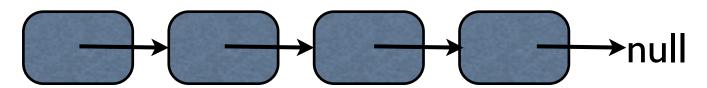
## Computer Science 210: Data Structures

Linked lists

### Arrays vs. Linked Lists

- We've seen arrays:
  - int[] a = new int[10];
  - a is a chunk of memory of size 10 x sizeof(int)
  - a has a fixed size

- A linked list is fundamentally different way of storing collections
  - each element stores a reference to the element after it



#### Arrays vs. Linked Lists

#### Arrays

- have a pre-determined fixed size
- easy access to any element a[i] in constant time
- no space overhead
  - Size = n x sizeof(element)

#### Linked lists

- no fixed size; grow one element at a time
- space overhead
  - each element must store an additional reference
  - Size = n x sizeof (element) + n x sizeof(reference)
- no easy access to i-th element wrt the head of the list
  - · need to hop through all previous elements

#### Linked-lists in Java

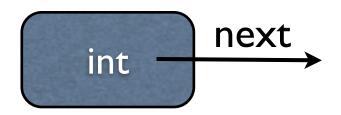
- Search for class Java LinkedList
- Has all expected methods and features
  - add(int index, Object element)
  - add(Object o)
  - addAll(Collection c)
  - addAll(int index, Collection c)
  - addFirst(Object o)
  - addLast(Object o)
  - contains(Object o)
  - get(int index)
  - getFirst()
  - getLast()
  - indexOf(Object o)
  - lastIndexOf(Object o)

- remove(int index)
- remove(Object o)
- removeFirst()
- removeLast()
- set(int index, Object element)
- size()

## Implementing a linked list

- We want to implement a linked list class, much like Java's LinkedList
- For simplicity, we can think of a linked list of integers

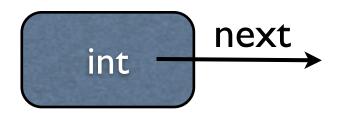
#### The Node class



We want to define the node in a list linked of integers.

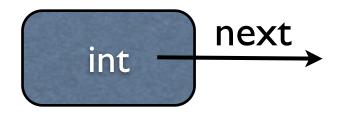
```
/** Node of a singly linked list of integers */
public class Node {
...
```

#### The Node class



We want to define the node in a list linked of integers.

#### The Node class



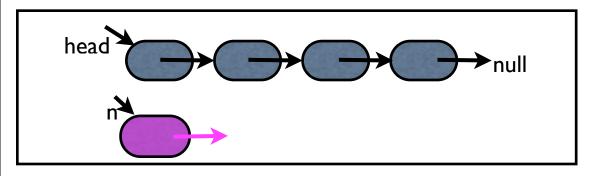
```
/** Node of a singly linked list of integers */
public class Node {
 private int element; // we assume elements are ints
 private Node next;
  /** Creates a node with the given element and next node. */
 public Node(int s, Node n) {
    element = s;
    next = n;
  /** Returns the element of this node. */
 public int getElement() { return element; }
  /** Returns the next node of this node. */
 public Node getNext() { return next; }
  // Modifier methods:
  /** Sets the element of this node. */
 public void setElement(int newElem) { element = newElem; }
  /** Sets the next node of this node. */
 public void setNext(Node newNext) { next = newNext; }
```

# A Single-Linked-List class head head head

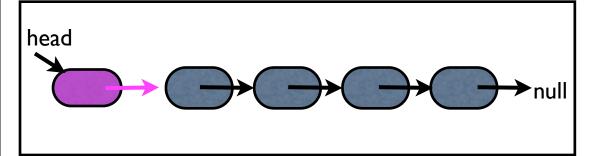
# A Single-Linked-List class head head head

- We'll discuss the following methods
  - addFirst(Node n)
  - addAfter(Node n)
  - Node get(int i)
  - Node removeFirst()
  - addLast(Node n)
  - removeLast(Node n)

#### Inserting at head



void addFirst(Node n)

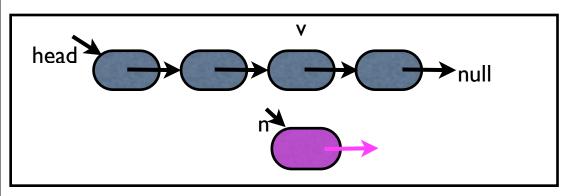


```
void addFirst(Node n) {
    n.setNext(head);
    head = n;
    size++;
}
```

