifier (see Fig. 3.11).
4. Implementing Cascading Classifiers: The cascade classifier is made up of a series of stages, where each stage is a collection of weak learners (see Fig. 3.12). Weak learners are trained using boosting, which allows for a highly accurate classifier from the mean prediction of all weak learners.  
Based on this prediction, the classifier either decides to indicate an object was found (positive) or move on to the next region (negative). Stages are designed to reject negative samples as fast as possible, because a majority of the windows do not contain anything of interest.  
It’s important to maximize a low false negative rate, because classifying an object as a non-object will severely impair your object detection algorithm, so it is critical to tune hyperparameters accordingly when training your model.

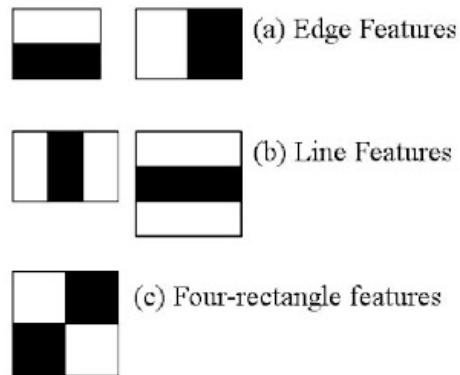


Figure 3.9.: Examples of Haar features (withdrawn from [43])

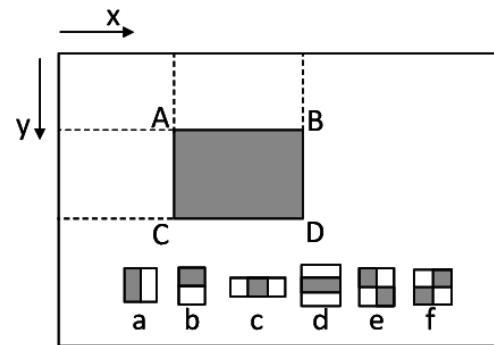


Figure 3.10.: Integral image creation illustration (withdrawn from [44])

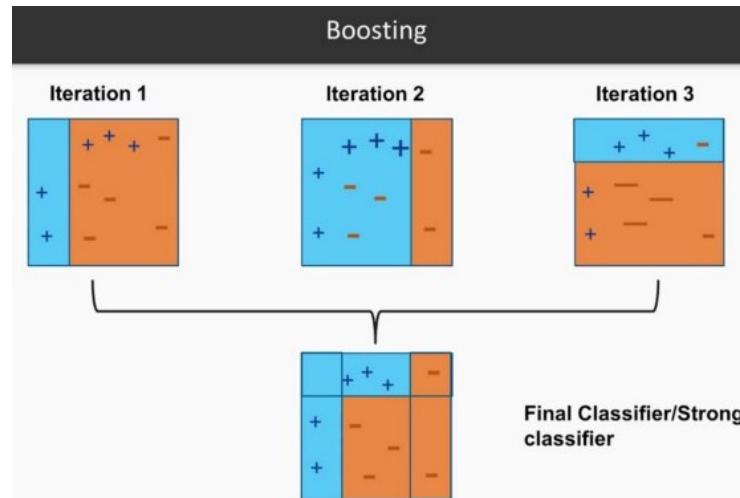


Figure 3.11.: Illustration of a boosting algorithm (withdrawn from [44])

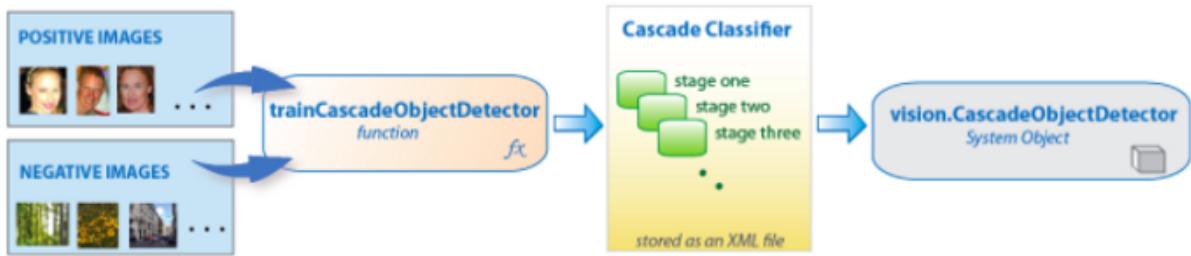


Figure 3.12.: Flowchart of a cascade classifiers (withdrawn from [44])

### 3.10.3. Hand gesture recognition

Human communication naturally relies on hand gestures for clarification of intentions, thus making a viable bridge for interchanging information between humans and computers [45]. The most notable applications of hand gesture recognition are: sign language comprehension, robotics control, assisting disabled people, creation of contactless user interfaces, game industry (e.g. Kinect), home automation, clinical and health, virtual environments, etc. [46].

However, the computational performance of recognizing hand gestures is affected by the environmental surroundings (such as light, background, distance, skin color) and the position and direction of hand [47]. Additionally, hand gestures applications require users to be well trained at employing and understanding the meaning of different gestures, as for each application, a different group of gestures is used to perform its operations [46].

In this section the hand gesture recognition process is explored, concerning its workflow and most relevant methods.

The basic hand gesture recognition procedure is illustrated in Fig. 3.13, comprising the following stages [46]:

- image frame acquisition: capture the human gesture image by computer, usually performed using a web camera or depth camera [48]. Special tools can be used such as wired or wireless gloves to detect hand movements, and motion sensing input devices (e.g., Kinect, Leap Motion, etc.) to capture the hand gestures and motions.
- Hand tracking: is the ability of the computer to trace the user hand and separate it from the background and from the surroundings objects [48]. This can be done using multi-scale color feature hierarchies that gives the users hand and the background different shades of colors to be able to identify and remove the background – background subtraction and thresholding, or by using clustering algorithms that are capable of treating each finger as a cluster and removing the empty spaces in-between them [46].
- Feature extraction: The features vary depending on the application, with the most common being:

### 3.11. RDBMS

fingers status, thumb status, skin color, alignments of fingers, and the palm position [48]. Several feature extraction methods can be used, such as Fourier descriptor method for capturing the palm, the fingers and the finger tips, or the centroid method which captures the essential structure of the hand [46].

