

# Automata As Input

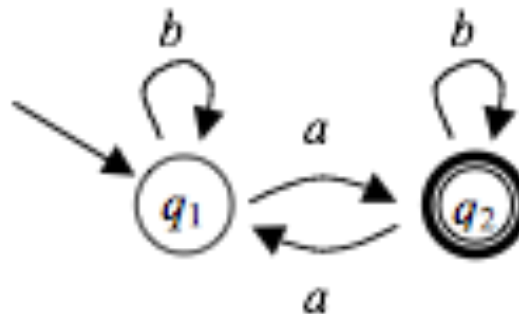
- Our goal is to make TMs that can simulate other automata, given as input
- TMs can only take strings as input, so we need a way to encode automata as strings
- We'll start with the simplest: DFAs...

# DFAs Encoded Using $\{0,1\}$

- The DFA's alphabet and strings:
  - Number  $\Sigma$  arbitrarily as  $\Sigma = \{\sigma_1, \sigma_2, \dots\}$
  - Use the string  $1^i$  to represent symbol  $\sigma_i$
  - Use 0 as a separator for strings
  - For example, if  $\Sigma = \{a, b\}$ , let  $a = \sigma_1$  and  $b = \sigma_2$ ; then *abba* is represented by 101101101
- The DFA's states:
  - Number  $Q = \{q_1, q_2, \dots\}$ , making  $q_1$  the start state and numbering the others arbitrarily
  - Use the string  $1^i$  to represent symbol  $q_i$

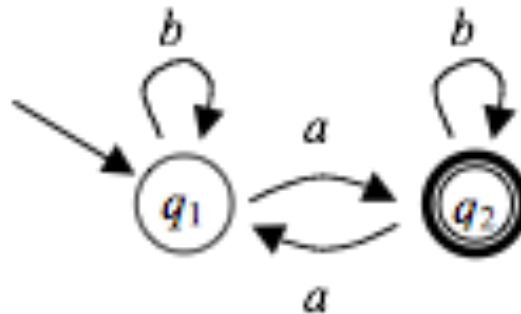
# DFA Encoding, Continued

- The DFA's transition function:
  - Encode each transition  $\delta(q_i, \sigma_j) = q_k$  as a string  $1^i 0 1^j 0 1^k$
  - Encode the entire transition function as a list of such transitions, in any order, using 0 as a separator
  - For example,



- Numbering  $a$  as  $\sigma_1$  and  $b$  as  $\sigma_2$ ,  $\delta$  is
$$\begin{array}{ll} \delta(q_1, \sigma_1) = q_2 & \delta(q_1, \sigma_2) = q_1 \\ \delta(q_2, \sigma_1) = q_1 & \delta(q_2, \sigma_2) = q_2 \end{array}$$
- That is encoded as:  
101011 0 101101 0 110101 0 11011011

# DFA Encoding, Continued



- The DFA's set of accepting states:
  - We already encode each state  $q_i$  as  $1^i$
  - Use a list of state codes, separated by 0s
- Finally, the complete DFA:
  - Transition-function string, 00, accepting-state string:  
101011 0 101101 0 110101 0 11011011 00 11

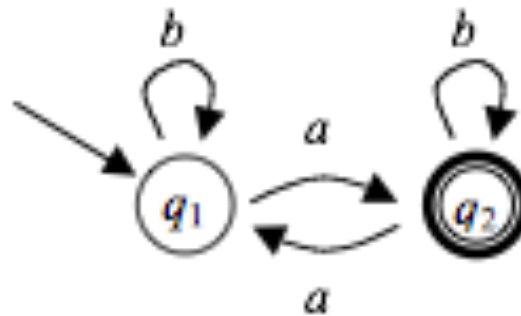
# Simulating a DFA

- We have a way to represent a DFA as a string over  $\{0,1\}$
- Now, we'll show how to construct a TM that simulates any given DFA
  - Given the encoded DFA as input, along with an encoded input string for it
  - Decide whether the given DFA accepts the given string
- We'll use a 3-tape TM...

