



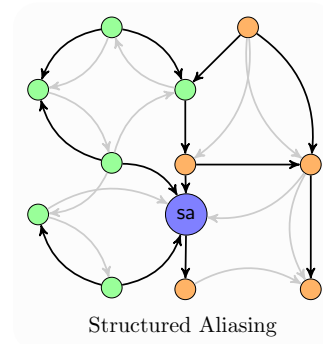
Actors without Borders: Amnesty for Imprisoned State

Elias Castegren, Tobias Wrigstad

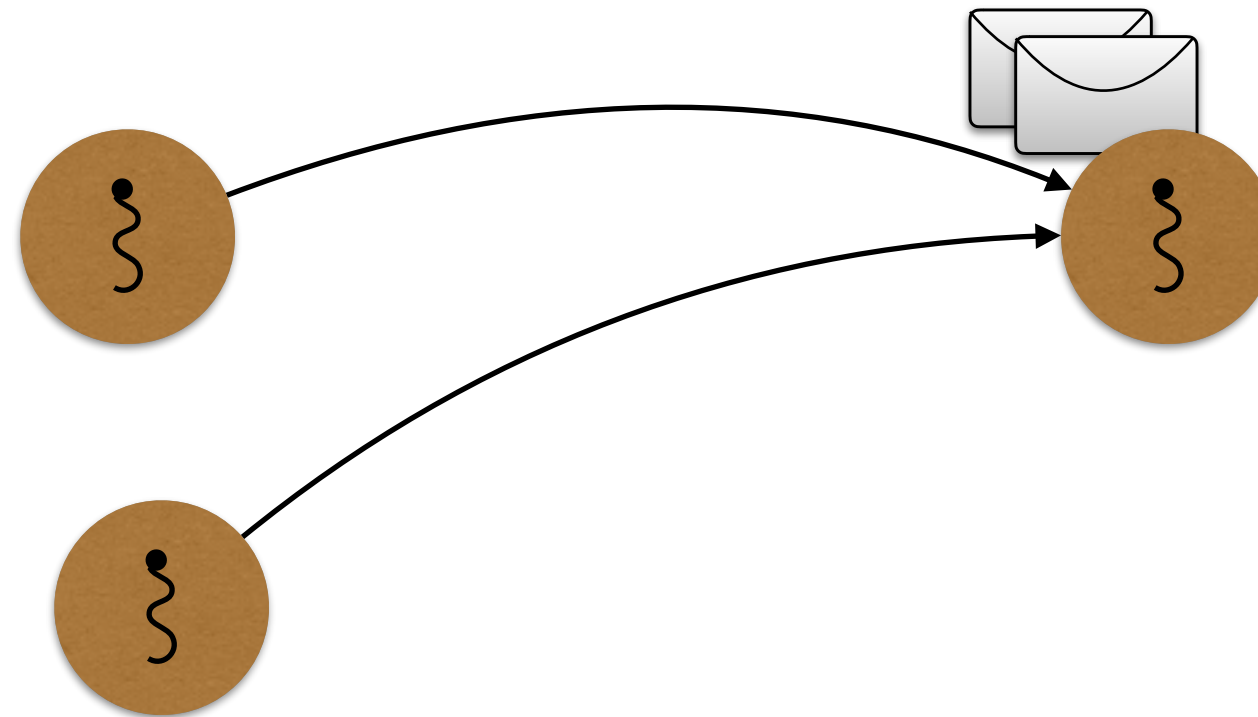
PLACES'17, Uppsala



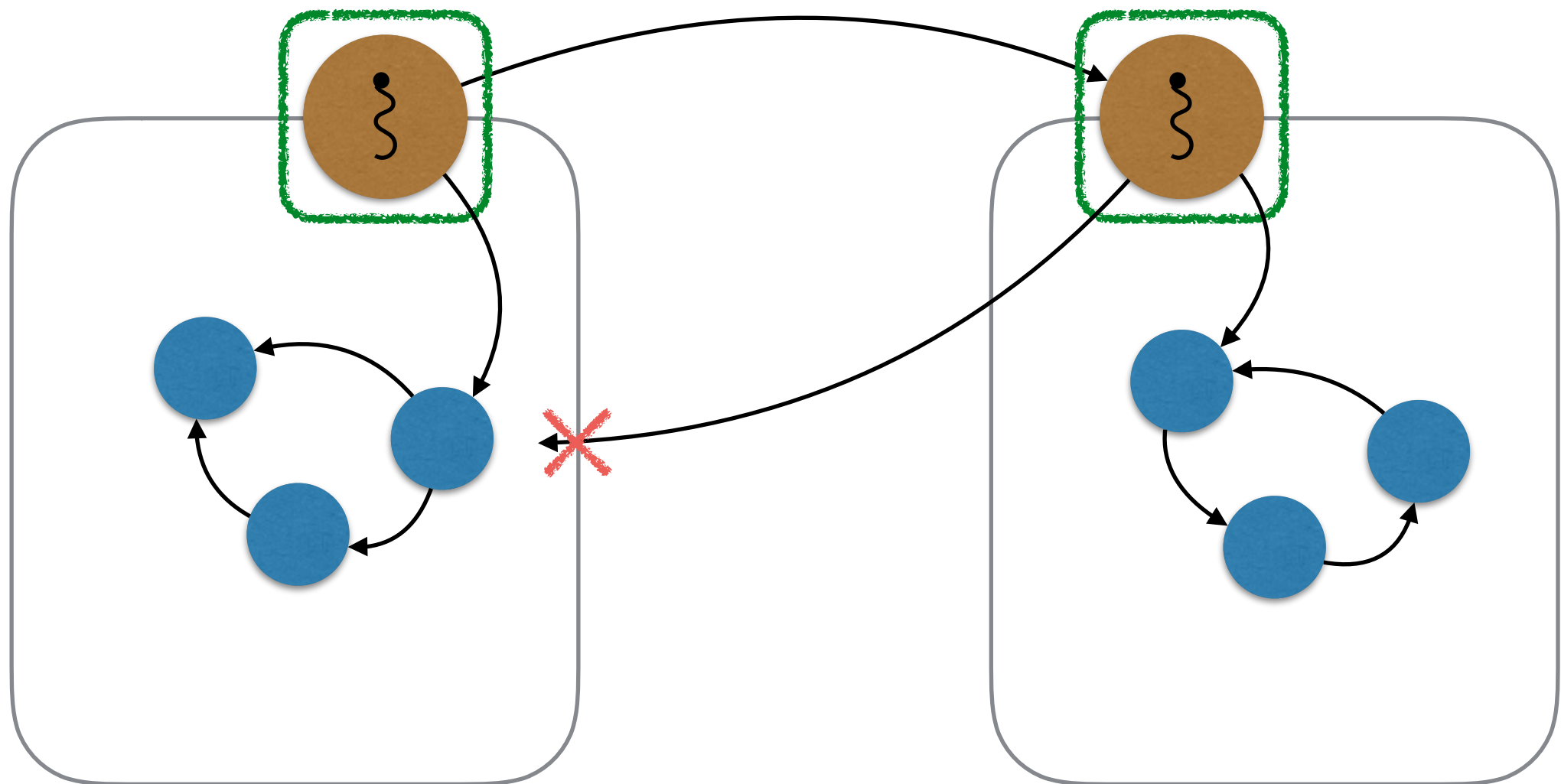
UP/MARC



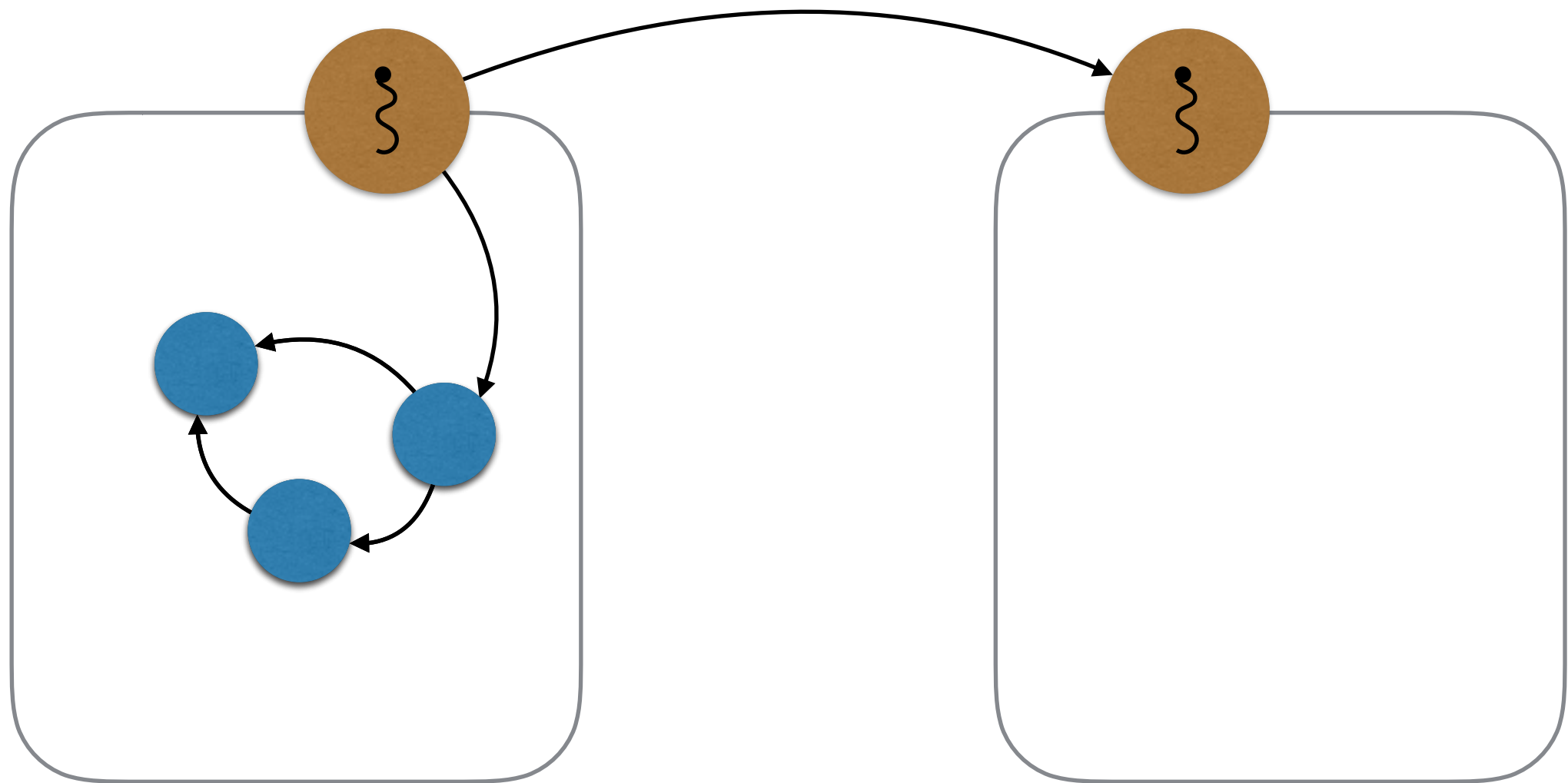
The Actor Model



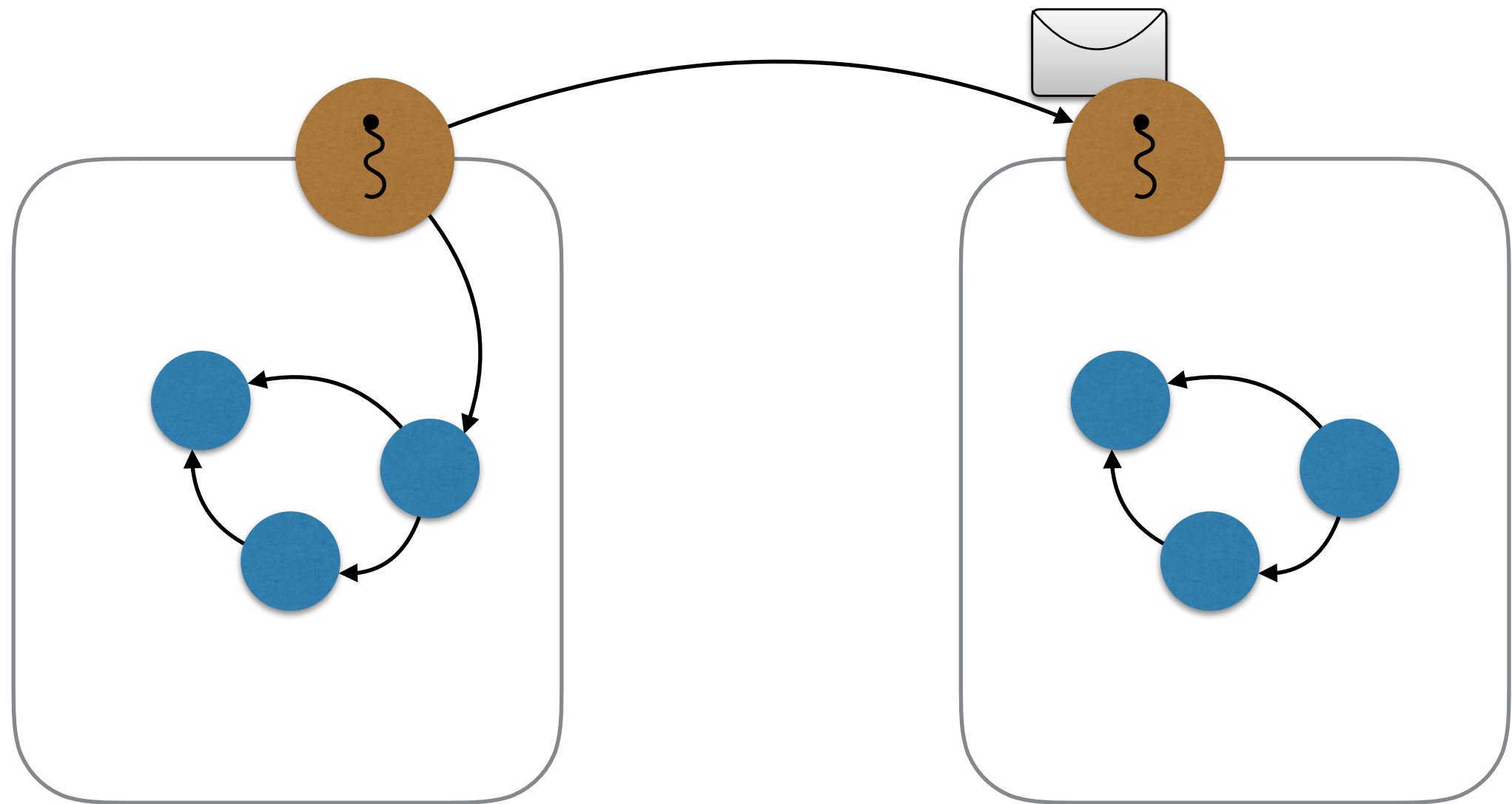
Actor Isolation allows Sequential Reasoning



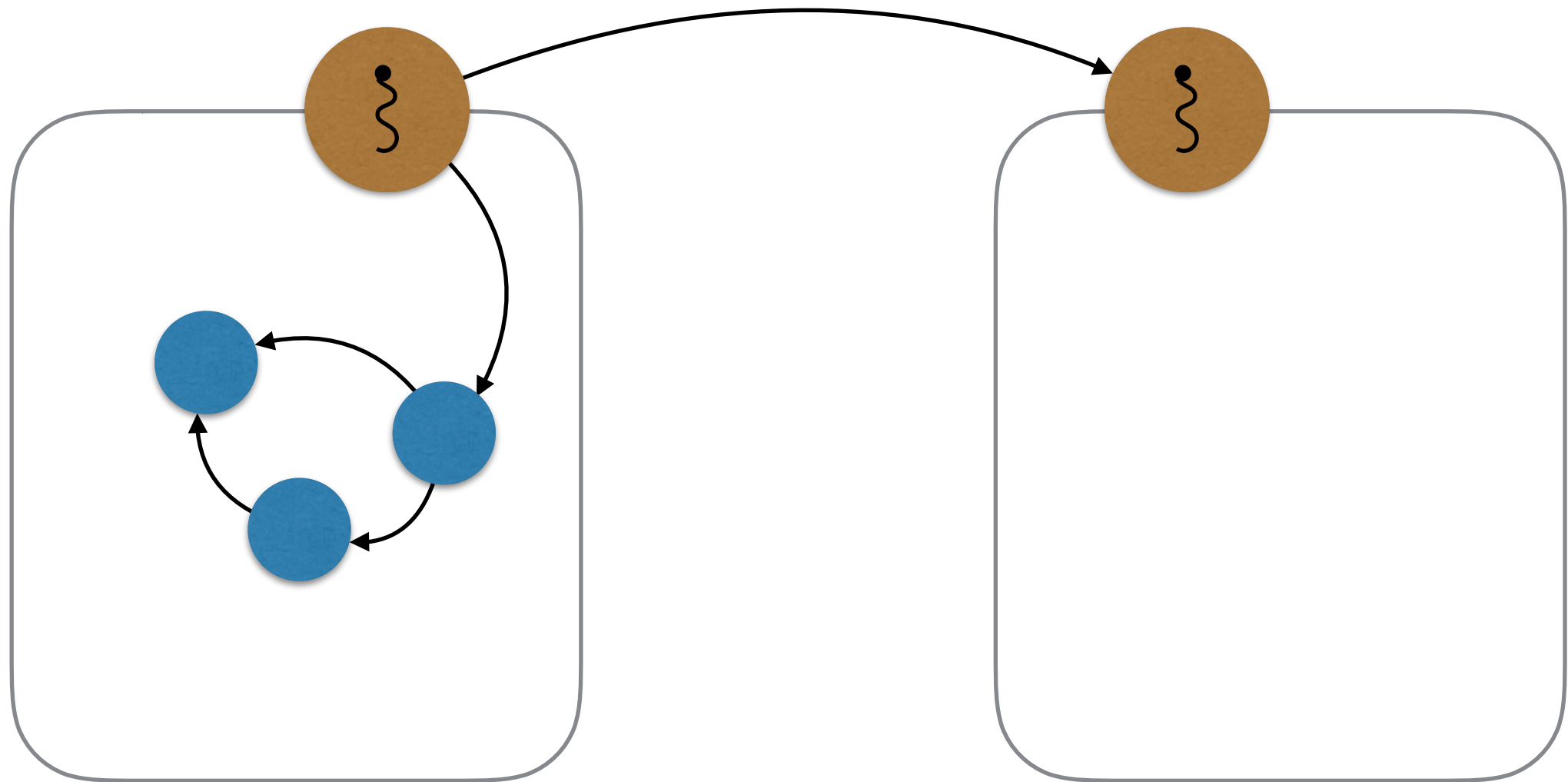
Passing Data Between Actors (by copy)



