

### Game Boy: Complete Technical Reference

gekkio https://gekkio.fi August 23, 2025

Revision 176



#### **Preface**

#### Caveat

IMPORTANT: This document focuses at the moment on 1st and 2nd generation devices (models before the Game Boy Color), and some hardware details are very different in later generations.

Be very careful if you make assumptions about later generation devices based on this document!

#### How to read this document

#### **\*** Speculation

This is something that hasn't been verified, but would make a lot of sense.

#### **i** Caveat

This explains some caveat about this documentation that you should know.

#### Warning

This is a warning about something.

#### 0.1 Formatting of numbers

When a single bit is discussed in isolation, the value looks like this: 0, 1.

Binary numbers are prefixed with 0b like this: 0b0101101, 0b11011, 0b00000000. Values are prefixed with zeroes when necessary, so the total number of digits always matches the number of digits in the value.

Hexadecimal numbers are prefixed with 0x like this: 0x1234, 0xDEADBEEF, 0xFF04. Values are prefixed with zeroes when necessary, so the total number of characters always matches the number of nibbles in the value.

#### Examples:

4-bit 8-bit 16-bit

Binary 0b0101 0b10100101 0b0000101010100101

Hexadecimal 0x5 0xA5 0x0AA5

#### 0.2 Register definitions

Register 0.1: 0x1234 - This is a hardware register definition

R/W-0	R/W-1	U-1	R-0	R-1	R-x	W-1	U-0
VALUE	<1:0>		BIGVAL<7:5>		FLAG		
bit 7	6	5	4	3	2	1	bit 0

#### Top row legend:

- **R** Bit can be read.
- **W** Bit can be written. If the bit cannot be read, reading returns a constant value defined in the bit list of the register in question.
- **U** Unimplemented bit. Writing has no effect, and reading returns a constant value defined in the bit list of the register in question.
- -n Value after system reset: 0, 1, or x.
- **1** Bit is set.
- Ø Bit is cleared.
- **x** Bit is unknown (e.g. depends on external things such as user input)

#### Middle row legend:

VALUE<1:0>	Bits 1 and 0 of VALUE
	Unimplemented bit
BIGVAL<7:5>	Bits 7, 6, 5 of BIGVAL
FLAG	Single-bit value FLAG

#### In this example:

- After system reset, VALUE is 0b01, BIGVAL is either 0b010 or 0b011, FLAG is 0b1.
- Bits 5 and 0 are unimplemented. Bit 5 always returns 1, and bit 0 always returns 0.
- Both bits of VALUE can be read and written. When this register is written, bit 7 of the written value goes to bit 1 of VALUE.
- FLAG can only be written to, so reads return a value that is defined elsewhere.
- BIGVAL cannot be written to. Only bits 5-7 of BIGVAL are defined here, so look elsewhere for the low bits 0-4.

### **Contents**

	Preface	. 2
	How to read this document	. 3
	0.1 Formatting of numbers	. 3
	0.2 Register definitions	. 4
C	Contents	. 5
	Game Boy console architecture	
	Introduction	
2	Clocks	
	2.1 System clock	
	System clock frequency	
	2.2 Clock periods, T-cycles, and M-cycles	12
IJ	I Sharp SM83 CPU core	14
3	Introduction	15
	3.1 History	
	Simple model	
5	6 CPU core timing	
	5.1 Fetch/execute overlap	
_	Fetch/execute overlap timing example	
6	Sharp SM83 instruction set	
	6.1 Overview	
	CB opcode prefix	
	Undefined opcodes	
	6.2 8-bit load instructions	
	LD r, r': Load register (register)	
	LD r, n: Load register (immediate)	
	LD (HL), r: Load from register (indirect HL)	
	LD (HL), n: Load from immediate data (indirect HL)	
	LD A, (BC): Load accumulator (indirect BC)	
	LD A, (DE): Load accumulator (indirect DE)	
	LD (BC), A: Load from accumulator (indirect BC)	
	LD (DE), A: Load from accumulator (indirect DE)	
	LD A, (nn): Load accumulator (direct)	
	LD (nn), A: Load from accumulator (direct)	
	LDH A, (C): Load accumulator (indirect 0xFF00+C)	
	LDH (C), A: Load from accumulator (indirect 0xFF00+C)	
	LDH A, (n): Load accumulator (direct 0xFF00+n)	
	LDH (n), A: Load from accumulator (direct 0xFF00+n)	
	LD A, (HL-): Load accumulator (indirect HL, decrement)	35
	LD (HL-), A: Load from accumulator (indirect HL, decrement)	36
	LD A, (HL+): Load accumulator (indirect HL, increment)	
	LD (HL+), A: Load from accumulator (indirect HL, increment)	
	6.3 16-bit load instructions	39
	LD rr, nn: Load 16-bit register / register pair	39

LD (nn), SP: Load from stack pointer (direct)			
LD SP, HL: Load stack pointer from HL	 		41
PUSH rr: Push to stack	 		42
POP rr: Pop from stack	 		43
LD HL, SP+e: Load HL from adjusted stack pointer	 		44
6.4 8-bit arithmetic and logical instructions	 		45
ADD r: Add (register)	 		45
ADD (HL): Add (indirect HL)	 		46
ADD n: Add (immediate)			
ADC r: Add with carry (register)	 		48
ADC (HL): Add with carry (indirect HL)			
ADC n: Add with carry (immediate)			
SUB r: Subtract (register)			
SUB (HL): Subtract (indirect HL)			
SUB n: Subtract (immediate)			
SBC r: Subtract with carry (register)			
SBC (HL): Subtract with carry (indirect HL)			
SBC n: Subtract with carry (immediate)			
CP r: Compare (register)			
CP (HL): Compare (indirect HL)			
CP n: Compare (immediate)			
INC r: Increment (register)			
INC (HL): Increment (indirect HL)			
DEC r: Decrement (register)			
DEC (HL): Decrement (indirect HL)			
AND (III) Bitwise AND (register)			
AND (HL): Bitwise AND (indirect HL)			
AND n: Bitwise AND (immediate)			
OR r: Bitwise OR (register)			
OR (HL): Bitwise OR (indirect HL)			
OR n: Bitwise OR (immediate)			
XOR r: Bitwise XOR (register)			
XOR (HL): Bitwise XOR (indirect HL)			
XOR n: Bitwise XOR (immediate)			
CCF: Complement carry flag			
SCF: Set carry flag			
DAA: Decimal adjust accumulator			
CPL: Complement accumulator			
6.5 16-bit arithmetic instructions			
INC rr: Increment 16-bit register			
DEC rr: Decrement 16-bit register			
ADD HL, rr: Add (16-bit register)			
ADD SP, e: Add to stack pointer (relative)			
6.6 Rotate, shift, and bit operation instructions			
RLCA: Rotate left circular (accumulator)			
RRCA: Rotate right circular (accumulator)			
RLA: Rotate left (accumulator)			
RRA: Rotate right (accumulator)			
RLC r: Rotate left circular (register)			86

RLC (HL): Rotate left circular (indirect HL)	 . 87
RRC r: Rotate right circular (register)	 . 89
RRC (HL): Rotate right circular (indirect HL)	 . 90
RL r: Rotate left (register)	
RL (HL): Rotate left (indirect HL)	 . 93
RR r: Rotate right (register)	
RR (HL): Rotate right (indirect HL)	
SLA r: Shift left arithmetic (register)	
SLA (HL): Shift left arithmetic (indirect HL)	
SRA r: Shift right arithmetic (register)	
SRA (HL): Shift right arithmetic (indirect HL)	
SWAP r: Swap nibbles (register)	
SWAP (HL): Swap nibbles (indirect HL)	
SRL r: Shift right logical (register)	
SRL (HL): Shift right logical (indirect HL)	
BIT b, r: Test bit (register)	
BIT b, (HL): Test bit (indirect HL)	
RES b, r: Reset bit (register)	
RES b, (HL): Reset bit (indirect HL)	
SET b, r: Set bit (register)	
SET b, (HL): Set bit (indirect HL)	
6.7 Control flow instructions	
JP nn: Jump	
JP HL: Jump to HL	
JP cc, nn: Jump (conditional)	
JR e: Relative jump	
JR cc, e: Relative jump (conditional)	
CALL nn: Call function	
CALL cc, nn: Call function (conditional)	
RET: Return from function	
RET cc: Return from function (conditional)	
RETI: Return from interrupt handler	
RST n: Restart / Call function (implied)	
6.8 Miscellaneous instructions	
HALT: Halt system clock	
STOP: Stop system and main clocks	 . 130
DI: Disable interrupts	
EI: Enable interrupts	 . 131
NOP: No operation	 . 132
III Game Boy SoC peripherals and features	
7 Boot ROM	
7.1 Boot ROM types	 . 136
DMG boot ROM	 . 136
MGB boot ROM	 136
SGB boot ROM	 . 136
SGB2 boot ROM	 . 136
Early DMG boot ROM ("DMG0")	 . 136
8 DMA (Direct Memory Access)	137

8.1 Object Attribute Memory (OAM) DMA	137
OAM DMA address decoding	138
OAM DMA transfer timing	138
OAM DMA bus conflicts	138
9 PPU (Picture Processing Unit)	139
10 Port P1 (Joypad, Super Game Boy communication)	
11 Serial communication	
IV Game Boy game cartridges	142
12 MBC1 mapper chip	
12.1 MBC1 registers	
12.2 ROM in the 0x0000–0x7FFF area	
ROM banking example 1	
ROM banking example 2	
12.3 RAM in the 0xA000–0xBFFF area	
RAM banking example 1	
12.4 MBC1 multicarts ("MBC1M")	
ROM banking example 1	
Detecting multicarts	
12.5 Dumping MBC1 carts	
13 MBC2 mapper chip	
13.1 MBC2 registers	149
13.2 ROM in the 0x0000–0x7FFF area	150
13.3 RAM in the 0xA000–0xBFFF area	150
13.4 Dumping MBC2 carts	151
14 MBC3 mapper chip	152
15 MBC30 mapper chip	153
16 MBC5 mapper chip	
16.1 MBC5 registers	
17 MBC6 mapper chip	
18 MBC7	
19 HuC-1 mapper chip	
20 HuC-3 mapper chip	
21 MMM01	
22 TAMA5	
ZZ TAIVIAS	101
Appendices	162
A Instruction set tables	
B Memory map tables	
C Game Boy external bus	
C.1 Bus timings	
D Chip pinouts	
D.1 CPU chips	
D.2 Cartridge chips	173
PSLP	4= -
Bibliography	174

# Part I Game Boy console architecture

### Chapter 1 Introduction

The original Game Boy and its successors were the most popular and financially successful handheld consoles in the 1990s and early 2000s with several millions units sold and a large catalogue of officially published games. Unlike many older consoles, Game Boys use only a single integrated System-on-a-Chip (SoC) for almost everything, and this SoC includes the processor (CPU) core, some memories, and various peripherals.

#### **1** Caveat

The Game Boy SoC is sometimes called the "CPU", even though it has a large amount of other peripherals as well. For example, the Game Boy Pocket SoC literally has the text "CPU MGB" on it, even though the CPU core takes only a small fraction of the entire chip area. This terminology is therefore misleading, and is like calling a computer mother-board and all connected expansion cards and storage devices the "CPU".

This document always makes a clear distiction between the entire chip (SoC) and the processor inside it (the CPU *core*).

Most Game Boy consoles are handhelds, starting from the original Game Boy in 1989, ending with the Game Boy Micro in 2005. In addition to handheld devices, Game Boy SoCs are also used in some accessories meant for other consoles, such as the Super Game Boy for the SNES/SFC.

Game Boy consoles and their SoCs can be categorized based on three supported technical architectures:

- GB: the original Game Boy architecture with a Sharp SM83 CPU core and 4-level grayscale graphics
- GBC: a mostly backwards compatible extension to the GB architecture that adds color graphics and small improvements
- GBA: a completely different architecture based on the ARM processor instruction set and a
  completely redesigned set of peripherals. This document does not cover GBA architecture,
  because it has little in common with GB/GBC. GBA-based consoles and chips are only
  mentioned for their backwards compatibility with GB/GBC architectures.

Table 1.1 lists all officially released Game Boy consoles, including handhelds and accessories for other consoles. Every model has an internal codename, such as original Game Boy's codename Dot Matrix Game (DMG), that is also present on the mainboard.

#### **6** Caveat

This document refers to different console models usually by their unique codename to prevent confusion. For example, using the abbreviation GBP could refer to either Game Boy Pocket or Game Boy Player, but there's no confusion when MGB and GBS are used instead.

In this document GBC refers to the technical architecture, while CGB refers to Game Boy Color consoles specifically. Likewise, GBA refers to the architecture and AGB to exactly one console model.

Console name	Codename	SoC type	GB	GBC	GBA
Handhelds					
Game Boy	DMG	DMG-CPU	1		
Game Boy Pocket	MGB	CPU MGB	1		
Game Boy Light	MGL	CPU MGB	1		
Game Boy Color	CGB	CPU CGB	1	1	
Game Boy Advance	AGB	CPU AGB	1	1	1
Game Boy Advance SP	AGS	CPU AGB	1	1	1
Game Boy Micro	OXY	CPU AGB			1
Accessories					
Super Game Boy	SGB	SGB-CPU	1		
Super Game Boy 2	SGB2	CPU SGB2	1		
Game Boy Player	GBS	CPU AGB	1	1	1

Table 1.1: Summary of Game Boy consoles

### Chapter 2 Clocks

#### 2.1 System clock

The system oscillator is the primary clock source in a Game Boy system, and it generates the **system clock**. Almost all other clocks are derived from the system clock using prescalers / clock dividers, but there are some exceptions:

- If a Game Boy is set up to do a serial transfer in secondary mode, the serial data register
  is directly clocked using the serial clock signal coming from the link port. Two Game Boys
  connected with a link cable never have precisely the same clock phase and frequency relative
  to each other, so the serial clock of the primary side has no direct relation to the system clock
  of the secondary side.
- The inserted game cartridge may use other clock(s) internally. A typical example in some official games is the Real Time Clock (RTC), which is based on a 32.768 kHz oscillator and a clock-domain crossing circuit so that RTC data can be read using the cartridge bus while the RTC circuit is ticking independently using its own clock.

The Game Boy SoC uses two pins for the system oscillator: XI and XO. These pins along with some external components can be used to form a Pierce oscillator circuit. Alternatively, the XI pin can be driven directly with a clock signal originating from somewhere else, and the XO pin can be left unconnected.

#### System clock frequency

In DMG and MGB consoles the system oscillator circuit uses an external quartz crystal with a nominal frequency of **4.194304 MHz** (=  $2^{22}$  MHz = 4 MiHz) to form a Pierce oscillator circuit. This frequency is considered to be the standard frequency of a Game Boy.

In SGB the system oscillator input is directly driven by the ICD2 chip on the SGB cartridge. The clock is derived via /5 division of the main SNES / SFC clock, which has a different frequency depending on the console region (21.447 MHz NTSC, 21.281 MHz PAL). The SNES / SFC clock does not divide into 4.194304 MHz with integer division, so the clock seen by the SGB SoC is not the same as in DMG and MGB consoles. The frequency is higher, so everything is sped up by a small amount and audio has a slightly higher pitch.

In SGB2, just like SGB, the system oscillator input is driven by the ICD2 chip, but instead of using the SNES / SFC clock, the ICD2 chip is driven by a Pierce oscillator circuit with a 20.971520 MHz crystal. ICD2 then divides this frequency by /5 to obtain the final frequency seen by the SGB2 SoC, which is 4.194304 MHz that matches the standard DMG / MGB frequency.

#### 2.2 Clock periods, T-cycles, and M-cycles

In digital logic, a clock switches between low and high states and every transition happens on a *clock edge*, which might be a rising edge (low  $\rightarrow$  high transition) or a falling edge (high  $\rightarrow$  low transition). A single *clock period* is measured between two edges of the same type, so that the clock goes through two opposing edges and returns to its original state after the clock period. The typical convention is that a clock period consists of a rising edge and a falling edge.

In addition to the system clock and other clocks derived from it, Game Boy systems also use *inverted clocks* in some peripherals, which means the rising edge of an inverted clock may happen at the same time as a falling edge of the original clock. Figure 2.1 shows two clock

periods of the system clock and an inverted clock derived from it, and how they are out of phase due to clock inversion.

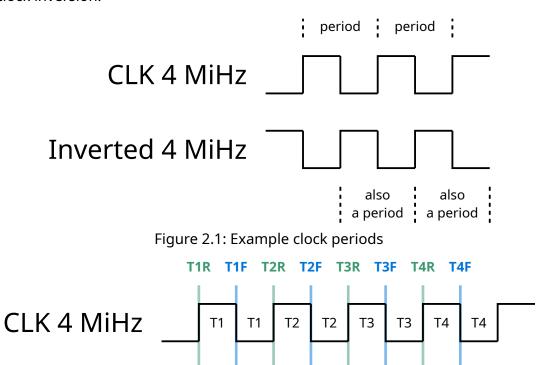


Figure 2.2: Clock edges in a machine cycle

PHI 1 MiHz

# Part II Sharp SM83 CPU core

### Chapter 3 Introduction

The CPU core in the Game Boy SoC is a custom Sharp design that hasn't publicly been given a name by either Sharp or Nintendo. However, using old Sharp datasheets and databooks as evidence, the core has been identified to be a Sharp **SM83** CPU core, or at least something that is 100% compatible with it. SM83 is a custom CPU core used in some custom Application Specific Integrated Chips (ASICs) manufactured by Sharp in the 1980s and 1990s.

#### Warning

Some sources claim Game Boy uses a "modified" Zilog Z80 or Intel 8080 CPU core. While the SM83 resembles both and has many identical instructions, it can't execute all Z80/8080 programs, and finer details such as timing of instructions often differ.

SM83 is an 8-bit CPU core with a 16-bit address bus. The Instruction Set Architecture (ISA) is based on both Z80 and 8080, and is close enough to Z80 that programmers familiar with Z80 assembly can quickly become productive with SM83 as well. Some Z80 programs may also work directly on SM83, assuming only opcodes supported by both are used and the program is not sensitive to timing differences.

#### **\*** Speculation

Sharp most likely designed SM83 to closely resemble Z80, so it would be easy for programmers already familiar with the widely popular Z80 to write programs for it. However, SM83 is not a "modified Z80" because the internal implementation is completely different. At the time Sharp also manufactured real Z80 chips such as LH0080 under a license from Zilog, so they were familiar with Z80 internals but did not directly copy the actual implementation of the CPU core. If you compare photos of a decapped Z80 chip and a GB SoC, you will see two very different-looking CPU cores.

#### 3.1 History

The first known mention of the SM83 CPU core is in Sharp Microcomputers Data Book (1990), where it is listed as the CPU core used in the SM8320 8-bit microcomputer chip, intended for inverter air conditioners [1]. The data book describes some details of the CPU core, such as a high-level overview of the supported instructions, but precise details such as full opcode tables are not included. Another CPU core called SM82 is also mentioned, but based on the details it's clearly a completely different one.

The SM83 CPU core later appeared in Sharp Microcomputer Data Book (1996), where it is listed as the CPU core in the SM8311/SM8313/SM8314/SM8315 8-bit microcomputer chips, meant for home appliances [2]. This data book describes the CPU core in much more detailed manner, and other than some mistakes in the descriptions, the details seem to match what is known about the GB SoC CPU core from other sources.

### Chapter 4 Simple model

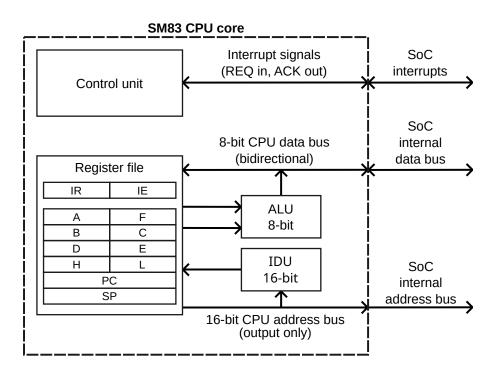


Figure 4.3: Simple model of the SM83 CPU core

Figure 4.3 shows a simplified model of the SM83 CPU core. The core interacts with the rest of the SoC using interrupt signals, an 8-bit bidirectional data bus, and a 16-bit address bus controlled by the CPU core.

The main subsystems of the CPU core are as follows:

**Control unit** The control unit decodes the executed instructions and generates control signals for the rest of the CPU core. It is also responsible for checking and dispatching interrupts.

