

Experiment No.5

Aim:

Implementation of Bellman ford , Johnson's algorithm for sparse graphs.

Problem Statement:

Analyzing and comparing the Bellman-Ford (single-source shortest path) algorithm and Johnson's (All-pairs shortest path) algorithm and observing their time-complexity behaviour and noting the key differences between the two algorithms over a sample of sparse graphs.

Bellman-Ford Algorithm

Objective:

- Define and Initialize the sparse graph using Python3 programming language.
- Select a vertex as a source vertex from the graph defined and pass the source vertex to Bellman-Ford algorithm as an argument for processing.
- Define and execute relaxation process to find shortest paths, and report if any negative weight cycle is found during this process.

Methodology:

1. Initialize the distances from source to all vertices as infinite and distance to source itself as 0. Create an array `dist[]` of size $|V|$ with all values as infinite, except `dist[src]` where `src` is source vertex.
2. Calculate the shortest distances. Do following $|V|-1$ times where $|V|$ is the number of vertices in given graph.
 - a) Do following for each edge (u,v) :
If $\text{dist}[v] > \text{dist}[u] + \text{weight of edge } (u,v)$, then:
 $\text{dist}[v] = \text{dist}[u] + \text{weight of edge } (u,v)$
3. Report if there is a negative weight cycle in graph.
 - a) Do following for each edge (u,v) :
If $\text{dist}[v] > \text{dist}[u] + \text{weight of edge } (u,v)$, then:
"Graph contains negative weight cycle"
4. The idea of step 3 is, step 2 guarantees shortest distances, if the graph doesn't contain negative weight cycle (as there can be maximum $|V| - 1$ edges in any simple path, that is why the outer loop runs $|V| - 1$ times). If we iterate through all edges one more time and still get a shorter path for any vertex, then there must be a negative weight cycle present.
5. Time Complexity : $O(V.E)$

Implementation:

1. To instantiate and initialize a Graph we define a *class Graph* with following member functions.
2. `__init__(self, vertices)` function to instantiate the object of class Graph with number of vertices.
3. `addEdge(self, u, v, w)` function is defined to add edges to the Graph Object, where u,v denote - edge u to v and w denotes weight of the edge.
4. `printArr(self, dist, src)` function is used to result on the terminal where `dist` is an array of integers representing shortest distance of vertices from source vertex `src`.
5. Finally, `BellmanFord(self, src)` function for calling our algorithm.

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Johnson's Algorithm

Objective:

- Define and Initialize the sparse graph using Python3 programming language.
- Define a subroutine that implements Bellman-Ford Algorithm using Python3 programming language.
- Define another subroutine that implements Dijkstra's Algorithm using Python3 programming language.
- Define a common subroutine for Johnson's Algorithm that uses above two subroutines, and implements re-weighting process to transform the graph into non-negative weight graph.

Methodology:

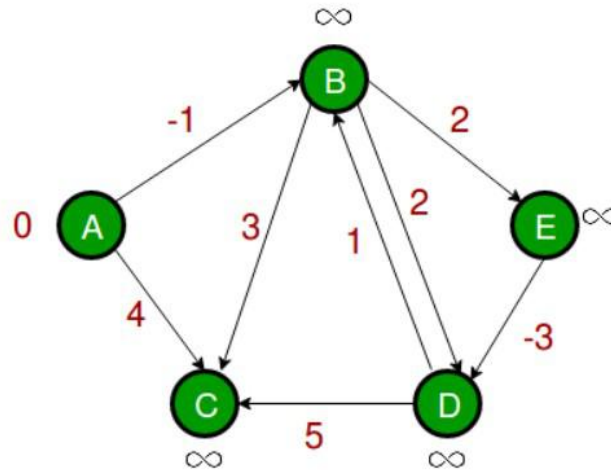
1. Let the given graph be G . Add a new vertex S to the graph, add edges from new vertex to all vertices of G . Let the modified graph be G' .
2. Run Bellman-Ford algorithm on G' with S as source. Let the distances calculated by Bellman-Ford be $dist[0], dist[1], \dots, dist[V-1]$. If we find a negative weight cycle, then return. Note that the negative weight cycle cannot be created by new vertex S as there is no incoming edge to S . All edges are outgoing edges from S .
3. Re-weight the edges of original graph. For each edge (u, v) , assign the new weight as :
“(original weight) + $dist[u] - dist[v]$ ”.
4. Remove the added vertex S and run Dijkstra's algorithm for every vertex.
5. Time Complexity :

Bellman Ford is $O(V.E)$ + Dijkstra is $O(V.LogV)$ for every vertex
so overall, $O(V^2.log V + V.E)$

Implementation:

1. We define *BellmanFord(edges, graph, num_vertices)* function to execute Bellman-ford algorithm on the given *graph{edges,num_vertices}*.
2. Define *Dijkstra(graph, modifiedGraph, src)* and *minDistance(dist, visited)* function to execute Dijkstra's single source shortest path algorithm for every vertex in the *modifiedGraph* for source vertex *src*.
3. Define *JohnsonAlgorithm(graph)* function which will :
 - a. call *BellmanFord(edges, graph, num_vertices)* to get distance of all vertices from source *src*.
 - b. Define the *modifiedGraph* by using re-weighting process to convert all negative weights into positive weights.
 - c. Call *Dijkstra(graph, modifiedGraph, src)* function on previously defined *modifiedGraph*.

