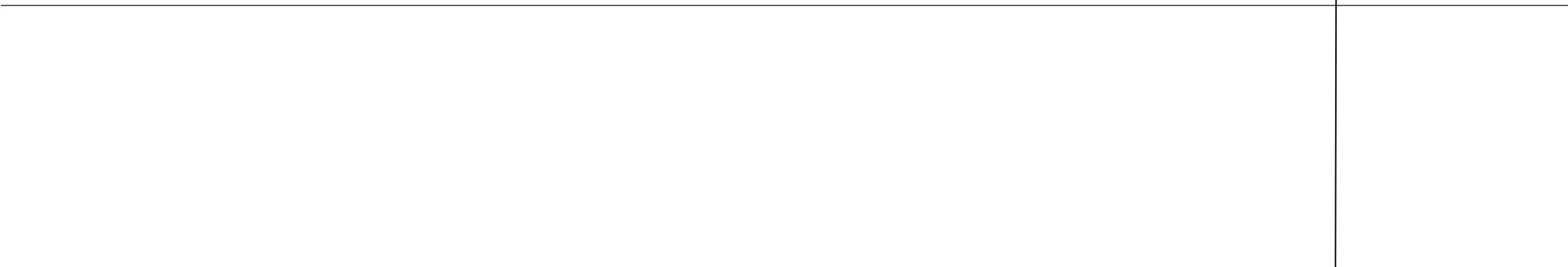


分而治之篇：最大子数组问题I



问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组为数组X中**连续**的一段序列

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为： $3 + 5 - 4 + 3 + 2 = 9$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$
 - 求和为: $-1 - 3 + 3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 12$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$
 - 求和为: $-1 - 3 + 3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 12$
- 子数组 $X[3..10]$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$
 - 求和为: $-1 - 3 + 3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 12$
- 子数组 $X[3..10]$
 - 求和为: $3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 16$

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$
 - 求和为: $-1 - 3 + 3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 12$
- 子数组 $X[3..10]$
 - 求和为: $3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 16$

问题：寻找数组X最大的非空子数组？

问题背景：子数组

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 子数组 $X[3..7]$
 - 求和为: $3 + 5 - 4 + 3 + 2 = 9$
- 子数组 $X[1..10]$
 - 求和为: $-1 - 3 + 3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 12$
- 子数组 $X[3..10]$
 - 求和为: $3 + 5 - 4 + 3 + 2 - 2 + 3 + 6 = 16$

问题：寻找数组X最大的非空子数组？

答案： $X[3..10] = 16$

问题定义

- 形式化定义

最大子数组问题

Max Continuous Subarray, MCS

输入

- 给定一个数组 $X[1..n]$ ，对于任意一对数组下标为 l, r ($l \leq r$)的非空子数组，其和记为

$$S(l, r) = \sum_{i=l}^r X[i]$$

输出

- 求出 $S(l, r)$ 的最大值，记为 S_{max}

蛮力枚举

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

蛮力枚举

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 数组 $X[1..n]$, 其所有的下标 $l, r (l \leq r)$ 组合分为以下两种情况

蛮力枚举

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 数组 $X[1..n]$, 其所有的下标 $l, r (l \leq r)$ 组合分为以下两种情况
 - 当 $l = r$ 时, 一共 $C_n^1 = n$ 种组合

蛮力枚举

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 数组 $X[1..n]$, 其所有的下标 $l, r (l \leq r)$ 组合分为以下两种情况
 - 当 $l = r$ 时, 一共 $C_n^1 = n$ 种组合
 - 当 $l < r$ 时, 一共 C_n^2 种组合

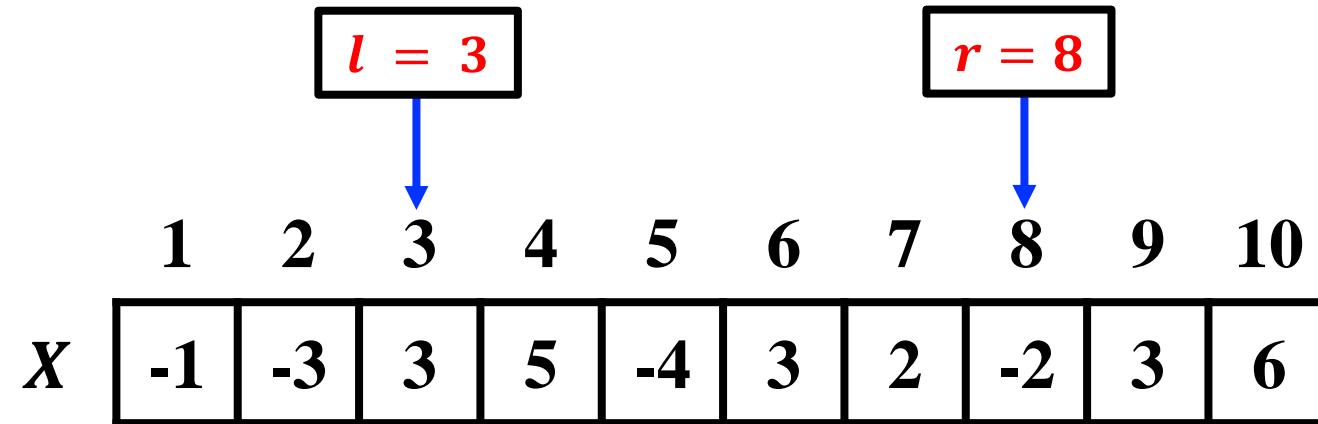
蛮力枚举

	1	2	3	4	5	6	7	8	9	10
X	-1	-3	3	5	-4	3	2	-2	3	6

- 数组 $X[1..n]$, 其所有的下标 $l, r (l \leq r)$ 组合分为以下两种情况
 - 当 $l = r$ 时, 一共 $C_n^1 = n$ 种组合
 - 当 $l < r$ 时, 一共 C_n^2 种组合
- 枚举 $n + C_n^2$ 种下标 l, r 组合, 求出最大子数组之和

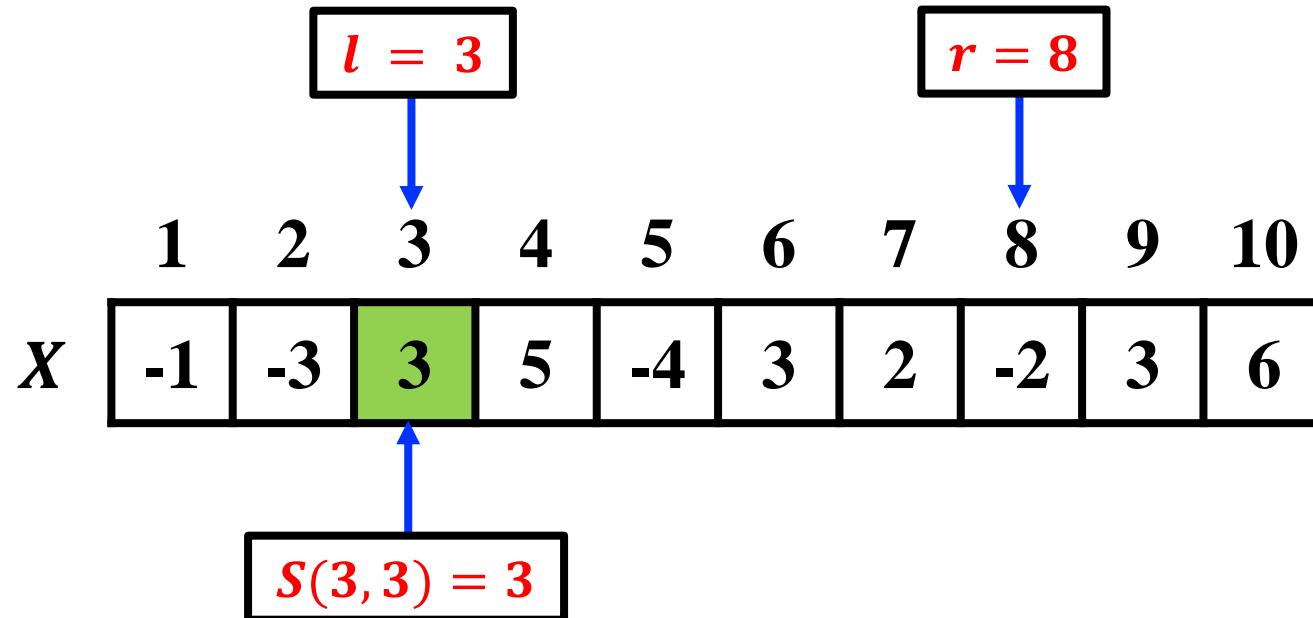
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



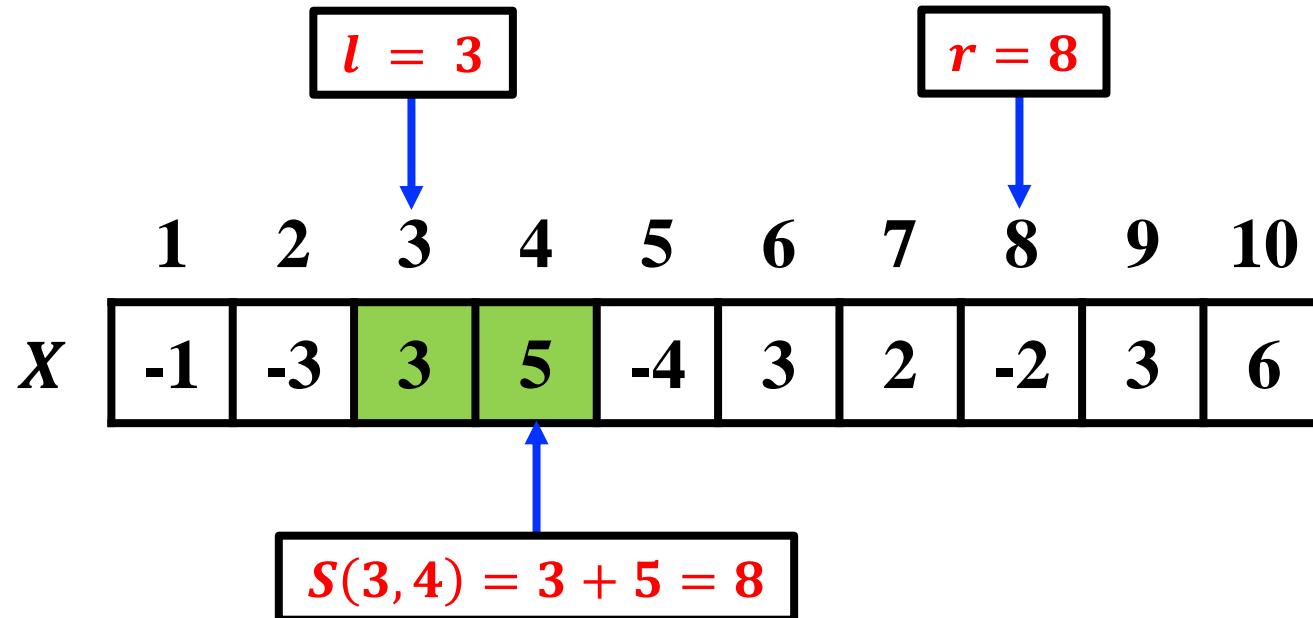
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



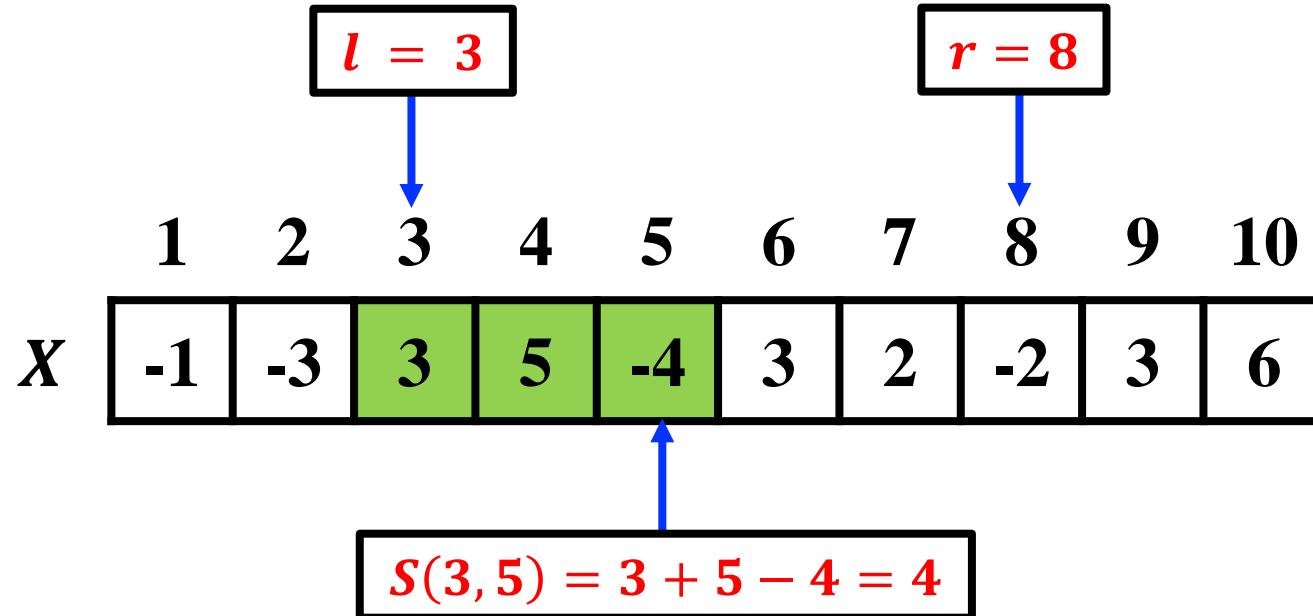
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



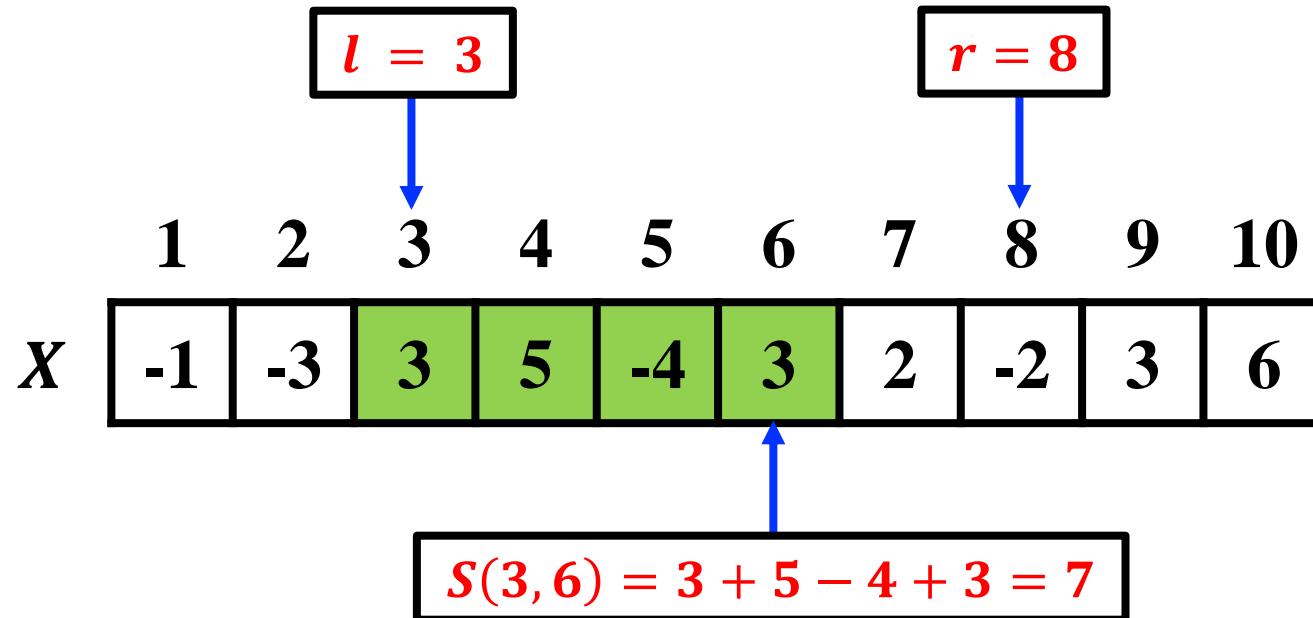
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



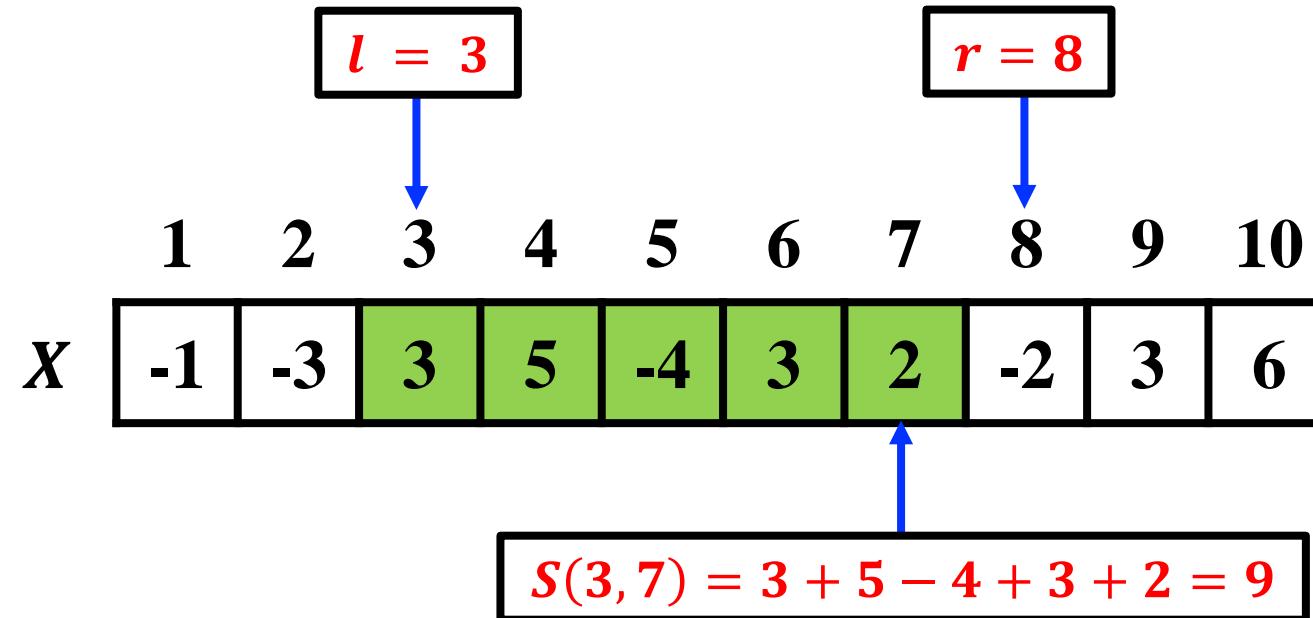
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



