

Sets and Data.

- $G = (V, E)$: undirected graph with vertex set V and edge set E .
- $K \in \mathbb{Z}_{\geq 1}$: number of blocks; $\Pi := \{1, 2, \dots, K\}$.
- For $v \in V$, $\text{CN}(v) \subseteq V$: (closed) neighborhood used in the code (given).
- Optional weights $\alpha_{v,i}$ and $\beta_{v,i}$ (aligned with ALPHA and BETA).

Decision Variables.

$$x_{vi} \in \{0, 1\} \quad (v \in V, i \in \Pi), \quad d_{vi} \in \{0, 1\} \quad (v \in V, i \in \Pi).$$

Here, $x_{vi} = 1$ if vertex v is assigned to block i , and $d_{vi} = 1$ if (by the model's logic) vertex v "dominates" block i .

Objective.

$$\min \sum_{v \in V} \sum_{i \in \Pi} \alpha_{vi} x_{vi} + \sum_{v \in V} \sum_{i \in \Pi} \beta_{vi} d_{vi}. \quad (1)$$

Core Constraints. (1) Assignment: each vertex to exactly one block

$$\sum_{i \in \Pi} x_{vi} = 1 \quad \forall v \in V. \quad (2)$$

(2) Nonempty blocks

$$\sum_{v \in V} x_{vi} \geq 1 \quad \forall i \in \Pi. \quad (3)$$

(3) Domination condition for non-edges

$$x_{ui} + d_{vi} \leq 1 \quad \forall i \in \Pi, \forall v, u \in V : v \neq u, \{v, u\} \notin E. \quad (4)$$

(4) Each vertex dominates exactly one block

$$\sum_{i \in \Pi} d_{vi} = 1 \quad \forall v \in V. \quad (5)$$

(5) Block usage order

$$\sum_{v \in V} x_{v,i} \geq \sum_{v \in V} x_{v,i+1} \quad \forall i \in \Pi \setminus \{K\}. \quad (6)$$

Valid Inequalities

VI-1: Minimum fill of the first block.

$$\sum_{v \in V} x_{v,1} \geq \left\lceil \frac{|V|}{K} \right\rceil. \quad (7)$$

VI-2: Per-block upper bound.

$$\sum_{v \in V} x_{v,i} \leq \left\lfloor \frac{|V| - K + i}{i} \right\rfloor \quad \forall i \in \Pi. \quad (8)$$

VI-3: Neighborhood-based lower link on d .

$$d_{vi} \geq \sum_{u \in \text{CN}(v)} x_{ui} - (|\text{CN}(v)| - 1) \quad \forall v \in V, \forall i \in \Pi. \quad (9)$$

VI-4: Neighborhood-based upper link on d .

$$d_{vi} \leq \sum_{u \in \text{CN}(v)} x_{ui} \quad \forall v \in V, \forall i \in \Pi. \quad (10)$$

VI-5: Maximal independent set (MIS) domination cap. Let \mathcal{I} be the family of given maximal independent sets. Then

$$\sum_{v \in S} d_{vi} \leq 1 \quad \forall S \in \mathcal{I}, \forall i \in \Pi. \quad (11)$$

VI-6: MIS-based cover inequality. For every $i \in \Pi$, $v \in V$, and every $S \in \mathcal{I}$ with $v \in S$,

$$\sum_{u \in \text{CN}(v)} x_{ui} \leq |\text{CN}(v)| - 1 + d_{vi} - \sum_{\substack{y \in S \\ y \neq v}} d_{yi}. \quad (12)$$

Variable Domains.

$$x_{vi} \in \{0, 1\}, \quad d_{vi} \in \{0, 1\} \quad \forall v \in V, \forall i \in \Pi. \quad (13)$$