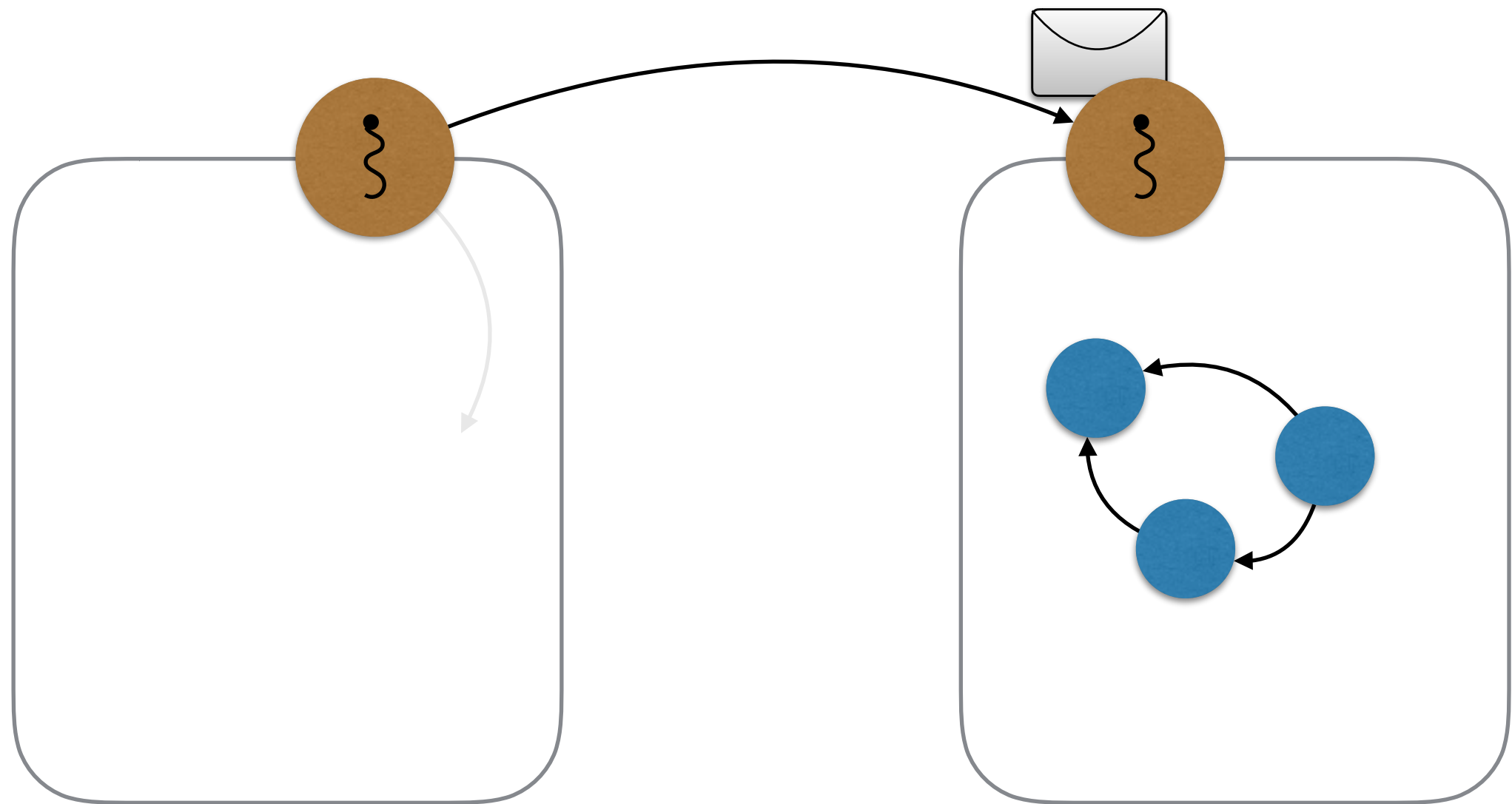
Passing Data Between Actors (by copy)



Passing Data Between Actors (by transfer)

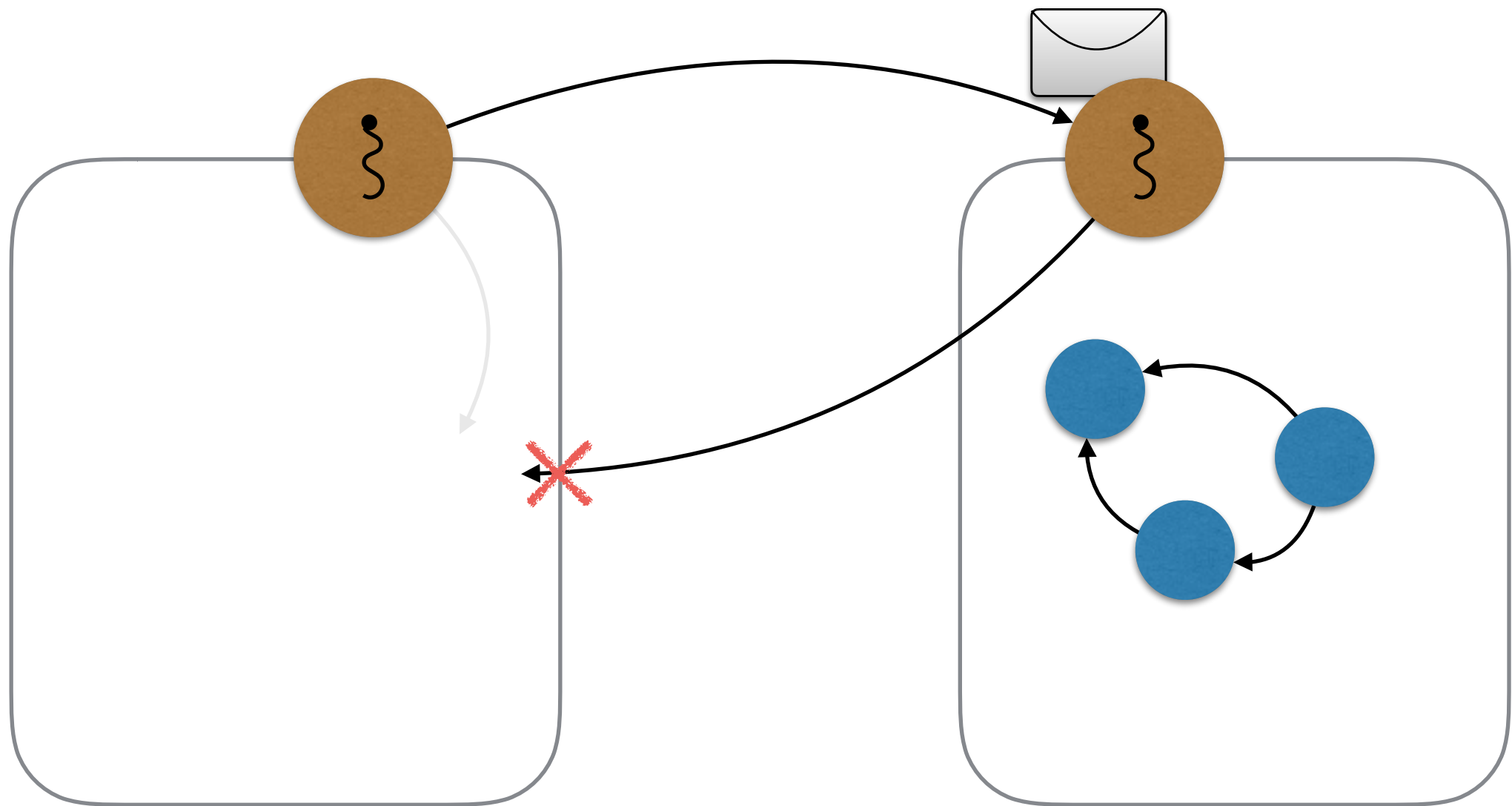


Passing Data Between Actors (by transfer)

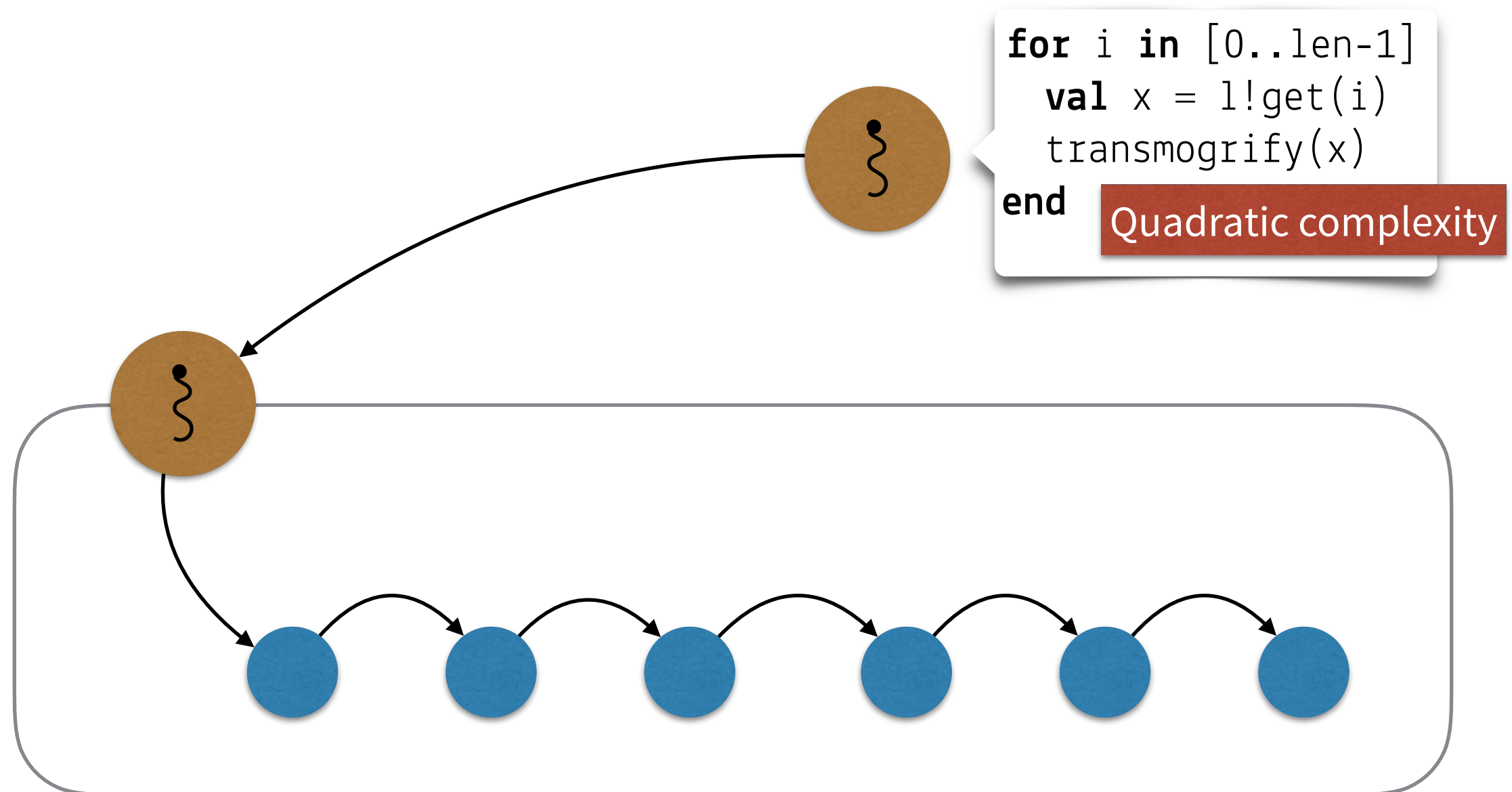


Passing Data Between Actors (by transfer)

