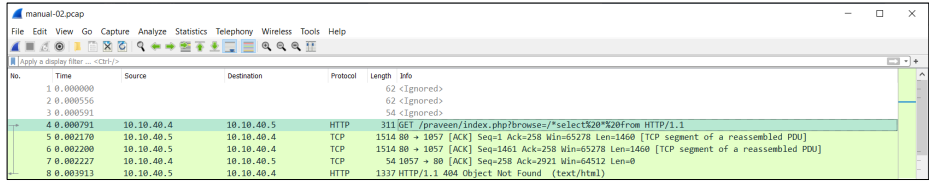




Research and Comments	There are several possibilities here as to the specific attack, but there are known vulnerabilities with overloading the buffer in IMAP in order to execute a denial of service. CVE-2004-1211 is one option here: "Multiple buffer overflows in the IMAP service in Mercury/32 4.01a allow remote authenticated users to cause a denial of service and possibly execute arbitrary code via long arguments to the ... (8) FETCH..." <sup>5</sup> The vulnerabilities in this IMAP configuration can allow for remote validated users to attempt a denial of service attack and to possibly inject code. <sup>4</sup>
Sources	<ol style="list-style-type: none"> <li>1. Busylog.net. "Telnet IMAP Commands Note," November 27, 2012. <a href="https://busylog.net/telnet-imap-commands-note/">https://busylog.net/telnet-imap-commands-note/</a></li> <li>2. Crispin &lt;mrc@cac.washington.edu&gt;, Mark R. "INTERNET MESSAGE ACCESS PROTOCOL - VERSION 4rev1." Accessed December 4, 2020. <a href="https://tools.ietf.org/html/rfc3501#section-6.4.5">https://tools.ietf.org/html/rfc3501#section-6.4.5</a></li> <li>3. "CVE - CVE-2004-1211." Accessed December 4, 2020. <a href="http://cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2004-1211">http://cve.mitre.org/cgi-bin/cvename.cgi?name=CVE-2004-1211</a></li> <li>4. FortiGuard "IMAP.Commands.Buffer.Overflow   IPS." Accessed December 4, 2020. <a href="https://fortiguard.com/encyclopedia/ips/11599">https://fortiguard.com/encyclopedia/ips/11599</a></li> <li>5. "NVD - CVE-2004-1211." Accessed December 4, 2020. <a href="https://nvd.nist.gov/vuln/detail/CVE-2004-1211">https://nvd.nist.gov/vuln/detail/CVE-2004-1211</a></li> </ol>

2	
Diagnosis	SQL injection in search bar
Wireshark Analysis	<p>Basic HTML traffic. Packet 4 [Figure 2.1] of the packet capture shows the attempted SQL injection through the search function of the website, using the SQL phrase "SELECT * FROM." The search then returned a 404 error page in response. Unclear (to this analyst) if the SQL injection was successful.</p>  <p>Figure 2.1</p>
Research and Comments	SQL injection is an exploitation method used to try and access information contained in a SQL database through unguarded access. Search bars and other input fields on a webpage are easy points of access, as webpages looking to collect some form of data from the user need to have those points of entry open for collection. Users with ill intent can attempt to inject SQL into these fields in hopes of finding a weakness. The user here attempted a basic SQL command and received a 404 error, so going by this packet capture alone, it seems the user was unsuccessful.
Sources	<ol style="list-style-type: none"> <li>1. Acunetix. "PHP Injection: Directory Traversal &amp; Code Injection." Accessed December 4, 2020. <a href="https://www.acunetix.com/websitesecurity/php-security-2/">https://www.acunetix.com/websitesecurity/php-security-2/</a></li> <li>2. "SQL Injection in Search Fields   Explore Security." Accessed December 4, 2020. <a href="https://www.exploresecurity.com/sql-injection-in-search-fields/">https://www.exploresecurity.com/sql-injection-in-search-fields/</a></li> <li>3. w3resource. "SQL Injection Tutorial." Accessed December 4, 2020. <a href="https://www.w3resource.com/sql/sql-injection/sql-injection.php">https://www.w3resource.com/sql/sql-injection/sql-injection.php</a></li> </ol>

## Diagnoses

## ICMP Flood Attack

## Wireshark Analysis

The capture contains a hair under a million packets. The 'capture file properties' [Figure 3.1] shows that the average PPM (packets per minute) were approximately 15,187. In addition, the capture contains nothing but pings, with requests and replies back and forth. The length of each packet [Figure 3.2] remains identical through the entire capture (verified using the filter: `frame.len!=60 && frame.len!=42`, which yielded no results). This presumably means the payload from each request was identical, and the purpose of which was to overload the network with requests in order to hinder or completely shut down network activity. The source IP indicated as being on the same network, so it is likely that the attacker spoofed their source in order to further befuddle the victim's network.

