



Introdução à Programação Paralela em Clusters Heterogêneos

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ERAD-SP 2018

UNESP Center for Scientific Computing

- Consolidates scientific computing resources for São Paulo State University (UNESP) researchers
 - It mainly uses Grid computing paradigm
- Main users
 - UNESP researchers, students, and software developers
 - SPRACE (São Paulo Research and Analysis Center)
 - physicists and students
- Caltech, Fermilab, CERN
- São Paulo CMS Tier-2 Facility



13th International Meeting on High Performance Computing for Computational Science

- Parallel Architectures
- Parallelism Levels
- Xeon and Xeon Phi
- Heterogeneous Cluster @ NCC
- Offload Programming
- Optimization Examples
- Hands-on

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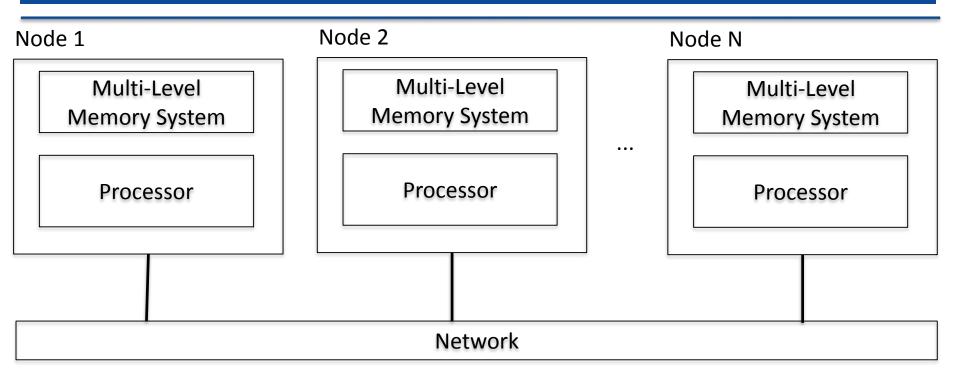
Parallel Processing

- □ A parallel computer is a computer system that uses multiple processing elements simultaneously in a cooperative manner to solve a computational problem
- Parallel processing includes techniques and technologies that make it possible to compute in parallel
 - ☐ Hardware, networks, operating systems, parallel libraries, languages, compilers, algorithms, tools, ...
- ☐ Parallel computing is an evolution of serial computing
 - □ Parallelism is natural
 - ☐ Computing problems differ in level / type of parallelism

Parallelism

- Concurrency = Different units of work that can be executed in parallel
 - Ex: loop Iterations
- Parallelism = concurrency + parallel hardware
 - Find concurrent execution opportunities
 - Develop application to execute in parallel
 - Run application on parallel hardware
- Motivations for parallelism
 - Faster time to solution
 - Solve bigger computing problems
 - Effective use of machine resources
- Serial machines have inherent limitations
 - Processor speed, memory bottlenecks, ...

Homogeneous Parallel Architectures

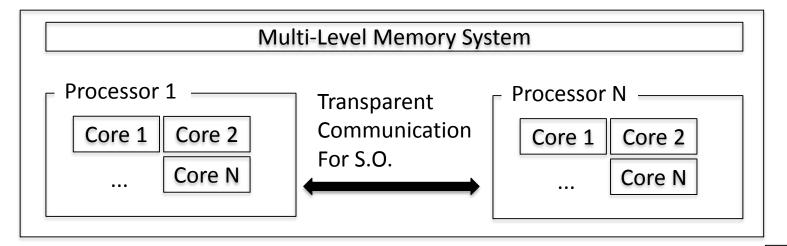


First initiatives:

- Identical Nodes interconnected;
- Parallel architectures is deployed as a group of serial computers;

