

Road map

◆ Midterm a'comin

Friday in class

Exams page on web site has info + practice problems

◆ Today's lecture

A first look at assembly

Where is our data stored?

The mov instruction and addressing modes

It's bits all the way down...

◆ Data representation so far

Integer (unsigned, 2's complement signed)

Char (ASCII)

Address (unsigned long)

Float/double (IEEE floating point)

Aggregates (arrays, structs)

◆ The code itself is binary too!

Instructions (machine encoding)



Compiling code, what happens?

simple.c

```
int find_max(int arr[], size_t n)
{
    int max = arr[0];
    for (size_t i = 1; i < n; i++)
        if (arr[i] > max)
            max = arr[i];
    return max;
}
...
```

```
myth> make
gcc simple.c -o simple
```

```
^ELF^B^A^A^@^@^@^@^@^@^@^@^@B^@
>^@^A^@^@^@^@300^D^@^@^@^@^@^@^@
^@^@^@^@^@^@^@370\225^@^@^@^@^@^@
^@^@^@^@^@^@8^@^@^@^@&^@#^@F^@^
^@^E^@^@^@^@^@^@^@^@^@^@^@^@^@
...
```

simple

Source file (in text form)
Compiler parses input
validates language rules,
generates assembly instructions
writes object file (in binary form)

What's in an object file?

objdump -d simple

00000000004005b6 <find_max>:

4005b6: 8b 07
4005b8: ba 01 00 00 00
4005bd: eb 0d
4005bf:
4005c2:
4005c4:
4005c6: 89 c8
4005c8: 48 83 c2 01
4005cc: 48 39 f2
4005cf: 72 ee
4005d1: f3 c3

Sequential
instructions are at
sequential addresses

machine code
each instruction
encoded in binary

Name of function,
memory address of code
(function pointer)

mov \$0x1,%edx
jmp 4005cc <find_max+0x16>
mov (%rdi,%rdx,4),%ecx
cmp %ecx,%eax
jge 4005c8 <find_max+0x12>
mov %ecx,%eax
add \$0x1,%rdx
cmp %rsi,%rdx
jb 4005bf <find_max+0x9>
repz retq

each machine instruction decoded
into human-readable
assembly

What is an assembly instruction?

```
4005c6: 89 c8
4005c8: 48 83 c2 01
4005cc: 48 39 f2
4005cf: 72 ee
```

\$0x1

is constant value
("immediate")

mov
add
cmp
jb

opcode
(instruction
name/type)

%ecx,%eax
\$0x1,%rdx
%rsi,%rdx
4005bf <find_max+0x9>

operands

(arguments to instruction)

