

Comp 250 Homework3

Question2

$$\begin{aligned}
 \text{for } n > 1. \quad T(n) &= 2T(n-1) + n \\
 &= 2(2T(n-2) + n-1) + n \\
 &= 4T(n-2) + 2n + n - 2 \\
 &= 4(2T(n-3) + n-2) + (2+1)n - 2 \\
 &= 8T(n-3) + (4+2+1)n - (4 \times 2 + 2 \times 1) \\
 &= 2^k T(n-k) + n \sum_{i=0}^{k-1} 2^i - \sum_{i=1}^{k-1} i 2^i \\
 &= 2^k T(n-k) + n(2^k - 1) - 2(1 + 2^{k-1}(k-2)) \\
 \text{now } k &= n-1 \\
 T(1) &= 1 \\
 &= 2^{n-1} + n(2^{n-1} - 1) - 2 - 2^{n-1}(n-3) \\
 &= 2^{n-1}(1 + n - n + 3) - n - 2 \\
 &= 4 \cdot 2^{n-1} - n - 2 \\
 &= 2^{n+1} - n - 2
 \end{aligned}$$

Question3:

Algorithm orderCards(n)

Input: An integer n

Output: Prints the correct card ordering.

//Pseudocode

//Use a double-linked list to implement the method.

for i ← n to 1 do{

 //add new node at top

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Node k ← new Node(i);
k.next=head;
head=k;

//move
Node k ← tail;
tail ← k.prev;
k.next ← head;
head ← k;
}
Node k ← new Node(1);
k.next=head;
head=k;

```

Question 4

$\log(n!) = \sum_{i=1}^n \log i \leq n \log n$ for $n \geq 1$.
 $\therefore \log(n!) \in O(n \log n)$

"Proof for $\log(n!) \in \Omega(n \log n)$.
~~Induct~~ \Downarrow

There are constant $c > 0$ and n_0 such that $\log(n!) \geq cn \log n$ for every $n \geq n_0$.

Induction hypothesis: $\log((n-1)!) \geq c(n-1) \log(n-1)$

$e^x = 1 + \frac{x}{1} + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots$
 for any $x > 0$, $e^x > 1 + x$.
 $e^{\frac{1}{n-1}} > 1 + \frac{1}{n-1} = \frac{n}{n-1}$
 $\frac{1}{n-1} \log e > \log\left(\frac{n}{n-1}\right)$
 $= \log n - \log(n-1)$
 $\log(n-1) > \log n - \log e / (n-1)$ for $c = \frac{1}{2}$, $\log n > \log e$.
 $n \geq 3$. $c(1-c) \log n - c \log e > 0$.
 $\therefore \log(n!) \geq cn \log n$.

Base case: $n=2$.
 $\log(2!) = 1 \geq \frac{1}{2} \log 2$
 \therefore for every $n \geq n_0 = 2$, $\log(n!) \in \Omega(n \log n)$.
 $\therefore \log(n!) \in O(n \log n)$ & $\log(n!) \in \Omega(n \log n)$.
 $\therefore \log(n!) \in \Theta(n \log n)$.