Notes for simulation of traffic flow on an arbitrary network of one-way single-lane roads with traffix lights at intersections.

Charles S. Peskin April 17, 2017 C= can index, nc = # f cars

i = i Wersechm inclex, ni = # fintersection

b = block mdex, nb = # of blocks

i1(b), i2(b) = indices of intersections connected by block b, ordered by the direction of traffic flow. (All blocks are one-way.)

nbin(i) = # f blocks eweriz intersection i

 $bin(i,j) = index f j^{th} block enterior intersection i$ <math>j=1...nbin(i)

nbout (i) = # fblocks leaving intersecon i

bout (i, i) = index of jh block leaving ontersechn i
j=1...nbout(i)

Nose that noin, bin can be derived from 12, and that nout, bout can be derived from i1, as follows:

for l=1:ni nbin(l)=Sum(i2==i) nbout(i)=Sum(i1==i)end nbinmax=max(nbin) nboutmax=max(nbout) bin=Zeros(ni,nbinmax) bout=2eros(ni,nboutmax) for i=1:ni bm(i,1:nbin(i))=ford(i2==i) bout(i,1:nbout(i))=ford(i1==i)end

As a check, it should be the care that

sum (nbin) = sum (nbout) = nb

Traffiz lights

At any given time, the traffiz light at Musechin is green for exactly one of the blocks that enter that intersection and red for all of the others entery that intersection

Let jgreenli) be an integer designating which block has the green light, where

 $1 \leq jgreen(i) \leq nbin(i)$

let S(b) be the state of the light at the end of block b, where S=0 denotes red and S=1 denotes green.

Grenthe array jgreen, s can be set

 $S = 2e_{10}s(1, nb)$ $f_{10} i = 1: ni$ b = bin(i, jgreen(i))s(b) = 1 Geometric Mormation court the network of neads

xi/i), yi/i) = coordinates of Mersechin (

L(b) = length of block b

(UX(b), uy(b)) = unit vector alm block b m direction of traffic flow

Gren Xi, yi, we can find L, ux, uy as follows

ux = xi(i2) - xi(i1) uy = yi(i2) - yi(i1)

L = Syrt(ux.12 + uy.12)

ux = ux./L uy = uy./L

Cars on blocks

Let p(c) be the position of can com whatever block it happens to be on, measured as distance from the Stant of the block. If can c is on block b, then

$$0 \le p(c) < L(b)$$

and the coordinates of can a are given by

$$\chi(c) = \chi_i(i1(b)) + p(c) * u\chi(b)$$

$$y(c) = y_i(i1(b)) + p(c) * uy(b)$$

To access all of the cars on a block in order of decreasing p, we use the following linked-list data structure

first can (b) = index of first can an block b next can (c) = index of can immediately behind can c on the same block last can (b) = index of last can on block b

In all cases, an entry f D means that there is no such can. Thus next can (last can (b) = 0, and if block b is empty then first can (b) = 1 ast (ast can (b) = 0.

Entry of cars and choice of their destinations

Cars enter the roadway (from parking garages or parking spaces) at random times and locations. Let R be the rate at which This occurs. Then R has units of 1/(time. length). Choose the time step dt small enough that R*Lmax *dt << 1, where Lmax is the largest length of any block. Then we can make the approximetry that at most one car enters the roadway per block per time step. To decide Whether this happens and to choose the location p on the block if it does, we can do the following for each block b:

if (rand < dt * R* L(b)) nc = nc + 1p(nc) = rand * L(b)

When a can enters the roadway, it is assigned a destruction. This can also be done randomly. Let bd(c) be the black on which the destruction fres and let pd(c) be the position in that block, expressed as distance from the

Stant of the block. A simple way to make this choice is

bd(c) = 1 + floor(rand*nb)pd(c) = rand*L(bd(c))

but note that this choice grues equal weight to any block regardless of its buyth. To make the probability of choosing a block be proportional to its buyth, we can use the method of rejection:

bd(c) = 1 + floor (rand*nb) pd(c) = rand * Lmax while (pd(c) >= L(bd(c))) bd(c) = 1 + floor (rand*nb) pd(c) = rand * Lmax

end

In this version we keep trying until we find a position that fits on the block, and this makes the block that is ultimately chosen be more likely to be a longer one. In fact, the probability of choosing a block is exactly proportional to its length, and pd is uniformly distributed over that length.

Steering a canto its destination (despite one-way streets!)

For this we need the Cartesian coordinates of the destination, which are given by

xd(c) = xi(i1(bd(c))) + pd(c)*ux(bd(c))yd(c) = yi(i1(bd(c))) + pd(c)*uy(bd(c))

When a can comes to an intersection, it can choose to enter any of the blocks kewing that intersection. The natural choice is the one that most nearly points towards the destination. To determine this, evaluate the vector from the intersection to the destination, and then the dot product of that vector with all of the unit vectors of the blocks leaving the intersection. The block that should be chosen is the one that maximizes this dot product (in the algebraic sense, i.e., choose the most positive or least negative vesult, not the one layest in magnitude).

According to the above prescription, if can c is at messection i, it should choose the next block b to enter in the following way

xdvec = xd(c) - xi(i)ydvec = yd(c) - yi(i)

dp=ux(bout(i, 1:nbout(i)))* xdvec +uy(bout(i, 1:nbout(i)))* ydvec

[dpmax,jb] = max(dp)

b = bout(i, jb)

In the above use of max, there are two outputs. The second one, ib, is the index of the element of dp that has the maximum value.

