

# Sequential Decision Making

## Lecture 5 : Dynamic Programming

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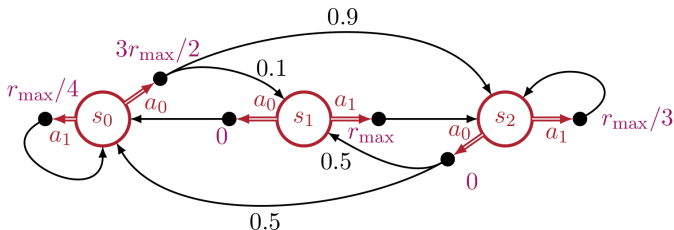


M2 Data Science, 2022/2023

# Reminder : Markov Decision Process

A MDP is parameterized by a tuple  $(\mathcal{S}, \mathcal{A}, R, P)$  where

- ▶  $\mathcal{S}$  is the **state space**
- ▶  $\mathcal{A}$  is the **action space** (or  $\mathcal{A}_s$  for each  $s \in \mathcal{S}$ )
- ▶  $R = (\nu_{(s,a)})_{(s,a) \in \mathcal{S} \times \mathcal{A}}$  where  $\nu_{(s,a)} \in \Delta(\mathbb{R})$  is the **reward distribution** for the state-action pair  $(s, a) \rightarrow r(s, a) = \mathbb{E}_{R \sim \nu_{(s,a)}}[R]$
- ▶  $P = (p(\cdot|s, a))_{(s,a) \in \mathcal{S} \times \mathcal{A}}$  where  $p(\cdot|s, a) \in \Delta(\mathcal{S})$  is the **transition kernel** associated to the state-action pair  $(s, a)$



## Reminder : Policy

A policy  $\pi = (\pi_0, \pi_1, \dots)$  is a sequence of mapping  $\pi_t : \mathcal{S} \rightarrow \Delta(\mathcal{A})$  that maps a state to a distribution over actions.

**Under policy  $\pi$** , at time  $t$ , the agent in state  $s_t$  selects

$$a_t \sim \pi_t(s_t),$$

receives the instantaneous reward

$$r_t \sim \nu_{(s_t, a_t)} \text{ such that } \mathbb{E}[r_t | s_t, a_t] = r(s_t, a_t)$$

and transits to the new state  $s_{t+1} \sim p(\cdot | s_t, a_t)$ .

→ a policy defines a probability model  $\mathbb{P}^\pi, \mathbb{E}^\pi$  over sequences of observations :

$$s_1, a_1, r_1, s_2, a_2, r_2 \dots$$

# Reminder : Value Function

## Definition

The value function of a policy  $\pi = (\pi_1, \pi_2, \dots)$  is  $V^\pi : \mathcal{S} \rightarrow \mathbb{R}$

- ① Finite-horizon criterion

$$V^\pi(s) = \mathbb{E}^\pi \left[ \sum_{t=1}^H r_t \mid s_1 = s \right]$$

- We want to compute the **optimal value**  $V^*(s) = \max_{\pi} V^\pi(s)$  and an **optimal policy**  $\pi_*$  such that  $V^* = V^{\pi_*}$ .
- We will be able to do so when  **$\mathcal{S}$  and  $\mathcal{A}$  are finite**

$$S := |\mathcal{S}| < \infty \quad \text{and} \quad A := |\mathcal{A}| < \infty$$

*(some optimality equation may extend to continuous state spaces)*

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- ② Infinite horizon with a discount  $\gamma$

$$V^\pi(s) = \mathbb{E}^\pi \left[ \sum_{t=1}^{\infty} \gamma^{t-1} r(s_t, a_t) \middle| s_1 = s \right]$$

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# Outline

- 1 Solving a Known MDP : Finite Horizon
- 2 Solving a Known MDP : the Discounted Case
  - Policy Evaluation
  - Computing the Optimal Policy
- 3 Value Iteration, Policy Iteration

# Value functions

Let  $H$  be the known time horizon.

