TMA4160 Cryptography – Fall 2010 – answers

Note: Several gcd-computations are omitted.

Problem 1

a. We compute

$$7k_1 + 4k_1^2 + 11k_1^3 + 15k_1^4 + k_2 \equiv 21 + 36 + 297 + 1215 + 21 \equiv 24 \pmod{29},$$

$$14k_1 + 10k_1^2 + k_2 \equiv 42 + 90 + 21 \equiv 8 \pmod{29}.$$

Alice sent the message HELP.

b. We see that $f(k_1, k_2, \text{KELP}) - f(k_1, k_2, \text{HELP}) = (10 - 7)k_1 = 3k_1 = 7 - 21 = 15$, and therefore $k_1 = 5$. We get

$$k_2 = 21 - (7k_1 + 4k_1^2 + 11k_1^3 + 15k_1^4) = 11.$$

Finally,

$$t \equiv 14k_1 + 10k_1^2 + k_2 \equiv 70 + 250 + 11 \equiv 12.$$

Problem 2.

a. We compute

and then because $650^2 \equiv 1 \pmod{1829}$,

$$650^{(1829-1)/2} = 650^{2^9}650^{2^8}650^{2^7}650^{2^4}650^2 = 1,$$

so 1829 is composite.

b. Following the algorithm in Stinson, we get:

$$x_0 = 2$$
 $x'_0 = 2^2 + 1 = 5$ $\gcd(5 - 2, 1829) = 1$
 $x_1 = 5$ $x'_1 = (5^2 + 1)^2 + 1 = 677$ $\gcd(677 - 5, 1829) = 1$
 $x_2 = 26$ $x'_2 = (677^2 + 1)^2 + 1 = 1080$ $\gcd(1080 - 26, 1829) = 31$

c. We multiply the three relations to get

$$(807 \cdot 1656 \cdot 1150)^2 = 5^4 \cdot 7^2 \cdot 19^2.$$

We get the square roots

$$807 \cdot 1656 \cdot 1150 \equiv 628 \pmod{1829}$$
 and $5^2 \cdot 7 \cdot 19 \equiv 1496 \pmod{1829}$,

and gcd(1496 - 628, 1829) = 31.

Problem 3

a. We compute

$$g^{m} = (1+n)^{m} + \langle n^{2} \rangle = \sum_{i=0}^{m} {m \choose i} 1^{i} n^{m-i} + \langle n^{2} \rangle$$
$$= 1 + {m \choose 1} n + \langle n^{2} \rangle = 1 + mn + \langle n^{2} \rangle,$$

since $\binom{m}{1} = m$. From this, it is clear that g has order n since $1 + n^2 + \langle n^2 \rangle = 1 + \langle n^2 \rangle$.

b. Let $x, y \in H$, so that for some $x_0, y_0 \in \mathbb{Z}_{n^2}^*$, $x = x_0^n$ and $y = y_0^n$. It is clear that $1 \in H$. We have

$$xy = (x_0^n)(y_0^n) = (x_0y_0)^n \in H,$$

and

$$x^{-1} = (x_0^n)^{-1} = (x_0^{-1})^n \in H.$$

Hence, H is a subgroup.

Let $x \in H$ and suppose $x = x_0^n$, $x_0 = a + \langle n^2 \rangle$. Then $x = (a + \langle n^2 \rangle)^n = a^n + \langle n^2 \rangle$ and $\phi(a + \langle n \rangle) = a^n + \langle n^2 \rangle = x$. Also, for any c relatively prime to n, $\phi(c + \langle n \rangle) = c^n + \langle n^2 \rangle = (c + \langle n^2 \rangle)^n \in H$. Hence, the image of ϕ is H.

c. Since H is the image of ϕ , we only need to show that it is an injective homomorphism. It is an homomorphism because

$$\phi((a+\langle n\rangle)(b+\langle n\rangle)) = \phi(ab+\langle n\rangle) = (ab)^n + \langle n^2\rangle = (a^n + \langle n^2\rangle)(b^n + \langle n^2\rangle)$$
$$= \phi(a+\langle n\rangle)\phi(b+\langle n\rangle).$$

It is injective because if $\phi(a + \langle n \rangle) = \phi(b + \langle n \rangle)$, then

$$a^n + \langle n^2 \rangle = b^n + \langle n^2 \rangle \Rightarrow a^n + \langle n \rangle \equiv b^n + \langle n \rangle \Rightarrow a + \langle n \rangle = b + \langle n \rangle,$$

which is true because n is invertible modulo (p-1)(q-1).

d. Exponentiation by un is the identity on \mathbb{Z}_n^* , hence it is the identity on H, because H is isomorphic to \mathbb{Z}_n^* . We compute

$$(xg^m)^{un} = x^{un}(g^n)^{mu} = x.$$

e. Let $c = \phi(r)g^m$. Then $c/c^{un} = \phi(r)g^m\phi(r)^{-1} = g^m$, and using the fact that $g^m = 1 + mn + \langle n^2 \rangle$, we can easily recover m.