

---

## TUTORIAL 3

---

### 1 Alternative FFT algorithm

Let  $P$  be a polynomial of degree at most  $2^k - 1$ , and write  $P = P_h X^{2^{k-1}} + P_l$ . Let  $\omega$  be a primitive  $2^k$ -th root of 1.

1. Prove that  $P(\omega^{2i}) = P_h(\omega^{2i}) + P_l(\omega^{2i})$  and  $P(\omega^{2i+1}) = -P_h(\omega^{2i+1}) + P_l(\omega^{2i+1})$
2. Deduce an alternative FFT algorithm. You will need to introduce the polynomial

$$Q(X) = P_l(\omega X) - P_h(\omega X).$$

### 2 The “binary splitting” method computation of $n!$

We want to compute  $n!$ . We denote by  $M(k)$  the cost (in terms of elementary operations) of the multiplication of two  $k$ -bit numbers, and we assume  $2M(k/2) \leq M(k)$  (we recall some typical values:  $M(k) = O(k^2)$  with naive multiplication,  $O(k^{\log(3)/\log(2)})$  with Karatsuba multiplication and  $O(k \log k \log \log k)$  with the FFT-in finite ring variant of the Schönhage & Strassen algorithm). Use the fact that  $\log n! \sim n \log n$ .

1. What is the cost of multiplying  $O(n)$ -digit integer by a  $O(1)$ -digit integer by the naive algorithm. Argue that it is essentially optimal.
2. We first consider the simplest approach:  $x_1 = 1$ ,  $x_2 = 2x_1$ ,  $x_3 = 3x_2$ ,  $\dots$ ,  $x_n = nx_{n-1}$ . Show that the cost of this approach is  $O(n^2(\log n)^2)$ .
3. We define

$$p(a, b) = (a+1)(a+2)\cdots(b-1)b = \frac{b!}{a!}.$$

Suggest a recursive method to compute  $n!$  with cost  $O(\log n M(n \log n))$ . Conclude on the complexity of your method under different values of  $M(k)$ .

### 3 Computing square root of $F \bmod X^n$

In class, you have seen how to compute  $\sqrt{F} \bmod X^n$  for a polynomial  $F$  of degree  $< n$  in time  $3M(n)$  using Newton’s iteration. In this exercise, we develop an algorithm to compute the square root of  $F \bmod X^n$ . For simplicity, assume  $n = 2^k$ .

1. Consider the polynomial

$$\Phi(y) = y^2 - F.$$

Give an algorithm to compute  $\sqrt{F} \bmod X^n$  and argue about its complexity (expressing this complexity as  $c \cdot M(n) + \Theta(n)$ , we are mostly interested to determine  $c$ ).

2. Can you improve your algorithm by considering

$$\Phi(y) = \frac{1}{y^2} - f ?$$

Give an algorithm and argue about its complexity.