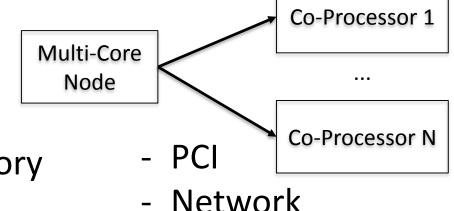
Heterogeneous Parallel Architectures

Multi-Core Node



Multi-Core Architecture

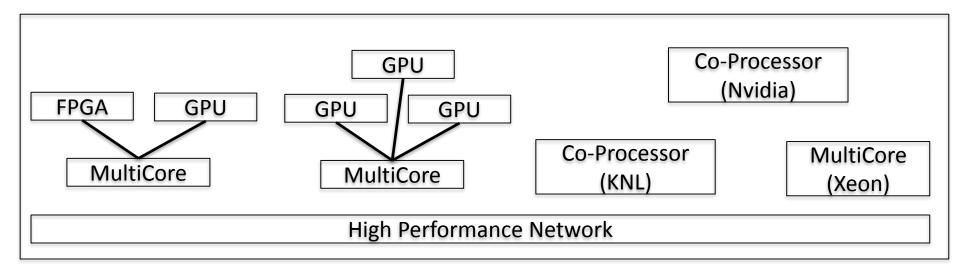
- Multi Processors
- Multiple Cores
- **NUMA** (Non-Uniform Memory Access)



Heterogeneous Parallel Architectures

- Several families and architectures:
- Hybrid Parallel nodes
 - Coprocessors and accelerators;

- Multi-level parallelism:
 - Processing core;
 - Chip multiprocessor;
 - Computing node;
 - Computing cluster;



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Exploiting the parallel universe

Instruction Level Parallelism

- Single thread (ST) performance
- Automatically exposed by HW/tools
- Effectively limited to a few instructions

Data Level Parallelism

• Single thread (ST) performance

Vectorization

- Exposed by tools and programming models
- Operate on 4/8/16 elements at a time

Task Level **Parallelism**

- Multi thread/task (MT) performance
- **OpenMP**
- Exposed by programming models
- Execute tens/hundreds/thousands task concurrently

Process Level Parallelism

• Multi Process (MP) performance

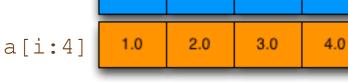
MPI

- Exposed by programming models
- Execute tens/hundreds/thousands of process concurrently across several nodes

Vectorization

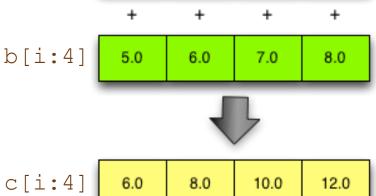
Vectorization

- Loading data into cache accordingly;
- Store elements on SIMD registers or vectors;
- Apply the same operation to a set of Data at the same time;
- Iterations need to be independent;
- Usually on inner loops.



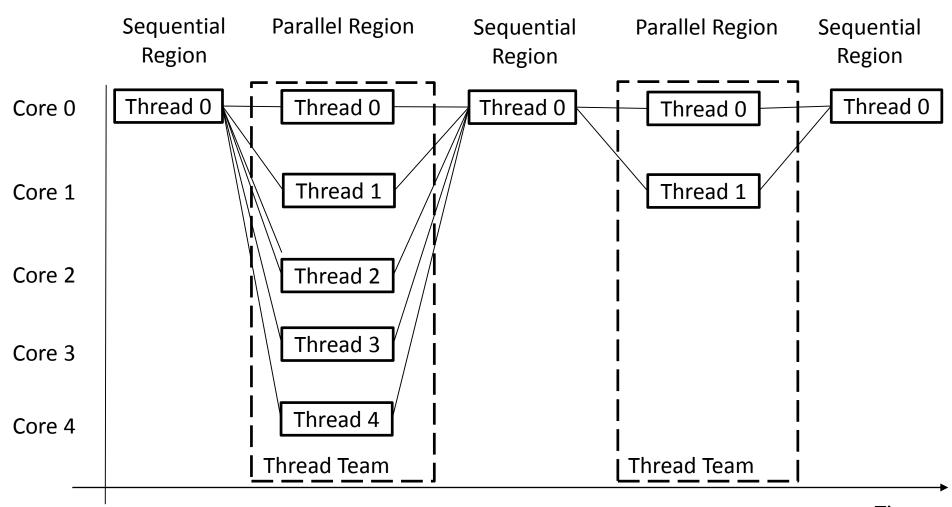
Scalar loop

for	(int	i	= 0;	i	< N; i++)
	c[i]	=	a[i]	+	b[i];