- Classification: The features extracted are then sent to training and testing the classification algorithm (such as Artificial Neural Networks (ANN), K-nearest neighbor (KNN), Naive Bayes (NB), SVM, etc.) to reach the output gesture. A simple case of an output gesture can contain two classes to detect only two gestures such as open and closed hand gestures [46].

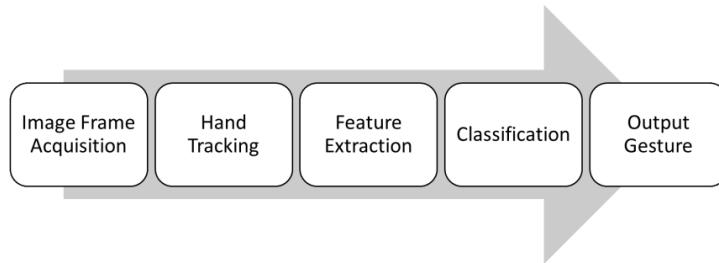


Figure 3.13.: Basic steps of hand gesture recognition (withdrawn from [46])

Fig. 3.14 illustrates an hand gesture recognition workflow, materializing the aforementioned procedure, namely: building a gesture dataset, tracking the hand and extracting the gesture by creating a Region Of Interest (ROI) and applying background subtraction and binary thresholding, training the model with the new dataset, and perform the hand gesture recognition using the predictive model.

## 3.11. RDBMS

A database is a collection of data, typically describing the activities of one or more related organizations. For example, a university database might contain information about the following [51]:

- Entities: such as students, faculty, courses and classrooms
- Relationships between entities: such as students' enrollment in courses, faculty teaching courses, and the use of rooms for courses.

An Database Management System (DBMS) is a software designed to assist in maintaining and utilizing large collections of data. A Relational Database Management System (RDBMS) is a subset of Database Management System (DBMS) with relationship between tables (entities) and rows (entities' attributes). It follows the relational model, introduced by E.F. Codd in 1970 [51], instead of navigational model, where in the data is stored in multiple tables. The tables are related to each other using primary and foreign keys. It is the most used database model widely used by enterprises and developers for storing complex and

### 3.11. RDBMS

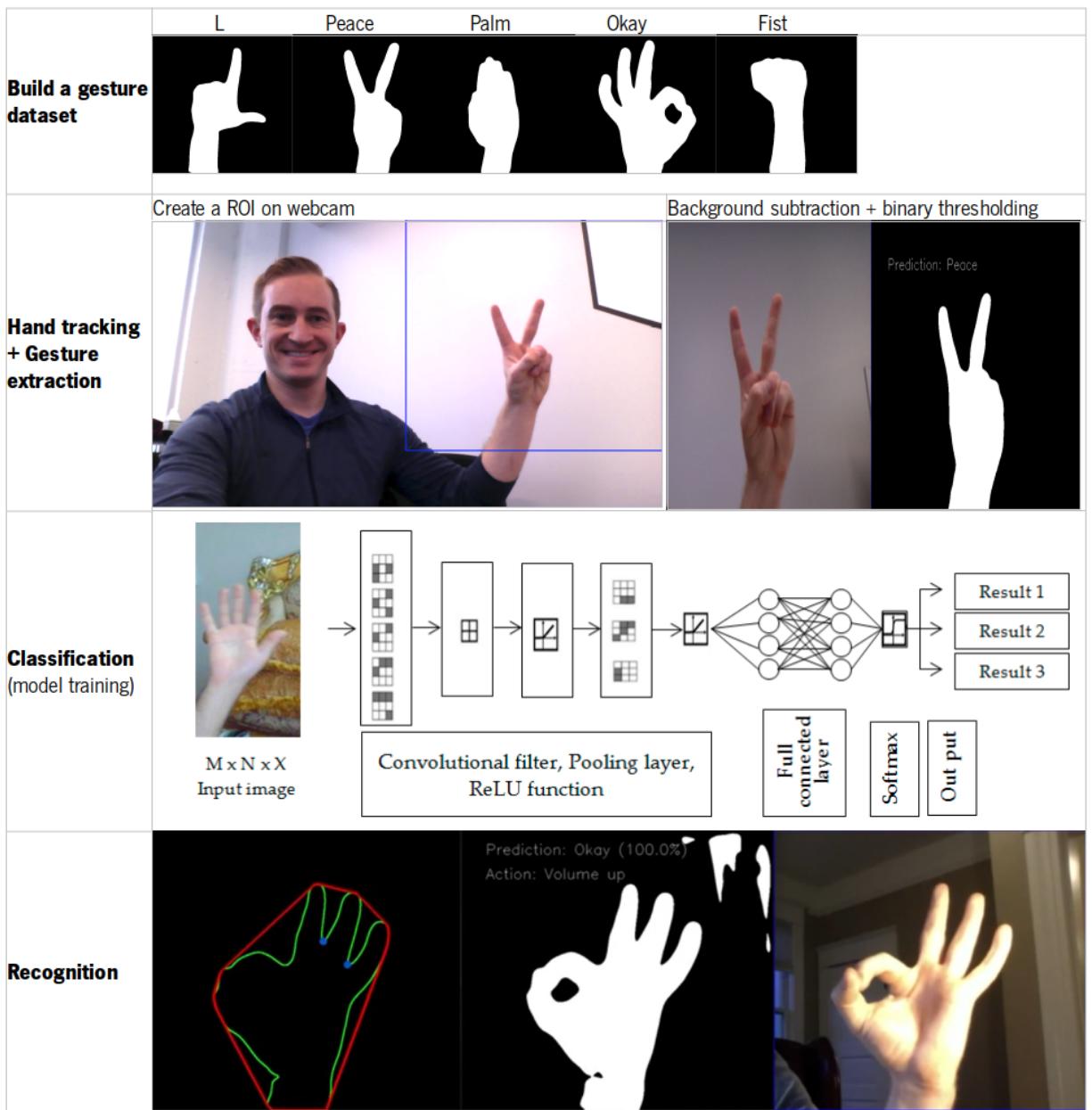


Figure 3.14.: Hand gesture recognition workflow illustrated: example – adapted from [49] and [50]

huge amounts of data [51]. Some examples of RDBMS are Oracle Database, MySQL, IBM DB2, SQLite, PostgreSQL, and MariaDB.

