

Dalarna

A Simplistic Capability-Based Dynamic
Language Design For Data Race Freedom

Kiko Fernandez Reyes¹

Erin Greenwood-Thressman²

James Noble²

Michael Homer²

Isaac Oscar Gariano²

Tobias Wrigstad¹

¹Uppsala University

²Victoria University of Wellington



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Motivation

- Object-oriented dynamic languages in the rise [1]:



- All programming languages have concurrency constructs

Goroutines

Async & Await

Task

Web Workers

Promise & Future

[1] Most Loved Language from Stackoverflow'20 Survey:

<https://insights.stackoverflow.com/survey/2020#technology-most-loved-dreaded-and-wanted-languages-loved>

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Legend

- Dynamic (Red)
- Unsafe (not data race free) (Black)
- Deep copy (Blue)

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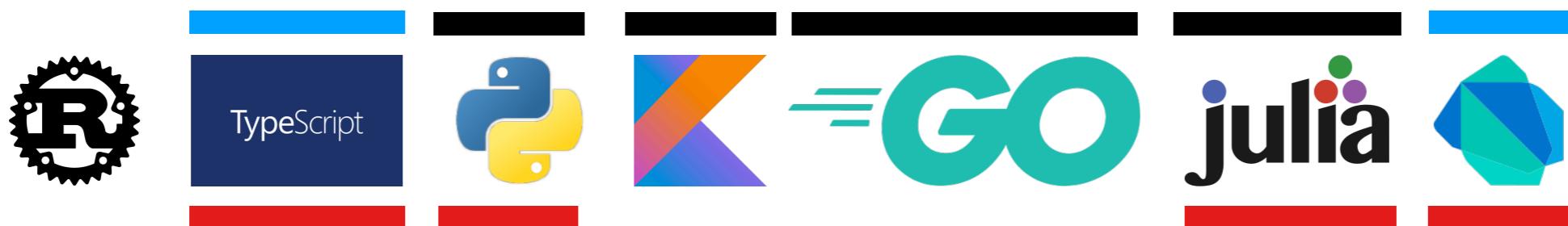
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- Unsafe (not data race free)
- Deep copy

Dalarna Programming Model



- Focused on *dynamic* languages
- Make programs *data race free* by mapping objects to *capabilities*:
 - Remove *data races* in ()
 - Remove *deep copying* where possible
- Embedded Dalarna into the Grace language

Dalarna Programming Model

<i>Programs</i>	$P ::= t$
<i>Objects</i>	$O ::= K \mathbf{obj} \{\overline{F} \overline{M}\}$
<i>Fields</i>	$F ::= f = w$
<i>Methods</i>	$M ::= \mathbf{method} m(x) \{t\}$
<i>Terms</i>	$t ::= w \mid \mathbf{let} x = e \mathbf{in} t$
<i>Expressions</i>	$e ::= w \mid x.f \mid x.f = w$ $\mid x.m(w) \mid K \mathbf{copy} x \mid O$ $\mid \mathbf{spawn}(x)\{t\} \mid \blacksquare_i \iota \mid v$ $\mid x \leftarrow w \mid \leftarrow x$
<i>Variables</i>	$w ::= x \mid \mathbf{consume} x \mid (K) w$
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Dalarna Programming Model

- Capabilities annotations

Object	Effects			
	Read	Write	Alias	Transfer
Imm	•	×	•	•
Iso	•	•	×	•
Local	•	•	•	×
Unsafe	•	•	•	•

Dalarna Programming Model

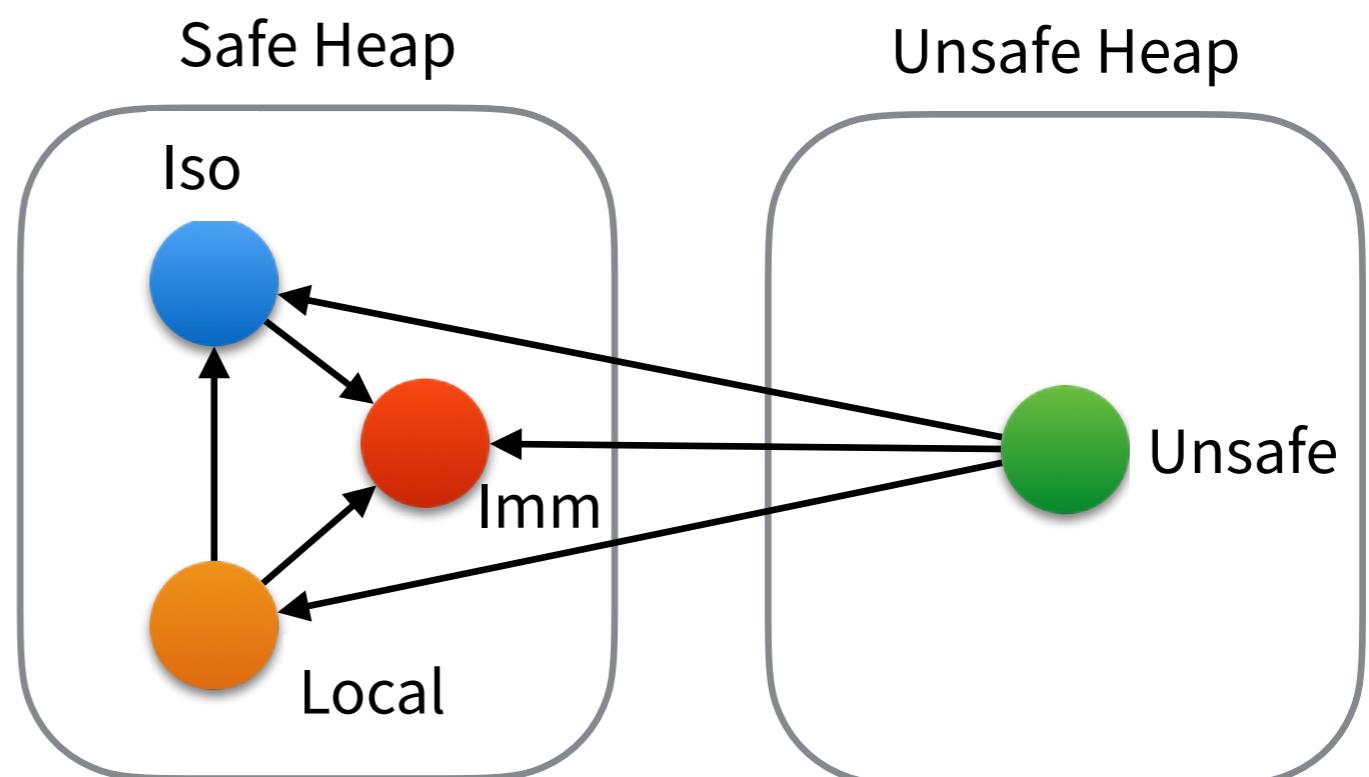
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	Read	Write	Alias	Transfer
Imm	•	✗	•	•
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Local & Imm = Read and Alias

Dalarna Programming Model

- Object capability containment relation



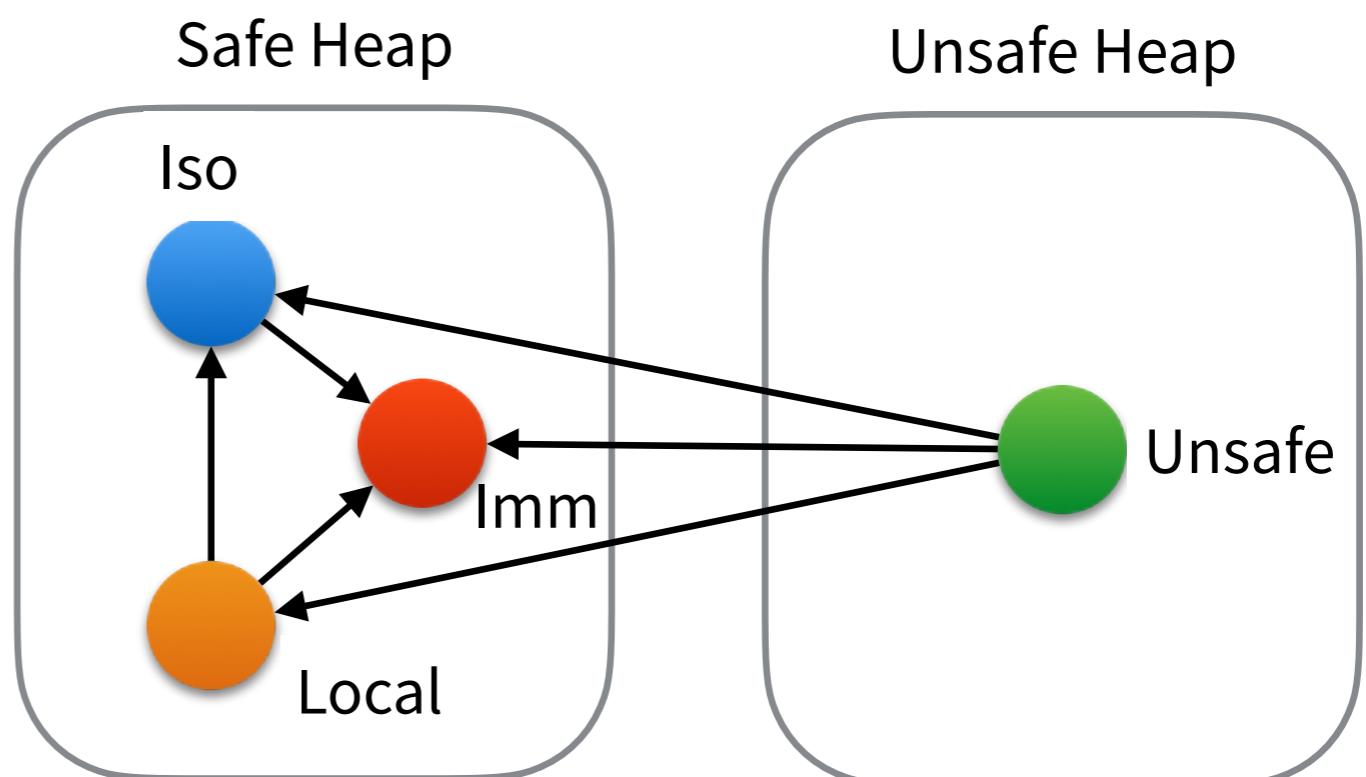
Dalarna Programming Model

- Object capability containment relation

This relation must **hold**:

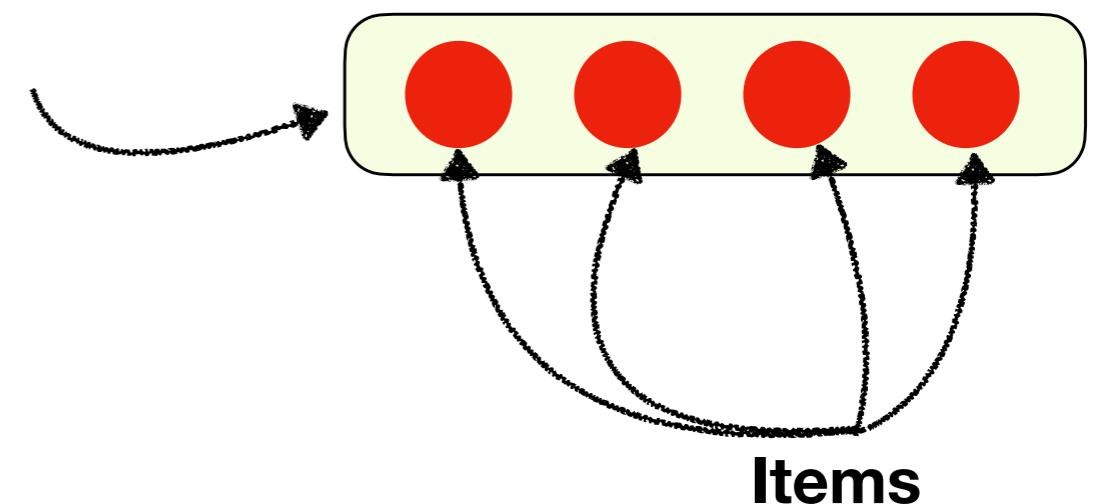
$$\text{unsafe} \leq \text{local} \leq \text{iso} \leq \text{imm}$$

(Reflexive & Transitive)