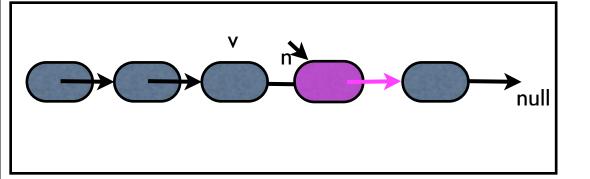
#### Notes

- Special cases: works when head is null, i.e. list is empty
- Efficiency: O(1) time

### Inserting in the middle



void insertAfter(Node v, Node n)



```
//insert node n after node v

void insertAfter(Node v, Node n)
    n.setNext(v.getNext());
    v.setNext(n);
    size++;
}
```

- Notes:
  - Efficiency: O(1)
  - Special cases
    - does not work if v or n are null
      - null pointer exception

#### Get the i-th element

```
//return the i-th node
Node get(int i) {
    ...
```

#### Get the i-th element

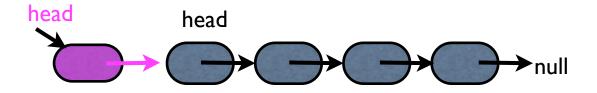
```
//return the i-th node

Node get(int i) {
   if (i >= size) print error message and return null
   Node ptr = head;
   for (int k=0; k<i; k++)
       ptr = ptr.getNext();
   return ptr;
}</pre>
```

#### Notes

- Special cases
  - does it work when list is empty?
- Efficiency: takes O(i) time
  - · constant time per element traversed
  - · unlike arrays, accessing i-th element is not constant time

#### Remove at head



#### Insert at tail

```
void addLast(Node n) {
   insertAfter (get(size), n);
}
```

- Notes
  - Special cases
    - does it work when list is empty?
      - Nope (first node in insertAfter is null).
      - How to fix it?
  - Efficiency: takes O(size) time

#### Delete at tail

- Remove at end: similar
  - need to get to the last element from the head
  - O(size) time

#### Linked lists

- Single-linked lists support insertions and deletions at head in Theta(1) time.
- Insertions and deletion at the tail can be supported in O(size) time.

addFirst: O(1) time

removeFirst: O(1) time

• addLast: O(size) time

removeLast: O(size) time

- Why? because we keep track of the head.
  - To access the tail in constant time, need to keep track of tail as well.

#### Linked-list with tail

```
/** Singly linked list .*/
public class SLinkedList {
 private Node head, tail; // head and tail nodes of the list
 private long size; // number of nodes in the list
void SLinkedList() {
   head = tail = null;
                                           all methods must update tail
   size = 0;
}
void addFirst(Node n) {...}
Node removeFirst() {...}
```

#### Insert at tail

```
void addLast(Node n) {
   //if list is empty the new element is head and tail
   if (tail == null) {
      n.setNext(null);
      head = tail = n;
   } else {
      //the list is not empty: link tail to n and n becomes the new
      tail
      tail.setNext(n);
      n.setNext(null);
      tail = n;
   //increment size
   size++
}
```

- Special cases: list is empty
- Efficiency: Theta(1)

#### Remove at tail

- · What we want: delete the last element and set the new tail
- Is that possible?

#### Remove at tail

- What we want: delete the last element and set the new tail
- Is that possible?
- Remove at tail
  - set the tail to the node BEFORE the tail
  - need the node before the tail: O(size)
- To remove an element from a list you need the node BEFORE it as well

```
remove(Node n) {
    //link n.before to n.next
}
```

To remove a node efficiently need to keep track of previous node

### Doubly-linked lists



```
/** Node of a doubly linked list of integers */
public class DNode {
 protected int element; //element stored by a node
 protected DNode next, prev; // Pointers to next and previous nodes
  /** Constructor that creates a node with given fields */
 public DNode(int e, DNode p, DNode n) {
    element = e;
   prev = p;
   next = n;
  /** Returns the element of this node */
 public int getElement() { return element; }
  /** Returns the previous node of this node */
 public DNode getPrev() { return prev; }
  /** Returns the next node of this node */
 public DNode getNext() { return next; }
  /** Sets the element of this node */
 public void setElement(Int newElem) { element = newElem; }
  /** Sets the previous node of this node */
 public void setPrev(DNode newPrev) { prev = newPrev; }
  /** Sets the next node of this node */
 public void setNext(DNode newNext) { next = newNext; }
```