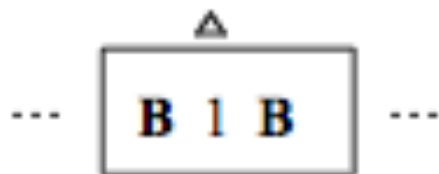
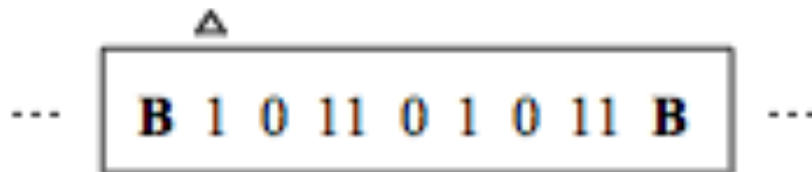
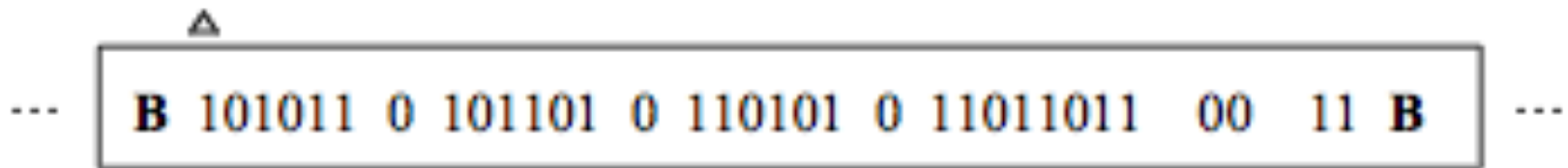
# 3-Tape DFA Simulator

- First tape holds the DFA being simulated
- Second tape holds the DFA's input string
- Third tape hold the DFA's current state  $q_i$ , encoded as  $1^i$  as usual

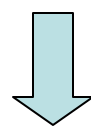
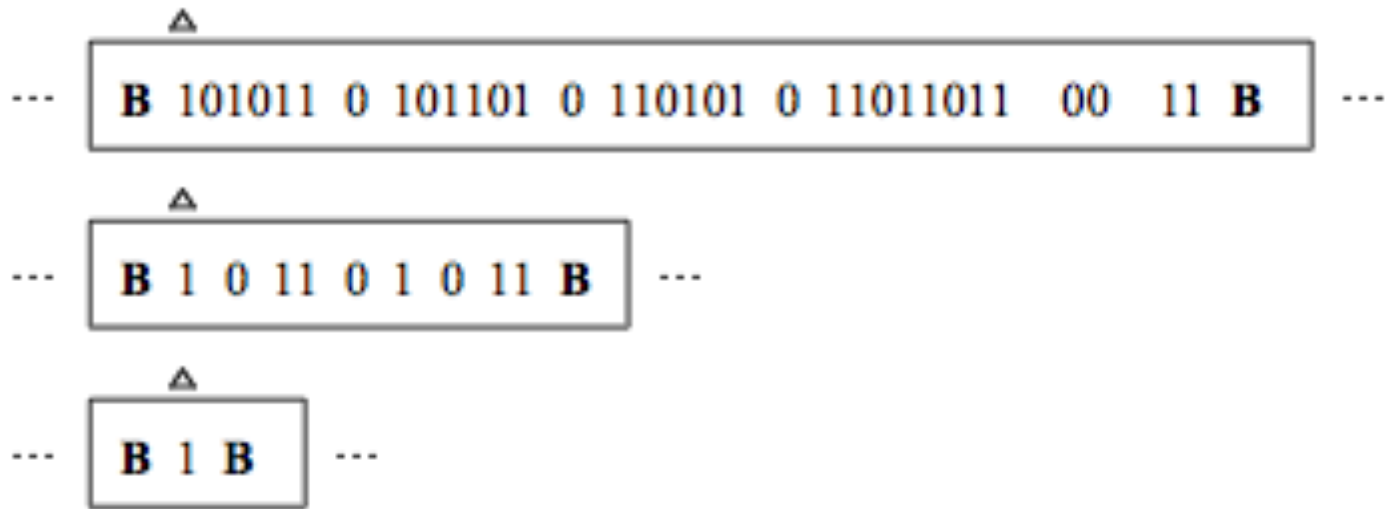
# Example:



