

Today

- LRVN
- RioVista
- Hg

Next week

- Internet scale  
Computing
- (Lesson 9)

# RVM Primitives

## Initialization

- `initialize(options)`
- `map(region, options)`
- `unmap(region)`

## GC to reduce log space

- `flush()`
  - `truncate()`
- done by LRVM automatically

⇓  
Provided for  
app flexibility

## Body of server code

- `begin_xact(tid, restore-mode)`
- `set_range(tid, addr, size)`
- `end_xact(tid, commit-mode)`
- `abort_xact(tid)`

## Miscellaneous

- `query_options(region)`
- `set_options(options)`
- `create_log(options, len, mode)`

⇒ **Simplicity** - small set of primitives

## How the Server uses the primitives

Initialize address space from Ext Segs

```
begin_xact(tid, mode);  
  set_range(tid, base_addr, #bytes);  
  write_metadata m1           // contained in range  
  write_metadata m2           // contained in range  
end_xact(tid, mode);          // or this can be abort
```

# How the Server uses the primitives

Initialize address space from Ext Segs

```
begin_xact(tid, mode);  
set_range(tid, base_addr, #bytes);
```

LRVM creates  
undo record

// contained in range

// contained in range

// or this can be about

restore from  
undo record

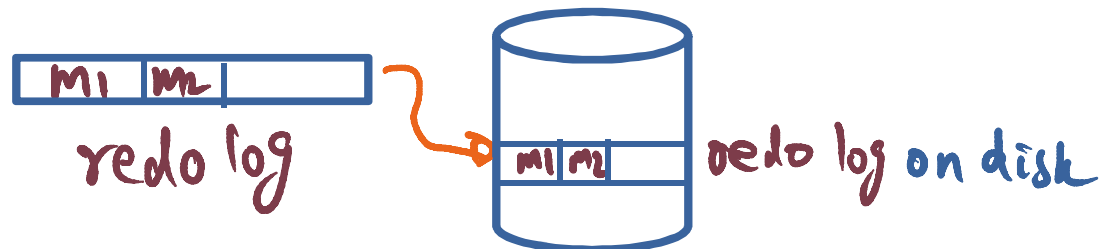
No action  
by  
LRVM

```
{ write metadata m1  
  write metadata m2  
}
```

LRVM → end\_xact(tid, mode);  
LRVM gets rid  
of undo  
record

LRVM creates redo log in memory

- flush to disk sync or later depending on mode



## Opportunities for Server to optimize transactions

no-restore mode in begin\_xact

— no need to create in-mem undo record

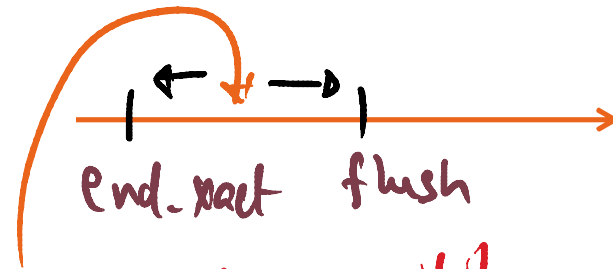
no-flush mode in end\_xact

— no need to sync flush redo log to disk

⇒ lazy persistence

⇒ upshot?