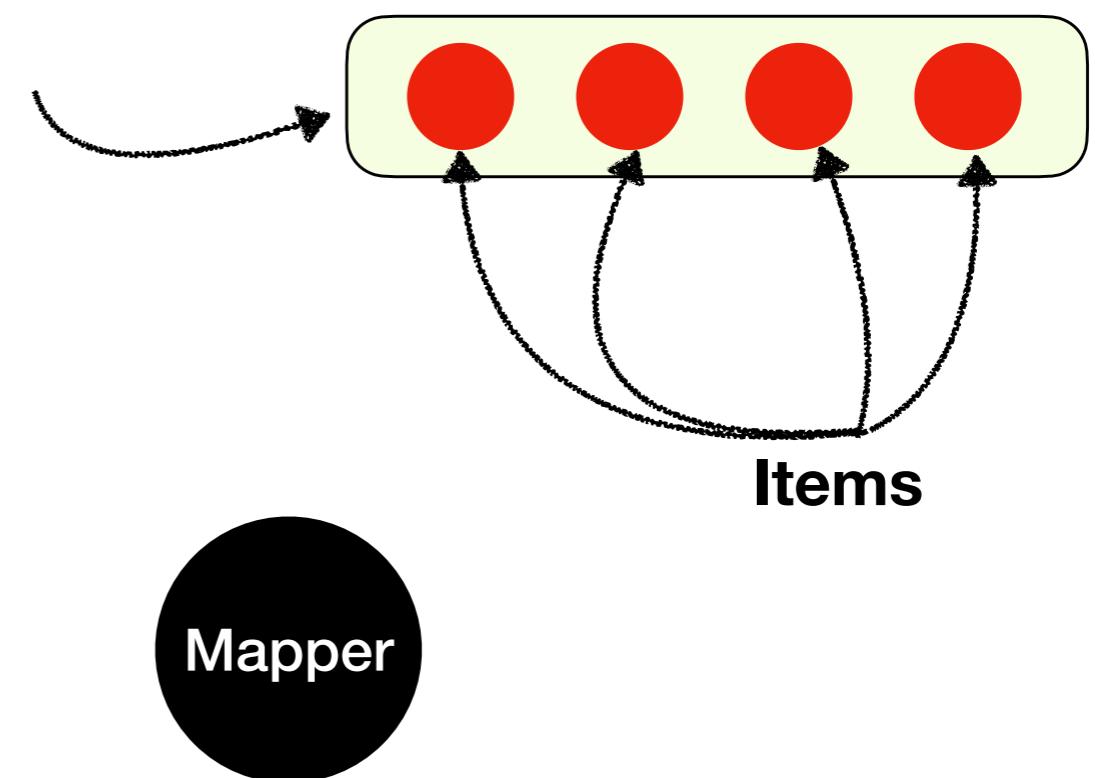
Dalarna Programming Model

Initial code



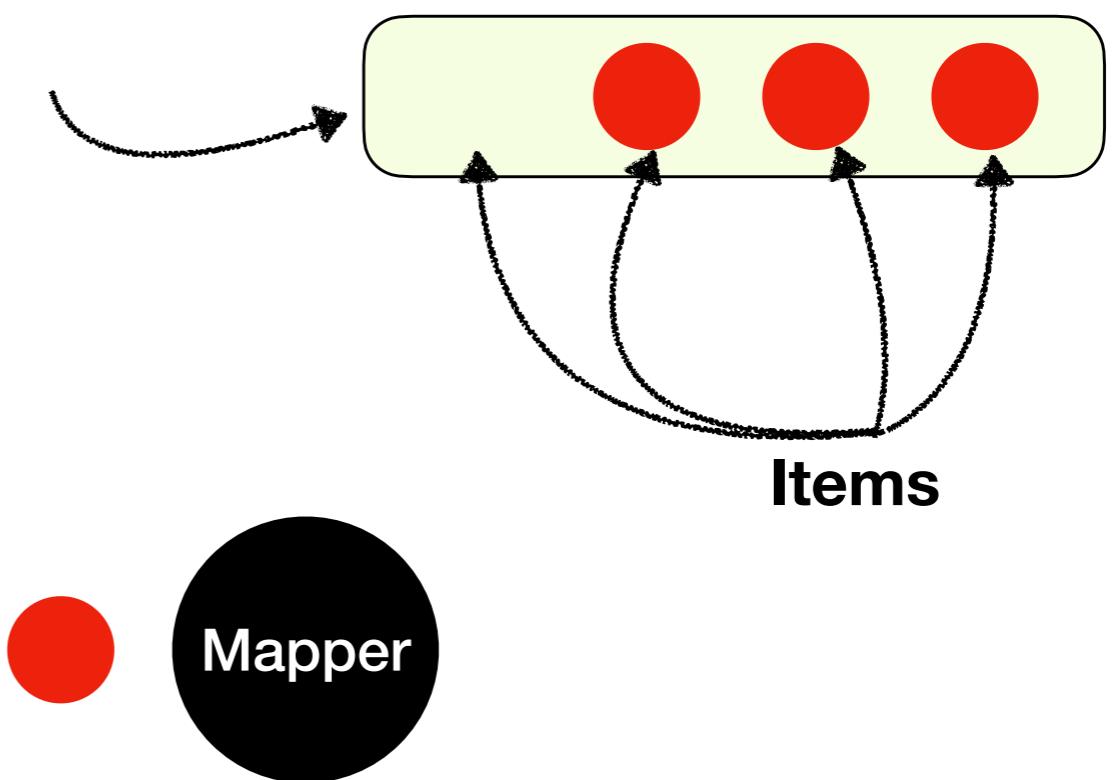
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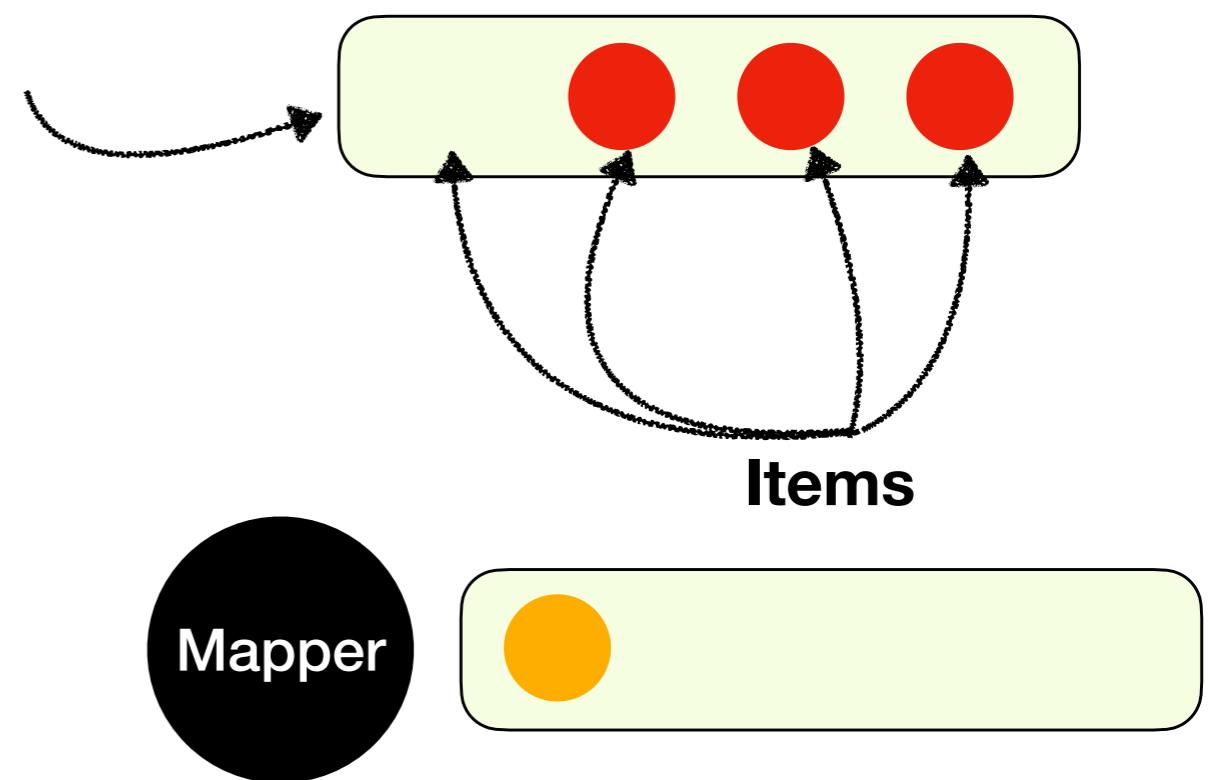
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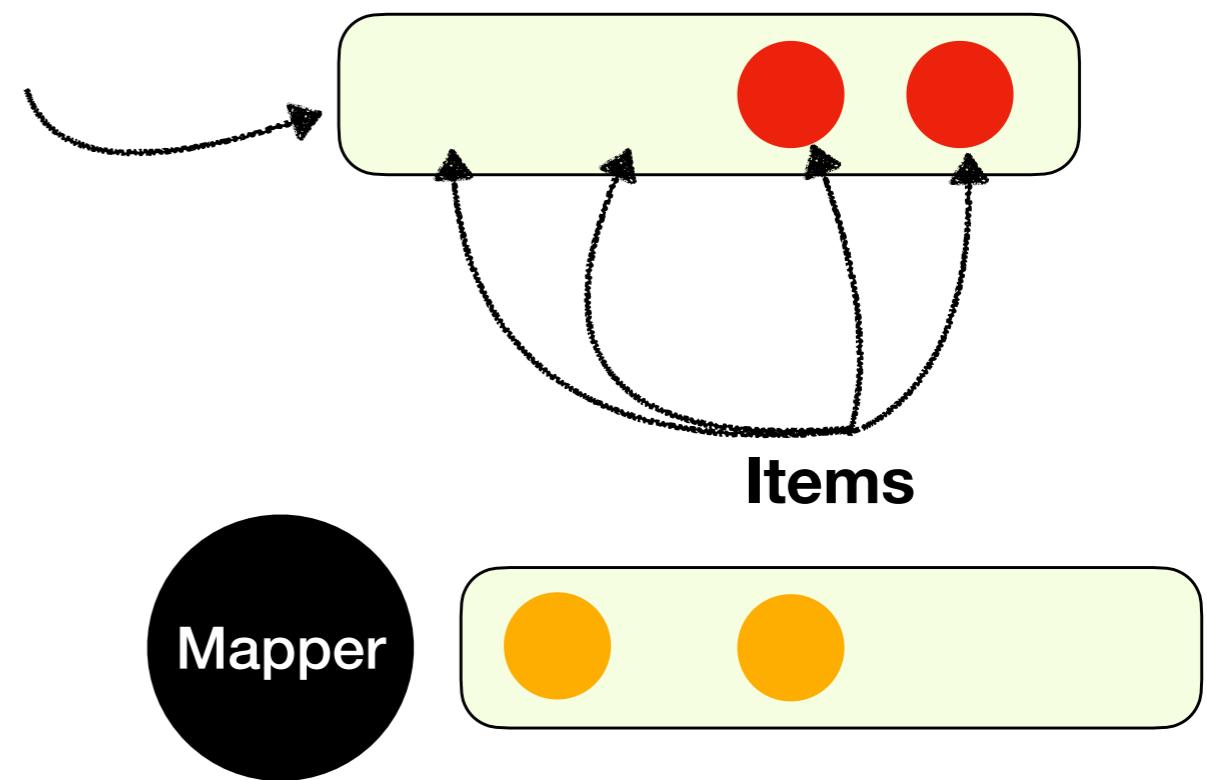
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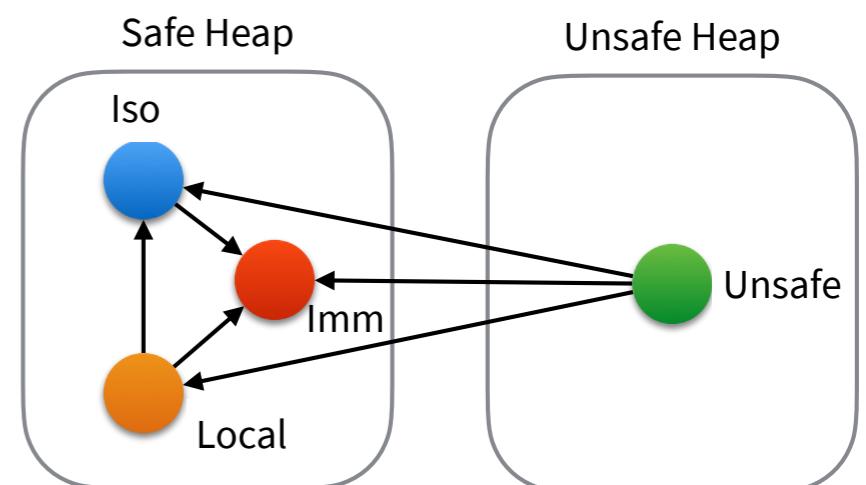
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| ls.append(object 0 {...} )  
  