**Register file** The register file holds most of the state of the CPU inside registers. It contains the 16-bit Program Counter (PC), the 16-bit Stack Pointer (SP), the 8-bit Accumulator (A), the Flags register (F), general-purpose register pairs consisting of two 8-bit halves such as BC, DE, HL, and the special-purpose 8-bit registers Instruction Register (IR) and Interrupt Enable (IE).

ALU An 8-bit Arithmetic Logic Unit (ALU) has two 8-bit input ports and is capable of performing various calculations. The ALU outputs its result either to the register file or the CPU data bus.

A dedicated 16-bit Increment/Decrement Unit (IDU) is capable of performing only simple increment/decrement operations on the 16-bit address bus value, but they can be performed independently of the ALU, improving maximum performance of the CPU core. The IDU always outputs its result back to the register file, where it can be written to a register pair or a 16-bit register.

## Chapter 5 CPU core timing

#### 5.1 Fetch/execute overlap

Sharp SM83 uses a microprocessor design technique known as *fetch/execute overlap* to improve CPU performance by doing opcode fetches in parallel with instruction execution whenever possible. Since the CPU can only perform one memory access per M-cycle, it is worth it to try to do memory operations as soon as possible. Also, when doing a memory read, the CPU cannot use the data during the same M-cycle so the true minimum effective duration of instructions is 2 machine cycles, not 1 machine cycle.

Every instruction needs one machine cycle for the fetch stage, and at least one machine cycle for the decode/execute stage. However, the fetch stage of an instruction always overlaps with the last machine cycle of the execute stage of the previous instruction. The overlapping execute stage cycle may still do some work (e.g. ALU operation and/or register writeback) but memory access is reserved for the fetch stage of the next instruction.

Since all instructions effectively last one machine cycle longer, fetch/execute overlap is usually ignored in documentation intended for programmers. It is much easier to think of a program as a sequence of non-overlapping instructions and consider only the execute stages when calculating instruction durations. However, when emulating a SM83 CPU core, understanding and emulating the overlap can be useful.

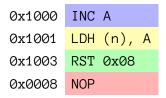
#### Warning

Sharp SM831x is a family of single-chip SoCs from Sharp that use the SM83 CPU core, and their datasheet [3] includes a description of fetch/execute overlap. However, the description is not completely correct and can in fact be misleading.

For example, the timing diagram includes an instruction that does not involve opcode fetch at all, and memory operations for two instructions are shown to happen at the same time, which is not possible.

#### Fetch/execute overlap timing example

Let's assume the CPU is executing a program that starts from the address 0x1000 and contains the following instructions:



The following timing diagram shows all memory operations done by the CPU, and the fetch and execute stages of each instruction:

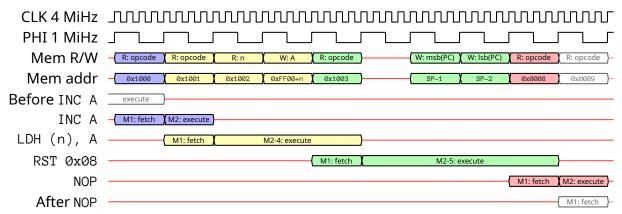


Figure 5.4: Fetch/execute overlap example

## **Chapter 6 Sharp SM83 instruction set**

6.1 Overview

**CB** opcode prefix

**Undefined opcodes** 

#### 6.2 8-bit load instructions

#### LD r, r': Load register (register)

Load to the 8-bit register r, data from the 8-bit register r'.

Opcode ObO1xxxyyy/various Duration 1 machine cycle Length 1 byte: opcode Flags - Simple timing and pseudocode

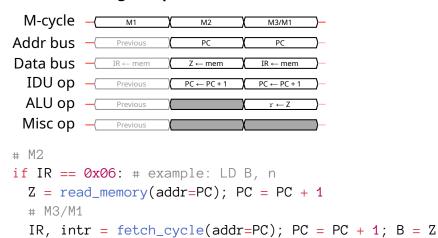
```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x41: # example: LD B, C
B = C
```

#### LD r, n: Load register (immediate)

Load to the 8-bit register r, the immediate data n.

```
Opcode 0b00xxx110/various
Length 2 bytes: opcode + n
Simple timing and pseudocode
```

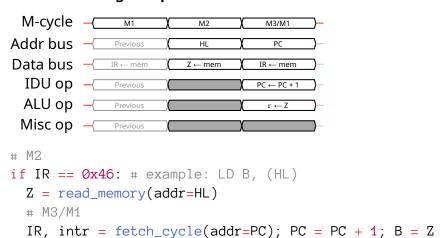
```
Duration 2 machine cycles Flags -
```



#### LD r, (HL): Load register (indirect HL)

Load to the 8-bit register r, data from the absolute address specified by the 16-bit register HL.

```
M-cycle — ( M1 ) M2 — Mem R/W — Opcode ( R:data ) — Opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x46: # example: LD B, (HL)
B = read_memory(addr=HL)
```



#### LD (HL), r: Load from register (indirect HL)

Load to the absolute address specified by the 16-bit register HL, data from the 8-bit register r.

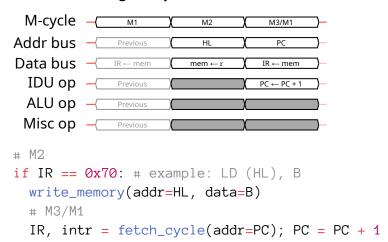
```
M-cycle — ( M1 ) M2 )—

Mem R/W — ( opcode ) ( W:data )—

opcode = read_memory(addr=PC); PC = PC + 1

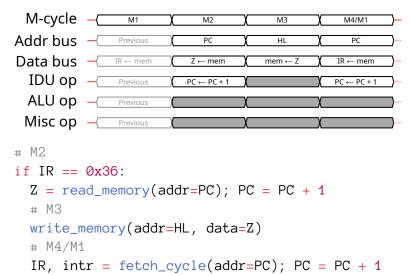
if opcode == 0x70: # example: LD (HL), B

write_memory(addr=HL, data=B)
```



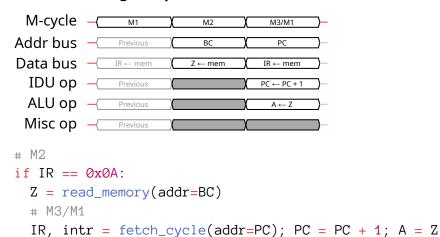
#### LD (HL), n: Load from immediate data (indirect HL)

Load to the absolute address specified by the 16-bit register HL, the immediate data n.



#### LD A, (BC): Load accumulator (indirect BC)

Load to the 8-bit A register, data from the absolute address specified by the 16-bit register BC.



#### LD A, (DE): Load accumulator (indirect DE)

Load to the 8-bit A register, data from the absolute address specified by the 16-bit register DE.

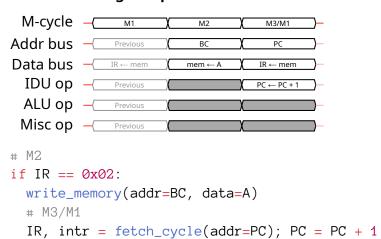
```
Opcode0b00011010/0x1ADuration2 machine cyclesLength1 byte: opcodeFlags-Simple timing and pseudocodeFlags-
```

```
M-cycle — (M1 ) (M2 ) (M3/M1) —
Addr bus — Previous (DE ) PC —
Data bus — IR — mem (Z — mem ) IR — mem —
IDU op — Previous (PC — PC + 1) —
ALU op — Previous (A — Z —
Misc op — Previous (A — Z —
Misc op — Previous (A — Z —

# M2
if IR == Øx1A:
Z = read_memory(addr=DE)
# M3/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1; A = Z
```

#### LD (BC), A: Load from accumulator (indirect BC)

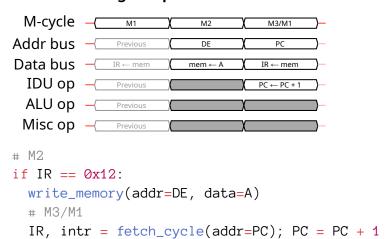
Load to the absolute address specified by the 16-bit register BC, data from the 8-bit A register.



#### LD (DE), A: Load from accumulator (indirect DE)

Load to the absolute address specified by the 16-bit register DE, data from the 8-bit A register.

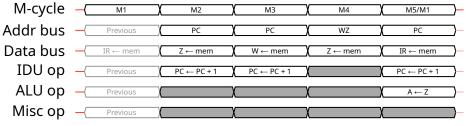
Opcode 0b00010010/0x12 Duration 2 machine cycles Length 1 byte: opcode Flags - Simple timing and pseudocode



#### LD A, (nn): Load accumulator (direct)

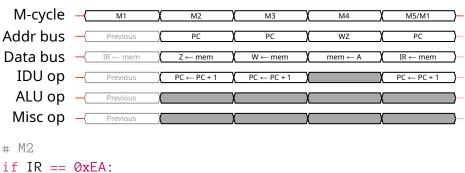
Load to the 8-bit A register, data from the absolute address specified by the 16-bit operand nn.

Opcode 0b11111010/0xFA Duration 4 machine cycles Length 3 bytes: opcode + LSB(nn) + MSB(nn) Flags - Simple timing and pseudocode



#### LD (nn), A: Load from accumulator (direct)

Load to the absolute address specified by the 16-bit operand nn, data from the 8-bit A register.



```
if IR == 0xEA:
   Z = read_memory(addr=PC); PC = PC + 1
# M3
W = read_memory(addr=PC); PC = PC + 1
# M4
write_memory(addr=WZ, data=A)
# M5/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### LDH A, (C): Load accumulator (indirect 0xFF00+C)

if IR == 0xF2:

# M3/M1

Load to the 8-bit A register, data from the address specified by the 8-bit C register. The full 16-bit absolute address is obtained by setting the most significant byte to 0xFF and the least significant byte to the value of C, so the possible range is 0xFF00-0xFFFF.

```
Opcode 0b11110010/0xF2
                                       Duration 2 machine cycles
Length 1 byte: opcode
                                       Flags
Simple timing and pseudocode
 M-cycle –( M1 )
Mem R/W — opcode C R: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xF2:
 A = read_memory(addr=unsigned_16(lsb=C, msb=0xFF))
Detailed timing and pseudocode
 M-cycle ( M1 ) M2
Addr bus Previous
                    0xFF00+C
Data bus — IR ← mem / IR ← mem / IR ← mem
 IDU op - Previous
 ALU op — Previous A ← Z
 Misc op - Previous
# M2
```

 $Z = read\_memory(addr=unsigned\_16(lsb=C, msb=0xFF))$ 

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1; A = Z

#### LDH (C), A: Load from accumulator (indirect 0xFF00+C)

Load to the address specified by the 8-bit C register, data from the 8-bit A register. The full 16-bit absolute address is obtained by setting the most significant byte to 0xFF and the least significant byte to the value of C, so the possible range is 0xFF00-0xFFFF.

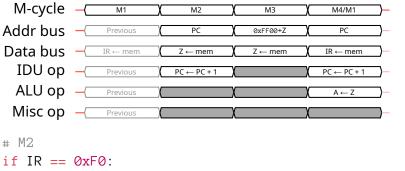
```
Opcode 0b11100010/0xE2
                                        Duration 2 machine cycles
Length 1 byte: opcode
                                        Flags
Simple timing and pseudocode
  M-cycle — M1 X
Mem R/W — opcode W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xE2:
  write_memory(addr=unsigned_16(lsb=C, msb=0xFF), data=A)
Detailed timing and pseudocode
 M-cycle — M1 X
                     M2
Addr bus — Previous
                     0xFF00+C
Data bus — IR ← mem / Mem ← A / IR ← mem
 IDU op - Previous
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xE2:
  write_memory(addr=unsigned_16(lsb=C, msb=0xFF), data=A)
 # M3/M1
```

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

#### LDH A, (n): Load accumulator (direct 0xFF00+n)

Load to the 8-bit A register, data from the address specified by the 8-bit immediate data n. The full 16-bit absolute address is obtained by setting the most significant byte to 0xFF and the least significant byte to the value of n, so the possible range is 0xFF00-0xFFFF.

Opcode 0b11110000/0xF0 Duration 3 machine cycles Length 2 bytes: opcode + n Flags - Simple timing and pseudocode

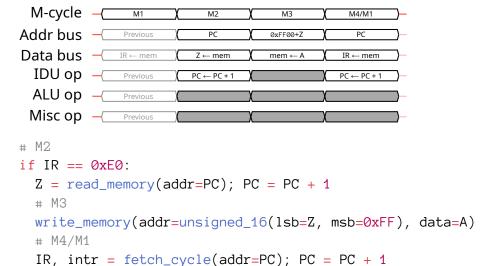


```
f IR == 0xF0:
Z = read_memory(addr=PC); PC = PC + 1
# M3
Z = read_memory(addr=unsigned_16(lsb=Z, msb=0xFF))
# M4/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1; A = Z
```

#### LDH (n), A: Load from accumulator (direct 0xFF00+n)

Load to the address specified by the 8-bit immediate data n, data from the 8-bit A register. The full 16-bit absolute address is obtained by setting the most significant byte to 0xFF and the least significant byte to the value of n, so the possible range is 0xFF00-0xFFFF.

Opcode 0b11100000/0xE0 Duration 3 machine cycles Length 2 bytes: opcode + n Flags - Simple timing and pseudocode



#### LD A, (HL-): Load accumulator (indirect HL, decrement)

Load to the 8-bit A register, data from the absolute address specified by the 16-bit register HL. The value of HL is decremented after the memory read.

**Opcode** 0b00111010/0x3A **Duration** 2 machine cycles Length 1 byte: opcode Flags Simple timing and pseudocode M-cycle (M1) (M2) Mem R/W \_ ( opcode ) R: data } opcode = read\_memory(addr=PC); PC = PC + 1 if opcode == 0x3A:  $A = read_memory(addr=HL); HL = HL - 1$ Detailed timing and pseudocode M-cycle –( M1 M2 M3/M1 Addr bus Previous HL PC Data bus — R←mem Z←mem IR←mem IDU op - Previous ( HL ← HL - 1 ( PC ← PC + 1 ) ALU op — Previous A ← Z Misc op Previous # M2 if IR == 0x3A:  $Z = read_memory(addr=HL); HL = HL - 1$ # M3/M1

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1; A = Z

#### LD (HL-), A: Load from accumulator (indirect HL, decrement)

Load to the absolute address specified by the 16-bit register HL, data from the 8-bit A register. The value of HL is decremented after the memory write.

**Opcode** 0b00110010/0x32 **Duration** 2 machine cycles Length 1 byte: opcode Flags Simple timing and pseudocode M-cycle (M1) (M2) Mem R/W \_ ( opcode ) W: data } opcode = read\_memory(addr=PC); PC = PC + 1 if opcode == 0x32: write\_memory(addr=HL, data=A); HL = HL - 1 **Detailed timing and pseudocode** M-cycle –( M1 ) M2 M3/M1 Addr bus Previous HL **X** PC Data bus — R←mem mem←A IR←mem IDU op - Previous ( HL ← HL - 1 ) PC ← PC + 1 ALU op Previous Misc op Previous # M2 if IR == 0x32: write\_memory(addr=HL, data=A); HL = HL - 1 # M3/M1

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

# LD A, (HL+): Load accumulator (indirect HL, increment)

Load to the 8-bit A register, data from the absolute address specified by the 16-bit register HL. The value of HL is incremented after the memory read.

**Opcode** 0b00101010/0x2A **Duration** 2 machine cycles Length 1 byte: opcode Flags Simple timing and pseudocode M-cycle (M1) (M2) Mem R/W \_ ( opcode ) R: data } opcode = read\_memory(addr=PC); PC = PC + 1 if opcode == 0x2A:  $A = read_memory(addr=HL); HL = HL + 1$ Detailed timing and pseudocode M-cycle — M1 X M2 M3/M1 Addr bus Previous HL PC Data bus — R←mem Z←mem IR←mem IDU op - Previous ( HL ← HL + 1 ( PC ← PC + 1 ) ALU op — Previous A ← Z Misc op — Previous # M2 if IR == 0x2A:  $Z = read_memory(addr=HL); HL = HL + 1$ # M3/M1

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1; A = Z

# LD (HL+), A: Load from accumulator (indirect HL, increment)

Load to the absolute address specified by the 16-bit register HL, data from the 8-bit A register. The value of HL is incremented after the memory write.

**Opcode** 0b00100010/0x22 **Duration** 2 machine cycles Length 1 byte: opcode Flags Simple timing and pseudocode M-cycle (M1) (M2) Mem R/W \_ ( opcode ) W: data } opcode = read\_memory(addr=PC); PC = PC + 1 if opcode == 0x22: write\_memory(addr=HL, data=A); HL = HL + 1 **Detailed timing and pseudocode** M-cycle –( M1 ) M2 M3/M1 Addr bus Previous HL PC Data bus — R←mem mem←A IR←mem IDU op - Previous  $HL \leftarrow HL + 1$   $PC \leftarrow PC + 1$ ALU op Previous Misc op — Previous # M2 if IR == 0x22: write\_memory(addr=HL, data=A); HL = HL + 1 # M3/M1

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

# 6.3 16-bit load instructions

# LD rr, nn: Load 16-bit register / register pair

Load to the 16-bit register rr, the immediate 16-bit data nn.

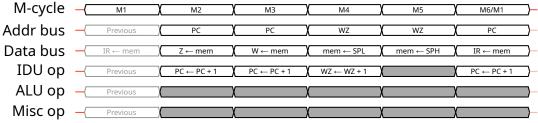
Opcode <a href="Oblight style="block">Oblight style="block">Obligh

```
M-cycle — ( M1 ) M2 ) M3
                                          M4/M1
Addr bus Previous PC
                                  PC
Data bus \sqrt{R \leftarrow mem} Z \leftarrow mem W \leftarrow mem R \leftarrow mem
 IDU op - Previous
                   PC \leftarrow PC + 1 PC \leftarrow PC + 1 PC \leftarrow PC + 1
 ALU op Previous
 MISC OD — Previous ( rr ← WZ )—
# M2
if IR == 0x01: # example: LD BC, nn
 Z = read_memory(addr=PC); PC = PC + 1
  # M3
 W = read_memory(addr=PC); PC = PC + 1
  # M4/M1
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1; BC = WZ
```

#### LD (nn), SP: Load from stack pointer (direct)

Load to the absolute address specified by the 16-bit operand nn, data from the 16-bit SP register.

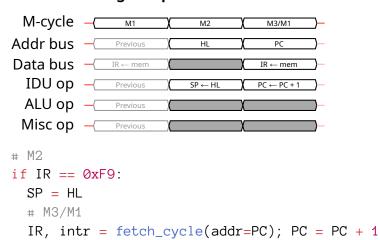
Opcode 0b00001000/0x08 Duration 5 machine cycles Length 3 bytes: opcode + LSB(nn) + MSB(nn) Flags - Simple timing and pseudocode



```
# M2
if IR == 0x08:
    Z = read_memory(addr=PC); PC = PC + 1
    # M3
W = read_memory(addr=PC); PC = PC + 1
    # M4
    write_memory(addr=WZ, data=lsb(SP)); WZ = WZ + 1
    # M5
    write_memory(addr=WZ, data=msb(SP))
# M6/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# LD SP, HL: Load stack pointer from HL

Load to the 16-bit SP register, data from the 16-bit HL register.