Z

OpenMP

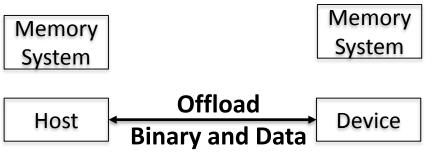


Time

Offload

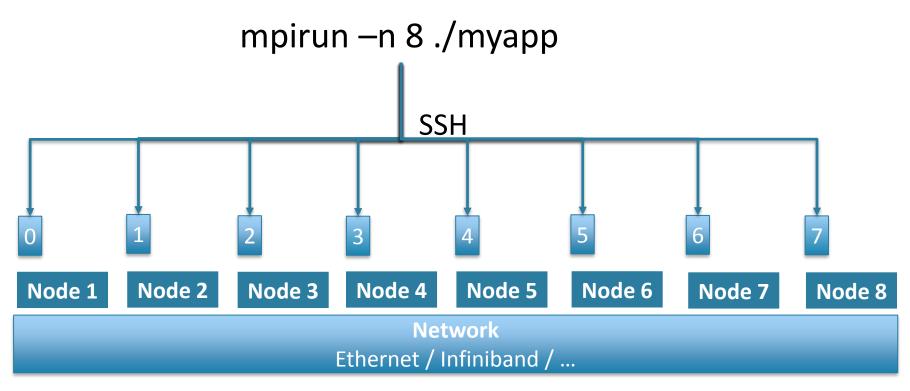
- Regions of code is marked to be executed in a device attached to the main processor
 - Communication using PCI Express or Network

- Main Processor and Device has independent Memory system
 - Developer need to control data transfer of data structures needed by application
 - Can Present Delay



Message Passing Interface (MPI)

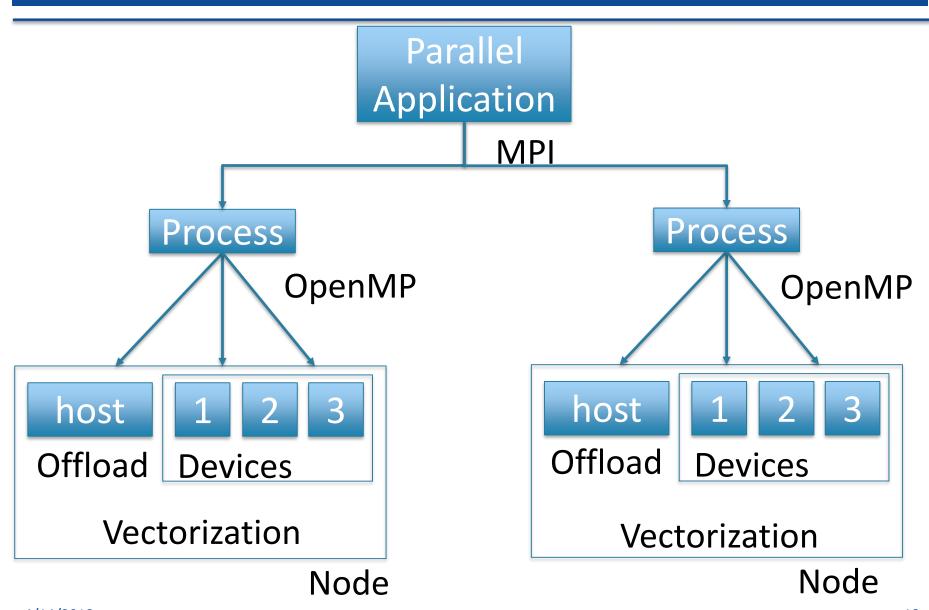
Execution of a MPI application across several nodes using SSH (Secure Shell)



Infiniband: computer-networking communications standard used in high-performance computing

- Very high throughput;
- Very low latency.

Combining Levels

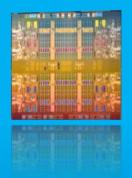


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Intel Xeon and Intel® Xeon Phi™ Overview

Intel® Multicore Architecture

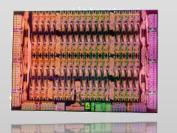




- Foundation of HPC Performance
- Suited for full scope of workloads
- Focus on fast single core/thread performance with "moderate" number of cores

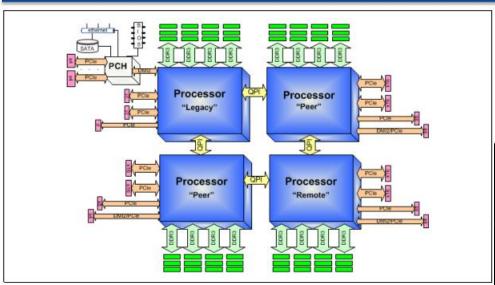
Intel[®] Many Integrated Core Architecture