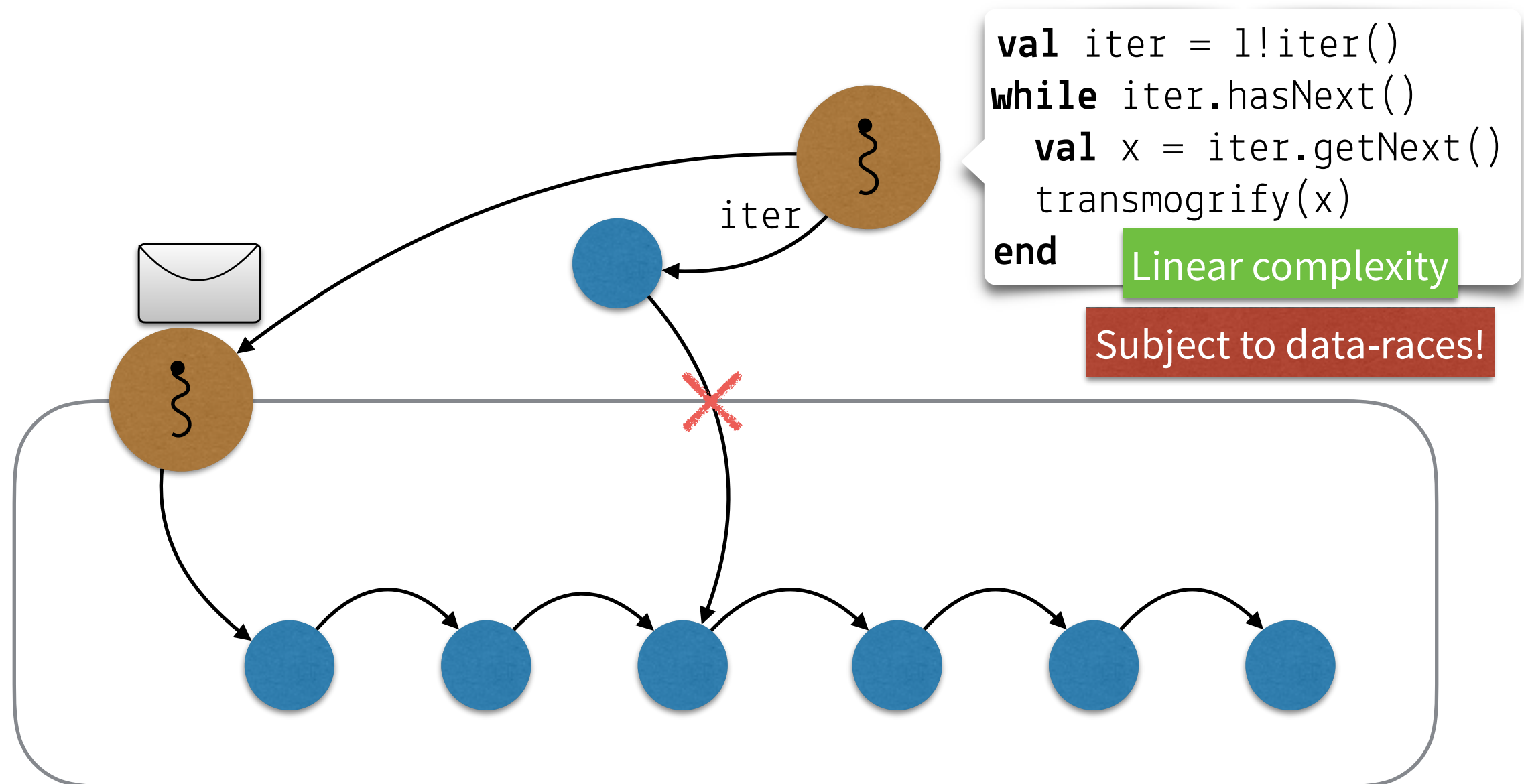
Strong Encapsulation



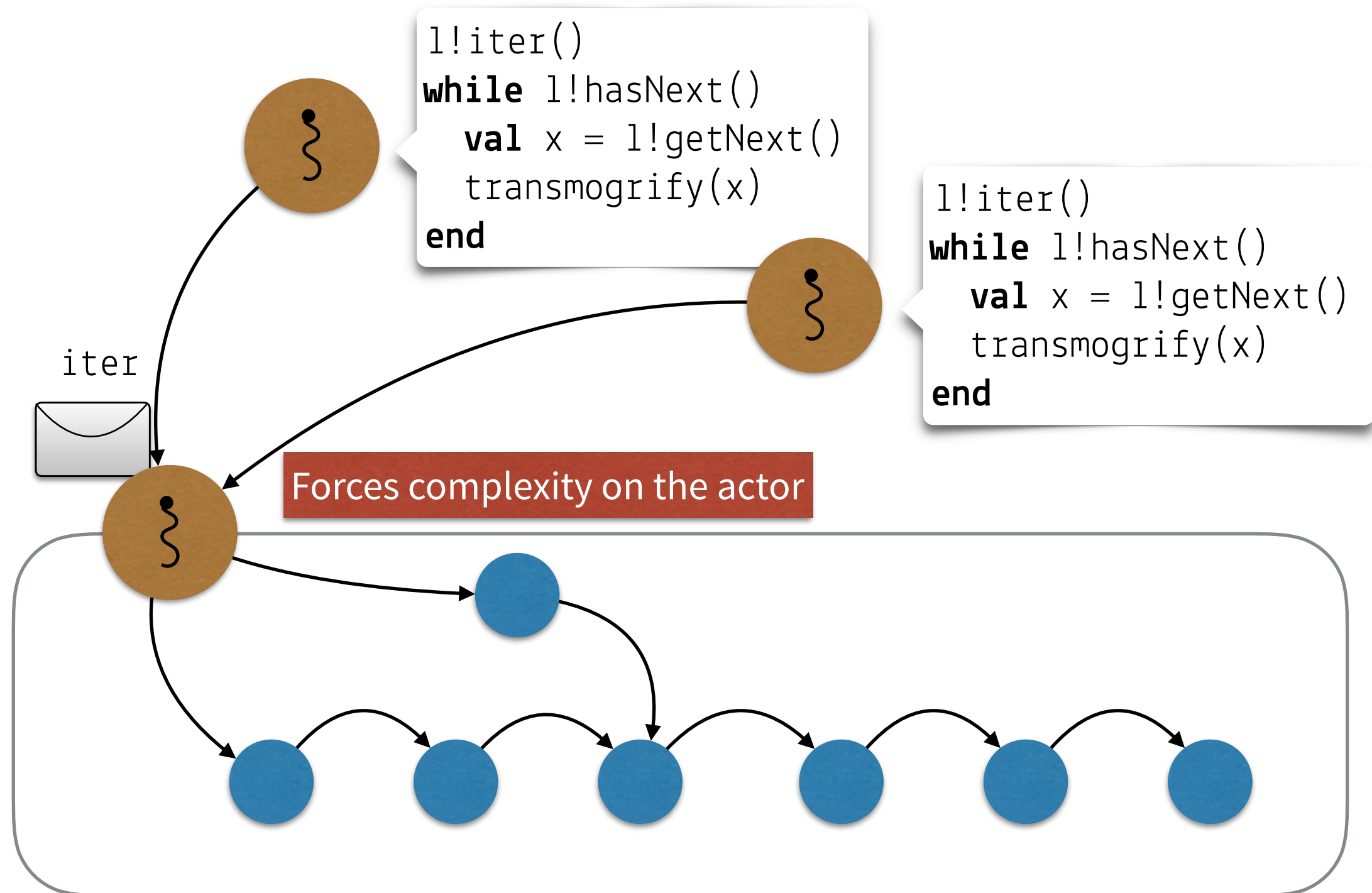
(Too) Strong Encapsulation



Solution? Breaking Actor Isolation

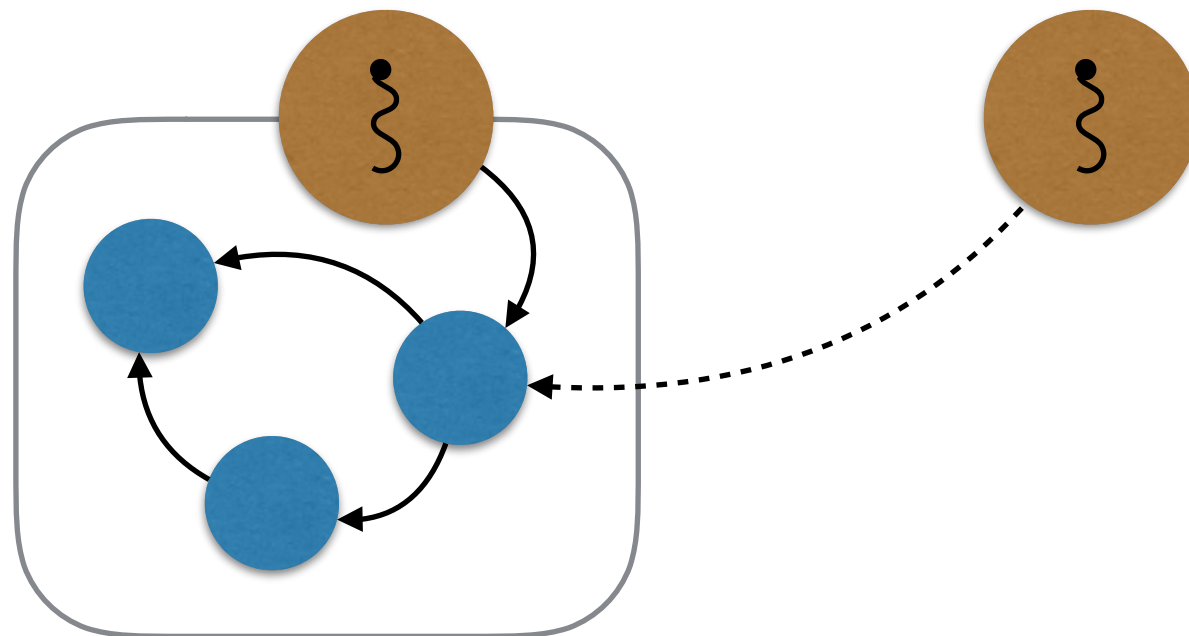


Solution? Keeping Complexity Isolated

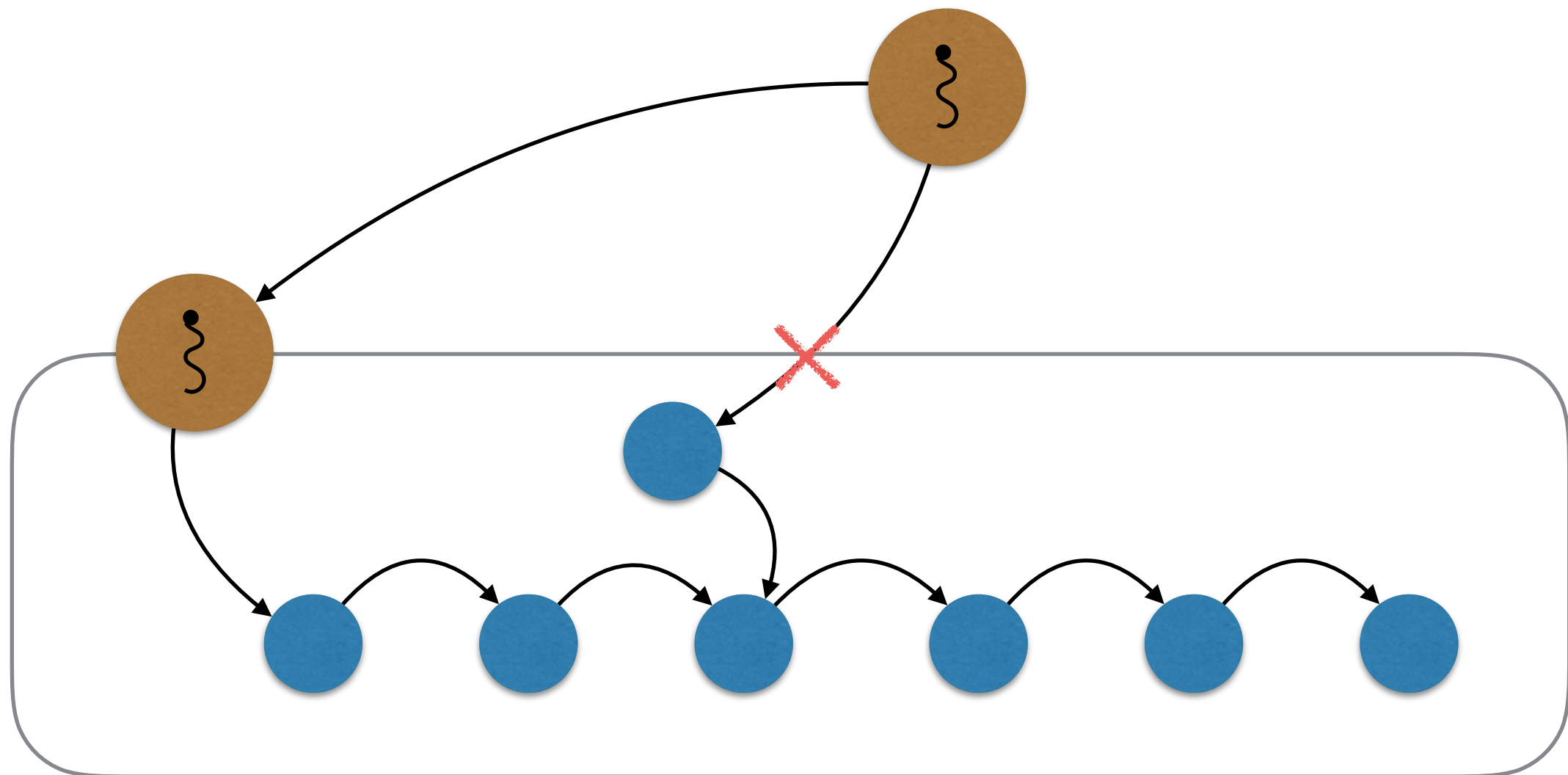


Problem Overview

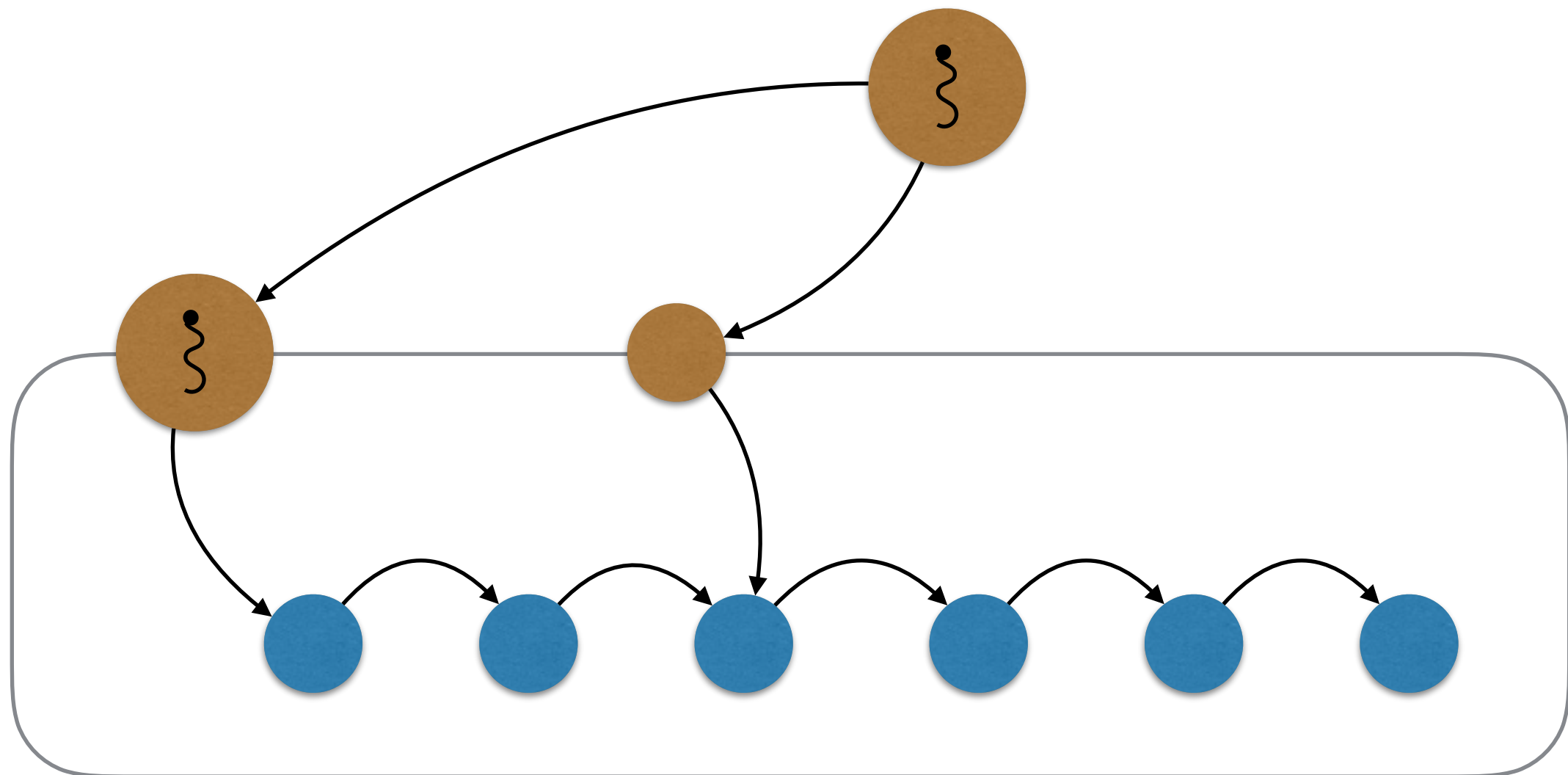
- Actors rely on isolation for sequential reasoning
 - Too strong for certain patterns
- Breaking actor isolation prevents sequential reasoning
- Extending actors to handle these patterns make them complex
- **Observation:**
Holding a reference to an isolated object is OK as long as only the owner accesses it



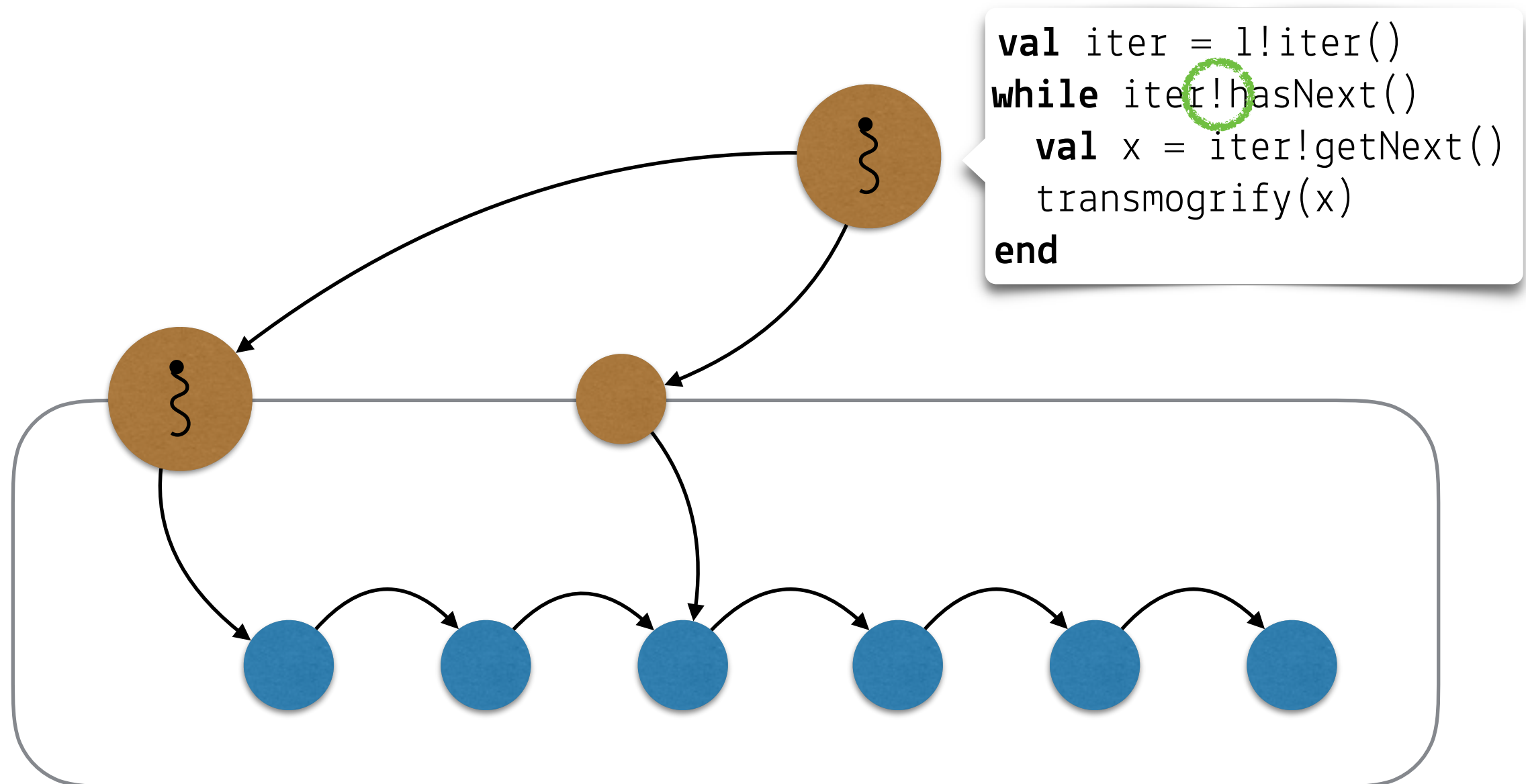
Relax Actor Isolation by Bestowing Activity



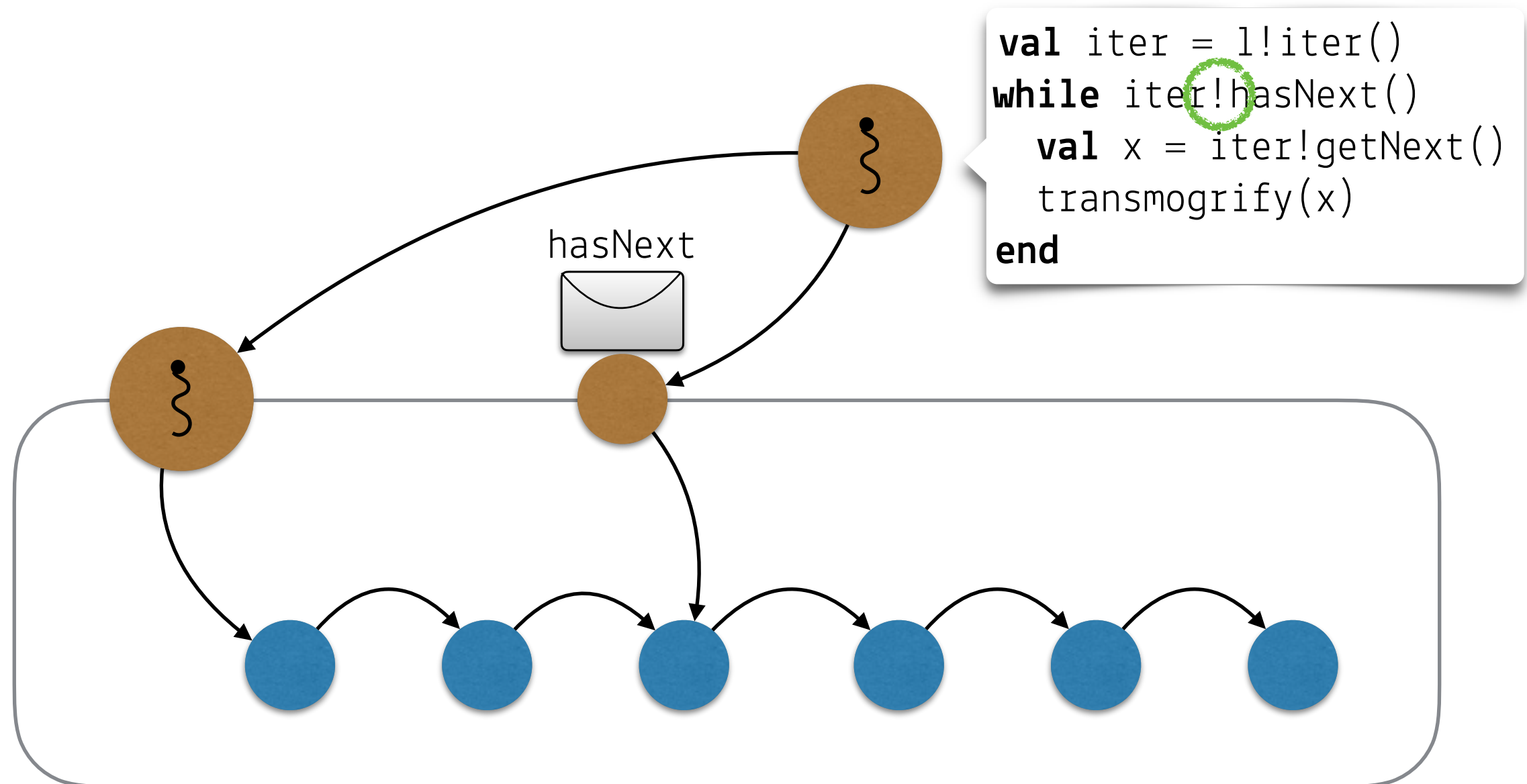
Relax Actor Isolation by Bestowing Activity