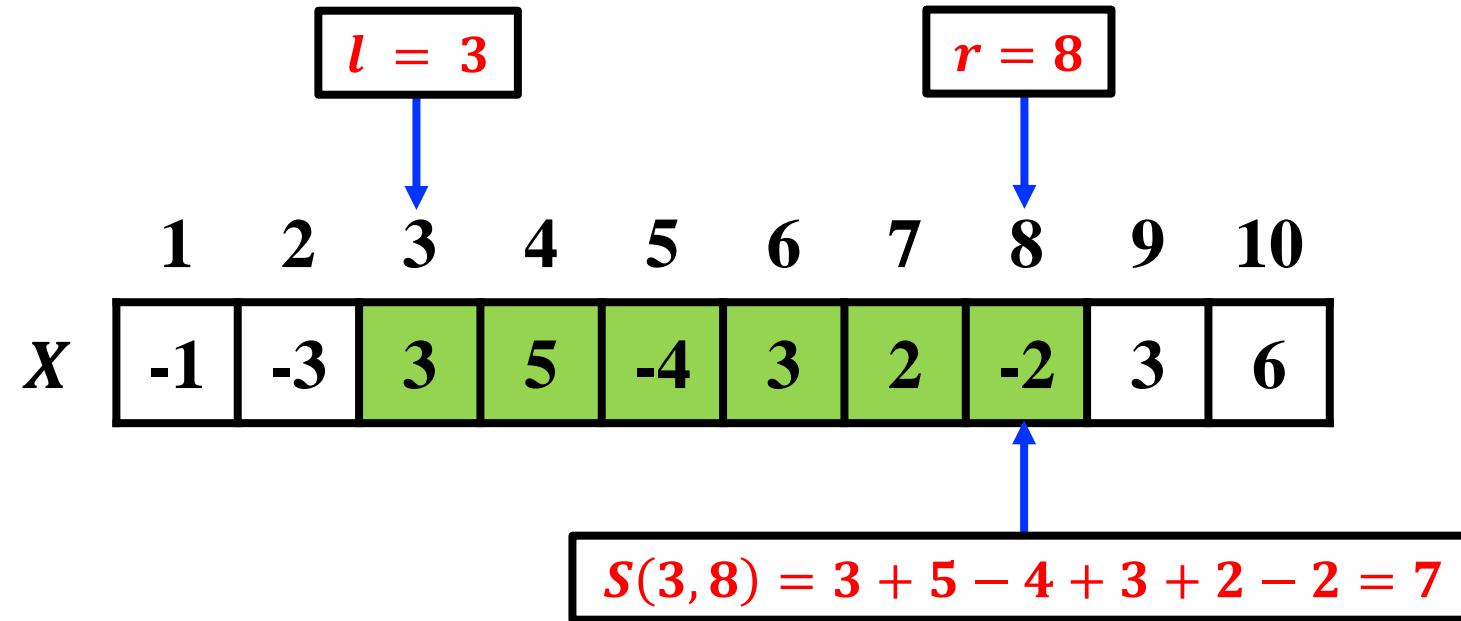
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



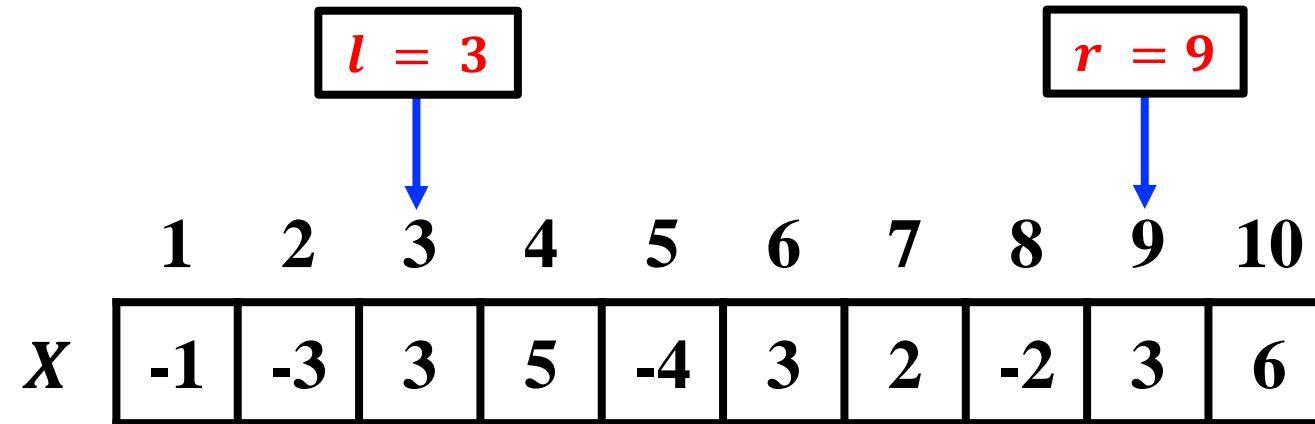
蛮力枚举：算法实例

- $l = 3, r = 8$
- 计算 $S(3, 8)$:



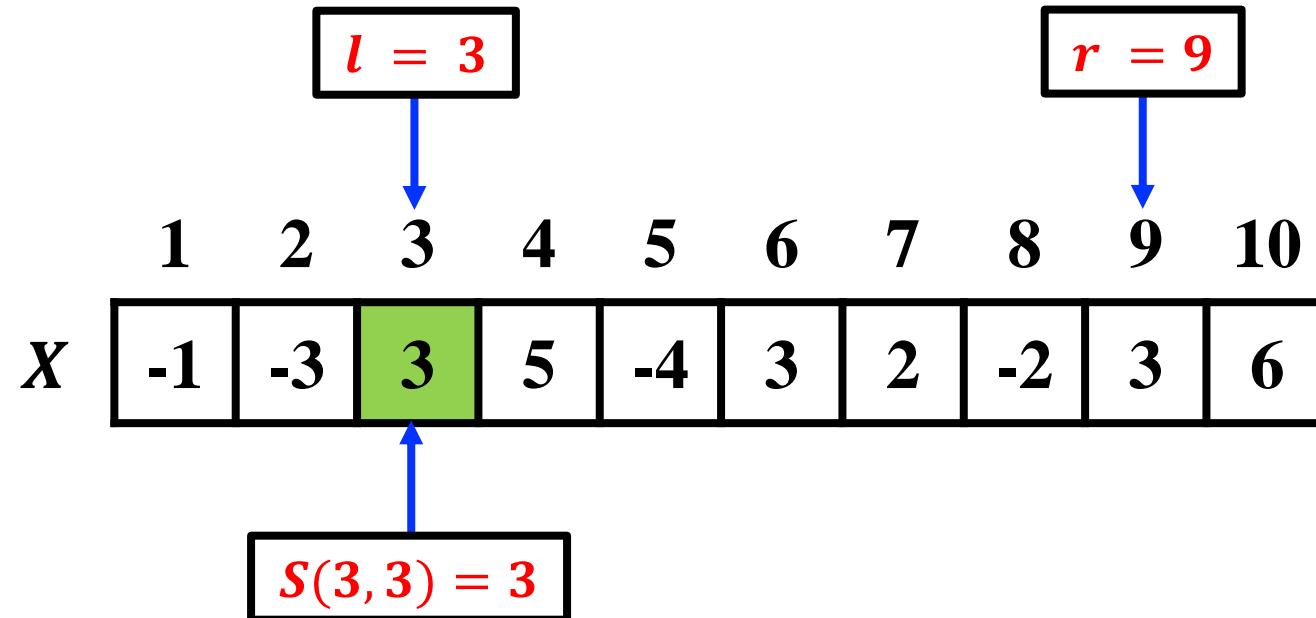
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



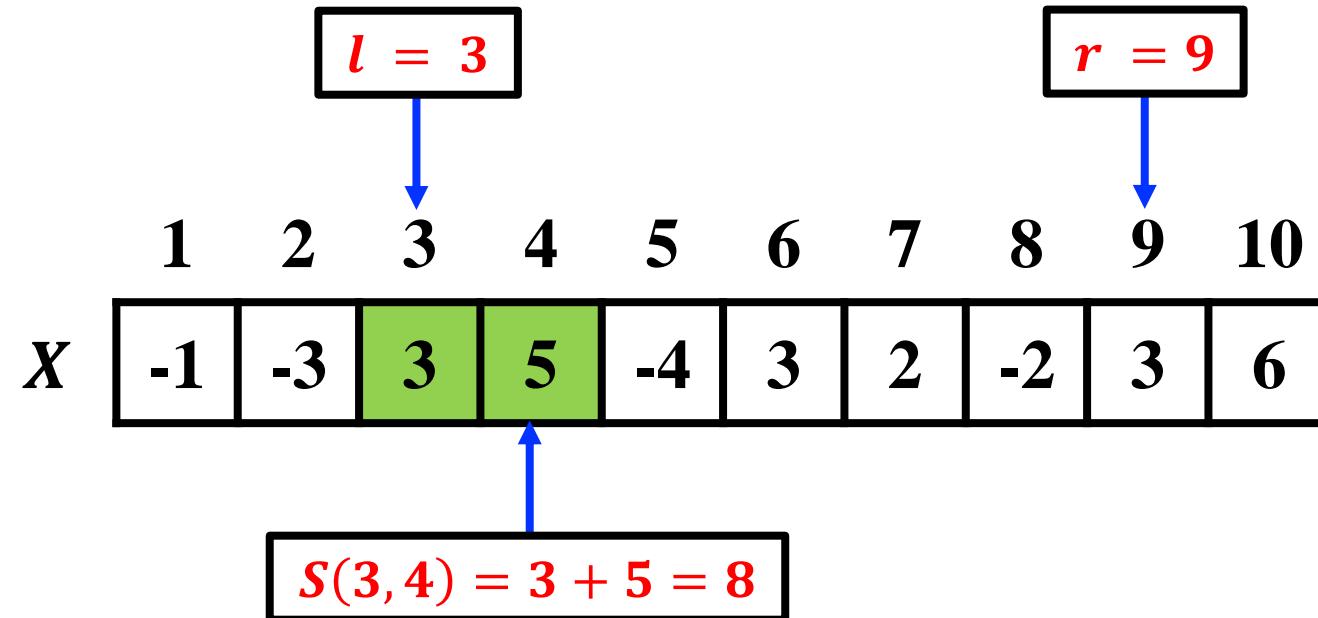
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



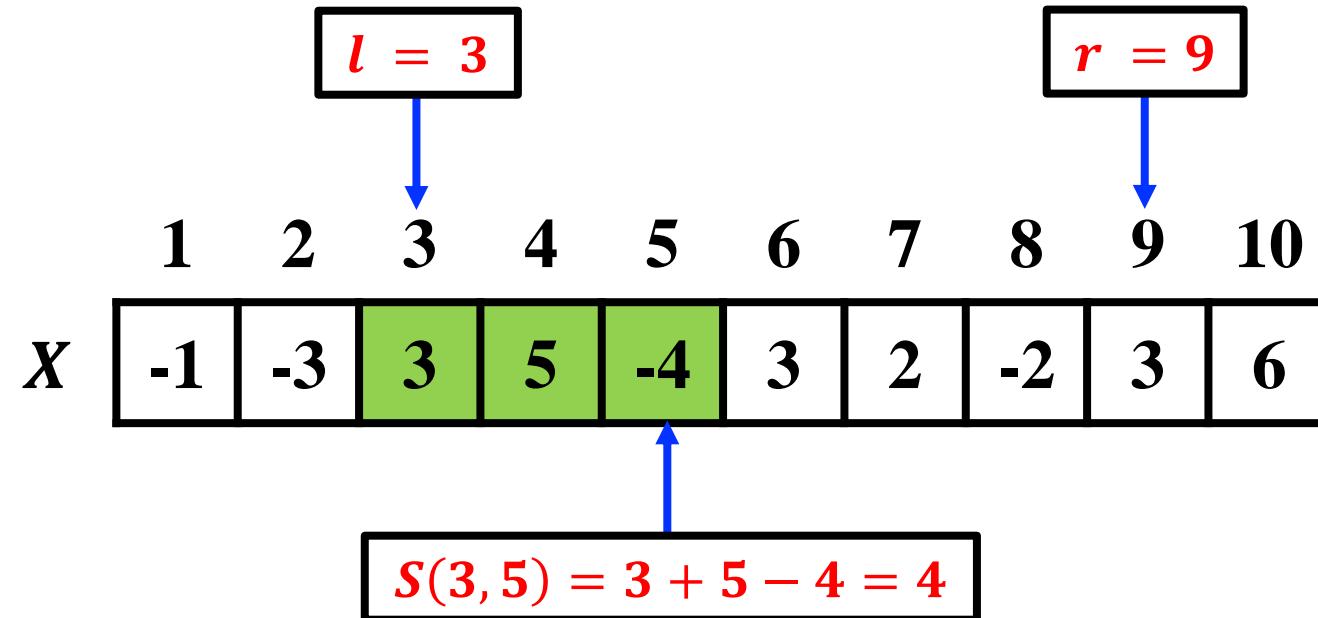
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



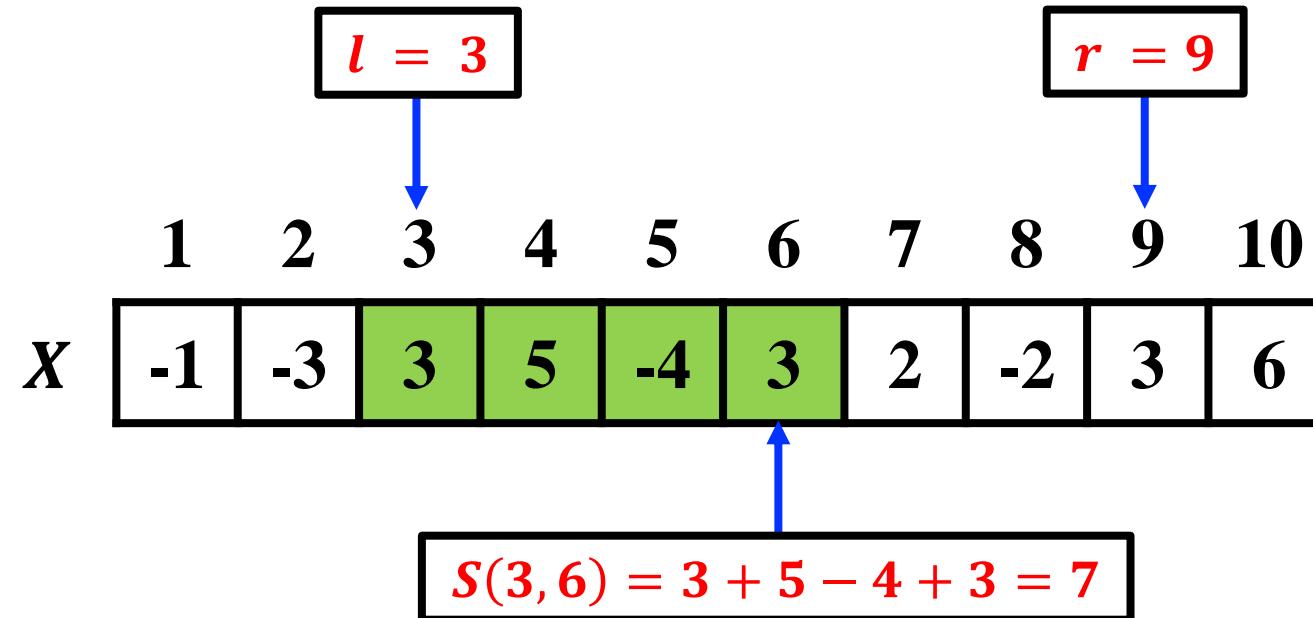
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



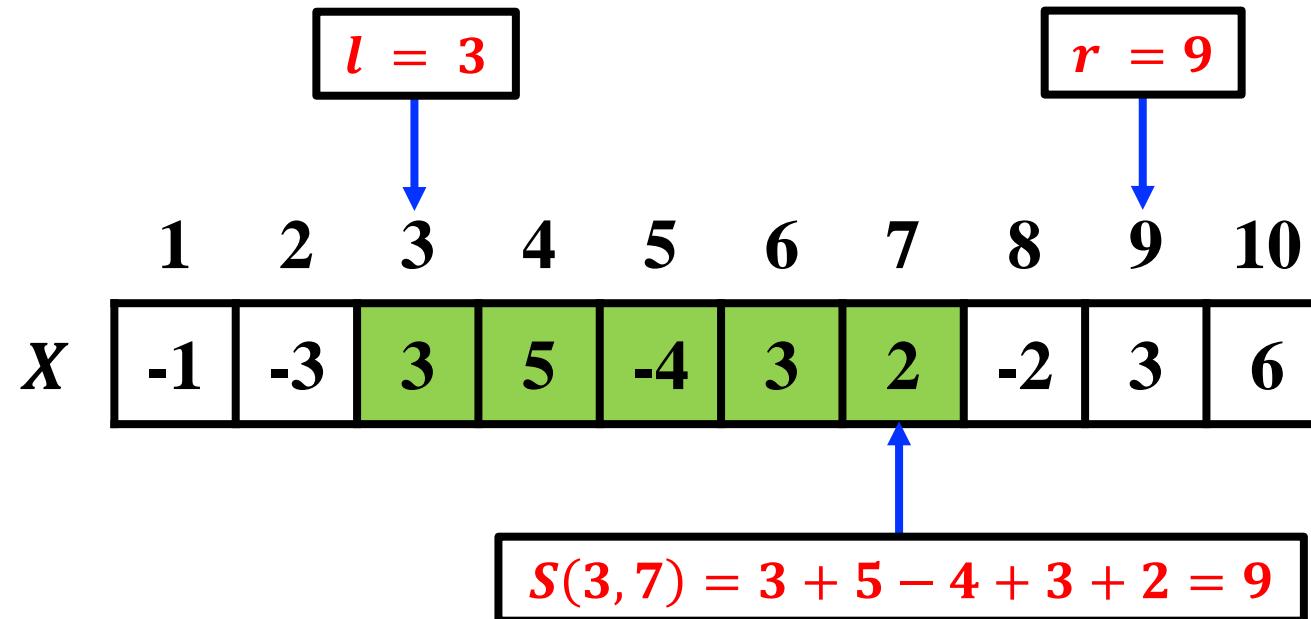
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