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
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}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
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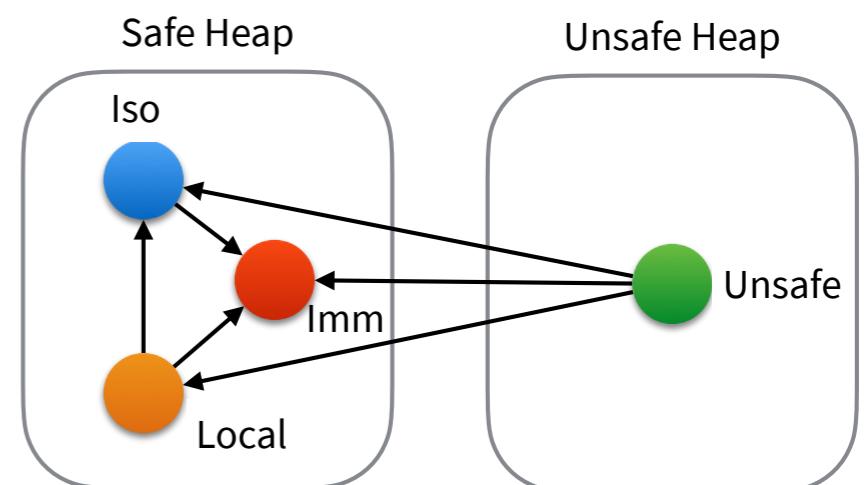
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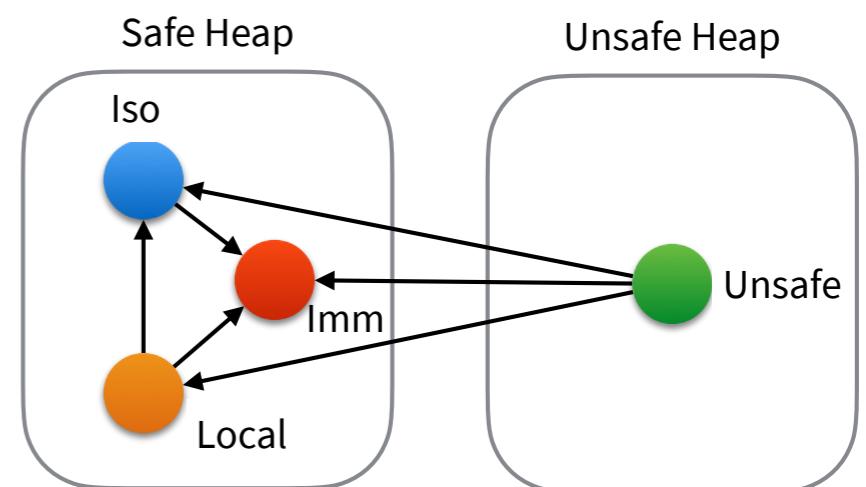
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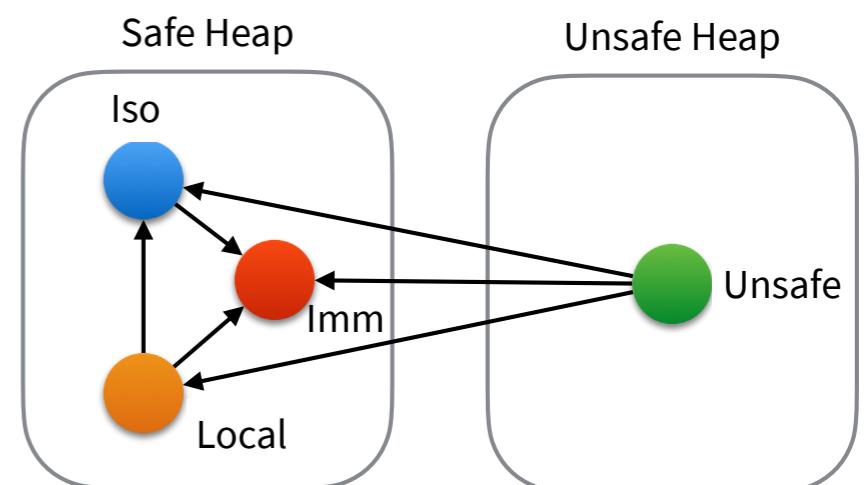
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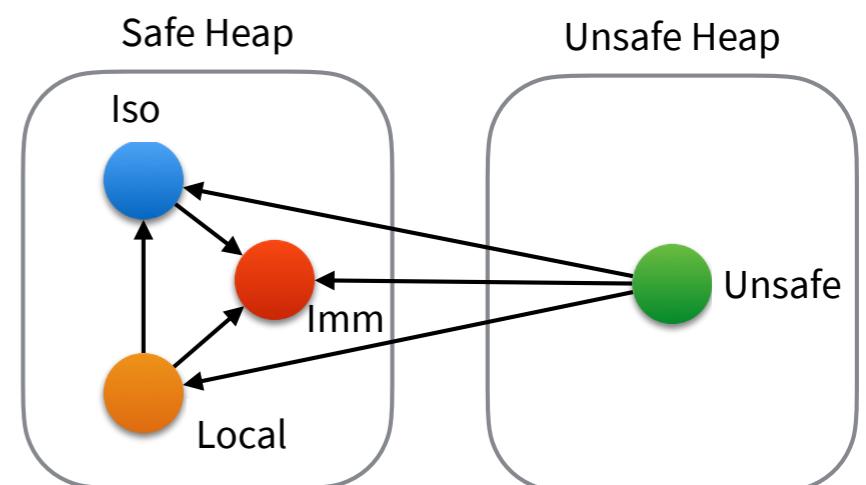
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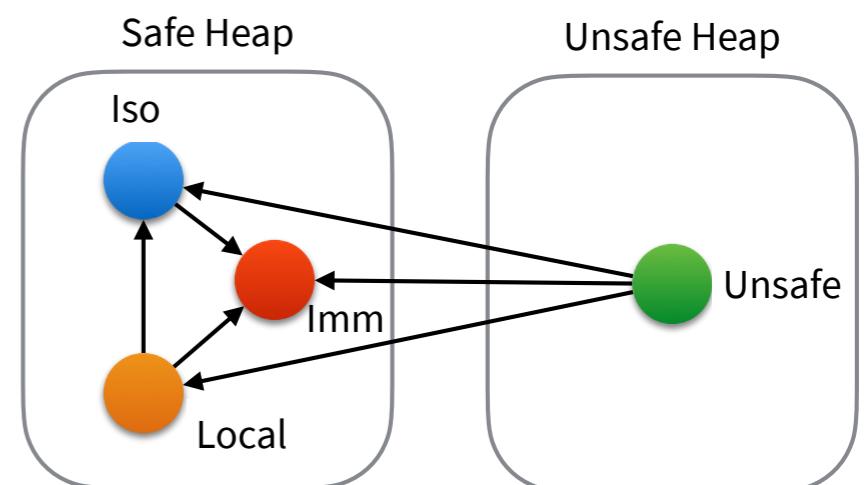
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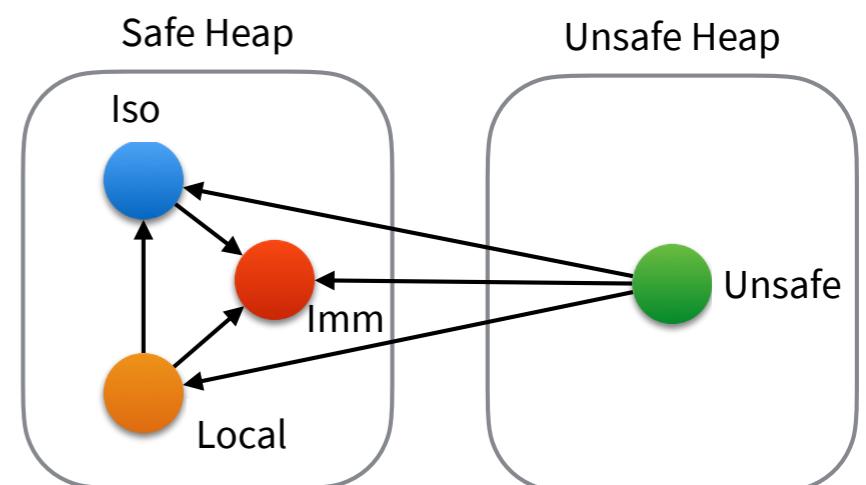
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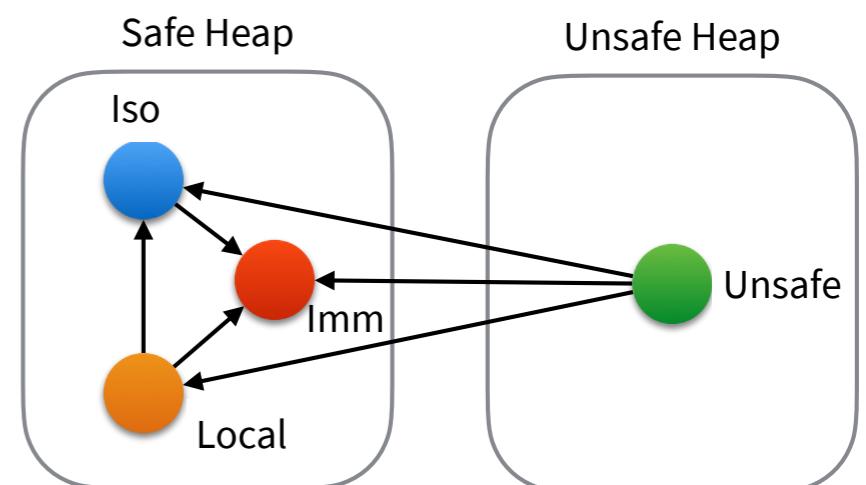
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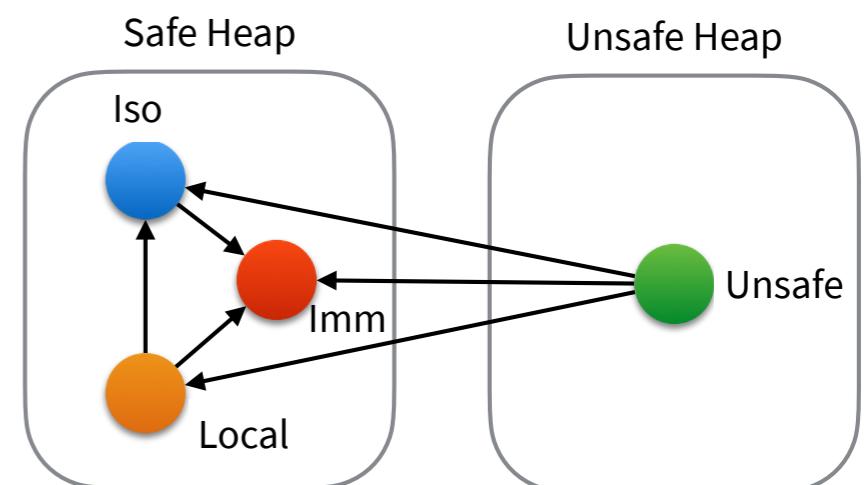
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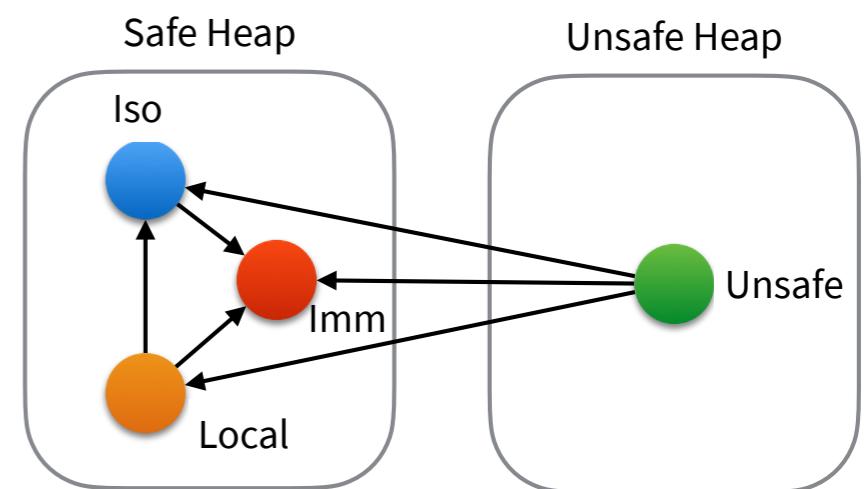
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This code is **not** data race free
Do not share state,
but there is no enforcement

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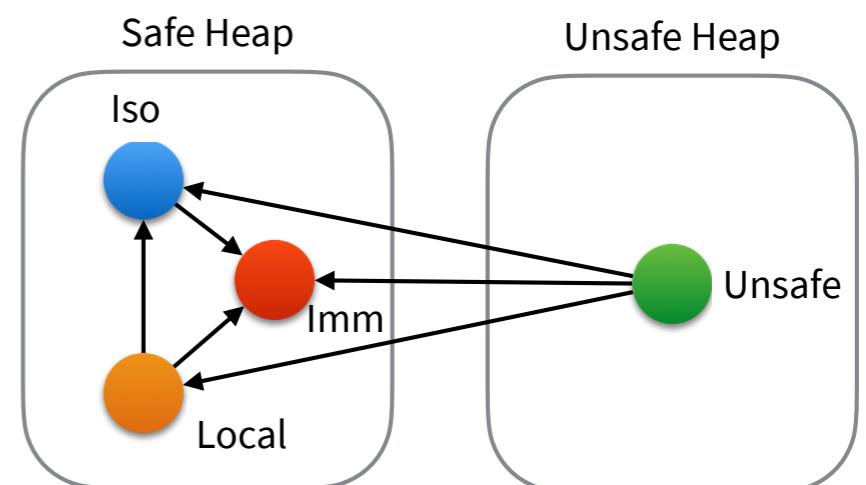
Dalarna Programming Model