## Value functions at step $h$

For a policy  $\pi = (\pi_1, \dots, \pi_H)$ ,

$$V_h^\pi(s) = \mathbb{E}^\pi \left[ \sum_{t=h}^H r_t \mid s_h = s \right]$$

and

$$V_h^*(s) = \max_{\pi_h, \dots, \pi_H} \mathbb{E}^\pi \left[ \sum_{t=h}^H r_t \mid s_h = s \right]$$

**Goal :** compute

$$V^\pi(s) = V_1^\pi(s), V^*(s) = V_1^*(s) \quad \text{and} \quad \pi^* = (\pi_1^*, \dots, \pi_H^*).$$

→ we will actually compute  $V_h^\pi(s)$  and  $V_h^*(s)$  for all  $h \leq H$ .



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# Value functions

Let  $H$  be the known time horizon.

## Value functions at step $h$

For a **deterministic** policy  $\pi = (\pi_1, \dots, \pi_H)$ ,

$$V_h^\pi(s) = \mathbb{E}^\pi \left[ \sum_{t=h}^H r(s_t, \pi_t(s_t)) \mid s_h = s \right]$$

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**How ?**

- Monte-Carlo estimation ? **only approximate**
- Develop the tree of all possible realizations ? **too complex**

# Bellman equations

## Proposition

The value functions of a deterministic policy  $\pi$  satisfies the following equations : for all  $h \in \{1, \dots, H\}$ ,

$$V_h^\pi(s) = r(s, \pi_h(s)) + \sum_{s' \in \mathcal{S}} p(s'|s, \pi_h(s)) V_{h+1}^\pi(s'),$$

with the convention that  $V_{H+1}^\pi(s) = 0$  for all  $s \in \mathcal{S}$ .

**Consequence** : for a finite state space  $\mathcal{S}$  such that  $|\mathcal{S}| = S$

- $V_1^\pi(s)$  can be computed using **backwards induction**
- space complexity :  $S \times H$
- time complexity :  $S \times (S + 1) \times H$

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**Proof.**

$$\begin{aligned} V_h^\pi(s) &= \mathbb{E}^\pi \left[ r(s_h, \pi_h(s_h)) + \sum_{t=h+1}^H r(s_t, \pi_t(s_t)) \middle| s_h = s \right] \\ &= r(s, \pi_h(s)) + \mathbb{E}^\pi \left[ \sum_{t=h+1}^H r(s_t, \pi_t(s_t)) \middle| s_h = s, a_h = \pi_h(s) \right] \\ &= r(s, \pi_h(s)) + \sum_{s' \in \mathcal{S}} \mathbb{P}(s_{h+1} = s' | s_h = s, a_h = \pi_h(s)) \mathbb{E}^\pi \left[ \sum_{t=h+1}^H r(s_t, \pi_t(s_t)) \middle| s_{h+1} = s' \right] \\ &= r(s, \pi_h(s)) + \sum_{s' \in \mathcal{S}} p(s'|s, \pi_h(s)) V_{h+1}^\pi(s') \end{aligned}$$

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with the convention that  $V_{H+1}^\pi(s) = 0$  for all  $s \in \mathcal{S}$ .

These equations may be generalized :

- ▶ to a possibly infinite state space

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with the convention that  $V_{H+1}^\pi(s) = 0$  for all  $s \in \mathcal{S}$ .

These equations may be generalized :

- ▶ to a possibly infinite state space
- ▶ to randomized policies

# Bellman equations for the optimal values

## Proposition

The **optimal values**  $V_h^*$  satisfy the **Bellman equations** :

$$V_h^*(s) = \max_{a \in \mathcal{A}} \left[ r(s, a) + \sum_{s' \in \mathcal{S}} p(s'|s, a) V_{h+1}^*(s') \right] \quad \text{for all } h \leq H,$$

with the convention that  $V_{H+1}^*(s) = 0$  for all  $s \in \mathcal{S}$ .