From the users' application standpoint, a RDBMS is a management system for databases, but is useless unless it provides an efficient and easy method to pose questions involving the data stored in the databases. These questions are called queries [51]. A DBMS provides a specialized language – query language – in which queries can be performed. The Structured Query Language (SQL) for relational databases, is now

the standard. Arguably, the most widely used form of concurrent programming is the concurrent execution of database programs (called transactions). Users write programs as if they are to be run by themselves, and the responsibility for running them concurrently is given to the DBMS [51].

In this section an overview is presented about RDBMS foundations: description and storage of data in a DBMS, relational model, levels of abstraction in a DBMS, transaction management, and the structure of a DBMS. Additionally, a brief overview over SQL and a C++ interface is presented.

### **3.11.1. Description and storage of data in a DBMS**

The user of a DBMS is ultimately concerned about the description of various aspects of some real-world enterprise in the form of data. There are two important data models used [51]:

- Data model: collection of high-level data description constructs that hide many low-level storage details. A DBMS allows a user to define the data to be stored in terms of a data model, such as the **relational data model**. It is closer to how the DBMS stores the data.
- Semantic data model: more abstract, high-level data model, closer to human thinking, serving as an useful starting point for the database design. The semantic data model is subsequently translated into a database design in terms of the data model the DBMS actually supports. An example is the Entity-Relationship (ER) model which allows the user to pictorially denote entities and the relationship among them. The semantic data

### **3.11.2. Relational model**

The central data description construct in the relational model is a relation, which can be thought as a set of records [51]. A schema is a description of data in terms of the data model. In the relation model, the schema for a relation specifies its name, the name of each field (or attribute or column), and the type of each field. As an example, student information in a university database may be stored in a relation with the following schema [51]:

```
Students(sid: string, name: string, login: string, age: integer, gpa: real)
```

The preceding schema states that each record in the Students relation has five fields, with field names and types as indicated. An example instance of the student relation appears in Table 3.1 [51].

Each row in the Students relation is a record that describes a student, following the schema of the Students relation. Thus, the schema can be thought as a template for describing a student. This description can be made more precise by specifying integrity constraints, i.e., the conditions that the records in a relation must satisfy [51]. For example, one could specify that every student had a unique sid value, thus making

### 3.11. RDBMS

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Table 3.1.: An instance of the students relation – withdrawn from [51]

sid	name	login	age	gpa
53666	Jones	jones@cs	18	3.4
53688	Smith	smith@ee	18	3.2
53650	Smith	smith@math	19	3.8
53831	Madayan	madayan@music	11	1.8
53932	Guldu	guldu@music	12	2.0

a potential candidate for a primary key, i.e., an unique identifier that univocally identifies each record in a relation. This information cannot be captured by simply adding another field to the Students schema, thus requiring integrity constraints to increase the expressiveness of the constructs of a data model [51].

#### 3.11.3. Levels of abstraction in a DBMS

The data in a DBMS is described at three levels of abstraction, as illustrated in Fig. 3.16, namely [51]:

- External schema: allow data access to be customized (and authorized) at the level of individual users or group of users. Any given database has exactly one conceptual schema and one physical schema because it has just one set of stored relations, but it may have several external schemas, each tailored to a particular group of users. Each external schema consists of a collection of one or more views and relations from the conceptual schema. A view is conceptually a relation, but the records in a view are not stored in the DBMS. Rather, they are computed using a definition for the view, in terms of relations stored in the DBMS.
- Conceptual schema: also known as the logical schema, describes the stored data in terms of the data model of the DBMS. In a RDBMS, the conceptual schema describes all relations that are stored in the database. The conceptual schema may be design using the Entity-Relationship (ER) model.
- Physical schema: translates how the relations described in the conceptual schema are actually stored on secondary storage devices such as disks and tapes. Decisions about the physical schema are based on the understanding of how the data is typically accessed, typically requiring the design to decide about the file organizations used to store the relations and to create auxiliary data structures called indexes to speed up data retrieval operations.

#### 3.11.4. Transaction management

An important task of a DBMS is to schedule concurrent accesses to data in a safe and seamless way to the user. Such accesses are named transactions, i.e., any one execution of a user program in a DBMS,

### 3.11. RDBMS

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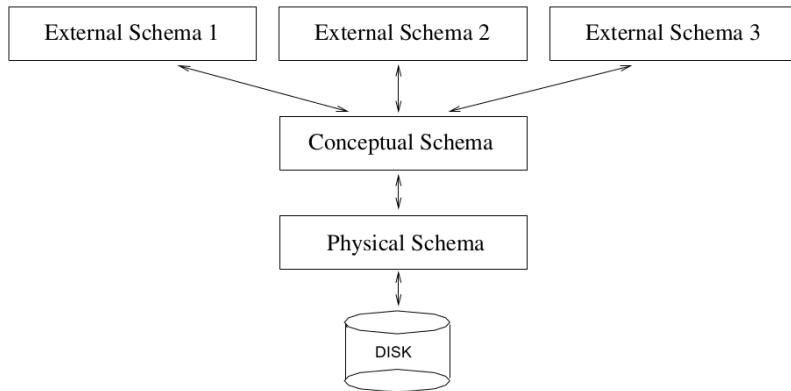


Figure 3.15.: Levels of abstraction in a DBMS (withdrawn from [51])

corresponding to the basic unit of change as seen by the DBMS. Partial transactions are not allowed, and the effect of a group of transactions is equivalent to some serial execution of all transactions [51].