~~Initial code~~ Isolated list

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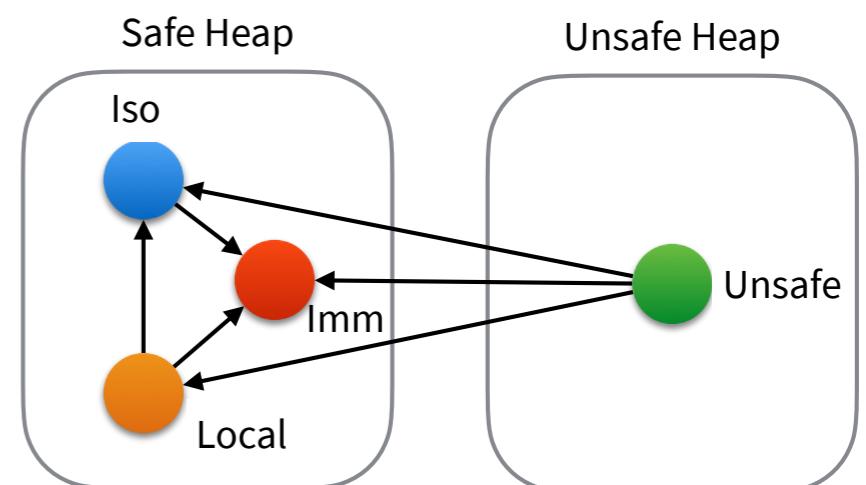
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```



Dalarna Programming Model

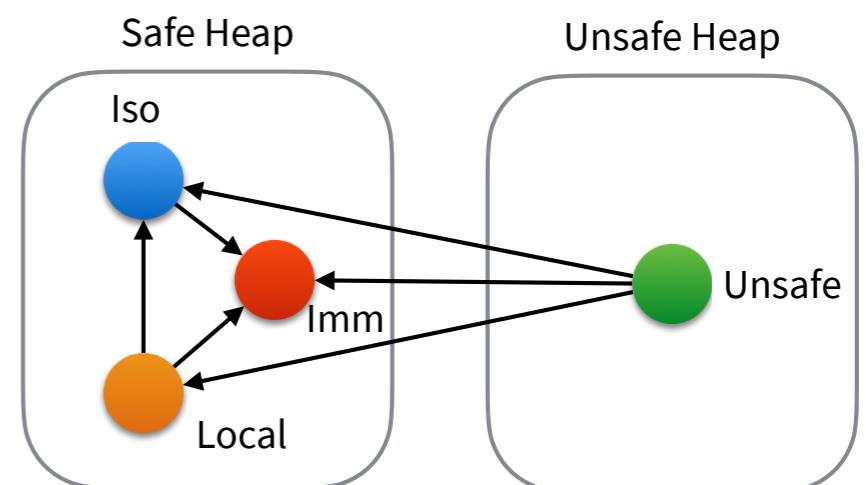
Initial code Isolate list

```
| ls = iso object List {...}  
| ls.append(object 0 ...) X  
  
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Broken containment relation

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Isolate list

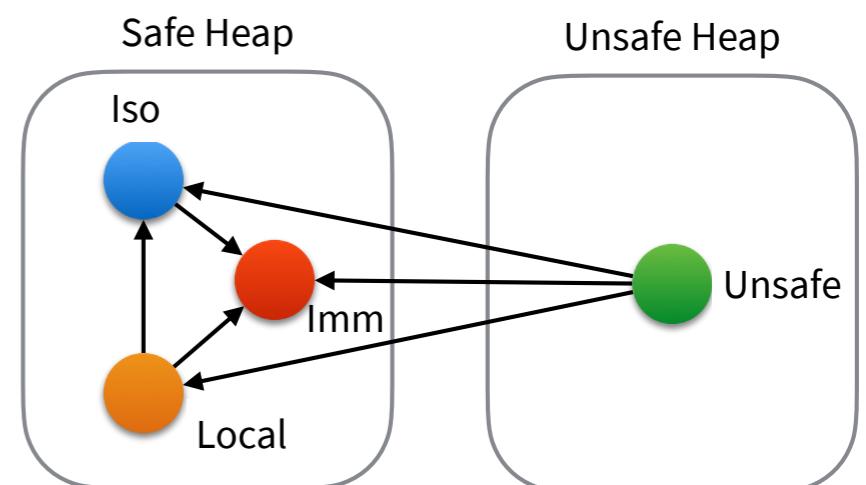
```
| ls = iso object List {...}  
| ls.append(imm object o {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Isolate list

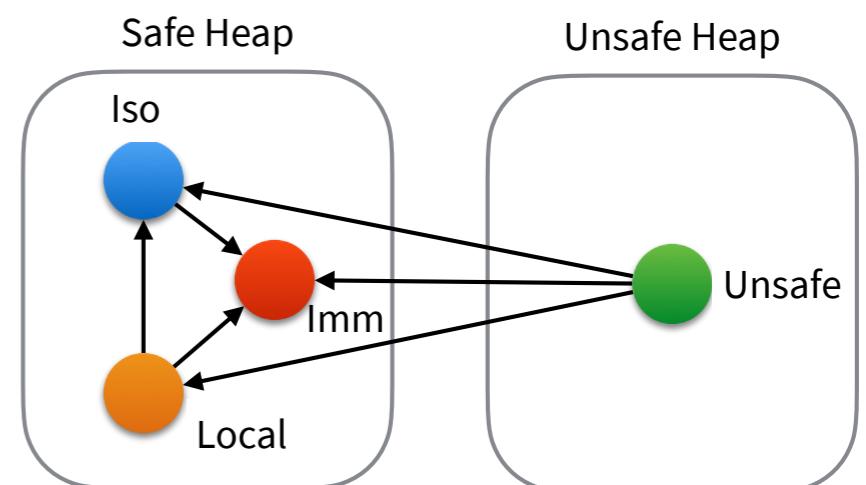
```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Isolate list

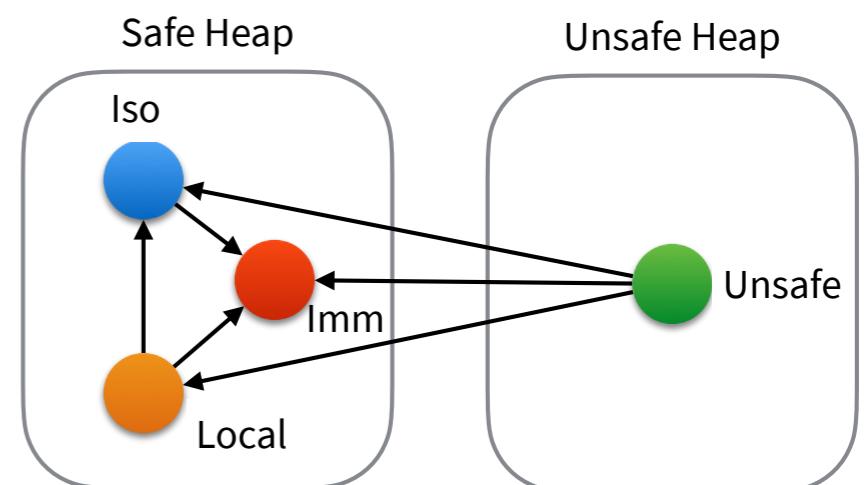
```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Isolate list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



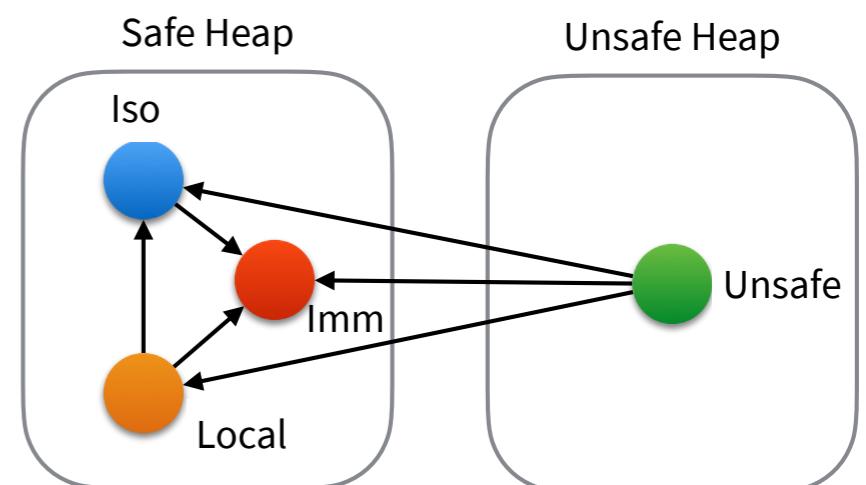
```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
ls.map(mapper)
```

Zero copying



Default object considered with
unsafe capability

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Isolate list

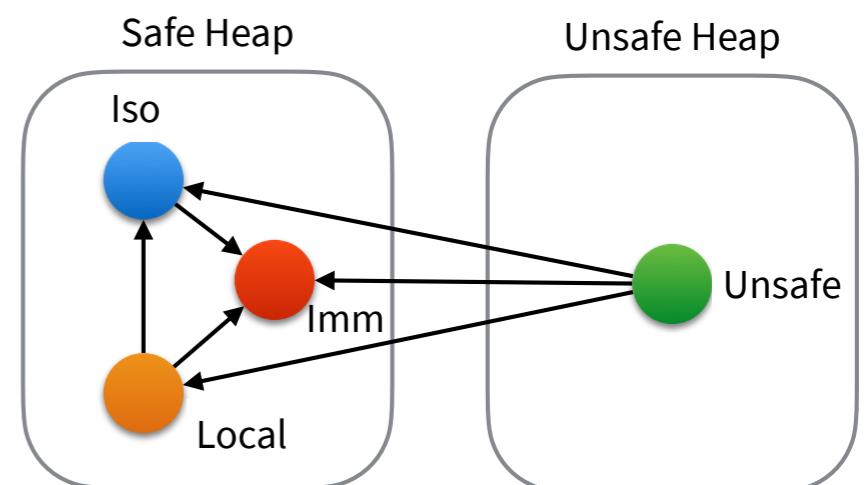
```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object X  
        }  
        ch <- o  
        o.f = ... -- mutate object X  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Isolate list

```
ls = iso object List {...}  
ls.append(imm object 0 {...})
```

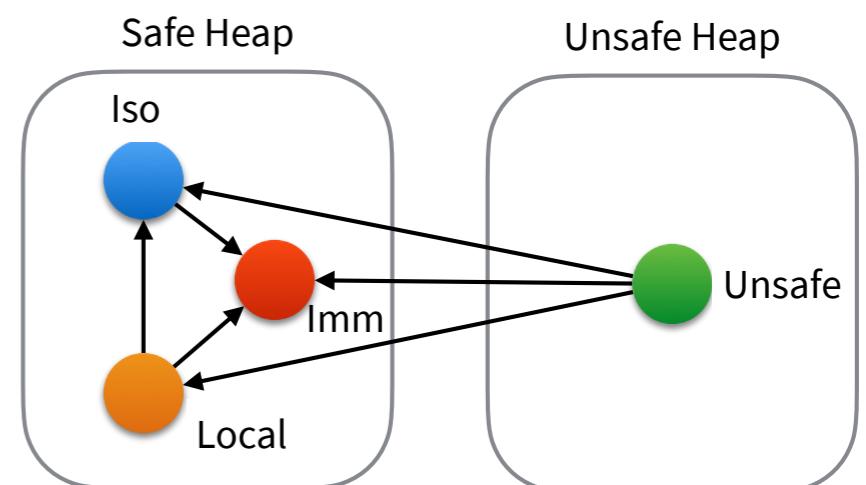


```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object X  
        }  
        ch <- o  
        o.f = ... -- mutate object X  
    }  
}  
  
ls.map(mapper)
```

Immutable objects cannot be mutated

Default object considered with
unsafe capability

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



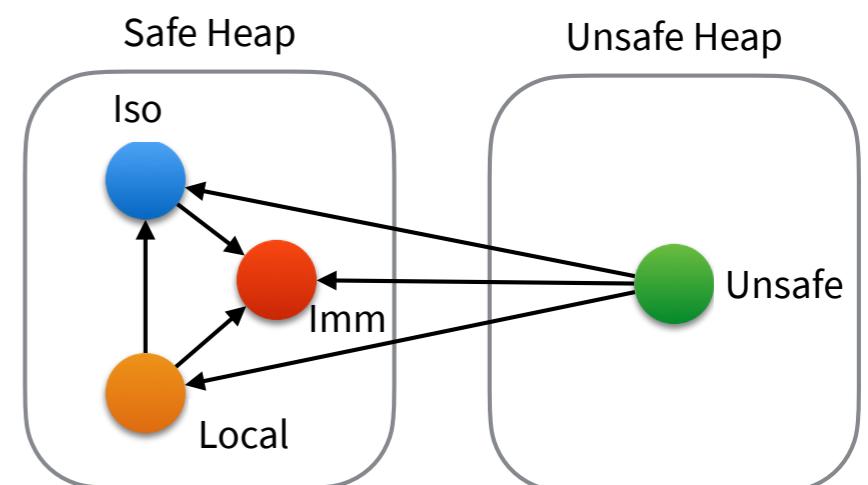
Dalarna Programming Model

Initial code Local list

