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Syntax.
                                       Term variables
                                                                                x, y \in Var Type variables t, u \in TVar
                                       Type constructors T^{\kappa} \in \mathcal{T}^{\kappa} where \{ -*^{\bullet}, -*^{\circ}, \rightarrow^{\circ} \} \subseteq \mathcal{T}^{\star \to \star \to \star}
                               \operatorname{Kinds}
                                                                                \kappa ::= \star \mid \kappa \to \kappa
                                                                             \tau^{\kappa} \, ::= t \stackrel{\cdot}{|} \, T^{\kappa} \mid \tau^{\kappa' \to \kappa} \, \tau^{\kappa'}
                               Types
                                                                                 \pi ::= \mbox{Un } \tau \mid \mbox{SeFun } \tau \mid \mbox{ShFun } \tau \mid \tau \geq \upsilon
                               Predicates
                               Qualified types
                                                                                 \rho ::= \tau^{\star} \mid \pi \Rightarrow \rho
                               Type schemes
                                                                                 \sigma ::= \rho \mid \forall t. \sigma
                               Environments
                                                                                H ::= \varepsilon \mid x : \sigma \mid H, H \mid H; H
                               Environment contexts \mathcal{H} := \varepsilon \mid x : \sigma \mid \Box \mid \mathcal{H}, H \mid H, \mathcal{H} \mid H; H
                               Expressions
                                                                        M,N ::= x \mid \lambda {*}x.M \mid \lambda \& x.M \mid M \ N \mid \underline{\mathtt{let}} \ x = M \ \underline{\mathtt{in}} \ N
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Figure 1: Types and terms.

$$(VAR) \frac{P \mid H \vdash M : \sigma}{P \mid H : \sigma \vdash X : \sigma} \qquad (LET) \frac{P \mid H \vdash M : \sigma \vdash P \mid H', x : \sigma \vdash N : \tau}{P \mid H, H' \vdash \text{let } x = M \text{ in } N : \tau}$$

$$(CTR-UN) \frac{P \mid \mathcal{H}(H', H') \vdash M : \sigma}{P \mid \mathcal{H}(H') \vdash M : \sigma} \qquad (WKN-UN) \frac{P \mid \mathcal{H}(E) \vdash M : \sigma}{P \mid \mathcal{H}(H') \vdash M : \sigma}$$

$$(CTR-SH) \frac{P \mid \mathcal{H}(H') \vdash M : \sigma}{P \mid \mathcal{H}(H') \vdash M : \sigma} \qquad (WKN-SH) \frac{P \mid \mathcal{H}(H') \vdash M : \sigma}{P \mid \mathcal{H}(H') \vdash M : \sigma}$$

$$(^{*}I) \frac{P \mid H, x : \tau \vdash M : v}{P \mid H \vdash \lambda * x . M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}I) \frac{P \mid H \vdash \lambda * x . M : \phi \tau v}{P \mid H \vdash \lambda * x . M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}I) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}I) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}I) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}I) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \rho}{P \mid H \vdash M : \phi \rho}$$

$$(^{*}VI) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}VI) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}VI) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}VI) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v}$$

$$(^{*}VI) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad (^{*}E) \frac{P \mid H \vdash M : \phi \tau v}{P \mid H \vdash M : \phi \tau v} \qquad ($$

Figure 2: Typing rules.

$$(* I) \frac{P \mid H \vdash M : \tau \quad P \mid H' \vdash N : \tau'}{P \mid H, H' \vdash (M, N) : \tau * \tau'} \qquad (* E) \frac{P \mid H \vdash M : \tau * \tau' \quad P \mid H', x : \tau, x' : \tau' \vdash N : v}{P \mid H, H' \vdash \underline{\text{let}} (x, x') = M \underline{\text{in}} N : v}$$

$$(\& I) \frac{P \mid H \vdash M : \tau \quad P \mid H' \vdash N : \tau'}{P \mid H; H' \vdash (M; N) : \tau \& \tau'} \qquad (\& E_1) \frac{P \mid H \vdash M : \tau \& \tau'}{P \mid H \vdash \text{fst } M : \tau} \qquad (\& E_2) \frac{P \mid H \vdash M : \tau \& \tau'}{P \mid H \vdash \text{snd } M : \tau'}$$

Figure 3: Derivable typing rules

$$(->1) \frac{(x:t,y:u);z:v\vdash(x,z):t*v}{x:t,y:u\vdash\lambda\&z.(x,z):v\to t*v}$$

$$(-*1) \frac{x:t+\lambda y\&z.(x,z):v\to t*v}{x:t\vdash\lambda y\&z.(x,z):u\to v\to t*v}$$

$$(-*1) \frac{x:t\vdash\lambda x*y\&z.(x,z):u\to v\to t*v}{x:t;y:u\vdash x:t}$$

$$(-*1) \frac{x:t\vdash x:t}{x:t;y:u\vdash x:t}$$

$$(-*1) \frac{(x:t;y:u),z:v\vdash(x,z):t*v}{x:t;y:u\vdash\lambda x:z.(x,z):v\to t*v}$$

$$(-*1) \frac{x:t;y:u\vdash\lambda x:z.(x,z):v\to t*v}{x:t;y:u\vdash\lambda x:z.(x,z):v\to t*v}$$

$$(-*1) \frac{x:t;y:u\vdash\lambda x*z.(x,z):v\to t*v}{x:t\vdash\lambda x*y.\lambda x*z.(x,z):u\to v\to t*v}$$

$$(-*1) \frac{(-*1)}{(-*1)} \frac{(-*1)}{(-*1$$

Figure 4: Sample derivations