For the concurrent execution of transactions to take place, a locking protocol is enforced by the DBMS, establishing a set of rules to be followed by each transaction, using a lock – a mechanism to control access to database objects. Two kinds of locks are commonly supported by a DBMS: shared locks on an object can be held by two different transactions at the same time, but an exclusive lock on an object ensures that no other transactions hold any lock on this object [51].

Transactions can be interrupted before running to completion for a variety of reasons, e.g., a system crash. A DBMS must ensure that the changes made by such incomplete transactions are removed from the database. To do so, the DBMS maintains a log of all writes to the database, even before the corresponding change is reflected in the database itself, enabling the DBMS to detect and undo the changes if a system crash occurs. This property is called Write-Ahead Log (WAL). To ensure this property, the DBMS must be able to selectively force a page in memory to disk [51].

The log is also used to ensure that the changes made by a successfully completed transaction are not lost due to a system crash, as explained. Bringing the database to a consistent state after a system crash can be a slow process, since the DBMS must ensure that the effects of all transactions that completed prior to the crash are restored, and that the effects of incomplete transactions are undone. The time required to recover from a crash can be reduced by periodically forcing some information to disk; this periodic operation is called a checkpoint [51].

Summarizing, the main takeaways for DBMS support for concurrency control and recovery are [51]:

- Locking: every object that is read or written by a transaction is first locked in shared or exclusive mode, respectively. Placing a lock on an object restricts its availability to other transactions and thereby affects performance.
- WAL: for efficient log maintenance, the DBMS must be able to selectively force a collection of pages

in main memory to disk. OS support for this operation is not always satisfactory.

- Periodic checkpoint: can reduce the time needed to recover from a crash. There is a trade-off between speed and system integrity, as checkpointing too often slows down normal execution.

### 3.11.5. Structure of a RDBMS

Fig. 3.16 shows the structure (with some simplification) of a typical DBMS based on the relational data model.

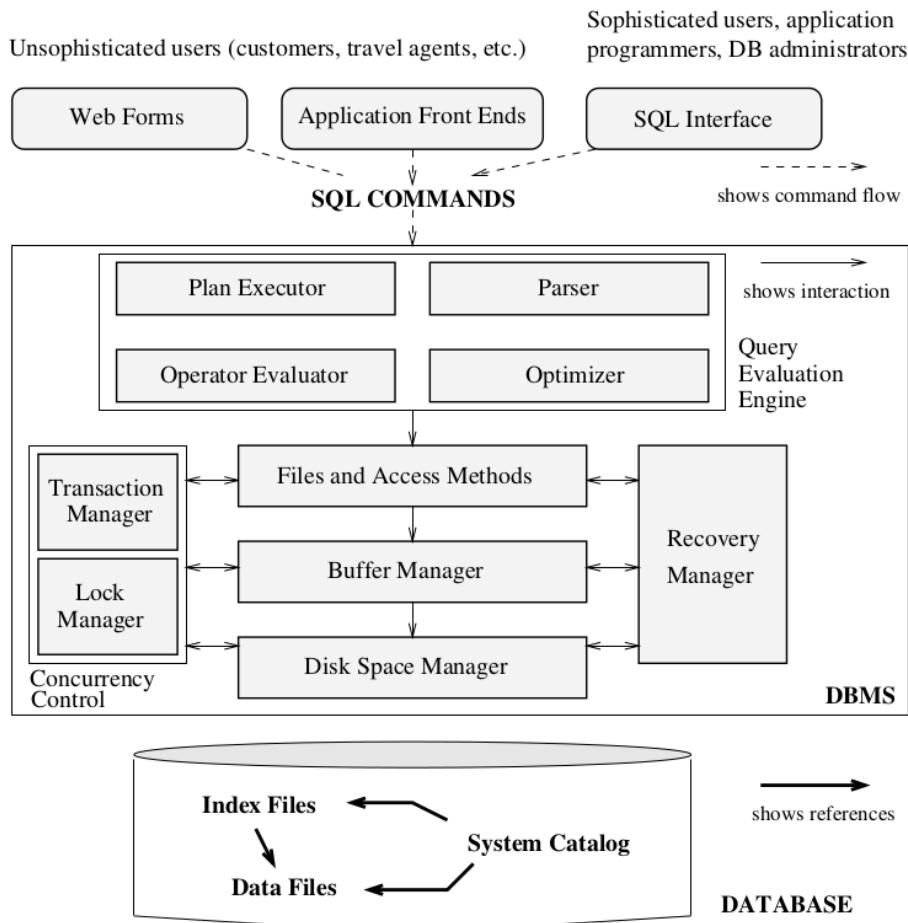


Figure 3.16.: Architecture of a DBMS (withdrawn from [51])

It is comprised of [51]:

- Interfaces: The DBMS accepts SQL commands generated from a variety of user interfaces, produces query evaluation plans, executes these plans against the database, and returns the answers. SQL commands can also be embedded in host-language application programs, e.g., Java or C++ programs,

but here one concentrates only on the core DBMS functionality.