# **PUSH rr: Push to stack**

Push to the stack memory, data from the 16-bit register rr.



```
# M2
if IR == 0xC5: # example: PUSH BC
  SP = SP - 1
# M3
  write_memory(addr=SP, data=msb(BC)); SP = SP - 1
# M4
  write_memory(addr=SP, data=lsb(BC))
# M5/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **POP rr: Pop from stack**

Pops to the 16-bit register rr, data from the stack memory.

This instruction does not do calculations that affect flags, but POP AF completely replaces the F register value, so all flags are changed based on the 8-bit data that is read from memory.

Opcode 0b11xx0001/various Length 1 byte: opcode Simple timing and pseudocode **Duration** 3 machine cycles**Flags** See the instruction description

```
M-cycle –( M1 X
                        M2 X
                                   М3
                                             M4/M1
Addr bus Previous
                         SP
                                    SP
Data bus — IR ← mem Z ← mem W ← mem IR ← mem
 IDU op — Previous SP \leftarrow SP + 1 SP \leftarrow SP + 1 PC \leftarrow PC + 1
 ALU op - Previous
 Misc op - Previous
if IR == 0xC1: # example: POP BC
  Z = read_memory(addr=SP); SP = SP + 1
  # M3
  W = read_memory(addr=SP); SP = SP + 1
  # M4/M1
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1; BC = WZ
```

#### LD HL, SP+e: Load HL from adjusted stack pointer

Load to the HL register, 16-bit data calculated by adding the signed 8-bit operand e to the 16-bit value of the SP register.

Opcode 0b11111000/0xF8Length 2 bytes: opcode + e
Simple timing and pseudocode

Duration 3 machine cycles  $Z = 0, N = 0, H = \bigstar, C = \bigstar$ 

# M-cycle — M1 X M2 X M3 — Mem R/W — Opcode X R:e X — Opcode = read\_memory(addr=PC); PC = PC + if opcode == 0xF8:

```
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xF8:
    e = signed_8(read_memory(addr=PC)); PC = PC + 1
    result, carry_per_bit = SP + e
    HL = result
    flags.Z = 0
    flags.N = 0
    flags.H = 1 if carry_per_bit[3] else 0
    flags.C = 1 if carry_per_bit[7] else 0
```

```
M-cycle ( M1 ) M2
                               М3
Addr bus Previous
                       PC
Data bus — R←mem Z←mem Z←mem
 IDU op — Previous PC ← PC + 1
 ALU op Previous
                   ( L ← SPL + Z
 Misc op — Previous
# M2
if IR == 0xF8:
  Z = read_memory(addr=PC); PC = PC + 1
  result, carry_per_bit = lsb(SP) + Z
  L = result
  flags.Z = 0
  flags.N = 0
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  Z_{sign} = bit(7, Z)
  # M4/M1
  adj = 0xFF if Z_sign else 0x00
  result = msb(SP) + adj + flags.C
  H = result
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### 6.4 8-bit arithmetic and logical instructions

# ADD r: Add (register)

Adds to the 8-bit A register, the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10000xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — PC + 1

if opcode == 0x80: # example: ADD B

result, carry_per_bit = A + B

A = result

flags.Z = 1 if result == 0 else 0

flags.N = 0

flags.H = 1 if carry_per_bit[3] else 0

flags.C = 1 if carry_per_bit[7] else 0
```

```
M-cycle – M1
                      M2/M1
Addr bus Previous
Data bus — R←mem R ← mem —
 IDU op — Previous PC ← PC + 1
 ALU op Previous A - A+r
 Misc op — Previous
# M2/M1
if IR == 0x80: # example: ADD B
 result, carry_per_bit = A + B
 A = result
 flags.Z = 1 if result == 0 else 0
 flags.N = 0
 flags.H = 1 if carry_per_bit[3] else 0
 flags.C = 1 if carry_per_bit[7] else 0
 IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### ADD (HL): Add (indirect HL)

Adds to the 8-bit A register, data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10000110/0x86 Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                       HL
Data bus — IR ← mem Z ← mem IR ← mem
 IDU op — Previous PC ← PC+1
 ALU op \overline{} Previous \overline{} A \leftarrow A + Z
 Misc op — Previous
# M2
if IR == 0x86:
 Z = read_memory(addr=HL)
 # M3/M1
  result, carry_per_bit = A + Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = 0
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### ADD n: Add (immediate)

Adds to the 8-bit A register, the immediate data n, and stores the result back into the A register.

Opcode 0b11000110/0xC6 Length 2 bytes: opcode + n **Duration** 2 machine cycles

Flags  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
# M2
if IR == 0xC6:
    Z = read_memory(addr=PC); PC = PC + 1
    # M3/M1
    result, carry_per_bit = A + Z
    A = result
    flags.Z = 1 if result == 0 else 0
    flags.N = 0
    flags.H = 1 if carry_per_bit[3] else 0
    flags.C = 1 if carry_per_bit[7] else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### ADC r: Add with carry (register)

Adds to the 8-bit A register, the carry flag and the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10001xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle Flags  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — PC + 1

if opcode == 0x88: # example: ADC B

result, carry_per_bit = A + B + flags.C

A = result

flags.Z = 1 if result == 0 else 0

flags.N = 0

flags.H = 1 if carry_per_bit[3] else 0

flags.C = 1 if carry_per_bit[7] else 0
```

#### ADC (HL): Add with carry (indirect HL)

Adds to the 8-bit A register, the carry flag and data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10001110/0x8E Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 )( M2 )( M3/M1
Addr bus Previous
                       HL
Data bus — IR ← mem Z ← mem / IR ← mem
 IDU op — Previous PC ← PC+1
 ALU op - Previous A \leftarrow A +_c Z
 Misc op — Previous
# M2
if IR == 0x8E:
 Z = read_memory(addr=HL)
 # M3/M1
 result, carry_per_bit = A + Z + flags.C
 A = result
 flags.Z = 1 if result == 0 else 0
 flags.N = 0
 flags.H = 1 if carry_per_bit[3] else 0
 flags.C = 1 if carry_per_bit[7] else 0
 IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **ADC n: Add with carry (immediate)**

Adds to the 8-bit A register, the carry flag and the immediate data n, and stores the result back into the A register.

Opcode 0b11001110/0xCE Length 2 bytes: opcode + n **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — ( M1 ) M2 ) M3/M1
Addr bus Previous
                        PC
Data bus — IR ← mem Z ← mem IR ← mem
 IDU op — Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op - Previous A \leftarrow A +_c Z
 Misc op — Previous
# M2
if IR == 0xCE:
  Z = read_memory(addr=PC); PC = PC + 1
  # M3/M1
  result, carry_per_bit = A + Z + flags.C
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = 0
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **SUB r: Subtract (register)**

Subtracts from the 8-bit A register, the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10010xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle Flags  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — PC + 1

if opcode == 0x90: # example: SUB B

result, carry_per_bit = A - B

A = result

flags.Z = 1 if result == 0 else 0

flags.N = 1

flags.H = 1 if carry_per_bit[3] else 0

flags.C = 1 if carry_per_bit[7] else 0
```

#### **SUB (HL): Subtract (indirect HL)**

Subtracts from the 8-bit A register, data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10010110/0x96 Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                      HL
Data bus — IR←mem Z←mem IR←mem
 IDU op — Previous (PC ← PC + 1)
 ALU op ← Previous ( A ← A - Z
 Misc op — Previous
# M2
if IR == 0x96:
 Z = read_memory(addr=HL)
 # M3/M1
 result, carry_per_bit = A - Z
 A = result
 flags.Z = 1 if result == 0 else 0
 flags.N = 1
 flags.H = 1 if carry_per_bit[3] else 0
 flags.C = 1 if carry_per_bit[7] else 0
 IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **SUB n: Subtract (immediate)**

Subtracts from the 8-bit A register, the immediate data n, and stores the result back into the A register.

Opcode 0b11010110/0xD6 Length 2 bytes: opcode + n **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                       PC
Data bus — IR←mem Z←mem IR←mem
 IDU op — Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op ← Previous ( A ← A - Z
 Misc op — Previous
# M2
if IR == 0xD6:
  Z = read_memory(addr=PC); PC = PC + 1
  # M3/M1
  result, carry_per_bit = A - Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = 1
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# **SBC** r: Subtract with carry (register)

Subtracts from the 8-bit A register, the carry flag and the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10011xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — PC + 1

if opcode == 0x98: # example: SBC B

result, carry_per_bit = A - B - flags.C

A = result

flags.Z = 1 if result == 0 else 0

flags.N = 1

flags.H = 1 if carry_per_bit[3] else 0

flags.C = 1 if carry_per_bit[7] else 0
```

#### SBC (HL): Subtract with carry (indirect HL)

Subtracts from the 8-bit A register, the carry flag and data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10011110/0x9E Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                       HL
Data bus — IR ← mem Z ← mem IR ← mem
 IDU op — Previous (PC ← PC + 1)
 ALU op \overline{} Previous \overline{} A \leftarrow A - cZ
 Misc op — Previous
# M2
if IR == 0x9E:
 Z = read_memory(addr=HL)
 # M3/M1
  result, carry_per_bit = A - Z - flags.C
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = 1
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **SBC** n: Subtract with carry (immediate)

Subtracts from the 8-bit A register, the carry flag and the immediate data n, and stores the result back into the A register.

Opcode 0b11011110/0xDE Length 2 bytes: opcode + n **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                        PC
Data bus — IR ← mem Z ← mem / IR ← mem
 IDU op — Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op - Previous A \leftarrow A - c Z
 Misc op — Previous
# M2
if IR == 0xDE:
  Z = read_memory(addr=PC); PC = PC + 1
  # M3/M1
  result, carry_per_bit = A - Z - flags.C
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = 1
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# **CP r: Compare (register)**

Subtracts from the 8-bit A register, the 8-bit register  $\mathbf{r}$ , and updates flags based on the result. This instruction is basically identical to SUB  $\mathbf{r}$ , but does not update the A register.

**Opcode** 0b10111xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — read_memory(addr=PC); PC = PC + 1

if opcode == 0xB8: # example: CP B

result, carry_per_bit = A - B

flags.Z = 1 if result == 0 else 0

flags.N = 1

flags.H = 1 if carry_per_bit[3] else 0

flags.C = 1 if carry_per_bit[7] else 0
```

```
M-cycle — M1
                      M2/M1
Addr bus - Previous
Data bus IR ← mem IR ← mem
 IDU op Previous PC ← PC + 1
 ALU op Previous
                      A-r
 Misc op - Previous
# M2/M1
if IR == 0xB8: # example: CP B
 result, carry_per_bit = A - B
  flags.Z = 1 if result == 0 else 0
  flags.N = 1
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **CP (HL): Compare (indirect HL)**

Subtracts from the 8-bit A register, data from the absolute address specified by the 16-bit register HL, and updates flags based on the result. This instruction is basically identical to SUB (HL), but does not update the A register.

Opcode 0b101111110/0xBE Length 1 byte: opcode **Duration** 2 machine cycles Flags  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

#### Simple timing and pseudocode

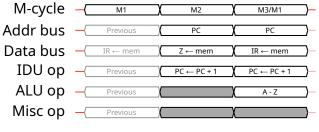
```
M-cycle — M1 X
                        M2
Addr bus — Previous
                       HL
                                 PC
Data bus — IR ← mem
                    Z ← mem IR ← mem
 IDU op - Previous
                            PC ← PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xBE:
 Z = read\_memory(addr=HL)
  # M3/M1
  result, carry_per_bit = A - Z
  flags.Z = 1 if result == 0 else 0
  flags.N = 1
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **CP n: Compare (immediate)**

Subtracts from the 8-bit A register, the immediate data n, and updates flags based on the result. This instruction is basically identical to SUB n, but does not update the A register.

Opcode 0b111111110/0xFE Length 2 bytes: opcode + n **Duration** 2 machine cycles Flags  $Z = \bigstar$ , N = 1,  $H = \bigstar$ ,  $C = \bigstar$ 

# Simple timing and pseudocode



```
# M2
if IR == 0xFE:
    Z = read_memory(addr=PC); PC = PC + 1
# M3/M1
    result, carry_per_bit = A - Z
    flags.Z = 1 if result == 0 else 0
    flags.N = 1
    flags.H = 1 if carry_per_bit[3] else 0
    flags.C = 1 if carry_per_bit[7] else 0
```

# **INC r: Increment (register)**

Increments data in the 8-bit register r.

**Opcode** 0b00xxx100/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ 

# Simple timing and pseudocode

# **Detailed timing and pseudocode**

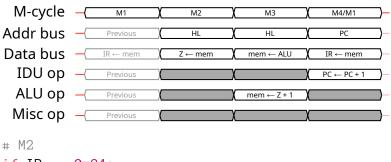
IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

# **INC (HL): Increment (indirect HL)**

Increments data at the absolute address specified by the 16-bit register HL.

Opcode 0b00110100/0x34 Length 1 byte: opcode **Duration** 3 machine cycles **Flags**  $Z = \bigstar$ , N = 0,  $H = \bigstar$ 

# Simple timing and pseudocode



```
# M2
if IR == 0x34:
    Z = read_memory(addr=HL)
    # M3
    result, carry_per_bit = Z + 1
    write_memory(addr=HL, data=result)
    flags.Z = 1 if result == 0 else 0
    flags.N = 0
    flags.H = 1 if carry_per_bit[3] else 0
    # M4/M1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# **DEC r: Decrement (register)**

Decrements data in the 8-bit register r.

**Opcode** 0b00xxx101/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ 

# Simple timing and pseudocode

#### **Detailed timing and pseudocode**

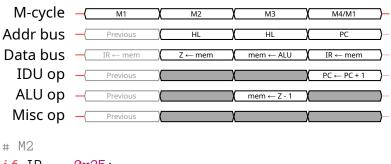
IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

# **DEC (HL): Decrement (indirect HL)**

Decrements data at the absolute address specified by the 16-bit register HL.

Opcode 0b00110101/0x35 Length 1 byte: opcode **Duration** 3 machine cycles **Flags**  $Z = \bigstar$ , N = 1,  $H = \bigstar$ 

#### Simple timing and pseudocode



```
# M2
if IR == 0x35:
    Z = read_memory(addr=HL)
    # M3
    result, carry_per_bit = Z - 1
    write_memory(addr=HL, data=result)
    flags.Z = 1 if result == 0 else 0
    flags.N = 1
    flags.H = 1 if carry_per_bit[3] else 0
    # M4/M1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### AND r: Bitwise AND (register)

Performs a bitwise AND operation between the 8-bit A register and the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10100xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle Flags  $Z = \bigstar$ , N = 0, H = 1, C = 0

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xA0: # example: AND B
result = A & B
A = result
flags.Z = 1 if result == 0 else 0
flags.N = 0
flags.H = 1
flags.C = 0
```

#### AND (HL): Bitwise AND (indirect HL)

Performs a bitwise AND operation between the 8-bit A register and data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10100110/0xA6 Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 1, C = 0

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                        HL
Data bus — R←mem Z←mem IR←mem
 IDU op \longrightarrow Previous PC \leftarrow PC +1
 ALU op \neg Previous \land A \leftarrow A and Z
 Misc op — Previous
# M2
if IR == 0xA6:
 Z = read_memory(addr=HL)
 # M3/M1
 result = A & Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = ∅
  flags.H = 1
  flags.C = ∅
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **AND n: Bitwise AND (immediate)**

Performs a bitwise AND operation between the 8-bit A register and immediate data n, and stores the result back into the A register.

Opcode 0b11100110/0xE6 Duration 2 machine cycles Length 2 bytes: opcode + n Flags  $Z = \frac{1}{100}$ , N = 0, H = 1, C = 0 Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                        PC
Data bus — R←mem Z←mem IR←mem
 IDU op — Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op \neg Previous \land A \leftarrow A and Z
 Misc op — Previous
# M2
if IR == 0xE6:
  Z = read_memory(addr=PC); PC = PC + 1
 # M3/M1
 result = A & Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = ∅
  flags.H = 1
  flags.C = ∅
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **OR r: Bitwise OR (register)**

Performs a bitwise OR operation between the 8-bit A register and the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10110xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle Flags  $Z = \bigstar$ , N = 0, H = 0, C = 0

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1 if opcode == 0xB0: # example: OR B result = A | B A = result flags.Z = 1 if result == 0 else 0 flags.N = 0 flags.C = 0
```

# **OR (HL): Bitwise OR (indirect HL)**

Performs a bitwise OR operation between the 8-bit A register and data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

Opcode 0b10110110/0xB6 Length 1 byte: opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0, C = 0

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                      HL
Data bus — R←mem Z←mem IR←mem
 IDU op - Previous PC \leftarrow PC + 1
 ALU op ← Previous A ← A or Z
 Misc op — Previous
# M2
if IR == 0xB6:
 Z = read_memory(addr=HL)
 # M3/M1
 result = A \mid Z
 A = result
 flags.Z = 1 if result == 0 else 0
 flags.N = ∅
 flags.H = ∅
 flags.C = ∅
 IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **OR n: Bitwise OR (immediate)**

Performs a bitwise OR operation between the 8-bit A register and immediate data n, and stores the result back into the A register.

Opcode 0b11110110/0xF6 Duration 2 machine cycles Length 2 bytes: opcode + n Flags  $Z = \frac{1}{100}$ , N = 0, H = 0, C = 0 Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                       PC
Data bus — R←mem Z←mem IR←mem
 IDU op - Previous PC \leftarrow PC+1 PC \leftarrow PC+1
 ALU op Previous A ← A or Z
 Misc op — Previous
# M2
if IR == 0xF6:
 Z = read_memory(addr=PC); PC = PC + 1
 # M3/M1
 result = A \mid Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = ∅
  flags.H = ∅
  flags.C = ∅
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **XOR r: Bitwise XOR (register)**

Performs a bitwise XOR operation between the 8-bit A register and the 8-bit register r, and stores the result back into the A register.

**Opcode** 0b10101xxx/various **Length** 1 byte: opcode

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , N = 0, H = 0, C = 0

# Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xA8: # example: XOR B
result = A ^ B
A = result
flags.Z = 1 if result == 0 else 0
flags.N = 0
flags.H = 0
flags.C = 0
```

#### **XOR (HL): Bitwise XOR (indirect HL)**

Performs a bitwise XOR operation between the 8-bit A register and data from the absolute address specified by the 16-bit register HL, and stores the result back into the A register.

**Opcode** 0b10101110/0xAE **Length** 1 byte: opcode

**Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0, C = 0

# Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                      HL
Data bus — R←mem Z←mem IR←mem
 IDU op - Previous PC \leftarrow PC + 1
 ALU op — Previous A ← A xor Z
 Misc op — Previous
# M2
if IR == 0xAE:
 Z = read_memory(addr=HL)
 # M3/M1
 result = A ^ Z
 A = result
 flags.Z = 1 if result == 0 else 0
 flags.N = ∅
 flags.H = ∅
 flags.C = ∅
 IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **XOR n: Bitwise XOR (immediate)**

Performs a bitwise XOR operation between the 8-bit A register and immediate data n, and stores the result back into the A register.

Opcode0b11101110/0xEEDuration2 machine cyclesLength2 bytes: opcode + nFlags $Z = \bigstar$ , N = 0, H = 0, C = 0Simple timing and pseudocode

```
M-cycle ( M1 ) M2 ) M3/M1
Addr bus Previous
                           PC
Data bus \longrightarrow IR \leftarrow mem \nearrow Z \leftarrow mem \nearrow IR \leftarrow mem
 IDU op Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op — Previous A ← A xor Z
 Misc op — Previous
# M2
if IR == 0xEE:
  Z = read_memory(addr=PC); PC = PC + 1
  # M3/M1
  result = A ^ Z
  A = result
  flags.Z = 1 if result == 0 else 0
  flags.N = ∅
  flags.H = 0
  flags.C = ∅
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# **CCF: Complement carry flag**

Flips the carry flag, and clears the N and H flags.

```
Opcode0b001111111/0x3FDuration1 machine cycleLength1 byte: opcodeFlagsN = 0, H = 0, C = \bigstar
```

### Simple timing and pseudocode

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1 if opcode == 0x3F:
flags.N = 0
flags.H = 0
flags.C = ~flags.C
```

## **Detailed timing and pseudocode**

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

### **SCF: Set carry flag**

Sets the carry flag, and clears the N and H flags.

```
Opcode 0b00110111/0x37 Duration 1 machine cycle 
Length 1 byte: opcode Flags N = 0, H = 0, C = 1 Simple timing and pseudocode
```

```
M-cycle — M1 — Mem R/W — Opcode — PC + 1

if opcode == 0x37:

flags.N = 0

flags.H = 0

flags.C = 1
```

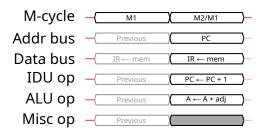
# **DAA: Decimal adjust accumulator**

TODO

Opcode 0b00100111/0x27 Length 1 byte: opcode

Simple timing and pseudocode

# **Detailed timing and pseudocode**



TODO

**Duration** 1 machine cycle **Flags**  $Z = \bigstar$ , H = 0,  $C = \bigstar$ 

# **CPL: Complement accumulator**

Flips all the bits in the 8-bit A register, and sets the N and H flags.

```
Opcode 0b001011111/0x2F
Length 1 byte: opcode
Simple timing and pseudocode

Duration 1 machine cycle
Flags N = 1, H = 1
```

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1 if opcode == 0x2F:

A = ~A
flags.N = 1
flags.H = 1
```

### 6.5 16-bit arithmetic instructions

## **INC rr: Increment 16-bit register**

Increments data in the 16-bit register rr.

