A workshop on monads with C++14



Meetup C/C++ de Madrid Joaquín Mª López Muñoz <joaquin@tid.es> Madrid, January 2015



Workshop styles

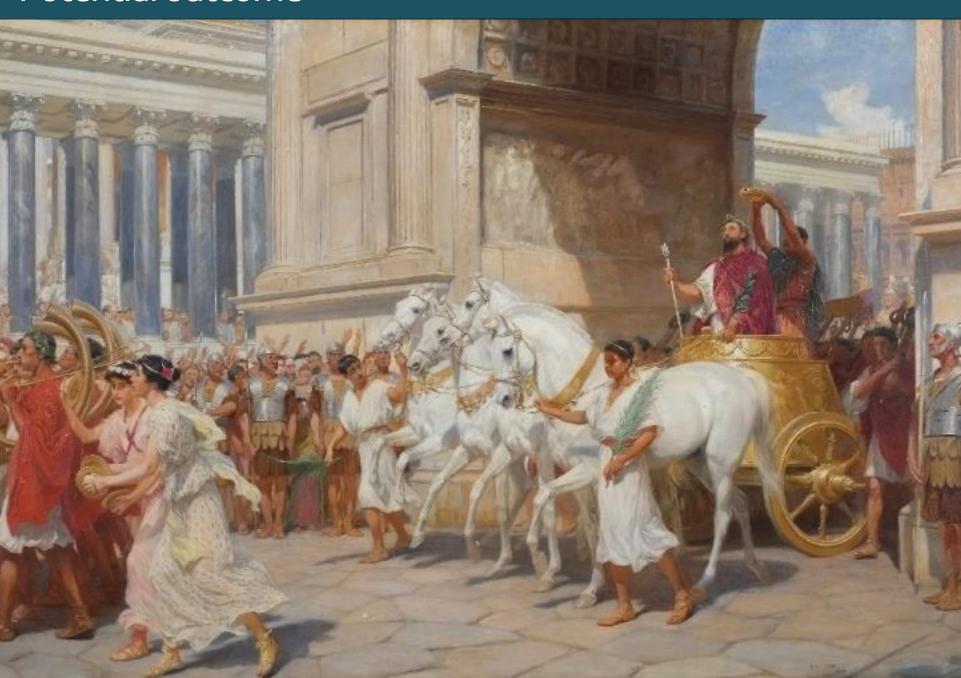


Workshop styles





Potential outcome



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Prerequisites



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- You are expected to be reasonably fluent with some post-2003 C++ stuff
 - auto
 - **std::function**, lambda functions, generic lambda functions
 - Function return type deduction, trailing return type declaration
 - Template template parameters (this is C++03, anyway)
 - decltype, std::declval
- You are encouraged to pre-read about functional programming and monads
 - For instance, go to github.com/joaquintides/usingstdcpp2014
- You are required to bring a computer with
 - Internet access (hopefully provided by host)
 - A C++14 compiler such as GCC 4.9 (-std=c++1y) with a recent Boost distro
 - Alternatively, you can use an online environment such as Coliru (coliru.stacked-crooked.com)

Maybe Monad

optional<T>

```
template<typename T>
struct optional
  optional(const T& x);
  optional(none t);
  const T& get()const;
  T&
           get();
  operator bool()const; // sort of
};
optional<double> inv(double x){
  if(x==0.0)return none;
  else
            return 1.0/x;
optional<double> sqr(double x){
  if(x<0.0)return none;</pre>
  else
           return std::sqrt(x);
}
optional<double> arcsin(double x){
  if(x<-1.0||x>1.0)return none;
  else
                   return std::asin(x);
```

Implement the following composition of functions

```
optional<double > ias(double x)
{
   // return inv(arcsin(sqrt(x))) where it makes sense
}
```

Implement this helper function

```
template<typename F>
optional<double> call(const optional<double>& x, F f)
{
   // call f if x is valid, none otherwise
}
```

Implement ias using call

■ Now **inv** is modified so that it truncates its result to **int**:

```
optional<int> inv(double x)
{
  if(x==0.0)return none;
  else     return int(1.0/x);
}
```

Modify call and ias to cope with type mismatches

```
template<template<typename> class M,typename T>
M<T> mreturn(const T& x)
{
   return x;
}

// must be overloaded
template<template<typename> class M,typename T,typename F>
auto operator>>=(const M<T>& m, F f)->decltype(f(std::declval<const T>()));
```

