

2. No: Regardless of whether the final serial order is, in order to get the T1 transactions together (or in other words, in order to get the T3 transactions together), we would have to switch the pairs R3(B) and W1(B) and/or W1(B) and W3(B), which are both conflicts (in the table below in red highlights). In other words, a cycle from T3 to T1 and back to T3 occurs, so the schedule is non conflict-serializable (switching pairs by hand avoiding conflicts as much as possible also results in the conflict at W1(B) and W3(B)).

| T1 | T2 | T3 |
|--------|--------|--------|
| R(A) | | |
| W(A) | | |
| | | R(A) |
| | | W(A) |
| | R(A) | |
| R(B) | | |
| | | R(B) |
| W(B) | | |
| | | W(B) |
| | R(B) | |
| | commit | |
| commit | | |
| | | commit |