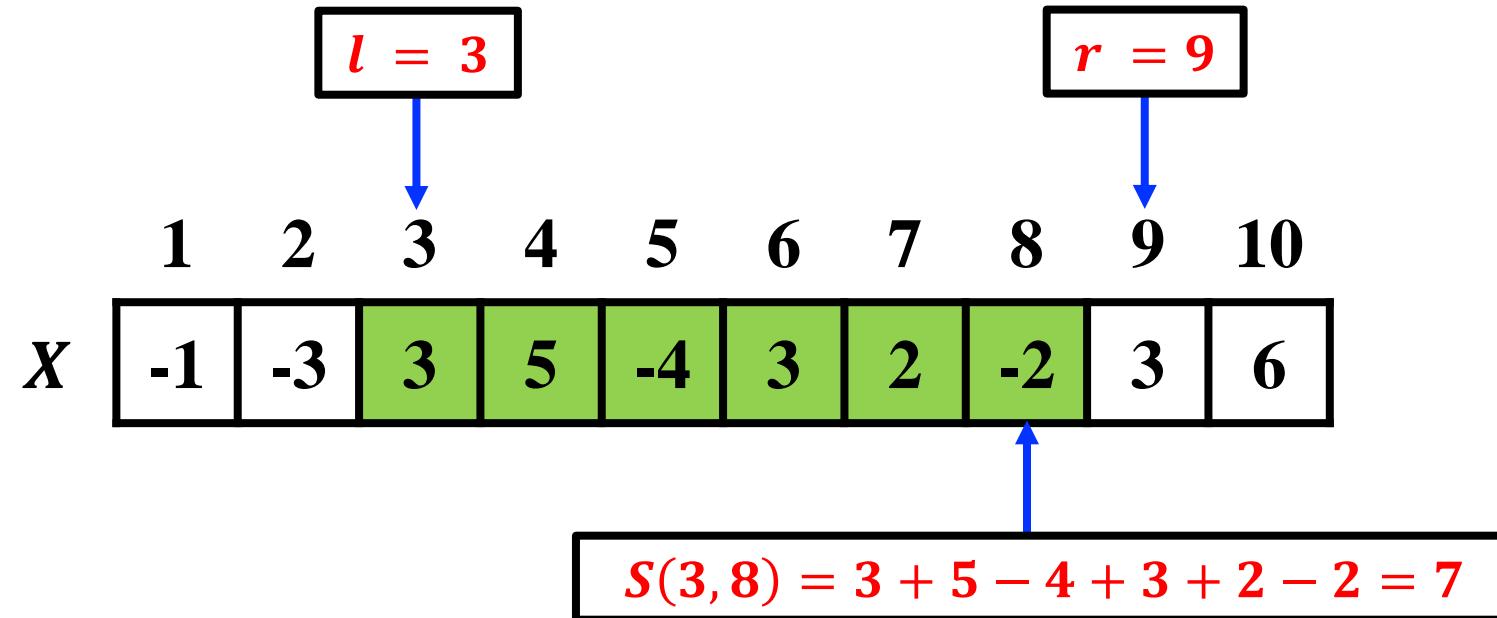
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



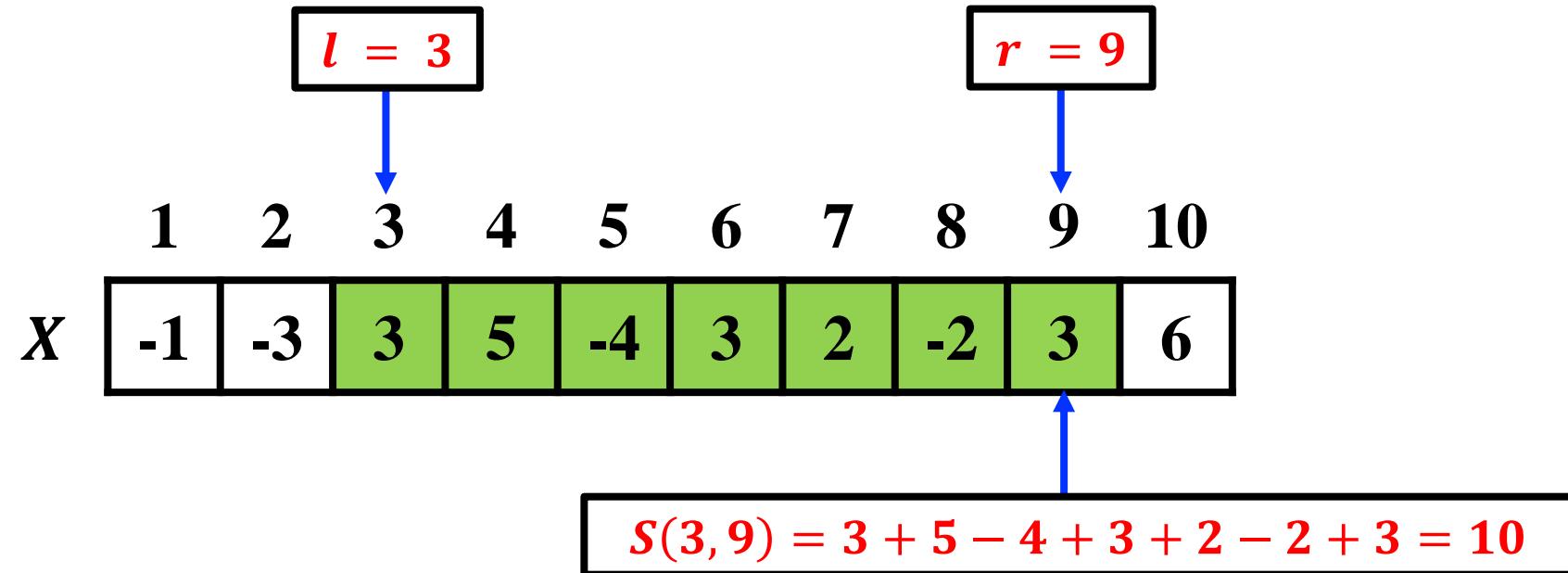
蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



蛮力枚举：算法实例

- $l = 3, r = 9$
- 计算 $S(3, 9)$:



蛮力枚举：伪代码

- 枚举 $n + C_n^2$ 种可能的区间 $[l, r]$ ($l \leq r$)，求解最大子数组之和 S_{max}

输入：数组 $X[1..n]$

输出：最大子数组之和 S_{max}

```
 $S_{max} \leftarrow -\infty$ 
for  $l \leftarrow 1$  to  $n$  do
    for  $r \leftarrow l$  to  $n$  do
         $S(l, r) \leftarrow 0$ 
        for  $i \leftarrow l$  to  $r$  do
            |  $S(l, r) \leftarrow S(l, r) + X[i]$ 
        end
         $S_{max} \leftarrow \max\{S_{max}, S(l, r)\}$ 
    end
end
return  $S_{max}$ 
```

枚举所有的起始下标 l

蛮力枚举：伪代码

- 枚举 $n + C_n^2$ 种可能的区间 $[l, r]$ ($l \leq r$)，求解最大子数组之和 S_{max}

输入：数组 $X[1..n]$

输出：最大子数组之和 S_{max}

```
 $S_{max} \leftarrow -\infty$ 
for  $l \leftarrow 1$  to  $n$  do
    for  $r \leftarrow l$  to  $n$  do
         $S(l, r) \leftarrow 0$ 
        for  $i \leftarrow l$  to  $r$  do
             $S(l, r) \leftarrow S(l, r) + X[i]$ 
        end
         $S_{max} \leftarrow \max\{S_{max}, S(l, r)\}$ 
    end
end
return  $S_{max}$ 
```

枚举所有的终止下标 r

蛮力枚举：伪代码

- 枚举 $n + C_n^2$ 种可能的区间 $[l, r]$ ($l \leq r$)，求解最大子数组之和 S_{max}

输入：数组 $X[1..n]$

输出：最大子数组之和 S_{max}

$S_{max} \leftarrow -\infty$

for $l \leftarrow 1$ to n **do**

for $r \leftarrow l$ to n **do**

$S(l, r) \leftarrow 0$

for $i \leftarrow l$ to r **do**

$S(l, r) \leftarrow S(l, r) + X[i]$

end

$S_{max} \leftarrow \max\{S_{max}, S(l, r)\}$

end

end
return S_{max}

迭代计算 $S(l, r)$ 并更新 S_{max}

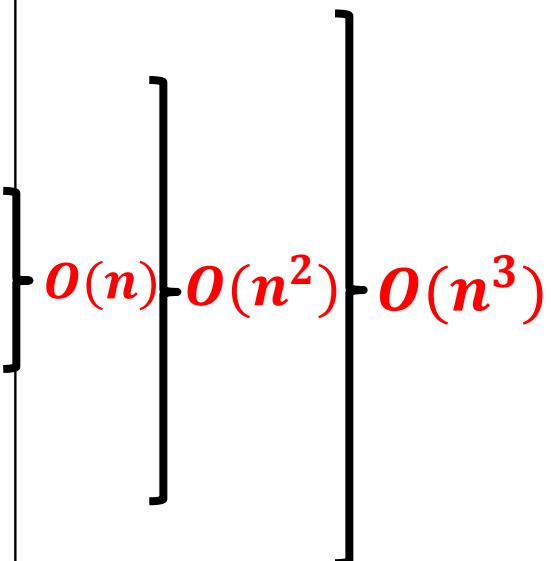
蛮力枚举：伪代码

- 枚举 $n + C_n^2$ 种可能的区间 $[l, r]$ ($l \leq r$)，求解最大子数组之和 S_{max}

输入：数组 $X[1..n]$

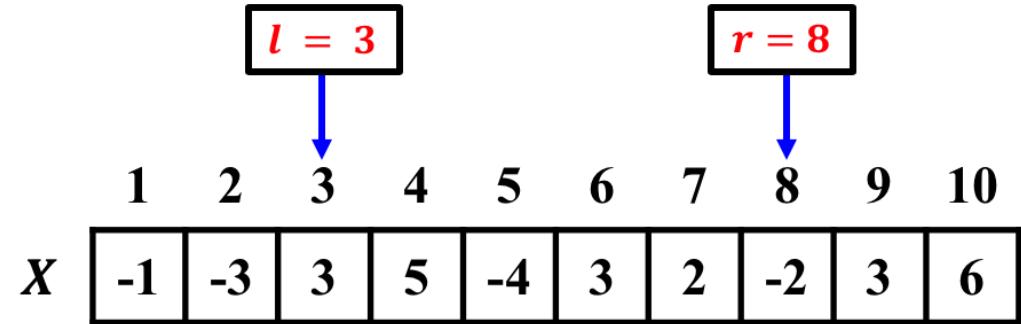
输出：最大子数组之和 S_{max}

```
 $S_{max} \leftarrow -\infty$ 
for  $l \leftarrow 1$  to  $n$  do
    for  $r \leftarrow l$  to  $n$  do
         $S(l, r) \leftarrow 0$ 
        for  $i \leftarrow l$  to  $r$  do
            |  $S(l, r) \leftarrow S(l, r) + X[i]$ 
        end
         $S_{max} \leftarrow \max\{S_{max}, S(l, r)\}$ 
    end
end
return  $S_{max}$ 
```

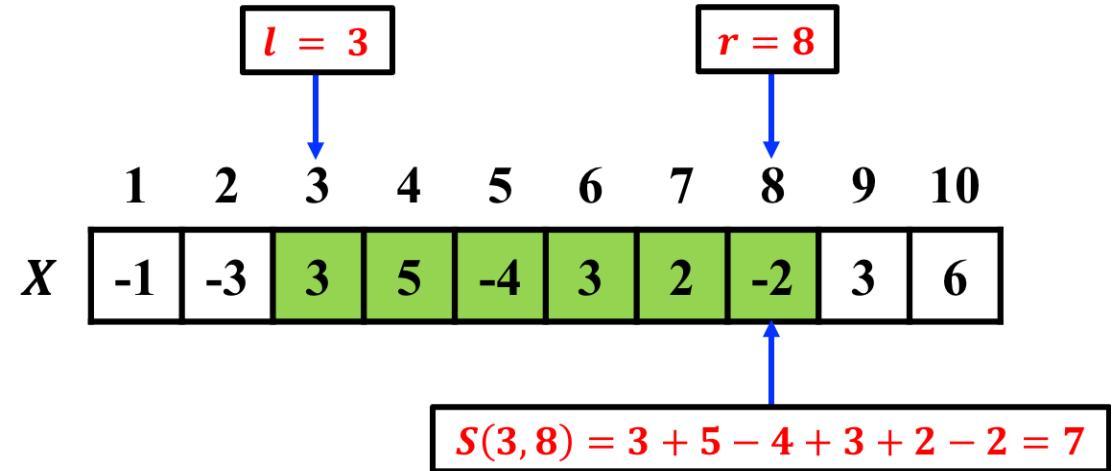
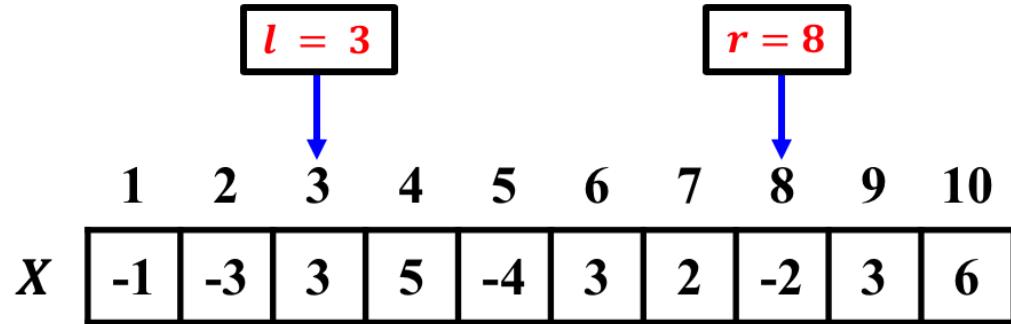


时间复杂度： $O(n^3)$

蛮力枚举：实例分析

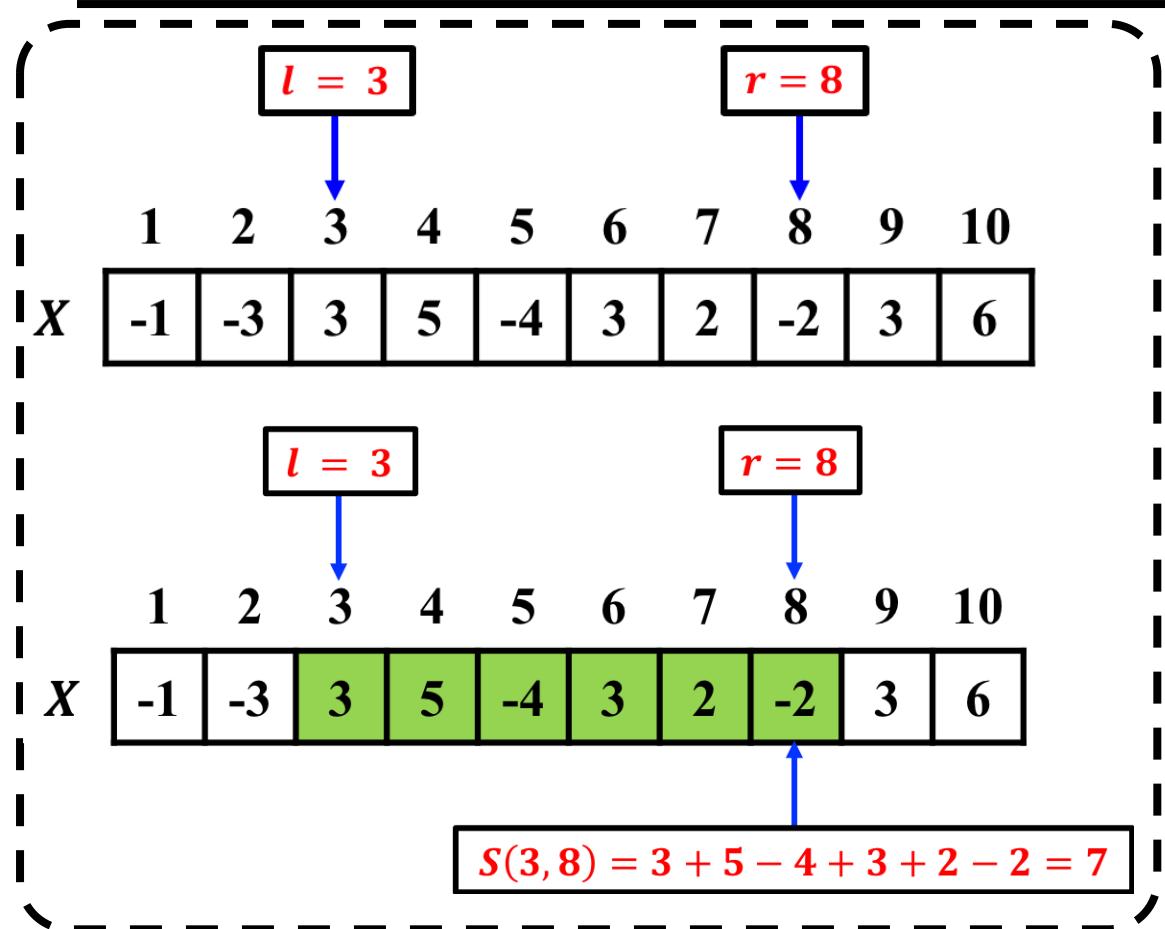


蛮力枚举：实例分析



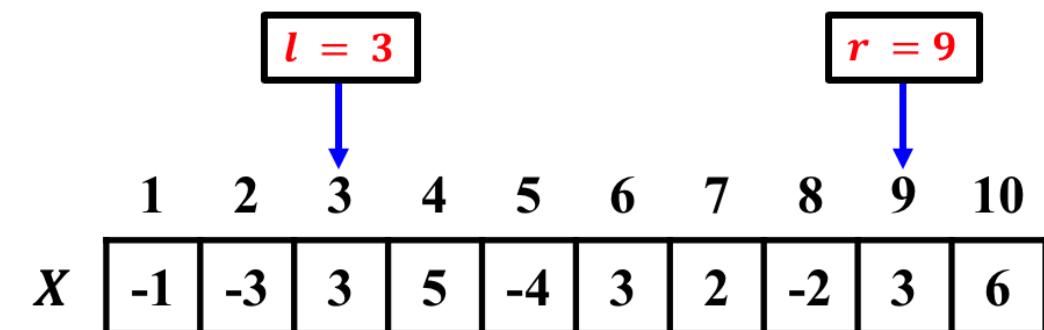
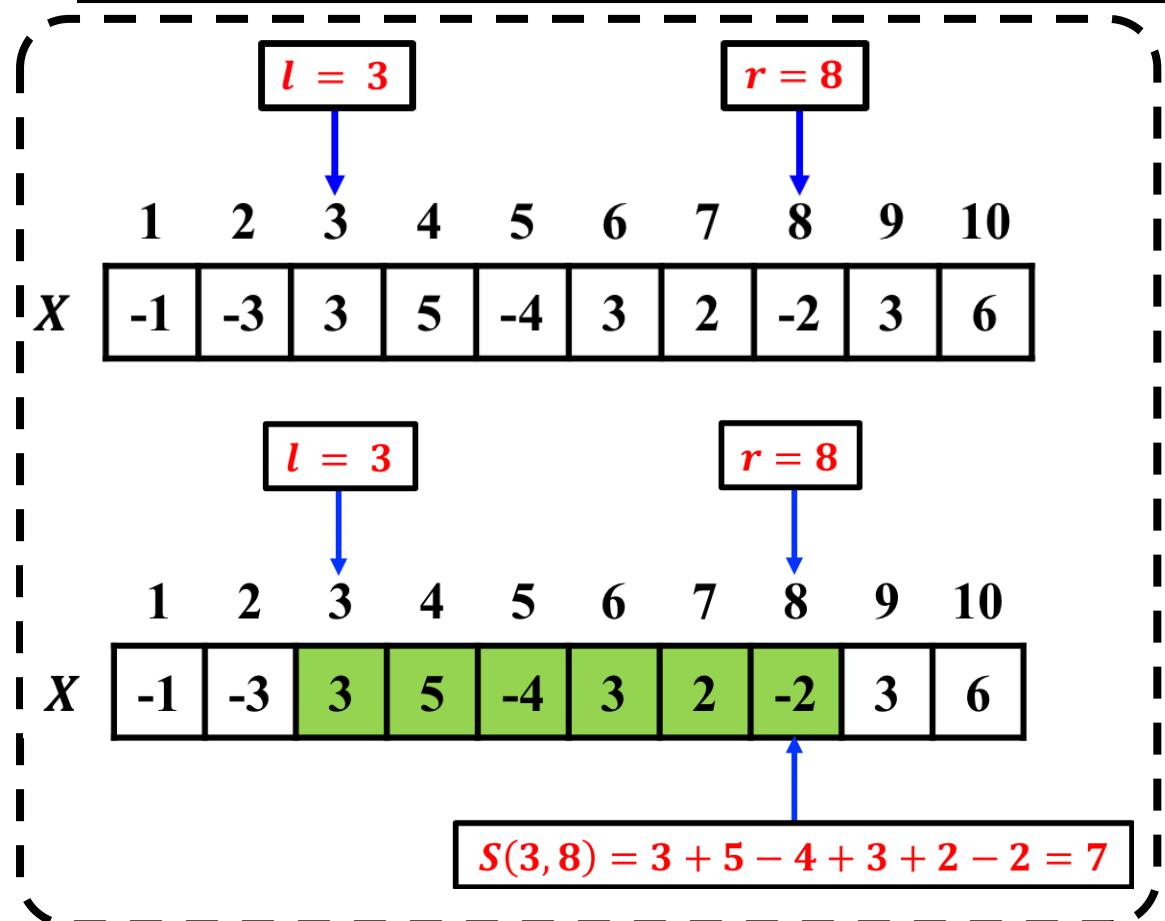
$$S(3, 8) = 3 + 5 - 4 + 3 + 2 - 2 = 7$$