Opcode 0b00xx0011/various Length 1 byte: opcode Simple timing and pseudocode **Duration** 2 machine cycles **Flags** -

# **DEC rr: Decrement 16-bit register**

Decrements data in the 16-bit register rr.

```
Opcode 0b00xx1011/various
Length 1 byte: opcode
Simple timing and pseudocode
```

**Duration** 2 machine cycles **Flags** -

#### ADD HL, rr: Add (16-bit register)

Adds to the 16-bit HL register pair, the 16-bit register rr, and stores the result back into the HL register pair.

N = 0, H = 4, C = 4

Opcode 0b00xx1001/various **Duration** 2 machine cycles **Length** 1 byte: opcode Flags

# Simple timing and pseudocode

```
M-cycle — M1 X M2
Mem R/W — opcode )
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x09: # example: ADD HL, BC
 result, carry_per_bit = HL + BC
 HL = result
 flags.N = 0
 flags.H = 1 if carry_per_bit[11] else 0
 flags.C = 1 if carry_per_bit[15] else 0
```

```
M-cycle — M1
                        M2
Addr bus — Previous
                       0x0000
Data bus 

IR ← mem

IR ← mem
 IDU op - Previous PC \leftarrow PC + 1
 ALU op Previous L \leftarrow L + Isb rr H \leftarrow H +_c msb rr
 Misc op - Previous
# M2
if IR == 0x09: # example: ADD HL, BC
 result, carry_per_bit = L + C
 L = result
  flags.N = ∅
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  result, carry_per_bit = H + B + flags.C
  H = result
  flags.N = ∅
  flags.H = 1 if carry_per_bit[3] else 0
  flags.C = 1 if carry_per_bit[7] else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### ADD SP, e: Add to stack pointer (relative)

Loads to the 16-bit SP register, 16-bit data calculated by adding the signed 8-bit operand e to the 16-bit value of the SP register.

Opcode 0b11101000/0xE8Length 2 bytes: opcode + e

Simple timing and pseudocode

Duration 4 machine cycles

Flags Z = 0, N = 0,  $H = \bigstar$ ,  $C = \bigstar$ 

# 

flags.H = 1 if carry\_per\_bit[3] else 0
flags.C = 1 if carry\_per\_bit[7] else 0

```
M-cycle –( M1 X
                         M2
                                    М3
                                                        M5/M1
Addr bus Previous
                          PC
Data bus — IR ← mem
                                    ALU
                     Z ← mem
                                               ALU
                                                       IR \leftarrow mem
 IDU op — Previous
                     PC ← PC + 1
 ALU op Previous
                      Z \leftarrow SPL + Z
                                          W ← SPH +c adj
 Misc op Previous
# M2
if IR == 0xE8:
  Z = read_memory(addr=PC); PC = PC + 1
```

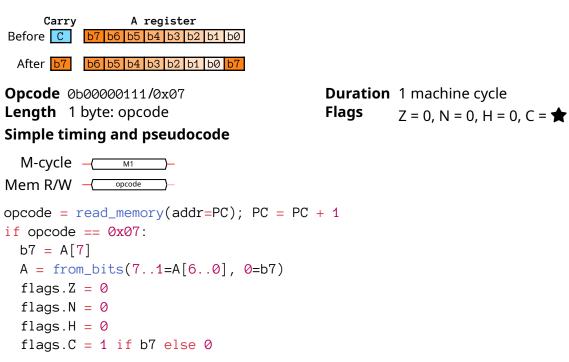
```
f IR == 0xE8:
Z = read_memory(addr=PC); PC = PC + 1
# M3
result, carry_per_bit = lsb(SP) + Z
Z = result
flags.Z = 0
flags.N = 0
flags.H = 1 if carry_per_bit[3] else 0
flags.C = 1 if carry_per_bit[7] else 0
# M4
result = msb(SP) + adj + flags.C
W = result
# M5/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1; SP = WZ
```

#### 6.6 Rotate, shift, and bit operation instructions

#### **RLCA: Rotate left circular (accumulator)**

Rotates the 8-bit A register value left in a circular manner (carry flag is updated but not used).

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). Bit 7 is copied both to bit 0 and the carry flag. Note that unlike the related RLC r instruction, RLCA always sets the zero flag to 0 without looking at the resulting value of the calculation.



```
M-cycle (M1) (M2/M1)
Addr bus - Previous
Data bus — IR ← mem IR ← mem
 IDU op - Previous PC \leftarrow PC + 1
 ALU op — Previous
                   A ← rlc A
 Misc op — Previous
# M2/M1
if IR == 0x07:
 b7 = A[7]
  result = from_bits(7..1=A[6..0], 0=b7)
  A = result
  flags.Z = 0
  flags.N = ∅
  flags.H = ∅
  flags.C = 1 if b7 else 0
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### RRCA: Rotate right circular (accumulator)

Rotates the 8-bit A register value right in a circular manner (carry flag is updated but not used).

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). Bit 0 is copied both to bit 7 and the carry flag. Note that unlike the related RRC r instruction, RRCA always sets the zero flag to 0 without looking at the resulting value of the calculation.

```
A register
                            Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C
 After b0 b7 b6 b5 b4 b3 b2 b1 b0
Opcode 0b00001111/0x0F
                                          Duration 1 machine cycle
Length 1 byte: opcode
                                          Flags
                                                    Z = 0, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle — M1
Mem R/W — opcode —
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x0F:
 b0 = A[0]
  A = from_bits(7=b0, 6..0=A[7..1])
  flags.Z = ∅
  flags.N = 0
  flags.H = 0
  flags.C = 1 if b0 else 0
```

```
M-cycle — M1
Addr bus - Previous
Data bus — IR ← mem IR ← mem
 IDU op - Previous
                   PC ← PC + 1
 ALU op Previous
                   A ← rrc A
 Misc op Previous
# M2/M1
if IR == 0x0F:
 b0 = A[0]
 result = from_bits(7=b0, 6..0=A[7..1])
 A = result
 flags.Z = 0
 flags.N = 0
 flags.H = ∅
 flags.C = 1 if b0 else 0
```

#### **RLA: Rotate left (accumulator)**

Rotates the 8-bit A register value left through the carry flag.

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). The carry flag is copied to bit 0, and bit 7 is copied to the carry flag. Note that unlike the related RL r instruction, RLA always sets the zero flag to 0 without looking at the resulting value of the calculation.

```
A register
Before C b7 b6 b5 b4 b3 b2 b1 b0
 After b7 b6 b5 b4 b3 b2 b1 b0 C
Opcode 0b00010111/0x17
                                          Duration 1 machine cycle
Length 1 byte: opcode
                                          Flags
                                                   Z = 0, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle — M1
Mem R/W — opcode —
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x17:
 b7 = A[7]
  A = from\_bits(7..1=A[6..0], 0=flags.C)
  flags.Z = ∅
  flags.N = 0
  flags.H = 0
  flags.C = 1 if b7 else 0
```

```
M-cycle — M1
Addr bus - Previous
Data bus — IR ← mem IR ← mem
 IDU op - Previous
                   PC ← PC + 1
 ALU op Previous
                   A ← rl A
 Misc op Previous
# M2/M1
if IR == 0x17:
 b7 = A[7]
 result = from_bits(7..1=A[6..0], 0=flags.C)
 A = result
 flags.Z = 0
 flags.N = 0
 flags.H = ∅
 flags.C = 1 if b7 else 0
```

#### RRA: Rotate right (accumulator)

Rotates the 8-bit A register value right through the carry flag.

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). The carry flag is copied to bit 7, and bit 0 is copied to the carry flag. Note that unlike the related RR r instruction, RRA always sets the zero flag to 0 without looking at the resulting value of the calculation.

```
A register
                            Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C
 After C b7 b6 b5 b4 b3 b2 b1 b0
Opcode 0b00011111/0x1F
                                          Duration 1 machine cycle
Length 1 byte: opcode
                                          Flags
                                                    Z = 0, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle — M1
Mem R/W — opcode —
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0x1F:
 b0 = A[0]
  A = from\_bits(7=flags.C, 6..0=A[7..1])
  flags.Z = ∅
  flags.N = 0
  flags.H = 0
  flags.C = 1 if b0 else 0
Detailed timing and pseudocode
```

```
M-cycle — M1
Addr bus - Previous
Data bus — IR ← mem IR ← mem
 IDU op - Previous
                    PC ← PC + 1
 ALU op Previous
                    A ← rr A
 Misc op — Previous
# M2/M1
if IR == 0x1F:
  b0 = A[0]
 result = from_bits(7=flags.C, 6..0=A[7..1])
  A = result
  flags.Z = 0
  flags.N = 0
  flags.H = ∅
  flags.C = 1 if b0 else 0
```

# **RLC** r: Rotate left circular (register)

Rotates the 8-bit register r value left in a circular manner (carry flag is updated but not used).

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). Bit 7 is copied both to bit 0 and the carry flag.

```
Carry
                 Register r
Before C b7 b6 b5 b4 b3 b2 b1 b0
 After b7 b6 b5 b4 b3 b2 b1 b0 b7
Opcode 0b00000xxx/various
                                            Duration 2 machine cycles
Length 2 bytes: CB prefix + opcode
                                            Flags
                                                      Z = \bigstar, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle \stackrel{\text{M1}}{\leftarrow}
                          M2
Mem R/W — CB prefix opcode
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
  opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x00: # example: RLC B
    b7 = B[7]
    B = from_bits(7..1=B[6..0], 0=b7)
    flags.Z = 1 if B == 0 else 0
    flags.N = 0
    flags.H = 0
    flags.C = 1 if b7 else 0
```

```
M-cycle – M1
Addr bus Previous
                        PC
                                  PC
Data bus — IR←mem
                    IR ← mem IR ← mem
 IDU op — Previous
                    PC ← PC + 1 PC ← PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
 cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x00: # example: RLC B
    b7 = B[7]
    result = from_bits(7..1=B[6..0], 0=b7)
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b7 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### RLC (HL): Rotate left circular (indirect HL)

Rotates, the 8-bit data at the absolute address specified by the 16-bit register HL, left in a circular manner (carry flag is updated but not used).

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). Bit 7 is copied both to bit 0 and the carry flag.

```
        Carry
        Data at address HL

        Before
        C

        b7 b6 b5 b4 b3 b2 b1 b0

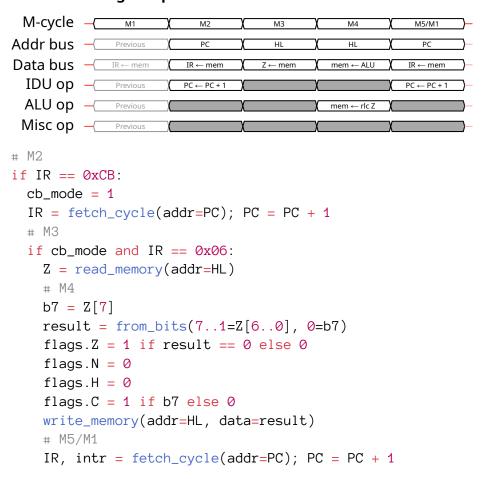
        After
        b7

        b6 b5 b4 b3 b2 b1 b0 b7
```

Opcode 0x06 Length 2 bytes: CB prefix + opcode

**Duration** 4 machine cycles C + Opcode  $C = \mathbf{flags}$   $C = \mathbf{flags}$   $C = \mathbf{flags}$   $C = \mathbf{flags}$ 

```
M-cycle — M1 X
Mem R/W → CB prefix
                                          W: data
                       opcode
                                 R: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
  opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x06:
    data = read_memory(addr=HL)
    b7 = data[7]
    result = from_bits(7..1=data[6..0], 0=b7)
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b7 else 0
    write_memory(addr=HL, data=result)
```



#### RRC r: Rotate right circular (register)

Rotates the 8-bit register r value right in a circular manner (carry flag is updated but not used).

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). Bit 0 is copied both to bit 7 and the carry flag.

```
Register r
                             Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C
 After b0 b7 b6 b5 b4 b3 b2 b1 b0
Opcode 0b00001xxx/various
                                            Duration 2 machine cycles
Length 2 bytes: CB prefix + opcode
                                            Flags
                                                      Z = \bigstar, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle \stackrel{\text{M1}}{\leftarrow}
                          M2
Mem R/W — CB prefix opcode
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
  opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x08: # example: RRC B
    b0 = B[0]
    B = from_bits(7..1=B[7..1], 0=b0)
    flags.Z = 1 if B == 0 else 0
    flags.N = 0
    flags.H = 0
    flags.C = 1 if b0 else 0
Detailed timing and pseudocode
```

```
M-cycle — M1
                                   M3/M1
Addr bus Previous
                          PC
                                    PC
Data bus — IR←mem
                     IR ← mem
                                  IR \leftarrow mem
 IDU op — Previous
                     PC ← PC + 1 PC ← PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x08: # example: RRC B
    b0 = B[0]
    result = from_bits(7..1=B[7..1], \emptyset=b\emptyset)
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b0 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# RRC (HL): Rotate right circular (indirect HL)

Rotates, the 8-bit data at the absolute address specified by the 16-bit register HL, right in a circular manner (carry flag is updated but not used).

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). Bit 0 is copied both to bit 7 and the carry flag.

```
      Data at address HL Carry

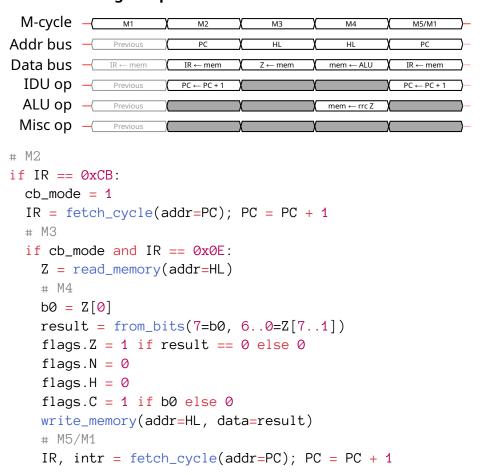
      Before
      b7 b6 b5 b4 b3 b2 b1 b0
      C

      After
      b0 b7 b6 b5 b4 b3 b2 b1
      b0
```

Opcode 0x0E Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

Flags  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

```
M-cycle — M1 X
Mem R/W → CB prefix
                       opcode
                                 R: data
                                           W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
  opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x0E:
    data = read_memory(addr=HL)
    b0 = data[0]
    result = from_bits(7=b0, 6..0=data[7..1])
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b0 else 0
    write_memory(addr=HL, data=result)
```



# RL r: Rotate left (register)

Rotates the 8-bit register r value left through the carry flag.

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). The carry flag is copied to bit 0, and bit 7 is copied to the carry flag.

```
Register r
Before C b7 b6 b5 b4 b3 b2 b1 b0
 After b7 b6 b5 b4 b3 b2 b1 b0 C
Opcode 0b00010xxx/various
                                          Duration 2 machine cycles
Length 2 bytes: CB prefix + opcode
                                          Flags
                                                    Z = \bigstar, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle –( M1 )
                        M2
Mem R/W — CB prefix opcode
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x10: # example: RL B
    b7 = B[7]
    B = from\_bits(7..1=B[6..0], 0=flags.C)
    flags.Z = 1 if B == 0 else 0
    flags.N = 0
    flags.H = 0
    flags.C = 1 if b7 else 0
```

```
M-cycle – M1
Addr bus Previous
                         PC
                                   PC
Data bus — IR←mem
                     IR ← mem IR ← mem
 IDU op — Previous
                     PC \leftarrow PC + 1 PC \leftarrow PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x10: # example: RL B
    b7 = B[7]
    result = from_bits(7..1=B[6..0], 0=flags.C)
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b7 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### RL (HL): Rotate left (indirect HL)

Rotates, the 8-bit data at the absolute address specified by the 16-bit register HL, left through the carry flag.

Every bit is shifted to the left (e.g. bit 1 value is copied from bit 0). The carry flag is copied to bit 0, and bit 7 is copied to the carry flag.

```
        Carry
        Data at address HL

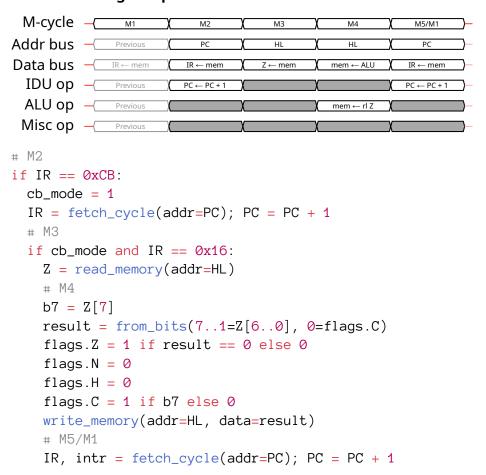
        Before C
        b7 b6 b5 b4 b3 b2 b1 b0

        After b7
        b6 b5 b4 b3 b2 b1 b0 C
```

Opcode 0x16 Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

Flags  $Z = ^{+}$ , N = 0, H = 0, C =  $^{+}$ 

```
M-cycle — M1 X M2
                                 M3
Mem R/W — CB prefix
                       opcode
                                 R: data
                                          W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x16:
    data = read_memory(addr=HL)
   b7 = data[7]
    result = from_bits(7..1=data[6..0], 0=flags.C)
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b7 else 0
    write_memory(addr=HL, data=result)
```



## RR r: Rotate right (register)

Rotates the 8-bit register r value right through the carry flag.

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). The carry flag is copied to bit 7, and bit 0 is copied to the carry flag.

```
Register r
                             Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C
 After C b7 b6 b5 b4 b3 b2 b1 b0
Opcode 0b00011xxx/various
                                            Duration 2 machine cycles
Length 2 bytes: CB prefix + opcode
                                            Flags
                                                      Z = \bigstar, N = 0, H = 0, C = \bigstar
Simple timing and pseudocode
  M-cycle \stackrel{\text{M1}}{\leftarrow}
                          M2
Mem R/W — CB prefix opcode
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
  opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x18: # example: RR B
    b0 = B[0]
    B = from\_bits(7=flags.C, 6..0=B[7..1])
    flags.Z = 1 if B == 0 else 0
    flags.N = 0
    flags.H = 0
    flags.C = 1 if b0 else 0
```

```
M-cycle – M1
Addr bus Previous
                         PC
                                   PC
Data bus — IR←mem
                     IR ← mem IR ← mem
 IDU op — Previous
                     PC \leftarrow PC + 1 PC \leftarrow PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x18: # example: RR B
    b0 = B[0]
    result = from_bits(7=flags.C, 6..0=B[7..1])
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b0 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

### RR (HL): Rotate right (indirect HL)

Rotates, the 8-bit data at the absolute address specified by the 16-bit register HL, right through the carry flag.

Every bit is shifted to the right (e.g. bit 1 value is copied to bit 0). The carry flag is copied to bit 7, and bit 0 is copied to the carry flag.

```
      Data at address HL Carry

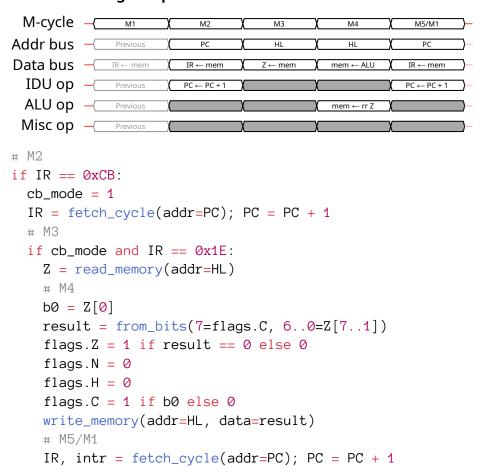
      Before
      b7 b6 b5 b4 b3 b2 b1 b0
      C

      After
      C b7 b6 b5 b4 b3 b2 b1
      b0
```

Opcode 0x1E Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

Flags  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

```
M-cvcle — M1 X M2
Mem R/W — CB prefix
                                          W: data
                       opcode
                                 R: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x1E:
    data = read_memory(addr=HL)
   b0 = data[0]
    result = from_bits(7=flags.C, 6..0=data[7..1])
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 1 if b0 else 0
    write_memory(addr=HL, data=result)
```



#### **SLA** r: Shift left arithmetic (register)

Shifts the 8-bit register r value left by one bit using an arithmetic shift.