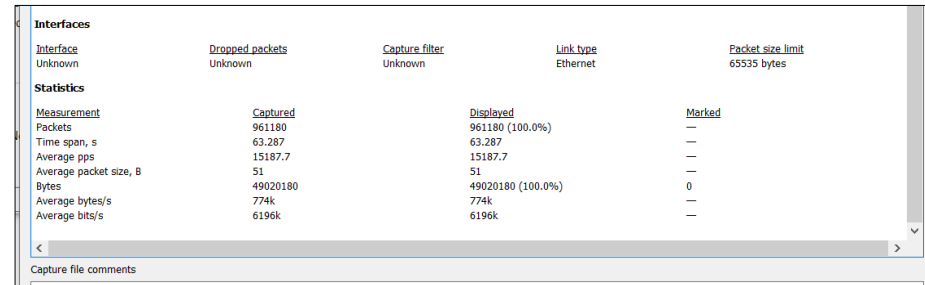


Figure 3.1

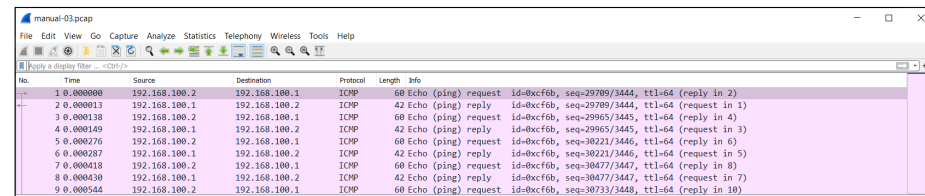


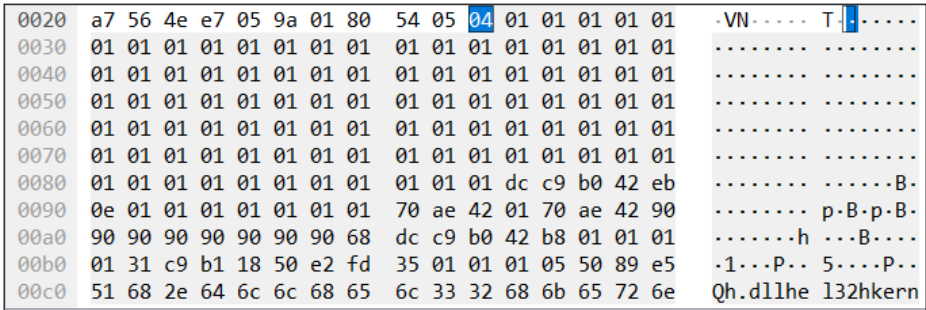
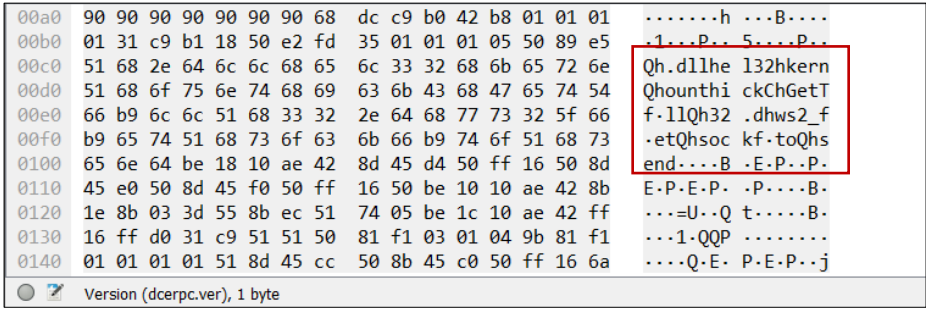
Figure 3.2