Moreover, an optimal policy is given by

$$\pi_h^*(s) \in \operatorname{argmax}_{a \in \mathcal{A}} \left[ r(s, a) + \sum_{s'=1}^S p(s'|s, a) V_{h+1}^*(s') \right].$$



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**Consequence** : for finite  $\mathcal{S}, \mathcal{A}$  such that  $|\mathcal{S}| = S, |\mathcal{A}| = A$

- $\pi^* = (\pi_1^*, \dots, \pi_H^*)$  can be computed using **backwards induction**
- space complexity :  $S \times H$
- time complexity :  $O(S^2AH)$

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This technique is known as **Dynamic Programming**

- ▶ term invented in the 50s by Bellman : an algorithmic principle for optimization in which solving an optimization problem of a given size reduces to solving (several) of the same optimization problem but of smaller size

# Proof

$$\begin{aligned}
 V_h^*(s) &= \max_{\pi_h, \pi_{h+1}, \dots} \mathbb{E}^{\pi} \left[ \sum_{t=h}^H r(s_t, a_t) \mid s_h = s \right] \\
 &= \max_{\pi_h, \pi_{h+1}, \dots} \sum_{a \in \mathcal{A}} \pi_h(a_h = a \mid s_h = s) \mathbb{E}^{\pi} \left[ r(s, a) + \sum_{t=h+1}^H r(s_t, a_t) \mid s_h = s, a_h = a \right] \\
 &= \max_{\pi_h, \pi_{h+1}, \dots} \sum_{a \in \mathcal{A}} \pi_h(a_h = a \mid s_h = s) \left[ r(s, a) + \sum_{s' \in \mathcal{S}} p(s' \mid s, a) \mathbb{E}^{\pi_{h+1}, \dots} \left[ \sum_{t=h+1}^H r(s_t, a_t) \mid s_{h+1} = s' \right] \right] \\
 &= \max_{a \in \mathcal{A}} \left[ r(s, a) + \sum_{s' \in \mathcal{S}} p(s' \mid s, a) \max_{\pi_{h+1}, \dots} V_{h+1}^{\pi_{h+1}, \dots}(s') \right] \\
 &= \max_{a \in \mathcal{A}} \left[ r(s, a) + \sum_{s' \in \mathcal{S}} p(s' \mid s, a) V_{h+1}^*(s') \right]
 \end{aligned}$$

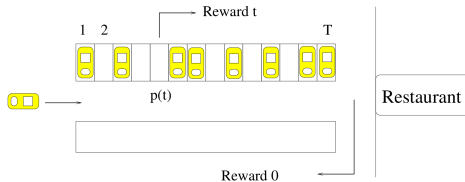
The maximizing policy is  $\pi_h^*, \pi_{h+1}^*, \dots$  with

$$V_{h+1}^{\pi_{h+1}^*, \dots} = V_{h+1}^* = \operatorname{argmax}_{\pi} V_{h+1}^{\pi}$$

and a **deterministic** mapping  $\pi_h^* : \mathcal{S} \rightarrow \mathcal{A}$  given by

$$\pi_h^*(s) \in \operatorname{argmax}_{a \in \mathcal{A}} \left[ r(s, a) + \sum_{s' \in \mathcal{S}} p(s' \mid s, a) V_{h+1}^*(s') \right].$$

# Example : Optimal Parking



## Exercise :

- model optimal parking as solving a MDP with a finite horizon
- write the optimal policy

# Outline

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# Values functions

Let  $\gamma \in (0, 1)$  be a known discount factor

## Value functions

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and

$$V^*(s) = \max_{\pi} \mathbb{E}^\pi \left[ \sum_{t=1}^{\infty} \gamma^{t-1} r_t \mid s_1 = s \right]$$

## How to compute them ?

→ We need to generalize Dynamic Programming to infinite horizon...

# Values functions

Let  $\gamma \in (0, 1)$  be a known discount factor

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For a **deterministic** policy  $\pi = (\pi_1, \pi_2, \dots)$ ,

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# Bellman equation for a **stationary** policy

## Proposition

Any **stationary** deterministic policy  $\pi$  satisfies, for all  $s \in \mathcal{S}$ ,