Bit 7 is shifted to the carry flag, and bit 0 is set to a fixed value of 0.

```
        Carry
        Register r

        Before C
        b7 b6 b5 b4 b3 b2 b1 b0

After b7 b6 b5 b4 b3 b2 b1 b0 0
```

**Opcode** 0b00100xxx/various **Length** 2 bytes: CB prefix + opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle –( M1 ) M2
                            M3/M1
Addr bus — Previous
                       PC
Data bus — IR ← mem / IR ← mem / IR ← mem
 IDU op - Previous
                   PC ← PC + 1 PC ← PC + 1
 ALU op - Previous
 Misc op Previous
# M2
if IR == 0xCB:
 cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x20: # example: SLA B
   b7 = B[7]
    result = from_bits(7..1=B[6..0], \emptyset=0)
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b7 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### SLA (HL): Shift left arithmetic (indirect HL)

Shifts, the 8-bit value at the address specified by the HL register, left by one bit using an arithmetic shift.

Bit 7 is shifted to the carry flag, and bit 0 is set to a fixed value of 0.

```
        Carry
        Data at address HL

        Before
        C

        b7 b6 b5 b4 b3 b2 b1 b0

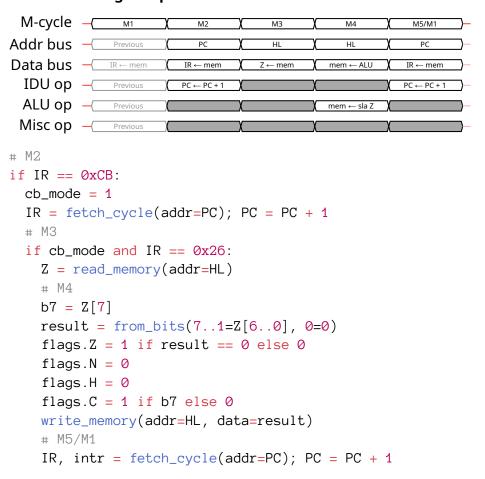
After b7

b6 b5 b4 b3 b2 b1 b0 0
```

Opcode 0x26 Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

Flags  $Z = \bigstar$ , N = 0, H = 0, C =  $\bigstar$ 

```
M-cycle (M1)
                       M2
                                 М3
                                          M4
Mem R/W — CB prefix Opcode R: data
                                         W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x26:
    data = read_memory(addr=HL)
   b7 = data[7]
   result = from_bits(7..1=data[6..0], 0=0)
    data = result
    flags.Z = 1 if B == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b7 else 0
```



#### SRA r: Shift right arithmetic (register)

Shifts the 8-bit register r value right by one bit using an arithmetic shift.

Bit 7 retains its value, and bit 0 is shifted to the carry flag.

```
      Register r
      Carry

      Before
      b7 b6 b5 b4 b3 b2 b1 b0
      C

      After
      b7 b7 b6 b5 b4 b3 b2 b1
      b0
```

**Opcode** 0b00101xxx/various **Length** 2 bytes: CB prefix + opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 X M2
                           M3/M1
Addr bus — Previous
                  PC
Data bus — IR ← mem / IR ← mem / IR ← mem
 IDU op - Previous
                  PC ← PC + 1 PC ← PC + 1
 ALU op Previous
 Misc op Previous
# M2
if IR == 0xCB:
 cb_{mode} = 1
 IR = fetch_cycle(addr=PC); PC = PC + 1
 # M3/M1
 if cb_mode and IR == 0x28: # example: SRA B
   b0 = B[0]
   result = from_bits(7=B[7], 6..0=B[7..1])
   B = result
   flags.Z = 1 if result == 0 else 0
   flags.N = ∅
   flags.H = ∅
   flags.C = 1 if b0 else 0
   IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# SRA (HL): Shift right arithmetic (indirect HL)

Shifts, the 8-bit value at the address specified by the HL register, right by one bit using an arithmetic shift.

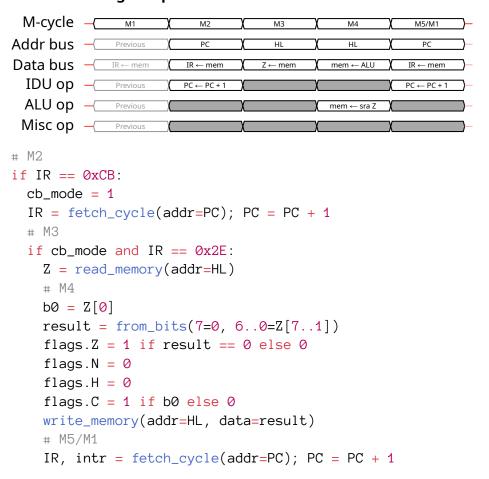
Bit 7 is set to a fixed value of 0, and bit 0 is shifted to the carry flag.

```
Data at address HL Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C

After 0 b7 b6 b5 b4 b3 b2 b1 b0
```

Opcode 0x2E Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

```
M-cycle (M1) (M2)
                                 М3
                                          M4
Mem R/W — CB prefix Opcode R: data
                                         W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x2E:
    data = read_memory(addr=HL)
   b0 = data[0]
   result = from_bits(7=0, 6..0=data[7..1])
    data = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b0 else 0
```



#### SWAP r: Swap nibbles (register)

Swaps the high and low 4-bit nibbles of the 8-bit register r.

```
Register r

Before b7 b6 b5 b4 b3 b2 b1 b0

After b3 b2 b1 b0 b7 b6 b5 b4
```

**Opcode** 0b00110xxx/various **Length** 2 bytes: CB prefix + opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0, C = 0

# Simple timing and pseudocode

```
M-cycle - M1
Addr bus - Previous
                                   PC
                         PC
Data bus — IR←mem
                    IR ← mem IR ← mem
 IDU op - Previous PC \leftarrow PC +1 PC \leftarrow PC +1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x30: # example: SWAP B
    result = from_bits(7..4=B[3..0], 3..0=B[7..4])
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **SWAP (HL): Swap nibbles (indirect HL)**

Swaps the high and low 4-bit nibbles of the 8-bit data at the absolute address specified by the 16-bit register HL.

```
Data at address HL

Before b7 b6 b5 b4 b3 b2 b1 b0

After b3 b2 b1 b0 b7 b6 b5 b4
```

Opcode 0x36 Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

**Flags**  $Z = \frac{1}{100}$ , N = 0, H = 0, C = 0

### Simple timing and pseudocode

```
M-cycle — M1
                                     МЗ
                                                M4
                                                          M5/M1
Addr bus - Previous
                           PC
                                     HL
                                                HL
                                                           PC
Data bus — IR← mem
                                    Z ← mem
                        IR \leftarrow mem
  IDU op — Previous
                       PC ← PC + 1
 ALU op — Previous
                                             mem ← swap Z
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3
  if cb_mode and IR == 0x36:
    Z = read_memory(addr=HL)
    result = from_bits(7..4=\mathbb{Z}[3..0], 3..0=\mathbb{Z}[7..4])
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = 0
    flags.C = 0
    write_memory(addr=HL, data=result)
    # M5/M1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### SRL r: Shift right logical (register)

Shifts the 8-bit register r value right by one bit using a logical shift.

Bit 7 is set to a fixed value of 0, and bit 0 is shifted to the carry flag.

```
      Register r
      Carry

      Before
      b7
      b6
      b5
      b4
      b3
      b2
      b1
      b0
      C

After

      0
      b7
      b6
      b5
      b4
      b3
      b2
      b1
      b0
```

**Opcode** 0b00111xxx/various **Length** 2 bytes: CB prefix + opcode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 0,  $C = \bigstar$ 

# Simple timing and pseudocode

```
M-cycle — M1 X
                      M2
                            M3/M1
Addr bus — Previous
                       PC
Data bus — IR ← mem / IR ← mem / IR ← mem
 IDU op — Previous
                   PC ← PC + 1 PC ← PC + 1
 ALU op - Previous
 Misc op Previous
# M2
if IR == 0xCB:
 cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x28: # example: SRL B
   b0 = B[0]
   result = from_bits(7=0, 6..0=B[7..1])
    B = result
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b0 else 0
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

# SRL (HL): Shift right logical (indirect HL)

Shifts, the 8-bit value at the address specified by the HL register, right by one bit using a logical shift.

Bit 7 is set to a fixed value of 0, and bit 0 is shifted to the carry flag.

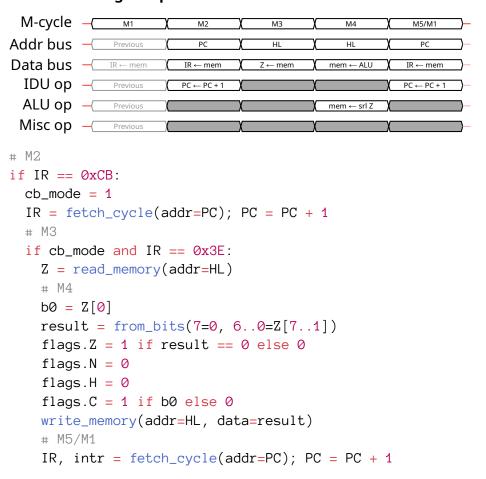
```
Data at address HL Carry
Before b7 b6 b5 b4 b3 b2 b1 b0 C

After 0 b7 b6 b5 b4 b3 b2 b1 b0
```

Opcode 0x3E Length 2 bytes: CB prefix + opcode **Duration** 4 machine cycles

**Flags** Z = - N = 0, H = 0, C = -

```
M-cycle (M1) (M2)
                                 М3
                                          M4
Mem R/W — CB prefix opcode X
                                R: data
                                         W: data
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xCB: # CB prefix
 opcode = read_memory(addr=PC); PC = PC + 1
  if opcode == 0x3E:
    data = read_memory(addr=HL)
    b0 = data[0]
   result = from_bits(7=0, 6..0=data[7..1])
    flags.Z = 1 if result == 0 else 0
    flags.N = ∅
    flags.H = ∅
    flags.C = 1 if b0 else 0
    write_memory(addr=HL, data=result)
```



#### BIT b, r: Test bit (register)

Tests the bit b of the 8-bit register r.

The zero flag is set to 1 if the chosen bit is 0, and 0 otherwise.

Opcode 0b01xxxxxx/various Length 2 bytes: CB prefix + opcode Simple timing and pseudocode **Duration** 2 machine cycles **Flags**  $Z = \bigstar$ , N = 0, H = 1

#### **Detailed timing and pseudocode**

```
M-cycle — M1 X
                         M2
Addr bus — Previous
                          PC
Data bus — IR←mem
                     IR ← mem IR ← mem
 IDU op — Previous
                     PC ← PC + 1 PC ← PC + 1
 ALU op - Previous
 Misc op - Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3/M1
  if cb_mode and IR == 0x40: # example: BIT 0, B
    flags.Z = 1 if B[\emptyset] == \emptyset else \emptyset
    flags.N = ∅
    flags.H = 1
```

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

#### BIT b, (HL): Test bit (indirect HL)

Tests the bit b of the 8-bit data at the absolute address specified by the 16-bit register HL.

The zero flag is set to 1 if the chosen bit is 0, and 0 otherwise.

**Opcode** 0b01xxx110/various **Duration** 3 machine cycles **Length** 2 bytes: CB prefix + opcode **Flags**  $Z = \bigstar$ , N = 0, H = 1 **Simple timing and pseudocode** 

```
M-cycle — M1
                                   М3
                                           M4/M1
Addr bus Previous
                         PC
                                  HL
Data bus 

— IR ← mem
                    IR ← mem Z ← mem
 IDU op - Previous
                    PC ← PC + 1
 ALU op - Previous
 Misc op — Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  if cb_mode and IR == 0x46: # example: BIT 0, (HL)
    Z = read_memory(addr=HL)
    # M4/M1
    flags.Z = 1 if Z[0] == 0 else 0
    flags.N = ∅
    flags.H = 1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### RES b, r: Reset bit (register)

Resets the bit b of the 8-bit register r to 0.

Opcode 0b10xxxxxx/various Length 2 bytes: CB prefix + opcode Simple timing and pseudocode

```
Duration 2 machine cycles Flags -
```

#### **Detailed timing and pseudocode**

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

#### RES b, (HL): Reset bit (indirect HL)

Resets the bit b of the 8-bit data at the absolute address specified by the 16-bit register HL, to 0.

Opcode 0b10xxx110/various Duration 4 machine cycles Length 2 bytes: CB prefix + opcode Flags - Simple timing and pseudocode

#### **Detailed timing and pseudocode**

```
M-cycle — M1
                                М3
Addr bus Previous
                      PC
                               HL
                                         HL
Data bus — IR ← mem / IR ← mem / Z ← mem / mem ← ALU
 IDU op — Previous
                  ( PC ← PC + 1 )
 ALU op - Previous
                                    mem ← res b, Z
 Misc op — Previous (
# M2
if IR == 0xCB:
 cb_{mode} = 1
 IR = fetch_cycle(addr=PC); PC = PC + 1
 # M3
 if cb_mode and IR == 0x86: # example: RES 0, (HL)
   Z = read_memory(addr=HL)
```

result = from\_bits(7..1=Z[7..1], 0=0)
write\_memory(addr=HL, data=result)

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

# M5/M1

#### SET b, r: Set bit (register)

Sets the bit b of the 8-bit register r to 1.

Opcode 0b11xxxxxx/various Length 2 bytes: CB prefix + opcode Simple timing and pseudocode

```
Duration 2 machine cycles Flags -
```

#### **Detailed timing and pseudocode**

IR, intr = fetch\_cycle(addr=PC); PC = PC + 1

#### SET b, (HL): Set bit (indirect HL)

Sets the bit b of the 8-bit data at the absolute address specified by the 16-bit register HL, to 1.

Opcode Ob11xxx110/various Duration 4 machine cycles
Length 2 bytes: CB prefix + opcode Flags Simple timing and pseudocode

```
M-cycle — M1
                                  М3
Addr bus Previous
                        PC
                                  HL
                                            HL
Data bus 	─ IR ← mem
                   IR \leftarrow mem Z \leftarrow mem mem \leftarrow ALU
 IDU op — Previous
                    ( PC ← PC + 1 )
 ALU op - Previous
 Misc op Previous
# M2
if IR == 0xCB:
  cb_{mode} = 1
  IR = fetch_cycle(addr=PC); PC = PC + 1
  # M3
  if cb_mode and IR == 0xC6: # example: SET 0, (HL)
    Z = read_memory(addr=HL)
    result = from_bits(7..1=\mathbb{Z}[7..1], \emptyset=1)
    write_memory(addr=HL, data=result)
    # M5/M1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### 6.7 Control flow instructions

#### JP nn: Jump

Unconditional jump to the absolute address specified by the 16-bit immediate operand nn.

```
f IR == 0xC3:
Z = read_memory(addr=PC); PC = PC + 1
# M3
W = read_memory(addr=PC); PC = PC + 1
# M4
PC = WZ
# M5/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### JP HL: Jump to HL

Unconditional jump to the absolute address specified by the 16-bit register HL.

Opcode <code>0b11101001/0xE9</code>
Length 1 byte: opcode
Simple timing and pseudocode

Duration 1 machine cycle
Flags -

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1 if opcode == 0xE9:

PC = HL
```

#### Warning

In some documentation this instruction is written as JP <code>[HL]</code>. This is very misleading, since brackets are usually used to indicate a memory read, and this instruction simply copies the value of HL to PC.

#### JP cc, nn: Jump (conditional)

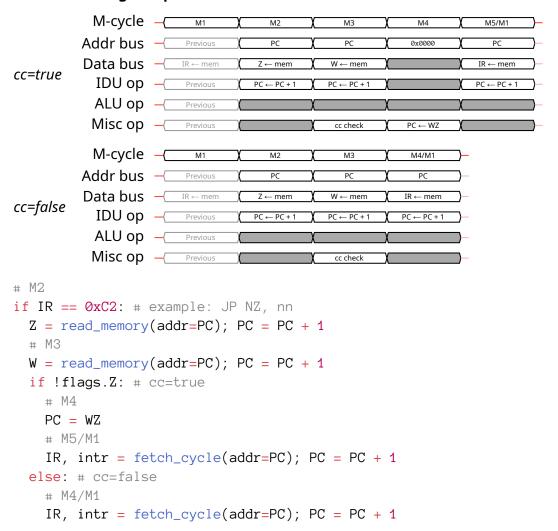
Conditional jump to the absolute address specified by the 16-bit operand nn, depending on the condition cc.

Note that the operand (absolute address) is read even when the condition is false!

Opcode 0b110xx010/variousDuration 4 machine cycles (cc=true)3 machine cycles (cc=false)

**Length** 3 bytes: opcode + LSB(nn) + MSB(nn) **Flags Simple timing and pseudocode** 

```
M-cycle – M1
                                M2
                                         М3
cc=true
        Mem R/W ─ opcode
                            R: Isb(nn) R: msb(nn)
          M-cycle — M1
                                M2
cc=false
        Mem R/W — opcode
                            R: Isb(nn)
                                     R: msb(nn)
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xC2: # example: JP NZ, nn
  nn_lsb = read_memory(addr=PC); PC = PC + 1
  nn_msb = read_memory(addr=PC); PC = PC + 1
  nn = unsigned_16(lsb=nn_lsb, msb=nn_msb)
  if !flags.Z: # cc=true
   PC = nn
```



#### JR e: Relative jump

Unconditional jump to the relative address specified by the signed 8-bit operand e.

#### **Detailed timing and pseudocode**

```
M-cycle — M1 X
                       M2
Addr bus Previous PC
                                 PCH
                                       WZ }-
Data bus — IR ← mem
                    Z ← mem
                                  ALU
 IDU op - Previous PC \leftarrow PC + 1 W \leftarrow adj PCH PC \leftarrow WZ + 1
 ALU op - Previous
                    Z \leftarrow PCL + Z
 Misc op - Previous
# M2
if IR == 0x18:
 Z = read_memory(addr=PC); PC = PC + 1
  # M3
  Z_{sign} = bit(7, Z)
  result, carry_per_bit = Z + lsb(PC)
  Z = result
  adj = 1 if carry_per_bit[7] and not Z_sign else
       -1 if not carry_per_bit[7] and Z_sign else
        0
 W = msb(PC) + adj
  # M4/M1
```

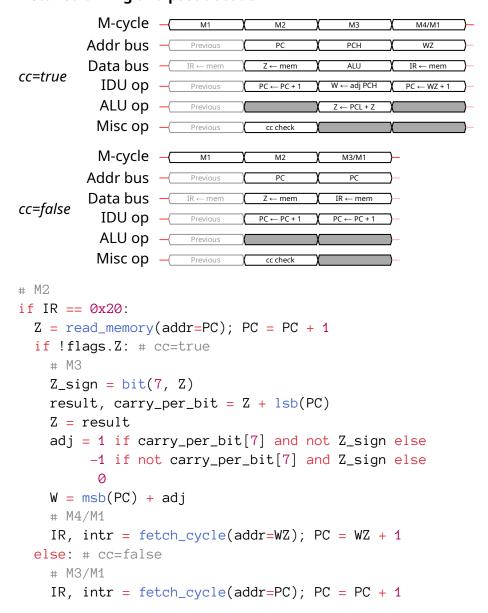
IR, intr = fetch\_cycle(addr=WZ); PC = WZ + 1

#### JR cc, e: Relative jump (conditional)

Conditional jump to the relative address specified by the signed 8-bit operand e, depending on the condition cc.

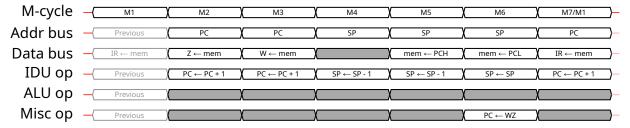
Note that the operand (relative address offset) is read even when the condition is false!

**Opcode** 0b001xx000/various **Duration** 3 machine cycles (cc=true) 2 machine cycles (cc=false) **Length** 2 bytes: opcode + e Flags Simple timing and pseudocode M-cycle — M1 cc=true Mem R/W — opcode X R: e M-cycle — M1 M2 cc=false Mem R/W — Opcode R:e opcode = read\_memory(addr=PC); PC = PC + 1 if opcode == 0x20: e = signed\_8(read\_memory(addr=PC)); PC = PC + 1 if !flags.Z: # cc=true PC = PC + e



#### **CALL nn: Call function**

Unconditional function call to the absolute address specified by the 16-bit operand nn.



```
# M2
if IR == 0xCD:
    Z = read_memory(addr=PC); PC = PC + 1
# M3
W = read_memory(addr=PC); PC = PC + 1
# M4
SP = SP - 1
# M5
write_memory(addr=SP, data=msb(PC)); SP = SP - 1
# M6
write_memory(addr=SP, data=lsb(PC)); PC = WZ
# M7/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **CALL cc, nn: Call function (conditional)**

Conditional function call to the absolute address specified by the 16-bit operand nn, depending on the condition cc.

Note that the operand (absolute address) is read even when the condition is false!