Signatures

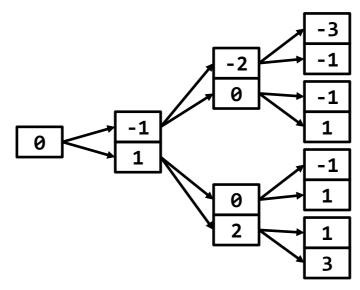
```
return: T \rightarrow M < T >
>>=: (M < T >, T \rightarrow M < T' >) \rightarrow M < T' >
```

List Monad

```
template<typename T> struct list:std::list<T>
{
   using base=std::list<T>;
   using base::base;
   list(const T& x):base(1,x){} // compatibility with mreturn
};

list<int> decinc(int x){return {x-1,x+1};}
```

How do we compose decinc with itself?



- Hint: returning lists is like having multiple potential results
- Hint: look at the signature of >>=

■ Define >>= for list

```
template<typename T,typename F>
auto operator>>=(const list<T>& 1, F f)
{
  // ...
}
```

Calculate

```
((decinc(0)>>=decinc)>>=decinc)
```

Calculate

list<<u>int</u>>{0,1,4,5}>>=decinc

Define transform in terms of >>= (generically!)

```
template<template<typename > class M,typename T,typename F>
auto transform(const M<T>& m,F f)
{
    // ...
}
```

- Hint: $F: T \rightarrow R$, but >>= expects $F': T \rightarrow M < R$ >
- Calculate

```
transform(list<int>{0,1,2},[](int x){return x+1;}
```

What will these be?

```
transform(optional<int>{1},[](int x){return x+1;})
transform(optional<int>{none},[](int x){return x+1;})
```

List comprehension:

```
{x+1 | x <- [0 1 2]}
implemented as

list<int>{0,1,2}>>=[](int x){return mreturn<list>(x+1);}

Implement

{x+y | x <- [0 1 2], y <- [1 2 3]} // [1 2 3 2 3 4 3 4 5]</pre>
```

do notation

```
\{x+y \mid x \leftarrow [0 \ 1 \ 2], y \leftarrow [1 \ 2 \ 3]\}
   Equivalent in do notation (Haskell, beware return is a monadic function, not
   the C++ keyword)
do x <- [0 1 2]
   y <- [1 2 3]
   return (x + y)
equivalent to (Haskell)
[0 \ 1 \ 2] >>= \ \ x \rightarrow
   [1 \ 2 \ 3] >>= \ \ y ->
     return (x + y)
translated to (C++)
list<int>{0,1,2}>>=[&](int x){}
  return list<int>{1,2,3}>>=[&](int y){
    return mreturn<list>(x+y);
  };
```

Do you hate macros?

```
#define DO(var, monad, body)
((monad)>>=[=](const auto& var){ \
  return body; \
})
```

Just for recreational purposes

Implement

```
{x+y | x <- [0 1 2], y <- [1 2 3], even (x+y)}

{x+y | x <- [0 1 2], even x, y <- [1 2 3]}

{x+y | x <- [0 1 2], y <- [1 2 3], even y}

{(x,y) | x <- [0 1 2], y <- [1 2 3]}

{(x,y) | x <- [0 1 2], y <- [1 2 3], (even x) || (even y)}
```

Hint: use a filter helper defined like this

```
template<typename Pred>
auto filter(Pred pred)
{
   // return function object that accepts some x and returns
   // mreturn<list>(x) if pred(x) or an empty list otherwise
}
```

Recap so far



Recap so far

- Monads are type constructors (aka class templates) along with operations return and >>= for $T \rightarrow M < T'$ > function composition
- Monads are a design pattern for composition of functions returning extended types
 - An extended type M < T > is a "container" with 0 or more T values
- Monad contents can **only** be read via >>=
 - This allows for controlled access and execution semantics
- What other value-holding types can be seen as monads?
- do notation emulates imperative style in a functional setting
- What about the **monadic laws**, Cap'n?
 - Trust your heart (just kidding)

Continuation Monad

 In CPS functions are equipped with a continuation or callback object to feed the result to

```
template<typename C>
auto add(int x,int y,C c){return c(x+y);}
template<typename C>
auto square(int x,C c){return c(x*x);}
template<typename C>
auto pyth(int x,int y,C c) // x*x + y*y
  return square(x,[=](int xx){
    return square(y,[=](int yy){
      return add(xx,yy,c);
    });
  });
int main()
 pyth(3,4,[](int r){std::cout<<r<<"\n";});</pre>
```