⇒ Window of Vulnerability

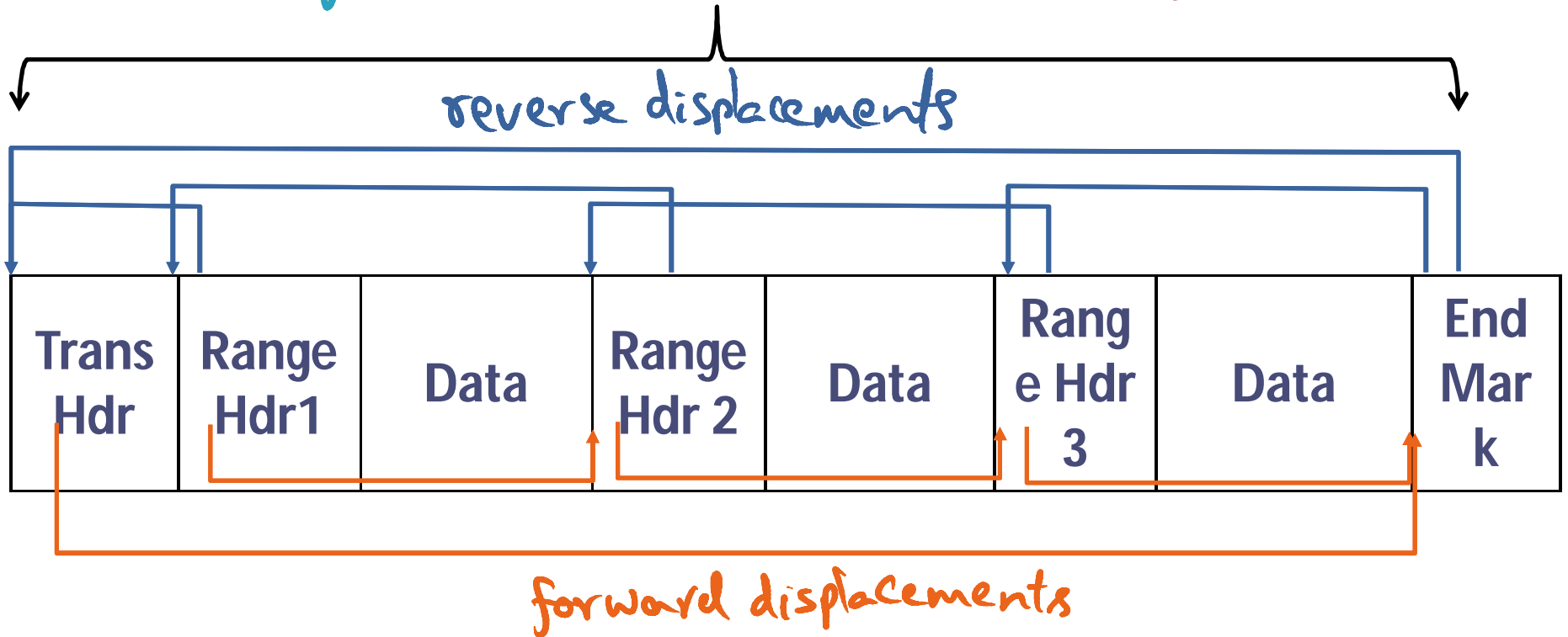


use xactions as insurance

# Implementation

Redo log

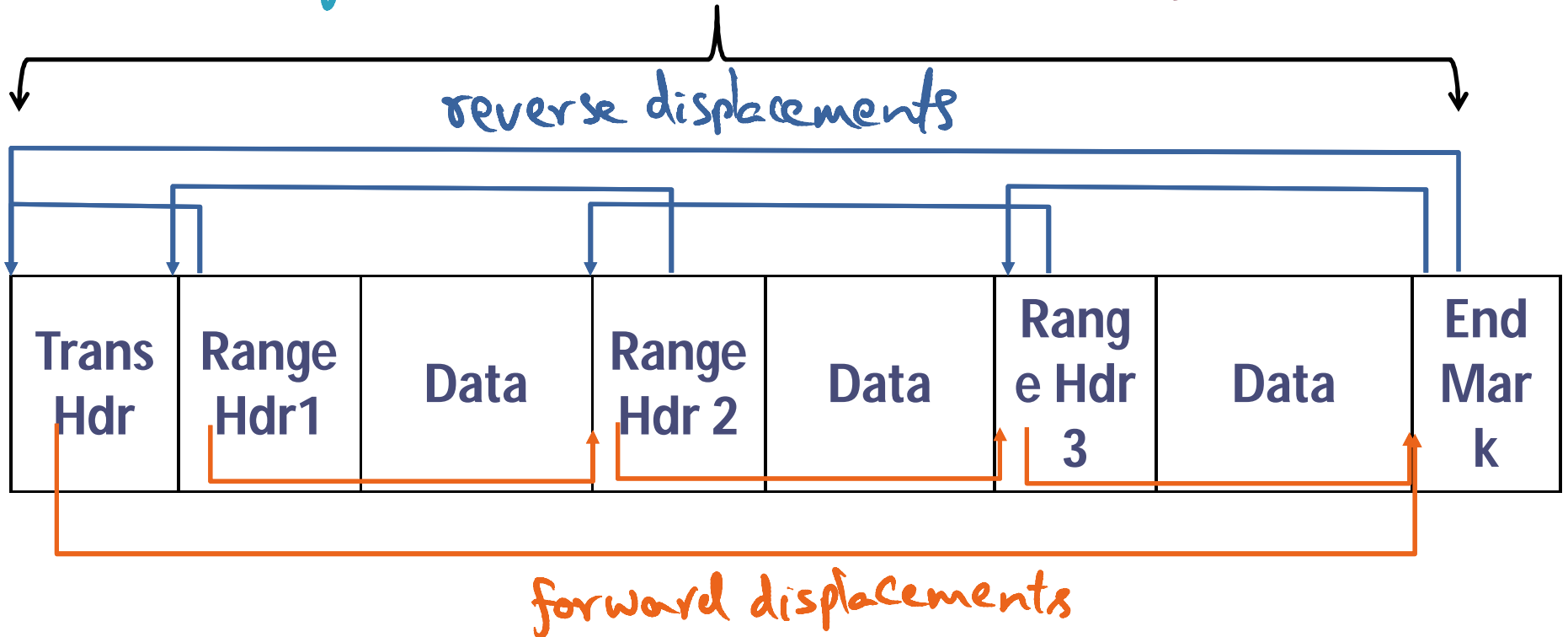
- All changes to different regions between begin and end xact



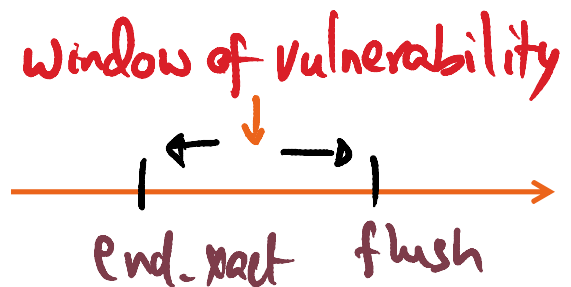
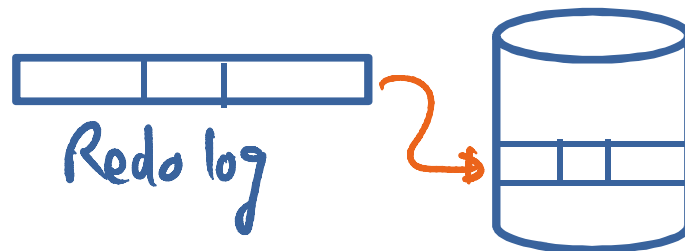
# Implementation