**Opcode** 0b110xx100/various **Duration** 6 machine cycles (cc=true)

3 machine cycles (cc=false)

**Length** 3 bytes: opcode + LSB(nn) + MSB(nn) **Flags Simple timing and pseudocode** 

```
M-cycle — M1
                                 M2
                                           М3
cc=true
         Mem R/W — opcode
                             R: Isb(nn) R: msb(nn)
                                                           W: msb(PC<sub>0</sub>+3) W: lsb(PC<sub>0</sub>+3)
           M-cvcle — M1
                                  M2
cc=false
         Mem R/W ─ opcode
                            R: Isb(nn) R: msb(nn)
opcode = read_memory(addr=PC); PC = PC + 1
if opcode == 0xC4: # example: CALL NZ, nn
  nn_lsb = read_memory(addr=PC); PC = PC + 1
  nn_msb = read_memory(addr=PC); PC = PC + 1
  nn = unsigned_16(lsb=nn_lsb, msb=nn_msb)
  if !flags.Z: # cc=true
    SP = SP - 1
    write_memory(addr=SP, data=msb(PC)); SP = SP - 1
    write_memory(addr=SP, data=lsb(PC))
    PC = nn
```

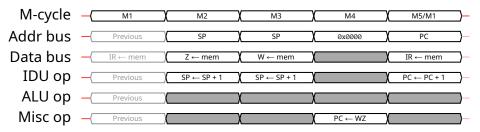


#### **RET: Return from function**

Unconditional return from a function.

Opcode 0b11001001/0xC9
Length 1 byte: opcode
Simple timing and pseudocode

**Duration** 4 machine cycles **Flags** -



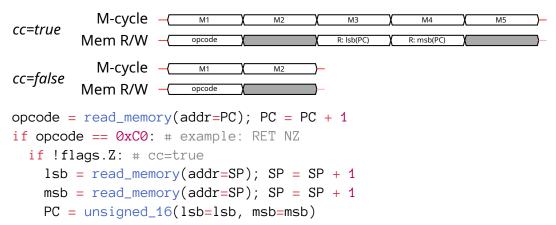
```
# M2
if IR == 0xC9:
    Z = read_memory(addr=SP); SP = SP + 1
    # M3
    W = read_memory(addr=SP); SP = SP + 1
    # M4
    PC = WZ
    # M5/M1
    IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

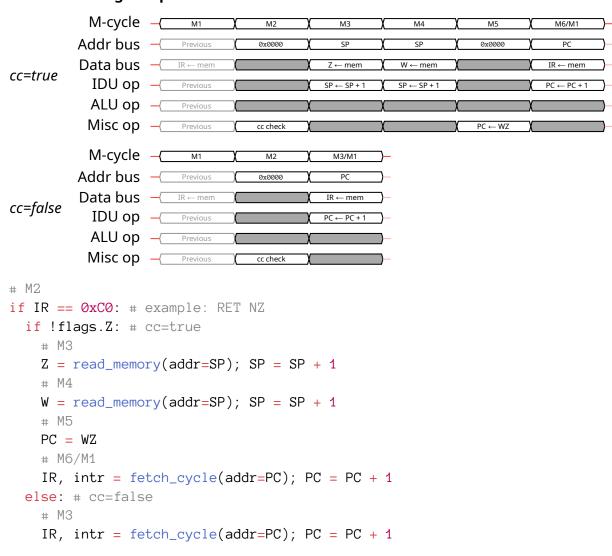
#### **RET cc: Return from function (conditional)**

Conditional return from a function, depending on the condition cc.

OpcodeØb110xx000/variousDuration5 machine cycles (cc=true)2 machine cycles (cc=false)Length1 byte: opcodeFlags-

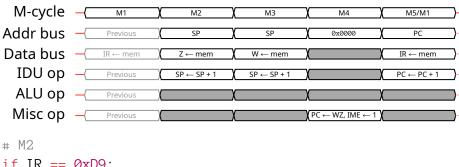
Simple timing and pseudocode





#### **RETI: Return from interrupt handler**

Unconditional return from a function. Also enables interrupts by setting IME=1.

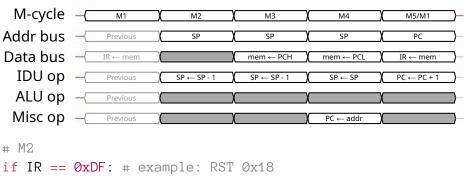


```
if IR == 0xD9:
    Z = read_memory(addr=SP); SP = SP + 1
# M3
W = read_memory(addr=SP); SP = SP + 1
# M4
PC = WZ; IME = 1
# M5/M1
IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### RST n: Restart / Call function (implied)

Unconditional function call to the absolute fixed address defined by the opcode.

Opcode Ob11xxx111/various Duration 4 machine cycles
Length 1 byte: opcode Flags Simple timing and pseudocode



```
if IR == 0xDF: # example: RST 0x18
  SP = SP - 1
# M3
  write_memory(addr=SP, data=msb(PC)); SP = SP - 1
# M4
  write_memory(addr=SP, data=lsb(PC)); PC = 0x0018
# M5/M1
  IR, intr = fetch_cycle(addr=PC); PC = PC + 1
```

#### **6.8 Miscellaneous instructions**

#### **HALT: Halt system clock**

**TODO** 

#### STOP: Stop system and main clocks

**TODO** 

 $IME = \emptyset$ 

#### **DI: Disable interrupts**

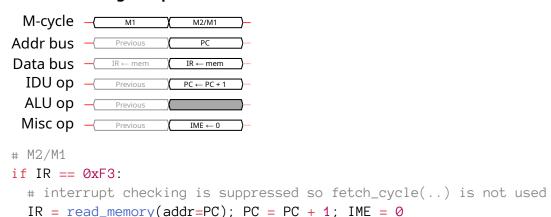
Disables interrupt handling by setting IME=0 and cancelling any scheduled effects of the EI instruction if any.

Opcode 0b11110011/0xF3
Length 1 byte: opcode
Simple timing and pseudocode

M-cycle — M1 — Mem R/W — Opcode — Opcode — PC + 1

if opcode == 0xF3:

Duration 1 machine cycle
Flags 
Plags 
Pla



#### **EI: Enable interrupts**

Schedules interrupt handling to be enabled after the next machine cycle.

Opcode <code>0b11111011/0xFB</code>
Length 1 byte: opcode
Simple timing and pseudocode

Duration 1 machine cycle
Flags -

```
M-cycle — M1 — Mem R/W — Opcode — opcode = read_memory(addr=PC); PC = PC + 1 if opcode == 0xFB:

IME_next = 1
```

```
M-cycle — (M1 ) (M2/M1) —
Addr bus — Previous (PC —
Data bus — IR — mem (IR — mem ) —
IDU op — Previous (PC — PC + 1) —
ALU op — Previous (IME — 1) —

# M2/M1

if IR == ØxFB:
IR, intr = fetch_cycle(addr=PC); PC = PC + 1; IME = 1
```

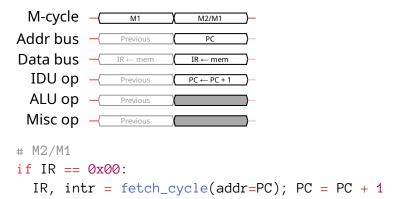
#### **NOP: No operation**

No operation. This instruction doesn't do anything, but can be used to add a delay of one machine cycle and increment PC by one.

Opcode 0b0000000/0x00 Duration 1 machine cycle
Length 1 byte: opcode Flags Simple timing and pseudocode

M-cycle — M1 — Mem R/W — Opcode — Opcode — PC + 1

if opcode == 0x00:
# nothing



# Part III Game Boy SoC peripherals and features

## Chapter 7 Boot ROM

The Game Boy SoC includes a small embedded boot ROM, which can be mapped to the 0x0000-0x00FF memory area. While mapped, all reads from this area are handled by the boot ROM instead of the external cartridge, and all writes to this area are ignored and cannot be seen by external hardware (e.g. the cartridge MBC).

The boot ROM is enabled by default, so when the system exits the reset state and the CPU starts execution from address  $0 \times 0000$ , it executes the boot ROM instead of instructions from the cartridge ROM. The boot ROM is responsible for showing the initial logo, and checking that a valid cartridge is inserted into the system. If the cartridge is valid, the boot ROM unmaps itself before execution of the cartridge ROM starts at  $0 \times 0100$ . The cartridge ROM has no chance of executing any instructions before the boot ROM is unmapped, which prevents the boot ROM from being read byte by byte in normal conditions.

#### Warning

Don't confuse the boot ROM with the additional SNES ROM in SGB/SGB2 that is executed by the SNES CPU.

Register 7.1: 0xFF50 - BOOT - Boot ROM lock register

U	U	U	U	U	U	U	R/W-0
							BOOT_OFF
bit 7	6	5	4	3	2	1	bit 0

#### bit 7-1 Unimplemented: Ignored during writes, reads are undefined

#### **bit 0 BOOT OFF**: Boot ROM lock bit

0b1 = Boot ROM is disabled and 0x0000-0x00FF works normally.

0b0 = Boot ROM is active and intercepts accesses to 0x0000-0x00FF.

BOOT\_OFF can only transition from 0b0 to 0b1, so once 0b1 has been written, the boot ROM is permanently disabled until the next system reset. Writing 0b0 when BOOT\_OFF is 0b0 has no effect and doesn't lock the boot ROM.

The 1-bit BOOT register controls mapping of the boot ROM. Once 0b1 has been written to it to unmap the boot ROM, it can only be mapped again by resetting the system.

#### 7.1 Boot ROM types

Туре	CRC32	MD5	SHA1
DMG	59c8598e	32fbbd84168d3482956eb3c5051637f5	4ed31ec6b0b175bb109c0eb5fd3d193da823339f
MGB	e6920754	71a378e71ff30b2d8a1f02bf5c7896aa	4e68f9da03c310e84c523654b9026e51f26ce7f0
SGB	ec8a83b9	d574d4f9c12f305074798f54c091a8b4	aa2f50a77dfb4823da96ba99309085a3c6278515
SGB2	53d0dd63	e0430bca9925fb9882148fd2dc2418c1	93407ea10d2f30ab96a314d8eca44fe160aea734
DMG0	c2f5cc97	a8f84a0ac44da5d3f0ee19f9cea80a8c	8bd501e31921e9601788316dbd3ce9833a97bcbc

Table 7.1: Summary of boot ROM file hashes

#### **DMG boot ROM**

The most common boot ROM is the DMG boot ROM used in almost all original Game Boy units. If a valid cartridge is inserted, the boot ROM scrolls a logo to the center of the screen, and plays a "di-ding" sound recognizable by most people who have used Game Boy consoles.

This boot ROM was originally dumped by neviksti in 2003 by decapping the Game Boy SoC and visually inspecting every single bit.

#### MGB boot ROM

This boot ROM was originally dumped by Bennvenn in 2014 by using a simple clock glitching method that only requires one wire.

#### **SGB boot ROM**

This boot ROM was originally dumped by Costis Sideris in 2009 by using an FPGA-based clock glitching method [4].

#### **SGB2 boot ROM**

This boot ROM was originally dumped by gekkio in 2015 by using a Teensy 3.1 -based clock glitching method [5].

#### Early DMG boot ROM ("DMG0")

Very early original Game Boy units released in Japan (often called "DMG0") included the launch version "DMG-CPU" SoC chip, which used a different boot ROM than later units.

This boot ROM was originally dumped by gekkio in 2016 by using a clock glitching method invented by BennVenn.

## **DMA (Direct Memory Access)**

#### 8.1 Object Attribute Memory (OAM) DMA

OAM DMA is a high-throughput mechanism for copying data to the OAM area (a.k.a. Object Attribute Memory, a.k.a. sprite memory). It can copy one byte per machine cycle without involving the CPU at all, which is much faster than the fastest possible memory routine that can be written with the SM83 instruction set. However, a transfer cannot be cancelled and the transfer length cannot be controlled, so the DMA transfer always updates the entire OAM area (= 160 bytes) even if you actually want to just update the first couple of bytes.

The Game Boy CPU chip contains a DMA controller that coordinates transfers between a **source area** and the **OAM area** independently of the CPU. While a transfer is in progress, it takes control of the source bus and the OAM area, so some precaution is needed with memory accesses (including instruction fetches) to avoid OAM DMA bus conflicts. OAM DMA uses a different address decoding scheme than normal memory accesses, so the source bus is always either the external bus or the video RAM bus, and the contents normally visible to the CPU in the <code>0xFE00-0xFFFF</code> address range cannot be used as a source for OAM DMA transfers.

The upper 8 bits of the OAM DMA source address are stored in the DMA register, while the lower 8 bits used by both the source and target address are stored in the DMA controller and are not accessible directly. A transfer always begins with 0x00 in the lower bits and copies exactly 160 bytes, so the lower bits are never in the 0xA0-0xFF range.

Writing to the DMA register updates the upper bits of the DMA source address and also triggers an OAM DMA transfer request, although the DMA transfer does not begin immediately.

Register 8.1: 0xFF46 - DMA - OAM DMA control register

R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x	R/W-x		
	DMA<7:0>								
bit 7	6	5	4	3	2	1	bit 0		

#### bit 7-0 DMA<7:0>: OAM DMA source address

Specifies the top 8 bits of the OAM DMA source address.

Writing to this register requests an OAM DMA transfer, but it's just a request and the actual DMA transfer starts with a delay.

Reading this register returns the value that was previously written to the register. The stored value is not cleared on reset, so the initial value before the first write is unknown and should not be relied on.

#### Warning

Avoid writing 0xE0-0xFF to the DMA register, because some poorly designed flash carts can trigger bus conflicts or other dangerous behaviour.

#### OAM DMA address decoding

The OAM DMA controller uses a simplified address decoding scheme, which leads to some addresses being unusable as source addresses. Unlike normal memory accesses, OAM DMA transfers interpret all accesses in the 0xA000-0xFFFF range as external RAM transfers. For example, if the OAM DMA wants to read 0xFF00, it will output 0xFF00 on the external address bus and will assert the external RAM chip select signal. The P1 register which is normally at 0xFF00 is not involved at all, because OAM DMA address decoding only uses the external bus and the video RAM bus. Instead, the resulting behaviour depends on several factors, including the connected cartridge. Some flash carts are not prepared for this unexpected scenario, and a bus conflict or worse behaviour can happen.

DMA register value	Used bus	Asserted chip select signal
0x00-0x7F	external bus	external ROM (A15)
0x80-0x9F	video RAM bus	video RAM (MCS)
0xA0-0xFF	external bus	external RAM (CS)

Table 8.1: OAM DMA address decoding scheme

#### **OAM DMA transfer timing**

**TODO** 

#### **OAM DMA bus conflicts**

**TODO** 

## **PPU (Picture Processing Unit)**

#### Register 9.1: 0xFF40 - LCDC - PPU control register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0
LCD_EN	WIN_MAP	WIN_EN	TILE_SEL	BG_MAP	OBJ_SIZE	OBJ_EN	BG_EN
bit 7	6	5	4	3	2	1	bit 0

#### Register 9.2: 0xFF41 - STAT - PPU status register

U	R/W-0	R/W-0	R/W-0	R/W-0	R-0	R-0	R-0
	INTR_LYC	INTR_M2	INTR_M1	INTR_MØ	LYC_STAT	LCD_MO[	DE <1:0>
bit 7	6	5	4	3	2	1	bit 0

#### Register 9.3: 0xFF42 - SCY - Vertical scroll register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0			
	SCY<7:0>									
bit 7 6 5 4 3 2 1 bit 0										

#### Register 9.4: 0xFF43 - SCX - Horizontal scroll register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0			
	SCX<7:0>									
bit 7 6 5 4 3 2 1 bit 0										

#### Register 9.5: 0xFF44 - LY - Scanline register

R-0	R-0	R-0	R-0	R-0	R-0	R-0	R-0			
	LY<7:0>									
bit 7	bit 7 6 5 4 3 2 1 bit 0									

#### Register 9.6: 0xFF45 - LYC - Scanline compare register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0			
	LYC<7:0>									
bit 7	6	5	4	3	2	1	bit 0			

## Port P1 (Joypad, Super Game Boy communication)

Register 10.1: 0xFF00 - P1 - Joypad/Super Game Boy communication register

U	U	W-0	W-0	R-x	R-x	R-x	R-x
		P15	P14	P13	P12	P11	P10
bit 7	6	5	4	3	2	1	bit 0

bit 7-6 Unimplemented: Ignored during writes, reads are undefined

bit 5 P15

bit 4 P14

bit 3 P13

bit 2 P12

bit 1 P11

bit 0 P10

## **Serial communication**

Register 11.1: 0xFF01 - SB - Serial data register

R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0	R/W-0			
	SB<7:0>									
bit 7 6 5 4 3 2 1 bit 0										

**bit 7-0 SB<7:0>**: Serial data

Register 11.2: 0xFF02 - SC - Serial control register

R/W-0	U	U	U	U	U	J	R/W-0
SIO_EN							SIO_CLK
bit 7	6	5	4	3	2	1	bit 0

bit 7 SIO\_EN

**bit 6-1 Unimplemented**: Ignored during writes, reads are undefined

bit 0 SIO\_CLK

# Part IV Game Boy game cartridges

## Chapter 12 MBC1 mapper chip

The majority of games for the original Game Boy use the MBC1 chip. MBC1 supports ROM sizes up to 16 Mbit (128 banks of 0x4000 bytes) and RAM sizes up to 256 Kbit (4 banks of 0x2000 bytes). The information in this section is based on my MBC1 research, Tauwasser's research notes [6], and Pan Docs [7].

#### 12.1 MBC1 registers

#### Caveat

These registers don't have any standard names and are usually referred to using their address ranges or purposes instead. This document uses names to clarify which register is meant when referring to one.

The MBC1 chip includes four registers that affect the behaviour of the chip. Of the cartridge bus address signals, only A13-A15 are connected to the MBC, so lower address bits don't matter when the CPU is accessing the MBC and all registers are effectively mapped to address ranges instead of single addresses. All registers are smaller than 8 bits, and unused bits are simply ignored during writes. The registers are not directly readable.

Register 12.1: 0x0000-0x1FFF - RAMG - MBC1 RAM gate register

U	U	U	U	W-0	W-0	W-0	W-0	
				RAMG<3:0>				
bit 7	6	5	4	3	2	1	bit 0	

**bit 7-4 Unimplemented**: Ignored during writes

**bit 3-0 RAMG<3:0>**: RAM gate register

0b1010 = enable access to chip RAM

All other values disable access to chip RAM

The RAMG register is used to enable access to the cartridge SRAM if one exists on the cartridge circuit board. RAM access is disabled by default but can be enabled by writing to the 0x0000-0x1FFF address range a value with the bit pattern 0b1010 in the lower nibble. Upper bits don't matter, but any other bit pattern in the lower nibble disables access to RAM.

When RAM access is disabled, all writes to the external RAM area 0xA000-0xBFFF are ignored, and reads return undefined values. Pan Docs recommends disabling RAM when it's not being accessed to protect the contents [7].

#### \* Speculation

We don't know the physical implementation of RAMG, but it's certainly possible that the 0b1010 bit pattern check is done at write time and the register actually consists of just a single bit.

Register 12.2: 0x2000-0x3FFF - BANK1 - MBC1 bank register 1

U	U	U	W-0	W-0	W-0	W-0	W-1	
			BANK1 < 4 : 0>					
bit 7	6	5	4	3	2	1	bit 0	

**bit 7-5 Unimplemented**: Ignored during writes

bit 4-0 BANK1<4:0>: Bank register 1

Never contains the value 0b00000.

If 0b00000 is written, the resulting value will be 0b00001 instead.

The 5-bit BANK1 register is used as the lower 5 bits of the ROM bank number when the CPU accesses the 0x4000-0x7FFF memory area.

MBC1 doesn't allow the BANK1 register to contain zero (bit pattern 0b00000), so the initial value at reset is 0b00001 and attempting to write 0b00000 will write 0b00001 instead. This makes it impossible to read banks 0x00, 0x20, 0x40 and 0x60 from the 0x4000-0x7FFF memory area, because those bank numbers have 0b000000 in the lower bits. Due to the zero value adjustment, requesting any of these banks actually requests the next bank (e.g. 0x21 instead of 0x20).