- DBMS: contains:
  - Query Evaluation Engine: When a user issues a query, the parsed query is presented to a query optimizer, which uses information about how the data is stored to produce an efficient execution plan for evaluating the query. An execution plan is a blueprint for evaluating a query, and is usually represented as a tree of relational operators (with annotations that contain additional detailed information about which access methods to use, etc.).
  - Concurrency control: The Transaction Manager ensures that transactions request and release locks according to a suitable locking protocol and schedules the execution transactions. The Lock Manager keeps tracks of request for locks and grants locks on database objects when they become available.
  - Recovery manager: responsible for maintaining a log, and restoring the system to a consistent state after a crash.
  - Disk access manager: The file and access methods layer includes a variety of software for supporting the concept of a file, which, in a DBMS, is a collection of pages or a collection of records. This layer typically supports a heap file, or file of unordered pages, as well as indexes. In addition to keeping track of the pages in a file, this layer organizes the information within a page. The Buffer manager handles pages as a response to read requests. The disk space manager deals with management of space on disk, where the data is stored. Higher layers allocate, deallocate, read, and write pages through (routines provided by) this layer.
- Database: contains the system catalog information consisting of the index files referencing the data files storing the actual data on physical memory.

#### 3.11.6. Database design overview

The database design process can be divided into six steps, namely [51]:

1. Requirements Analysis: the requirements for the database application are elicited and analyzed assessing what data is to be stored in the database, what applications must be built on top of it, and what operations are most frequent and subject to performance requirements. Several methodologies have been proposed for organizing and presenting the information gathered in this step, and some automated tools have been developed to support this process.
2. Conceptual Database Design: the information gathered in the previous step is used to develop a high-level description of the data to be stored in the database, along with the constraints that are known

to hold over this data. This step is often carried out using the ER model, or a similar high-level data model.

3. Logical Database Design: A DBMS must be chosen to implement the database design, and convert the conceptual database design – ER schema – into a database schema in the data model of the chosen DBMS – relational schema.
4. Schema Refinement: Next, the collection of relations in the relation database schema are analyzed to identify potential problems, and to refine it. This is performed by normalizing relations, restructuring them to ensure some desirable properties.
5. Physical Database Design: In this step, the typical expected workloads that the database must support are analyzed and further refine the database design to ensure that it meets desired performance criteria. This may simply involve building indexes on some tables and clustering some tables, or it may involve a substantial redesign of parts of the database schema obtained from the earlier design steps.
6. Security Design: Lastly, the different user groups and different roles played by various users are identified (e.g., the development team for a product, the customer support representatives, the product manager). For each role and user group, the permitted and forbidden parts of the database are identified and the policies are enforced to ensure this. A DBMS provides several mechanisms to assist in this step.

### **3.11.7. Entity-Relationship model**

The Entity-Relationship (ER) data model enables the description of the data involved in a real-world enterprise in terms of entities and their relationships and is widely used to develop an initial database design. The ER model is most relevant to the first three steps of the database design [51] (see Section 3.10.6). In this section are presented the key concepts for the ER model as a database design modeling tool.

#### **Key concepts**

It is important to understand some key concepts for the ER model, namely [51]:

- Entity: is an object in the real world that is distinguishable from other objects, e.g., the Pokemon toy, the toy department, the manager of the toy department, etc.
- Entity set: collection of entities. They do not need to be disjoint, i.e., entities can simultaneously belong to different entity sets. For example, one can define an entity set called **Employees** that contain both the toy and appliance department employee sets. An entity set is represented by a rectangle in the ER model.

- Attributes: an entity is described using a set of attributes. All entities in a given entity set have the same attributes. For example, the **Employees** entity set could use **name**, social security number (**ssn**) and parking lot (**lot**) as attributes. An attribute is represented by an oval in the ER model.
- Domain: for each attribute associated with a set, one must identify a domain of possible values. For example, the domain associated with the attribute **name** of **Employees** might be the set of 20-character strings. The domain information can be listed along the attribute name.
- Key: minimal set of attributes whose values uniquely identify an entity in the set. There could be more than one candidate key: if so, one designate one of them as the primary key. Each attribute in the primary key is underlined in the ER model. The foreign key is(are) the attribute(s) which in a relationship one-to-many is(are) primary key(s) in the entity of the side **one** and integrates the set of attributes of the entity in the side **many** of that relationship
- Relationship: association between two or more entities. For example, one may have the relationship that Joe works in the pharmacy department. A relationship is represented by a straight line in the ER model and may also have descriptive attributes.
- Relationship set: collection of similar relationships.
- Key constraints: restrictions over entities. For example, the restriction that each department has at most one manager, and is denoted by using an arrow from the entity to the relationship.
- Participation constraints: states the participation of an entity set in the relationship. It can be partial or total. For example, if every department is required to have a manager, the participation of the entity set **Departments** in the relationship set **Manages** is said to be total.
- Class hierarchies: classification of entities in an entity set into subclasses, using the relationship is a. For example, a **Car** is a **Vehicle** and a **Truck** is a **Vehicle** too.
- Aggregation: indicates that a relationship set (identified through a dashed box) participates in another relationship set.

The ERDs use a graphical conventional to quickly and clearly depict the entities involved and how they relate to each other. In a ERD entities are represented by rectangles, attributes by ellipses, and the relationships as lines between entities. In the rectangles and ellipses are placed the names of the different entities and attributes. The relationships have cardinalities – 1:1 (one-to-one), 1:M (one-to-many), and M:N (many-to-many) – and may be mandatory or optional – e.g., a vehicle may not have any parking space assigned, but to each parking space is assigned one, and one only, vehicle.

Several notations can be used, namely, crow's foot, UML, Chen, Barker, etc [52]. In crow's foot notation:

- A multiplicity of one and a mandatory relationship is represented by a straight line perpendicular to the relationship line.
- A multiplicity of many is represented by the three-pronged 'crow-foot' symbol.

### 3.11. RDBMS

- An optional relationship is represented by an empty circle.

Fig. 3.17 illustrates an example of an Entities-Relationships diagram (ERD) using crow's foot notation for a company. There are three entity sets – Customers, Orders, and Shipments. Within each of these are the attributes, with the primary key being underlined. Additionally it also indicates the foreign key that resolves the one-to-many relationship. Thus, a customer can place 0 or many orders, which, in turn, can have 0 or many shipment methods.