## Research and Comments

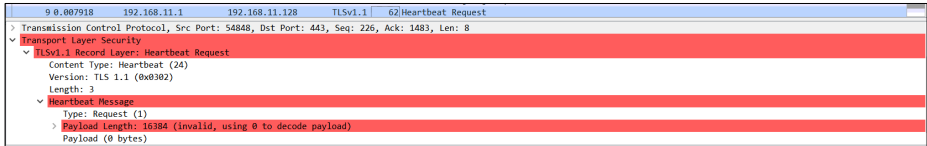
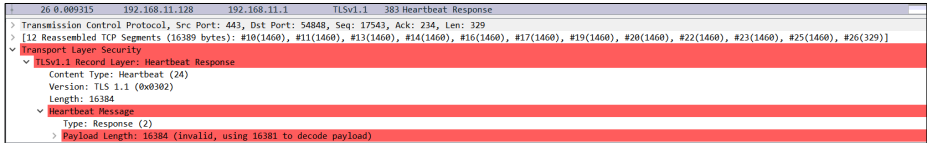
Ping attacks, or ICMP attacks work by sending a literal flood of request packets, as the network is required to send a response to each incoming request. This causes strain on the network, which cannot keep up with the requests, and can lead a complete denial of service. By spoofing the target's IP address, the attacker's own sources are freed to just send packets, while the source is forced to receive and send.<sup>1</sup>

## Sources

1. "Crisuolo - Distributed Denial of Service - a396999.pdf." Accessed December 4, 2020. <https://apps.dtic.mil/dtic/tr/fulltext/u2/a396999.pdf>
2. Learning Center. "What Is a Ping Flood | ICMP Flood | DDoS Attack Glossary | Imperva." Accessed December 4, 2020. <https://www.imperva.com/learn/ddos/ping-icmp-flood/>

1	
Diagnosis	MS-SQL Slammer Worm
Snort Output Analysis	<p>A single packet was analyzed. It was a UDP packet, with no fragmentation, only containing 418 bits. Nothing else of interest. The Slammer Worm uses a single UDP packet to gain entry.</p> <p>The alert file recognized the packet as "MS-SQL Slammer Worm Activity," with an additional warning about potentially bad traffic. It also indicated that traffic was going to UDP port 1434, a recognized trait of the Slammer Worm.</p>
Wireshark Analysis	<p>Identified the packet as using the DCERPC protocol, a protocol used for procedure calls on remote machines, using UDP. Looking at the hex output, there are several features that further confirm the diagnoses. It starts with '04', which indicates to the SQL server that the following data will be name of the online database it is looking. It is then padded with series of '01.' [Figure 4.1], another indicator that this is indeed a Slammer Worm. The payload also contains the strings ""h.dllhel32hkernQhounthickChGetTf", "hws2", "Qhsockf" and "toQhsend"<sup>2</sup>, which are yet more traits of the Slammer Worm. [Figure 4.2]</p> <div>  <p>Figure 4.1</p> </div> <div>  <p>Figure 4.2</p> </div>
Research and Comments	The SQL Slammer Worm works by exploiting the small space where SQL attempts to perform the normal operation of reading in the data. When it encounters the overlong payload, the SQL server rewrites its own stack, unwittingly

	reprogramming itself with instructions from the worm. The worm then self-propagates by targeting random IP addresses and sending itself to them. This leads to a huge volume of traffic, with the intent to perform a DOS.
Sources	<ol style="list-style-type: none"> <li>1. "2003 - Attack of Slammer Worm - A Practical Case Study.Pdf." Accessed December 4, 2020. <a href="https://www.giac.org/paper/gcih/414/attack-slammer-worm-practical-case-study/103632">https://www.giac.org/paper/gcih/414/attack-slammer-worm-practical-case-study/103632</a></li> <li>2. "2003 - MS SQL SlammerSapphire Worm.Pdf." Accessed December 4, 2020. <a href="https://www.giac.org/paper/gsec/3091/ms-sql-slammer-sapphire-worm/105136">https://www.giac.org/paper/gsec/3091/ms-sql-slammer-sapphire-worm/105136</a></li> <li>3. "What Is DCE RPC?" Accessed December 4, 2020. <a href="http://www.doublersolutions.com/docs/dce/OSFdocs/htmls/overview/Dceint63.htm">http://www.doublersolutions.com/docs/dce/OSFdocs/htmls/overview/Dceint63.htm</a></li> </ol>