$$V^\pi(s) = r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) V^\pi(s')$$

**Proof.**

$$\begin{aligned} V^\pi(s) &= \mathbb{E}^\pi \left[ r(s, \pi(s)) + \sum_{t=2}^{\infty} \gamma^{t-1} r(s_t, \pi(s_t)) \middle| s_1 = s \right] \\ &= r(s, \pi(s)) + \gamma \mathbb{E}^\pi \left[ \sum_{t=2}^{\infty} \gamma^{t-2} r(s_t, \pi(s_t)) \middle| s_1 = s, a_1 = \pi(s) \right] \\ &= r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} \mathbb{P}(s_2 = s' | s_1 = s, a_1 = \pi(s)) \mathbb{E}^\pi \left[ \sum_{t=2}^{\infty} \gamma^{t-2} r(s_t, \pi(s_t)) \middle| s_2 = s' \right] \\ &= r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) \underbrace{\mathbb{E}^\pi \left[ \sum_{t=1}^{\infty} \gamma^{t-1} r(s_t, \pi(s_t)) \middle| s_1 = s' \right]}_{V^\pi(s')} \end{aligned}$$

# Bellman equation for a stationary policy

## Proposition

Any stationary deterministic policy  $\pi$  satisfies, for all  $s \in \mathcal{S}$ ,

$$V^\pi(s) = r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) V^\pi(s')$$

More general statement :

$$V^\pi(s) = \mathbb{E}_{a \sim \pi(s)} [r(s, a) + \gamma \mathbb{E}_{s' \sim p(\cdot|s, a)} [V^\pi(s')]]$$

(also applies to infinite state space and randomized policies)

# Solving the Bellman equations

Fix a stationary, deterministic policy  $\pi$ .

## Proposition

$$\forall s \in \mathcal{S}, \quad V^\pi(s) = r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) V^\pi(s')$$

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Introducing the vectors

$$\begin{aligned} V^\pi &= (V^\pi(s))_{s=1}^S \in \mathbb{R}^S \\ r^\pi &= (r(s, \pi(s)))_{s=1}^S \in \mathbb{R}^S \end{aligned}$$

and the matrix

$$P^\pi = \left( p(s'|s, \pi(s)) \right)_{\substack{1 \leq s \leq S \\ 1 \leq s' \leq S}} \in \mathbb{R}^{S \times S},$$

the Bellman equations rewrite

$$V^\pi = r^\pi + \gamma P^\pi V^\pi$$

# Solving the Bellman equations

$$V^\pi = r^\pi + \gamma P^\pi V^\pi$$

The vector  $V^\pi \in \mathbb{R}^S$  satisfies

$$\begin{aligned}(I - \gamma P^\pi) V^\pi &= r^\pi \\ V^\pi &= (I - \gamma P^\pi)^{-1} r^\pi\end{aligned}$$

provided that the matrix  $I - \gamma P^\pi$  is invertible.

## Proposition

The eigenvalues of the *stochastic*<sup>a</sup> matrix  $P^\pi$  all belong to  $[0, 1]$ . As a consequence,  $-\gamma \notin \text{sp}(P^\pi)$  thus  $I - \gamma P^\pi$  is invertible.

---

a. the entries in its rows sum to 1

→  $V^\pi$  can be computed by inverting a  $S \times S$  matrix

## An alternative : Exploiting the Bellman operator

### Definition

The **Bellman operator** associated to a policy  $\pi$  is defined by

$$\begin{aligned} T^\pi : \mathbb{R}^S &\longrightarrow \mathbb{R}^S \\ V &\mapsto T^\pi(V) \end{aligned}$$

where

$$T^\pi(V)(s) = r(s, \pi(s)) + \gamma \sum_{s' \in S} p(s'|s, \pi(s)) V(s')$$

The Bellman equation for policy  $\pi$  rewrites

$$V^\pi = T^\pi V^\pi$$

→ the vector  $V^\pi$  is a **fixed point** of the Bellman operator  $T^\pi$

## Intermezzo : Fixed Point Theorem

### Definition

Let  $(\mathcal{X}, |\cdot|)$  be a normed vector space.

A mapping  $f : \mathcal{X} \rightarrow \mathcal{X}$  is called  $L$ -Lipschitz if

$$\forall (x, y) \in \mathcal{X}^2, |f(x) - f(y)| \leq L|x - y|.$$

If  $L < 1$ ,  $f$  is called a **contraction**.