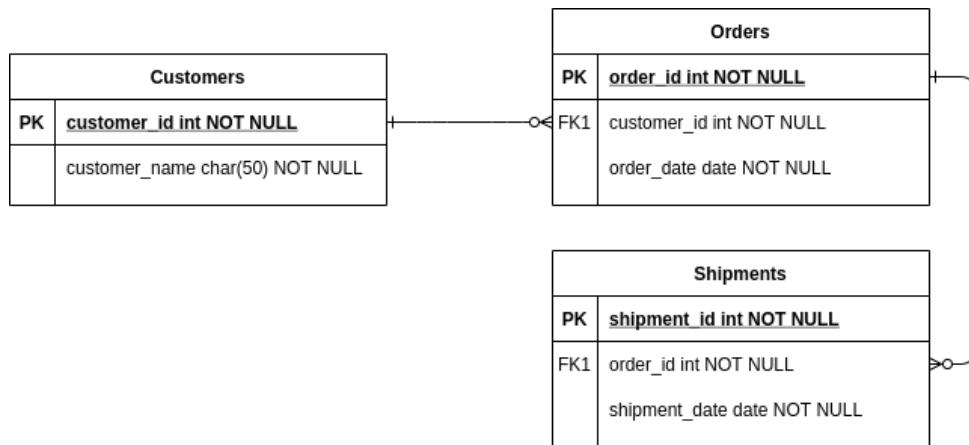


Figure 3.17.: Example of an ERD

#### 3.11.8. Choice of the RDBMS

In this section are presented the most relevant RDBMSs, its advantages and disadvantages, and the best use case, which can help the database designer to appropriately choose a RDBMS. Special focus is given to the free systems, namely MySQL and SQLite, due to budget restrictions.

The most relevant RDBMSs are [53]:

- **Oracle Database:** Oracle has provided high-quality database solutions since the 1970s. The most recent version of Oracle Database was designed to integrate with cloud-based systems, and it allows you to manage massive databases with billions of records.
  - Advantages: the most advanced technology and a wide range of solutions.
  - Disadvantages: an expensive solution and system upgrades might be required — many businesses have to upgrade their hardware before using Oracle solutions.
  - Best use case: if you're a large organization that needs to manage a massive amount of data, Oracle could be the ideal choice.

- **Microsoft SQL Server:** it is a database engine that is compatible with, both, on-site and cloud-based servers, and supports Windows and Linux OSes.
  - Advantages: it is mobile: this database engine allows you to access dashboard graphics and visuals via mobile devices. It integrates with Microsoft products. It is fast and stable.
  - Disadvantages: an expensive solution and requires a lot of HW resources.
  - Best use case: if you're an enterprise-level corporation that relies heavily on Microsoft products, the speed, agility, and reliability of Microsoft SQL Server could be an excellent choice.
- **MySQL:** MySQL is a free, open-source RDBMS solution that Oracle owns and manages. Even though it's freeware, MySQL benefits from frequent security and features updates. Large enterprises can upgrade to paid versions of MySQL to benefit from additional features and user support.
  - Advantages: it is open-source, free of charge (freeware) and highly compatible with many other database systems.
  - Disadvantages: lacking features common to other RDBMSs: because MySQL prioritizes speed and agility over features, some of the standard features found in other solutions may be missing, e.g., the ability to create incremental backups. Challenges getting quality support: The free version of MySQL does not come with on-demand support. However, MySQL does have an active volunteer community, useful forums, and a lot of documentation.
  - Best use case: MySQL is a particularly valuable RDBMS solution for businesses that need a solution with enterprise-level capabilities, but are operating under strict budget constraints. It is an extremely powerful and reliable modern RDBMS with a free tier.
- **SQLite [54]**: it is a C-language library that implements a small, fast, self-contained SQL database engine — an embedded DB — which means the DB engine runs as a part of the app.
  - Advantages: it is open-source, free of charge (freeware), a server-less and file-based database, and self-contained. It has a small storage footprint (the SQLite library is 250 KB in size, while the MySQL server is about 600 MB). It directly stores information into a single file, making it easy to copy, and no configuring is required.
  - Disadvantages: it lacks user management (not suitable for multiple users) and security features (the database can be accessed by anyone). It is not easily scalable and cannot be customized.
  - Best use case: SQLite is best suited for developing small standalone apps or smaller projects which do not require much scalability.

## Summary

Oracle Database and Microsoft SQL server are proprietary solutions, specially suited for large organizations, and can be very expensive. On the other hand, MySQL and SQLite are open-source and freeware solutions, which can be used to quickly test and iterate the DB design before moving into production. However, SQLite does not have user authentication or security features, and it is not easily scalable. Thus, MySQL arises as the best option, suited for distributed architectures, as a trade-off between cost, ease of use, scalability, security and user management.

### 3.11.9. SQL

Ideally, a database language allows the creation of a database and table structures, the execution of basic data management tasks (add, delete, and modify), and the execution of complex queries designed to transform the raw data into useful information. Moreover, it must provide a clear and easy syntax, it must be portable and conform to some basic standard. SQL complies well to these requirements [55].

SQL functions fit into two broad categories [55]:

1. It is a Data Definition Language (DDL): SQL includes commands to create database objects such as tables, indexes, and views, as well as commands to define access rights to those database objects (see Fig. 3.18).
2. It is a Data Manipulation Language (DML): SQL includes commands to insert, update, delete, and retrieve data within the database tables (see Fig. 3.19).