Redo log

- All changes to different regions between begin and end xact

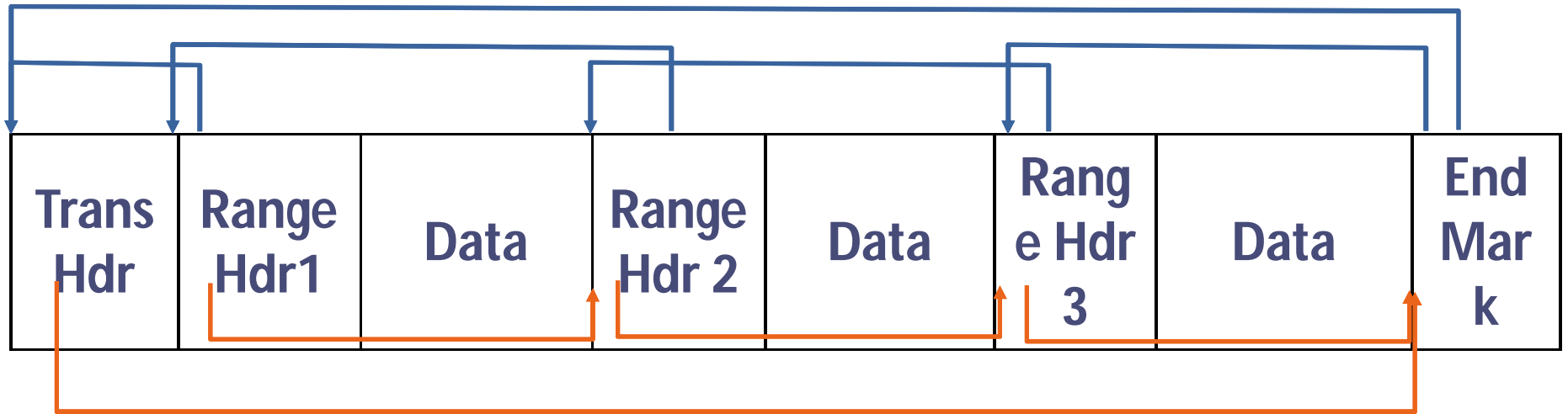


Commit



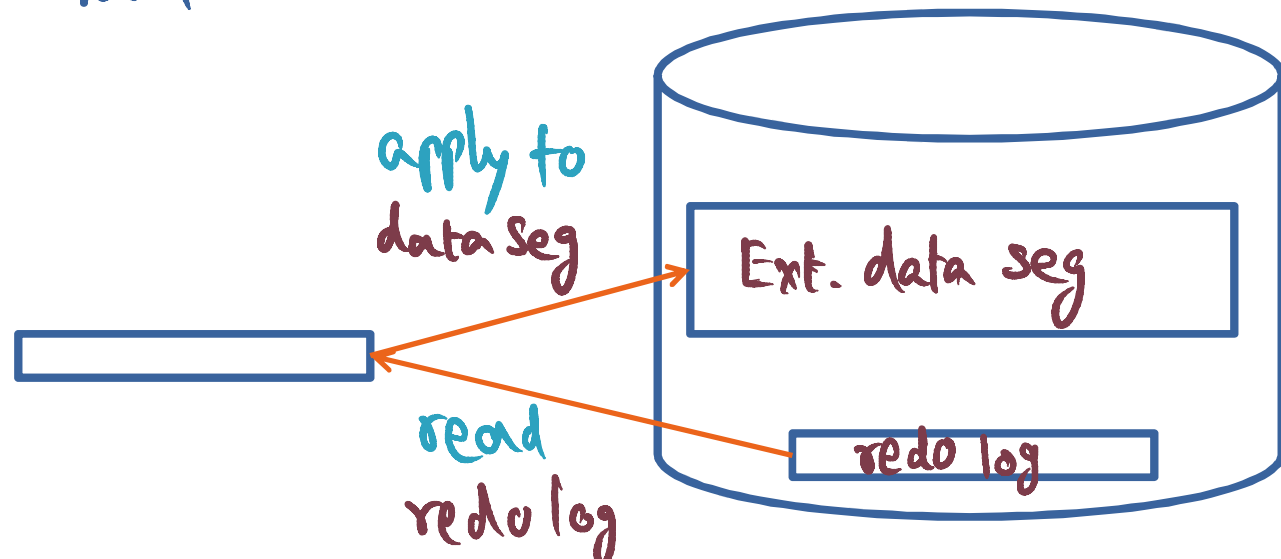
# Crash Recovery

reverse displacements



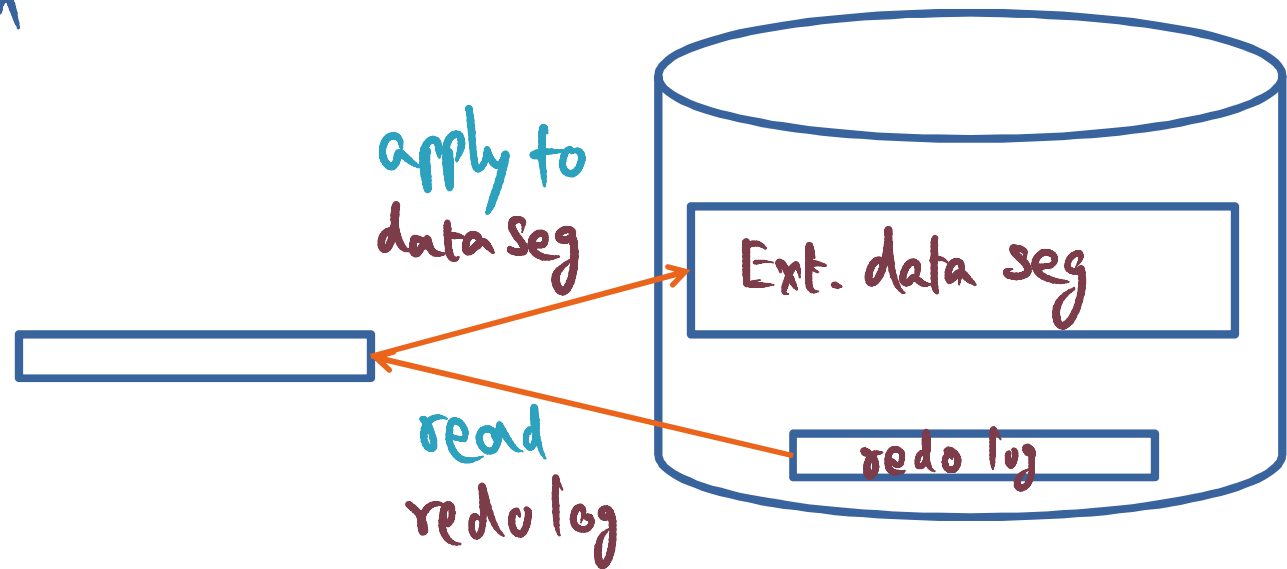
forward displacements

Resume from Crash