### Banach's fixed point theorem

Let  $(\mathcal{X}, |\cdot|)$  be a *complete* normed vector space and  $f : \mathcal{X} \rightarrow \mathcal{X}$  be a contraction. Then

- ▶ there exists a **unique fixed point**  $x^*$  satisfying  $f(x^*) = x^*$
- ▶ for every  $x_0 \in \mathcal{X}$ , the sequence defined by  $x_{n+1} = f(x_n)$  for all  $n \geq 0$  satisfies

$$\lim_{n \rightarrow \infty} x_n = x_*$$

# Exploiting the Bellman Operator

## Proposition

The operator

$$\begin{aligned} T^\pi: (\mathbb{R}^S, \|\cdot\|_\infty) &\longrightarrow (\mathbb{R}^S, \|\cdot\|_\infty) \\ V &\mapsto T^\pi(V)(s) = r(s, \pi(s)) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) V(s') \end{aligned}$$

is a  $\gamma$ -contraction.

**Proof.**

$$\begin{aligned} \|T^\pi(V) - T^\pi(V')\|_\infty &= \sup_{s \in \mathcal{S}} |T^\pi(V)(s) - T^\pi(V')(s)| \\ &= \sup_{s \in \mathcal{S}} |\gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) (V(s') - V'(s'))| \\ &\leq \gamma \sum_{s' \in \mathcal{S}} p(s'|s, \pi(s)) \|V - V'\|_\infty \\ &= \gamma \|V - V'\|_\infty. \end{aligned}$$



# Exploiting the Bellman Operator

## Proposition

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is a  $\gamma$ -contraction.

**Consequence :**

- ▶  $V^\pi$  is the unique fixed point of  $T^\pi$
- ▶  $V^\pi$  can be approximated by an iterative scheme :

$$V^\pi = \lim_{k \rightarrow \infty} V_k$$

where

$$\begin{cases} V_0 & \in \mathbb{R} \\ V_{k+1} & = T^\pi(V_k) \text{ for all } k \geq 0. \end{cases}$$

# Summary : Policy Evaluation

Two methods for computing  $V^\pi(s)$  for all  $s$  :

- ▶ solving linear equations (**matrix inversion**)
- ▶ **iterating the Bellman operator  $T^\pi$**

**Other possibility** : Monte-Carlo simulation

- provides only an **approximation**
- ... but **doesn't require the knowledge of  $r(s, a)$  and  $p(\cdot|s, a)$** ...

# Back to Retail Store Management

- Constant policy :  $\pi(s) = \max(M - s, c)$
- Threshold policy :  $\pi(s) = \mathbb{1}_{(s \leq m_1)}(m_2 - s)$

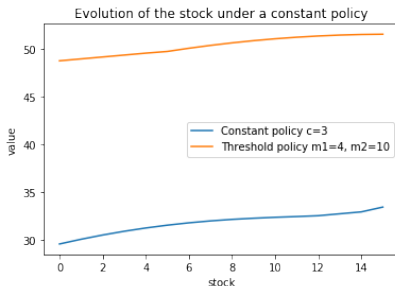


FIGURE –  $V^\pi$  for two different policies,  $\gamma = 0.97$

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# Bellman equations for the optimal values

## Proposition

$V^*(s) = \max_{\pi} V^{\pi}(s)$  satisfy the **Bellman equations** :

$$V^*(s) = \max_{a \in \mathcal{A}} \left[ r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V^*(s') \right]$$

Moreover, an optimal policy is given by  $\pi^* = (\pi^*, \pi^*, \dots)$  where

$$\pi^*(s) \in \operatorname{argmax}_{a \in \mathcal{A}} \left[ r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V^*(s') \right].$$

→  $\pi^*$  is the **greedy policy** with respect to  $V^*$  :

## Definition

The **greedy policy** with respect to a value  $V$ ,  $\pi = \text{greedy}(V)$  is

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→  $\pi^*$  is the **greedy policy** with respect to  $V^*$  :

**Intuition** : greedy( $V$ ) is the policy that “improves” a policy with value  $V$  by taking the best possible first action and then following the policy

# Solving the Bellman equations

## Proposition

The optimal value function  $V^*$  satisfies

$$V^*(s) = \max_{a \in \mathcal{A}} \left[ r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V^*(s') \right]$$

- ▶ a system of  $S$  non-linear equations for computing  $(V^*(s))_{s \in \mathcal{S}}$ .
- no hope for a simple “matrix inversion” technique...