2	
Diagnosis	Heartbleed / CVE-2014-0160
Snort Output Analysis	<p>The capture contained 30 packets, and logged two alerts.</p> <p>The alert log recognized a "heartbeat read overrun attempt", which is indicative of a Heartbleed attempt. This is the user asking for data. The second alert of "large heartbeat response - possible ssl heartbleed attempt" is also an alert for CVE-2014-0160, where the server attempted to send the request back.</p>
Wireshark Analysis	<p>The capture opens and closes with a standard exchange of signatures. Packet 9 of the capture contains the heartbeat request – the payload itself is very small, and only contains a few bytes [Figure 5.1]. Packet 26 contains the response, and shows a payload that is considerably higher than the request, at 16,384 bytes. [Figure 5.2]</p>  <p>Figure 5.1</p>  <p>Figure 5.2</p>
Research and Comments	<p>An SSL heartbeat allows two computers to communicate and inform each other that they are still connected. The Heartbleed vulnerability took advantage of the flaws in earlier implementations of OpenSSL, which failed to implement any sort of check to ensure that the incoming heartbeat requests were the length they claimed to be. This allowed attackers to send heartbeat request that claimed to be some small amount but claimed to much larger. The other computer would send a heartbeat back, but since the original message was so small, there was a lot of empty space – so the computer would automatically fill the space with whatever followed the original message in it's memory. By taking advantage of this flaw, attackers could steal up to 64k of memory.</p>

Sources	<ol style="list-style-type: none"> <li>1. "CVE - CVE-2014-0160." Accessed December 4, 2020. <a href="https://cve.mitre.org/cgi-bin/cvename.cgi?name=cve-2014-0160">https://cve.mitre.org/cgi-bin/cvename.cgi?name=cve-2014-0160</a></li> <li>2. Fruhlinger, Josh. "What Is the Heartbleed Bug, How Does It Work and How Was It Fixed?" CSO Online, September 13, 2017. <a href="https://www.csoonline.com/article/3223203/what-is-the-heartbleed-bug-how-does-it-work-and-how-was-it-fixed.html">https://www.csoonline.com/article/3223203/what-is-the-heartbleed-bug-how-does-it-work-and-how-was-it-fixed.html</a></li> <li>3. "NVD - CVE-2014-0160." Accessed December 4, 2020. <a href="https://nvd.nist.gov/vuln/detail/CVE-2014-0160">https://nvd.nist.gov/vuln/detail/CVE-2014-0160</a></li> <li>4. OpenSSL Security Advisory [07 Apr 2014]. Accessed December 4, 2020. <a href="https://www.openssl.org/news/secadv/20140407.txt">https://www.openssl.org/news/secadv/20140407.txt</a></li> </ol>
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### 3

Diagnosis	UDP Port Scanning																																																																						
Snort Output Analysis	<p>Output file shows 1000 UDP packets, with one alert, no fragmentation, and 588 UDP sessions created and deleted.</p> <p>There was a single alert for “UDP filtered portscan.”</p>																																																																						
Wireshark Analysis	<p>The packet capture shows that all traffic was between the exact same source and destination. About half of the packets are malformed. All the malformed packets also indicate that they used the HART-IP protocol, which is a diagnostic protocol. Every single packet was also the same length. The key here is that the source port never changes, but the destination port is constantly changing [Figure 6.1], a good indicator of scanning activity.</p> <table><thead><tr><th>No.</th><th>Time</th><th>Source</th><th>Destination</th><th>Protocol</th><th>Length</th><th>Info</th></tr></thead><tbody><tr><td>871</td><td>38.044792</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 29514 Len=14</td></tr><tr><td>780</td><td>36.714478</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 29514 Len=14</td></tr><tr><td>631</td><td>34.396305</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 9801 Len=14</td></tr><tr><td>630</td><td>34.369262</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 48986 Len=14</td></tr><tr><td>629</td><td>34.351838</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 20445 Len=14</td></tr><tr><td>628</td><td>34.340754</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 31301 Len=14</td></tr><tr><td>627</td><td>34.320231</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 43727 Len=14</td></tr><tr><td>626</td><td>34.311354</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 14468 Len=14</td></tr><tr><td>625</td><td>34.288512</td><td>10.0.32.25</td><td>80.237.98.132</td><td>UDP</td><td>46</td><td>10126 → 52202 Len=14</td></tr></tbody></table> <p>Figure 6.1</p>	No.	Time	Source	Destination	Protocol	Length	Info	871	38.044792	10.0.32.25	80.237.98.132	UDP	46	10126 → 29514 Len=14	780	36.714478	10.0.32.25	80.237.98.132	UDP	46	10126 → 29514 Len=14	631	34.396305	10.0.32.25	80.237.98.132	UDP	46	10126 → 9801 Len=14	630	34.369262	10.0.32.25	80.237.98.132	UDP	46	10126 → 48986 Len=14	629	34.351838	10.0.32.25	80.237.98.132	UDP	46	10126 → 20445 Len=14	628	34.340754	10.0.32.25	80.237.98.132	UDP	46	10126 → 31301 Len=14	627	34.320231	10.0.32.25	80.237.98.132	UDP	46	10126 → 43727 Len=14	626	34.311354	10.0.32.25	80.237.98.132	UDP	46	10126 → 14468 Len=14	625	34.288512	10.0.32.25	80.237.98.132	UDP	46	10126 → 52202 Len=14
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Research and Comments	UDP port scans target a specific IP address in order to find open or responsive services. They have internal patterns to randomize which ports they check. Once they receive a response from a port, it means it is open and accessible. This can give attackers access to even further information.																																																																						
Sources	<ol style="list-style-type: none"><li>1. Blog, WhatsUp Gold. “Port Scanning 101: What It Is, What It Does and Why Hackers Love It - WhatsUp Gold.” Accessed December 4, 2020. <a href="https://www.whatsupgold.com/blog/port-scanning-101-what-it-is-what-it-does">https://www.whatsupgold.com/blog/port-scanning-101-what-it-is-what-it-does</a></li><li>2. “Configuring UDP Port Scan Attack Screen - TechLibrary - Juniper Networks.” Accessed December 4, 2020. <a href="https://www.juniper.net/documentation/en_US/junos-cc19.2/topics/reference/general/security/19.2r1-cc/configuring-udp-port-scan-attack.html">https://www.juniper.net/documentation/en_US/junos-cc19.2/topics/reference/general/security/19.2r1-cc/configuring-udp-port-scan-attack.html</a></li></ol>																																																																						

