

1. Review the following terms.

- External sort-merge
- Runs
- Simple nested-loop join
- Page-Oriented Nested Loops Join
- Block nested-loop join
- Sort-merge join
- Hash join
- Equivalence rules
- Statistics estimation
- Size estimation
- Histograms

2. Let relations r_1 and r_2 have the following properties: r_1 has 20,000 tuples, r_2 has 45,000 tuples, 25 tuples of r_1 fit on one page, and 30 tuples of r_2 fit on one page.

(Hint: r_1 needs 800 pages, and r_2 needs 1500 pages. Let us assume M pages of memory and $M \leq 800$.)

1) Estimate the I/O cost required using each of the following join strategies for $r_1 \bowtie r_2$:

- a. Simple nested-loop join.
- b. Page-oriented nested-loop join.
- c. Sort-merge join. (Assume that one buffer page is needed to hold the evolving output page.)
- d. Hash join. (Assume that there is no need for recursive partitioning.)

2) Say we have $M = 100 + 2$ memory buffers, estimate the I/O cost required using the following join strategy for $r_1 \bowtie r_2$:

- e. Block nested-loop join.

3. Given a relation $r(A, B, C)$ with $n_r = 10000$ and $V(A, r) = 500$.

($V(A, r)$ means the number of distinct values that appear in the relation r for attribute A .)

- a. Estimate the size of the selection operation $\sigma_{A=10}(r)$.
- b. Assume the range of values for an attribute C is $[7, 59]$ and the values are uniformly distributed. Estimate the size of the selection operation $\sigma_{C < 10}(r)$.