蛮力枚举：实例分析



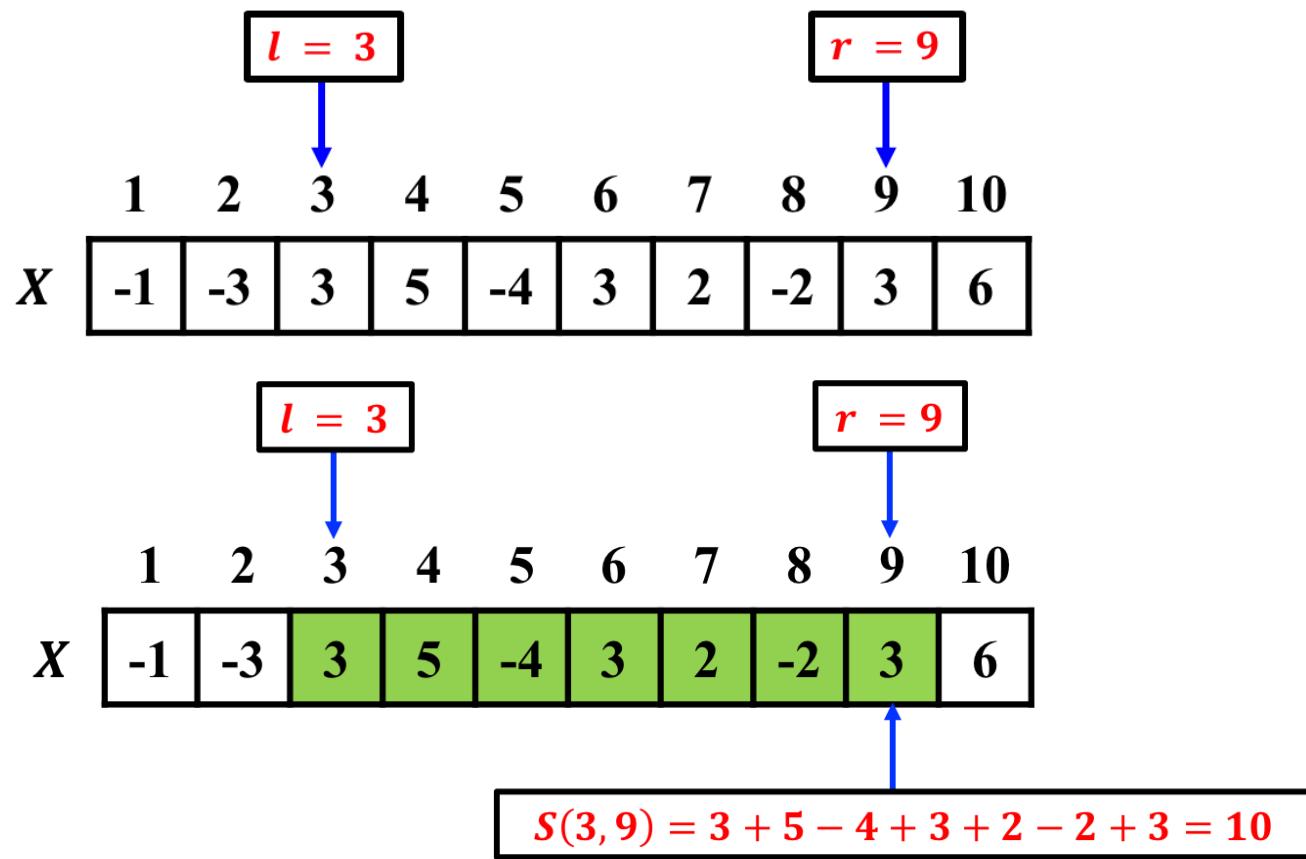
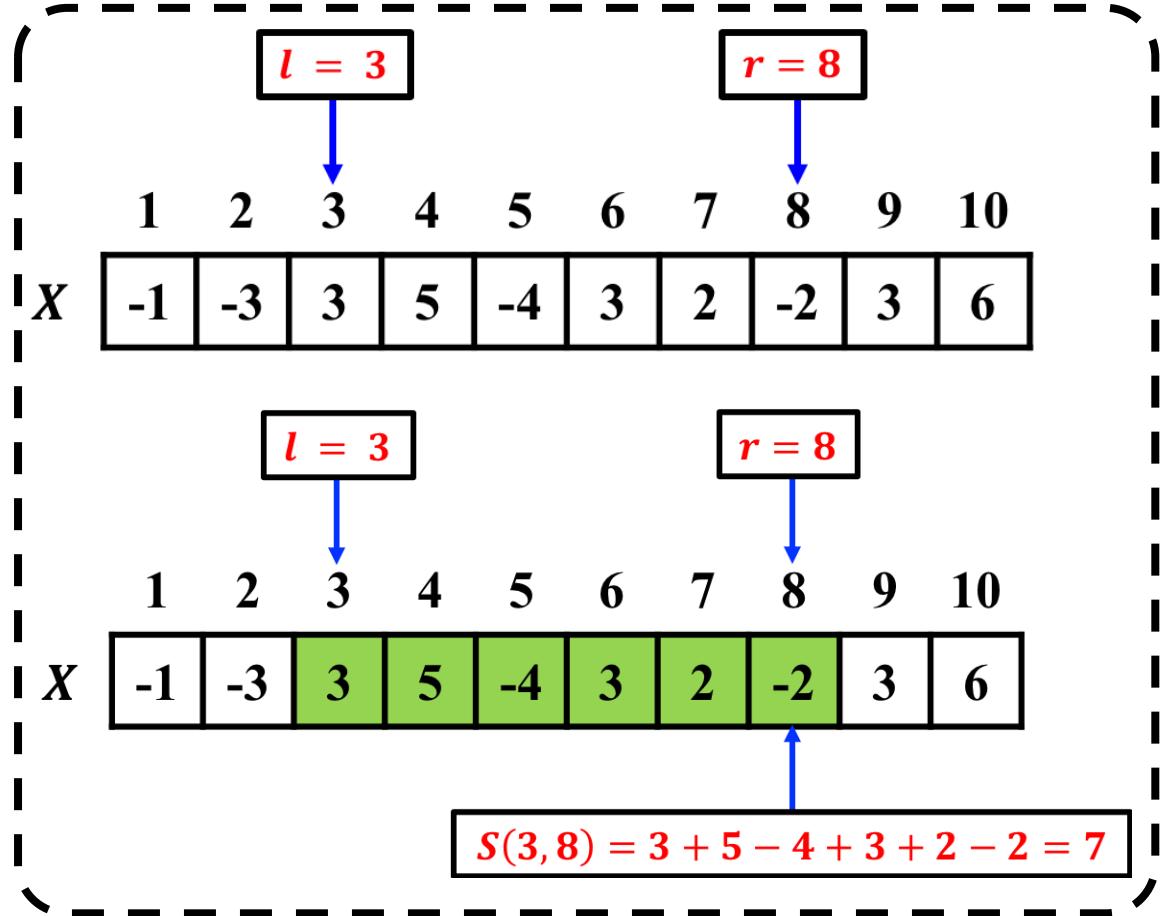
计算 $S(3, 8)$ 循环 6 次

蛮力枚举：实例分析



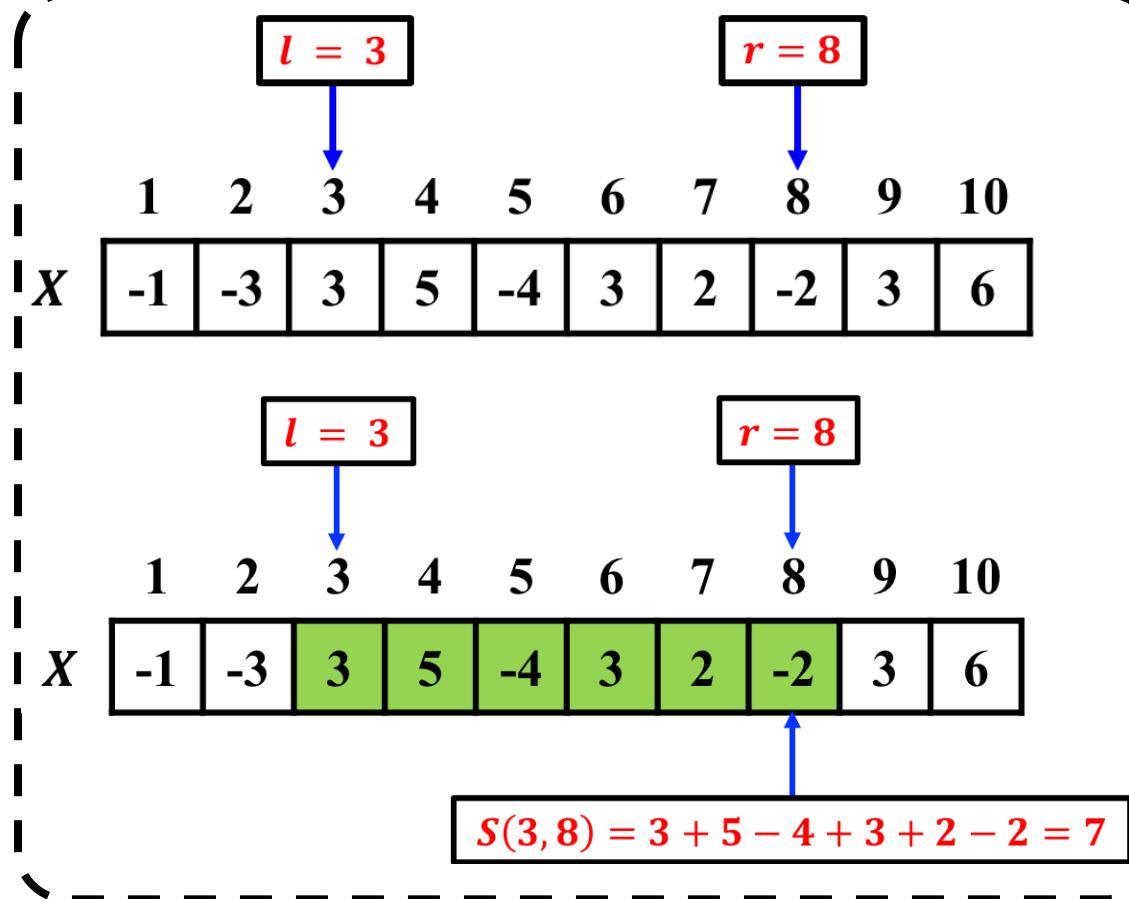
计算 $S(3, 8)$ 循环 6 次

蛮力枚举：实例分析

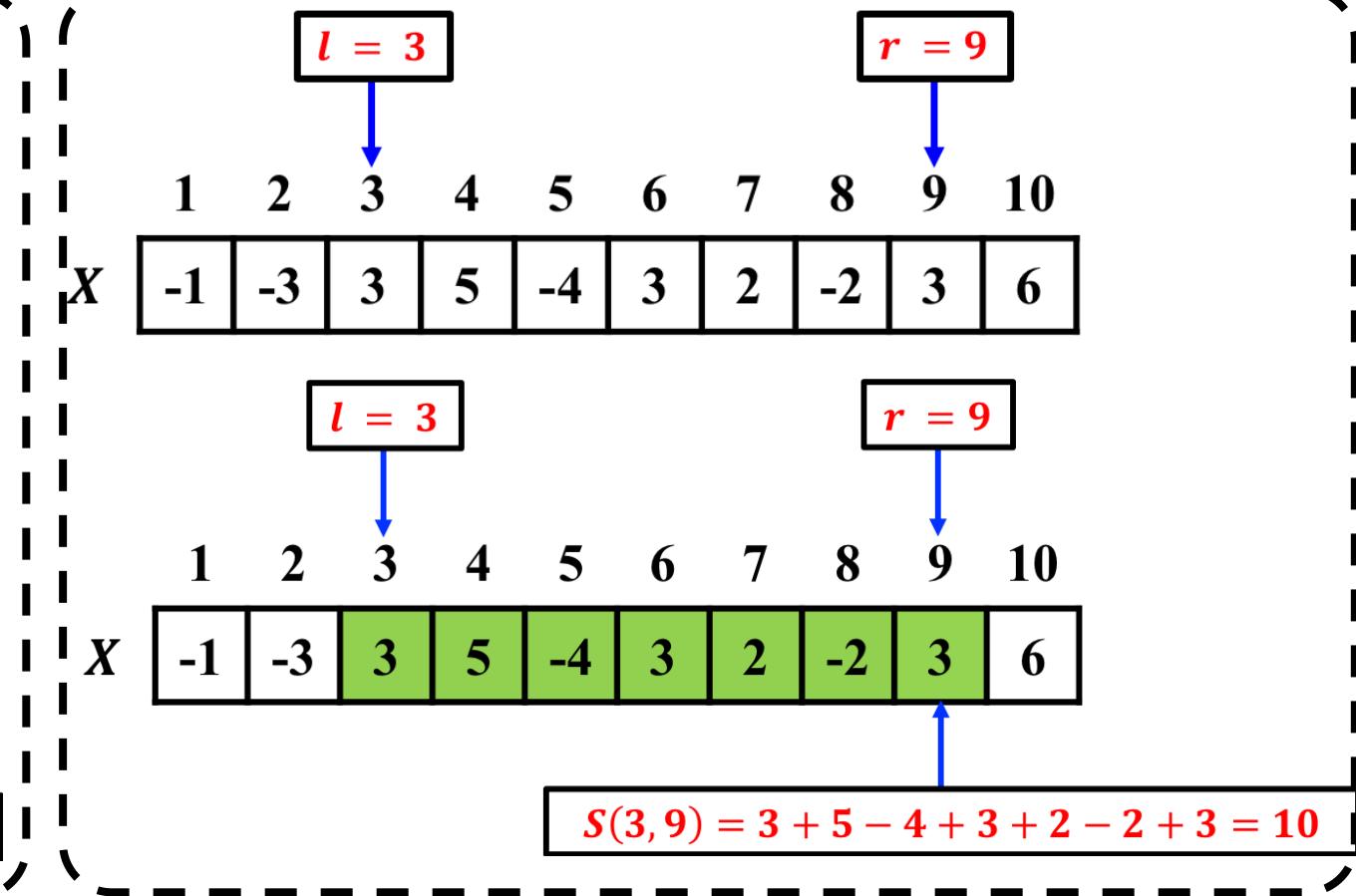


计算 $S(3, 8)$ 循环6次

蛮力枚举：实例分析

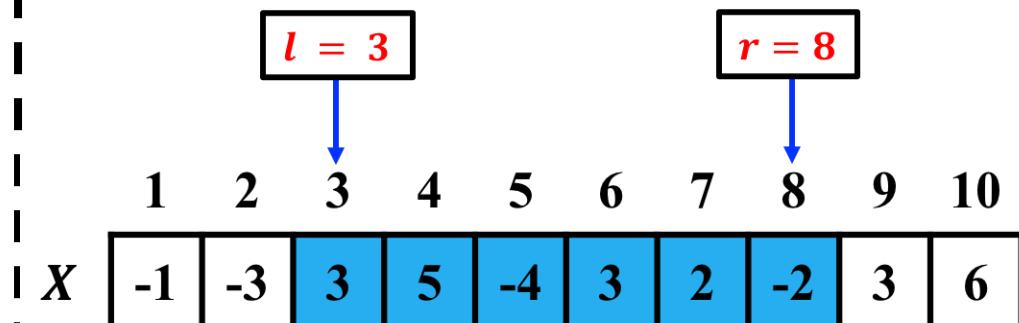
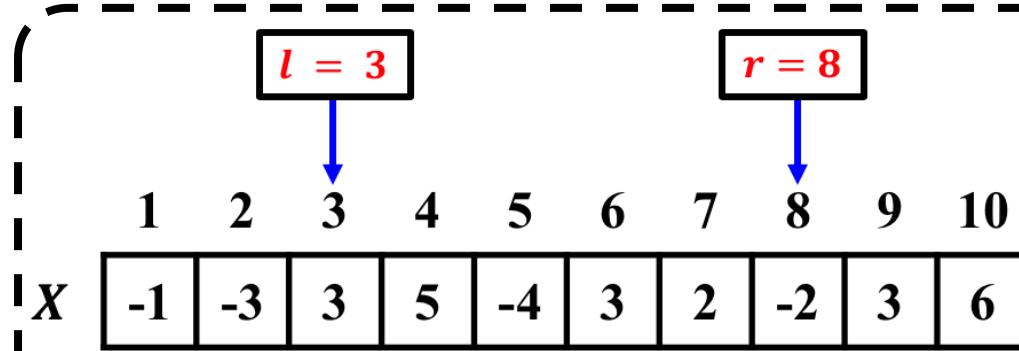