```
| ls = local object List {...}  
| ls.append(local object 0 {...})  
  
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

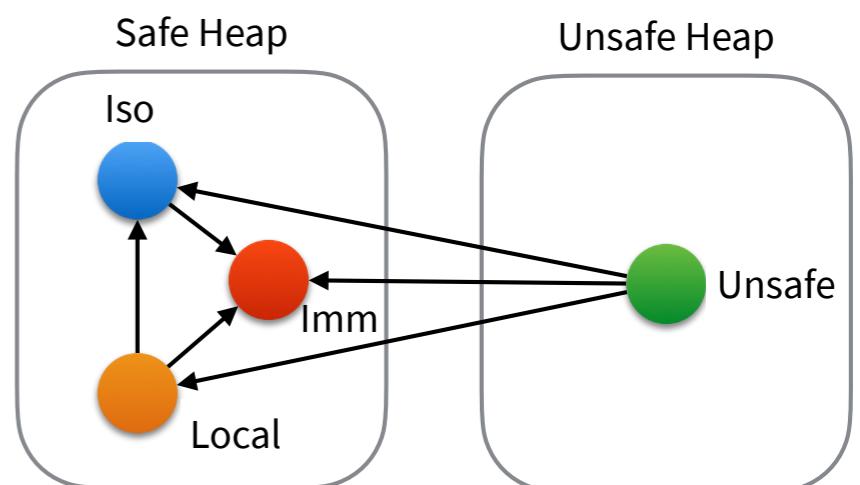
```
| ls = local object List {...}  
| ls.append(local object 0 {...})
```



```
mapper = object Mapper(...){  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

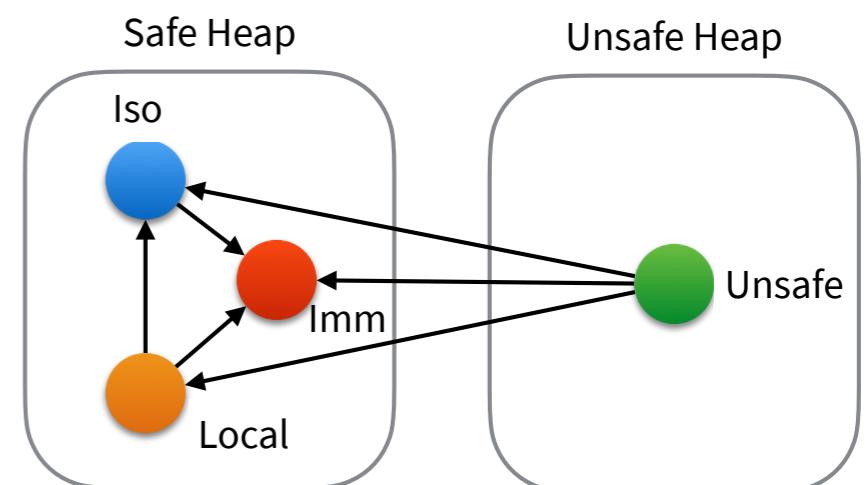
```
ls = local object List {...}  
ls.append(local object 0 {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

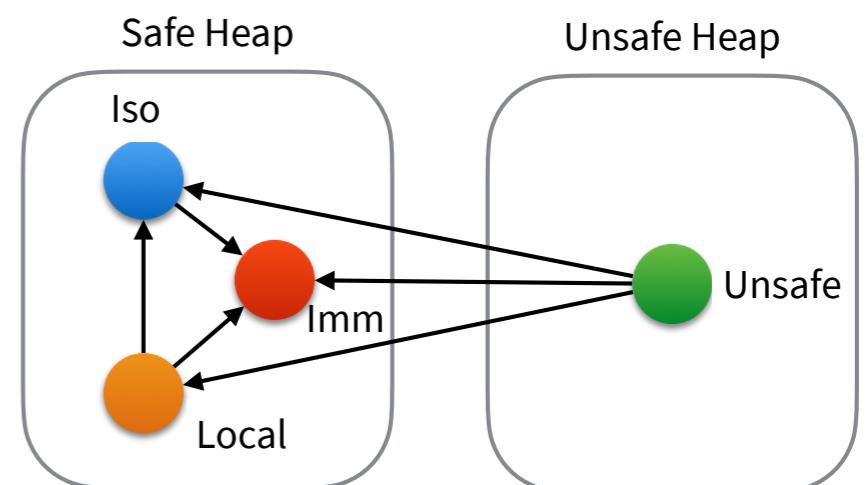
```
ls = local object List {...}  
ls.append(local object 0 {...})
```



```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

```
ls = local object List {...}  
ls.append(local object 0 {...})
```



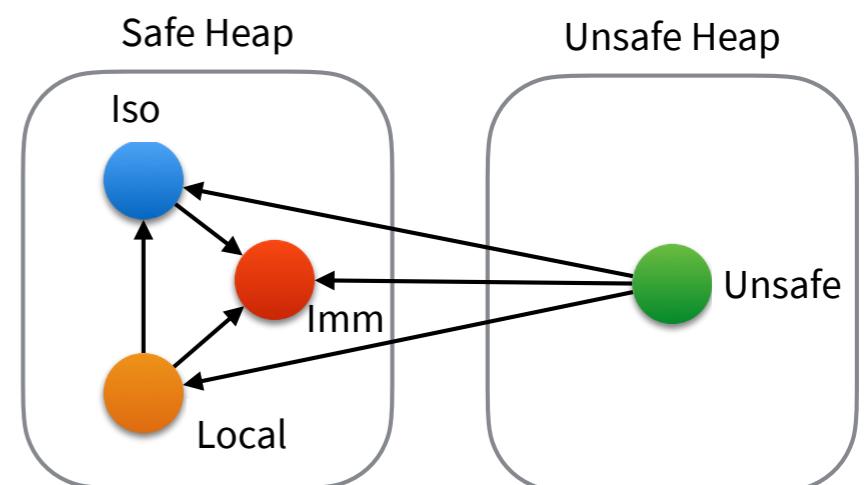
Default object considered with
unsafe capability

```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- o  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```



Local objects cannot move across threads

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

```
ls = local object List {...}  
ls.append(local object 0 {...})
```

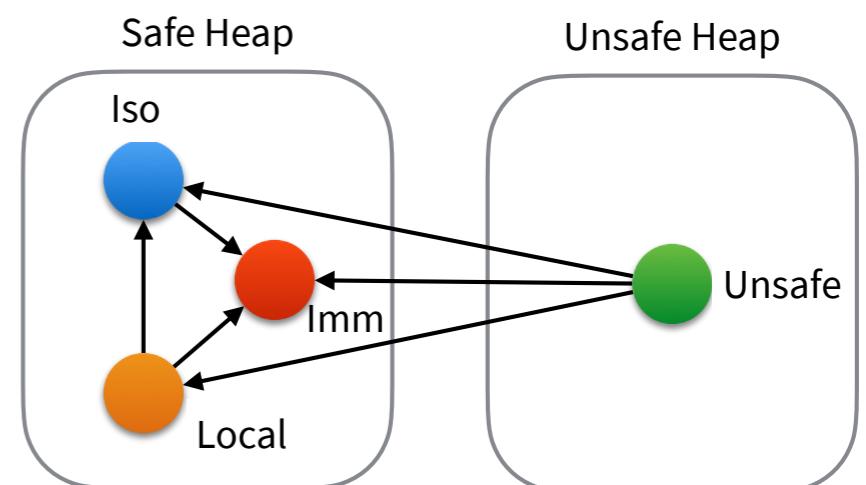


```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- copy(o)  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```



Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



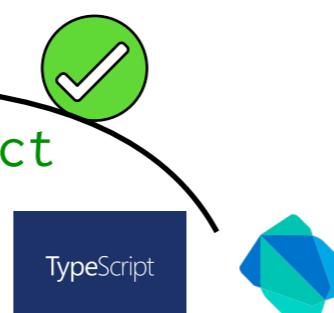
Dalarna Programming Model

Initial code Local list

```
ls = local object List {...}  
ls.append(local object 0 {...})
```

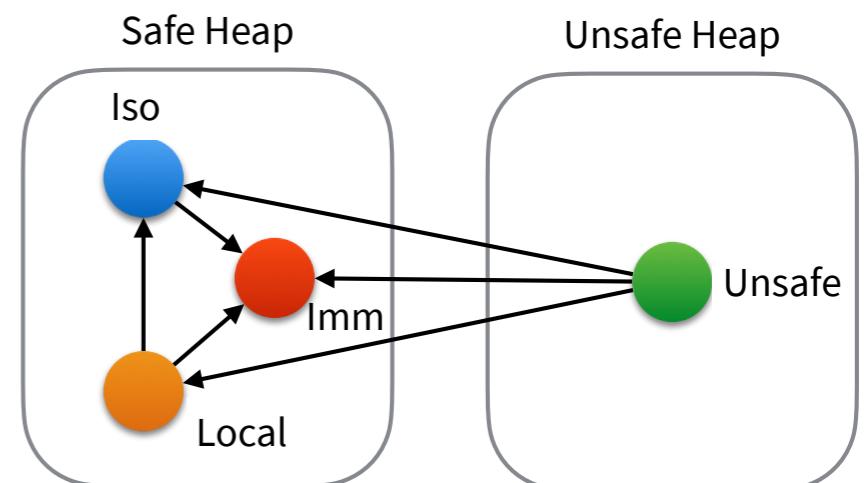


```
mapper = object Mapper(...){  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- copy(o) -- mutate object  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```



Default object considered with
unsafe capability

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

```
ls = local object List {...}  
ls.append(local object 0 {...})
```



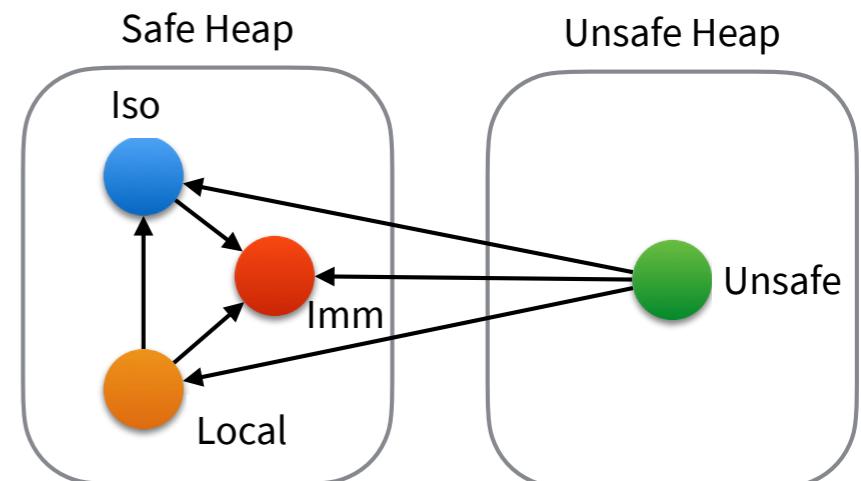
```
mapper = object Mapper(...) {  
    method map(o) {  
        let ch = spawn(ch) {  
            oo <- ch  
            ... -- mutate object  
        }  
        ch <- unsafe copy(o)  
        o.f = ... -- mutate object  
    }  
}  
  
ls.map(mapper)
```

TypeScript



Default object considered with
unsafe capability

```
ls = object List {...}  
      =  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code **Unsafe Local** list

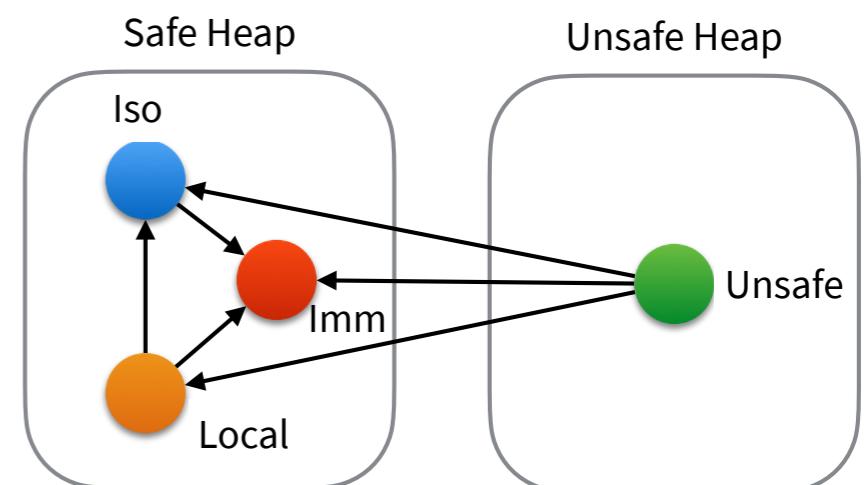
```
| ls = unsafe object List {...}  
| ls.append(local object 0 {...} )
```