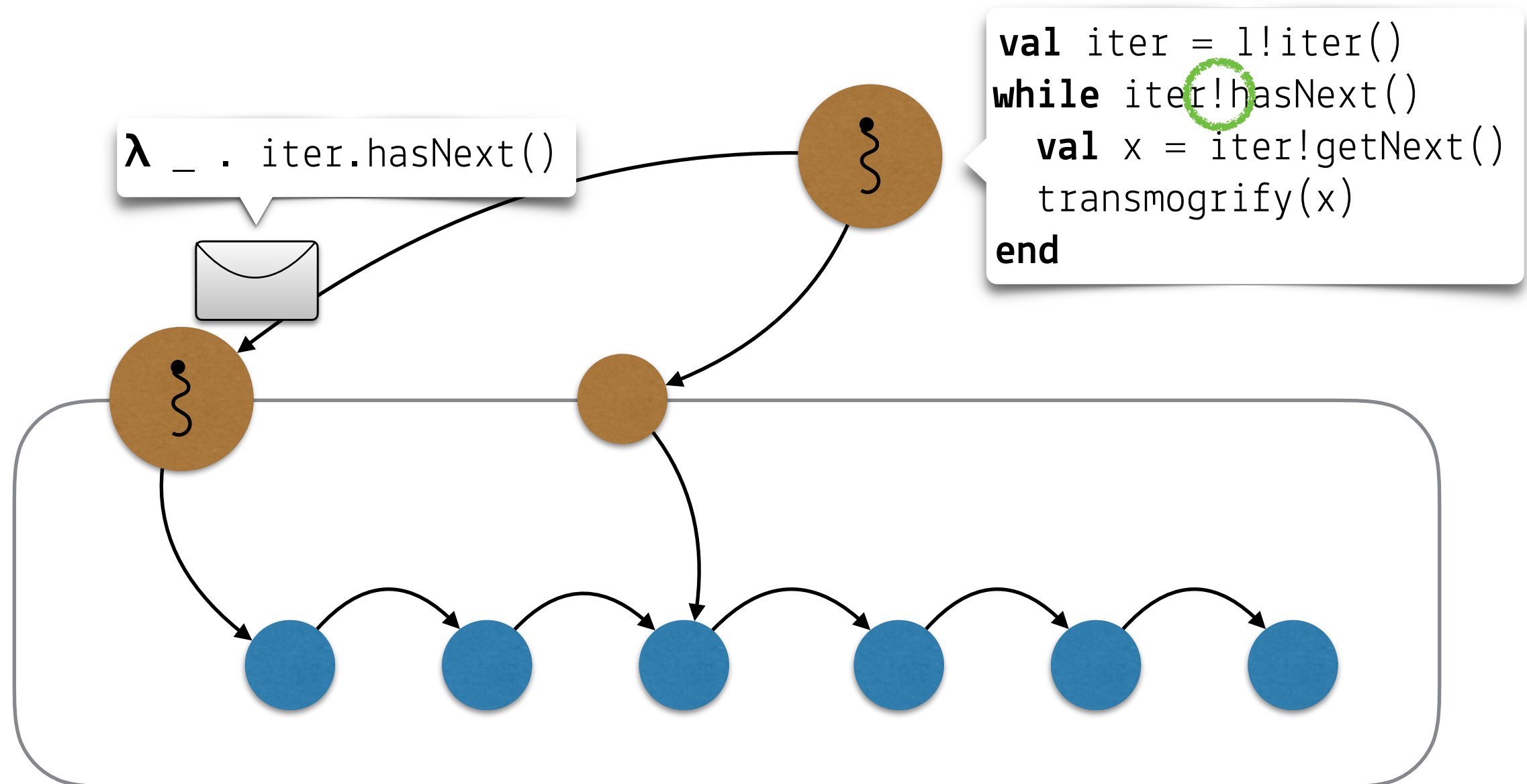
Relax Actor Isolation by Bestowing Activity



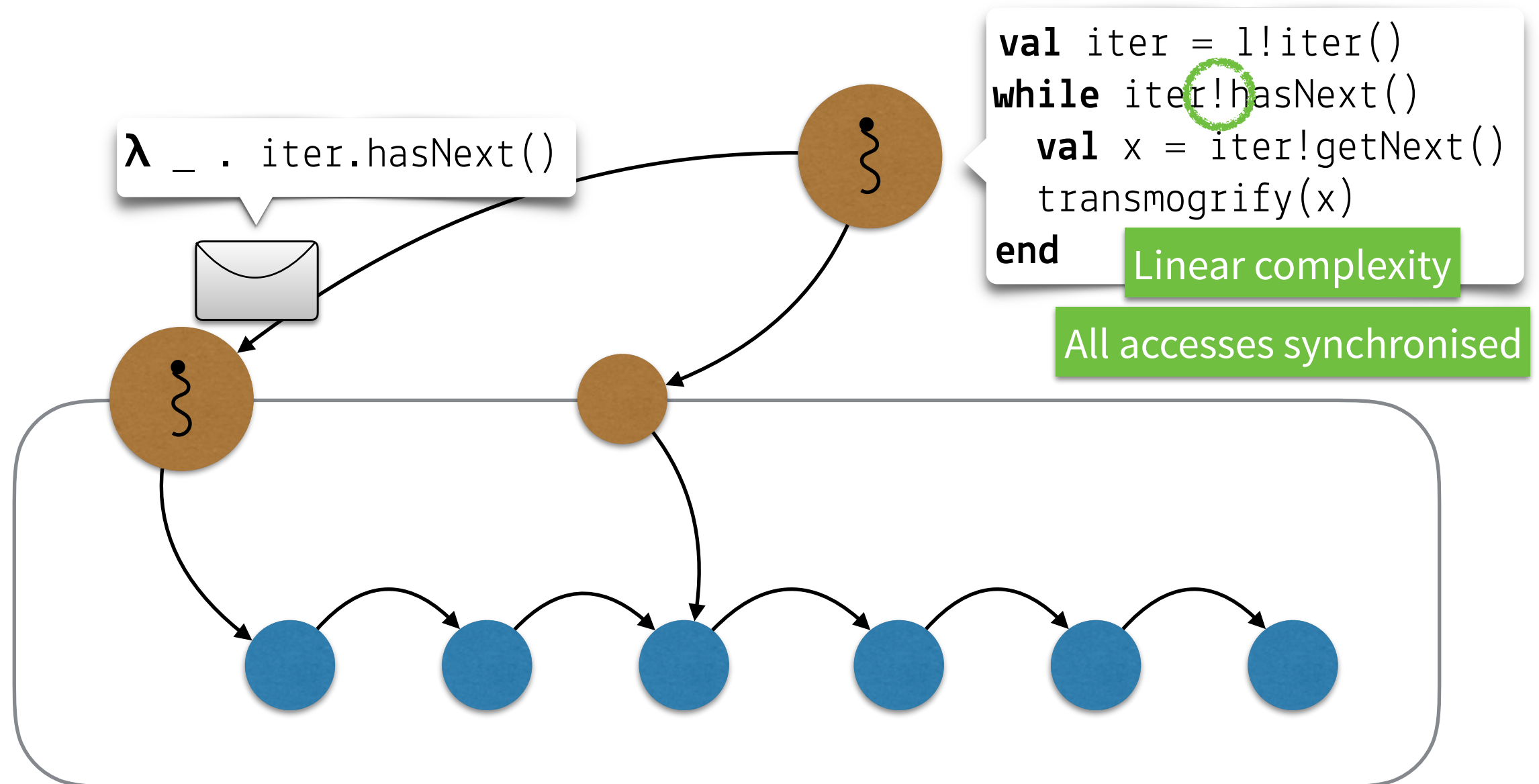
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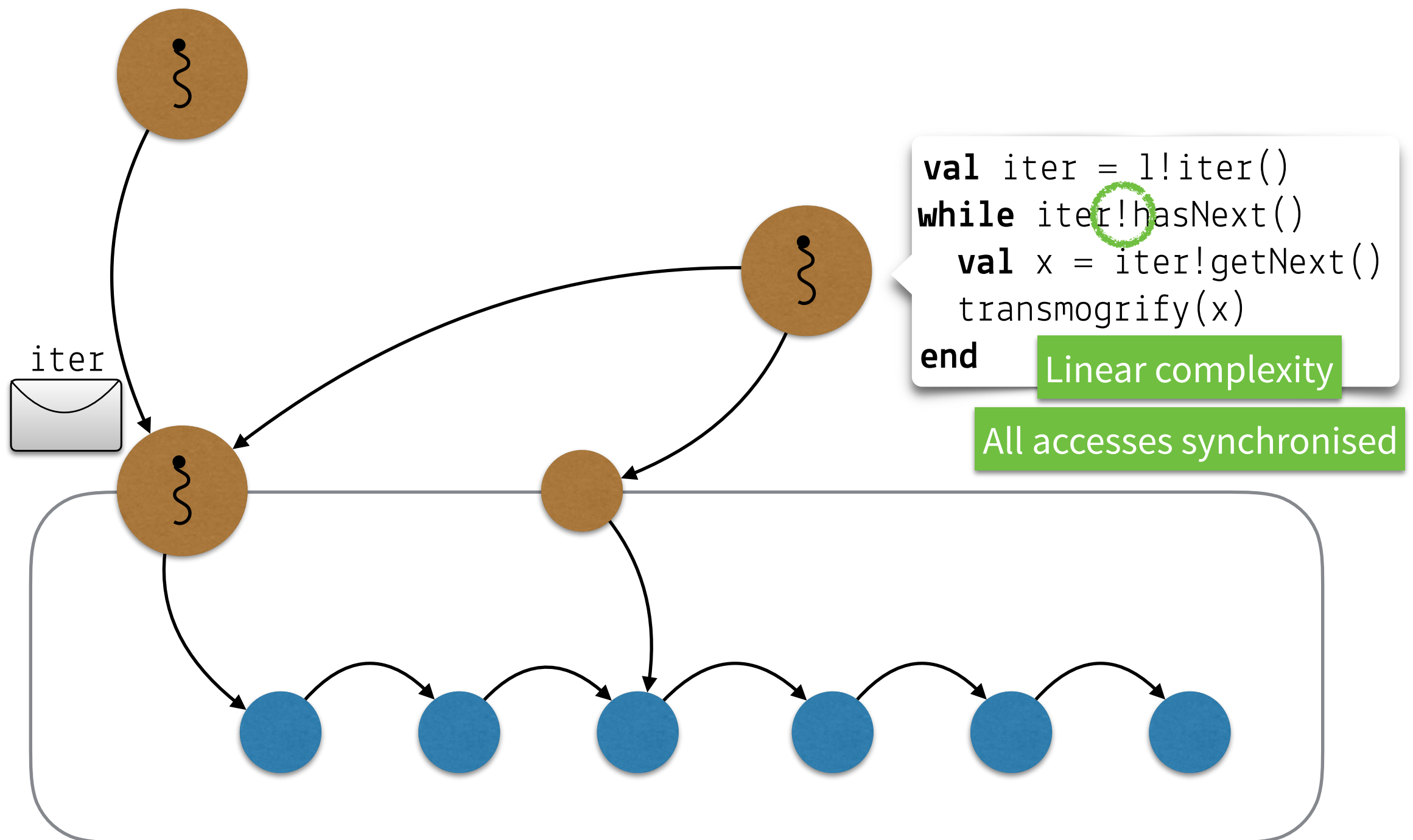
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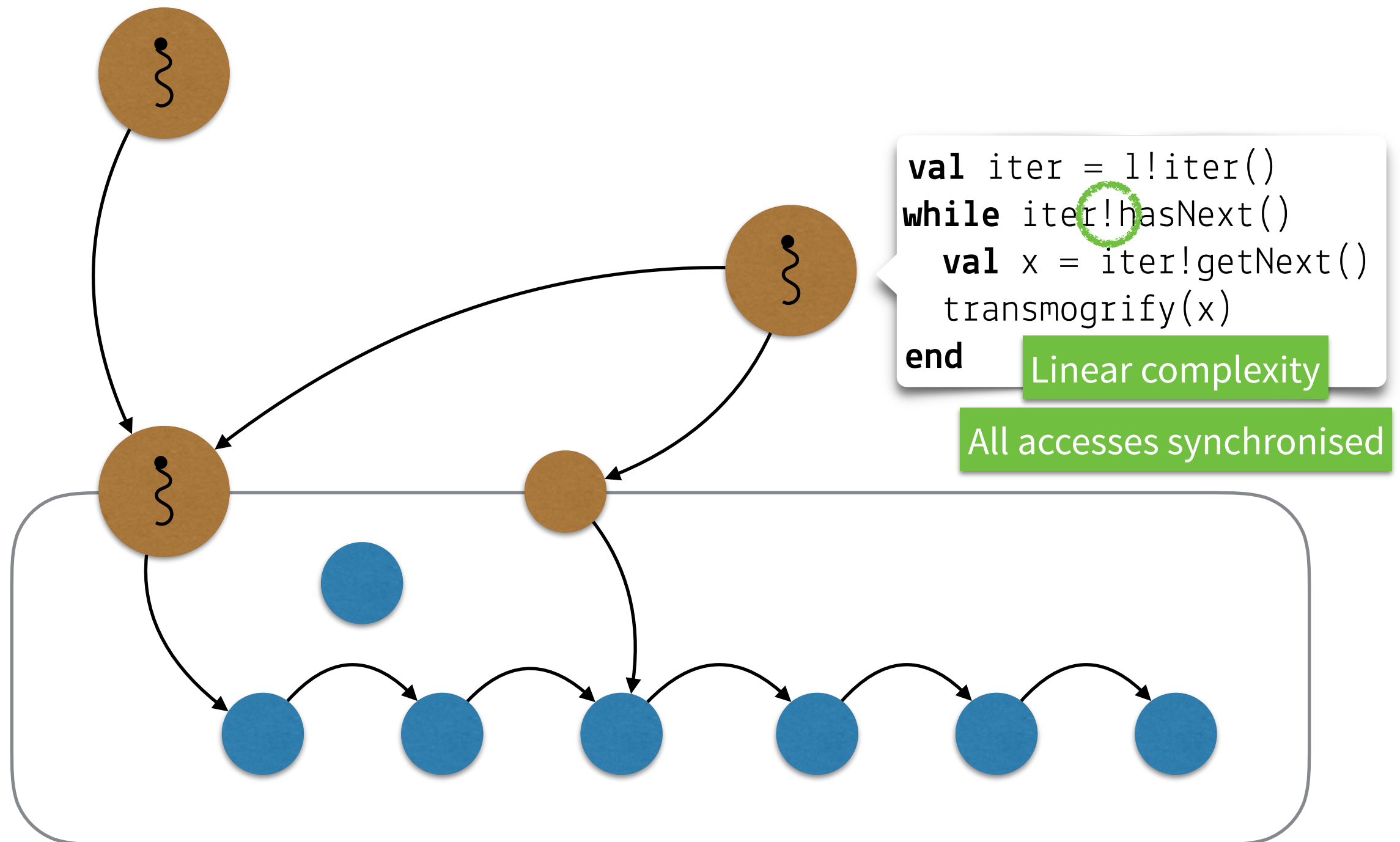
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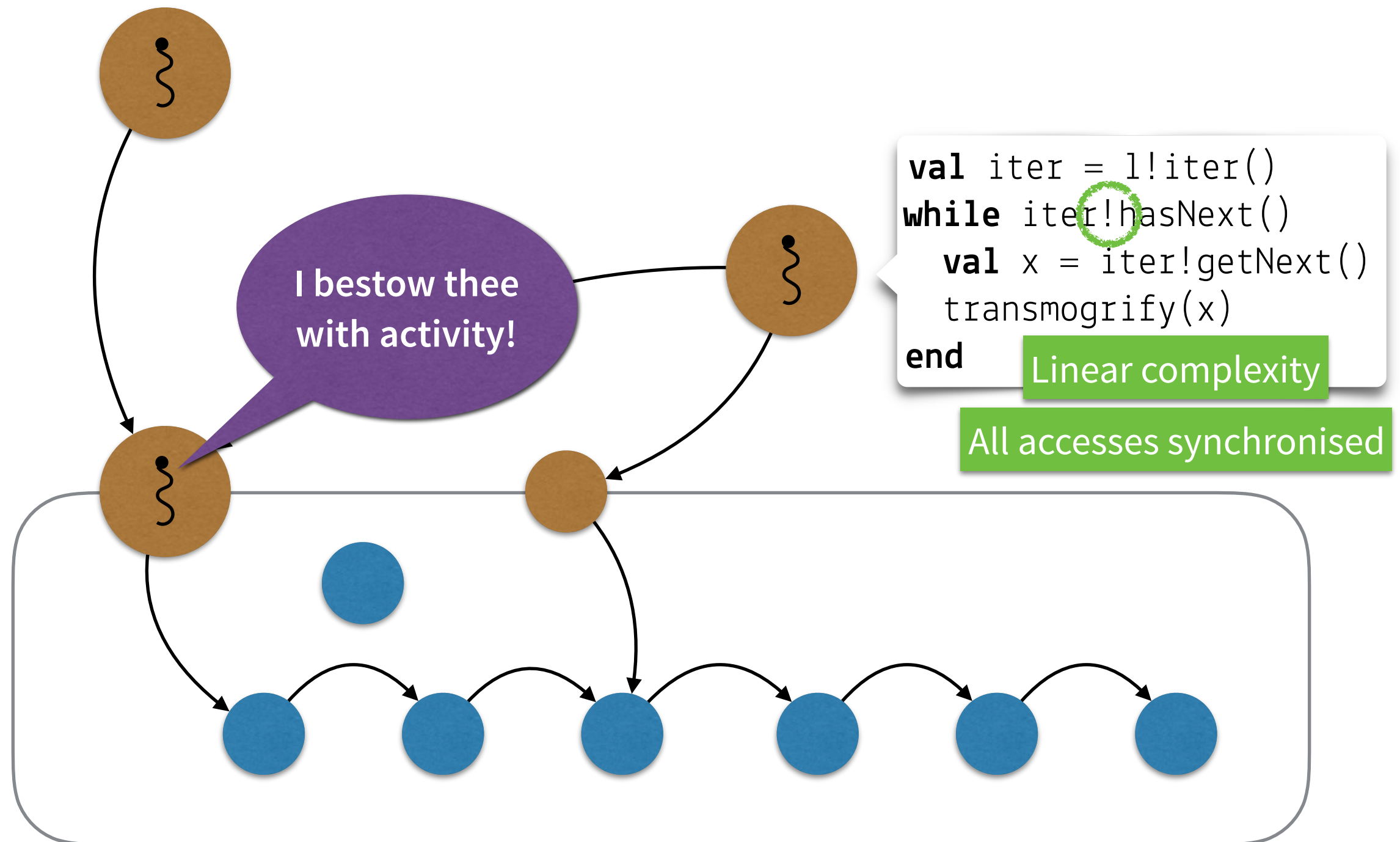
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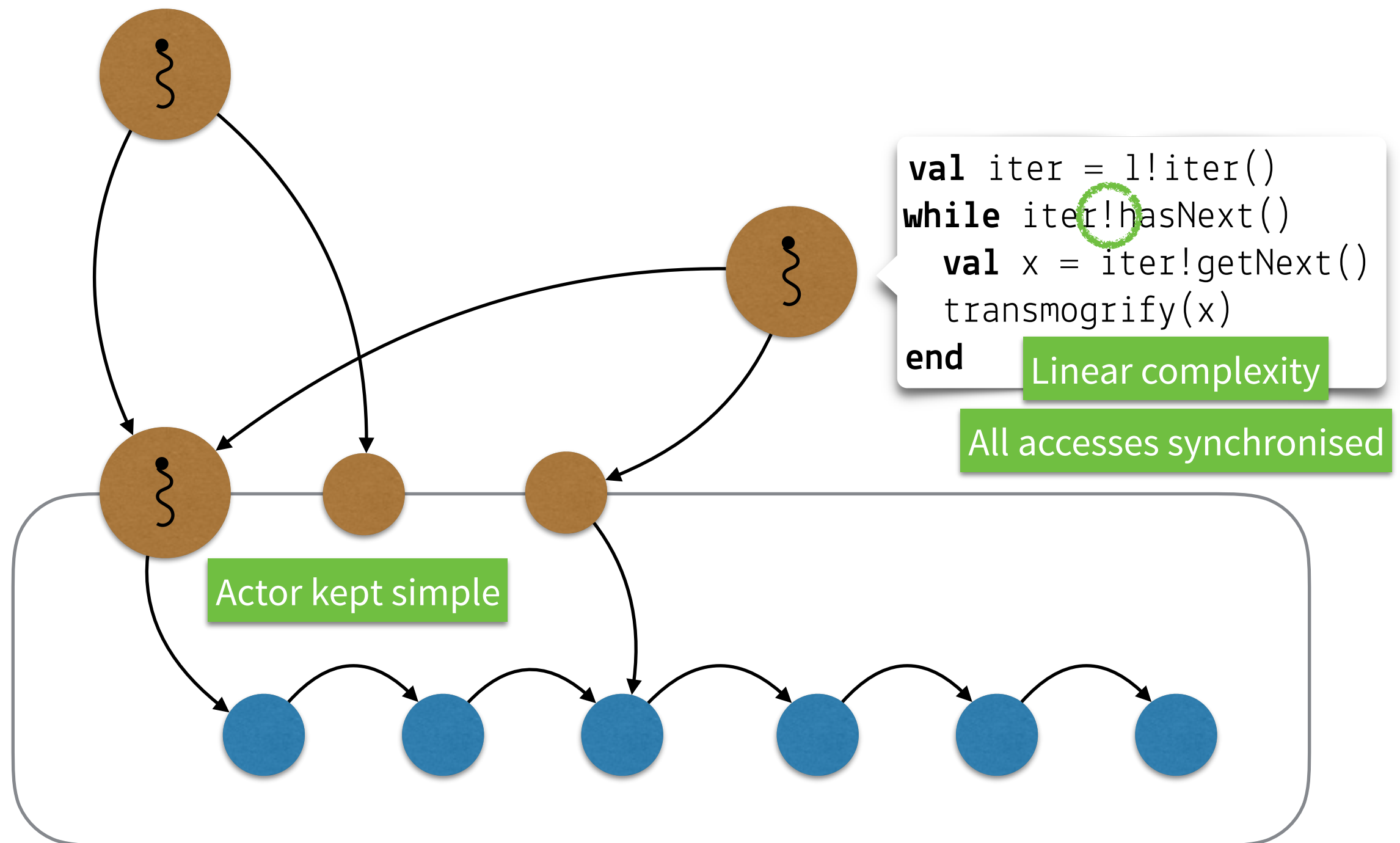
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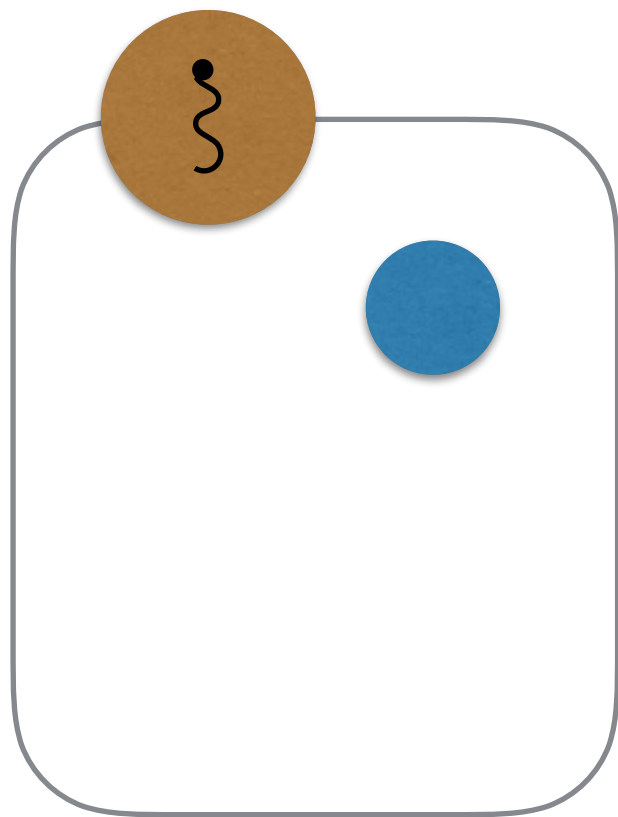
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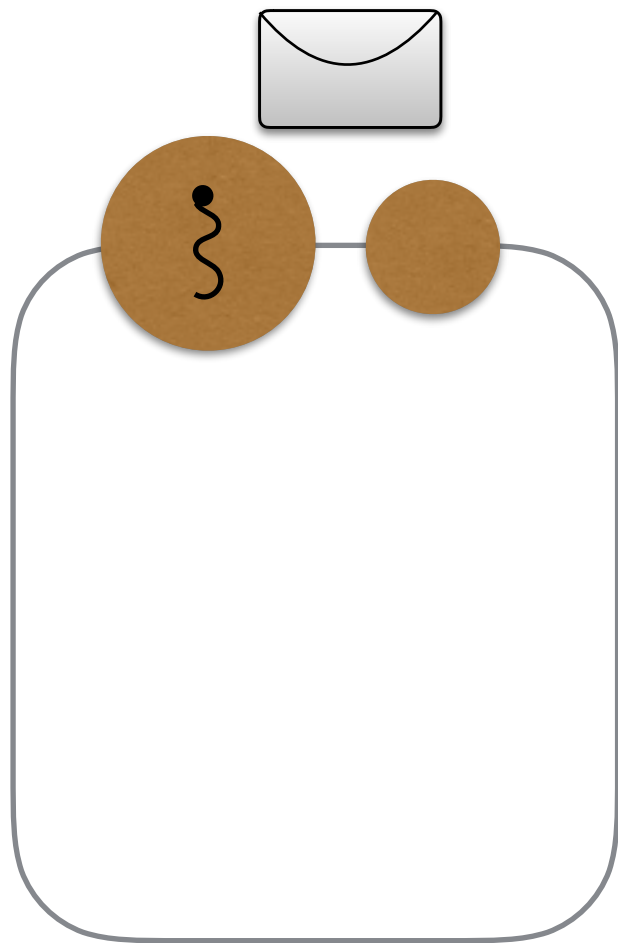
Delegating Parts of an Actor's Interface



```
actor A
  var c : C
  def foo() : unit
    ...
  end
  def bar() : unit
    ...
  end
  def getC : bestowed[C]
    bestow this.c
  end
end
```

```
class C
  def beep() : unit
    ...
  end
  def boop() : unit
    ...
  end
end
```