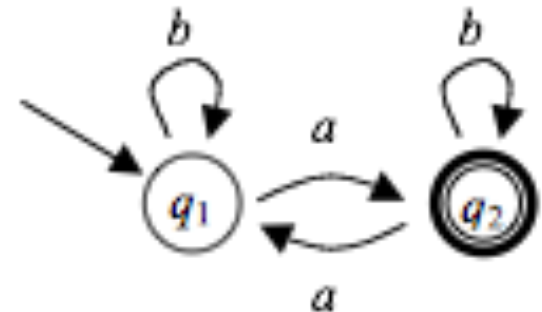
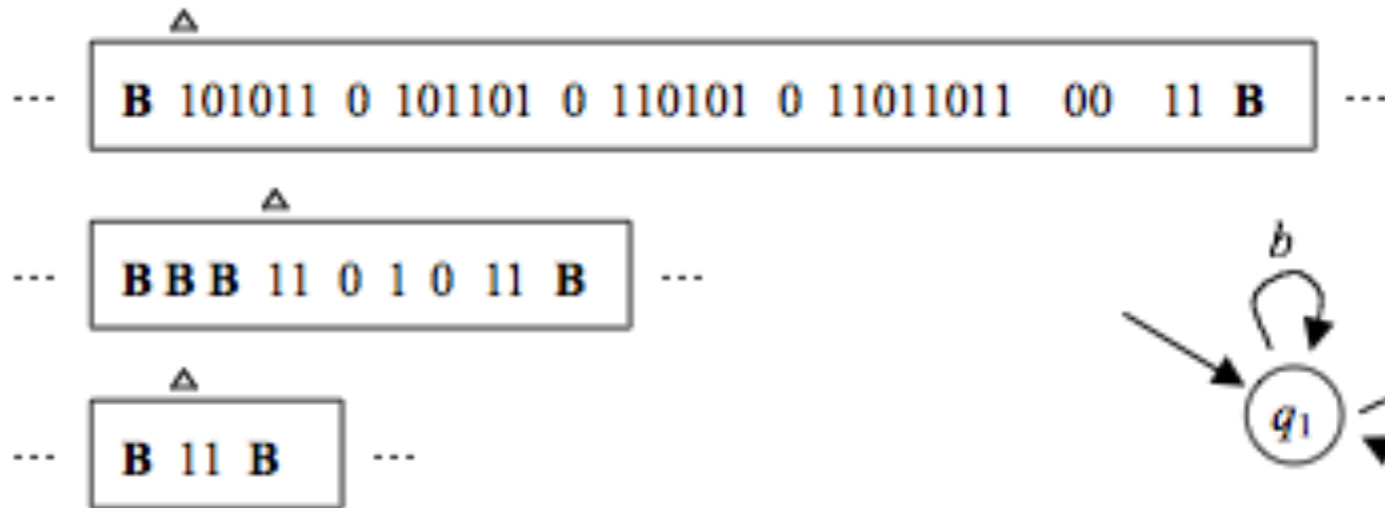
- Initial configuration, in the start state, on input *abab*:



- Each simulated move performs one state transition and erases one encoded input symbol...



First move on *abab*: read *a*,  
go to state  $q_2$





# Strategy

- Step 1: handle termination:
  - If the second tape is not empty, go to step 2
  - If it is empty, the DFA is done; search the list of accepting states (tape 1) for a match with the final state (tape 3)
  - If found, halt and accept; if not, halt and reject
- Step 2: look up move:
  - Search tape 1 for the move  $1^i 0 1^j 0 1^k$  that applies now, where  $1^i$  matches the current state (tape 3) and  $1^j$  matches the current input symbol (tape 2)
- Step 3: execute move:
  - Replace the  $1^i$  on the tape 3 with  $1^k$
  - Write **B** over the  $1^j$  (and any subsequent 0) on tape 2
  - Go to step 1