```
let ch = spawn(ch) {  
    oo <- ch  
    ...     -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code **Unsafe Local** list

```
| ls = unsafe object List {...}  
| ls.append(local object 0 {...})
```

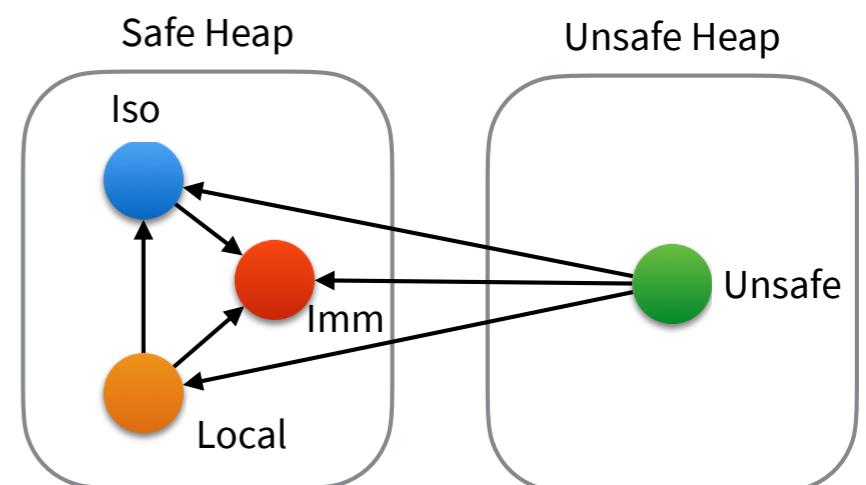


```
let ch = spawn(ch) {  
    oo <- ch  
    ... -- mutate object  
}
```

```
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

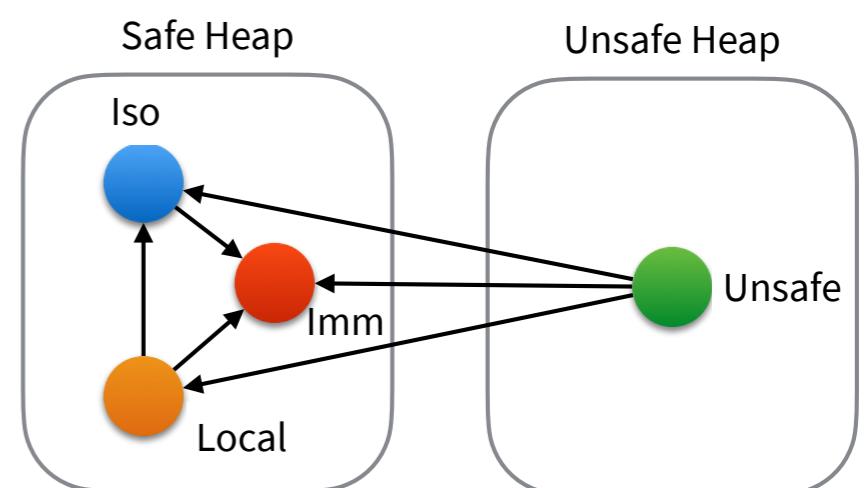
```
ls = unsafe object List {...}  
ls.append(local object 0 {...})
```



```
let ch = spawn(ch) {  
    oo <- ch  
    ... -- mutate object  
}  
  
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

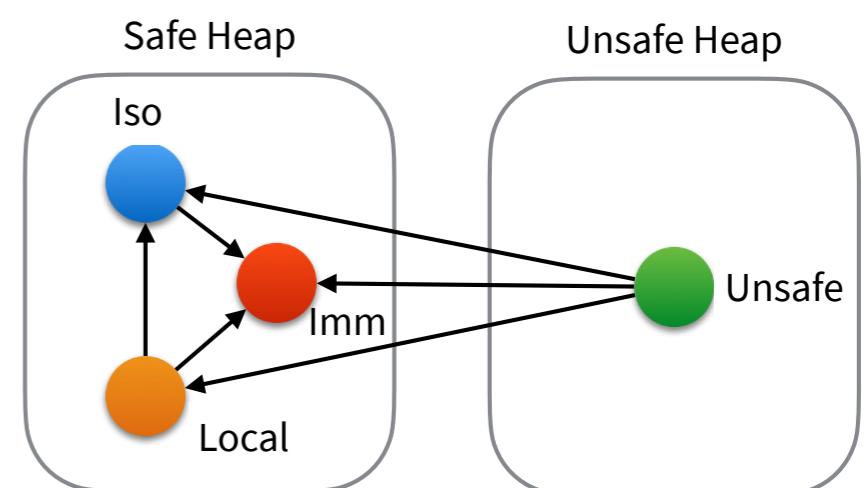
```
ls = unsafe object List {...}  
ls.append(local object 0 {...})
```



```
let ch = spawn(ch) {  
    oo <- ch  
    ... -- mutate object  
}  
  
ch <- ls
```

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

Initial code Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...})
```



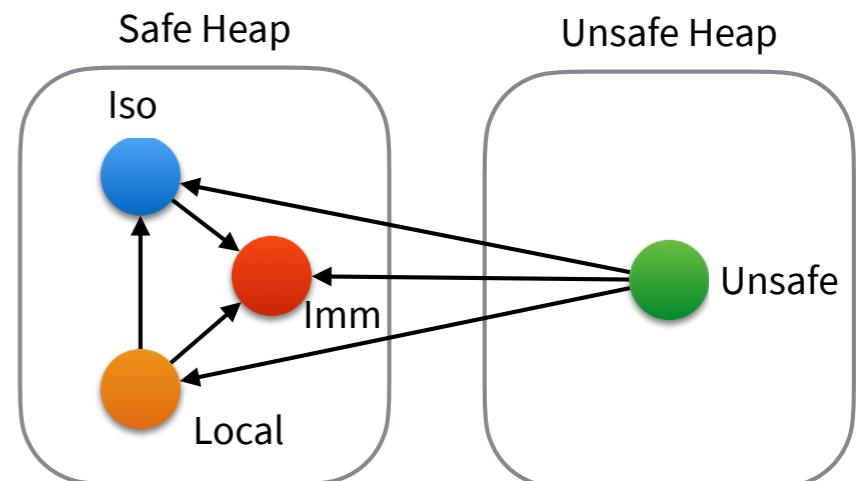
```
let ch = spawn(ch) {  
    oo <- ch  
    ... -- mutate object  
}  
  
ch <- ls
```



Local objects cannot move across threads

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```



Dalarna Programming Model

~~Initial code~~ Local list

```
ls = unsafe object List {...}  
ls.append(local object 0 {...})
```



```
let ch = spawn(ch) {  
    oo <- ch  
    ... -- mutate object  
}  
  
ch <- ls
```



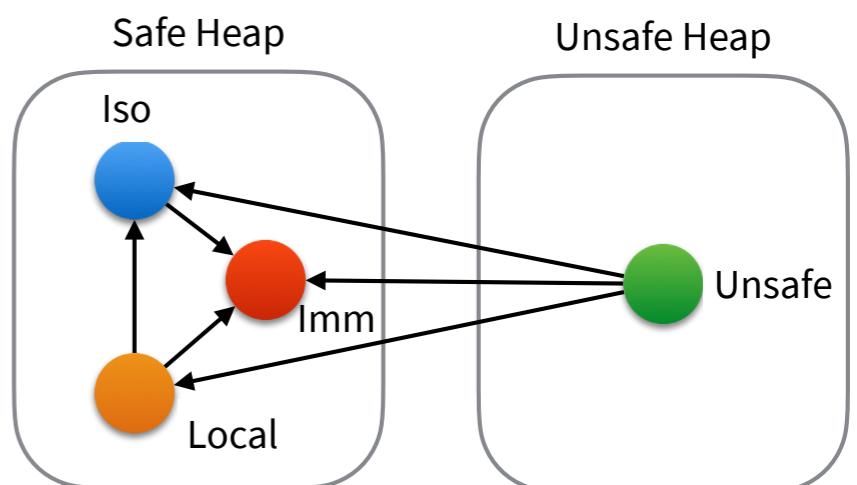
Local objects cannot move across threads

Default object considered with
unsafe capability

