

Name: _____

By writing my name I promise to abide by the Honor Code

Section number: _____

Read the following:

- **RIGHT NOW!** Write your name on the top of your exam.
- You are allowed one $8\frac{1}{2} \times 11$ in sheet of **handwritten** notes (both sides). No magnifying glasses!
- You may use a calculator provided that it cannot access the internet or store large amounts of data.
- You may **NOT** use a smartphone as a calculator.
- Clearly mark answers to multiple choice questions on the provided answer line.
- Mark only one answer for multiple choice questions. If you think two answers are correct, mark the answer that **best** answers the question. No justification is required for multiple choice questions.
- If you do not know the answer to a question, skip it and come back to it later.
- For free response questions you must clearly justify all conclusions to receive full credit. A correct answer with no supporting work will receive no credit.
- If you need more space for free-response questions, there are blank pages at the end of the exam. If you choose to use the extra pages, make sure to **clearly** indicate which problem you are continuing.
- You have **90 minutes** for this exam.

Page	Points	Score
2	12	
3	12	
4	16	
5	20	
6	20	
7	20	
Total	100	



1. (3 points) Consider the pseudocode for procedure `parity_party(A,B,n)`, given below. The input to the procedure is two matrices A and B , and scalar n . $A[i][j]$ and $B[i][j]$ (the element in row i , column j of matrices A and B , respectively) are all integers. n gives the number of rows and columns of A and B (they are both $n \times n$ square matrices). Give an estimate of the complexity of this procedure, where complexity is measured by the number of **additions and subtractions** needed. Justify your answer.

```

procedure parity_party(A, B, n):
    par_ctr = 0
    for i from 1 to n:
        for j from 1 to n:
            if (A[i][j] + B[i][j]) % 2 == 0,
                then par_ctr = par_ctr + 1
            else
                then par_ctr = par_ctr - 1
    return odd_ctr

```

- A. `parity_party` is order n
- B. `parity_party` is order n^2
- C. `parity_party` is order n^3
- D. `parity_party` is order n^4

1. **B**

2. (3 points) What is the **smallest** integer p such that $f(n) = 3n^2 + \log(n^2) + n \log(n^4)$ is $\mathcal{O}(n^p)$?

- A. $p = 2$
- B. $p = 3$
- C. $p = 4$
- D. $p = 5$

2. **A**

3. (3 points) Select the answer that is a closed form solution to this recurrence relation:

$$a_n = 2a_{n-1} + 3, a_0 = 1$$

- A. $a_n = 2^n - 3$
- B. $a_n = 2^{n+2} - 3$
- C. $a_n = 4n + 1$
- D. $a_n = 2n^2 + 2n + 1$

3. **B**

4. (3 points) Suppose b and q are integers, p and m are positive integers, and that $mb + pq \equiv 3 \pmod{m}$. Which of the following necessarily must be true?

- A. $pq = 1$
- B. $pq \equiv 1 \pmod{m}$
- C. $pq \equiv 2 \pmod{m}$
- D. $pq \equiv 3 \pmod{m}$

4. **D**

9. (8 points) Use either **Chinese Remainder Theorem** or **Back Substitution** to find all solutions x to the system of congruences:

CRT:

$$x \equiv 2 \pmod{4}$$

$$x \equiv 4 \pmod{7}$$

$$x \equiv a_1 M_1 y_1 + a_2 M_2 y_2 \pmod{m}$$

$$m = 4 \times 7 = 28$$

$$M_1 = \frac{m}{m_1} = 7 \quad \& \quad M_2 = \frac{m}{m_2} = 4$$

$$y_1 = \text{inverse of } M_1 \pmod{m_1} \\ 7 \pmod{4}$$

$$\underline{y_1 = 3 \text{ works!}} \quad (7 \cdot 3 = 21 \equiv 1 \pmod{4})$$

$$y_2 = \text{inverse of } M_2 \pmod{m_2} \\ 4 \pmod{7}$$

$$\underline{y_2 = 2 \text{ works!}} \quad (4 \cdot 2 = 8 \equiv 1 \pmod{7})$$

$$\Rightarrow x \equiv 2 \cdot 7 \cdot 3 + 4 \cdot 4 \cdot 2 \pmod{28} \\ \equiv 42 + 32 \pmod{28}$$

$$\boxed{x \equiv 18 \pmod{28}}$$

Back Sub:

$$x \equiv 2 \pmod{4} \rightarrow x = 2 + 4k \quad (k \in \mathbb{Z})$$

$$\rightarrow 2 + 4k \equiv 4 \pmod{7}$$

$$\rightarrow 4k \equiv 2 \pmod{7}$$

$$\rightarrow \text{need inverse of } 4 \pmod{7}$$

$$\rightarrow \underline{2 \text{ works!}} \quad (\text{see left})$$

$$\rightarrow k \equiv 2 \cdot 2 \pmod{7}$$

$$\rightarrow k = 4 + 7l \quad (l \in \mathbb{Z})$$

$$\rightarrow x = 2 + 4(4 + 7l) \\ = 18 + 28l$$

$$\rightarrow \boxed{x \equiv 18 \pmod{28}}$$

10. (8 points) Come up with a formula (in terms of n) for the number of **trit-strings** of length n that are **NOT** palindromic. Reminders:

- You must show work to receive credit
- Trit-strings are made up of 0s, 1s and 2s
- A palindrome is the same forwards and backwards. For example: 102201.