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Formalising Bestowed References

- Syntax

$$\begin{array}{ll} e & ::= x \mid e \ e \mid e!v \mid e.\text{mutate}() \mid \mathbf{new} \ \tau \mid \mathbf{bestow} \ e \mid v \\ v & ::= \lambda x : \tau. e \mid () \mid id \mid \iota \mid \iota_{id} \end{array} \qquad \begin{array}{ll} \tau & ::= \alpha \mid \mathbf{p} \mid \tau \rightarrow \tau \mid \text{Unit} \\ \alpha & ::= \mathbf{c} \mid \mathbf{B}(\mathbf{p}) \end{array}$$

Formalising Bestowed References

- Syntax

Message send

Bestowing

Passive type

$$e ::= x \mid e e \mid e!v \mid e.\text{mutate}() \mid \text{new } \tau \mid \text{bestow } e \mid v$$

$$v ::= \lambda x:\tau.e \mid () \mid id \mid \iota \mid \iota_{id}$$

$$\tau ::= \alpha \mid \mathbf{p} \mid \tau \rightarrow \tau \mid \text{Unit}$$

$$\alpha ::= \mathbf{c} \mid \mathbf{B}(\mathbf{p})$$

Bestowed type

- Static semantics

$\Gamma \vdash e : \tau$

Bestowed value

Actor type

(Expressions)

E-VAR
 $\frac{\Gamma(x) = \tau}{\Gamma \vdash x : \tau}$

E-APPLY
 $\frac{\Gamma \vdash e : \tau' \rightarrow \tau \quad \Gamma \vdash e' : \tau'}{\Gamma \vdash e e' : \tau}$

E-NEW-PASSIVE
 $\frac{}{\Gamma \vdash \text{new } \mathbf{p} : \mathbf{p}}$

E-NEW-ACTOR
 $\frac{}{\Gamma \vdash \text{new } \mathbf{c} : \mathbf{c}}$

E-MUTATE
 $\frac{\Gamma \vdash e : \mathbf{p}}{\Gamma \vdash e.\text{mutate}() : \text{Unit}}$

E-BESTOW
 $\frac{\Gamma \vdash e : \mathbf{p}}{\Gamma \vdash \text{bestow } e : \mathbf{B}(\mathbf{p})}$

E-SEND
 $\frac{\Gamma \vdash e : \alpha \quad \Gamma_{\alpha, x : \mathbf{p}} \vdash e' : \tau' \quad \nexists \iota . \iota \in e'}{\Gamma \vdash e! \lambda x : \mathbf{p}. e' : \text{Unit}}$

E-FN
 $\frac{\Gamma, x : \tau \vdash e : \tau'}{\Gamma \vdash (\lambda x : \tau. e) : \tau \rightarrow \tau'}$

E-UNIT
 $\frac{}{\Gamma \vdash () : \text{Unit}}$

E-LOC
 $\frac{}{\Gamma \vdash \iota : \mathbf{p}}$

E-ID
 $\frac{}{\Gamma \vdash id : \mathbf{c}}$

E-BESTOWED
 $\frac{}{\Gamma \vdash \iota_{id} : \mathbf{B}(\mathbf{p})}$

Formalising Bestowed References

- Syntax

Message send

Bestowing

Passive type

$$\begin{aligned} e &::= x \mid e e \mid e!v \mid e.\text{mutate}() \mid \text{new } \tau \mid \text{bestow } e \mid v \\ v &::= \lambda x:\tau.e \mid () \mid id \mid \iota \mid \iota_{id} \end{aligned}$$

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E-ID
 $\frac{}{\Gamma \vdash id : \mathbf{c}}$

E-BESTOWED
 $\frac{}{\Gamma \vdash \iota_{id} : \mathbf{B}(\mathbf{p})}$

Formalising Bestowed References

- Dynamic semantics

$$\boxed{id \vdash \langle H, e \rangle \hookrightarrow \langle H', e' \rangle}$$

(Evaluation of expressions)

EVAL-SEND-ACTOR

$$\frac{\begin{array}{l} H(id') = (\iota, L, Q, e) \\ H' = H[id' \mapsto (\iota, L, v \ Q, e)] \end{array}}{id \vdash \langle H, id' ! v \rangle \hookrightarrow \langle H', () \rangle}$$

EVAL-SEND-BESTOWED

$$\frac{\begin{array}{l} H(id') = (\iota', L, Q, e) \\ H' = H[id' \mapsto (\iota', L, (\lambda x : \mathbf{p}. v \ \iota) \ Q, e)] \end{array}}{id \vdash \langle H, \iota_{id'} ! v \rangle \hookrightarrow \langle H', () \rangle}$$

EVAL-APPLY

$$\frac{e' = e[x \mapsto v]}{id \vdash \langle H, (\lambda x : \tau. e) \ v \rangle \hookrightarrow \langle H, e' \rangle}$$

EVAL-MUTATE

$$id \vdash \langle H, \iota. \text{mutate}() \rangle \hookrightarrow \langle H, () \rangle$$

EVAL-BESTOW

$$id \vdash \langle H, \text{bestow } \iota \rangle \hookrightarrow \langle H, \iota_{id} \rangle$$

EVAL-NEW-PASSIVE

$$\frac{\begin{array}{l} H(id) = (\iota, L, Q, e) \quad \iota' \text{ fresh} \\ H' = H[id \mapsto (\iota, L \cup \{\iota'\}, Q, e)] \end{array}}{id \vdash \langle H, \text{new p} \rangle \hookrightarrow \langle H', \iota' \rangle}$$

EVAL-NEW-ACTOR

$$\frac{\begin{array}{l} id' \text{ fresh} \quad \iota' \text{ fresh} \\ H' = H[id' \mapsto (\iota', \{\iota'\}, \epsilon, ())] \end{array}}{id \vdash \langle H, \text{new } \alpha \rangle \hookrightarrow \langle H', id' \rangle}$$

EVAL-CONTEXT

$$\frac{id \vdash \langle H, e \rangle \hookrightarrow \langle H', e' \rangle}{id \vdash \langle H, E[e] \rangle \hookrightarrow \langle H', E[e'] \rangle}$$

Formalising Bestowed References

- Dynamic semantics

$id \vdash \langle H, e \rangle \hookrightarrow \langle H', e' \rangle$

(Evaluation of expressions)

EVAL-SEND-ACTOR

$$\frac{H(id') = (\iota, L, Q, e) \quad H' = H[id' \mapsto (\iota, L, v \ Q, e)]}{id \vdash \langle H, id' ! v \rangle \hookrightarrow \langle H', () \rangle}$$

EVAL-SEND-BESTOWED

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EVAL-MUTATE

$$id \vdash \langle H, \iota. \text{mutate}() \rangle \hookrightarrow \langle H, () \rangle$$

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$$\frac{id' \text{ fresh} \quad \iota' \text{ fresh} \quad H' = H[id' \mapsto (\iota', \{\iota'\}, \epsilon, ())]}{id \vdash \langle H, \text{new } \alpha \rangle \hookrightarrow \langle H', id' \rangle}$$

EVAL-CONTEXT

$$\frac{id \vdash \langle H, e \rangle \hookrightarrow \langle H', e' \rangle}{id \vdash \langle H, E[e] \rangle \hookrightarrow \langle H', E[e'] \rangle}$$