```
ls = object List {...}  
      ≡  
ls = unsafe object List {...}
```

Main Point:

*Any code that can lead
to a data race is a runtime error*



Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$

Thread $T ::= t \mid Err$

Thread Err $Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$

Thread $T ::= t \mid Err$

Thread Err $Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

$\frac{(\text{R-LET})}{x \notin \text{dom}(H)}$ $H; \text{let } x = v \text{ in } t \rightsquigarrow H, x \mapsto v; t$	$\frac{(\text{R-VAR})}{H(x) = v \quad \neg \text{isIso}(H, v)}$ $H; E[x] \rightsquigarrow H; E[v]$	$\frac{(\text{R-CONSUME})}{H(x) = v \quad H' = H[x \mapsto \top]}$ $H; E[\text{consume } x] \rightsquigarrow H'; E[v]$	$\frac{(\text{R-FIELD})}{H(H(x)) = _ \mathbf{obj} \{ _ f = v \bar{M} \} \quad \neg \text{isIso}(H, v)}$ $H; E[x.f] \rightsquigarrow H; E[v]$	
$\frac{(\text{R-UPDATE})}{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ \bar{f} = \bar{v} \} f = v' \bar{M}}$ $\neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)$ $H' = H[\iota \mapsto K \mathbf{obj} \{ \bar{f} = \bar{v} \} f = v' \bar{M}]$	$\frac{(\text{R-CASTLoc})}{H(\iota) = K \mathbf{obj} \{ _ _ \}}$ $H; E[(K) \iota] \rightsquigarrow H; E[\iota]$	$\frac{(\text{R-NEW})}{\forall f = v \in \bar{f} = \bar{v}. \text{OkRef}(H, K, v)}$ $\iota \text{ fresh } \quad H' = H, \iota \mapsto K \mathbf{obj} \{ \bar{f} = \bar{v} \bar{M} \}$ $H; E[K \mathbf{obj} \{ \bar{f} = \bar{v} \bar{M} \}] \rightsquigarrow H'; E[\iota]$	$\frac{(\text{R-COPY})}{\mathbf{iso} \notin K}$ $\text{OkDup}(H, K, H(x)) = (H', \iota)$ $H; E[K \mathbf{copy} x] \rightsquigarrow H'; E[\iota]$	
$\frac{(\text{R-CALL})}{x', y' \text{ fresh } \quad H(x) = \iota}$ $H(\iota) = _ \mathbf{obj} \{ _ \bar{M} \mathbf{method} m(y) \{ t \} \}$	$\frac{(\text{R-SPAWN})}{\iota, i \text{ fresh } \quad x \notin \text{dom}(H)}$ $H' = H, x \mapsto \iota, \iota \mapsto \mathbf{chan} \{ i, \emptyset \}$ $H; E[\mathbf{spawn} (x) \{ t \}] \rightsquigarrow H'; E[\iota] t$	$\frac{(\text{R-RECV})}{H(\iota) = \mathbf{chan} \{ i, i' \}}$ $H' = H[\iota \mapsto \mathbf{chan} \{ i, \emptyset \}]$ $H; E[\leftarrow \iota] \rightsquigarrow H'; E[\iota']$		
$\frac{(\text{R-SENDBLOCK})}{H(\iota) = \mathbf{chan} \{ _, \emptyset \} \quad \text{OkSend}(H, v)}$ $i \text{ fresh } \quad H' = H[\iota \mapsto \mathbf{chan} \{ i, v \}]$	$\frac{(\text{R-SENDUNBLOCK})}{H(\iota) = \mathbf{chan} \{ i', v \}}$ $v = \emptyset \vee i \neq i'$	$\frac{(\text{REFCHECK})}{H(\iota) = K' \mathbf{obj} \{ _ _ \}}$ $\text{OkField}(K, K')$	$\frac{(\text{HELPER-OKSEND})}{H(\iota) = K \mathbf{obj} \{ _ _ \}}$ $\neg \exists \iota' \in \text{ROG}(H, \iota). \text{isLocal}(H, \iota')$	$\frac{(\text{HELPER-OKFIELD})}{K_3 = \text{Max}(K_1)}$ $\forall K_4 \in K_2. K_3 \leq K_4$
$H; E[\iota \leftarrow v] \rightsquigarrow H'; E[\blacksquare_i \iota]$	$H; E[\blacksquare_i \iota] \rightsquigarrow H; E[\iota]$	$\text{OkRef}(H, K, \iota)$	$\text{OkSend}(H, \iota)$	$\text{OkField}(K_1, K_2)$

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$

Thread $T ::= t \mid Err$

Thread Err $Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

$\frac{\text{(R-LET)} \\ x \notin \text{dom}(H)}{H; \mathbf{let} \ x = v \ \mathbf{in} \ t \rightsquigarrow H, x \mapsto v; t}$	$\frac{\text{(R-VAR)} \\ H(x) = v \quad \neg \text{isIso}(H, v)}{H; E[x] \rightsquigarrow H; E[v]}$	$\frac{\text{(R-CONSUME)} \\ H(x) = v \quad H' = H[x \mapsto \top]}{H; E[\mathbf{consume} \ x] \rightsquigarrow H'; E[v]}$	$\frac{\text{(R-FIELD)} \\ H(H(x)) = _ \mathbf{obj} \ \{ _ f = v \ \bar{M} \} \quad \neg \text{isIso}(H, v)}{H; E[x.f] \rightsquigarrow H; E[v]}$
$\frac{\text{(R-UPDATE)} \\ H(x) = \iota \quad H(\iota) = K \ \mathbf{obj} \ \{ \bar{f} = \bar{v} \ f = v' \ \bar{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \ \mathbf{obj} \ \{ \bar{f} = \bar{v} \ f = v' \ \bar{M} \}]}{H; E[x.f = v] \rightsquigarrow H'; E[v']}$	$\frac{\text{(R-CASTLoc)} \\ H(\iota) = K \ \mathbf{obj} \ \{ _ _ \}}{H; E[(K) \ i] \rightsquigarrow H; E[\iota]}$	$\frac{\text{(R-NEW)} \\ \forall f = v \in \bar{f} = \bar{v}. \ \text{OkRef}(H, K, v) \\ \iota \ fresh \quad H' = H, \iota \mapsto K \ \mathbf{obj} \ \{ \bar{f} = \bar{v} \ \bar{M} \}}{H; E[K \ \mathbf{obj} \ \{ \bar{f} = \bar{v} \ \bar{M} \}] \rightsquigarrow H'; E[\iota]}$	$\frac{\text{(R-COPY)} \\ \mathbf{iso} \notin K \\ \text{OkDup}(H, K, H(x)) = (H', \iota)}{H; E[K \ \mathbf{copy} \ x] \rightsquigarrow H'; E[\iota]}$
$\frac{\text{(R-CALL)} \\ x', y' \ fresh \quad H(x) = \iota \\ H(\iota) = _ \mathbf{obj} \ \{ _ \bar{M} \ \mathbf{method} \ m(y) \ \{ t \} \}}{H; E[x.m(v)] \rightsquigarrow H, x' \mapsto \iota, y' \mapsto v; E[t[\mathbf{self} = x'][y = y']]}$	$\frac{\text{(R-SPAWN)} \\ \iota, i \ fresh \quad x \notin \text{dom}(H) \\ H' = H, x \mapsto \iota, \iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}}{H; E[\mathbf{spawn} \ (x) \ \{ t \}] \rightsquigarrow H'; E[\iota] \ t}$	$\frac{\text{(R-RECV)} \\ H(\iota) = \mathbf{chan} \ \{ i, i' \} \\ H' = H[\iota \mapsto \mathbf{chan} \ \{ i, \emptyset \}]}{H; E[\leftarrow \ i] \rightsquigarrow H'; E[i']}$	
$\frac{\text{(R-SENDBLOCK)} \\ H(\iota) = \mathbf{chan} \ \{ _, \emptyset \} \quad \text{OkSend}(H, v) \\ \iota \ fresh \quad H' = H[\iota \mapsto \mathbf{chan} \ \{ i, v \}]}{H; E[\iota \leftarrow v] \rightsquigarrow H'; E[\blacksquare_i \ i]}$	$\frac{\text{(R-SENDUNBLOCK)} \\ H(\iota) = \mathbf{chan} \ \{ i', v \} \\ v = \emptyset \vee i \neq i'}{H; E[\blacksquare_i \ i] \rightsquigarrow H; E[\iota]}$	$\frac{\text{(REFCHECK)} \\ H(\iota) = K' \ \mathbf{obj} \ \{ _ _ \} \\ \text{OkField}(K, K')}{\text{OkRef}(H, K, \iota)}$	$\frac{\text{(HELPER-OKSEND)} \\ H(\iota) = K \ \mathbf{obj} \ \{ _ _ \} \\ \neg \exists i' \in \text{ROG}(H, \iota). \ \text{isLocal}(H, i')}{\text{OkSend}(H, \iota)}$
			$\frac{\text{(HELPER-OKFIELD)} \\ K_3 = \text{Max}(K_1) \\ \forall K_4 \in K_2. \ K_3 \leq K_4}{\text{OkField}(K_1, K_2)}$

Dalarna Dynamic Semantics

$$\frac{\begin{array}{c} (\text{R-UPDATE}) \\ H(x) = \iota \quad H(\iota) = K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \} \\ \neg \text{islImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}
 \qquad
 \frac{\begin{array}{c} (\text{REFCHECK}) \\ H(\iota) = K' \text{ obj } \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

$$\frac{\begin{array}{c} (\text{HELPER-OKFIELD}) \\ K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}}{\text{OkField}(K_1, K_2)}$$

Dalarna Dynamic Semantics

$$\frac{\begin{array}{c} (\text{R-UPDATE}) \\ H(x) = \iota \quad H(\iota) = K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ \hline H' = H[\iota \mapsto K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']} \qquad \frac{\begin{array}{c} (\text{REFCHECK}) \\ H(\iota) = K' \text{ obj } \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

$$\frac{\begin{array}{c} (\text{HELPER-OKFIELD}) \\ K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ \hline K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}}{\text{OkField}(K_1, K_2)}$$

Dalarna Dynamic Semantics

$$\frac{\begin{array}{c} \text{(R-UPDATE)} \\ H(x) = \iota \quad H(\iota) = K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \} \\ \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']}
 \qquad
 \frac{\begin{array}{c} \text{(REFCHECK)} \\ H(\iota) = K' \text{ obj } \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

$$\frac{\begin{array}{c} \text{(HELPER-OKFIELD)} \\ K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}}{\text{OkField}(K_1, K_2)}$$

Dalarna Dynamic Semantics

$$\frac{\begin{array}{c} (\text{R-UPDATE}) \\ H(x) = \iota \quad H(\iota) = K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \} \\ \neg \text{isLImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']} \qquad \frac{\begin{array}{c} (\text{REFCHECK}) \\ H(\iota) = K' \text{ obj } \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

$$\frac{\begin{array}{c} (\text{HELPER-OKFIELD}) \\ K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}}{\text{OkField}(K_1, K_2)}$$

Dalarna Dynamic Semantics

$$\frac{\begin{array}{c} (\text{R-UPDATE}) \\ H(x) = \iota \quad H(\iota) = K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \} \\ \neg \text{isLImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\ H' = H[\iota \mapsto K \text{ obj } \{ \overline{f = v} \mid f = v' \overline{M} \}] \end{array}}{H; E[x.f = v] \rightsquigarrow H'; E[v']} \qquad \frac{\begin{array}{c} (\text{REFCHECK}) \\ H(\iota) = K' \text{ obj } \{ _ _ \} \\ \text{OkField}(K, K') \end{array}}{\text{OkRef}(H, K, \iota)}$$

$$\frac{\begin{array}{c} (\text{HELPER-OKFIELD}) \\ K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\ K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\ K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}} \end{array}}{\text{OkField}(K_1, K_2)}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')}
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isLImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{_ _ \} \quad \text{OkField}(K, K')}{\text{OkRef}(H, K, \iota)}
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\} \\
 \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\
 \hline
 H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\}]
 \end{array}$$

$H; E[x.f = v] \rightsquigarrow H'; E[v']$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 H(\iota) = K' \text{ obj } \{__ \} \\
 \text{OkField}(K, K') \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\
 K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\
 \hline
 K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}
 \end{array}$$

OkField(K_1, K_2)

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\} \\
 \neg \text{islImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\
 \hline
 H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\}]
 \end{array}$$

$H; E[x.f = v] \rightsquigarrow H'; E[v']$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 H(\iota) = K' \text{ obj } \{__ \} \\
 \text{OkField}(K, K') \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \\
 K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \\
 \hline
 K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}
 \end{array}$$

$\text{OkField}(K_1, K_2)$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\} \\
 \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v) \\
 \hline
 H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \ f = v' \overline{M}\}]
 \end{array}$$

$H; E[x.f = v] \rightsquigarrow H'; E[v']$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 H(\iota) = K' \text{ obj } \{__ \} \\
 \hline
 \text{OkField}(K, K')
 \end{array}$$

$\text{OkRef}(H, K, \iota)$

(HELPER-OKFIELD)

$$K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1)$$

$$K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2)$$

$$K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}$$

$\text{OkField}(K_1, K_2)$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isLImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{_ _ \} \quad \text{OkField}(K, K')}{\text{OkRef}(H, K, \iota)}
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2) \quad K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}}{\text{OkField}(K_1, K_2)}
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2)}{K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}} \\
 \hline
 \text{OkField}(K_1, K_2)
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2)}{K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}} \\
 \hline
 \text{OkField}(K_1, K_2)
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

unsafe \leq local \leq iso \leq imm

Dalarna Dynamic Semantics

$$\begin{array}{c}
 (\text{R-UPDATE}) \\
 \frac{H(x) = \iota \quad H(\iota) = K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\} \quad \neg \text{isImm}(H, \iota) \quad \text{OkRef}(H, K, v)}{H' = H[\iota \mapsto K \text{ obj } \{\overline{f = v} \quad f = v' \overline{M}\}]} \\
 \hline
 H; E[x.f = v] \rightsquigarrow H'; E[v']
 \end{array}$$

$$\begin{array}{c}
 (\text{REFCHECK}) \\
 \frac{H(\iota) = K' \text{ obj } \{__ \}}{\text{OkField}(K, K')} \\
 \hline
 \text{OkRef}(H, K, \iota)
 \end{array}$$

$$\begin{array}{c}
 (\text{HELPER-OKFIELD}) \\
 \frac{K_3^{\text{Max}}, K_3^{\text{Min}} = \text{Max}(K_1), \text{Min}(K_1) \quad K_4^{\text{Max}}, K_4^{\text{Min}} = \text{Max}(K_2), \text{Min}(K_2)}{K_3^{\text{Max}} \leq K_4^{\text{Max}} \quad K_3^{\text{Min}} \leq K_4^{\text{Min}}} \\
 \hline
 \text{OkField}(K_1, K_2)
 \end{array}$$

$$\begin{array}{c}
 \text{local, iso} = \text{Max(local iso)}, \quad \text{Min(local iso)} \\
 \text{local, imm} = \text{Max(local imm)}, \quad \text{Min(local imm)} \\
 \text{local} \leq \text{local} \quad \text{iso} \leq \text{imm} \quad \checkmark \\
 \hline
 \text{OkField}(\text{local iso}, \text{local imm})
 \end{array}$$

unsafe \leq local \leq iso \leq imm

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$

Thread $T ::= t \mid Err$

Thread Err $Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

(E-NoSuchField)	(E-NoSuchMethod)	(E-NoSuchFieldAssign)	(E-SendBadTargetOrArgument)	
$H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \bar{M} \}$	$m \notin names(\bar{M})$	$H(x.f) = \perp$	$H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}$	
$H(x.f) = \perp$		$H; E[x.f = v] \rightsquigarrow H; Err_N$	$H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N$	
$H; E[x.f] \rightsquigarrow H; Err_N$	$H; E[x.m(v)] \rightsquigarrow H; Err_N$			
(E-RecvBadTarget)	(E-CastError)	(E-AbsentVar)	(E-Consume)	(E-AbsentTarget)
$H(\iota) = _ \mathbf{obj} \{ _ _ \}$	$H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K$	$H(x) = \top$	$H(x) = \top$	$H(x) = \top$
$H; E[\iota \leftarrow \iota] \rightsquigarrow H; Err_N$	$H; E[(K) \iota] \rightsquigarrow H; Err_C$	$H; E[x] \rightsquigarrow H; Err_A$	$H; E[\mathbf{consume} x] \rightsquigarrow H; Err_A$	$H; E[x.m(v)] \rightsquigarrow H; Err_A$
(E-AbsentTargetAccess)	(E-AbsentFieldAssign)	(E-AbsentCopyTarget)	(E-BadFieldAssign)	
$H(x) = \top$	$H(x) = \top$	$H(x) = \top$	$H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \}$	
$H; E[x.f] \rightsquigarrow H; Err_A$	$H; E[x.f = v] \rightsquigarrow H; Err_A$	$H; E[K \mathbf{copy} x] \rightsquigarrow H; Err_A$	$isImm(H, \iota) \vee \neg OkRef(H, K, v)$	
			$H; E[x.f = v] \rightsquigarrow H; Err_P$	
(E-BadInstantiation)	(E-SendingLocal)	(E-AliasIso)	(E-IsoField)	
$\neg \forall v \in \bar{v}. OkRef(H, K, v)$	$H(v) = \mathbf{chan} \{ _ \} \quad \neg OkSend(H, \iota)$	$H(x) = \iota \quad isIso(H, \iota)$	$H(x.f) = \iota \quad isIso(H, \iota)$	
$H; E[K \mathbf{obj} \{ \bar{f} = v \bar{M} \}] \rightsquigarrow H; Err_P$	$H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P$	$H; E[x] \rightsquigarrow H; Err_P$	$H; E[x.f] \rightsquigarrow H; Err_P$	

Dalarna Dynamic Semantics

Configuration $Cfg ::= H, \bar{T}$

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Thread Err $Err ::= Err_N \mid Err_A \mid Err_P \mid Err_C$

(E-NoSuchField)	(E-NoSuchMethod)	(E-NoSuchFieldAssign)	(E-SENDBADTARGETORARGUMENT)	
$H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{ _ \bar{M} \}$	$m \notin names(\bar{M})$	$H(x.f) = \perp$	$H(\iota) = _ \mathbf{obj} \{ _ _ \} \vee H(\iota') = \mathbf{chan} \{ _ \}$	
$H(x.f) = \perp$		$H; E[x.f = v] \rightsquigarrow H; Err_N$		$H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N$
$H; E[x.f] \rightsquigarrow H; Err_N$	$H; E[x.m(v)] \rightsquigarrow H; Err_N$			
(E-RecvBadTarget)	(E-CastError)	(E-AbsentVar)	(E-Consume)	(E-AbsentTarget)
$H(\iota) = _ \mathbf{obj} \{ _ _ \}$	$H(\iota) = K' \mathbf{obj} \{ _ _ \} \quad K' \neq K$	$H(x) = \top$	$H(x) = \top$	$H(x) = \top$
$H; E[\iota \leftarrow \iota] \rightsquigarrow H; Err_N$	$H; E[(K) \iota] \rightsquigarrow H; Err_C$	$H; E[x] \rightsquigarrow H; Err_A$	$H; E[\mathbf{consume} x] \rightsquigarrow H; Err_A$	$H; E[x.m(v)] \rightsquigarrow H; Err_A$
(E-AbsentTargetAccess)	(E-AbsentFieldAssign)	(E-AbsentCopyTarget)	(E-BadFieldAssign)	
$H(x) = \top$	$H(x) = \top$	$H(x) = \top$	$H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{ _ _ \}$	
$H; E[x.f] \rightsquigarrow H; Err_A$	$H; E[x.f = v] \rightsquigarrow H; Err_A$	$H; E[K \mathbf{copy} x] \rightsquigarrow H; Err_A$	$\text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)$	
			$H; E[x.f = v] \rightsquigarrow H; Err_P$	
(E-BadInstantiation)	(E-SendingLocal)	(E-AliasIso)	(E-IsoField)	
$\neg \forall v \in \bar{v}. \text{OkRef}(H, K, v)$	$H(v) = \mathbf{chan} \{ _ \} \quad \neg \text{OkSend}(H, \iota)$	$H(x) = \iota \quad \text{isIso}(H, \iota)$	$H(x.f) = \iota \quad \text{isIso}(H, \iota)$	
$H; E[K \mathbf{obj} \{ \bar{f} = v \bar{M} \}] \rightsquigarrow H; Err_P$	$H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P$	$H; E[x] \rightsquigarrow H; Err_P$	$H; E[x.f] \rightsquigarrow H; Err_P$	

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(E-NoSuchField)

$$\frac{H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$$

(E-NoSuchMethod)

$$\frac{H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{_ \bar{M}\} \quad m \notin names(\bar{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$$

(E-NoSuchFieldAssign)

$$\frac{H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$$

(E-SENDBADTARGETORARGUMENT)

$$\frac{H(\iota) = _ \mathbf{obj} \{__ \} \vee H(\iota') = \mathbf{chan} \{_ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$$

(E-RecvBadTarget)

$$\frac{H(\iota) = _ \mathbf{obj} \{__ \}}{H; E[\iota \leftarrow \iota] \rightsquigarrow H; Err_N}$$

(E-CastError)

$$\frac{H(\iota) = K' \mathbf{obj} \{__ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$$

(E-AbsentVar)

$$\frac{H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$$

(E-Consume)

$$\frac{H(x) = \top}{H; E[\mathbf{consume} \ x] \rightsquigarrow H; Err_A}$$

(E-AbsentTarget)

$$\frac{}{H(x) = \top}$$

(E-AbsentTargetAccess)

$$\frac{H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$$

(E-AbsentFieldAssign)

$$\frac{H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$$

(E-AbsentCopyTarget)

$$\frac{H(x) = \top}{H; E[K \mathbf{copy} \ x] \rightsquigarrow H; Err_A}$$

(E-BadFieldAssign)

$$\frac{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{__ \} \quad \text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$$

(E-BadInstantiation)

$$\frac{\neg \forall v \in \bar{v}. \text{OkRef}(H, K, v)}{H; E[K \mathbf{obj} \{\bar{f} = v \bar{M}\}] \rightsquigarrow H; Err_P}$$

(E-SendingLocal)

$$\frac{H(v) = \mathbf{chan} \{_ \} \quad \neg \text{OkSend}(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$$

(E-AliasIso)

$$\frac{H(x) = \iota \quad \text{isIso}(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$$

(E-IsoField)

$$\frac{H(x.f) = \iota \quad \text{isIso}(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$$

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(E-NoSuchField)

$$\frac{H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$$

(E-NoSuchMethod)

$$\frac{H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{_ \bar{M}\} \quad m \notin names(\bar{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$$

(E-NoSuchFieldAssign)

$$\frac{H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$$

(E-SENDBADTARGETORARGUMENT)

$$\frac{H(\iota) = _ \mathbf{obj} \{__ \} \vee H(\iota') = \mathbf{chan} \{_ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$$

(E-RECVBADTARGET)

$$\frac{H(\iota) = _ \mathbf{obj} \{__ \}}{H; E[\iota \leftarrow \iota] \rightsquigarrow H; Err_N}$$

(E-CASTERROR)

$$\frac{H(\iota) = K' \mathbf{obj} \{__ \} \quad K' \neq K}{H; E[(K) \iota] \rightsquigarrow H; Err_C}$$

(E-ABSENTVAR)

$$\frac{H(x) = \top}{H; E[x] \rightsquigarrow H; Err_A}$$

(E-CONSUME)

$$\frac{H(x) = \top}{H; E[\mathbf{consume} \ x] \rightsquigarrow H; Err_A}$$

(E-ABSENTTARGET)

$$\frac{}{H(x) = \top}$$

(E-ABSENTTARGETACCESS)

$$\frac{H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$$

(E-ABSENTFIELDASSIGN)

$$\frac{H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$$

(E-ABSENTCOPYTARGET)

$$\frac{H(x) = \top}{H; E[K \mathbf{copy} \ x] \rightsquigarrow H; Err_A}$$

(E-BADFIELDASSIGN)

$$\frac{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{__ \} \quad \text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$$

(E-BADINstantiation)

$$\frac{\neg \forall v \in \bar{v}. \text{OkRef}(H, K, v)}{H; E[K \mathbf{obj} \{\bar{f} = v \bar{M}\}] \rightsquigarrow H; Err_P}$$

(E-SENDINGLOCAL)

$$\frac{H(v) = \mathbf{chan} \{_ \} \quad \neg \text{OkSend}(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$$

(E-ALIASISO)

$$\frac{H(x) = \iota \quad \text{isIso}(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$$

(E-IsoFIELD)

$$\frac{H(x.f) = \iota \quad \text{isIso}(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$$

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$$\frac{H(x.f) = \perp}{H; E[x.f] \rightsquigarrow H; Err_N}$$

(E-NoSuchMethod)

$$\frac{H(x) = \iota \quad H(\iota) = _ \mathbf{obj} \{_ \bar{M}\} \quad m \notin names(\bar{M})}{H; E[x.m(v)] \rightsquigarrow H; Err_N}$$

(E-NoSuchFieldAssign)

$$\frac{H(x.f) = \perp}{H; E[x.f = v] \rightsquigarrow H; Err_N}$$

(E-SENDBADTARGETORARGUMENT)

$$\frac{H(\iota) = _ \mathbf{obj} \{__ \} \vee H(\iota') = \mathbf{chan} \{_ \}}{H; E[\iota \leftarrow \iota'], T \rightsquigarrow H; Err_N}$$

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(E-AbsentTargetAccess)

$$\frac{H(x) = \top}{H; E[x.f] \rightsquigarrow H; Err_A}$$

(E-AbsentFieldAssign)

$$\frac{H(x) = \top}{H; E[x.f = v] \rightsquigarrow H; Err_A}$$

(E-AbsentCopyTarget)

$$\frac{H(x) = \top}{H; E[K \mathbf{copy} \ x] \rightsquigarrow H; Err_A}$$

(E-BadFieldAssign)

$$\frac{H(x) = \iota \quad H(\iota) = K \mathbf{obj} \{__ \} \quad \text{isImm}(H, \iota) \vee \neg \text{OkRef}(H, K, v)}{H; E[x.f = v] \rightsquigarrow H; Err_P}$$

(E-BadInstantiation)

$$\frac{\neg \forall v \in \bar{v}. \text{OkRef}(H, K, v)}{H; E[K \mathbf{obj} \{\bar{f} = v \bar{M}\}] \rightsquigarrow H; Err_P}$$

(E-SendingLocal)

$$\frac{H(v) = \mathbf{chan} \{_ \} \quad \neg \text{OkSend}(H, \iota)}{H; E[v \leftarrow \iota], T \rightsquigarrow H; Err_P}$$

(E-AliasIso)

$$\frac{H(x) = \iota \quad \text{isIso}(H, \iota)}{H; E[x] \rightsquigarrow H; Err_P}$$

(E-IsoField)

$$\frac{H(x.f) = \iota \quad \text{isIso}(H, \iota)}{H; E[x.f] \rightsquigarrow H; Err_P}$$

Dalarna Properties

Theorem 4.4 (Dalarna is Data-Race Free Modulo Unsafe Objects). *Any data race in Dalarna directly or indirectly involves an unsafe object. A data race is defined as a read/write or write/write access to an object from different threads without any interleaving synchronisation, which in our case means a transfer of the object from one thread to the other.*

Theorem 4.5 (Preservation). *Given a well-formed configuration $\Gamma; H \vdash t \bar{T}$ and $H; t \bar{T} \rightsquigarrow H'; \bar{T}' \bar{T}$ then, there exists a Γ' s.t. $\Gamma' \supseteq \Gamma$ and $\Gamma'; H' \vdash \bar{T}' \bar{T}$*

Theorem 4.6 (Progress). *Given a well-formed configuration $\Gamma; H \vdash \bar{T}$, then either $\Gamma; H \vdash \bar{T}$ is a terminal configuration (see below) or $H; \bar{T} \rightsquigarrow H'; \bar{T}'$.*

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Future Work

- Addition of capability-based gradual type system to statically reject ill-capable programs

