

23 Sept 2019

Dynamic Programming

- related to DTC: we solve a problem by breaking down into simpler problems & then combining the simpler problems.
- DP helps when we have independent overlap between the subproblems / simpler problems

DP Stock Exchange (p , n)

$\Theta(n)$ { soln \leftarrow array of size n
~~so that~~
minidx $\leftarrow 1$ } the soln for $p[1 \dots 1]$
sol[1] $\leftarrow 0$ }

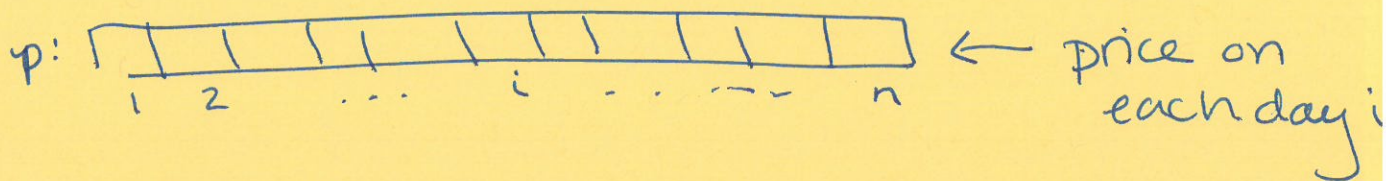
for $i = 2 \dots n$

$\Theta(1)$ { repeated n -times
 \Downarrow
 $\Theta(n)$ for the for loop. }

if (soln[i-1] < $p[i] - p[\text{minidx}]$)
| sol'n[i] $\leftarrow p[i] - p[\text{minidx}]$
else
| soln[i] \leftarrow soln[i-1]
endelse
if ($p[i] < p[\text{minidx}]$)
| minidx $\leftarrow i$
~~end if~~
endif
end for
return soln[n]

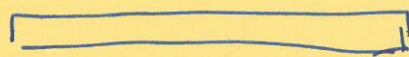
Runtime: $\Theta(n)$ Linear!

Recall the stock market problem:

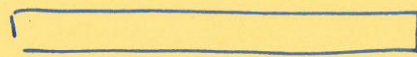


goal: find $i < j$ to max $p(j) - p(i)$

Sol'n:



↑
solve (L)



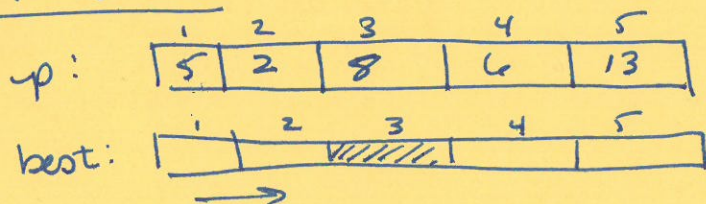
↑
solve (R)

$$\max(\text{solve}(L), \text{solve}(R), \max(R) - \min(L))$$

time complexity :

$$T(n) = 2T\left(\frac{n}{2}\right) + \Theta(n)$$
$$= \Theta(n \log n)$$

Another Sol'n



$i = 3$
best(i) holds
sol'n for $p[1 \dots i]$
and any info needed
to solve $p[1 \dots i+1]$

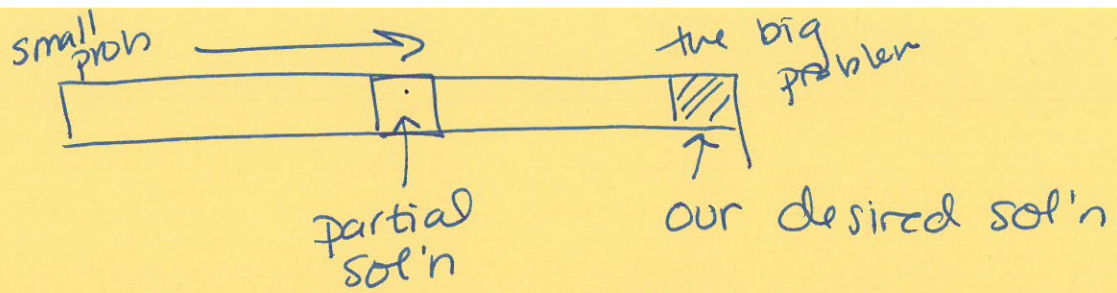
best[3] = { solution = 8 - 2 = 6
pair (2, 3)
minidx = 2

$\text{best}[4] = \text{sol'n} = \max(\text{best}[3].\text{sol'n}, \text{up}[4] - p[\text{minidx}])$