计算 $S(3, 8)$ 循环6次



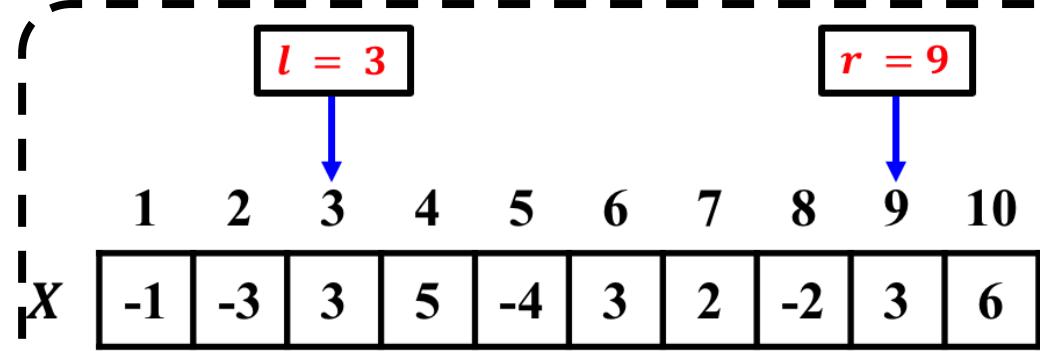
计算 $S(3, 9)$ 循环7次

蛮力枚举：实例分析



$$S(3, 8) = 3 + 5 - 4 + 3 + 2 - 2 = 7$$

计算 $S(3, 8)$ 循环 6 次

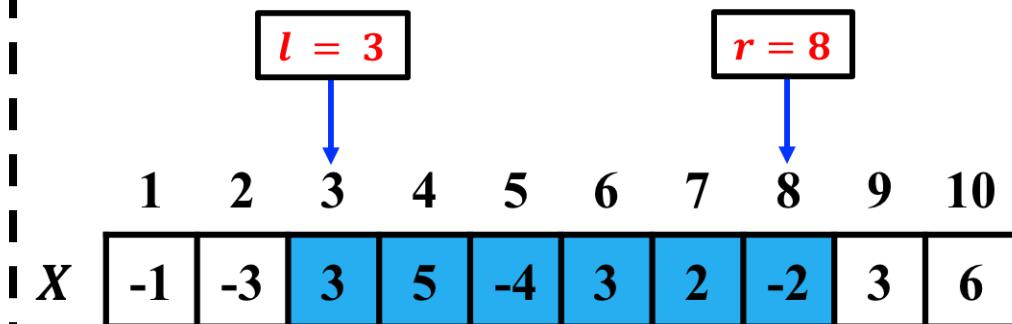
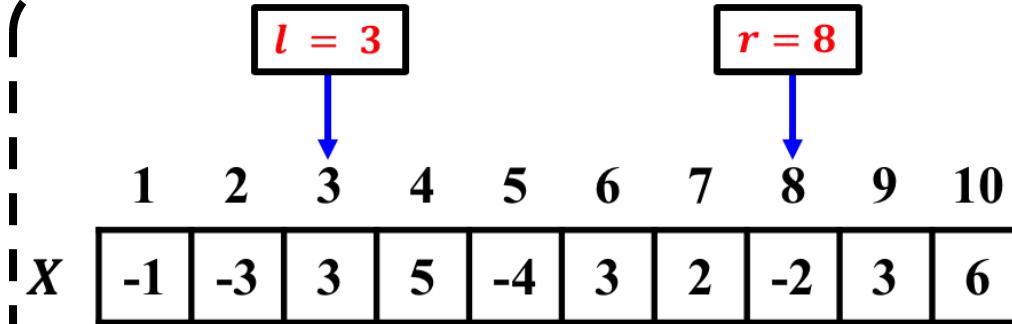


$$S(3, 9) = 3 + 5 - 4 + 3 + 2 - 2 + 3 = 10$$

计算 $S(3, 9)$ 循环 7 次

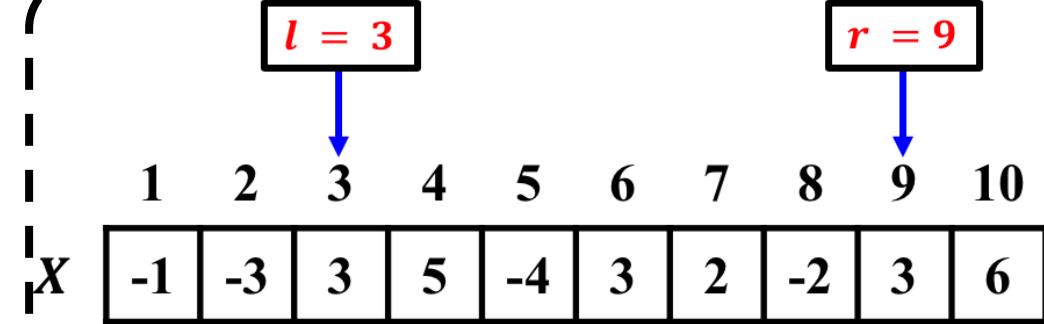
存在重复计算

蛮力枚举：实例分析



$$S(3, 8) = 3 + 5 - 4 + 3 + 2 - 2 = 7$$

计算 $S(3, 8)$ 循环 6 次



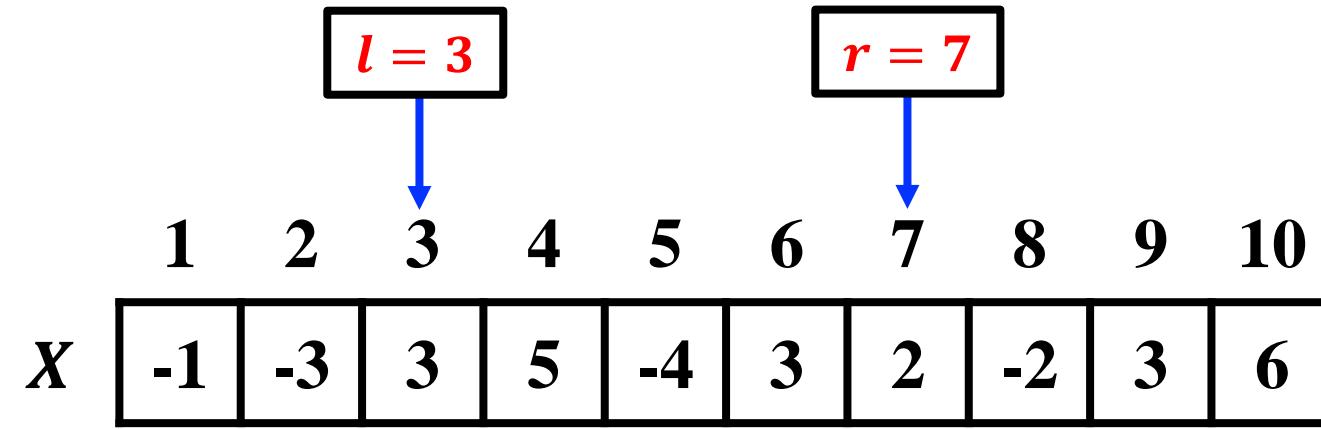
$$S(3, 9) = 3 + 5 - 4 + 3 + 2 - 2 + 3 = 10$$