The above steering algorithm works well for reasonable road networks, including cases in which it is necessary to go around the block to reach the destination because of one-way streets, but it is not guaranteed to work. For some roadway sayouts and some destinations, a can can get trapped and go through a cycle of

blocks repeatedly by follows the above algorithm without ever reaching its destination. One way to avoid this is the flee can to remember the intersections it has been to and the choices it has made those, and never make the same choice twice at any given whersection. Another way that is easier to program is for the can to decide randomly at each intersection whether to follow the above algorithm whether to follow the above algorithm. Or to choose a random block. This can be programed as block.

if (rand < prchoice) jb = 1 + floor (rand * nbout(i)) b = bout(i, jb)

else

choose b by the method of maximizing the dot product as described above

Here prehoice is the probability that a random choice will be made.

```
% man program: traffic.m
mitialize
for clock = 1: clock max
   t = clock * dt
   Setlighes
    morecars
plot cans
% Setlights. m if t > t/c
    for i=1:ni
       jgreen(i) = jgreen(i) + 1
if jgreen(i) > nbin(i)
        jgreen(i) = 1
end
    tlc = tlc + tlcstep
S = Zeros(1, nb)
```

S = Zeros(1, nb) f(i) = 1:ni b = bin(i, jgreen(i)) s(b) = 1end % mitialization for setlights

jgreen = ones(1, ni) tlcstep = % time interval between light changes tlc = tlcstep

16 greatecars. M for h=1:nb if (rand < dt * R*L(b)) nc = nc + 1p(nc) = rand * L(b) x(nc) = xi(i1(b)) + p(nc) * ux(b) y(nc) = yi(i1(b)) + p(nc) * uy(b) onroad(nc) = 1insertnewcan choose destruction nextb(nc) = btenter (nc) = t benter (nc) = b ponser (nc) = p(nc) end

% insertnewcan.
$$m$$
 $C = first can(b)$
 $if(C = = 0 | | p(nc) = p(c))$
 $next can(nc) = C$
 $first can(b) = nC$
 $if(C = = 0)$
 $lost can(b) = nC$
 $else if p(nc) <= p(last can(b)) = nC$
 $last can(last can(b)) = nC$
 $else$
 $Ca = C$
 $C = rext can(c)$
 $while(p(nc) <= p(c))$
 $Ca = C$
 $C = next can(c)$
 end
 $next can(nc) = C$
 end

% choosedeshination. M
% use method of rejection to choose a
% block with probability proprisind to
% its kength, and with p uniformly
% distributed in that block.

bd(nc) = 1 + floor(rand * nb) pd(nc) = rand * Lmax while(pd(nc) >= L(bd(nc))) bd(nc) = 1 + floor(vand * nb) pd(nc) = rand * Lmax end xd(nc) = xi(il(bd(nc))) + pd(nc) * ux(bd(nc)) yd(nc) = yi(il(bd(nc))) + pd(nc) * uy(bd(nc))

% Lmax = max(L)

```
% movecars. m
for b=1:nb
  C = fixstcan(b)
   while (C>0)
      if(c = firstcan(b))
         if (bd(c) == b) 44 (pd(c) > p(c))
             d = dmax
         Cleif(S(b)==0)
             d = L(b) - p(c)
        else
            decidenext block
            if (lastcar (nextb(c)) > 0
              d = (L(b) - p(c)) + p(lastcan(nextb(c)))
            else
             d = dnax
           end
         cnd
      else
         d = p(ca) - p(c)
      end
     pz = p(c); nextc = nextcan(c)

p(c) = p(c) + dt * V(d)
```

if(b==bd(c))ff(pz < pd(c)) ff(pd(c) <= p(c)) removecar elseif(L(b) <= p(c)) p(c) = p(c) - L(b) if(nex+b(c) == bd(c)) ff(pd(c) <= p(c)) removecan else cartonextblock

end

else

 $\chi(c) = \chi i (i1(b)) + p(c) * u \chi(b)$ y(c) = y i (i1(b)) + p(c) * u y(b)ca = x i (i1(b)) + p(c) * u y(b)

end

C = NextC % saved value of next can(c)

end % while loop over cars on block

end % for loop over blocks

```
% decidenextblock.m
% only do this if decision is not already made
if nextb(c) == b
    i = i2(b)
    if rand < prohoice

jnext = 1 + floor (rand of nbout (i))
         Nextb(c) = bout(i, jnext)
    else
        xdvec = xd(c) - xi(i)
        ydvec = yd(c) - yi(i)

dp = ux(bout(i, 1: nbout(i))) * xdvec ...

+ uy(bout(i, 1: nbout(i))) * ydvec
        [dpmax, jnext] = max(dp)
   nextb(c) = bout (i, jnext)
end
end
```

16 renovecas. m onroad(c)=0; texit(c)=tif (c == firstcan(b)) firstcar(b) = nextcar(c) if (c == lastcar (b)) [astca1(b)=0

else

nextcan(ca) = nextcan(c) if(c = = lastcan(b))lastear (b) = ca

ond

10 not really needed, but ... x(c) = xd(c) y(c) = yd(c) nextcay(c) = 0% recall that we provisions set next (= next can(c)

$$\chi(c) = \chi i \left(i \frac{1}{(nextb(c))} + p(c) * u \times (nextb(c))$$

 $\chi(c) = \chi i \left(i \frac{1}{(nextb(c))} + p(c) * u \times (nextb(c)) \right)$

% plotears .m

if (nc > 0 & & Sum (onroad) > 0)

Set (hcars, 'xdata', x(find (onroad)), ...

'ydata', y(find(onroad)))

End