#### Doubly-linked lists

```
/** Doubly linked list with nodes of type DNode */
public class DList {
  protected int size; // number of elements
  protected DNode head, tail;

  void addFirst(Node n);
  void addLast(Node n);
  Node deleteFirst();
  Node deleteLast();
  void delete(Node n);
}
```

Operations on doubly linked lists

```
addFirst(): O(1) time
addLast(): O(1) time
deleteFirst(): O(1) time
deleteLast(): O(1) time
delete(): O(1) time
get(i): O(i) time
```

#### Insert at head

```
void addFirst(Node n) {
    n.setNext(head);
    n.setprev(null);
    head.setPrev(n);
    head = n;
    size++;
}
```

Does this work?

#### Insert at head

```
void addFirst(Node n) {
                                         void addFirst(Node n) {
                                           if (head==null) {
     n.setNext(head);
     n.setprev(null);
                                              /* this is the first
                                              element: set both head
     head.setPrev(n);
                                              and tail to it */
     head = n;
                                              head = tail = n;
                                              n.setPrev(null);
     size++;
                                              n.setNext(null);
Special cases?
                                           else {

    empty list: head is null; need to

                                           n.setNext(head);
     set tail too
                                           n.setprev(null);
                                           head.setPrev(n);
 Efficiency?
                                           head = n;
   • O(1)
                                           size++;
```

}

#### Insert at tail

```
void addLast(Node n) {
   tail.setNext(n);
   n.setprev(tail);
   n.setNect(null);
   tail = n;
   size++;
}
```

Does this work?

#### Insert at tail

```
void addLast(Node n) {
                                             void addLast(Node n) {
                                                if (tail == null) {
     tail.setNext(n);
                                                 head = tail = n;
     n.setprev(tail);
                                                 n.setPrev(null);
     n.setNect(null);
                                                 n.setNext(null);
     tail = n;
                                                else {
     size++;
                                                 tail.setNext(n);
  }
                                                 n.setprev(tail);
                                                 n.setNect(null);
                                                 tail = n;
Special cases?
                                                size++;

    empty list: tail is null; need to set head

     too
```

• Efficiency: O(1)

## Doubly-linked lists

 Class work: Sketch the following methods for doubly-linked lists, and analyze their efficiency.

Node removeFirst()

Node removeLast()

void remove(Node n)

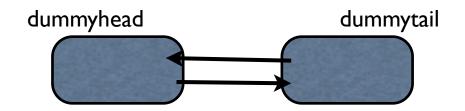
• Node search(int k)

#### Sentinels

- Sentinels for singly-linked list: keep a dummy head
  - an empty list is one node: the dummy head
- Sentinels for doubly-linked lists
  - dummy head and dummy tail
- · Why? elegant. Unifies special cases when head or tail are null

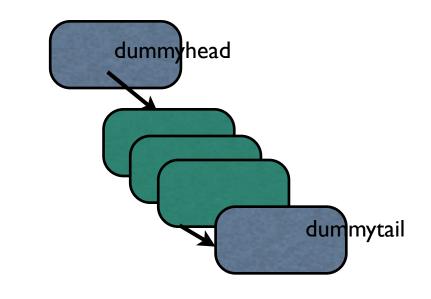
#### DLLists with Sentinels

- the empty list:
  - size = 0

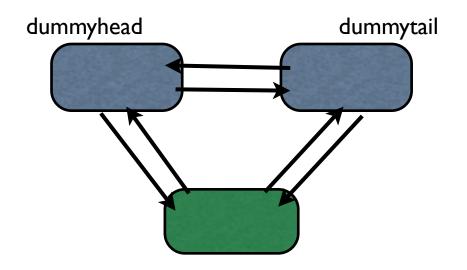


#### DLLists with sentinels

```
insertFirst(Node n) {
    n.setNext(dummyHead.getNext());
    dummyHead.getNext().setPrev(n);
    dummyHead.setNext(n);
    n.setPrev(dummyhead);
    size++;
}
```



- Special cases: none
  - works for empty list



#### Extensions

Circular lists: make last node point to the first (instead of null)

```
class CircularList {
    SNode head;
    int size;
}
```

Let's say we want to insert at head

```
insertAtHead(Node n) {
    n.setNext(head.getNext());
    head.setNext(n);
}
```

If head is null?

```
if (head ==null) {
    n.setNext(n);
    head = n;
}
```