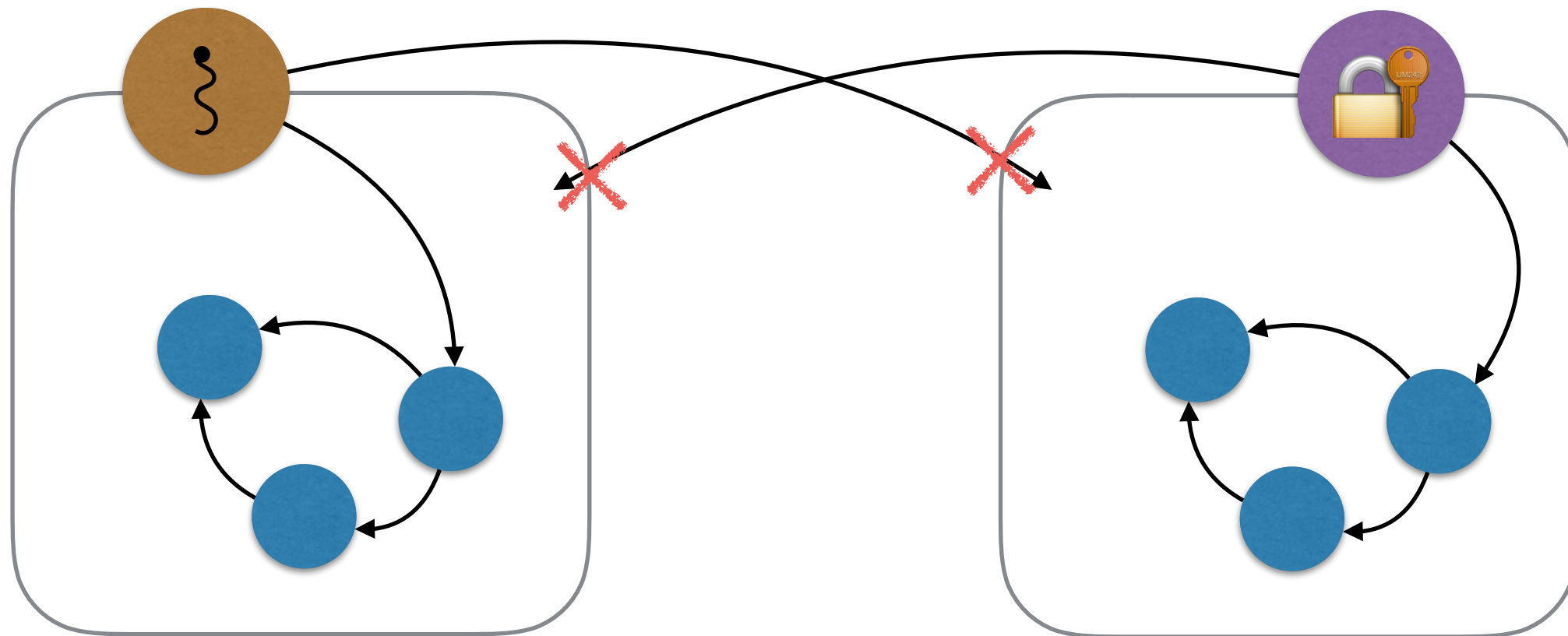
Properties of the Formalism

- **Progress:** $\vdash H \implies (\exists H' . H \hookrightarrow H') \vee (\forall id \in \mathbf{dom}(H) . H(id) = (\iota, L, \epsilon, v))$
- **Preservation:** $\vdash H \wedge H \hookrightarrow H' \implies \vdash H'$
- **Data-race freedom:** Two actors will never mutate the same passive object

$$\left(\begin{array}{c} id_1 \neq id_2 \\ \wedge H(id_1) = (\iota_1, L_1, Q_1, \iota.\text{mutate}()) \\ \wedge H(id_2) = (\iota_2, L_2, Q_2, \iota'.\text{mutate}()) \end{array} \right) \implies \iota \neq \iota'$$

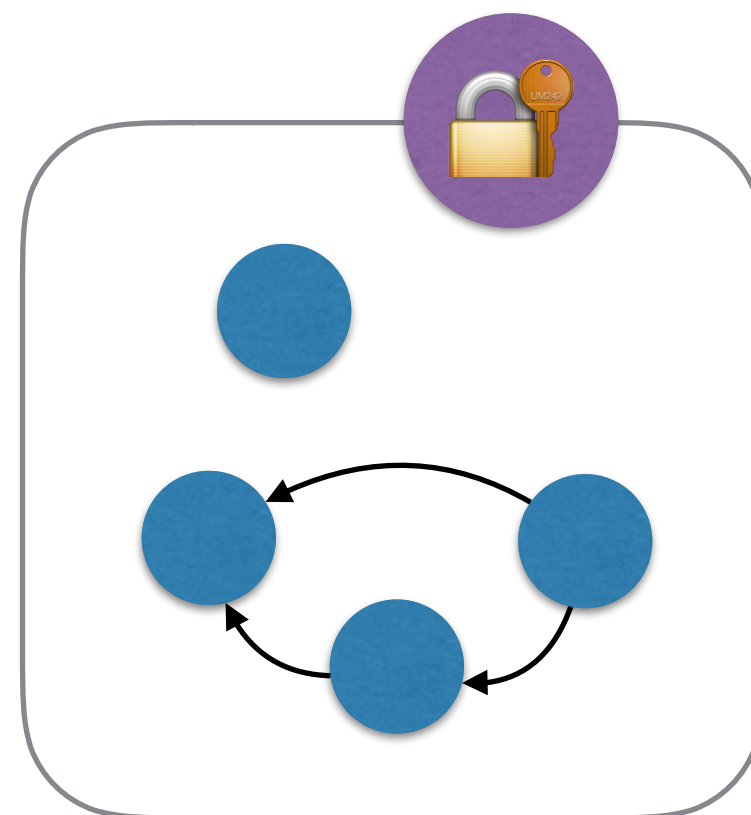
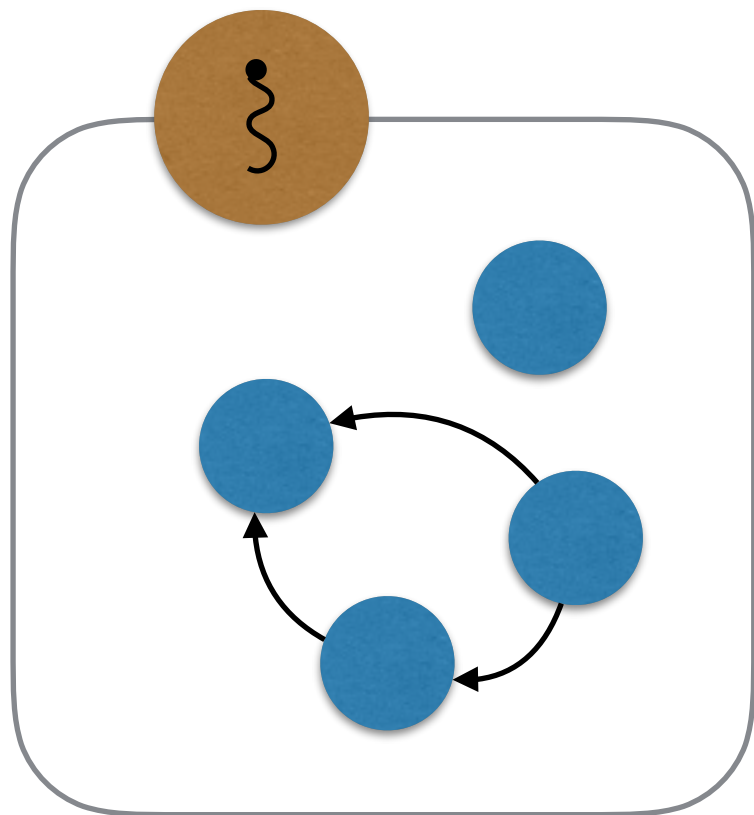
The Big Picture

- Kappa is a capability based type-system for concurrent OO-programming [ECOOP'16]
 - Tracks the boundaries of objects to achieve strong encapsulation



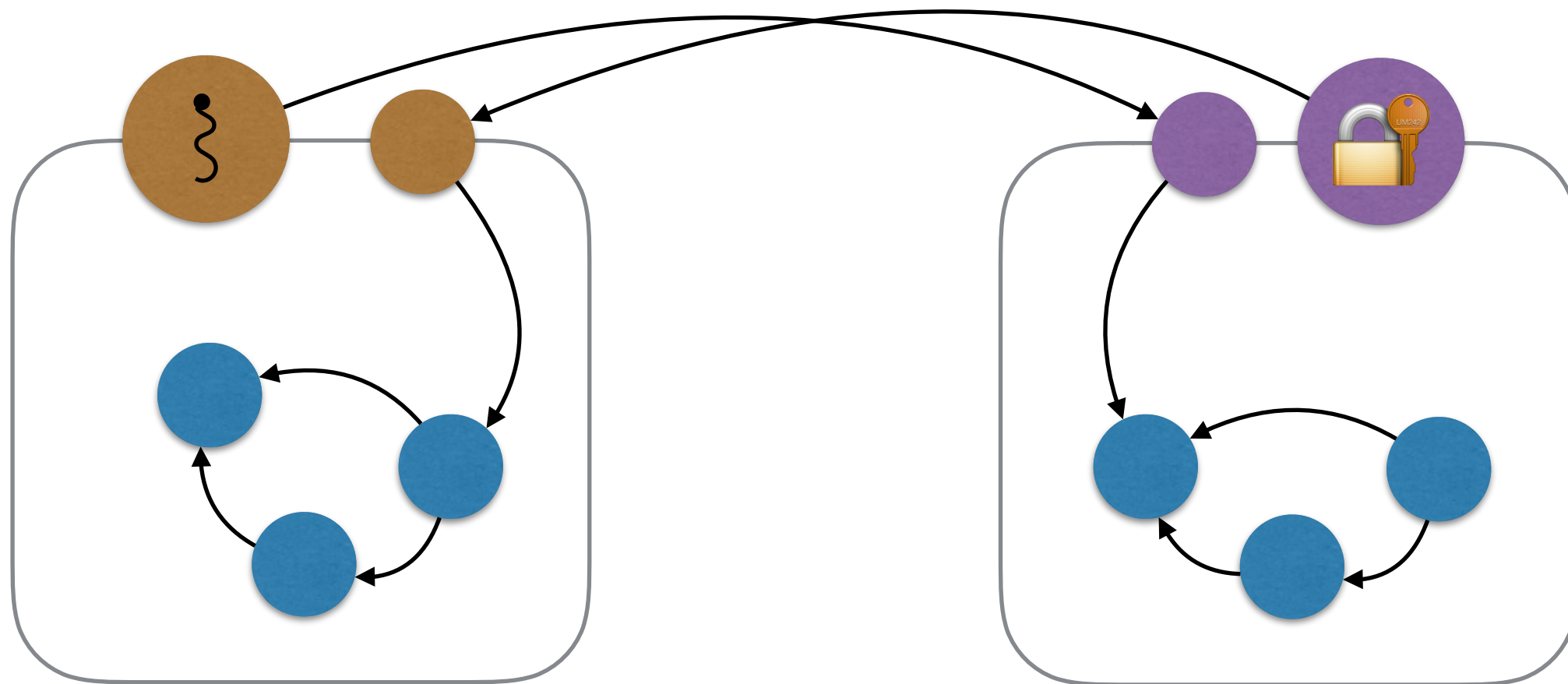
Generalizing Bestowed References

- Bestowed references provide a thread-safe way to break encapsulation



Generalizing Bestowed References

- Bestowed references provide a thread-safe way to break encapsulation



Future Work

- Implementation in Encore



```
active class Foo
  var box : IntBox
  def bestowBox() : Bestowed[IntBox]
    bestow this.box
  end
end

active class Main
  def main() : unit
    var b = new Foo()
    var box = get(b ! bestowBox())
    box ! inc()
  end
end
```

- Enriched formalism with copying and ownership transfer

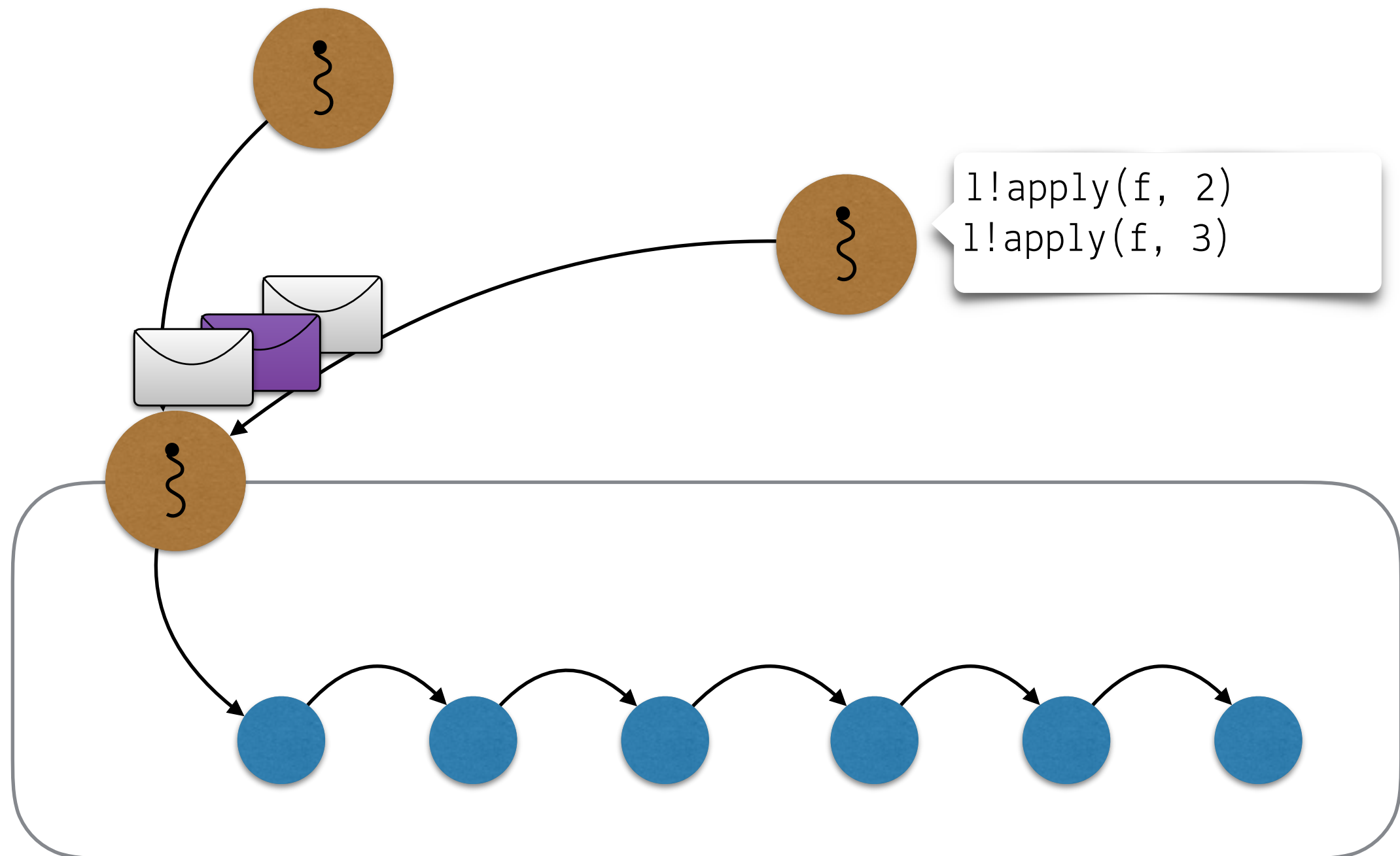
Summary

- Actor isolation can be relaxed by **bestowing** encapsulated objects with activity
 - All accesses will be synchronised via the message queue of the owning actor
 - Actors do not need to know the implementation of their bestowed objects
 - A bestowed object does not need to know that it is bestowed
- The same kind of relaxed encapsulation works for locks
 - All accesses will be synchronised via the lock of the owning object
- **Also in the paper:**
 - Atomic blocks to group operations
 - Implementation sketches
 - Polymorphic concurrency control

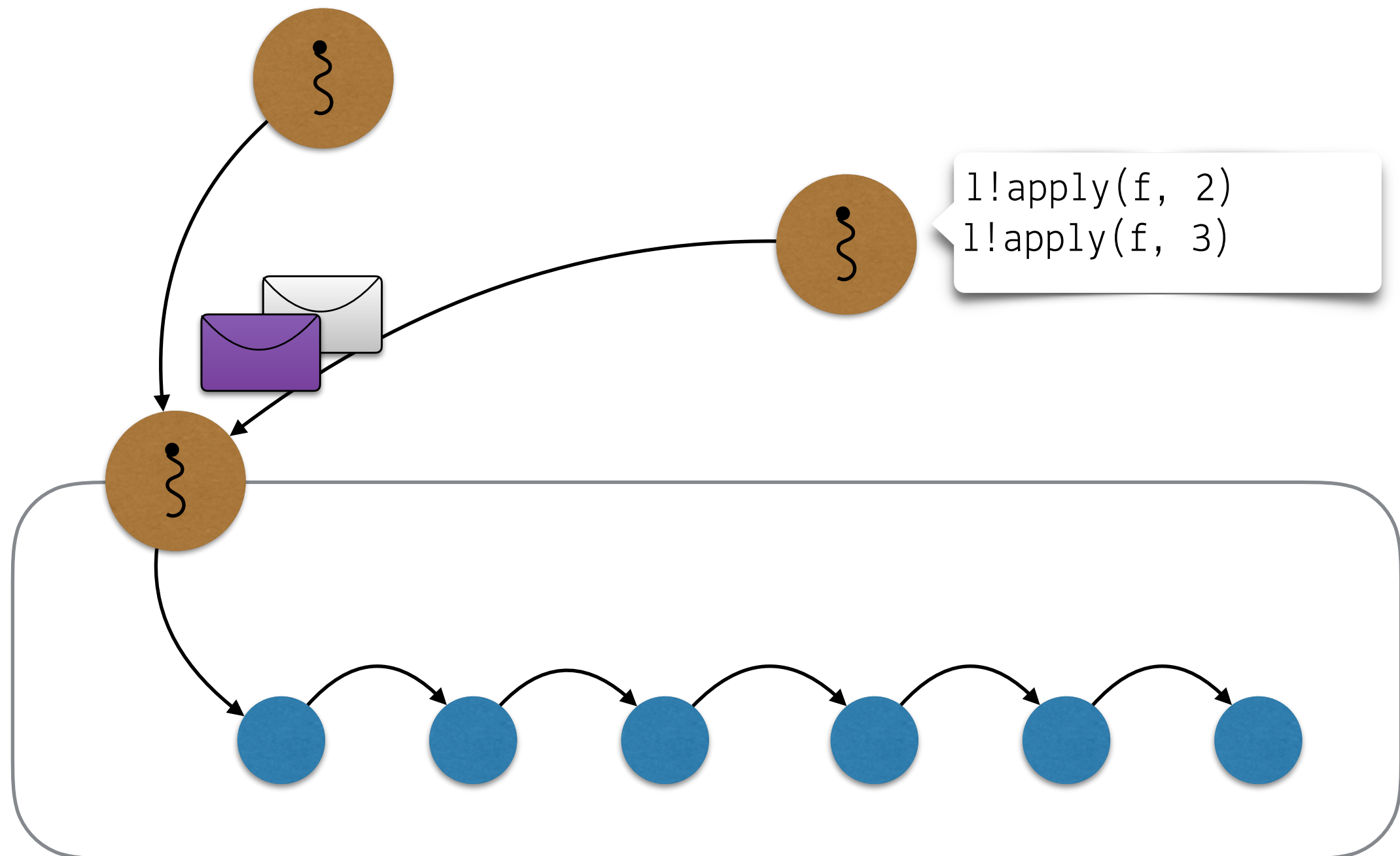
Tack! Frågor?



