

CS 5/7343 Operating System

Programming Assignment 2

Due: Friday, November 22, 2024 at 11:59 PM

You are required to complete Programming Projects 2, 3, and/or 4 from Chapter 7.

- **CS 5343 students:** Complete any **2** projects.
- **CS 7343 students:** Complete all **3** projects.

Project Folder Setup:

- Create a separate folder for each project.
- Name each folder as follows:
 - **SleepingTA**
 - **DiningPhilosophers**
 - **ProducerConsumer**

For each project, include the following:

1. The source code with comments (**25 points**)
2. A brief description of how to test your program (**12.5 points**)
3. A sample run (including input and output), which can be a screenshot or video of the program running (**12.5 points**)

Submission:

- Compress (zip) each project's folder individually and upload the zip files to Canvas.
 - **Undergraduate students (CS 5343):** Submit **2 zip files**.
 - **Graduate students (CS 7343):** Submit **3 zip files**.

Project 2—The Sleeping Teaching Assistant

A university computer science department has a teaching assistant (TA) who helps undergraduate students with their programming assignments during regular office hours. The TA's office is rather small and has room for only one desk with a chair and computer. There are three chairs in the hallway outside the office where students can sit and wait if the TA is currently helping another student. When there are no students who need help during office hours, the TA sits at the desk and takes a nap. If a student arrives during office hours and finds the TA sleeping, the student must awaken the TA to ask for help. If a student arrives and finds the TA currently helping another student, the student sits on one of the chairs in the hallway and waits. If no chairs are available, the student will come back at a later time.

Using POSIX threads, mutex locks, and semaphores, implement a solution that coordinates the activities of the TA and the students. Details for this assignment are provided below.

The Students and the TA

Using Pthreads (Section 4.4.1), begin by creating n students where each student will run as a separate thread. The TA will run as a separate thread as well. Student threads will alternate between programming for a period of time and seeking help from the TA. If the TA is available, they will obtain help. Otherwise, they will either sit in a chair in the hallway or, if no chairs are available, will resume programming and will seek help at a later time. If a student arrives and notices that the TA is sleeping, the student must notify the TA using a semaphore. When the TA finishes helping a student, the TA must check to see if there are students waiting for help in the hallway. If so, the TA must help each of these students in turn. If no students are present, the TA may return to napping.

Perhaps the best option for simulating students programming—as well as the TA providing help to a student—is to have the appropriate threads sleep for a random period of time.

Coverage of POSIX mutex locks and semaphores is provided in Section 7.3. Consult that section for details.

DiningPhilosophers

Project 3—The Dining-Philosophers Problem

In Section 7.1.3, we provide an outline of a solution to the dining-philosophers problem using monitors. This project involves implementing a solution to this problem using either POSIX mutex locks and condition variables or Java condition variables. Solutions will be based on the algorithm illustrated in Figure 7.7.

Both implementations will require creating five philosophers, each identified by a number 0 . . 4. Each philosopher will run as a separate thread. Philosophers alternate between thinking and eating. To simulate both activities, have each thread sleep for a random period between one and three seconds.

I. POSIX

Thread creation using Pthreads is covered in Section 4.4.1. When a philosopher wishes to eat, she invokes the function

```
pickup_forks(int philosopher_number)
```

where `philosopher_number` identifies the number of the philosopher wishing to eat. When a philosopher finishes eating, she invokes

```
return_forks(int philosopher_number)
```

Your implementation will require the use of POSIX condition variables, which are covered in Section 7.3.

II. Java

When a philosopher wishes to eat, she invokes the method `takeForks(philosopherNumber)`, where `philosopherNumber` identifies the number of the philosopher wishing to eat. When a philosopher finishes eating, she invokes `returnForks(philosopherNumber)`.

Your solution will implement the following interface:

```
public interface DiningServer
{
    /* Called by a philosopher when it wishes to eat */
    public void takeForks(int philosopherNumber);

    /* Called by a philosopher when it is finished eating */
    public void returnForks(int philosopherNumber);
}
```

It will require the use of Java condition variables, which are covered in Section 7.4.4.

