

1. Review the following terms.

External sort-merge  
Runs  
Simple nested-loop join  
Page-Oriented Nested Loops Join  
Block nested-loop join  
Sort-merge join  
Hash join  
Equivalence rules  
Statistics estimation  
Size estimation  
Histograms

2. Let relations  $r_1$  and  $r_2$  have the following properties:  $r_1$  has 20,000 tuples,  $r_2$  has 45,000 tuples, 25 tuples of  $r_1$  fit on one page, and 30 tuples of  $r_2$  fit on one page.

(Hint:  $r_1$  needs 800 pages, and  $r_2$  needs 1500 pages. Let us assume M pages of memory and  $M \leq 800$ .)

1) Estimate the I/O cost required using each of the following join strategies for  $r_1 \bowtie r_2$ :

- Simple nested-loop join.
- Page-oriented nested-loop join.
- Sort-merge join. (Assume that one buffer page is needed to hold the evolving output page.)
- Hash join. (Assume that there is no need for recursive partitioning.)

2) Say we have  $M = 100 + 2$  memory buffers, estimate the I/O cost required using the following join strategy for  $r_1 \bowtie r_2$ :

- Block nested-loop join.

3. Given a relation  $r(A, B, C)$  with  $n_r = 10000$  and  $V(A, r) = 500$ .

( $V(A, r)$  means the number of distinct values that appear in the relation  $r$  for attribute  $A$ .)

- Estimate the size of the selection operation  $\sigma_{A=10}(r)$ .
- Assume the range of values for an attribute  $C$  is  $[7, 59]$  and the values are uniformly distributed. Estimate the size of the selection operation  $\sigma_{c<10}(r)$ .