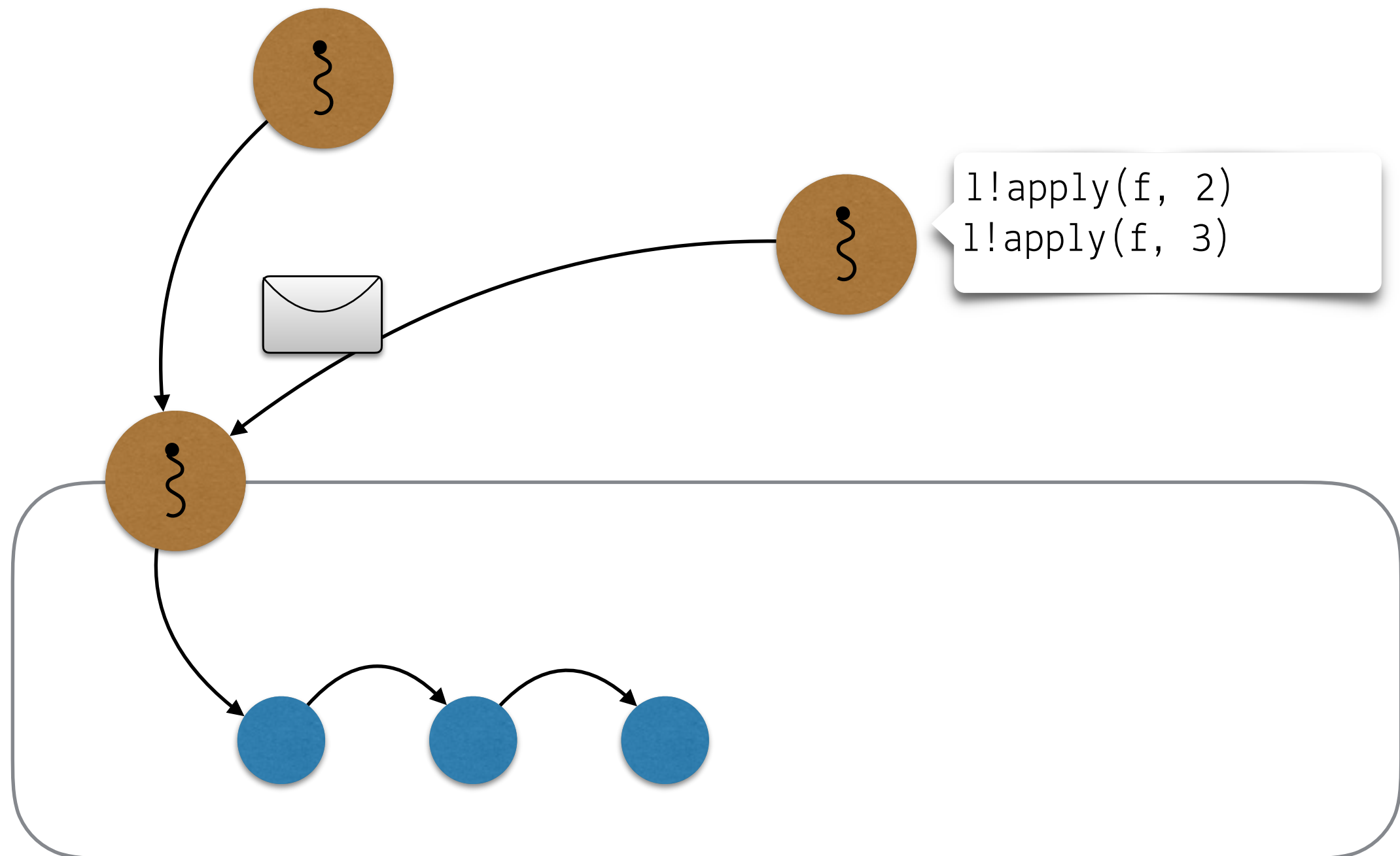
Incompatible Interleaving



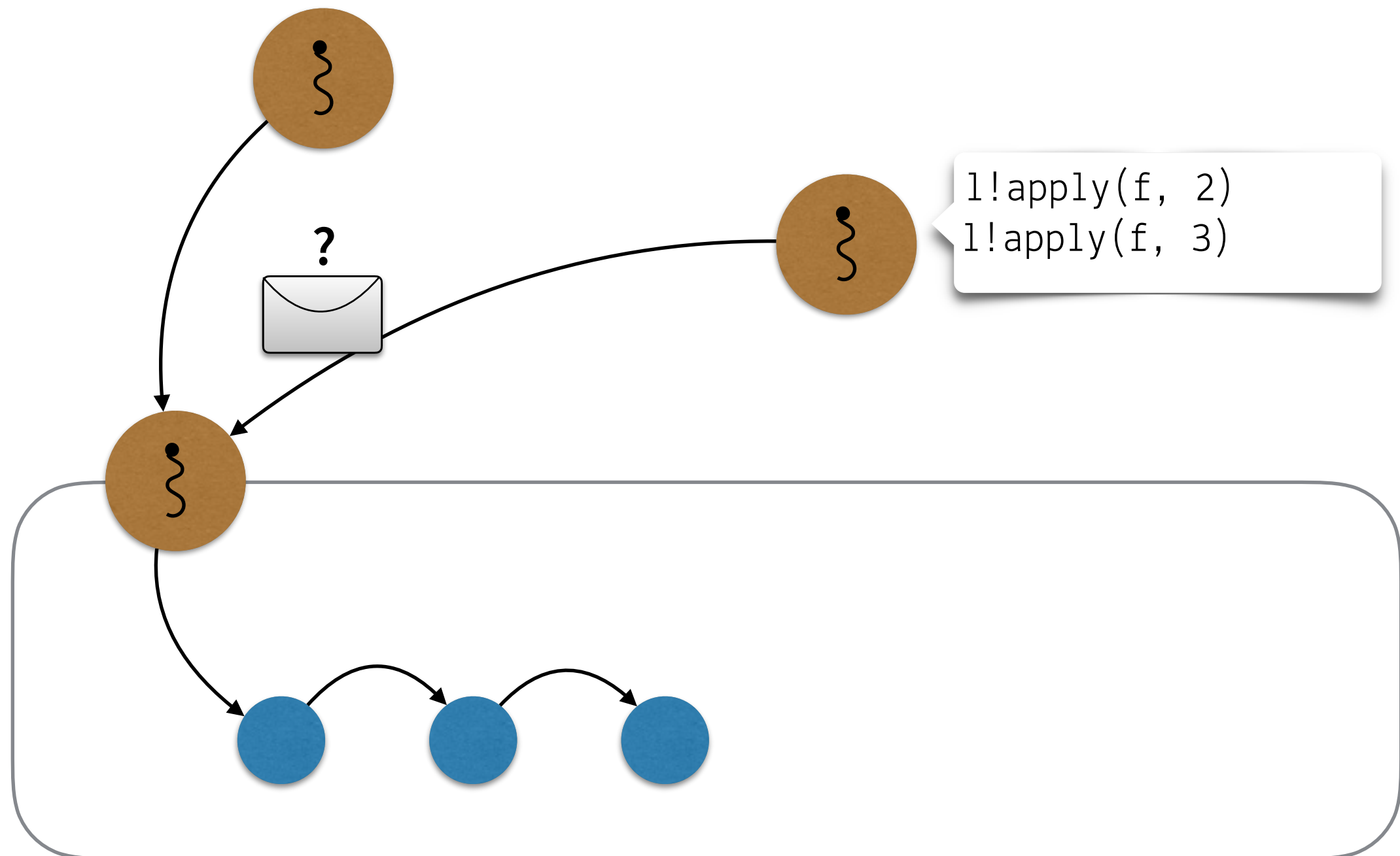
Incompatible Interleaving



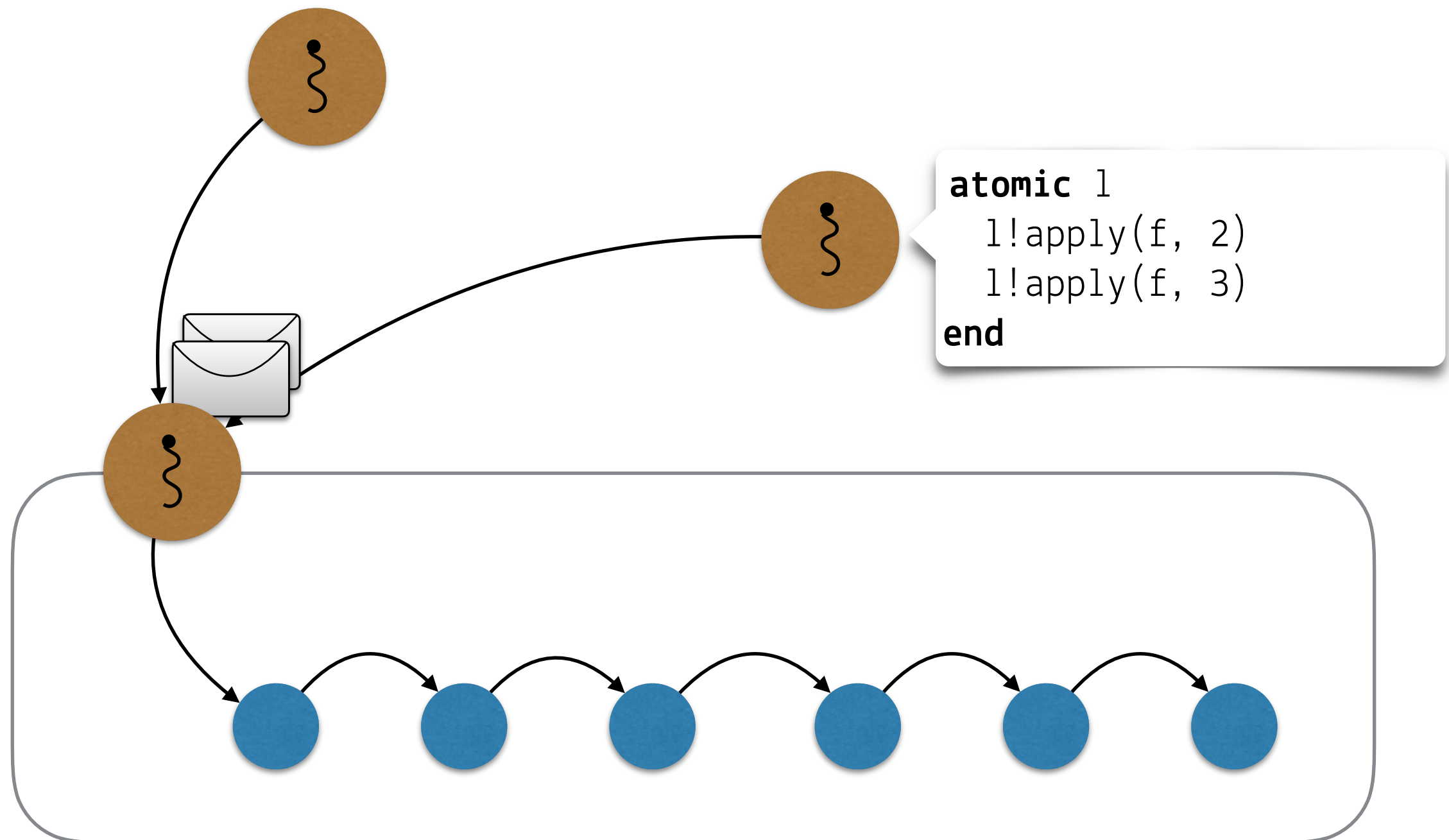
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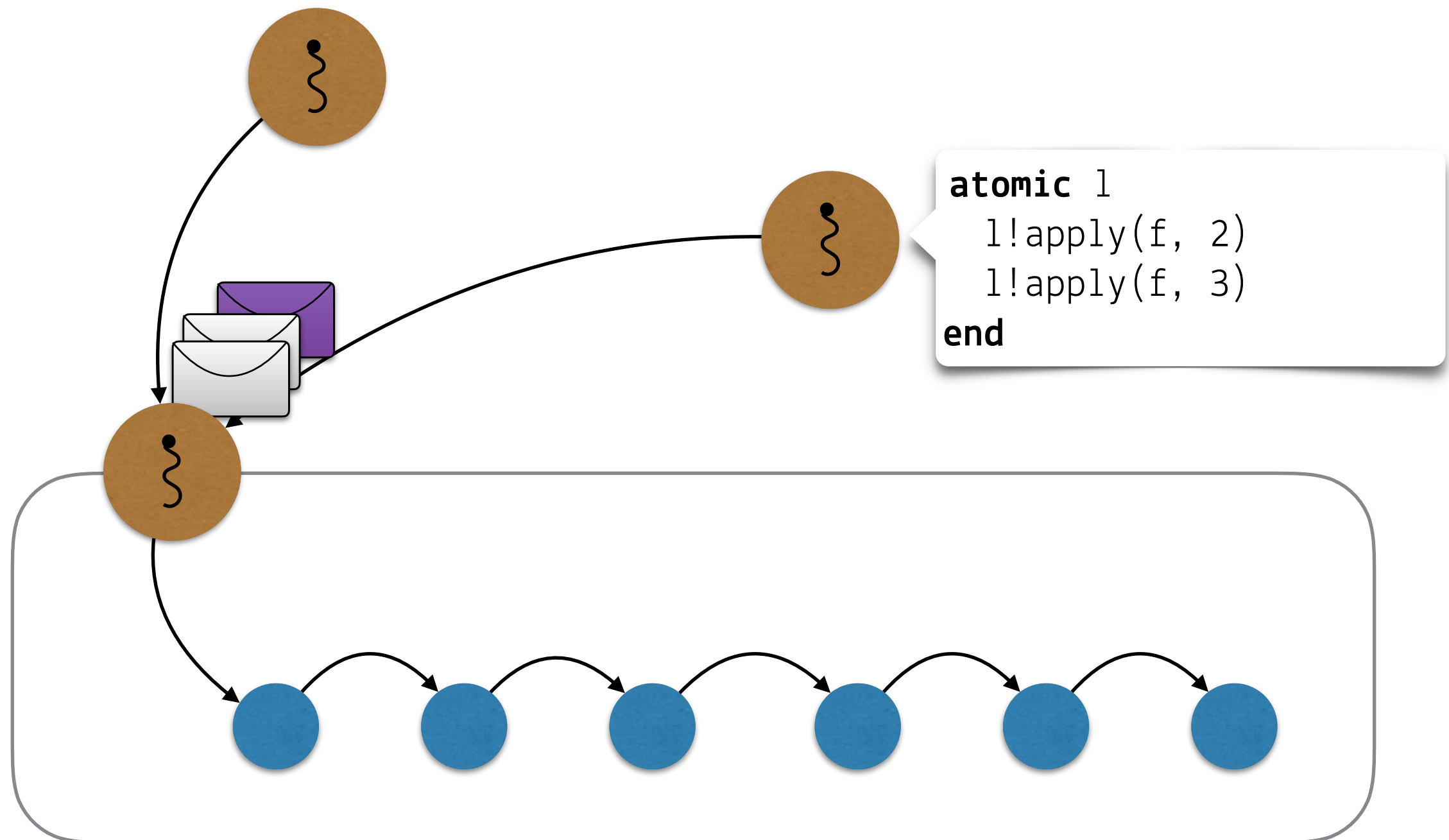
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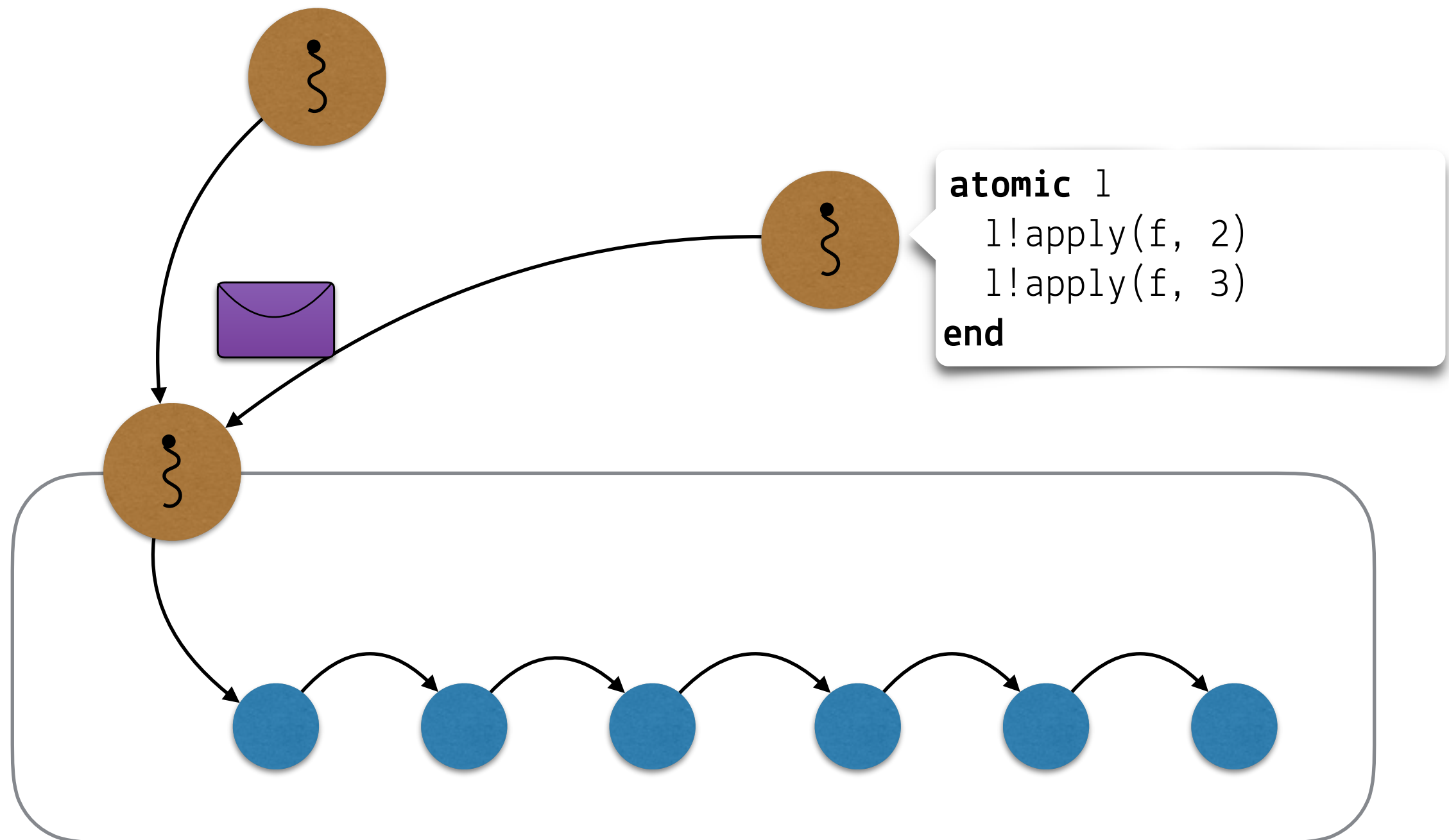
Grouping Messages



