

# 3D Mesh Crypto-Compression

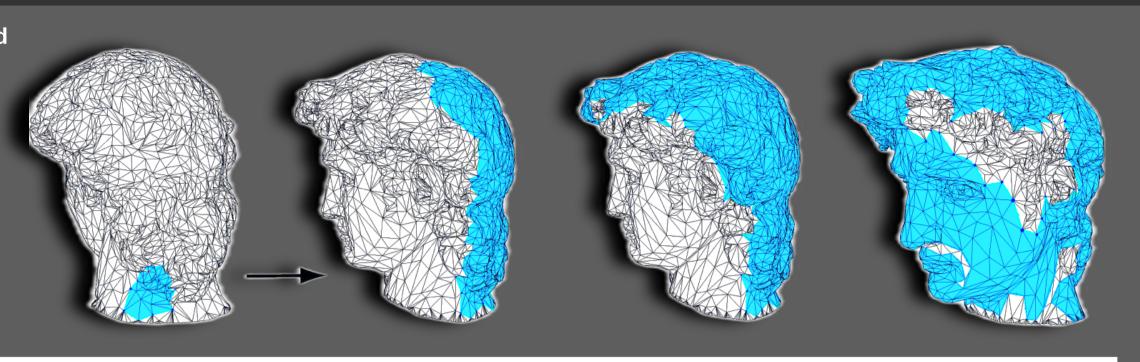
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3D meshes can be considered the most popular representation of volumes and surfaces and are used on a variety of devices with different performance constraints.

In a need for realism, these are becoming more and more detailed and complex over time. This is accompanied by an increase in the volume of data and can cause problems (memory, speed, performance, etc.) due to the constraints of the different media on which they are displayed.

To solve these problems, 3D mesh compression methods have been developed to reduce the amount of data required to represent them, making them easier to store and transfer.



#### **VALENCE-DRIVEN APPROACH**

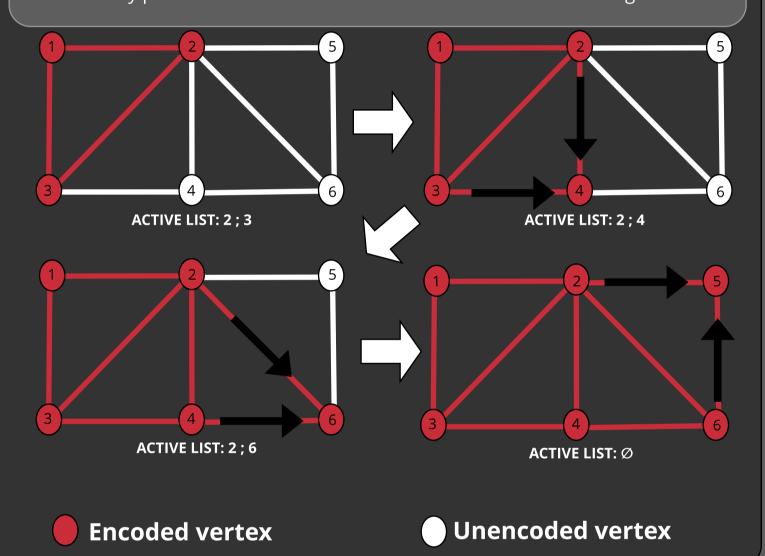
Among the large number of algorithms that exist to encode connectivity, we decided to go with a valence-based approach. It is one of the most optimized algorithms since it encodes our vertex with an average of 3 bits. The principle relies on the conquest of the mesh edges by expanding an initial boundary and encoding the conquered vertices along the edges.

#### Algorithm:

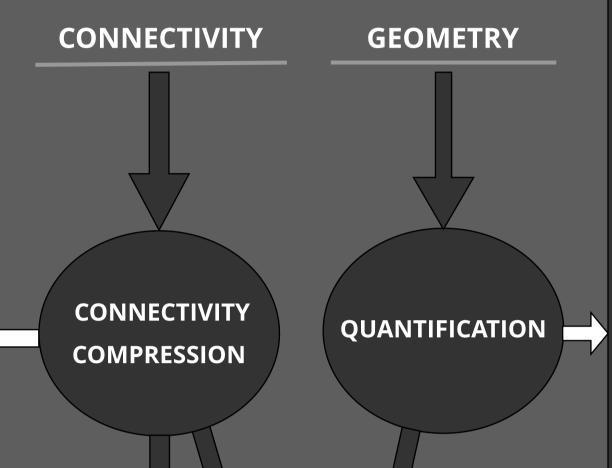
- start from a random triangle whose vertices are added in an active list (AL)
   representing the boundary to expand
- pick a vertex pivot from the AL
- ► try to conquer its unencoded edges in counter-clockwise order
- the vertices along each edge are put into the active list and encoded as a symbol that depends on their situation ( see below )
- when the vertex has no more unencoded edges, the next vertex in the active list becomes the next pivot and we go on until the AL is empty

