

```
double answer = 1, d = 0.00001
while (answer * answer <= num) {
    answer += d;
}
if (answer * answer > num) {
    answer -= d;
}
return (int) (answer * 1000000000000000000.0);
}
```

Subtract 11

loop invariant

~~1. $\frac{1}{2} \leq d \leq 1$~~

true when $\angle num$

answer
 $\text{answer} \leftarrow \max \{ 0 \leq k \leq \text{answer} + d \}$
 answer - d }

$$\text{answer} \pm d = \text{answer} = \text{sqrt}(C/n)$$

loop is valid

answer $1.1 \approx 1 \quad 1 - 0.000001 \approx 1 \quad \approx 591 + \text{num}$

Inductive Property Assm

At some point since answer ~~grows~~^{may} reach 10^9 ,
it will overflow equalnum so terminates loop
and returns $\text{sqrt}(1) \approx 1$

~~After loop is done~~

2 continued on other
page

2 continued

Prove total correctness
since code has a loop maintainer
So invariant

$$\text{answer} = (\text{answer} - d) \quad \text{AND } \text{answer} > 0$$

So it can terminate

Satisfies criteria for loop correctness

$$P \Rightarrow I$$

$$(I \text{ and } B) \Rightarrow S$$

$$(I \text{ and } (\text{not } B)) \Rightarrow Q$$

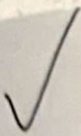
loop terminator

Satisfies pos L condition of being

$$S \Rightarrow (I \text{ and } B)$$

It also follows the format of

$\{P\} \text{ while } \underline{B} \text{ do } \underline{S} \text{ end } \{Q\}$



Worst P (10)

1)

If $(x > y) \rightarrow y = 2x + 1$ else $y = 3x - 1$ ($y > 3$)
 When 1st case $x > y$ $y = 2x + 1$ and $y > 3, x > 1$
 if not $y = 3x - 1$
 $y > 3, x > 2$

→

W.P. $\{x > 1 \text{ or } x > 2\}$

2) $x = y + 10$ ($y > 10$)

when $y = 10$ $x = 10 + 10 = 20 > 20$

W.P. $\{x > 20\}$

3)

$y = y + 1$; $z = x + y$; $\{x < 2\}$

if last condition

$z = x + y$ and when $y > 1 \rightarrow x < 2$

W.P. $\{y > 1\}$

4) $x = 3y + 1$; $y = x - 4$; $\{y < 0\}$

$$3 \cdot (x - 4) + 1$$

$$3x - 12 + 1 = 3x - 11$$

this case y is negative so < 0

Answer
logic

So $x - 4 < 0$
 ~~$3x - 12 + 1 = 11/3$~~
 $3x - 12 + 1 = 11/3$
 W.P. $\{y < \frac{11}{3}\}$

~~W.P. $\{y < 1\}$~~
 So $y < 0$

W.P. $\{y < 1\}$

if $y < 0$ $y = (3y + 1) - 4 \rightarrow 3y - 3 = 3y = 3 \rightarrow y = 1$

5 w.p for $(x > 0)$ then $y = y + 2$ else $y = y + 3$ & $y > 0$

$$y = y + 2$$

$$y = y + 3$$

w.p. $\{y \geq -2 \text{ or } y \geq -3\}$

5

Java do while

do {

code to be executed }

while (condition)

so if \rightarrow loop: code to be executed

if (condition) = false goto out

goto loop

out: ...

Ada for

for I in start \rightarrow complete loop

I = first

loop: If I < last go out

\rightarrow so I = I + 1

go to loop

out: ...