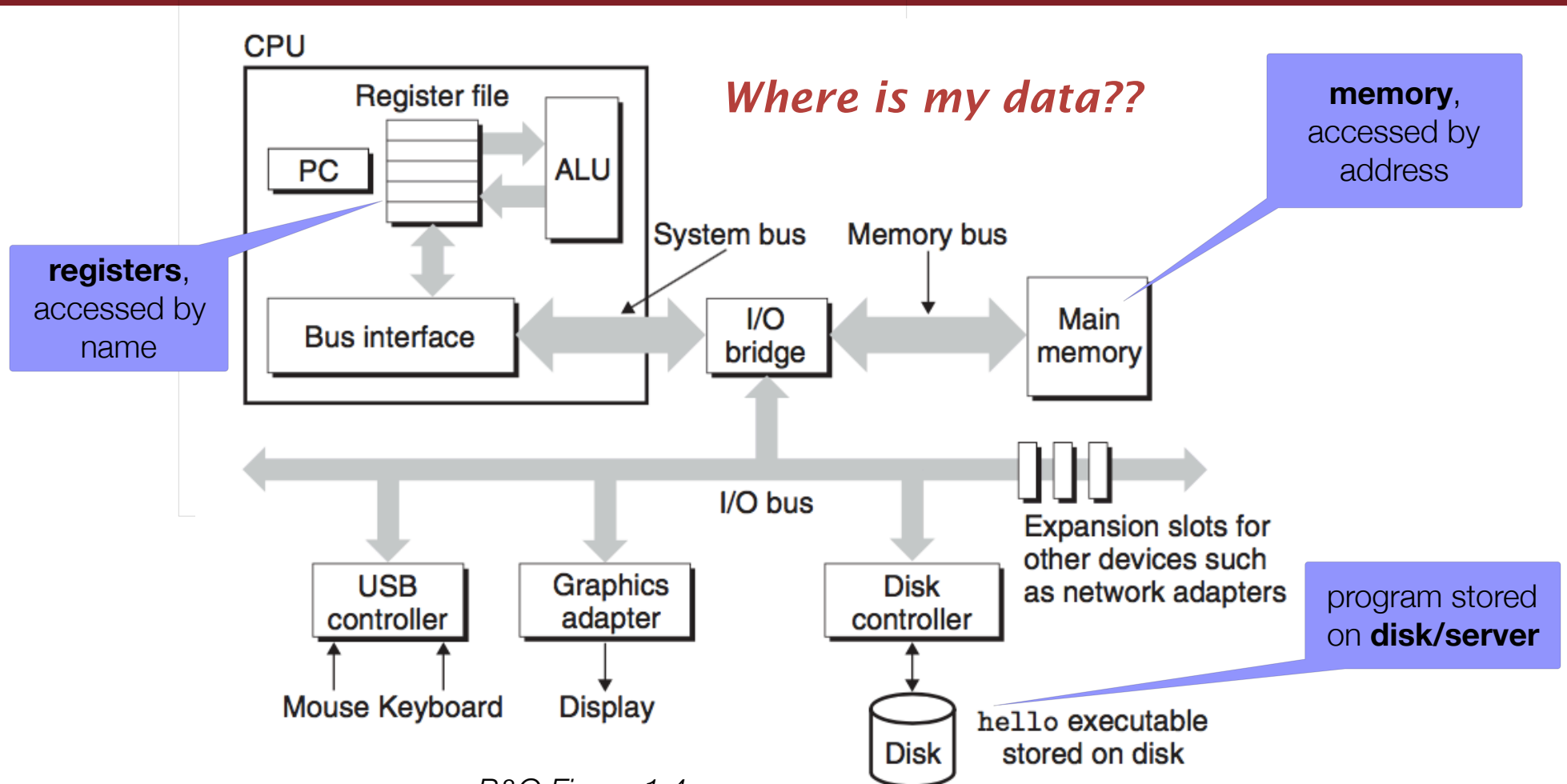
%eax

is register name,
(storage location on CPU)

4005bf

is direct address

Computer anatomy



B&O Figure 1.4

Instruction set architecture

◆ The ISA defines

Operations that the processor can execute

Data transfer operations, how to access data

Control mechanisms like branch, jump (think loops and if-else)

Contract between programmer/compiler and hardware

◆ Layer of abstraction

Above: programmer/compiler emits instructions as allowed in ISA

Below: hardware implements what is described in ISA

◆ ISAs have incredible inertia!

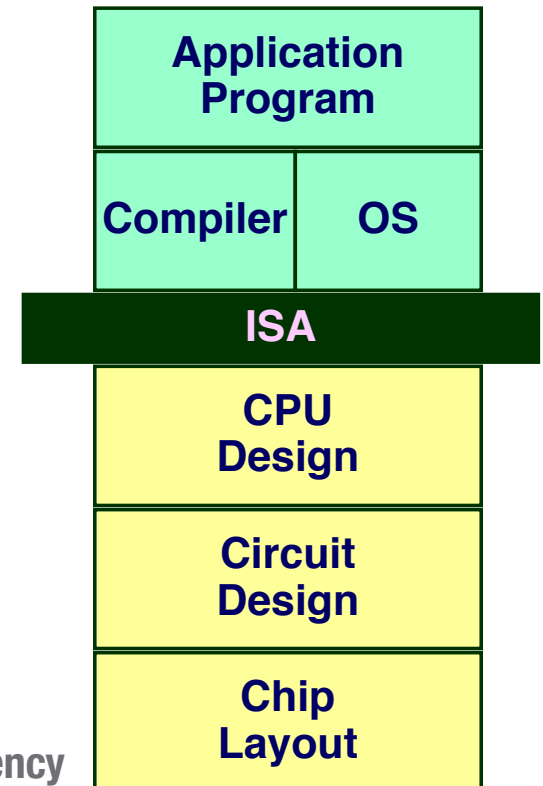
Legacy support is a huge issue for x86-64

◆ CISC vs RISC

(CISC, x86) Large set of specialized/expressive instructions, slower frequency

(RISC, ARM) Small set of simple instructions, higher frequency

Pres. Hennessy Turing Award!