# Bellman operator to the rescue

## Optimal Bellman operator

The **optimal Bellman operator** (or dynamic programming operator) is

$$\begin{aligned} T^* : \mathbb{R}^S &\longrightarrow \mathbb{R}^S \\ V &\mapsto T^*(V) \end{aligned}$$

where

$$T^*(V)(s) = \max_{a \in \mathcal{A}} \left[ r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V(s') \right]$$

The optimal value function satisfies

$$V^* = T^*(V^*)$$

→  $V^*$  is a **fixed point** of the optimal Bellman operator  $T^*$ .



# Optimal Bellman Operator

## Properties

The optimal Bellman operator is a  $\gamma$ -contraction :

$$\forall V, V' \in \mathbb{R}^S, \quad \|T^*(V) - T^*(V')\|_\infty < \gamma \|V - V'\|_\infty.$$

As a consequence :

- ▶  $T^*$  admits a unique fixed point,  $V^*$
- ▶ for every  $V_0$ , the sequence  $V_{n+1} = T^*(V_n)$  converges to  $V^*$

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**Proof.** Uses that for all functions  $|\max f - \max g| \leq \max |f - g|$ .



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- ▶ for every  $V_0$ , the sequence  $V_{n+1} = T^*(V_n)$  converges to  $V^*$

→ provides a method for approximating  $V^*$

# Outline

- 1 Solving a Known MDP : Finite Horizon
- 2 Solving a Known MDP : the Discounted Case
  - Policy Evaluation
  - Computing the Optimal Policy
- 3 Value Iteration, Policy Iteration

# Value Iteration

- **Idea** : Iterate the operator  $T^*$  until  $V$  doesn't change much

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## Algorithm 1: Value Iteration

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**Input** :  $\epsilon$  = stopping parameter

$V_0$  = any function (e.g.  $V_0 \leftarrow 0_S$ )

```
1  $V \leftarrow V_0$ 
2 while  $\|V - T^*(V)\| \geq \epsilon$  do
3    $V \leftarrow T^*(V)$ 
4 end
```

**Return:**  $\pi = \text{greedy}(V)$

---

## Theorem

Value iteration converges in at most  $\log \left( \frac{\|T^*(V_0) - V_0\|_\infty}{\epsilon(1-\gamma)} \right) / \log(1/\gamma)$  iterations and outputs a policy  $\pi$  satisfying  $\|V^\pi - V^*\| \leq \epsilon$ .

**Proof.** contraction property + triangular inequality

$$\|V^* - V_k\| \leq \|V^* - T^*(V_k)\| + \|T^*(V_k) - V_k\|$$

# Policy Iteration

- **Idea** : alternate between **policy evaluation** and **policy improvement**

## Greedy policy

Recall that  $\pi' = \text{greedy}(V)$  is the policy defined as

$$\pi'(s) \in \operatorname{argmax}_{a \in \mathcal{A}} \left[ r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V(s') \right]$$

## Policy improvement theorem

For any policy  $\pi$ ,  $\pi' = \text{greedy}(V^\pi)$  improves over  $\pi$  :  $V^{\pi'} \geq V^\pi$ .

**Proof.**



# Policy Iteration

- **Idea** : alternate between **policy evaluation** and **policy improvement** .

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## Algorithm 2: Policy Iteration

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**Input** :  $\pi_0 =$  any policy (e.g. chosen at random)