C++ if then else

if (condition) {

Statement 1;

Statement 2; }

} \rightarrow code

else {

statement 3 }

}

\rightarrow

So if condition == false goto else code

goto out

else:

statement 3

C For

```
for (expressions 1, 2, 3)
    eval (expression 1);
loop = control = eval (expression 2)
if (control == 0)
    goto out
    evaluate (expression 3)
    goto loop
out: ...
```

C Switch

```
Switch (i) {
    case 1:
        Statement-1 x;
        break;
```

```
    case 2:
        Statement y;
        break;
```

```
    default }
```

→ if matched with case 1 then $i = 1$
if case 2 then $i = 2$
if none

→ goto default

6)

Denotational Semantics

→ taken from Dynamic Semantics
power point

Python for

$M(\text{For } f, \text{state } \text{state}) = M_f(f, M(f.\text{initial}, \text{state}))$
 $M_f(\text{For } f, \text{state } \text{state}) =$

$\{M(f, M_f(f.\text{pos}, M(f.\text{body}, \text{state})))$

if $(M(f.\text{test}, \text{state}) \text{ is true })$

{ state otherwise }

Java class

Add Simple Varmap

$M_e() = \text{VARMAP}(i, s)$
 then $S = \{ \langle i_1, v_1 \rangle \langle i_2, v_2 \rangle \dots \}$

C++ function

Varmap

if $M_s(B, S)$

then

S

C++ ~~for~~ switch

when switch := true;

~~switch~~

$m, (while\ B\ do\ L, s)\ \Delta =$

if $m_b(B, s) == error$

then error

else if $m_b(B, s)$

$== false$

then s

else if $m_s(L, s) == error$

then error

else $m, (while\ B\ do\ L,$

$m_s, (L, s))$.

then a switch (not switch) in core where stop

end loop

java boolean expression

$m_b(B, s)$ if $VARMAP(l, s) = \text{undefined}$

for some l in B

\hookrightarrow error

else

B

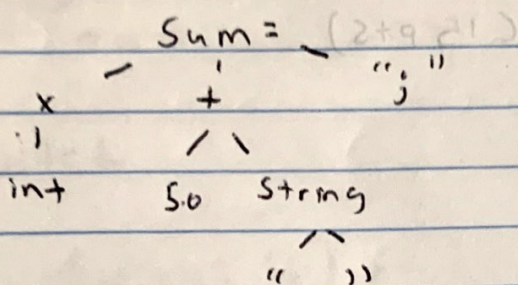
$VARMAP$

Decorated tree

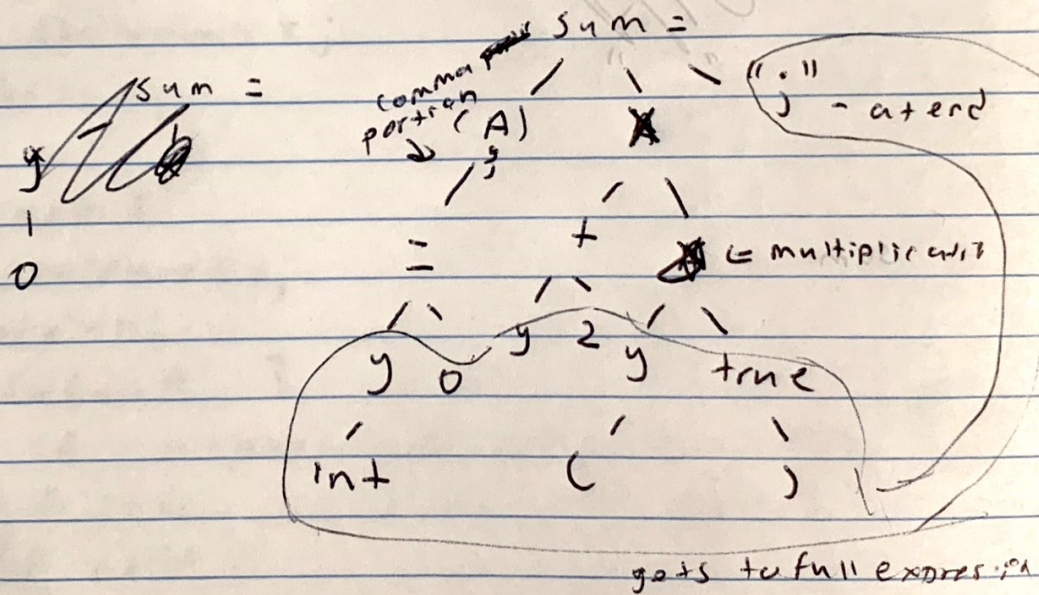
1)

decorated tree (15)

a)



b)



c)