(example adapted from en.wikibooks.org/wiki/Haskell/Continuation_passing_style)

A twist of syntax

Say we return some suitable continuation object rather than accept one

```
template<typename T>
class cont
public:
 cont(const T& x):run ([=](){return x;}){}
 T run()const{return run ();}
private:
 std::function<T()> run_;
};
auto add(int x,int y){return cont<int>(x+y);}
auto square(int x){return cont<int>(x*x);}
auto pyth(int x,int y){ /* ?? */ }
```

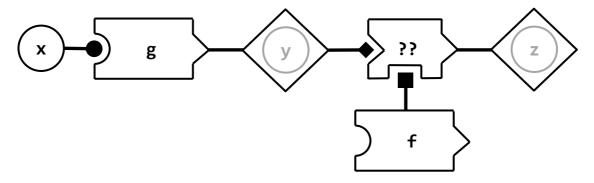
- How do we compose continuation objects?
- Let's look at the compositional interface of cont

Exercise

■ Figure out the implementation of **cont** composition ctor

```
template<typename T>
class cont
{
public:
    // ...
    template<typename F>
    cont(const cont& c,F f):run_(/* figure this out */){}
    //...

private:
    std::function<T()> run_;
};
```



■ I guess you saw it coming ③

```
template<typename T,typename F>
auto operator>>=(const cont<T>& c, F f)
{
  return cont<T>{c,f};
}
auto pyth(int x,int y)
{
  return
  DO(xx,square(x),
  DO(yy,square(y),
    add(xx,yy)
  ));
}
```

- Not a real monad: we've made a simplifying assumption (which?)
- Not the Haskell standard Continuation Monad, either
 - We map T to () $\to T$, Continuation Monad maps T to $(T \to R) \to R$
 - More like a "Future" Monad, actually

Recursively construct a continuation object that computes factorial

```
cont<int> fac(int n);
int main()
  auto f=fac(5);
  std::cout<<"running f\n";</pre>
  auto r=f.run();
  std::cout<<"fac(5)="<<r<<"\n";
```

- Bonus points: do it lazily (i.e. continuation chain is constructed at f.run() time)
- Extra bonus points: do it tail-recursively

Make our monad general

```
template<typename T,typename F>
auto operator>>=(const cont<T>& c, F f);
// f: T -> cont<T'> with T' not necessarily the same as T
```

Extra bonus points: Implement a (restricted) monad with stepwise computation

```
template<typename T>
class cont
{
public:
    // ...
    std::pair<T,boost::optional<cont>> step()const;
    // ...
};
```

Discussion

- Why is this monad interesting?
- The mother of all monads
- The Continuation Monad captures execution flow
- Imperative style in non-imperative settings
- Name application scenarios
- Which C++ entities resemble the Continuation Monad?

I see monads everywhere!

I see monads everywhere!

- Monads are a design pattern for composition of functions returning extended types
- In Haskell, monads + **do** emulate side-effects and imperative style
- In C++, monads are increasingly showing up
 - std::expected proposal
 - then-like extensions to std::future
- Have a look at "Resumable Functions (rev 3)" proposal (N4286)
 - await could potentially be extended to act as a monadic >>=
- My dream: C++ monadic GUI based on functional reactive programming



If you want to learn more about monads...

- "Monad tutorials online" compiles all relevant Internet stuff by time of publication! haskell.org/haskellwiki/Monad_tutorials_timeline
- "All about monads" haskell.org/haskellwiki/All_About_Monads is very comprehensive but assumes you know Haskell
 - Take a look at the monad catalog it provides
- Gabriel Gonzalez's "Haskell for all" haskellforall.com
 - Very readable even if you don't know the language, monads and more
- "Bartosz Milewski's Programming Cafe" bartoszmilewski.com has some monadic stuff as applied to C++
 - "Rediscovering Monads in C++" has many interesting ideas slideshare.net/sermp/rediscovering-monads
 - Bartosz's writing a book on Category Theory for programmers, due 2015
- My blog bannalia.blogspot.com has some (maybe not too friendly) articles on monads for C++ and C++ template metaprogramming

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Thank you

github.com/joaquintides/cpp14monadworkshop

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