#### 3 types of symbols are possible :

- ADD < vertex valence > : Occurs when a vertex is added to the active list.
- **SPLIT < valence offset >:** Occurs when the traversed vertex is already present in the active list. The active list is then split in half at the vertex occurrence.
- MERGE <AL index> <vertex offset>: Occurs when the traversed vertex is already present in another active list. The two lists are then merged



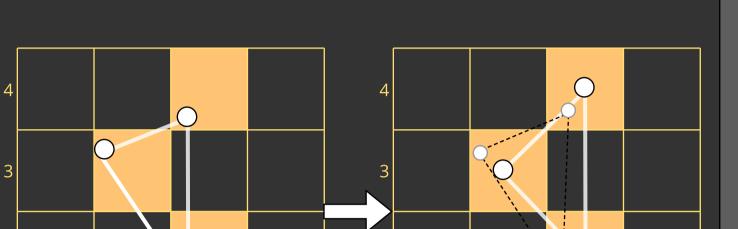
# **ALGORITHM SCHEME**



# QUANTIFICATION

First, the coordinates of each vertex are quantified. To do this, we normalize the coordinates between [0;1] and multiply the values by a quantification value.

This allows us to remove the superfluous decimals and thus to encode the vertices on fewer bits.





#### Parallelogram prediction:

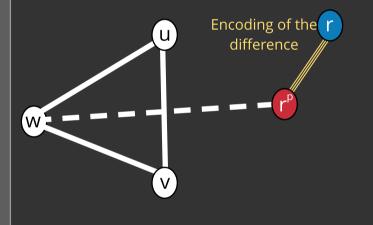
( Xi ; Yi

 Predict the position of vertex r using a triangle (u, v, w).

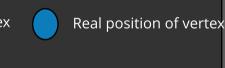
Compute r, the vertex that forms a parallelogram with the triangle.

 $\mathbf{r}^{\mathbf{p}} = \mathbf{u} + \mathbf{v} - \mathbf{w}$ 

► Encode the difference between r and r



( Xi ; Yi )



(3;2)

### **XOR ENCRYPTION** MESH **KEY** Quantified and predicted vertex 729 471 69 330 0101001010 | 0010010000 | 0111000001 1011011001 0111010111 0001000101 **XOR** Ζ Χ 915 327

# SHUFFLING ENCRYPTION

**PREDICTIVE** 

**ENCODING** 

XOR ENCRYPTION

#### **MESH KEY** Quantified and predicted vertex 365 145 2 145 | 354 | 142 | **ALGORITHM** [ 47 | 78 | 256 SHUFFLING 3 SHUFFLING 2 145 145 142 145 78 256 145 145 256 365 354 2 47 354 256 365 354 142 365 78 2 47 145 142 47 78 2

SHUFFLING ENCRYPTION

<u>RESULTS</u>

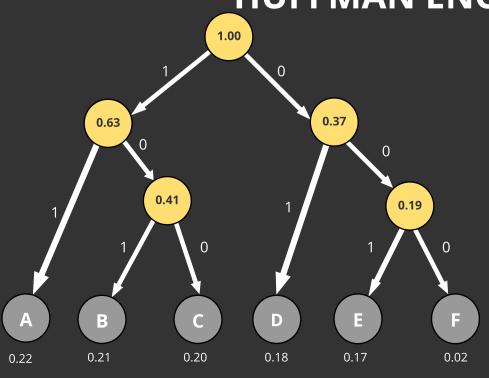
 NEW VERTEX 1
 NEW VERTEX 2
 NEW VERTEX 3

 X
 Y
 Z
 X
 Y
 Z
 X
 Y
 Z

 145
 78
 256
 365
 354
 2
 47
 145
 142

## **HUFFMAN ENCODING**

1110010011 0101000111 0110000100



The Huffman encoding algorithm builds a binary tree in which each leaf represents an element of the message to encode. The tree will have a shape such that the leaves corresponding to the most frequent elements will be located near the root, while the very infrequent elements will require a complex path from the root to access them. Therefore, the more frequently an element appears, the shorter its code will be when encoded.

# HUFFMAN ENCODING



## RESULTS



——640 vertices ——2562 vertices ——10242 vertices