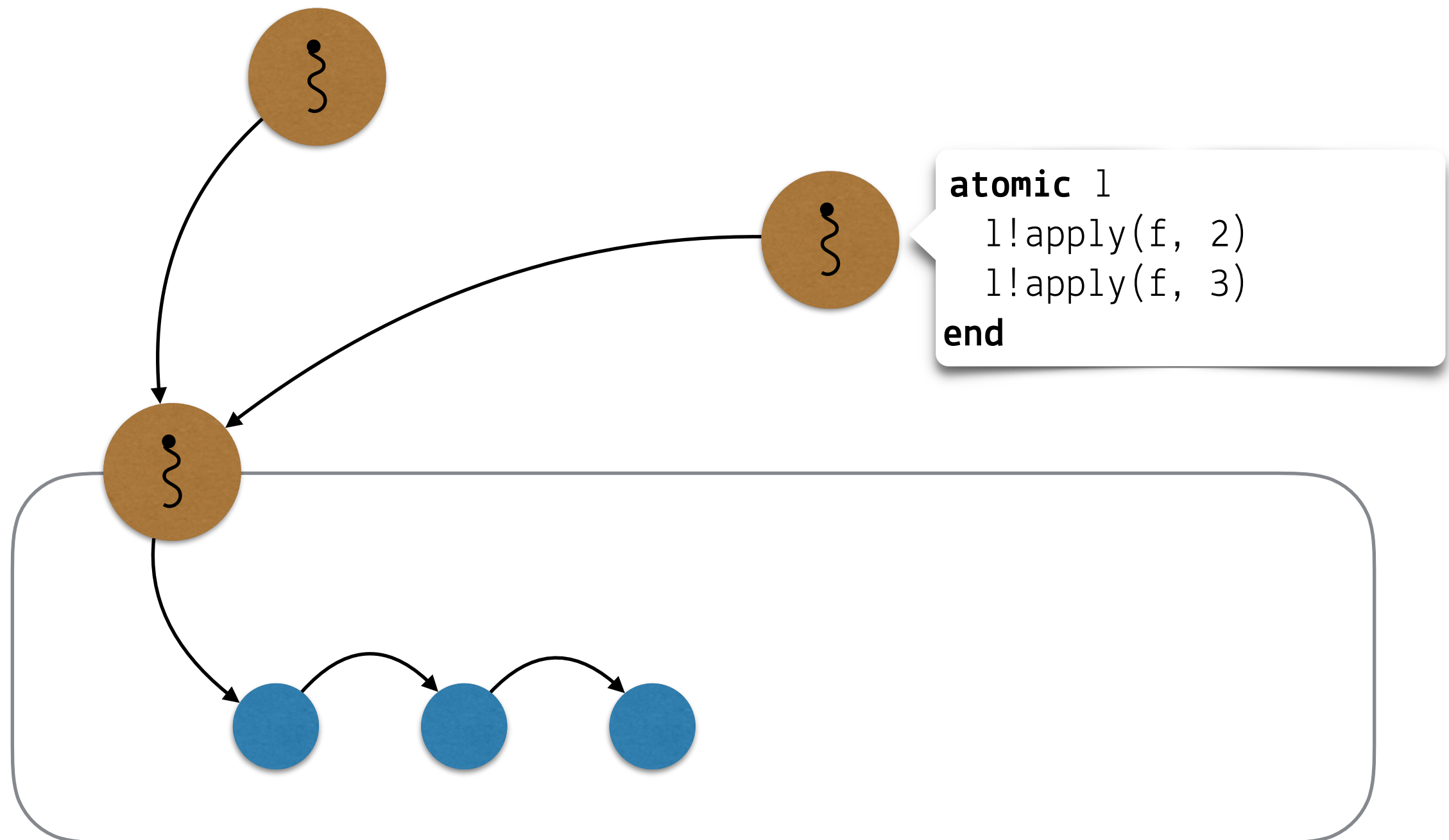
Grouping Messages



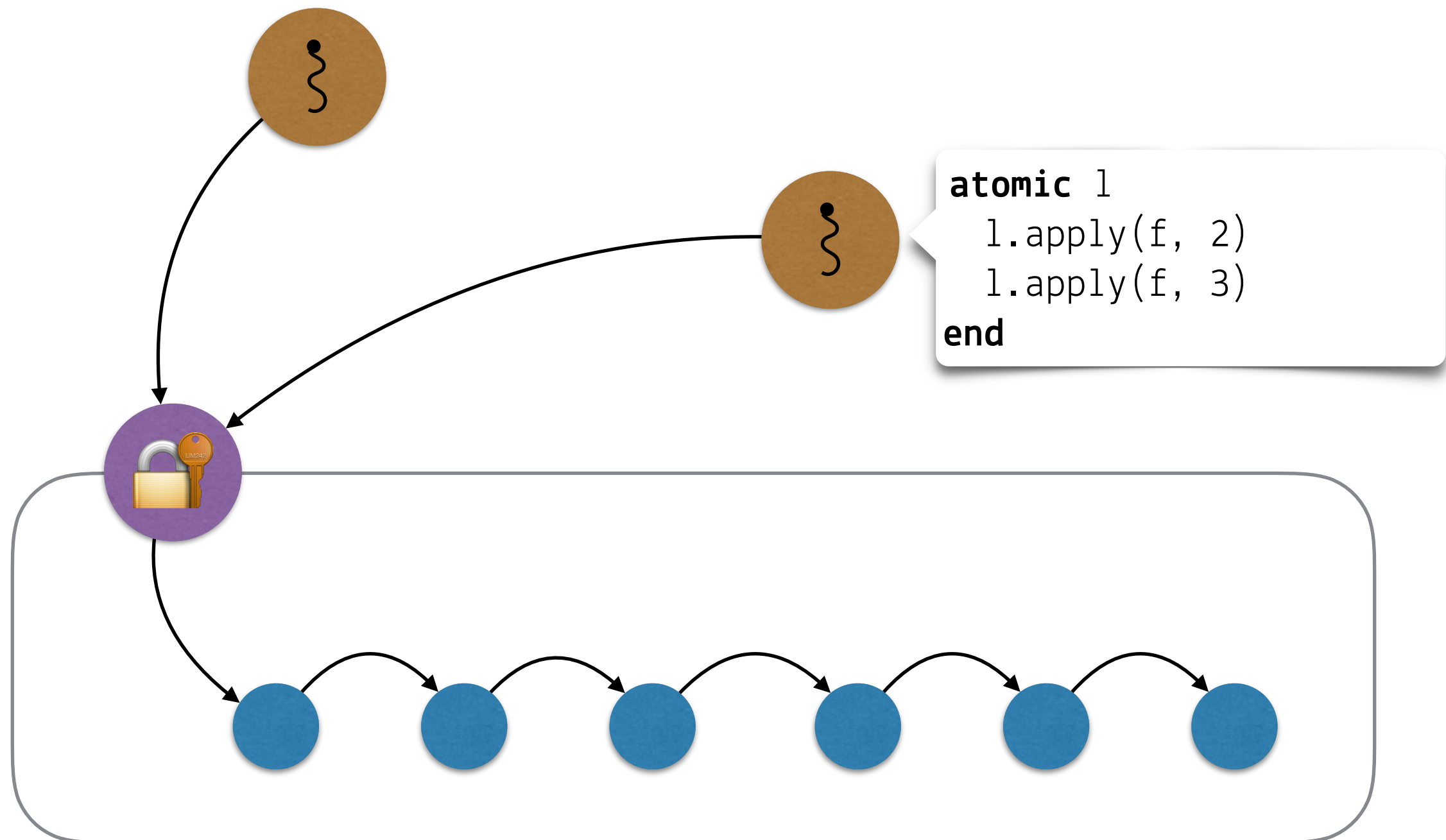
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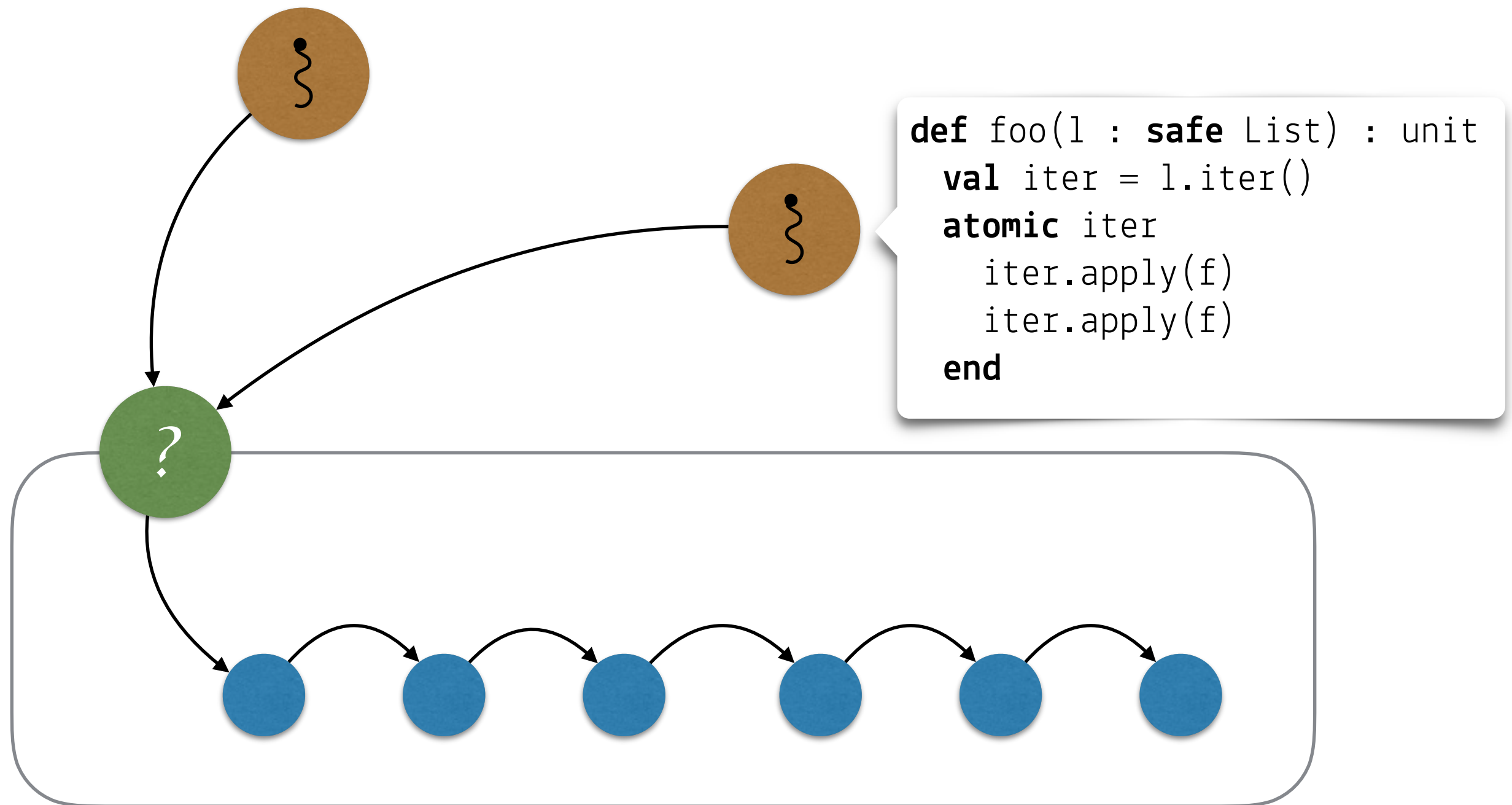
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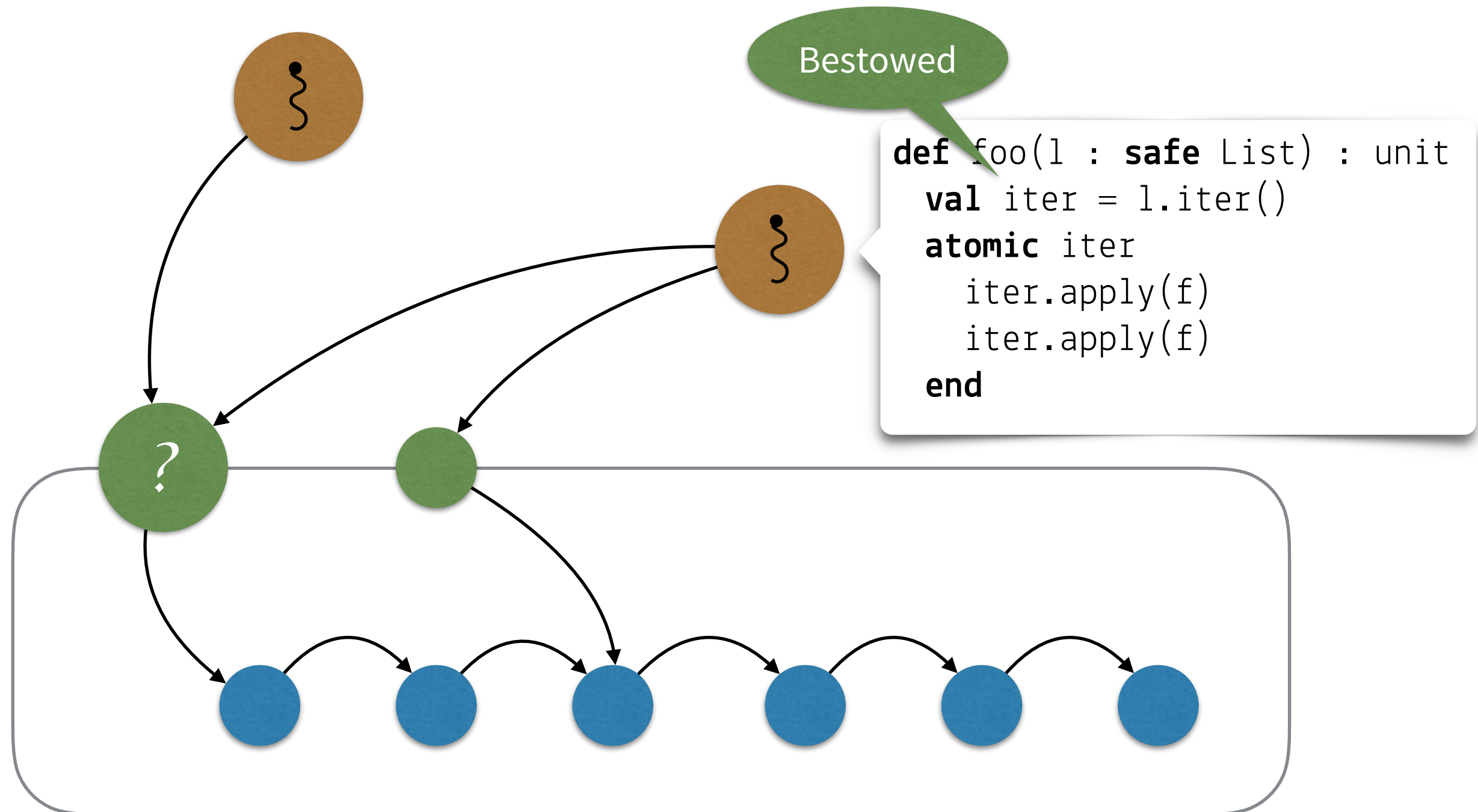
Grouping Locked operations



Polymorphic Concurrency Control



Polymorphic Concurrency Control



Tack! Frågor?



Summary

- Actor isolation can be relaxed by **bestowing** encapsulated objects with activity
 - All accesses will be synchronised via the message queue of the owning actor
 - Actors do not need to know the implementation of their bestowed objects
 - A bestowed object does not need to know that it is bestowed
- The same kind of relaxed encapsulation works for locks
 - All accesses will be synchronised via the lock of the owning object

- **Also in the paper:**

Atomic blocks to group operations

Implementation sketches

Polymorphic concurrency control