Register 12.3: 0x4000-0x5FFF - BANK2 - MBC1 bank register 2

U	U	U	U	U	U	W-0	W-0
						BANK2<1:0>	
bit 7	6	5	4	3	2	1	bit 0

bit 7-2 Unimplemented: Ignored during writes

bit 1-0 BANK2<1:0>: Bank register 2

The 2-bit BANK2 register can be used as the upper bits of the ROM bank number, or as the 2-bit RAM bank number. Unlike BANK1, BANK2 doesn't disallow zero, so all 2-bit values are possible.

Register 12.4: 0x6000-0x7FFF - MODE - MBC1 mode register

U	U	U	U	U	U	U	W-0
							MODE
bit 7	6	5	4	3	2	1	bit 0

#### **bit 7-1 Unimplemented**: Ignored during writes

#### **bit 0 MODE**: Mode register

0b1 = BANK2 affects accesses to 0x0000-0x3FFF, 0x4000-0x7FFF, 0xA000-0xBFFF

0b0 = BANK2 affects only accesses to 0x4000-0x7FFF

The MODE register determines how the BANK2 register value is used during memory accesses.

### Warning

Most documentation, including Pan Docs [7], calls value 0b0 ROM banking mode, and value 0b1 RAM banking mode. This terminology reflects the common use cases, but "RAM banking" is slightly misleading because value 0b1 also affects ROM reads in multicart cartridges and cartridges that have a 8 or 16 Mbit ROM chip.

#### 12.2 ROM in the 0x0000-0x7FFF area

In MBC1 cartridges, the A0-A13 cartridge bus signals are connected directly to the corresponding ROM pins, and the remaining ROM pins (A14-A20) are controlled by the MBC1. These remaining pins form the ROM bank number.

When the 0x0000-0x3FFF address range is accessed, the effective bank number depends on the MODE register. In MODE 0b0 the bank number is always 0, but in MODE 0b1 it's formed by shifting the BANK2 register value left by 5 bits.

When the 0x4000-0x7FFF addess range is accessed, the effective bank number is always a combination of BANK1 and BANK2 register values.

If the cartridge ROM is smaller than 16 Mbit, there are less ROM address pins to connect to and therefore some bank number bits are ignored. For example, 4 Mbit ROMs only need a 5-bit bank number, so the BANK2 register value is always ignored because those bits are simply not connected to the ROM.

	ROM address bits					
Accessed address	Bank r	number	Address within bank			
	20-19	18-14	13-0			
0x0000-0x3FFF, MODE = 0b0	0b00	0b00000	A<13:0>			
0x0000-0x3FFF, MODE = 0b1	BANK2	0b00000	A<13:0>			
0x4000-0x7FFF	BANK2	BANK1	A<13:0>			

Table 12.1: Mapping of physical ROM address bits in MBC1 carts

#### **ROM banking example 1**

Let's assume we have previously written 0x12 to the BANK1 register and 0b01 to the BANK2 register. The effective bank number during ROM reads depends on which address range we read and on the value of the MODE register:

Value of the BANK 1 register 0b10010

Value of the BANK 2 register ∅b<mark>∅1</mark>

Effective ROM bank number

(reading 0x0000-0x3FFF, MODE = 0b00000000 (= 0 = 0x00)

**Effective ROM bank number** 

#### **ROM banking example 2**

Let's assume we have previously requested ROM bank number 68, MBC1 mode is 0b0, and we are now reading a byte from 0x72A7. The actual physical ROM address that will be read is going to be 0x1132A7 and is constructed in the following way:

Value of the BANK 1 register 0b00100

Value of the BANK 2 register 0b10

**ROM bank number**  $0b_{1000100}$  (= 68 = 0x44)

**Address being read** 0b0111 0010 1010 0111 (= 0x72A7)

Actual physical ROM address 0b1 0001 0011 0010 1010 0111 (= 0x1132A7)

#### 12.3 RAM in the 0xA000-0xBFFF area

Some MBC1 carts include SRAM, which is mapped to the 0xA000-0xBFFF area. If no RAM is present, or RAM is not enabled with the RAMG register, all reads return undefined values and writes have no effect.

On boards that have RAM, the A0-A12 cartridge bus signals are connected directly to the corresponding RAM pins, and pins A13-A14 are controlled by the MBC1. Most of the time the RAM size is 64 Kbit, which corresponds to a single bank of 0x2000 bytes. With larger RAM sizes the BANK2 register value can be used for RAM banking to provide the two high address bits.

In MODE 0b0 the BANK2 register value is not used, so the first RAM bank is always mapped to the 0xA000–0xBFFF area. In MODE 0b1 the BANK2 register value is used as the bank number.

	RAM address bits				
Accessed address	Bank number	Address within bank			
	14-13	12-0			
0xA000-0xBFFF, MODE = 0b0	0b00	A<12:0>			
0xA000-0xBFFF, MODE = 0b1	BANK2	A<12:0>			

Table 12.2: Mapping of physical RAM address bits in MBC1 carts

#### **RAM banking example 1**

Let's assume we have previously written 0b10 to the BANK2 register, MODE is 0b1, RAMG is 0b1010 and we are now reading a byte from 0xB123. The actual physical RAM address that will be read is going to be 0x5123 and is constructed in the following way:

Value of the BANK 2 register 0b10

Address being read 0b1011 0001 0010 0011 (= 0xB123)

Actual physical RAM address 0b<mark>10</mark>1 0001 0010 0011 (= 0x5123)

#### 12.4 MBC1 multicarts ("MBC1M")

MBC1 is also used in a couple of "multicart" cartridges, which include more than one game on the same cartridge. These cartridges use the same regular MBC1 chip, but the circuit board is wired a bit differently. This alternative wiring is sometimes called "MBC1M", but technically the mapper chip is the same. All known MBC1 multicarts use 8 Mbit ROMs, so there's no definitive wiring for other ROM sizes.

In MBC1 multicarts bit 4 of the BANK1 register is not physically connected to anything, so it's skipped. This means that the bank number is actually a 6-bit number. In all known MBC1 multicarts the games reserve 16 banks each, so BANK2 can actually be considered "game number", while BANK1 is the internal bank number within the selected game. At reset BANK2 is 0600, and the "game" in this slot is actually a game selection menu. The menu code selects MODE 0b1 and writes the game number to BANK2 once the user selects a game.

From a ROM banking point of view, multicarts simply skip bit 4 of the BANK1 register, but otherwise the behaviour is the same. MODE 0b1 guarantees that all ROM accesses, including accesses to 0x0000-0x3FFF, use the BANK2 register value.

	ROM address bits				
Accessed address	Bank number		Address within bank		
	19-18 17-14		13-0		
0x0000-0x3FFF, MODE = 0b0	0b00	0b0000	A<13:0>		
0x0000-0x3FFF, MODE = 0b1	BANK2	0b0000	A<13:0>		
0x4000-0x7FFF	BANK2	BANK1<3:0>	A<13:0>		

Table 12.3: Mapping of physical ROM address bits in MBC1 multicarts

#### **ROM banking example 1**

Let's assume we have previously requested "game number" 3 (= 0b11) and ROM bank number 29 (= 0x1D), MBC1 mode is 0b1, and we are now reading a byte from 0x6C15. The actual physical ROM address that will be read is going to be 0xF6C15 and is constructed in the following way:

0b11

Value of the BANK 1 register 0b11101

Value of the BANK 2 register

**ROM bank number**  $0b_{111101}$  (= 61 = 0x3D)

Address being read 0b0110 1100 0001 0101 (= 0x6C15)

Actual physical ROM address 0b1111 0110 1100 0001 0101 (= 0xF6C15)

#### **Detecting multicarts**

MBC1 multicarts are not detectable by simply looking at the ROM header, because the ROM type value is just one of the normal MBC1 values. However, detection is possible by going through BANK2 values and looking at "bank 0" of each multicart game and doing some heuristics based on the header data. All the included games, including the game selection menu, have proper header data. One example of a good heuristic is logo data verification.

So, if you have a 8 Mbit cart with MBC1, first assume that it's a multicart and bank numbers are 6-bit values. Set BANK1 to zero and loop through the four possible BANK2 values while checking the data at  $0 \times 0.0104 - 0 \times 0.0133$ . In other words, check logo data starting from physical ROM locations  $0 \times 0.0104$ ,  $0 \times 4.0104$ ,  $0 \times 8.0104$ , and  $0 \times 0.0104$ . If proper logo data exists with most of the BANK2 values, the cart is most likely a multicart. Note that multicarts can just have two actual games, so one of the locations might not have the header data in place.

#### 12.5 Dumping MBC1 carts

MBC1 cartridge dumping is fairly straightforward with the right hardware. The total number of banks is read from the header, and each bank is read one byte at a time. However, BANK1 register zero-adjustment and multicart cartridges need to be considered in ROM dumping code.

Banks 0x20, 0x40 and 0x60 can only be read from the 0x0000-0x3FFF memory area and only when MODE register value is 0b1. Using MODE 0b1 has no undesirable effects when doing ROM dumping, so using it at all times is recommended for simplicity.

Multicarts should be detected using the logo check described earlier, and if a multicart is detected, BANK1 should be considered a 4-bit register in the dumping code.

```
BANK1 = 0x2000
BANK2 = 0x4000
MODE = 0x6000
write_byte(MODE, 0x01)
for bank in range(0, num_banks):
    write_byte(BANK1, bank)
    if is_multicart:
        write_byte(BANK2, bank >> 4)
        bank_start = 0x4000 if bank & 0x0f else 0x0000
else:
        write_byte(BANK2, bank >> 5)
        bank_start = 0x4000 if bank & 0x1f else 0x0000
for addr in range(bank_start, bank_start + 0x4000):
        buf += read_byte(addr)
Listing 12.1: Python pseudo-code for MBC1 ROM dumping
```

### **Chapter 13**

### **MBC2** mapper chip

MBC2 supports ROM sizes up to 2 Mbit (16 banks of 0x4000 bytes) and includes an internal 512x4 bit RAM array, which is its unique feature. The information in this section is based on my MBC2 research, Tauwasser's research notes [8], and Pan Docs [7].

#### **\*** Speculation

MBC1 is strictly more powerful than MBC2 because it supports more ROM and RAM. This raises a very important question: why does MBC2 exist? It's possible that Nintendo tried to integrate a small amount of RAM on the MBC chip for cost reasons, but it seems that this didn't work out very well since all later MBCs revert this design decision and use separate RAM chips.

#### 13.1 MBC2 registers

#### Caveat

These registers don't have any standard names and are usually referred to using one of their addresses or purposes instead. This document uses names to clarify which register is meant when referring to one.

The MBC2 chip includes two registers that affect the behaviour of the chip. The registers are mapped a bit differently compared to other MBCs. Both registers are accessible within 0x0000-0x3FFF, and within that range, the register is chosen based on the A8 address signal. In practice, this means that the registers are mapped to memory in an alternating pattern. For example, 0x0000, 0x2000 and 0x3000 are RAMG, and 0x0100, 0x2100 and 0x3100 are ROMB. Both registers are smaller than 8 bits, and unused bits are simply ignored during writes. The registers are not directly readable.

Register 13.1: 0x0000-0x3FFF when A8=0b0 - RAMG - MBC2 RAM gate register

U	U	U	U	W-0	W-0	W-0	W-0
					RAMG	<3:0>	
bit 7	6	5	4	3	2	1	bit 0

**bit 7-4 Unimplemented**: Ignored during writes

**bit 3-0 RAMG<3:0>**: RAM gate register

0b1010 = enable access to chip RAM

All other values disable access to chip RAM

The 4-bit MBC2 RAMG register works in a similar manner as MBC1 RAMG, so the upper bits don't matter and only the bit pattern 0b1010 enables access to RAM.

When RAM access is disabled, all writes to the external RAM area 0xA000-0xBFFF are ignored, and reads return undefined values. Pan Docs recommends disabling RAM when it's not being accessed to protect the contents [7].

#### **\*** Speculation

We don't know the physical implementation of RAMG, but it's certainly possible that the <code>0b1010</code> bit pattern check is done at write time and the register actually consists of just a single bit.

Register 13.2: 0x0000-0x3FFF when A8=0b1 - ROMB - MBC2 ROM bank register

U	U	U	U	W-0	W-0	W-0	W-1
				ROMB<3:0>			
bit 7	6	5	4	3	2	1	bit 0

#### bit 3-0 ROMB<3:0>: ROM bank register

Never contains the value 0b0000.

If 0b0000 is written, the resulting value will be 0b0001 instead.

The 4-bit ROMB register is used as the ROM bank number when the CPU accesses the 0x4000-0x7FFF memory area.

Like MBC1 BANK1, the MBC2 ROMB register doesn't allow zero (bit pattern 0b0000) in the register, so any attempt to write 0b0000 writes 0b0001 instead.

#### 13.2 ROM in the 0x0000-0x7FFF area

In MBC2 cartridges, the A0-A13 cartridge bus signals are connected directly to the corresponding ROM pins, and the remaining ROM pins (A14-A17) are controlled by the MBC2. These remaining pins form the ROM bank number.

When the 0x0000-0x3FFF address range is accessed, the effective bank number is always 0.

When the 0x4000-0x7FFF address range is accessed, the effective bank number is the current ROMB register value.

	ROM address bits				
Accessed address	Bank number	Address within bank			
	17-14	13-0			
0x0000-0x3FFF	0b0000	A<13:0>			
0x4000-0x7FFF	ROMB	A<13:0>			

Table 13.1: Mapping of physical ROM address bits in MBC2 carts

#### 13.3 RAM in the 0xA000-0xBFFF area

All MBC2 carts include SRAM, because it is located directly inside the MBC2 chip. These cartridges never use a separate RAM chip, but battery backup circuitry and a battery are optional. If RAM is not enabled with the RAMG register, all reads return undefined values and writes have no effect.

MBC2 RAM is only 4-bit RAM, so the upper 4 bits of data do not physically exist in the chip. When writing to it, the upper 4 bits are ignored. When reading from it, the upper 4 data signals are not driven by the chip, so their content is undefined and should not be relied on.

MBC2 RAM consists of 512 addresses, so only A0-A8 matter when accessing the RAM region. There is no banking, and the 0xA000-0xBFFF area is larger than the RAM, so the addresses wrap around. For example, accessing 0xA000 is the same as accessing 0xA200, so it is possible to write to the former address and later read the written data using the latter address.

	RAM address bits
Accessed address	
	8-0
0xA000-0xBFFF	A<8:0>

Table 13.2: Mapping of physical RAM address bits in MBC2 carts

#### 13.4 Dumping MBC2 carts

MBC2 cartridges are very simple to dump. The total number of banks is read from the header, and each bank is read one byte at a time. ROMB zero adjustment must be considered in the ROM dumping code, but this only means that bank 0 should be read from 0x0000-0x3FFF and not from 0x4000-0x7FFF like other banks.

```
ROMB = 0x2100
for bank in range(0, num_banks):
    write_byte(ROMB, bank)
    bank_start = 0x4000 if bank > 0 else 0x0000
    for addr in range(bank_start, bank_start + 0x4000):
        buf += read_byte(addr)
    Listing 13.2: Python pseudo-code for MBC2 ROM dumping
```

# Chapter 14 MBC3 mapper chip

MBC3 supports ROM sizes up to 16 Mbit (128 banks of 0x4000 bytes), and RAM sizes up to 256 Kbit (4 banks of 0x2000 bytes). It also includes a real-time clock (RTC) that can be clocked with a quartz crystal on the cartridge even when the Game Boy is powered down. The information in this section is based on my MBC3 research, and Pan Docs [7].

# Chapter 15 MBC30 mapper chip

MBC30 is a variant of MBC3 used by Japanese Pokemon Crystal to support a larger ROM chip and a larger RAM chip. Featurewise MBC30 is almost identical to MBC3, but supports ROM sizes up to 32 Mbit (256 banks of 0x4000 bytes), and RAM sizes up to 512 Kbit (8 banks of 0x2000 bytes). Information in this section is based on my MBC30 research.

### Warning

The circuit board of Japanese Pokemon Crystal includes a 1 Mbit RAM chip, but MBC30 is limited to 512 Kbit RAM. One of the RAM address pins is unused, so half of the RAM is wasted and is inaccessible without modifications. So, the game only uses 512 Kbit and there is a mismatch between accessible and the physical amounts of RAM.

### **Chapter 16**

### MBC5 mapper chip

The majority of games for Game Boy Color use the MBC5 chip. MBC5 supports ROM sizes up to 64 Mbit (512 banks of 0x4000 bytes), and RAM sizes up to 1 Mbit (16 banks of 0x2000 bytes). The information in this section is based on my MBC5 research, and The Cycle-Accurate Game Boy Docs [9].

#### 16.1 MBC5 registers

Register 16.1: 0x0000-0x1FFF - RAMG - MBC5 RAM gate register

W-0	W-0	W-0	W-0	W-0	W-0	W-0	W-0	
RAMG<7:0>								
bit 7	6	5	4	3	2	1	bit 0	

#### bit 7-0 RAMG<7:0>: RAM gate register

0b00001010 = enable access to cartridge RAM All other values disable access to cartridge RAM

The 8-bit MBC5 RAMG register works in a similar manner as MBC1 RAMG, but it is a full 8-bit register so upper bits matter when writing to it. Only 0b00001010 enables RAM access, and all other values (including 0b10001010 for example) disable access to RAM.

When RAM access is disabled, all writes to the external RAM area 0xA000-0xBFFF are ignored, and reads return undefined values. Pan Docs recommends disabling RAM when it's not being accessed to protect the contents [7].

### **\*** Speculation

We don't know the physical implementation of RAMG, but it's certainly possible that the 0b00001010 bit pattern check is done at write time and the register actually consists of just a single bit.

Register 16.2: 0x2000-0x2FFF - ROMB0 - MBC5 lower ROM bank register

W-0	W-0	W-0	W-0	W-0	W-0	W-0	W-1
ROMBØ<7:0>							
bit 7	6	5	4	3	2	1	bit 0

bit 7-0 ROMB0<7:0>: Lower ROM bank register

The 8-bit ROMB0 register is used as the lower 8 bits of the ROM bank number when the CPU accesses the 0x4000-0x7FFF memory area.

Register 16.3: 0x3000-0x3FFF - ROMB1 - MBC5 upper ROM bank register

U	U	U	U	U	U	U	W-0
							ROMB1
bit 7	6	5	4	3	2	1	bit 0

#### **bit 7-1 Unimplemented**: Ignored during writes

#### **bit 0 ROMB1**: Upper ROM bank register

The 1-bit ROMB1 register is used as the most significant bit (bit 9) of the ROM bank number when the CPU accesses the 0x4000-0x7FFF memory area.

Register 16.4: 0x4000-0x5FFF - RAMB - MBC5 RAM bank register

U	U	U	U	W-0	W-0	W-0	W-0
				RAMB<3:0>			
bit 7	6	5	4	3	2	1	bit 0

#### bit 7-4 Unimplemented: Ignored during writes

#### **bit 3-0 RAMB<3:0>**: RAM bank register

The 4-bit RAMB register is used as the RAM bank number when the CPU accesses the 0xA000-0xBFFF memory area.

# Chapter 17 MBC6 mapper chip

MBC6 supports ROM sizes up to 16 Mbit (256 banks of 0x2000 bytes), and RAM sizes up to 4 Mbit (128 banks of 0x1000 bytes). The information in this section is based on my MBC6 research.

# Chapter 18 MBC7

TODO.

# Chapter 19 **HuC-1 mapper chip**

HuC-1 supports ROM sizes up to 8 Mbit (64 banks of 0x4000 bytes), and RAM sizes up to 256 Kbit (4 banks of 0x2000 bytes). It also includes a sensor and a LED for infrared communication. The information in this section is based on my HuC-1 research.

### Chapter 20 HuC-3 mapper chip

HuC-3 supports ROM sizes up to 16 Mbit (128 banks of 0x4000 bytes), and RAM sizes up to 1 Mbit (16 banks of 0x2000 bytes). Like HuC-1, it includes support for infrared communication, but also includes a real-time-clock (RTC) and output pins used to control a piezoelectric buzzer. The information in this section is based on my HuC-3 research.

# Chapter 21 MMM01

TODO.

# Chapter 22 **TAMA5**

TODO.

### **Appendices**

## Appendix A Instruction set tables

These tables include all the opcodes in the Sharp SM83 instruction set. The style and layout of these tables was inspired by the opcode tables available at pastraiser.com [10].

