

Decentralized Privacy-Preserving Proximity Tracing

Based on the document

<https://github.com/DP-3T/documents>

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Goal

How to identify people who have been in contact with an infected person?

Academics issued a proposal of a design of an app aiming at **proximity tracing** – determining who has been in close physical proximity to an infected person.

Main purpose: quick notification of contact people

Ideally the app should ensure **security** and **privacy** for users.

This mini-lecture: overview of the design proposal

Please visit <https://github.com/DP-3T/documents> for details

High level idea

- the app generates a pseudorandom ID (using a PRG) that represents a user
- This ID is continually broadcasted via BlueTooth
- the app records all IDs observed in the proximity
- if ID* is diagnosed for the virus, the user ID* uploads all its recorded IDs to a central server
- the app once in a while downloads data from the servers and locally checks if any of the recorded IDs has been diagnosed for the virus

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Privacy considerations:

- the central server only stores IDs of infected people (no proximity data is uploaded)
- Nobody can track non-infected users seeing only their IDs
- No entity should be able to use the known data for malicious purposes (e.g.,)

More details I

1. Secret Key Generation:

For the day t , the smartphone generates a random key S_t :

$$S_t \leftarrow \text{KeyGen}(t)$$

using, for example, a PRG (see lecture 2).

For the next days, the secret keys S_{t+i} are generated using a **cryptographic hash function** \mathcal{H} :

$$S_{t+i} = \mathcal{H}(S_{t+i-1})$$

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2. Ephemeral ID generation:

Each user changes his **EphID** n times per day (tunable parameter).

At the beginning of day t , the app generates n ephemeral IDs **EphID** as

$$\text{EphID}_1 || \dots || \text{EphID}_n = \text{PRG}(\mathcal{H}(S_t, \text{'broadcast key'})),$$

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Each $|\text{EphID}_i|$ is 16 bytes (Bluetooth Low Energy beacons payload)

More details II

3. Local storage of observed data: "Seeing" any EphID in proximity, each smartphone stores

- EphID
- duration of contact
- proximity
- time indication (April, 13)

This whole data occupies approx. 32 bytes.

More details II

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4. If a user is infected: on the day t :

- The user sends S_t to the central server
- The sever computes all $EphID_i$ from the day t up until today (knowing S_t enable to do that)
- The server additionally checks if any of the $EphID_i$ has been recorded *before* S_t has been published
- The infected user chooses a completely new secret key

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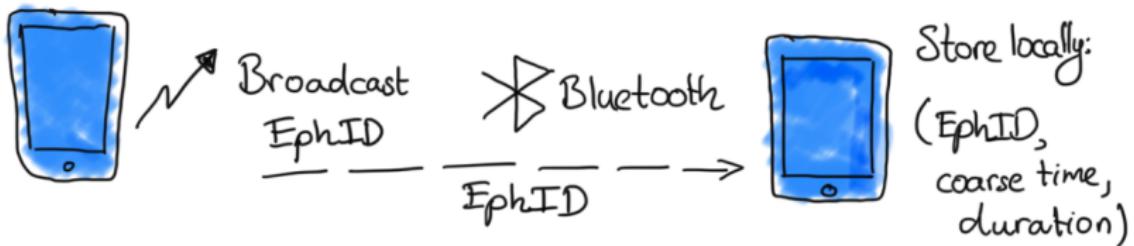
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5. Notification of risk:

- The server broadcasts secret keys (S_t, t) of infected users
- Each user loads (S_t, t) and recomputes the corresponding EphID's
- And checks these EphID' against the locally stored ones.

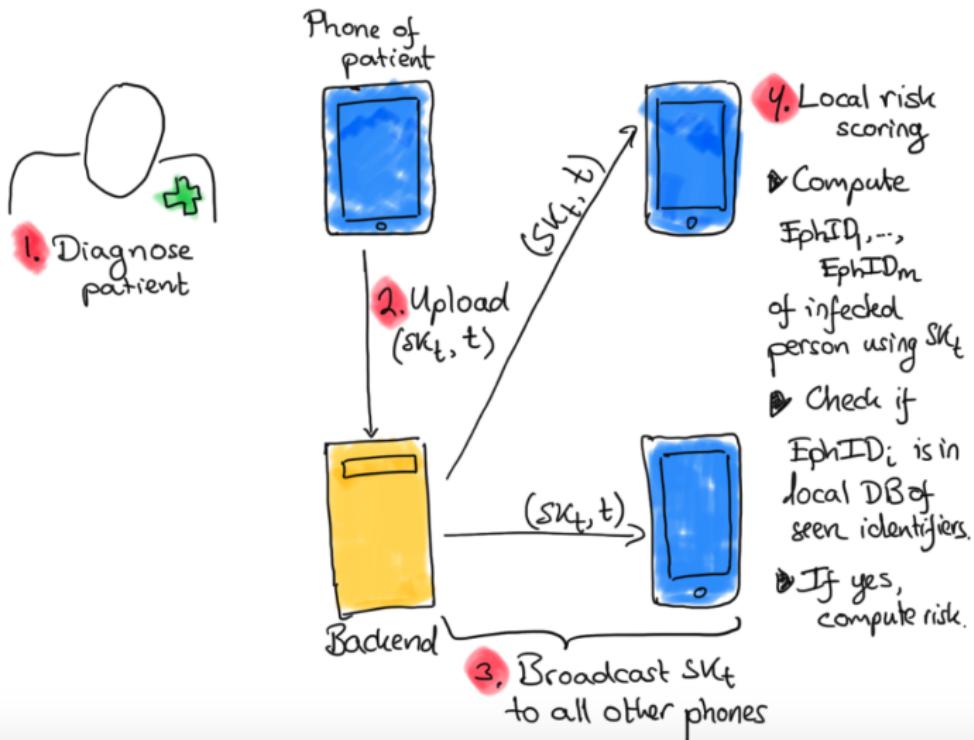
In picture I



The picture is taken from DP3T White Paper.pdf

<https://github.com/DP-3T/documents/blob/master/DP3T%20White%20Paper.pdf>

In picture II



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