COMMAND OR OPTION	DESCRIPTION
CREATE SCHEMA AUTHORIZATION	Creates a database schema
CREATE TABLE	Creates a new table in the user's database schema
NOT NULL	Ensures that a column will not have null values
UNIQUE	Ensures that a column will not have duplicate values
PRIMARY KEY	Defines a primary key for a table
FOREIGN KEY	Defines a foreign key for a table
DEFAULT	Defines a default value for a column (when no value is given)
CHECK	Validates data in an attribute
CREATE INDEX	Creates an index for a table
CREATE VIEW	Creates a dynamic subset of rows/columns from one or more tables
ALTER TABLE	Modifies a table's definition (adds, modifies, or deletes attributes or constraints)
CREATE TABLE AS	Creates a new table based on a query in the user's database schema
DROP TABLE	Permanently deletes a table (and its data)
DROP INDEX	Permanently deletes an index
DROP VIEW	Permanently deletes a view

Figure 3.18.: SQL data definition commands – withdrawn from [55]

### 3.11. RDBMS

COMMAND OR OPTION	DESCRIPTION
INSERT	Inserts row(s) into a table
SELECT	Selects attributes from rows in one or more tables or views
WHERE	Restricts the selection of rows based on a conditional expression
GROUP BY	Groups the selected rows based on one or more attributes
HAVING	Restricts the selection of grouped rows based on a condition
ORDER BY	Orders the selected rows based on one or more attributes
UPDATE	Modifies an attribute's values in one or more table's rows
DELETE	Deletes one or more rows from a table
COMMIT	Permanently saves data changes
ROLLBACK	Restores data to their original values
<b>COMPARISON OPERATORS</b>	
=, <, >, <=, >=, <>	Used in conditional expressions
<b>LOGICAL OPERATORS</b>	
AND/OR/NOT	Used in conditional expressions
<b>SPECIAL OPERATORS</b>	
BETWEEN	Checks whether an attribute value is within a range
IS NULL	Checks whether an attribute value is null
LIKE	Checks whether an attribute value matches a given string pattern
IN	Checks whether an attribute value matches any value within a value list
EXISTS	Checks whether a subquery returns any rows
DISTINCT	Limits values to unique values
<b>AGGREGATE FUNCTIONS</b>	
COUNT	Used with SELECT to return mathematical summaries on columns
MIN	Returns the number of rows with non-null values for a given column
MAX	Returns the minimum attribute value found in a given column
SUM	Returns the maximum attribute value found in a given column
Avg	Returns the sum of all values for a given column
	Returns the average of all values for a given column

Figure 3.19.: SQL data manipulation commands – withdrawn from [55]

#### 3.11.10. MySQL Interfaces

MySQL works under the client–server paradigm. It has several client interfaces that can interact with the server, through connectors and APIs, i.e., the drivers and libraries that one can use to connect applications in different programming languages to MySQL database servers. The application and database server can be on the same machine, or communicate across the network [56]. The following interfaces are available: Java, Python, JavaScript, C++, C, C#, PHP, OBDC, NBD Cluster, MySQL Shell, and X DevAPI.

##### C++ connector

From the list of available interfaces, the most well suited to interface the RDBMS are the C API and the C++ connector, as they are the most well known programming languages by the authors and the better ones in terms of performance. However, the C++ connector was chosen because it offers the following benefits over the MySQL C API provided by the MySQL client library [57]:

- Convenience of pure C++.

- Support for the object-oriented programming paradigm.
- Support for these application programming interfaces: X DevAPI, X DevAPI for C, Legacy JDBC 4.0-based API.
- Reduced development time.
- Licensed under the GPL with the FLOSS License Exception.
- Available under a commercial license upon request.

Listing 3.2 illustrates an application to connect to the MySQL server using the C++ connector [58]. Starting from the line 24 the application tries to connect to the MySQL database and execute a query and, if any exception is thrown, it is caught by the catch block at line 50. Looking into more detail at the try block, the pointers to objects sql::Driver, sql::Connection, sql::Statement and sql::ResultSet are instantiated for later usage. A connection to the MySQL server is established at line 32 – via driver – using the TCP/IP protocol at localhost address and port 3306 (the default for MySQL server), and passing the username and the password. Next, the specific database – test – is defined for access.

Then, at line 36, a statement is created to execute a query (line 37), with the result being captured by res (the sql::ResultSet object). A loop is performed to retrieve all data, and lines 41 and 44 illustrate different methods to access the column data – by column name or by numeric offset, respectively. Finally, all objects dynamically allocated are released (lines 46-48) and, if everything works well, the application returns successfully (line 61).

```

1 /* Standard C++ includes */
2 #include <stdlib.h>
3 #include <iostream>
4 /*
5  * Include directly the different
6  * headers from cppconn/ and mysql_driver.h + mysql_util.h
7  * (and mysql_connection.h). This will reduce your build time!
8  */
9 #include "mysql_connection.h"

11 #include <cppconn/driver.h>
12 #include <cppconn/exception.h>
13 #include <cppconn/resultset.h>
14 #include <cppconn/statement.h>

16 using std::cout;
17 using std::endl;

19 int main(void)
{
    cout << endl;

```

### 3.11. RDBMS

---

```
cout << "Running 'SELECT 'Hello World!' AS _message ...' << endl;

try {
    25     sql::Driver *driver;
    26     sql::Connection *con;
    27     sql::Statement *stmt;
    28     sql::ResultSet *res;

    /* Create a connection */
    29     driver = get_driver_instance();
    30     con = driver->connect("tcp://127.0.0.1:3306", "root", "root");
    31     /* Connect to the MySQL test database */
    32     con->setSchema("test");

    stmt = con->createStatement();
    34     res = stmt->executeQuery("SELECT 'Hello World!' AS _message");
    35     while (res->next()) {
        36         cout << "\t... MySQL replies: ";
        37         /* Access column data by alias or column name */
        38         cout << res->getString("_message") << endl;
        39         cout << "\t... MySQL says it again: ";
        40         /* Access column data by numeric offset, 1 is the first column */
        41         cout << res->getString(1) << endl;
        42     }
        43     delete res;
        44     delete stmt;
        45     delete con;

} catch (sql::SQLException &e) {
    46     cout << "# ERR: SQLException in " << __FILE__;
    47     cout << "(" << __FUNCTION__ << ") on line " <<
    48         << __LINE__ << endl;
    49     cout << "# ERR: " << e.what();
    50     cout << " (MySQL error code: " << e.getErrorCode();
    51     cout << ", SQLState: " << e.getSQLState() << " )" << endl;
    52 }

    53     cout << endl;

    54     return EXIT_SUCCESS;
}
```