	ж0	<b>x1</b>	<b>x</b> 2	<b>x</b> 3	х4	ж5	жб	ж7	ж8	ж9	ха	жb	жc	жđ	хe	хf
0x	NOP	LD BC,nn	LD (BC),A	INC BC	INC B	DEC B	LD B,n	RLCA	LD (nn),SP	ADD HL,BC	LD A,(BC)	DEC BC	INC C	DEC C	LD C,n	RRCA
1x	STOP	LD DE,nn	LD (DE),A	INC DE	INC D	DEC D	LD D,n	RLA	JR e	ADD HL, DE	LD A,(DE)	DEC DE	INC E	DEC E	LD E,n	RRA
2x	JR NZ,e	LD HL,nn	LD (HL+),A	INC HL	INC H	DEC H	LD H,n	DAA	JR Z,e	ADD HL,HL	LD A,(HL+)	DEC HL	INC L	DEC L	LD L,n	CPL
3x	JR NC,e	LD SP,nn	LD (HL-),A	INC SP	INC (HL)	DEC (HL)	LD (HL),n	SCF	JR C,e	ADD HL,SP	LD A,(HL-)	DEC SP	INC A	DEC A	LD A,n	CCF
4x	LD B,B	LD B,C	LD B,D	LD B,E	LD B,H	LD B,L	LD B,(HL)	LD B,A	LD C,B	LD C,C	LD C,D	LD C,E	LD C,H	LD C,L	LD C,(HL)	LD C,A
5x	LD D,B	LD D,C	LD D,D	LD D,E	LD D,H	LD D,L	LD D,(HL)	LD D,A	LD E,B	LD E,C	LD E,D	LD E,E	LD E,H	LD E,L	LD E,(HL)	LD E,A
6x	LD H,B	LD H,C	LD H,D	LD H,E	LD H,H	LD H,L	LD H,(HL)	LD H,A	LD L,B	LD L,C	LD L,D	LD L,E	LD L,H	LD L,L	LD L,(HL)	LD L,A
7x	LD (HL),B	LD (HL),C	LD (HL),D	LD (HL),E	LD (HL),H	LD (HL),L	HALT	LD (HL),A	LD A,B	LD A,C	LD A,D	LD A,E	LD A,H	LD A,L	LD A,(HL)	LD A,A
8x	ADD B	ADD C	ADD D	ADD E	ADD H	ADD L	ADD (HL)	ADD A	ADC B	ADC C	ADC D	ADC E	ADC H	ADC L	ADC (HL)	ADC A
9x	SUB B	SUB C	SUB D	SUB E	SUB H	SUB L	SUB (HL)	SUB A	SBC B	SBC C	SBC D	SBC E	SBC H	SBC L	SBC (HL)	SBC A
ax	AND B	AND C	AND D	AND E	AND H	AND L	AND (HL)	AND A	XOR B	XOR C	XOR D	XOR E	XOR H	XOR L	XOR (HL)	XOR A
bx	OR B	OR C	OR D	OR E	OR H	OR L	OR (HL)	OR A	CP B	CP C	CP D	CP E	CP H	CP L	CP (HL)	CP A
сx	RET NZ	POP BC	JP NZ,nn	JP nn	CALL NZ,nn	PUSH BC	ADD n	RST 0x00	RET Z	RET	JP Z,nn	СВ ор	CALL Z,nn	CALL nn	ADC n	RST 0x08
dx	RET NC	POP DE	JP NC,nn	-	CALL NC,nn	PUSH DE	SUB n	RST 0x10	RET C	RETI	JP C,nn	-	CALL C,nn	-	SBC n	RST 0x18
ex	LDH (n),A	POP HL	LDH (C),A	-	-	PUSH HL	AND n	RST 0x20	ADD SP,e	JP HL	LD (nn),A	-	-	-	XOR n	RST 0x28
fx	LDH A,(n)	POP AF	LDH A,(C)	DI	-	PUSH AF	OR n	RST 0x30	LD HL,SP+e	LD SP,HL	LD A,(nn)	EI	-	-	CP n	RST 0x38

Table A.1: Sharp SM83 instruction set

### Legend:

	8-bit loads	16-bit loads	8-bit arithmetic/logical	16-bit arithmetic	Rotates, shifts, and bit operations	Control flow	Miscellaneous
[	Undefined						

- **n** unsigned 8-bit immediate data
- **nn** unsigned 16-bit immediate data
- e signed 8-bit immediate data

	ж0	<b>x1</b>	ж2	ж3	x4	x5	ж6	<b>x</b> 7	ж8	<b>x</b> 9	ха	хb	жc	хd	хe	хf
0x	RLC B	RLC C	RLC D	RLC E	RLC H	RLC L	RLC (HL)	RLC A	RRC B	RRC C	RRC D	RRC E	RRC H	RRC L	RRC (HL)	RRC A
1x	RL B	RL C	RL D	RL E	RL H	RL L	RL (HL)	RL A	RR B	RR C	RR D	RR E	RR H	RR L	RR (HL)	RR A
2x	SLA B	SLA C	SLA D	SLA E	SLA H	SLA L	SLA (HL)	SLA A	SRA B	SRA C	SRA D	SRA E	SRA H	SRA L	SRA (HL)	SRA A
3x	SWAP B	SWAP C	SWAP D	SWAP E	SWAP H	SWAP L	SWAP (HL)	SWAP A	SRL B	SRL C	SRL D	SRL E	SRL H	SRL L	SRL (HL)	SRL A
4x	BIT 0,B	BIT 0,C	BIT 0,D	BIT 0,E	BIT 0,H	BIT 0,L	BIT 0,(HL)	BIT 0,A	BIT 1,B	BIT 1,C	BIT 1,D	BIT 1,E	BIT 1,H	BIT 1,L	BIT 1,(HL)	BIT 1,A
5x	BIT 2,B	BIT 2,C	BIT 2,D	BIT 2,E	BIT 2,H	BIT 2,L	BIT 2,(HL)	BIT 2,A	BIT 3,B	BIT 3,C	BIT 3,D	BIT 3,E	BIT 3,H	BIT 3,L	BIT 3,(HL)	BIT 3,A
6x	BIT 4,B	BIT 4,C	BIT 4,D	BIT 4,E	BIT 4,H	BIT 4,L	BIT 4,(HL)	BIT 4,A	BIT 5,B	BIT 5,C	BIT 5,D	BIT 5,E	BIT 5,H	BIT 5,L	BIT 5,(HL)	BIT 5,A
7x	BIT 6,B	BIT 6,C	BIT 6,D	BIT 6,E	BIT 6,H	BIT 6,L	BIT 6,(HL)	BIT 6,A	BIT 7,B	BIT 7,C	BIT 7,D	BIT 7,E	BIT 7,H	BIT 7,L	BIT 7,(HL)	BIT 7,A
8x	RES 0,B	RES Ø,C	RES Ø,D	RES Ø,E	RES Ø,H	RES 0,L	RES 0,(HL)	RES Ø,A	RES 1,B	RES 1,C	RES 1,D	RES 1,E	RES 1,H	RES 1,L	RES 1,(HL)	RES 1,A
9x	RES 2,B	RES 2,C	RES 2,D	RES 2,E	RES 2,H	RES 2,L	RES 2,(HL)	RES 2,A	RES 3,B	RES 3,C	RES 3,D	RES 3,E	RES 3,H	RES 3,L	RES 3,(HL)	RES 3,A
ax	RES 4,B	RES 4,C	RES 4,D	RES 4,E	RES 4,H	RES 4,L	RES 4,(HL)	RES 4,A	RES 5,B	RES 5,C	RES 5,D	RES 5,E	RES 5,H	RES 5,L	RES 5,(HL)	RES 5,A
bx	RES 6,B	RES 6,C	RES 6,D	RES 6,E	RES 6,H	RES 6,L	RES 6,(HL)	RES 6,A	RES 7,B	RES 7,C	RES 7,D	RES 7,E	RES 7,H	RES 7,L	RES 7,(HL)	RES 7,A
сх	SET 0,B	SET 0,C	SET 0,D	SET 0,E	SET 0,H	SET 0,L	SET 0,(HL)	SET 0,A	SET 1,B	SET 1,C	SET 1,D	SET 1,E	SET 1,H	SET 1,L	SET 1,(HL)	SET 1,A
dx	SET 2,B	SET 2,C	SET 2,D	SET 2,E	SET 2,H	SET 2,L	SET 2,(HL)	SET 2,A	SET 3,B	SET 3,C	SET 3,D	SET 3,E	SET 3,H	SET 3,L	SET 3,(HL)	SET 3,A
ex	SET 4,B	SET 4,C	SET 4,D	SET 4,E	SET 4,H	SET 4,L	SET 4,(HL)	SET 4,A	SET 5,B	SET 5,C	SET 5,D	SET 5,E	SET 5,H	SET 5,L	SET 5,(HL)	SET 5,A
fx	SET 6,B	SET 6,C	SET 6,D	SET 6,E	SET 6,H	SET 6,L	SET 6,(HL)	SET 6,A	SET 7,B	SET 7,C	SET 7,D	SET 7,E	SET 7,H	SET 7,L	SET 7,(HL)	SET 7,A

Table A.2: Sharp SM83 CB-prefixed instructions

# Appendix B Memory map tables

	bit 7	6	5	4	3	2	1	bit 0					
0xFF00 P1			P15 buttons	P14 d-pad	P13 <b>⊙</b> start	P12 <b>⊙</b> select	Р11 <b>⊙</b> в	P10 <b>⊘</b> A					
0xFF01 SB		•	•	SB<	7:0>	•		•					
0xFF02 SC	SIO_EN						SIO_FAST	SIO_CLK					
0xFF03		•						•					
0xFF04 DIV				DIVH	<7:0>								
0xFF05 TIMA	TIMA<7:0>												
0xFF06 TMA				TMA	<7:0>								
0xFF07 TAC						TAC_EN	TAC_C	LK<1:0>					
0xFF08													
0xFF09													
0xFF0A													
0xFF0B													
0xFF0C													
0xFF0D													
0xFF0E													
0xFF0F IF				IF_JOYPAD	IF_SERIAL	IF_TIMER	IF_STAT	IF_VBLANK					
0xFF10 NR10													
0xFF11 NR11													
0xFF12 NR12													
0xFF13 NR13													
0xFF14 NR14													
0xFF15													
0xFF16 NR21													
0xFF17 NR22													
0xFF18 NR23													
0xFF19 NR24													
0xFF1A NR30													
0xFF1B NR31													
0xFF1C NR32													
0xFF1D NR33													
0xFF1E NR34													
0xFF1F													
	bit 7	6	5	4	3	2	1	bit 0					

Table B.3: 0xFFxx registers: 0xFF00-0xFF1F

	bit 7	6	5	4	3	2	1	bit 0
0xFF20 NR41								
0xFF21 NR42								
0xFF22 NR43								
0xFF23 NR44								
0xFF24 NR50								
0xFF25 NR51								
0xFF26 NR52								
0xFF27		•						
0xFF28								
0xFF29								
0xFF2A								
0xFF2B								
0xFF2C								
0xFF2D								
0xFF2E								
0xFF2F								
0xFF30 WAV00								
0xFF31 WAV01								
0xFF32 WAV02								
0xFF33 WAV03								
0xFF34 WAV04								
0xFF35 WAV05								
0xFF36 WAV06								
0xFF37 WAV07								
0xFF38 WAV08								
0xFF39 WAV09								
0xFF3A WAV10								
0xFF3B WAV11								
0xFF3C WAV12								
0xFF3D WAV13								
0xFF3E WAV14								
0xFF3F WAV15								
	bit 7	6	5	4	3	2	1	bit 0

Table B.4: 0xFFxx registers: 0xFF20-0xFF3F

	bit 7	6	5	4	3	2	1	bit 0
0xFF40 LCDC	LCD_EN	WIN_MAP	WIN_EN	TILE_SEL	BG_MAP	OBJ_SIZE	OBJ_EN	BG_EN
0xFF41 STAT		INTR_LYC	INTR_M2	INTR_M1	INTR_M0	LYC_STAT	LCD_M(	DDE<1:0>
0xFF42 SCY								
0xFF43 SCX								
0xFF44 LY								
0xFF45 LYC								
0xFF46 DMA		•		DMA	<7:0>	•		•
0xFF47 BGP								
0xFF48 OBP0								
0xFF49 OBP1								
0xFF4A WY								
0xFF4B WX								
0xFF4C ????								
0xFF4D KEY1	KEY1_FAST							KEY1_EN
0xFF4E	_			•	•	•	•	
0xFF4F VBK							VBK	<1:0>
0xFF50 BOOT								BOOT_OFF
0xFF51 HDMA1								
0xFF52 HDMA2								
0xFF53 HDMA3								
0xFF54 HDMA4								
0xFF55 HDMA5								
0xFF56 RP								
0xFF57								
0xFF58								
0xFF59								
0xFF5A								
0xFF5B								
0xFF5C								
0xFF5D								
0xFF5E								
0xFF5F		Ţ ,		•				
	bit 7	6	5	4	3	2	1	bit 0

Table B.5: 0xFFxx registers: 0xFF40-0xFF5F

	bit 7	6	5	4	3	2	1	bit 0		
0xFF60 ????										
0xFF61										
0xFF62										
0xFF63										
0xFF64										
0xFF65										
0xFF66										
0xFF67										
0xFF68 BCPS										
0xFF69 BPCD										
0xFF6A OCPS										
0xFF6B OCPD										
0xFF6C ????										
0xFF6D			•							
0xFF6E										
0xFF6F										
0xFF70 SVBK							SVB	K<1:0>		
0xFF71										
0xFF72 ????										
0xFF73 ????										
0xFF74 ????										
0xFF75 ????										
0xFF76 PCM12		PCM1	2_CH2	!		PCM1	2 CH1			
0xFF77 PCM34			 34_CH4		PCM34_CH3					
0xFF78										
0xFF79										
0xFF7A										
0xFF7B										
0xFF7C										
0xFF7D										
0xFF7E										
0xFF7F										
0xFFFF IE		IE_UNUSED<2:0>		IE_JOYPAD	IE_SERIAL	IE_TIMER	IE_STAT	IE_VBLANK		
	bit 7	6	5	4	3	2	1	bit 0		

Table B.6: 0xFFxx registers: 0xFF60-0xFF7F, 0xFFFF

## Appendix C Game Boy external bus

#### **C.1 Bus timings**

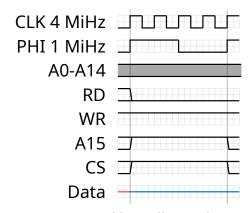


Figure C.5: External bus idle machine cycle

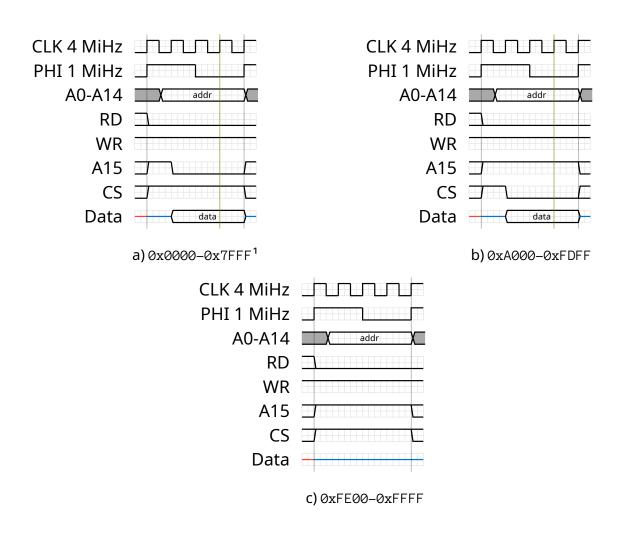


Figure C.6: External bus CPU read machine cycles

 $<sup>^{1}</sup>$ Does not apply to 0x0000-0x00FF accesses while the boot ROM is enabled. Boot ROM accesses do not affect the external bus, so it is in the idle state.

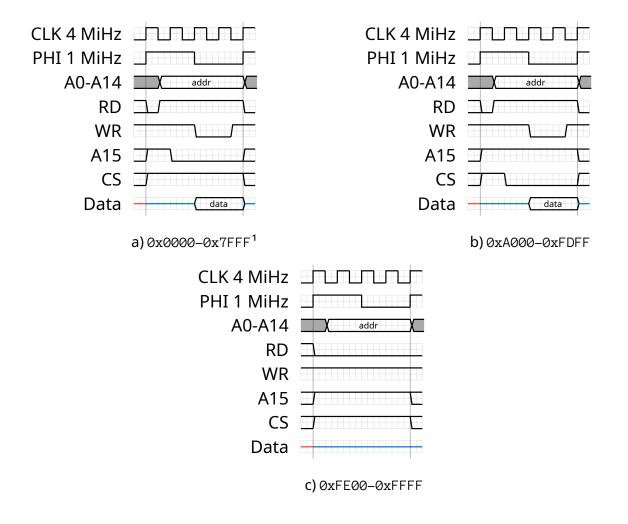


Figure C.7: External bus CPU write machine cycles

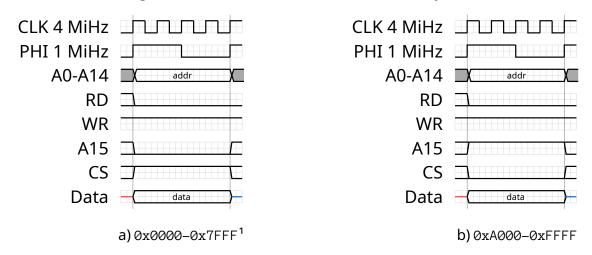


Figure C.8: External bus timings for OAM DMA read machine cycles

## Appendix D Chip pinouts

#### **D.1 CPU chips**

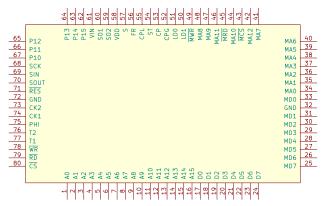


Figure D.9: DMG/SGB CPU (Sharp QFP080-P-1420)

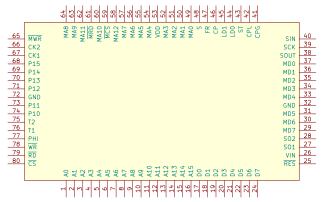


Figure D.10: MGB/SGB2 CPU (Sharp QFP080-P-1420)

#### **D.2 Cartridge chips**

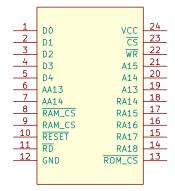


Figure D.11: MBC1 (Sharp SOP24-P-450) [6]

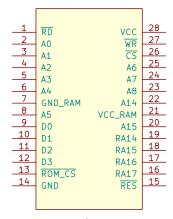


Figure D.12: MBC2 (Sharp SOP28-P-450) [8]

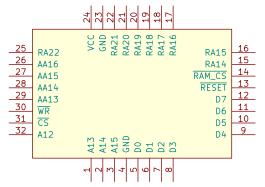


Figure D.13: MBC5 (Sharp QFP32-P-0707)

### **Bibliography**

- [1] "Microcomputers Data Book (1990)." Sharp Corporation. [Online]. Available: https://archive.org/details/bitsavers\_sharpdataBomputersDataBook\_31076011
- [2] "Microcomputer Data Book (1996)." Sharp Corporation. [Online]. Available: https://archive.org/details/bitsavers\_sharpdataBomputerDataBook\_13840187
- [3] "SM8311/SM8313/SM8314/SM8315 8-Bit Single-Chip Microcomputers (Controllers For Home Appliances)." Sharp Corporation.
- [4] Costis Sideris, "The quest for dumping GameBoy Boot ROMs!." [Online]. Available: http://www.its.caltech.edu/~costis/sgb\_hack/
- [5] gekkio, "Dumping the Super Game Boy 2 boot ROM." [Online]. Available: https://gekkio. fi/blog/2015/dumping-the-super-game-boy-2-boot-rom/
- [6] Tauwasser, "MBC1 Tauwasser's Wiki." [Online]. Available: https://wiki.tauwasser.eu/view/MBC1
- [7] Pan of ATX, M. Fayzullin, P. Felber, P. Robson, and M. Korth, "Pan Docs Everything You Always Wanted To Know About GAMEBOY." [Online]. Available: http://bgb.bircd.org/pandocs.htm
- [8] Tauwasser, "MBC2 Tauwasser's Wiki." [Online]. Available: https://wiki.tauwasser.eu/view/MBC2
- [9] Antonio Niño Díaz (AntonioND), "The Cycle-Accurate Game Boy Docs." [Online]. Available: https://github.com/AntonioND/giibiiadvance/tree/master/docs
- [10] "Gameboy CPU (LR35902) instruction set." [Online]. Available: http://www.pastraiser.com/cpu/gameboy/gameboy\_opcodes.html