计算 $S(3, 9)$ 循环 7 次

问题：如何减少重复计算？

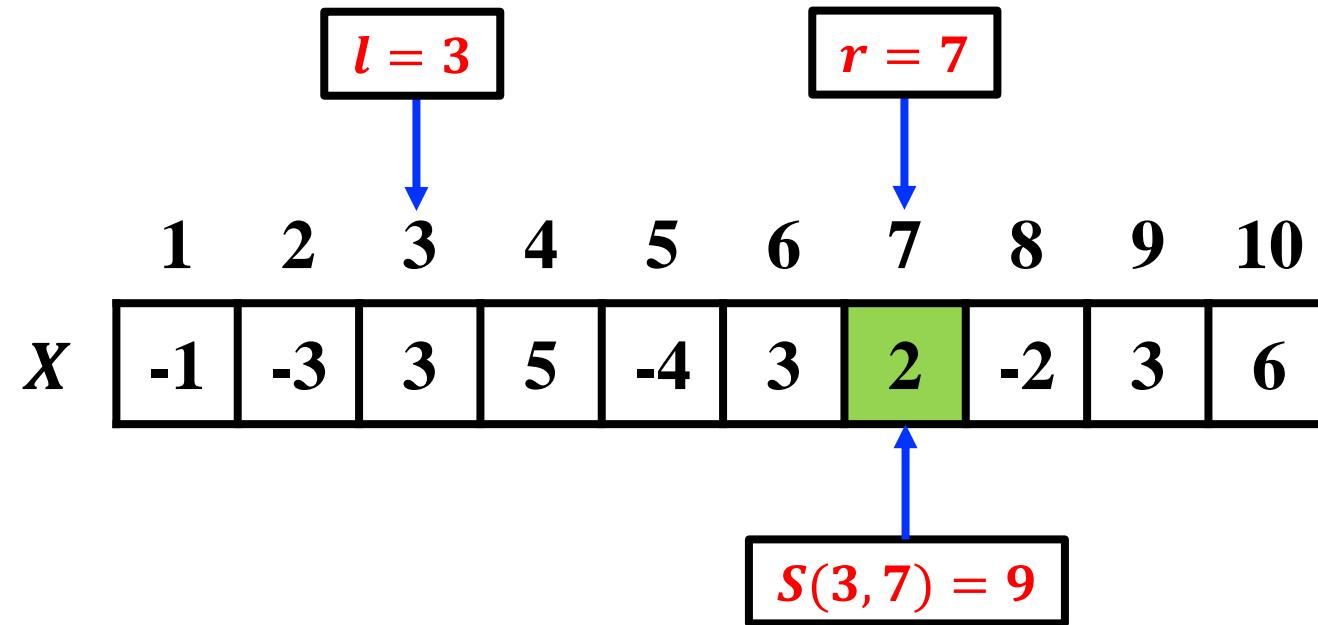
优化枚举：算法实例

- $l = 3, r = 7$



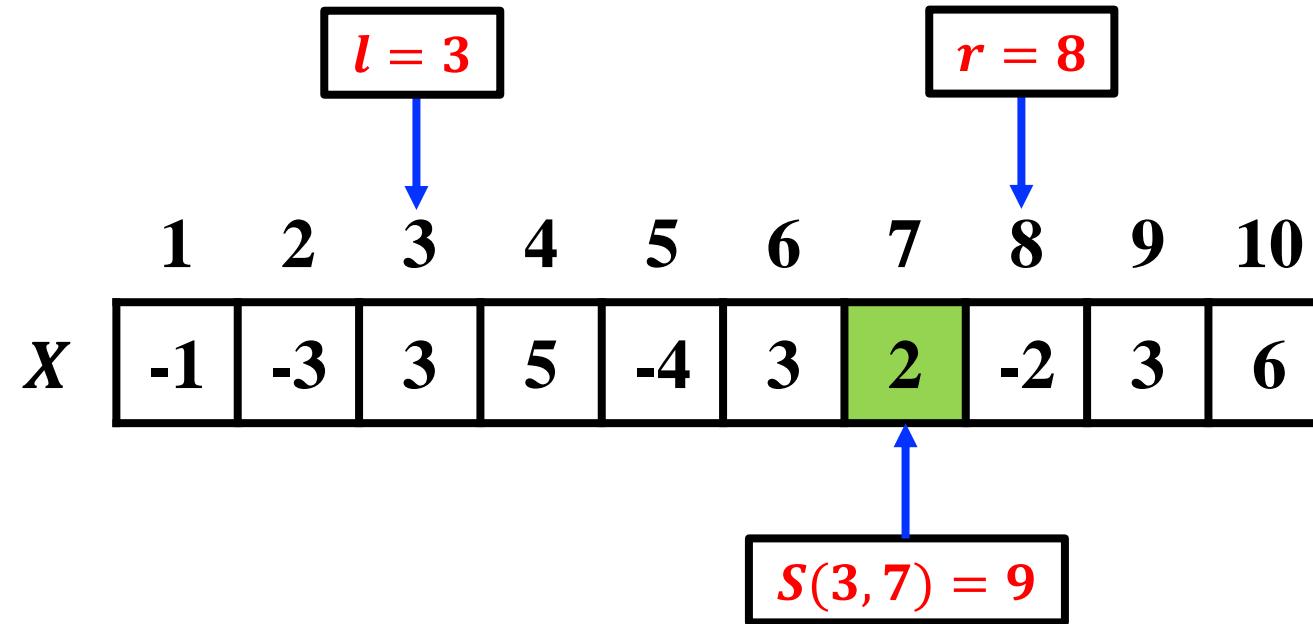
优化枚举：算法实例

- $l = 3, r = 7, S(3, 7) = 9$



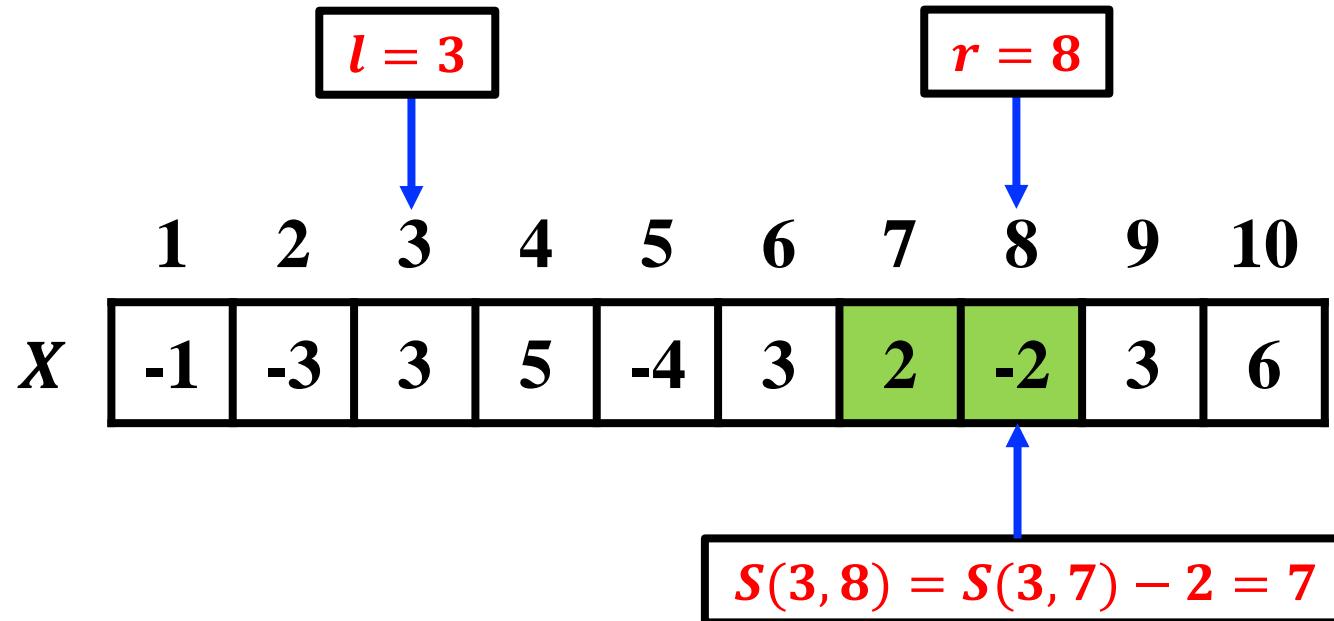
优化枚举：算法实例

- $l = 3, r = 8, S(3, 8) = ?$



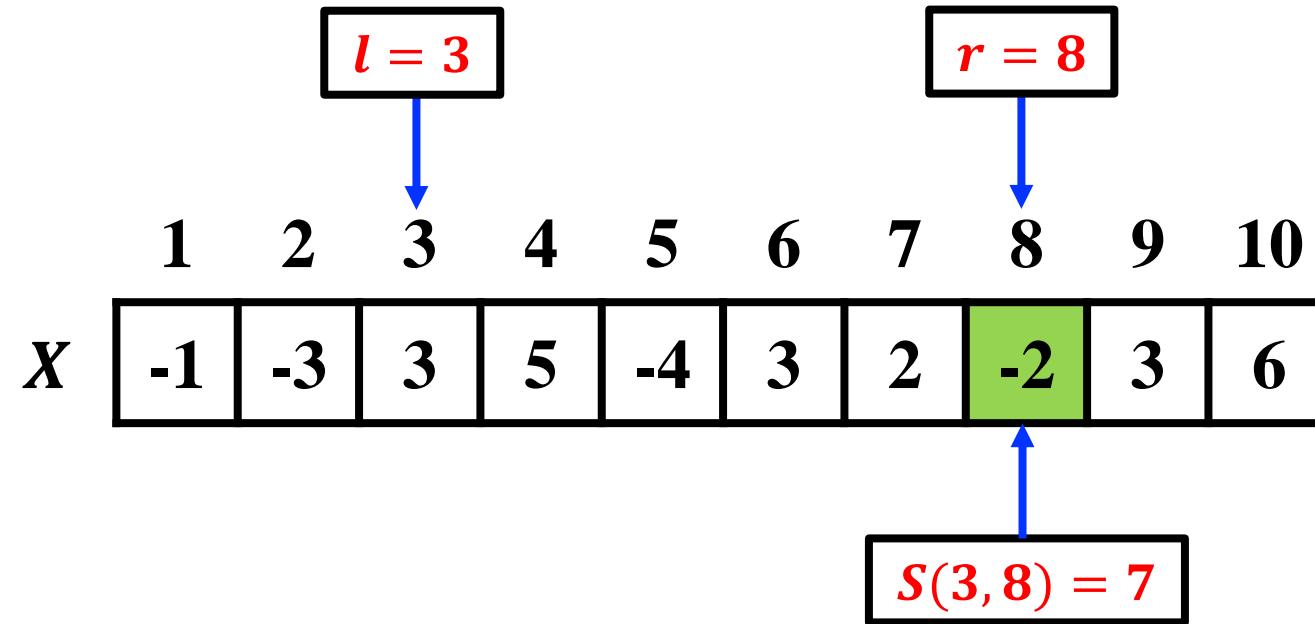
优化枚举：算法实例

- $l = 3, r = 8, S(3, 8) = 7$



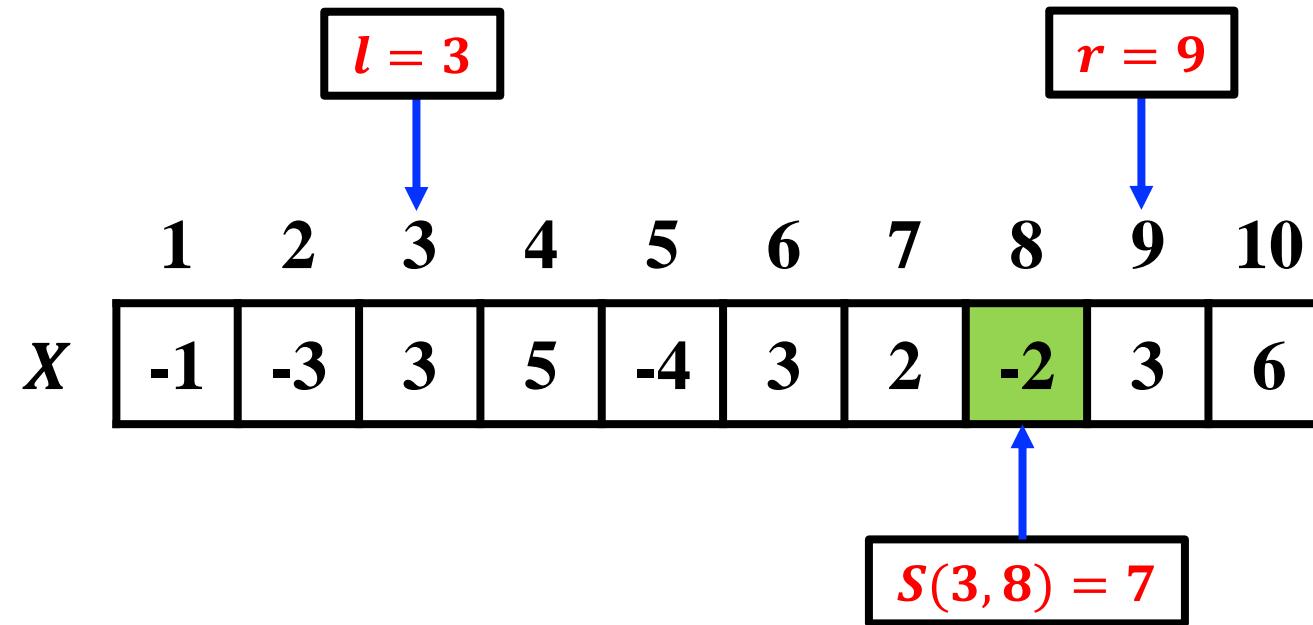
优化枚举：算法实例

- $l = 3, r = 8, S(3, 8) = 7$



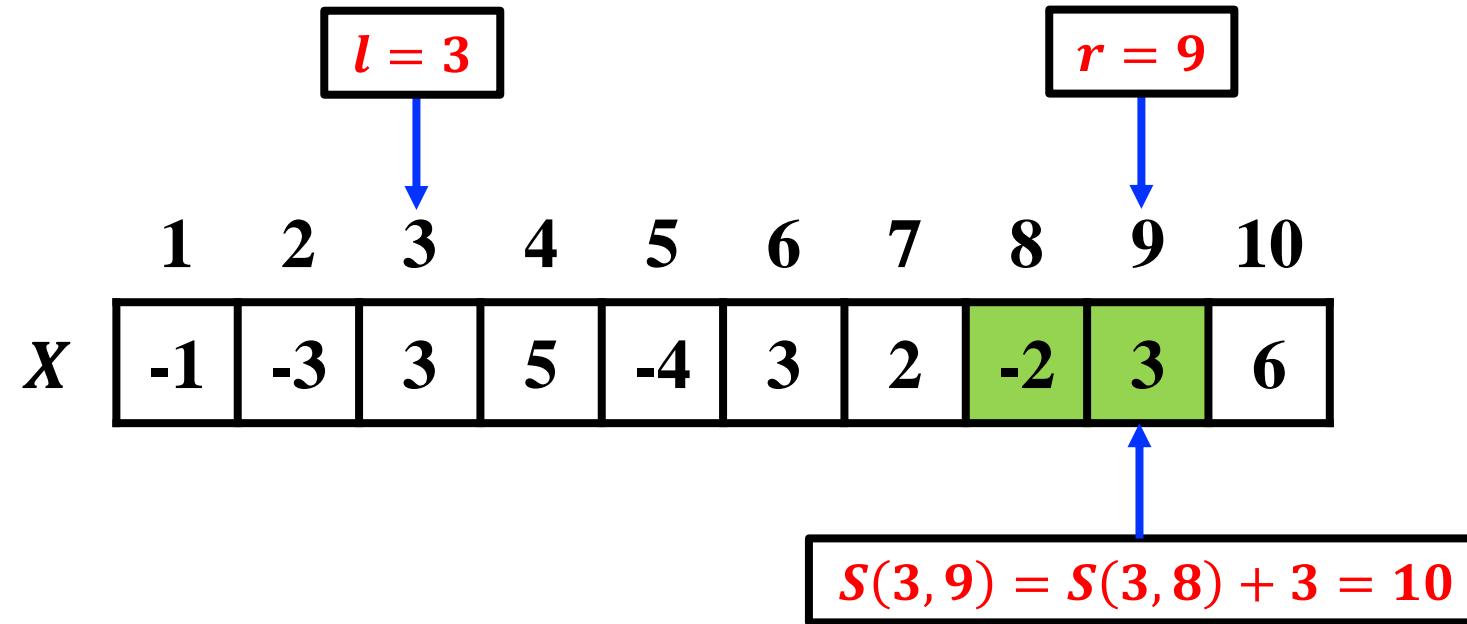
优化枚举：算法实例

- $l = 3, r = 9, S(3, 9) = ?$

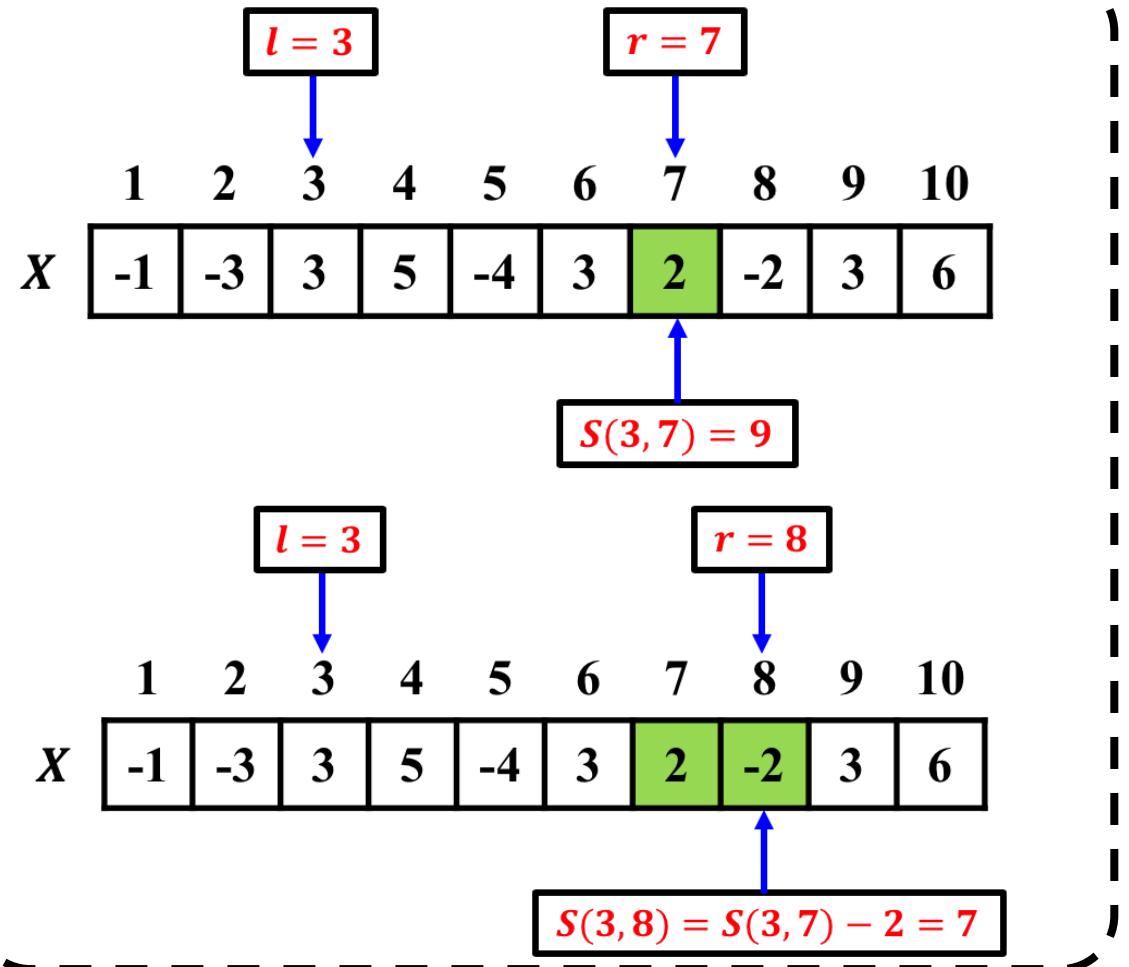


优化枚举：算法实例

- $l = 3, r = 9, S(3, 9) = ?$

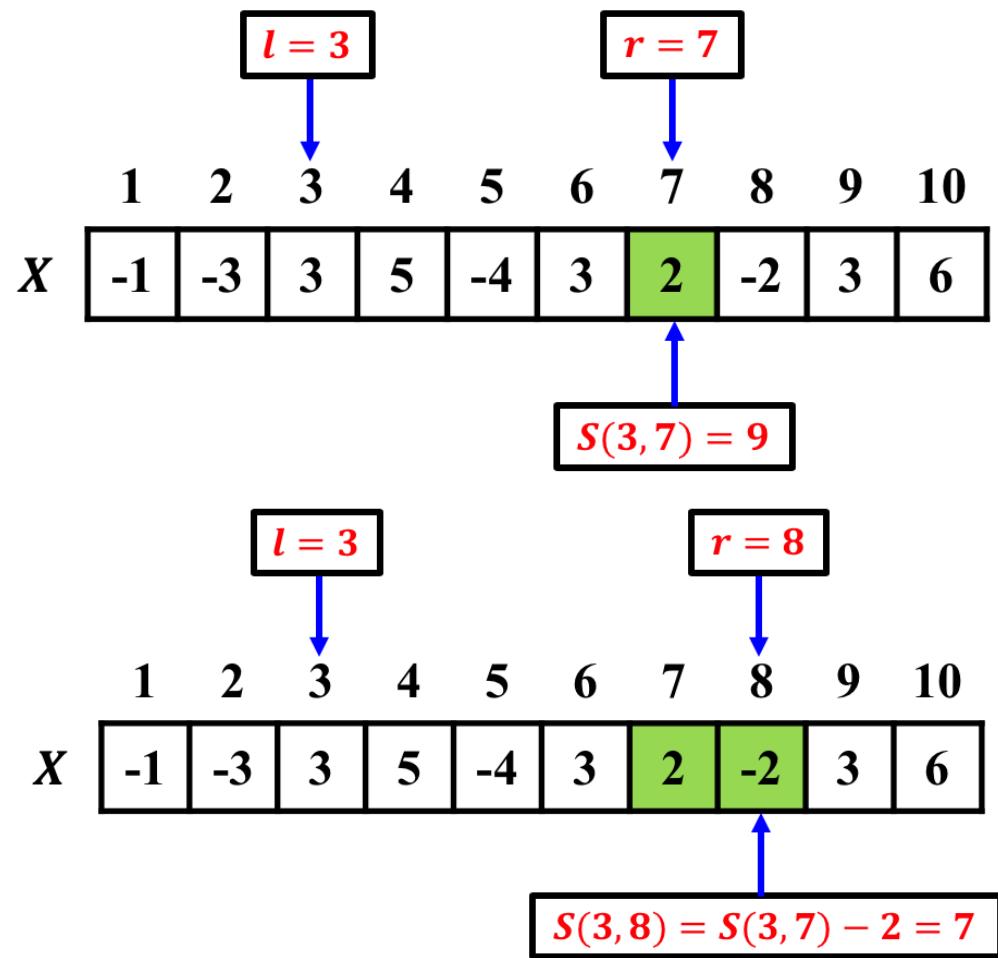


优化枚举：实例分析

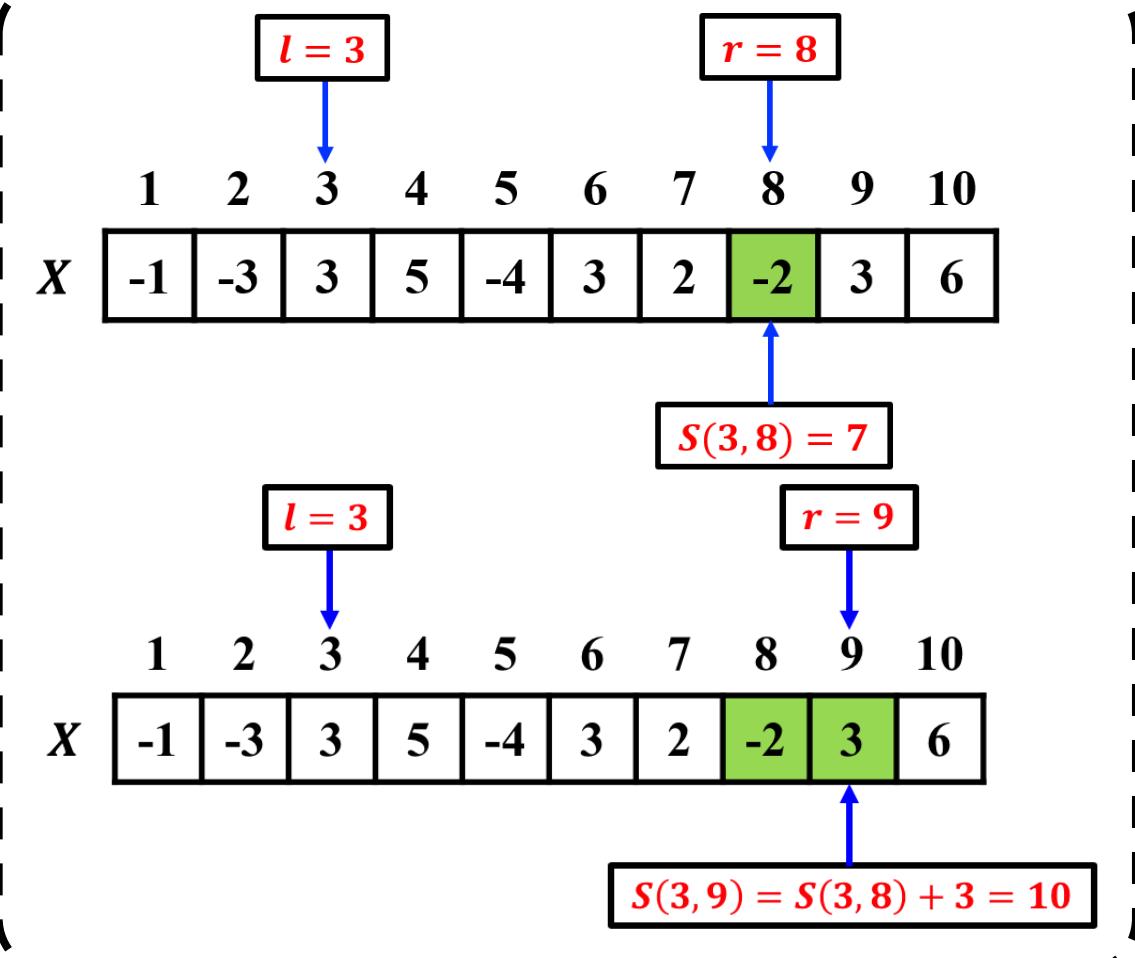


计算 $S(3, 8)$ 操作1次

优化枚举：实例分析

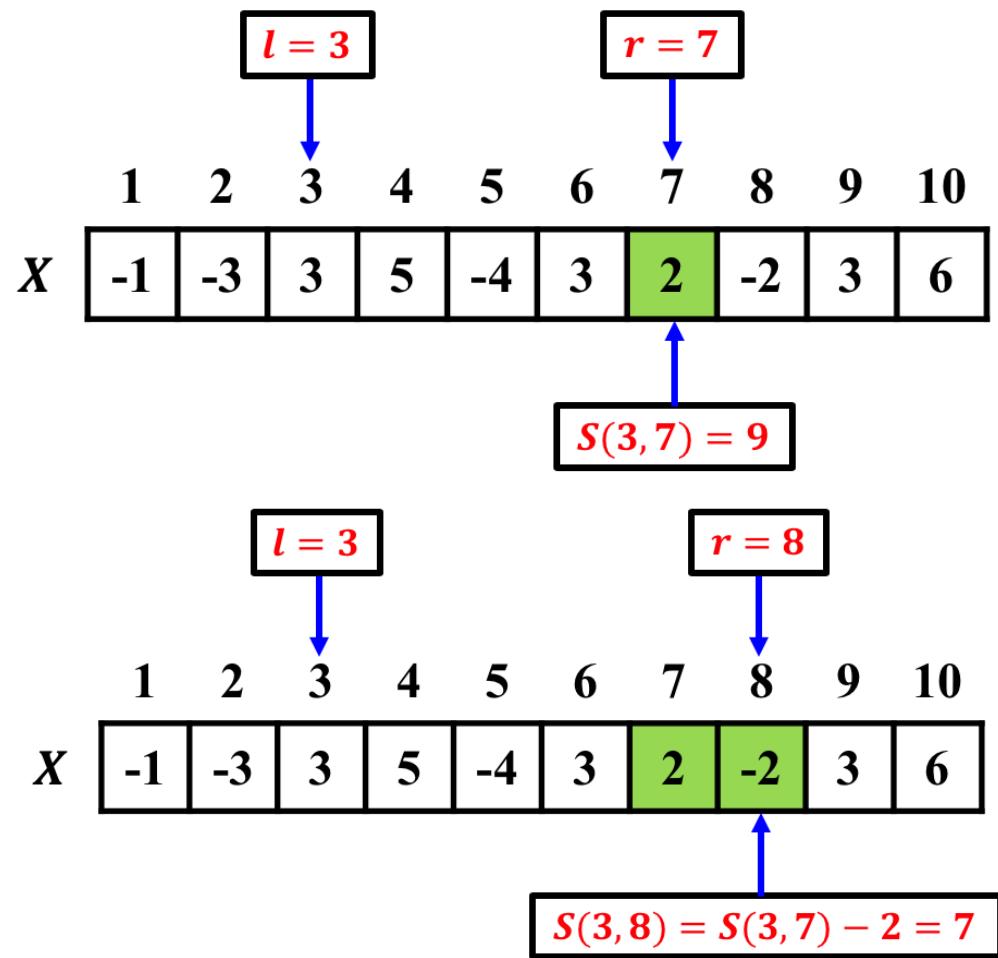


计算 $S(3, 8)$ 操作1次

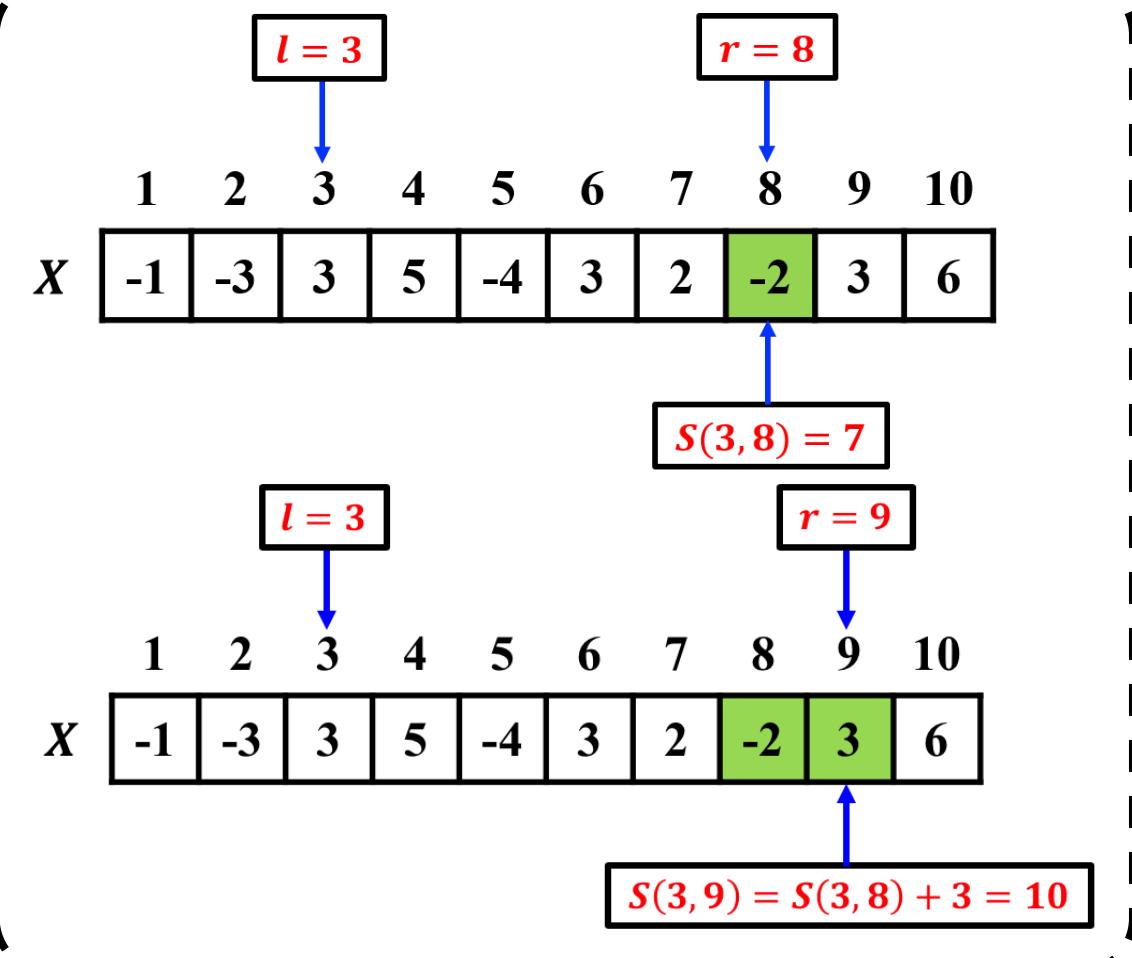


计算 $S(3, 9)$ 操作1次

优化枚举：实例分析



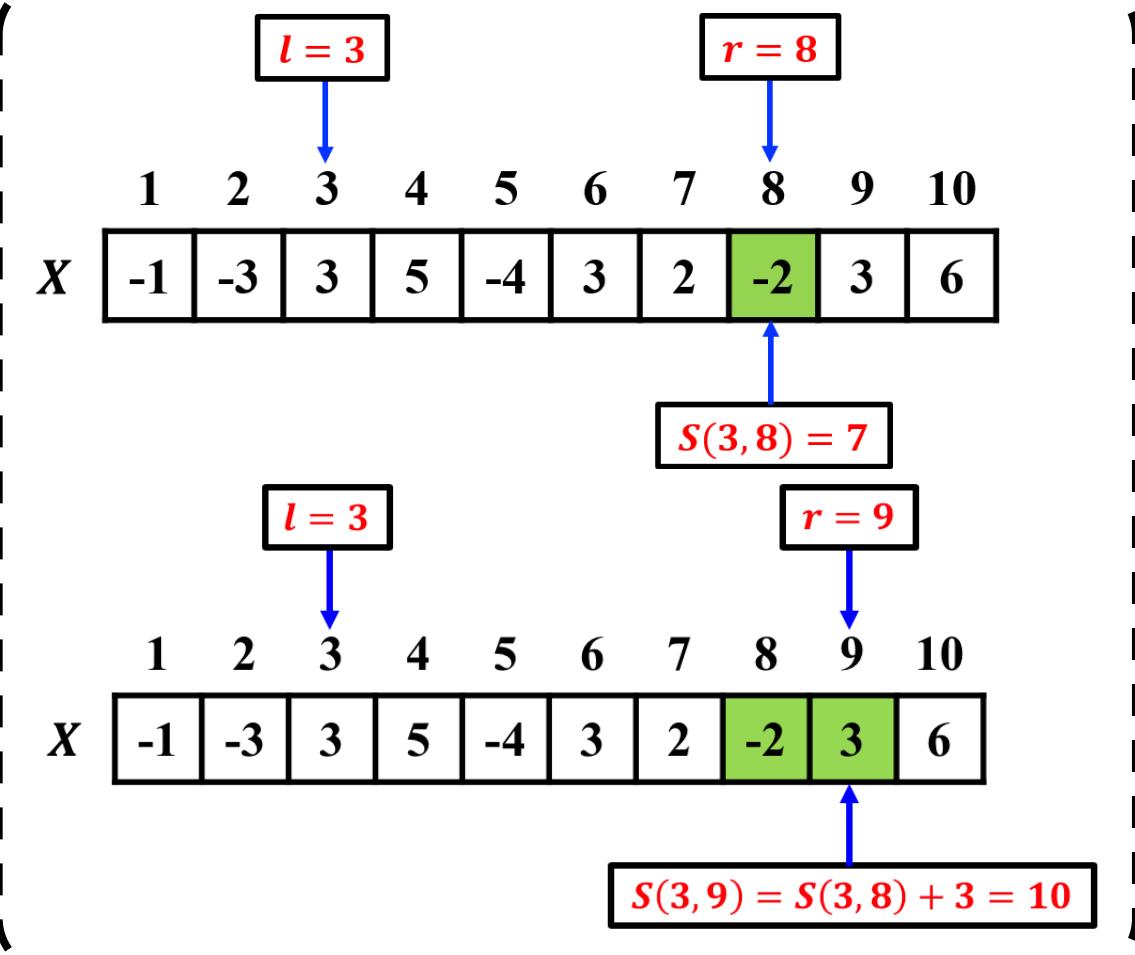
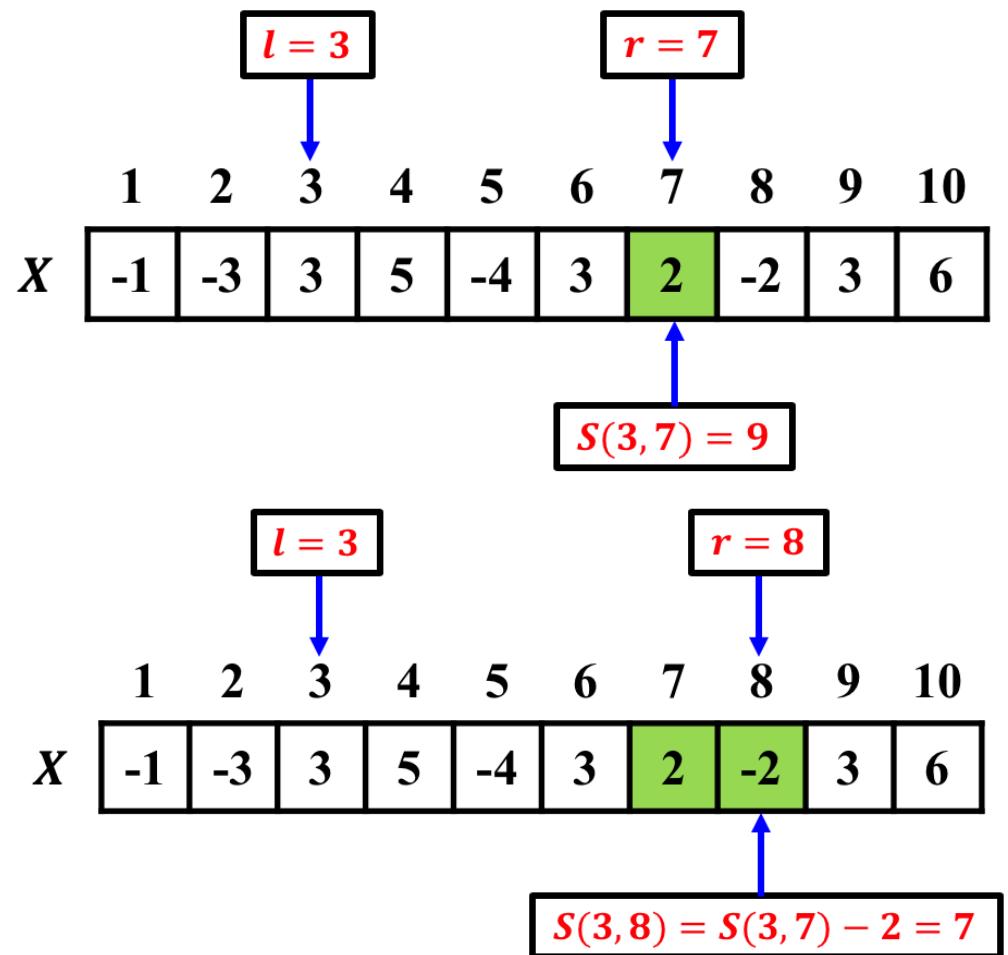
计算 $S(3, 8)$ 操作1次



计算 $S(3, 9)$ 操作1次

重复利用之前已经计算的数据

优化枚举：实例分析



$$S(l, r) = \sum_{i=l}^r X[i] = S(l, r-1) + X[r]$$

优化枚举：伪代码

- 核心思想: $S(l, r) = \sum_{i=l}^r X[i] = S(l, r-1) + X[r]$

输入: 数组 $X[1..n]$

输出: 最大子数组之和 S_{max}

$S_{max} \leftarrow -\infty$

for $l \leftarrow 1$ to n **do**

$S \leftarrow 0$

for $r \leftarrow l$ to n **do** -

$S \leftarrow S + X[r]$

 - $S_{max} \leftarrow \max\{S_{max}, S\}$

end

end

return S_{max}

迭代求解 S

优化枚举：伪代码

- 核心思想: $S(l, r) = \sum_{i=l}^r X[i] = S(l, r-1) + X[r]$