Another ex: $p = [2, 8, 1, 4, \dots]$

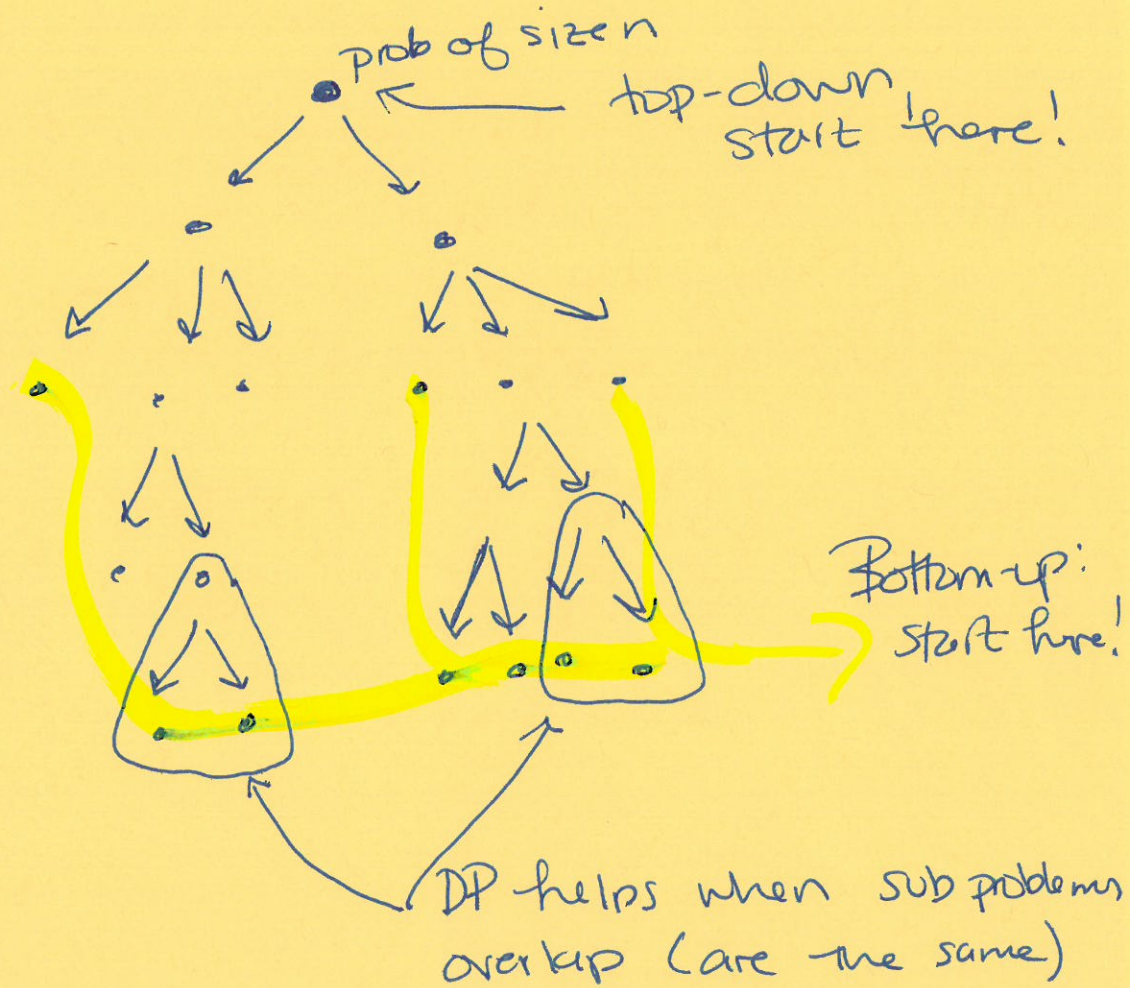
this update
is now
constant time!