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```
1  $\pi \leftarrow \pi_0$ 
2  $\pi' \leftarrow \text{NULL}$ 
3 while  $\pi \neq \pi'$  do
4    $\pi' \leftarrow \pi$ 
5   Evaluate policy  $\pi$  : compute  $V^\pi$ 
6   Improve policy  $\pi$  :  $\pi \leftarrow \text{greedy}(V^\pi)$ 
7 end
```

**Return:**  $\pi$

---

## Theorem

Policy iteration terminates after a **finite number of steps** and outputs the **optimal policy**  $\pi^*$ .

# Policy Iteration

## Why is that ?

Policy iteration generates a sequence of policies  $\pi_0, \pi_1, \dots$  such that

$$\pi_{k+1} = \text{greedy}(V^{\pi_k}).$$

By the policy improvement theorem,

$$V^{\pi_{k+1}} \geq V^{\pi_k}$$

and if  $\pi_{k+1} \neq \pi_k$  there must exist  $s$  such that  $V^{\pi_{k+1}}(s) > V^{\pi_k}(s)$   
(otherwise  $V^{\pi_k} = V^{\pi_{k+1}} = T^{\pi_{k+1}}(V^{\pi_{k+1}}) = T^*(V^{\pi_k})$  thus  $\pi_k = \pi^*$ )

→ as there is a finite number of possible values of  $V^\pi$  (finite number of policies), the sequence must be stationary at some point.



# Implementation of VI and PI

Both algorithm can be implemented using **Q-Values** .

## Definitions

The **Q-value** of a stationary policy is the expected cumulative reward when performing action  $a$  in state  $s$  and then following policy  $\pi$  :

$$Q^{\pi}(s, a) = \mathbb{E}^{\pi} \left[ \sum_{t=1}^{\infty} \gamma^{t-1} r(s_t, a_t) \middle| s_1 = s, a_1 = a \right]$$

The **optimal Q-value** is  $Q^*(s, a) = \max_{\pi} Q^{\pi}(s, a) = Q^{\pi^*}(s, a)$ .

## Properties :

- ▶  $V^{\pi}(s) = Q^{\pi}(s, \pi(s))$
- ▶  $V^*(s) = Q^*(s, \pi^*(s))$

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# Implementation of VI and PI

Q-Values are convenient for policy improvement.

Q-value associated to a value  $V$

To each value function  $V$ , we can associate the corresponding Q-value

$$Q(s, a) = r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V(s')$$

**Property** :  $\pi' = \text{greedy}(V)$  can be rewritten  $\pi'(s) = \operatorname{argmax}_a Q(s, a)$ .

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**Remark** :  $\pi^* = \text{greedy}(Q^*)$

# Value Iteration versus Policy Iteration

In these implementations, we propose to store **Q-values**.

## Value Iteration

Initialize  $Q_0$ .

At iteration  $k$  :

$$V_k(s) = \max_a Q_{k-1}(s, a) \quad (\text{apply } T^*)$$

$$Q_k(s, a) = r(s, a) + \gamma \sum_{s' \in \mathcal{S}} p(s'|s, a) V_k(s')$$

Output :  $\pi_K(s) = \operatorname{argmax}_{a \in \mathcal{A}} Q_K(s, a)$

## Policy iteration

Initialize  $\pi_0$

At iteration  $k$  :

$$Q_{k-1}(s, a) = Q^{\pi_{k-1}}(s, a) \quad (\text{policy evaluation})$$

$$\pi_k(s) = \operatorname{argmax}_{a \in \mathcal{A}} Q_{k-1}(s, a)$$

Output :  $\pi_K$

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**Space Complexity** :  $O(SA)$  in both cases

- ▶ VI : Storing Q Values + Values
- ▶ PI : Storing Q values + Policy

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**Per Iteration Time Complexity** :  $O(S^2A) + O(S^3)$

- ▶ VI : Compute Q values + compute  $S$  max
- ▶ PI : Compute Q values + compute  $S$  argmax + Policy Evaluation

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**Number of Iterations ?** : PI often requires **few iterations**

- ▶ VI : wait for  $V_{k+1} \simeq V_k$  (requires a termination criterion)
- ▶ PI : wait for  $\pi_{k+1} = \pi_k$  (finite number of iterations)

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Output :  $\pi_K$

## Guarantees :

- ▶ VI : a policy with value very close to  $V^*$  (often  $\pi^*$ )
- ▶ PI : an optimal policy  $\pi^*$

# Summary

We learned how to find the optimal policy in an MDP with finite state and action spaces :

- ▶ using backwards induction for a finite horizon  $H$
- ▶ using Policy and Value iteration for an infinite horizon with a discount  $\gamma \in (0, 1)$

Those two types of techniques are often indifferently referred to as

## Dynamic Programming.

We are now ready to propose **reinforcement learning algorithms**, that :

- ▶ operate without the knowledge of  $r(s, a)$  and  $p(\cdot|s, a)$
- ▶ or in very large state spaces in which standard Dynamic Programming is intractable