ProducerConsumer

Project 4—The Producer–Consumer Problem

In Section 7.1.1, we presented a semaphore-based solution to the producer–consumer problem using a bounded buffer. In this project, you will design a programming solution to the bounded-buffer problem using the producer and consumer processes shown in Figures 5.9 and 5.10. The solution presented in Section 7.1.1 uses three semaphores: `empty` and `full`, which count the number of empty and full slots in the buffer, and `mutex`, which is a binary (or mutual-exclusion) semaphore that protects the actual insertion or removal of items in the buffer. For this project, you will use standard counting semaphores for `empty` and `full` and a mutex lock, rather than a binary semaphore, to represent `mutex`. The producer and consumer—running as separate threads—will move items to and from a buffer that is synchronized with the `empty`, `full`, and `mutex` structures. You can solve this problem using either Pthreads or the Windows API.

The Buffer

Internally, the buffer will consist of a fixed-size array of type `buffer_item` (which will be defined using a `typedef`). The array of `buffer_item` objects will be manipulated as a circular queue. The definition of `buffer_item`, along with the size of the buffer, can be stored in a header file such as the following:

```
/* buffer.h */
typedef int buffer_item;
#define BUFFER_SIZE 5
```

The buffer will be manipulated with two functions, `insert_item()` and `remove_item()`, which are called by the producer and consumer threads, respectively. A skeleton outlining these functions appears in Figure 7.14.

The `insert_item()` and `remove_item()` functions will synchronize the producer and consumer using the algorithms outlined in Figure 7.1 and Figure 7.2. The buffer will also require an initialization function that initializes the mutual-exclusion object `mutex` along with the `empty` and `full` semaphores.

The `main()` function will initialize the buffer and create the separate producer and consumer threads. Once it has created the producer and consumer threads, the `main()` function will sleep for a period of time and, upon awakening, will terminate the application. The `main()` function will be passed three parameters on the command line:

1. How long to sleep before terminating
2. The number of producer threads
3. The number of consumer threads

A skeleton for this function appears in Figure 7.15.

```
#include "buffer.h"

/* the buffer */
buffer_item buffer[BUFFER_SIZE];

int insert_item(buffer_item item) {
    /* insert item into buffer
       return 0 if successful, otherwise
       return -1 indicating an error condition */
}

int remove_item(buffer_item *item) {
    /* remove an object from buffer
       placing it in item
       return 0 if successful, otherwise
       return -1 indicating an error condition */
}
```

Figure 7.14 Outline of buffer operations.

The Producer and Consumer Threads

The producer thread will alternate between sleeping for a random period of time and inserting a random integer into the buffer. Random numbers will be produced using the `rand()` function, which produces random integers between 0 and `RAND_MAX`. The consumer will also sleep for a random period of time and, upon awakening, will attempt to remove an item from the buffer. An outline of the producer and consumer threads appears in Figure 7.16.

```
#include "buffer.h"

int main(int argc, char *argv[]) {
    /* 1. Get command line arguments argv[1],argv[2],argv[3] */
    /* 2. Initialize buffer */
    /* 3. Create producer thread(s) */
    /* 4. Create consumer thread(s) */
    /* 5. Sleep */
    /* 6. Exit */
}
```

Figure 7.15 Outline of skeleton program.

```
#include <stdlib.h> /* required for rand() */
#include "buffer.h"

void *producer(void *param) {
    buffer_item item;

    while (true) {
        /* sleep for a random period of time */
        sleep(...);
        /* generate a random number */
        item = rand();
        if (insert_item(item))
            fprintf("report error condition");
        else
            printf("producer produced %d\n",item);
    }

    void *consumer(void *param) {
        buffer_item item;

        while (true) {
            /* sleep for a random period of time */
            sleep(...);
            if (remove_item(&item))
                fprintf("report error condition");
            else
                printf("consumer consumed %d\n",item);
        }
    }
}
```

Figure 7.16 An outline of the producer and consumer threads.