

Imperative programming

Expressions

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Table of contents

- 1 Lexical
- 2 Syntax
- 3 Semantics
- 4 Expression evaluation

Examples

$n + 1$

$3.14 * r * r$

$3 * v[0]$

$x < 3.14$

$3 * (r1 + r2) == \text{factorial}(x)$



Lexical elements

- Literals
- Operators
- Identifiers
- Braces
- Other signs, e.g.: comma

Syntax of function call

```
<function-call> ::= <identifier>()  
                  | <identifier>(<argument-list>)
```

```
<argument-list> ::= <expression>  
                   | <expression>, <argument-list>
```

```
pi()
```

```
factorial(n + m)
```

```
min(0, x + y)
```



Operator arity

- Unary, e.g.: $-x$, $c++$
- Binary, e.g.: $x - y$
- Ternary, e.g.: $x < 0 ? 0 : x$



Operator fixity

- Prefix, e.g.: $++c$
- Postfix, e.g.: $c++$
- Infix, e.g.: $x + y$
- Mixfix, e.g.: $x < 0 ? 0 : x$

Meaningful expressions

- Contained identifiers are declared
($n + 1$ is not meaningful)
- Well typed
(Make sense in C: $3 + 3.14$, "hello" + 1,
Doesn't make sense in C: "hello" * 42)
- Well typed, but still makes no sense
(e.g.: $1 / 0$)



Pure and impure languages

- Impure (e.g. C)
 - Determining an expression's value
 - Side-effect
- Pure (e.g. Haskell)
 - Determining an expression's value
 - Referential transparency

Rules of evaluation

- Fully bracketed expression
 $3 + ((12 - 3) * 4)$
- Precedence: $*$ binds stronger than $+$
 $12 - 3 * 4$
- Left and right associativity
 - In case of operators with the same precedence
 - $3 * n / 2$ means $(3 * n) / 2$ (left assoc. op.)
 - $n = m = 1$ means $n = (m = 1)$ (right assoc. op.)



Assignment

Assignment statement

```
n = 1;
```

Expression with side-effect

```
n = 1
```

Value of expression with side-effect

```
(n = 1) the value is 1
```

Value propagation

```
m = (n = 1)
```



Side-effects

```
printf("%d", n)
n = 1
i *= j
i++
++i
```



Increment/decrement operators

```
int n = 5;
```

```
int i;
```

```
i = ++n;
```

```
n == 6
```

```
i == 6
```

```
int m = 5;
```

```
int j;
```

```
j = m++;
```

```
m == 6
```

```
j == 5
```



Meaning of expressions

- “Normal” value
- Runtime error, e.g.: $5 / 0$ in many languages
- Infinite computation



Lazyness, greedyness

- Greedy: expression in form of $A + B$
- Lazy: expression in form of $A \ \&\& \ B$ (furthermore: $||$, and $?:$)

```
int f() { printf("f"); return 1; }  
int g() { printf("g"); return 2; }
```

```
int i = f() + g();  
/* Output: fg, Result: i == 3 */
```

```
bool b1 = (f() == 0) && (g() == 2);  
/* Output: f, Result: b1 == false */
```

```
bool b2 = (f() == 1) || (g() == 3);  
/* Output: f, Value: b2 == true */
```



Result of lazy and greedy “and” expression

Let \uparrow , \downarrow , \perp and ∞ denote the four possible results for the evaluation of a logical expression: true, false, exception, non-terminating computation. The value of expression $\alpha \wedge \beta$ depending on the values of α and β :

$\alpha \wedge_{\text{lazy}} \beta$	$\beta = \uparrow$	$\beta = \downarrow$	$\beta = \perp$	$\beta = \infty$
$\alpha = \uparrow$	\uparrow	\downarrow	\perp	∞
$\alpha = \downarrow$	\downarrow	\downarrow	\downarrow	\downarrow
$\alpha = \perp$	\perp	\perp	\perp	\perp
$\alpha = \infty$	∞	∞	∞	∞

$\alpha \wedge_{\text{greedy}} \beta$	$\beta = \uparrow$	$\beta = \downarrow$	$\beta = \perp$	$\beta = \infty$
$\alpha = \uparrow$	\uparrow	\downarrow	\perp	∞
$\alpha = \downarrow$	\downarrow	\downarrow	\downarrow	\downarrow
$\alpha = \perp$	\perp	\perp	\perp	\perp
$\alpha = \infty$	∞	∞	∞	∞



Example

Find the index of the first negative element of an array

```
/* Right */
```

```
for (i = 0; i < LENGTH && array[i] >= 0; ++i);
```

```
/* Wrong */
```

```
for (i = 0; array[i] >= 0 && i < LENGTH; ++i);
```

Side-effect in the operands of a lazy operator

```
(n = 1) + (m = 1)
```

```
(n = 1) || (m = 1)
```



Evaluation order of function parameters and operands

Evaluation order is not specified

```
int f() { printf("f"); return 1; }
int g() { printf("g"); return 2; }

void printSum(int a, int b) { printf("%d\n", a + b); }

printf("%d\n", f() + g());      /* Output: fg3 or gf3 */
printSum(f(), g());           /* Output: fg3 or gf3 */
```

Don't do this!

```
int i = 2;
int j = i-- - --i;
printf("%d\n", j);             /* Output: 2 or 0 */
```



Sequence point

- At the end of the whole expression
- At the end of the evaluation of a function's actual parameters
- After the first parameter of a lazy operator
- At a comma operator



Comma operator

```
<expression> ::= ...  
                | <expression>, <expression>
```

- Its result is the same as the last parameter's value
- Low precedence

```
int f() { printf("f"); return 1; }  
int g() { printf("g"); return 2; }  
int h() { printf("h"); return 3; }
```

```
printf("%d\n", (f(), g(), h()));      /* Output: fgh3 */
```