Assembly characteristics

◆ Data

"integer" data, 1/2/4/8 bytes

Char, int, long, pointer, signed/unsigned

Floating point data, 4/8/(10) bytes

Special-purpose registers and instructions

No aggregates

Arrays and structs are just contiguously located bytes in memory

No names, no types

Refer to data by where stored (register/memory), size in bytes

◆ Operations

Perform arithmetic/logical ops on register or memory data

Transfer data between memory and register

Load/store

Control flow

Unconditional jump to/from other functions

Conditional branch

The almighty mov instruction

◆ Programs manipulate data

Where is that data stored? registers, memory (also: disk, server, network, ...)

◆ **mov instruction is the assembly equivalent of assignment**

Most common instruction of all

◆ **Key insight: no access to variables by name/type**

High-level language had descriptive names, type information

Assembly accesses variable by identifying where it is stored (register/memory)

◆ **General form:** `movx src, dst`

Copy bytes from one place to another

Source can be memory, registers, constants

Destination can be memory, registers

Mov operands: imm/reg

Op	Src	Dst	Comments
movl	\$0,	%eax	src is immediate
movb	\$0x41,	%al	Virtual sub-register
mov	%rax,	%rdx	Register to register

movx suffix is how many bytes to move

b for byte (1), **w** for word (2), **l** for long (4), **q** for quad (8)

(suffixes show legacy...)

Elided if can be inferred from operands

63	31	15	8	7	0
%rax	%eax %ax	%ah %al			
%rbx	%ebx %bx	%bh %bl			
%rcx	%ecx %cx	%ch %cl			
%rdx	%edx %dx	%dh %dl			
%rsi	%esi %si	%sil			
%rdi	%edi %di	%dil			
%rbp	%ebp %bp	%bpl			
%rsp	%esp %sp	%spl			
%r8	%r8d %r8w	%r8b			
%r9	%r9d %r9w	%r9b			
%r10	%r10d %r10w	%r10b			
%r11	%r11d %r11w	%r11b			
%r12	%r12d %r12w	%r12b			
%r13	%r13d %r13w	%r13b			
%r14	%r14d %r14w	%r14b			
%r15	%r15d %r15w	%r15b			

Mov operands: direct/indirect

Op	Src	Dst	Comments
movl	\$0,	0x605428	Store, direct address (Note no prefix on address literal)
movl	\$0,	(%rsp)	Store, indirect address (address in register, dereference)
movl	0x605428,	%edx	Load
movl	(%rsp),	%edx	Load

Load = read from memory location

Store = write to memory location

No mem-to-mem transfer

Either src or dst is memory, not both

Direct: Data at fixed location

Indirect: Register holds pointer

Addressing modes

Op	Src	Dst	Comments
movl	\$0,	0x605428	Direct address
movl	\$0,	(%rsp)	Indirect address
movl	\$0,	20(%rsp)	Indirect with displacement

Displacement
is any constant
(negative or positive)

Base

Target address = base + displacement

Addressing modes

Op	Src	Dst	Comments
movl	\$0,	0x605428	Direct address
movl	\$0,	(%rsp)	Indirect address
movl	\$0,	20(%rsp)	Indirect with displacement
movl	\$0,	20(%rsp, %rax, 4)	Indirect with scaled-index

Displacement
is any constant
(negative or positive)
If missing, =0

Base
register
if missing, = 0

Index
register

Scale
must be 1, 2, 4, or 8
if missing, =1

Target address =
base + displacement + index*scale

Load effective address

◆ **lea** = "load effective address"

Basically a **mov** without the dereference

Used for address calculation, e.g. `&arr[x]`

Also arithmetic expressions of form $x + ky$ (faster than sequenced mul/add)

where $k = 1, 2, 4, 8$

◆ **Examples**

`leal (%rax, %rsi, 4), %rax`

Computes base + scaled-index, e.g. address of array elem

`leal 7(%rdx, %rdx, 4), %rdx`

Computes $x = 5x + 7$ (assuming x stored in `%rdx`)