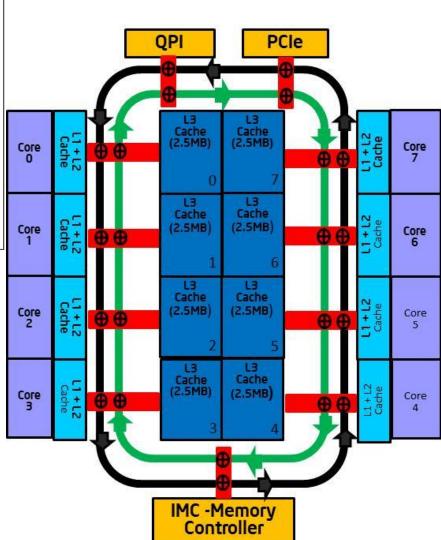


- Performance and performance/watt optimized for highly parallelized compute workloads
- ❖ IA extension to Manycore
- Many cores/threads with wide SIMD

Intel Xeon Architecture Overview



- •Socket: mechanical component that provides mechanical and electrical connections between a microprocessor and a printed circuit board (PCB).
- •QPI (Intel QuickPath Interconnect): high speed, packetized, point-to-point interconnection, that stitch together processors in distributed shared memory and integrated I/O platform architecture.



Knights Landing (KNL)

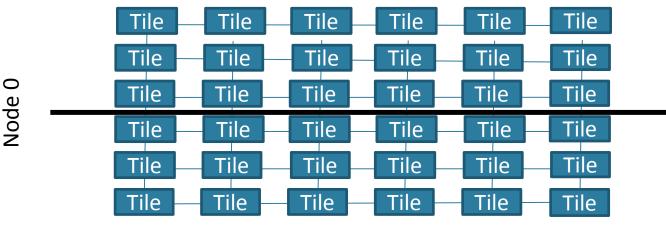
Up to 16GB high-bandwidth on-package Over 3 TF DP peak memory (MCDRAM) Full Xeon ISA compatibility through AVX-512 ~3x single-thread vs. compared to Knights Exposed as NUMA node Corner ~500 GB/s sustained BW **Up to 72** MCDRAM MCDRAM MCDRAM cores DDR4 Tile DDR4 2x 512b VPU 2D mesh 2 VPU H Up to per core DDR4 VPU architecture DDR4 Core Expression Core (Vector 72 cores DDR4 DDR4 **Processing** Units) MCDRAM MCDRAM MCDRAM **MCDRAM** 6 channels HEL DDR4 Based on Intel® Atom Silvermont processor with Connector Up to many HPC enhancements 384GB Twinax Cable Deep out-of-order buffers Gather/scatter in hardware 2 ports Intel Omni-Path Fabric On-Improved branch predition package 4 threads/core 50 GB/s bi-directional High cache bandwidth

Cluster modes

One single space address

Hemisphere:

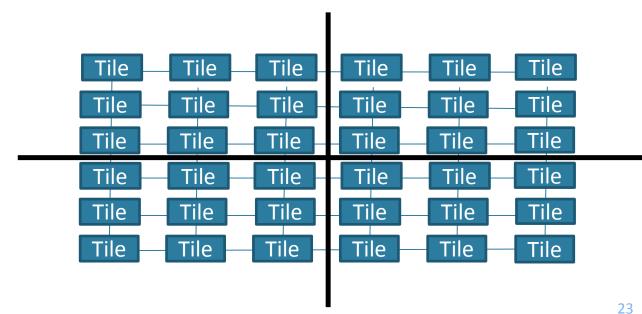
the tiles are divided into two parts called hemisphere



Quadrant:

tiles are divided into two parts called hemisphere or into four parts called qudrants