输入: 数组 $X[1..n]$

输出: 最大子数组之和 S_{max}

$S_{max} \leftarrow -\infty$

for $l \leftarrow 1$ to n **do**

$S \leftarrow 0$

for $r \leftarrow l$ to n **do**

$S \leftarrow S + X[r]$

$S_{max} \leftarrow \max\{S_{max}, S\}$

end

end

return S_{max}

更新 S_{max}

优化枚举：伪代码

- 核心思想: $S(l, r) = \sum_{i=l}^r X[i] = S(l, r-1) + X[r]$

输入: 数组 $X[1..n]$

输出: 最大子数组之和 S_{max}

```
 $S_{max} \leftarrow -\infty$ 
for  $l \leftarrow 1$  to  $n$  do
     $S \leftarrow 0$ 
    for  $r \leftarrow l$  to  $n$  do
         $S \leftarrow S + X[r]$ 
         $S_{max} \leftarrow \max\{S_{max}, S\}$ 
    end
end
return  $S_{max}$ 
```

$O(n)$ $O(n^2)$

时间复杂度: $O(n^2)$

分而治之：一般步骤

分解原问题

原问题分解成多个子问题

解决子问题

递归地求解各个子问题

合并问题解

将结果合并为原问题解

分而治之：一般步骤

$X[1..n]$

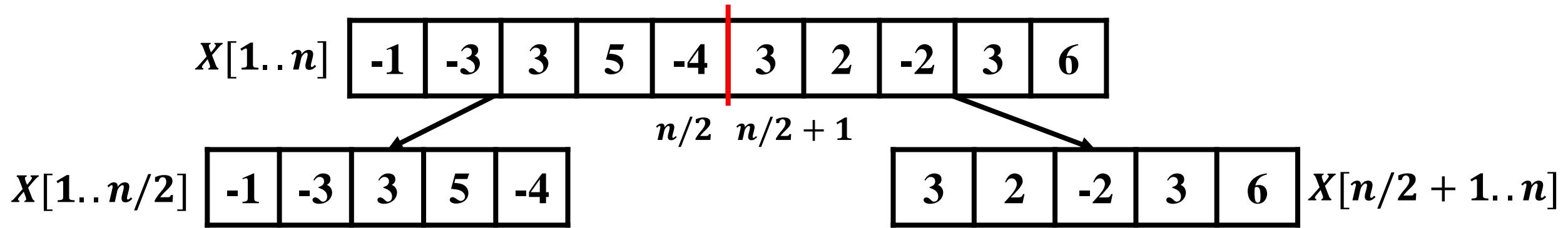
-1	-3	3	5	-4	3	2	-2	3	6
----	----	---	---	----	---	---	----	---	---

分解原问题

解决子问题

合并问题解

分而治之：一般步骤



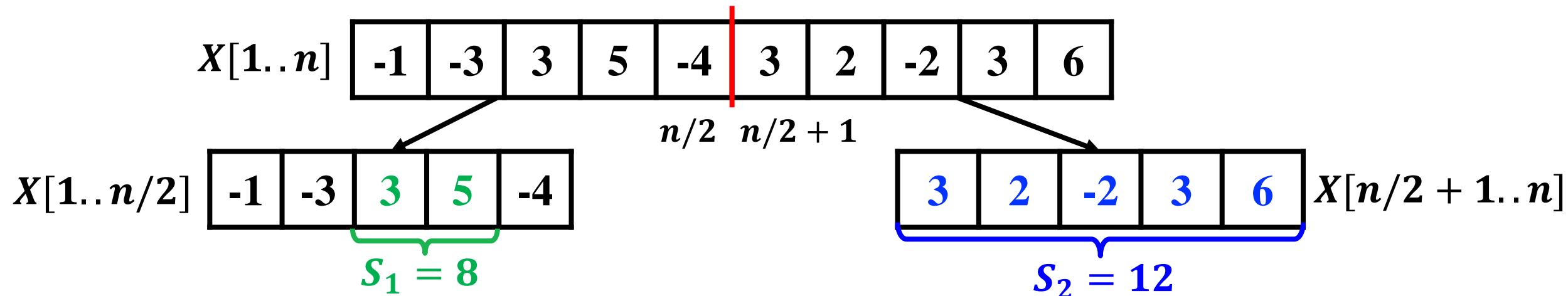
- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$

分解原问题

解决子问题

合并问题解

分而治之：一般步骤



- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$
- 递归求解子问题

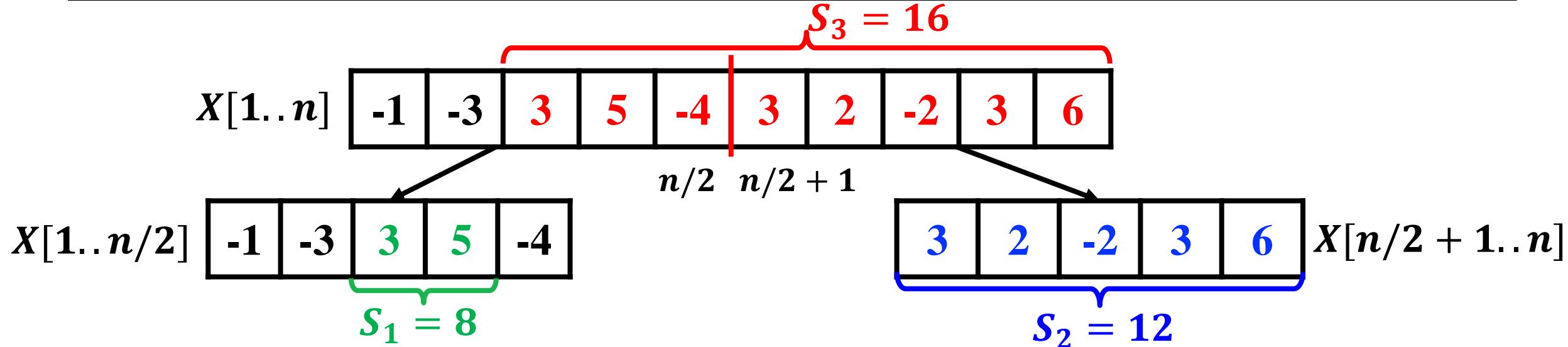
- S_1 : 数组 $X[1..n/2]$ 的最大子数组
- S_2 : 数组 $X[n/2 + 1..n]$ 的最大子数组

分解原问题

解决子问题

合并问题解

分而治之：一般步骤



- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$

分解原问题

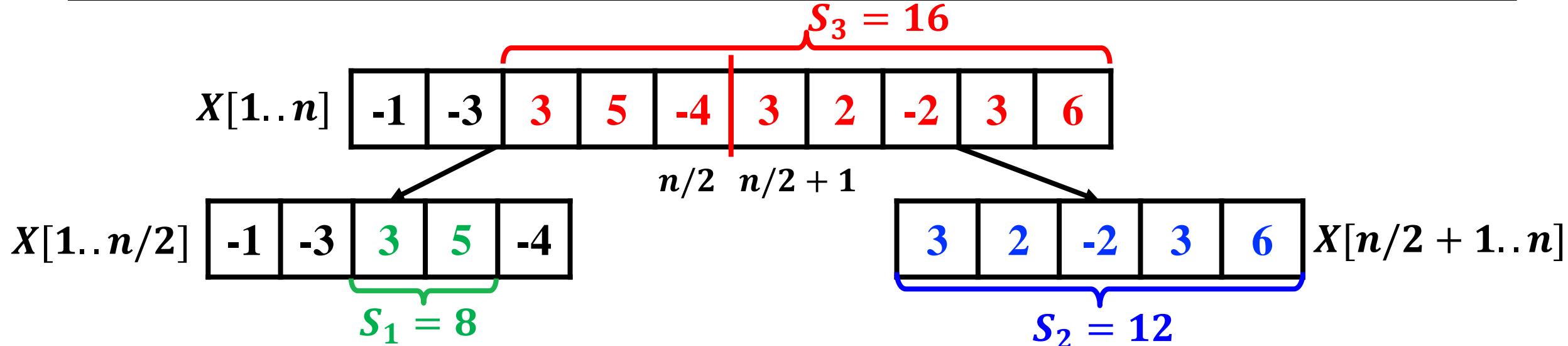
- 递归求解子问题

解决子问题

- S_1 : 数组 $X[1..n/2]$ 的最大子数组
 - S_2 : 数组 $X[n/2 + 1..n]$ 的最大子数组
- 合并子问题，得到 S_{max}
 - S_3 : 跨中点的最大子数组