Array



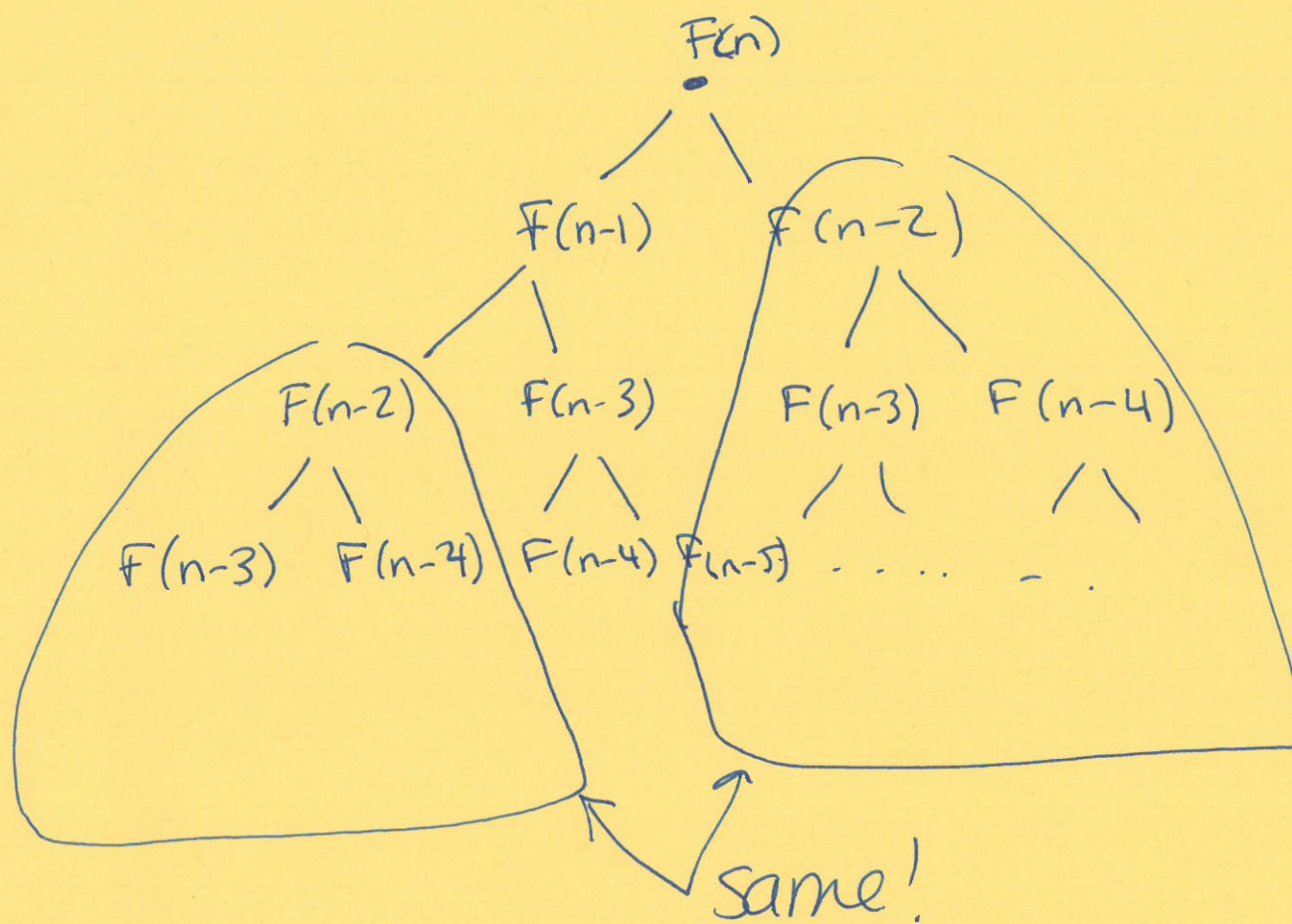
"bottom-up" approach: start w/ the small + grow until you solve the big problem.

Schematically
Think of recursion
tree

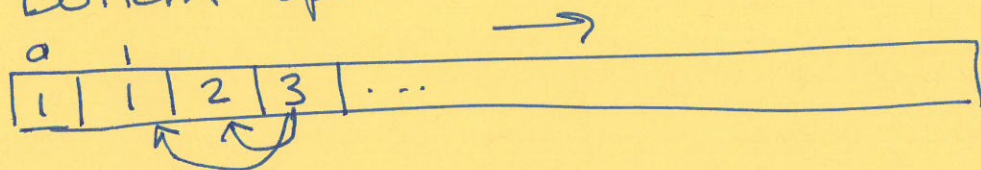


$$F(0) = F(1) = 1$$

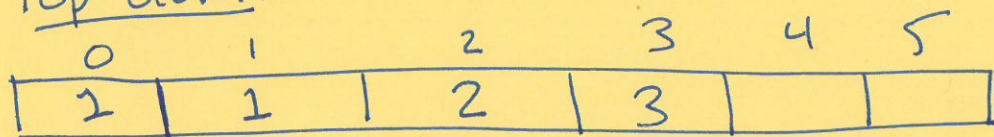
$$F(n) = F(n-1) + F(n-2), \quad \forall n \geq 2$$



Bottom-up:



Top-down:



$$F(5) = F(4) + F(3)$$

$$F(4) = F(3) + F(2)$$

$$F(3) = F(2) + F(1)$$

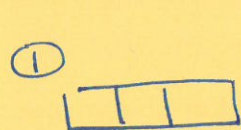
$$F(2) = F(1) + F(0)$$

Rod Cutting

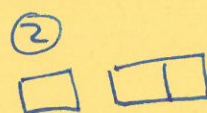
- given: a rod of length n , profit array
- want: cut into pieces (of unit multiples) to maximize profit

e.g., $n=3$
profit = $[6, 2, 17]$

- profit[i] = profit of having a rod of length = i
- profit is length n



17



$$6 + 2 = 8$$



$$6 + 6 + 6 = 18$$

The optimal sol'n
(note: sometimes opt is not unique!)

in groups: ① what is a DP sol'n to this?

② what is the $\Theta(2^n)$ brute force approach?

(Answers next time)