Node 0

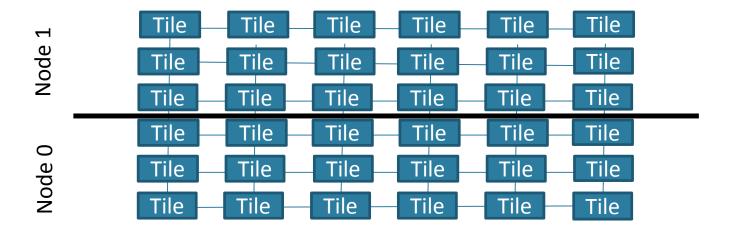


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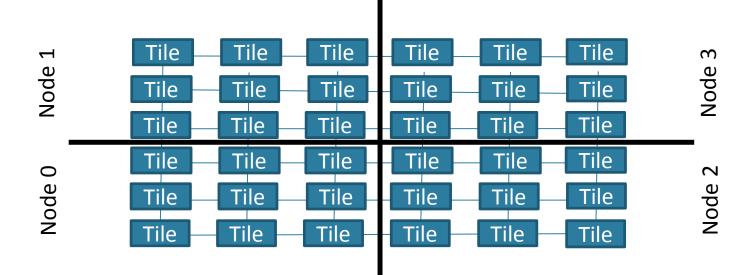
Cluster modes

Cache data are isolated in each sub numa domain

SNC-2: the tiles are divided into two Numa Nodes

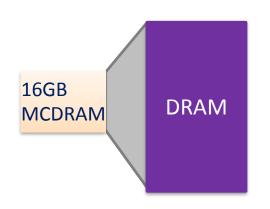


SNC-4: the tiles are divided into two Numa Nodes

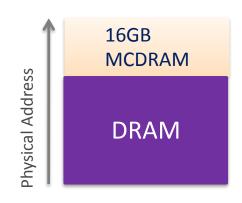


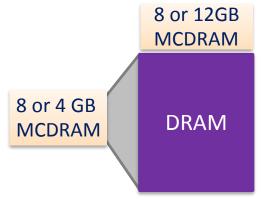
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Integrated On-Package Memory Usage Models



memory access behavior





Split Options: 25/75% or 50/50%

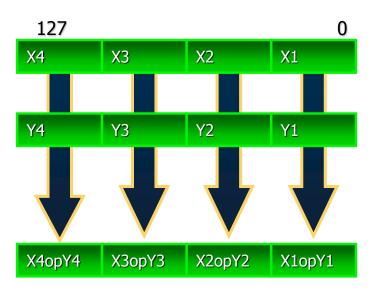
be configured as Cache

Cache Model	Flat Model	Hybrid Model	
Hardware automatically manages the MCDRAM as a "L3 cache" between CPU and ext DDR memory	Manually manage how the app uses the integrated on-package memory and external DDR for peak perf	Harness the benefits of both Cache and Flat models by segmenting the integrated on- package memory	
 App and/or data set is very large and will not fit into MCDRAM Unknown or unstructured 	 App or portion of an app or data set that can be, or is needed to be "locked" into MCDRAM so it doesn't get 	 Need to "lock" in a relatively small portion of an app or data set via the Flat model Remaining MCDRAM can then 	

flushed out

Vectorial processing Intel® Architecture

MMX	Intel® SSE	Intel® AVX / AVX2	Intel® MIC / AVX-512
Vector size: 64bit	Vector size: 128bit	Vector size: 256bit	Vector size: 512bit
Data types: 8, 16 and	Data types:	Data types: 32 and	Data types:
32 bit integers	8,16,32,64 bit	64 bit floats	32 and 64 bit
VL: 2,4,8	integers	VL: 4, 8, 16	integers
	32 and 64bit float	Sample: Xi, Yi 32 bit	32 and 64bit float
	VL: 2,4,8,16	int or float	VL: 8,16



AVX -512

Exponential & Reciprocal Instructions (ERI)

Prefetch Instructions (PFI)

Foundation instructions (FI)

Conflict Detection Instructions (CDI)

Byte & Word Instructions (BWI)

Double-/Quad-word Instructions (DQI)

Vector Length Extensions (VLE)