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Results:

Bellman-Ford Algorithm

```
(base) C:\Users\Vishal Ramane\OneDrive\College\AAC Lab\Code\Expt-5>python -u "c:\Users\Vishal Ramane\OneDrive\College\AAC Lab\Code\Expt-5\BellmanFord.py"

For source vertex:4

Vertex      Distance from Source
0           inf
1           -2
2           1
3           -3
4           0

(base) C:\Users\Vishal Ramane\OneDrive\College\AAC Lab\Code\Expt-5>
```

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Johnson's Algorithm

```
(base) C:\Users\Vishal Ramane\OneDrive\College\AAC Lab\Code\Expt-5>python -u "c:\Users\Vishal Ramane\OneDrive\College\AAC Lab\Code\Expt-5\JohnsonAlgo.py"  
Modified Graph: [[0, 1, 4, 0, 0], [0, 0, 1, 3, 0], [0, 0, 0, 0, 0], [0, 0, 2, 0, 0], [0, 0, 0, 0, 0]]
```

Shortest Distance with vertex 0 as the source:

```
Vertex 0: 0  
Vertex 1: 1  
Vertex 2: 2  
Vertex 3: 1  
Vertex 4: 1
```

Shortest Distance with vertex 1 as the source:

```
Vertex 0: inf  
Vertex 1: 0  
Vertex 2: 1  
Vertex 3: 0  
Vertex 4: 0
```

Shortest Distance with vertex 2 as the source:

```
Vertex 0: inf  
Vertex 1: inf  
Vertex 2: 0  
Vertex 3: inf  
Vertex 4: inf
```

Shortest Distance with vertex 3 as the source:

```
Vertex 0: inf  
Vertex 1: 0  
Vertex 2: 1  
Vertex 3: 0  
Vertex 4: 0
```

Shortest Distance with vertex 4 as the source:

```
Vertex 0: inf  
Vertex 1: 0  
Vertex 2: 1  
Vertex 3: 0  
Vertex 3: 0  
Vertex 4: 0
```

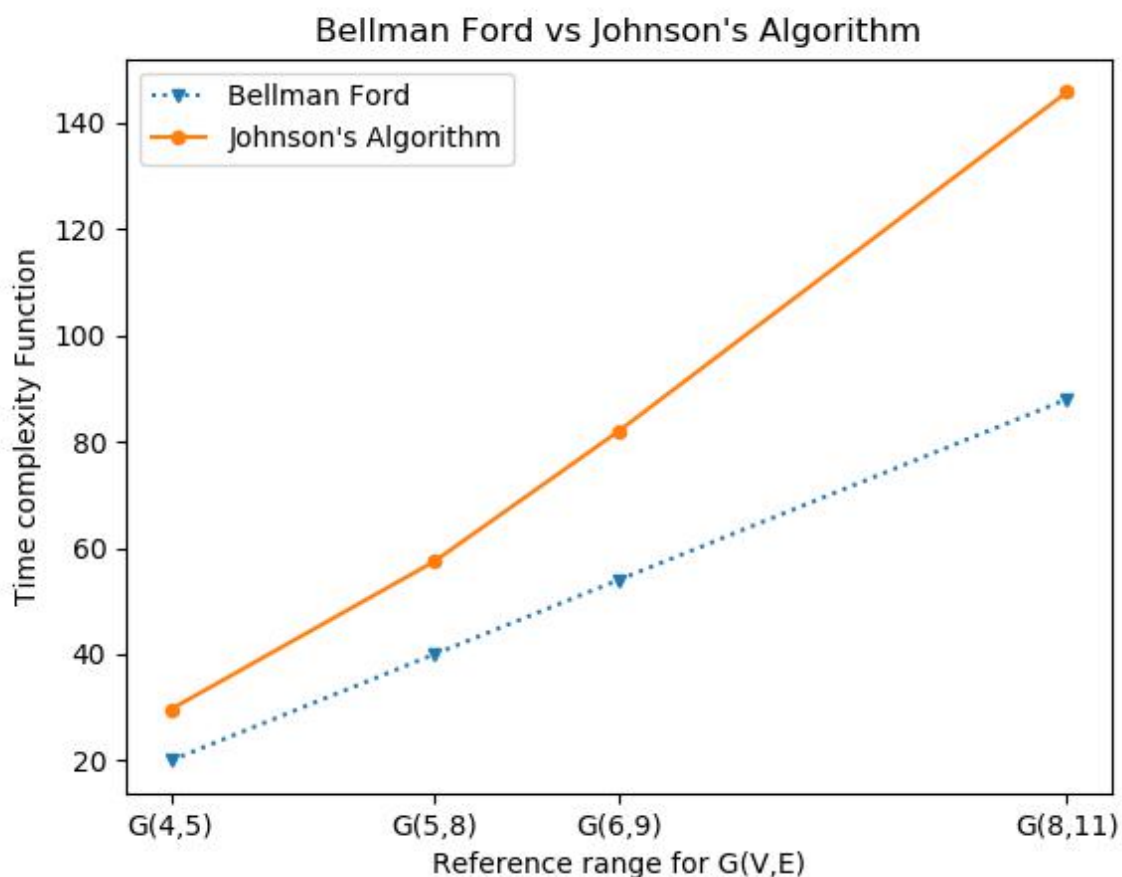
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Observations:

- Differences:

Bellman-Ford Algorithm	Johnson's Algorithm
Single-source shortest path algorithm.	All -pairs shortest path algorithm.
Based on Dynamic Programming design technique.	Consists of two subroutines Bellman-Ford Algorithm (based on Dy.Programming) and Dijkstra's Algorithm (based on Greedy Algorithm).
Uses Relaxation Process.	Uses Re-weighting process to convert negative weights to positive weights.
Time Complexity : $O(V.E)$	Time Complexity : $O(V^2.\log V + V.E)$

- Time Complexity :



Conclusion:

Thus Bellman Ford Single source shortest path algorithm and Johnson's All-pairs shortest path algorithm are successfully implemented, studied and analyzed for sparse graphs.