NOTE: other answers exist!

Even n : 102201

$\frac{n}{2}$ choices to make, & 3 options for each

$$\rightarrow \boxed{f(n) = 3^{n/2} \quad \text{if } n \text{ even}}$$

Odd n : 10201

$\frac{n+1}{2}$ choices to make, 3 options for each

$$\rightarrow \boxed{f(n) = 3^{\frac{n+1}{2}} \quad \text{if } n \text{ odd}}$$

11. (20 points) Let $x > 0$ be some fixed real number.

Use induction to prove that $(1+x)^n > 1+nx$ for all $n \geq 2$. Be sure to mention whether you are using strong or weak induction.

Solution: Base case : $n=2$
Strong induction $(1+x)^2 = 1+2x+x^2 > 1+2x$ since $x > 0$

IH: assume for all $2 \leq k \leq n$ $(1+x)^k > 1+kx$

will show: for $k=n+1$: $(1+x)^{n+1} > 1+(n+1)x$

$$\begin{aligned} \cdot (1+x)^{n+1} &= (1+x)^n (1+x) \stackrel{\text{IH}}{>} (1+nx)(1+x) \\ &= 1+x+nx+nx^2 = 1+(n+1)x + nx^2 > \\ &\quad 1+(n+1)x \end{aligned}$$

~~□~~

12. (20 points) After a very successful surfing season, the Interstellar Surfing Association (ISA) is putting on surfing demonstrations. There are 8 demonstrations total, which are split up between the two champion surfers (Alex and Tony, of course). Neither surfer can perform 4 demonstrations in a row, but either of them can perform more (or less) than 4 demonstrations total.

For example, one possible arrangement is for Alex to do 2 demonstrations, then Tony do 3, then Alex do 3.

Determine how many possible ways are there for the ISA to organize those 8 demonstrations. You do not need to simplify your answer.

Solution: The problem is equivalent to asking how many length 8 binary strings are there with no 4 consecutive 1's or 0's.

[Let Alex performing = 0 Tony Performing = 1
Then arrangements of surfing demonstrations
= binary strings]

of length 8 binary strings with no 4 consecutive 1's or 0's = # of length 8 bin. strings - # of length 8 bin strings with 4 consecutive 1's or 4 cons. 0's

$|A| = |B| - |C|$

• $|A| = 2^8$

• $|C|$: We calculate # strings with 4 consecutive 0's
strings with 4 consecutive 1's is same.

5 positions that 0000 can start:

1) 0000XXXX $\Rightarrow 2^4$

2) 10000XXX $\Rightarrow 2^3$

3) X10000XX $\Rightarrow 2^3$

4) XX10000X $\Rightarrow 2^3$

5) XXX10000 $\Rightarrow 2^3$

• Total = $2^4 + 4 \cdot 2^3 = 16 + 4 \cdot 8 = 16 + 32 = 48$

• same for 1's = 48

• Two strings with 4 cons. 0's & 1's
00001111 } need to subtract
11110000

$|C| = 2 \cdot 48 - 2 = 94$

$|A| = 2^8 - 94 = 256 - 94 = 162$

13. (20 points) Consider the function $f(n) = 3n^2 + \log(n^3) - n$.

- (a) Find a tight big- \mathcal{O} bound for $f(n)$. Be sure to specify your values of C and k in the definition of big- \mathcal{O} .
- (b) Find a tight big- Ω bound for $f(n)$. Be sure to specify your values of C and k in the definition of big- Ω .
- (c) Can you state that $f(n)$ is $\Theta(h(n))$ for some function $h(n)$? If so, state $h(n)$ and briefly justify your reasoning.

Solution:

a.) $f(n) = 3n^2 + \log(n^3) - n = 3n^2 + 3\log n - n$

$$3\log n \leq 3n \quad \text{for all } n > 1$$
$$\leq 3n^2$$
$$\text{and } -n < 0 \quad \text{for all } n > 0$$
$$\Rightarrow f(n) \leq 3n^2 + 3n^2 + 0 = 6n^2 \quad \text{for } n > 1$$
$$\Rightarrow \boxed{f \text{ is } O(n^2) \text{ w/ } C=6 \text{ \& } k=1}$$

b.) $n \leq n^2 \quad \text{for } n > 1$

$$\Rightarrow \underline{-n \geq -n^2}$$
$$\Rightarrow f(n) \geq 3n^2 + 3\log n - n^2$$
$$\geq 3n^2 - n^2 = 2n^2 \quad \text{for } n > 1$$
$$\Rightarrow \boxed{f \text{ is } \Omega(n^2) \text{ w/ } C=2 \text{ \& } k=1}$$

c.) Yes - $\boxed{f \text{ is } \Theta(n^2)}$ b/c it is both $O(n^2)$ \& $\Omega(n^2)$