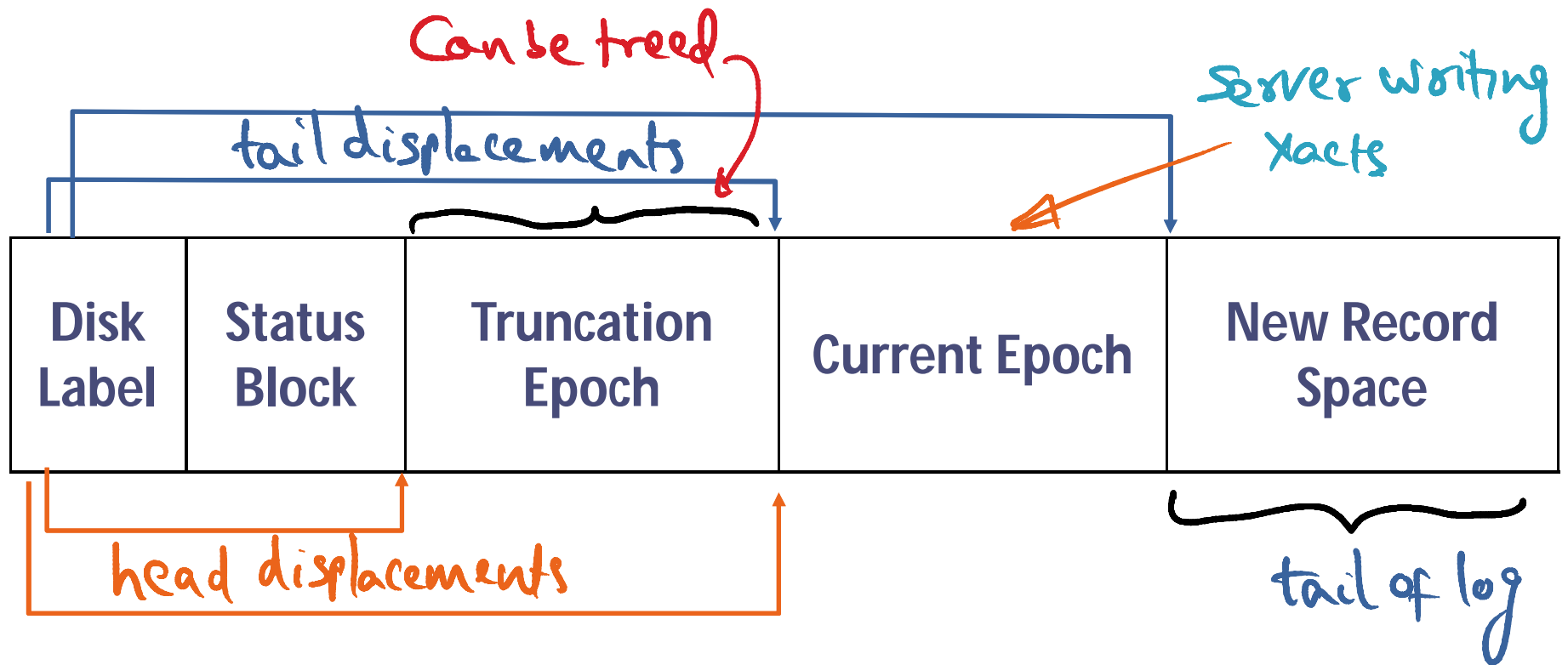
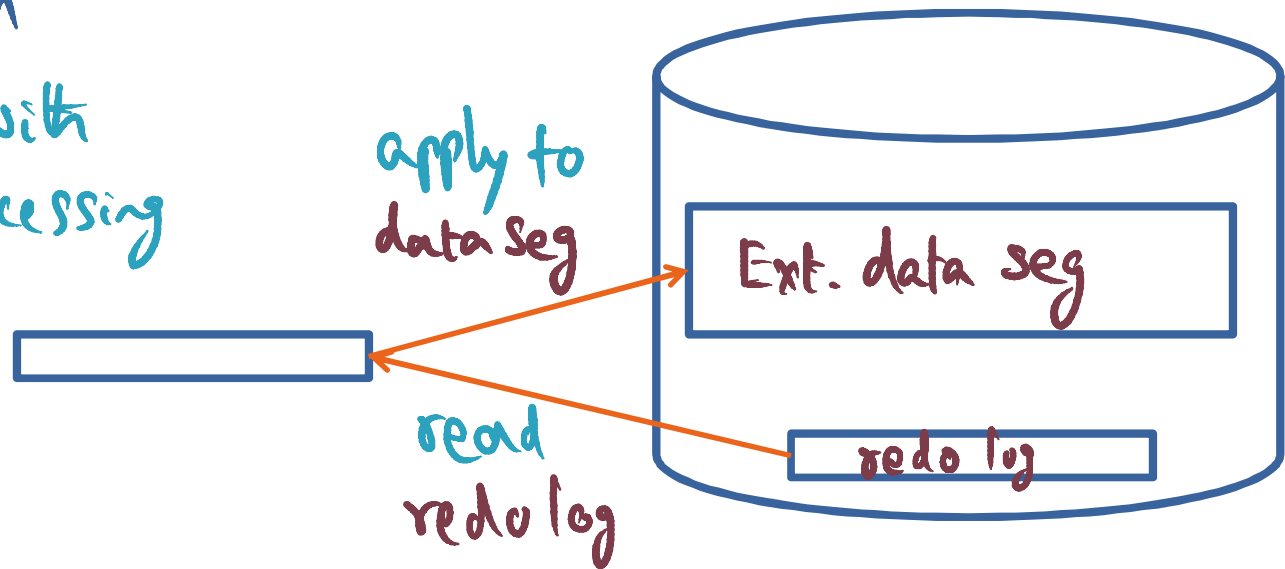


# Log truncation



# Log truncation

- in parallel with forward processing



1) Difference between

Traditional transaction

$V_s$

LRVM transaction

2) Why the difference

# Key Takeaways

- Classic systems research.
  - Understand **pain point** for developers
  - Solution to solve that pain point
- **Managing persistence** for critical data structures is the pain point identified by LRVM.
  - LRVM proposes using “**light weight**” transactions (without all the ACID properties)