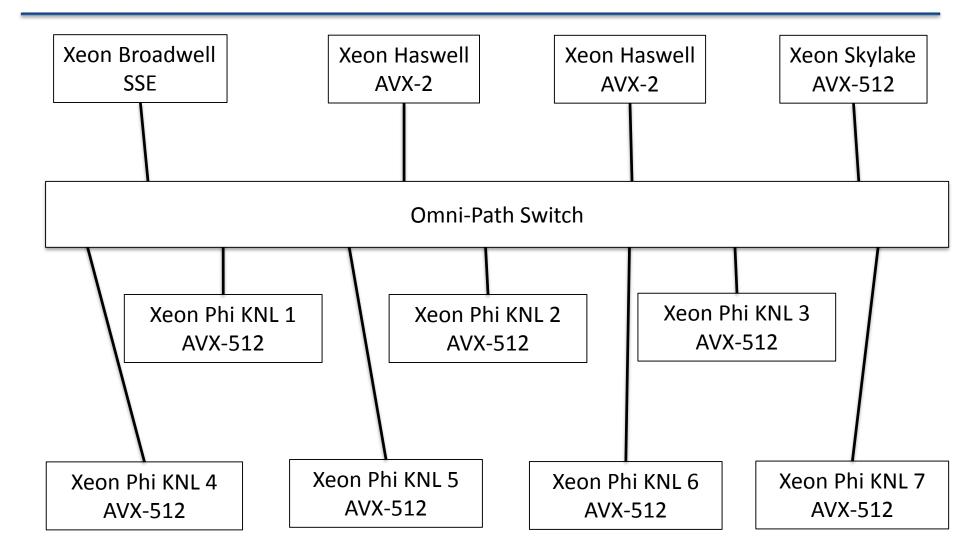
AVX-512 - Foundation

 Implements new version of AVX-2 functions using 512bit registers (ZMM);

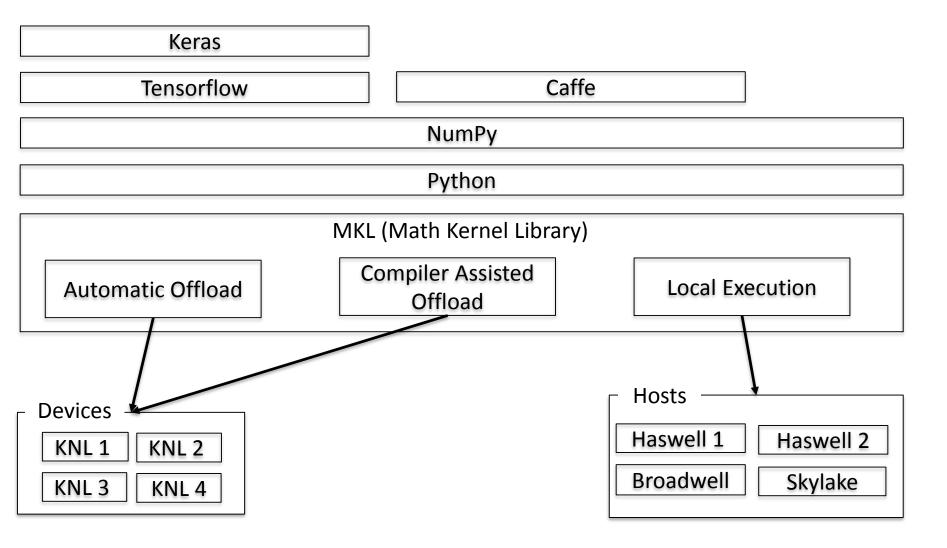
Xmm0Xmm7	Ymm0Ymm15	Zmm0Zmm31
SSE		
	→ AVX	
		AVX-512

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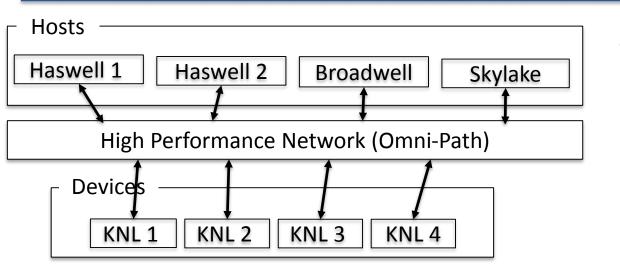
Heterogeneous Cluster @ NCC



Parallelism Model - MKL

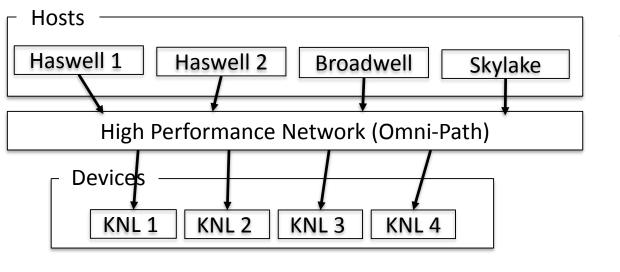


Parallelism Model – MPI x Offload



MPI (Message Passing Interface)

- High Performance Communication using MPI Functions
- Every node send and receive messages



Offload

- One Host node can offload code to one or more devices
- Hosts only send code regions for execution
- Devices only execute code regions

Offload Over Fabric

 Offload over Fabric (OOF) enables offloading from servers with Intel Xeon processor to servers with Intel Xeon Phi x200 processor within a high-speed network, such as Intel® Omni-Path Fabric.