```
ls = iso object List {...}  
ls.append(unsafe object O {...})  Statically rejected
```

- Improve the prototype and benchmark performance

Transient Typechecks Are (Almost) Free. [ECOOP 2019](#)

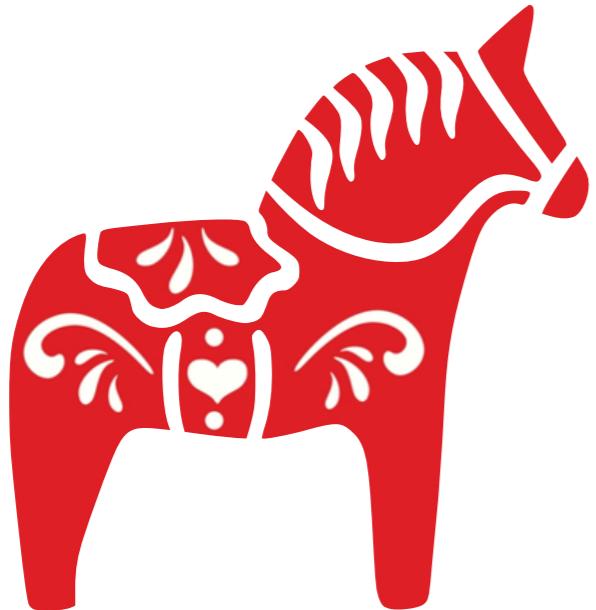
Richard Roberts, Stefan Marr, Michael Homer, James Noble

Conclusions



- Dalarna programming model
 - guarantees data race free programs
 - safe shared memory without deep copying
 - can be embedded in existing languages (Grace)

Image Source



"Dala Horse" Art Prints by JoniandCo | Redbubble

<https://dlpng.com/png/1244869>