### 4

Diagnosis	Eternal Blue Trojan
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Snort Output Analysis	<p>The output file registered 577 packets: TCP and IP4 both registered 573 packets, with 4 left over registered as ARP. It also shows in the Preprocessor Statistics that there were 47 total sessions, with 43 aborted, leaving 4.</p> <p>The alert file registered 8 alerts – 4 requests and 4 responses. All were identified as “MALWARE-CNC Win.Trojan.Eternalblue variant echo [either request or response]”</p>
Wireshark Analysis	<p>The alerts in the Snort alert file correspond with 4 SMB echo requests and 4 SMB echo responses. The other SMB packets are establishing the connection to the computer. After initiating a connection, the trojan attempted to send a large NT Transfer packet [packet 14]. Because it is so large, it triggered additional Trans2 Requests. When sending the additional requests, the Trojan deliberately tried to send malformed ones in order to try to exploit vulnerabilities. [e.g. packet 165]. It kept restarting in an attempt to gain entry; 4 times it attempted to request a Trans 2 session setup [packets 12, 257, 270, 565].</p>
Research and Comments	<p>Eternal Blue took advantage of a vulnerability in SMB traffic that allowed attackers to try and disguise the trojan as legitimate traffic being sent back to the computer. The vulnerability becomes available when triggered by malformed Trans2 Requests. Eternal Blue, and its more recent brother WannaCry, are ransomware trojans. It is a trojan specific to Windows, and while Windows released patches and updates to address the issue, users still had to download the updates themselves, and businesses are notorious for running on ancient software.</p>
Sources	<ol style="list-style-type: none"> <li>1. “NVD - CVE-2017-0144.” Accessed December 4, 2020. <a href="https://nvd.nist.gov/vuln/detail/CVE-2017-0144">https://nvd.nist.gov/vuln/detail/CVE-2017-0144</a></li> <li>2. FireEye. “SMB Exploited: WannaCry Use of ‘EternalBlue.’” Accessed December 4, 2020. <a href="https://www.fireeye.com/blog/threat-research/2017/05/smb-exploited-wannacry-use-of-eternalblue.html">https://www.fireeye.com/blog/threat-research/2017/05/smb-exploited-wannacry-use-of-eternalblue.html</a></li> <li>3. What Is EternalBlue and Why Is the MS17-010 Exploit Still Relevant? “What Is EternalBlue and Why Is the MS17-010 Exploit Still Relevant?” Accessed December 4, 2020. <a href="https://www.avast.com/c-eternalblue">https://www.avast.com/c-eternalblue</a></li> </ol>