- OFFLOAD_NODES environment variable define nodes available for the offloading operation from the offload host.
 - OFFLOAD_NODES is a comma-separated list of nodes' names or IP addresses.

Offload Over Fabric

icc -qopenmp testHW.c -o testHW

export OFFLOAD_NODES=10.0.0.5 ./testHW

Running on host: phi01.ncc.unesp.br

Running on target: phi05

export OFFLOAD_NODES=10.0.0.6 ./testHW

Running on host: phi01.ncc.unesp.br

Running on target: phi06

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OpenMP 4.0 Offload

- target: transfers the control flow to the target device
 - Transfer is sequential and synchronous
 - Transfer clauses control data flow
- target data: creates a scoped device data environment
 - Does not include a transfer of control
 - Transfer clauses control data flow
 - The device data environment is valid through the lifetime of the target data region
- target update: request data transfers from within a target data region
- **omp declare target:** creates a structured-block of functions that can be offloaded.

OpenMP 4.0 Offload Report

OFFLOAD REPORT:

- Measures the amount of time it takes to execute an offload region of code;
- Measures the amount of data transferred during the execution of the offload region;
- Turn on the report: export OFFLOAD REPORT=2
- [Var] The name of a variable transferred and the direction(s) of transfer.
- [CPU Time] The total time measured for that offload directive on the host.
- [MIC Time] The total time measured for executing the offload on the target.
- [CPU->MIC Data] The number of bytes of data transferred from the host to the target.
- [MIC->CPU Data] The number of bytes of data transferred from the target to the host.

Pragma omp declare target

Creates a structured-block of functions that can be offloaded.

- Syntax
 - #pragma omp declare target [clause[[,] clause],...] declaration of functions
 - #pragma omp end declare target

Pragma omp target

Transfer control [and data] from host to device

Syntax

- #pragma omp target [data] [clause[[,] clause],...]
 structured-block

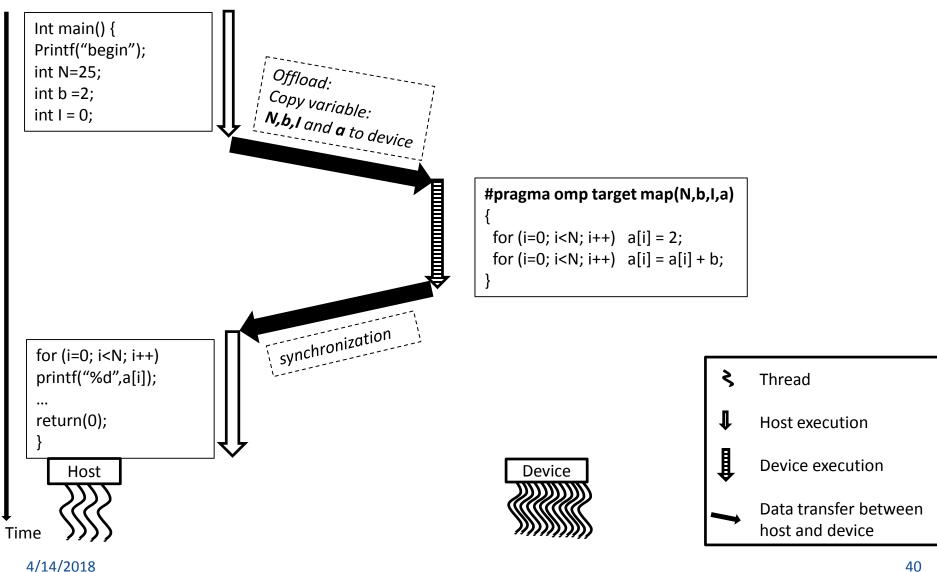
Clauses

Pragma omp target

Map clauses:

- alloc : allocate memory on device;
- to: transfer a variable from host to device;
- from: transfer a variable from device to host;
- tofrom :
 - □ transfer a variable from host to device before start execution;
 - □ transfer a variable from device to host after finish execution;

Offloading - omp target



Pragma omp target example

```
#pragma omp target device(0) map(a[0:NUM][0:NUM])
map(b[0:NUM][0:NUM]) map(c[0:NUM][0:NUM])
  #pragma omp parallel for collapse (2)
  for(i=0; i<msize; i++) {
    for(k=0; k<msize; k++) {
      #pragma omp simd
      for(j=0; j<msize; j++) {
        c[i][j] = c[i][j] + a[i][k] * b[k][j];
```

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N-Body Simulation

- An N-body simulation [1] aims to approximate the motion of particles that interact with each other according to some physical force;
- Used to study the movement of bodies such as satellites, planets, stars, galaxies, etc., which interact with each other according to the gravitational force;
- Newton's second law of motion can be used in a N-body simulation to define the bodies' movement.

[1] AARSETH, S. J. Gravitational n-body simulations. [S.I.]: Cambridge University Press, 2003. Cambridge Books Online.

N-Body Algorithm

Bodies struct:

- 3 matrix represents velocity (x,y and z)
- 3 matrix represents position (x,y and z)
- 1 matrix represent mass

A loop calculate temporal steps:

 At each temporal step new velocity and position are calculated to all bodies according to a function that implements Newton's second law of motion

N-Body - Parallel version (host only)

```
function Newton(step)
  #pragma omp for
  for each body[x] {
    #pragma omp simd
    for each body[y]
      calc force exerted from body[y] to body[x];
    calc new velocity of body[x]
  #pragma omp simd
  for each body[x]
     calc new position of body[x]
Main() {
  for each temporal step
    Newton(step)
```

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N-Body - Parallel version (Load balancing)

The temporal step loop remains sequential

 The N-bodies are divided among host and devices to be executed using Newton

- OpenMP offload pragmas are used to
 - Newton function offloading to devices
 - Transfer data (bodies) between host and devices

N-Body - Parallel version (Load balancing)

```
function Newton(step, begin_body, end_body, deviceId)
  #pragma omp target device (deviceId) {
    #pragma omp for
    for each body[x] from subset(begin_body, end_body) {
      #pragma omp simd
      for each body[y] from subset(begin_body, end_body)
        calc force exerted from body[y] to body[x];
      calc new velocity of body[x]
    #pragma omp simd
    for each body[x]
       calc new position of body[x]
```

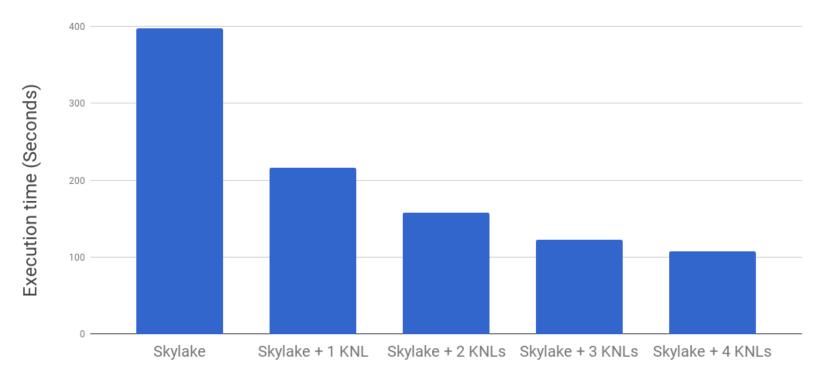
N-Body - Parallel version (Load balancing)

```
for each temporal step
  Divide the amount of bodies among host and devices;
  #pragma omp parallel
    #pragma omp target data device (tid) to(bodies[begin_body:
end body])
      Newton(step, begin_body, end_body, deviceId)
      #pragma omp target update device (tid) (from:bodies)
      #pragma omp barrier
      #pragma omp target data device (tid)
to(bodies[begin body: end body])
```

N-Body

N-Body Evaluation using offload over fabric (Xeon - Xeon Phi and Omni-Path)

Amount of Particles in Simulation: 30000



Amount of Skylake and KNLs

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