Listing 3.2: MySQL Connector example — adapted from [58]

## **3.12. Motion detection**

aaaaaaaaa

## **3.13. Camera recording and codecs**

aaaaaaaaa

## **3.14. Image filtering**

aaaaaaaaa

## **3.15. GIF generation**

aaaaaaaaa

## **3.16. Social media sharing APIs**

aaaaaaaaa

## **3.17. UI framework**

For the development of the Remote Client, it is necessary to use a UI Framework, in order to develop a UI to the users, making it more user friendly and interactive. There are several frameworks that can be used, such as:

- Qt;
- Sciter;
- Noesis GUI;
- wxWidgets;
- GTK+;
- and so on.

---

### **3.18. File transfer protocols**

For this project, Qt was chosen due to the following reasons:

- First of all, it is possible to ‘develop graphical user interfaces and cross-platform applications, both desktop and embedded’ [59]. The framework operates on different types of software and hardware. Also, it has an easy-to-read code and it is possible to make an attractive UI.
- Secondly, this framework is cost-friendly, not only because it has a free license, but also because its software development takes less time to develop due to its highly productive features.
- Lastly, it is implemented in C++, which means that it is possible to use many libraries. The wide choice of modules allow the project to have rich functionality and as a result, the software will have a Graphical User Interface (GUI) similar to a native one.

## **3.18. File transfer protocols**

aaaaaaaa

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# **Appendices**

## **A. Project Planning – Gantt diagram**

In Fig. A.1 is illustrated the Gantt chart for the project, containing the tasks' descriptions.

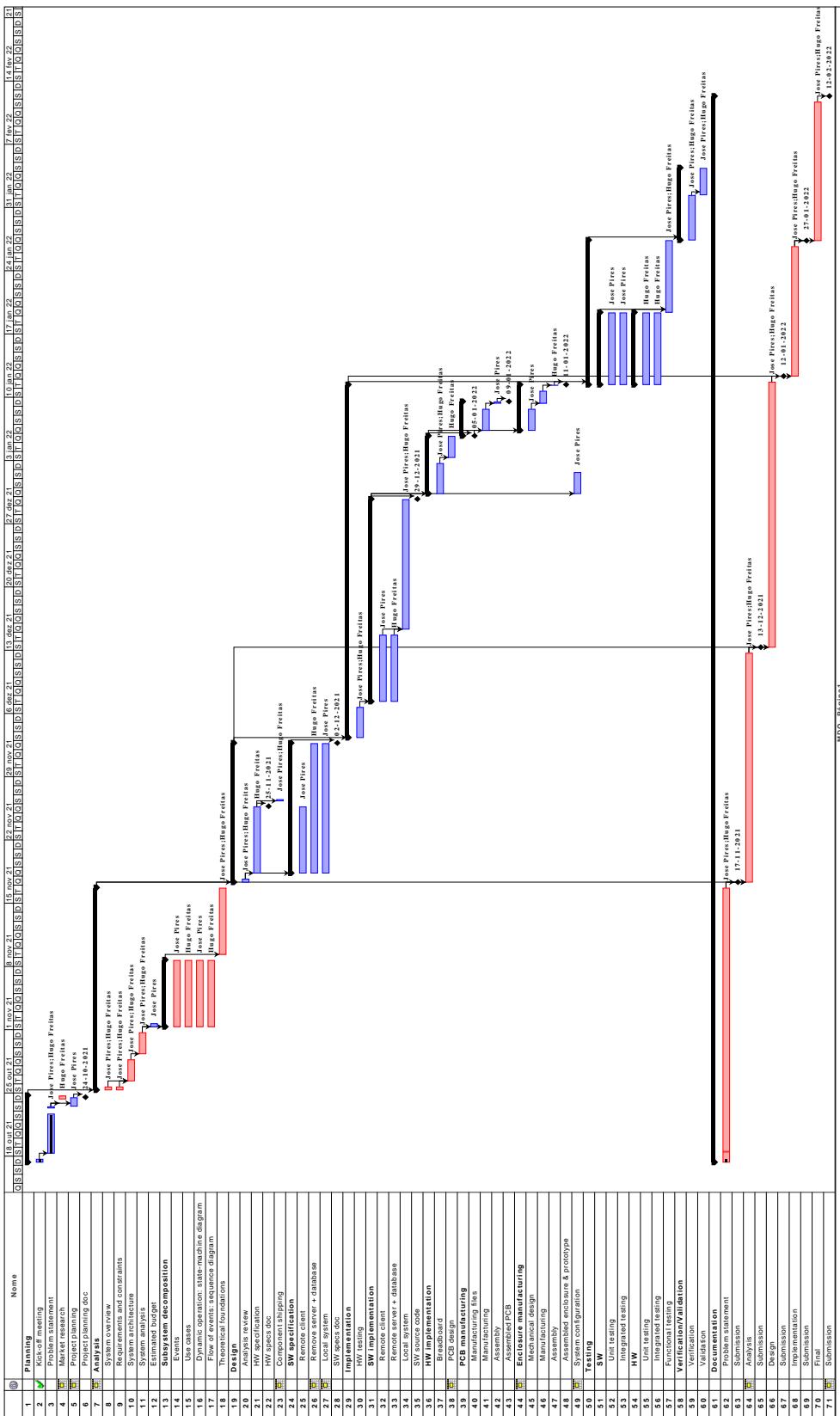


Figure A.1.: Project planning – Gantt diagram