合并问题解

分而治之：一般步骤



- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$

分解原问题

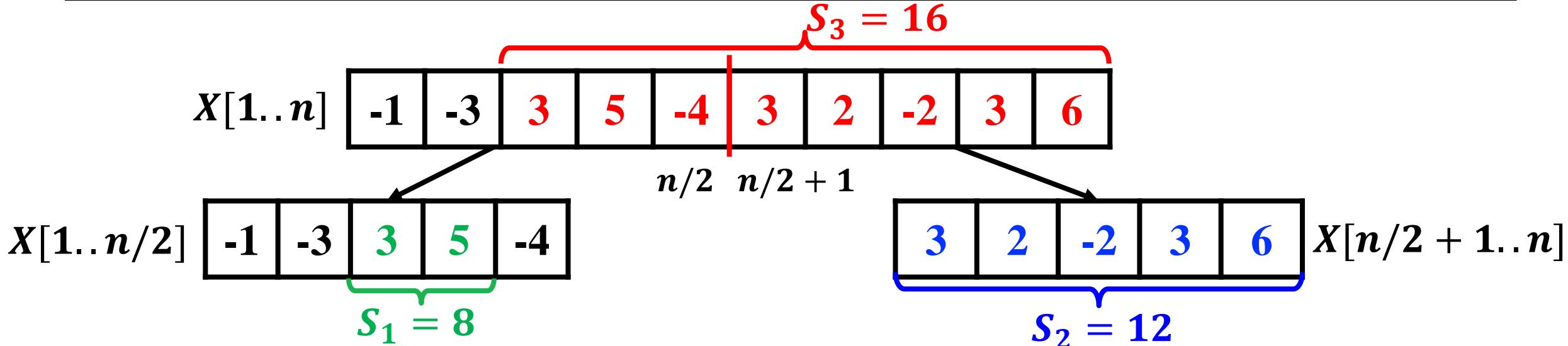
- 递归求解子问题

解决子问题

- S_1 : 数组 $X[1..n/2]$ 的最大子数组
 - S_2 : 数组 $X[n/2 + 1..n]$ 的最大子数组
- 合并子问题，得到 S_{max}
- S_3 : 跨中点的最大子数组
 - 数组 X 的最大子数组之和 $S_{max} = \max\{S_1, S_2, S_3\}$

合并问题解

分而治之：一般步骤



- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$

分解原问题

- 递归求解子问题

- S_1 : 数组 $X[1..n/2]$ 的最大子数组
- S_2 : 数组 $X[n/2 + 1..n]$ 的最大子数组

可递归求解

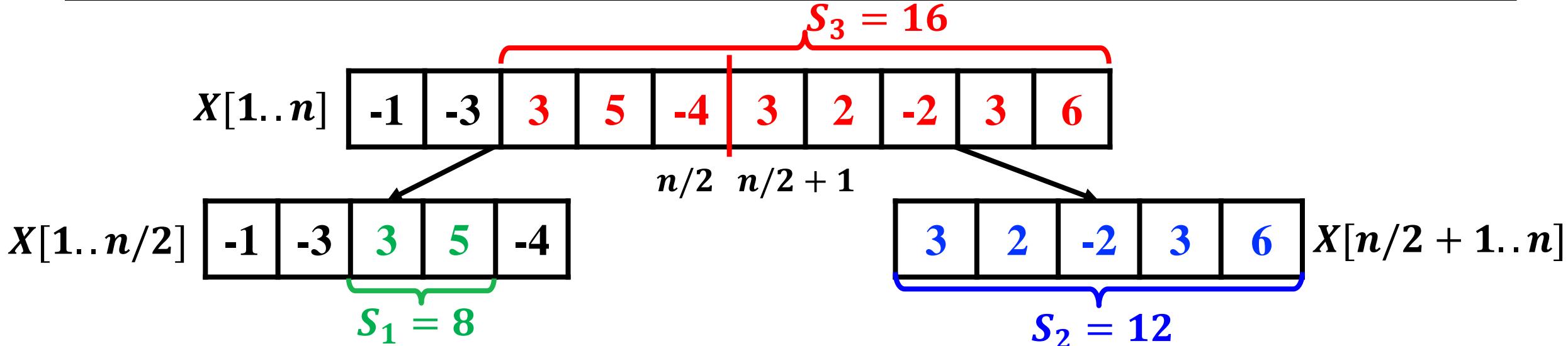
解决子问题

- 合并子问题，得到 S_{max}

合并问题解

- S_3 : 跨中点的最大子数组
- 数组 X 的最大子数组之和 $S_{max} = \max\{S_1, S_2, S_3\}$

分而治之：一般步骤



- 将数组 $X[1..n]$ 分为 $X[1..n/2]$ 和 $X[n/2 + 1..n]$

分解原问题

- 递归求解子问题

解决子问题

- S_1 : 数组 $X[1..n/2]$ 的最大子数组
- S_2 : 数组 $X[n/2 + 1..n]$ 的最大子数组

可递归求解

- 合并子问题，得到 S_{max}

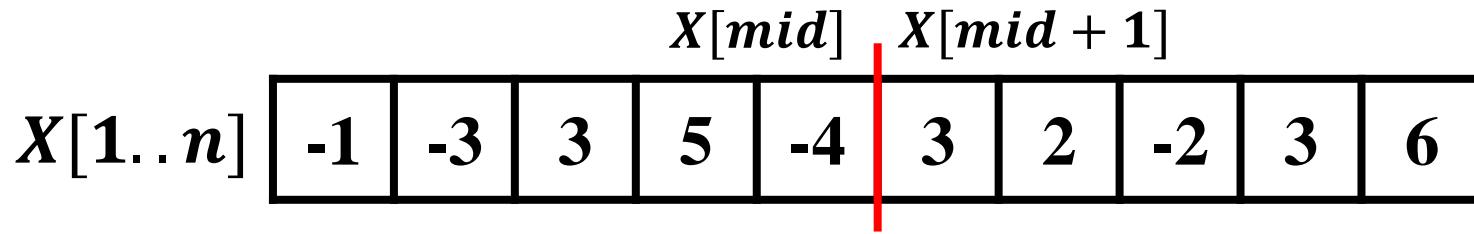
合并问题解

- S_3 : 跨中点的最大子数组

问题：如何求解 S_3 ？

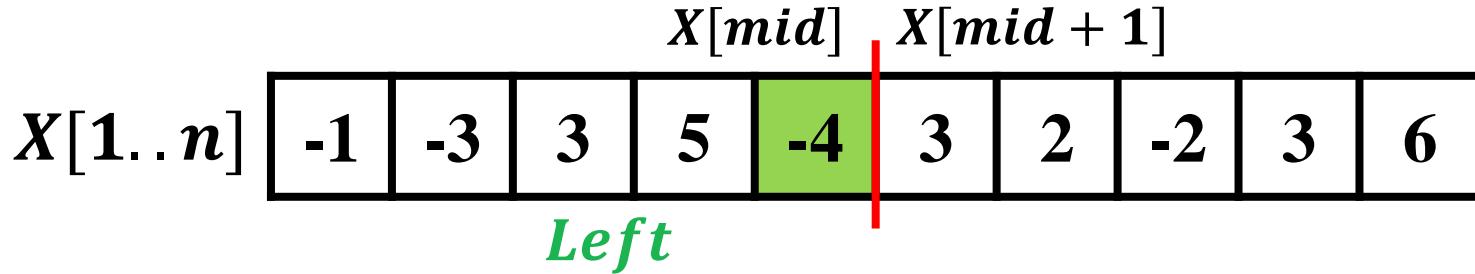
- 数组 X 的最大子数组之和 $S_{max} = \max\{S_1, S_2, S_3\}$

合并问题解：求解 S_3



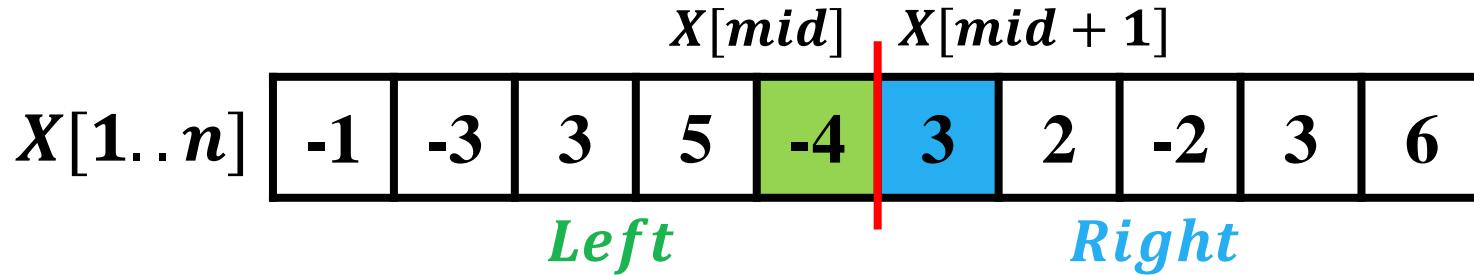
- 记 $mid = n/2$

合并问题解：求解 S_3



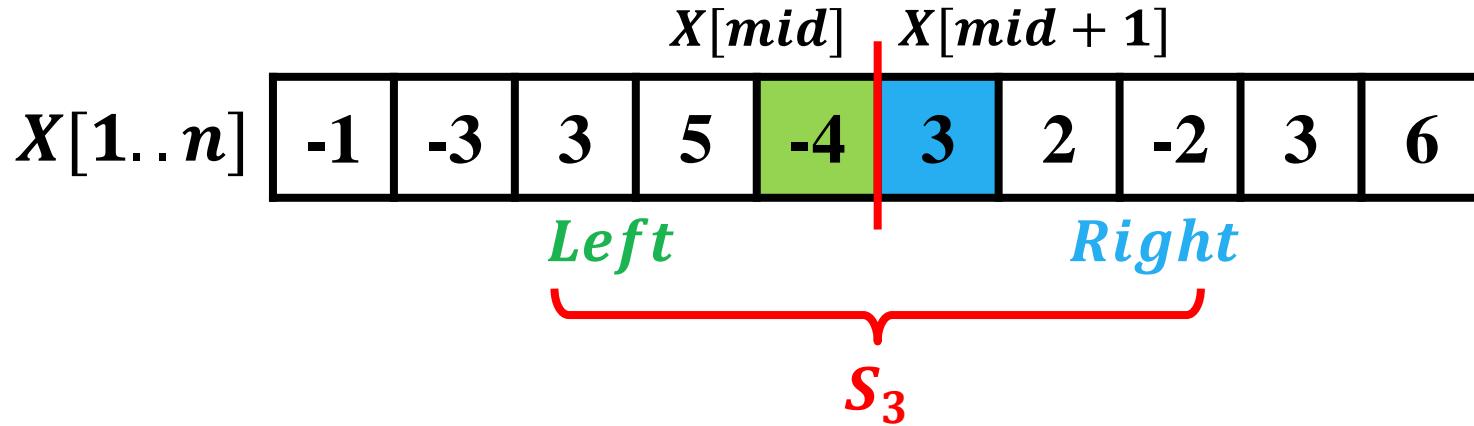
- 记 $mid = n/2$
- S_3 可分为左右两部分
 - *Left*: 以 $X[mid]$ 为结尾的最大子数组之和

合并问题解：求解 S_3



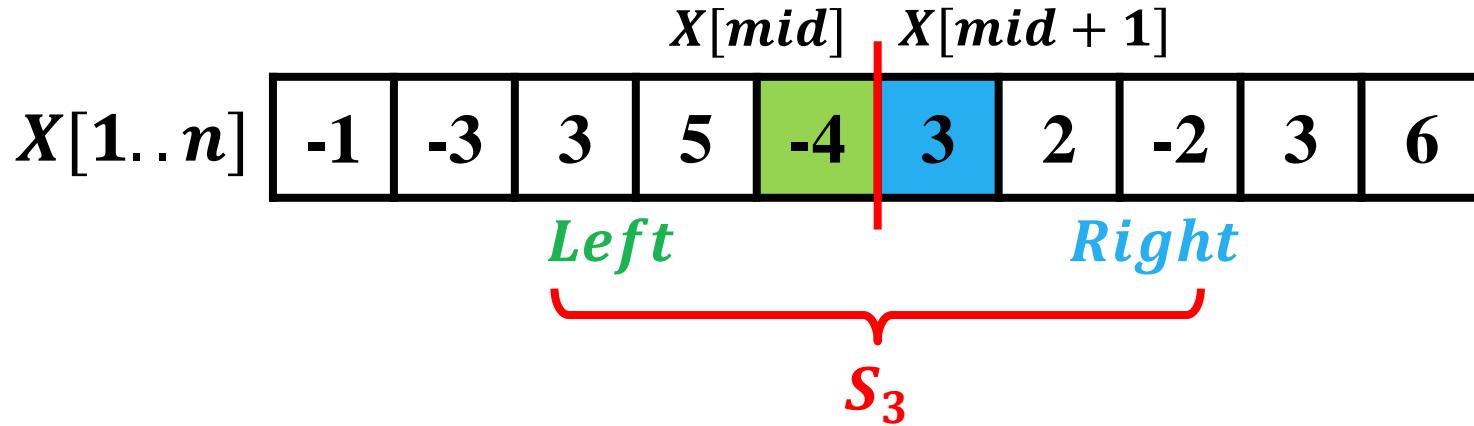
- 记 $mid = n/2$
- S_3 可分为左右两部分
 - *Left*: 以 $X[mid]$ 为结尾的最大子数组之和
 - *Right*: 以 $X[mid + 1]$ 为开头的最大子数组之和

合并问题解：求解 S_3



- 记 $mid = n/2$
- S_3 可分为左右两部分， $S_3 = Left + Right$
 - $Left$: 以 $X[mid]$ 为结尾的最大子数组之和
 - $Right$: 以 $X[mid + 1]$ 为开头的最大子数组之和

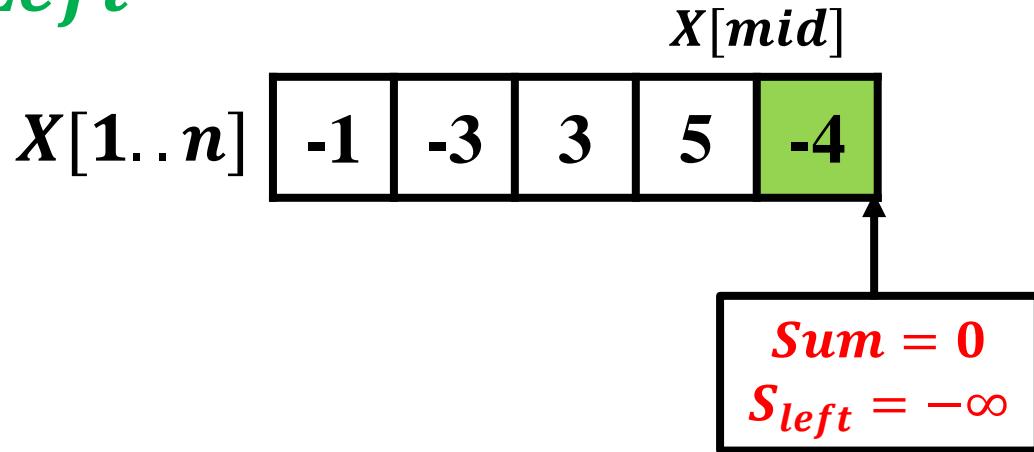
合并问题解：求解 S_3



- 记 $mid = n/2$
- S_3 可分为左右两部分， $S_3 = Left + Right$
 - $Left$: 以 $X[mid]$ 为结尾的最大子数组之和
 - $Right$: 以 $X[mid + 1]$ 为开头的最大子数组之和
 - 分别求出 $Left$ 和 $Right$ ，便可求出 S_3

合并问题解：求解 S_3

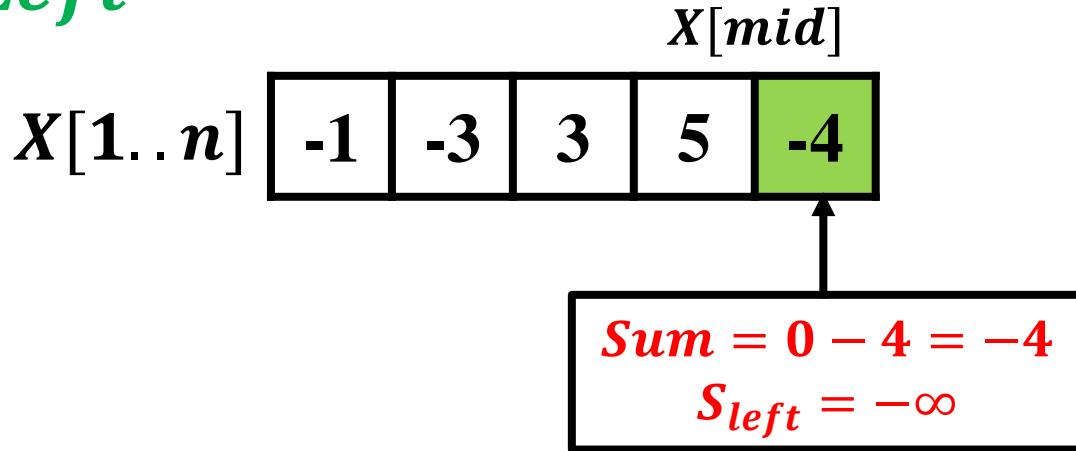
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

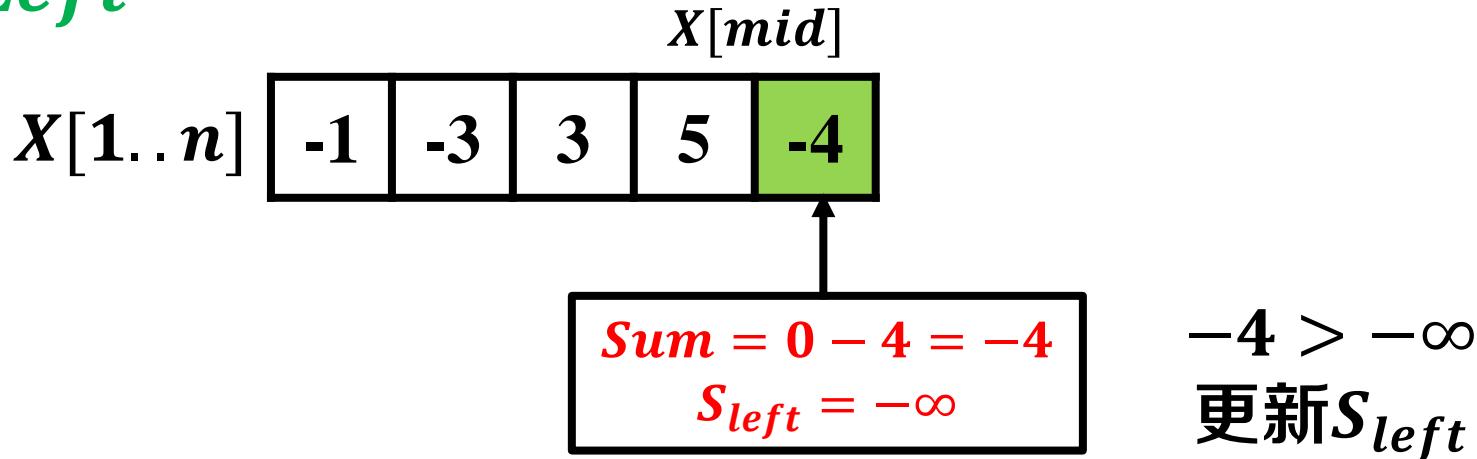
- 求解*Left*



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

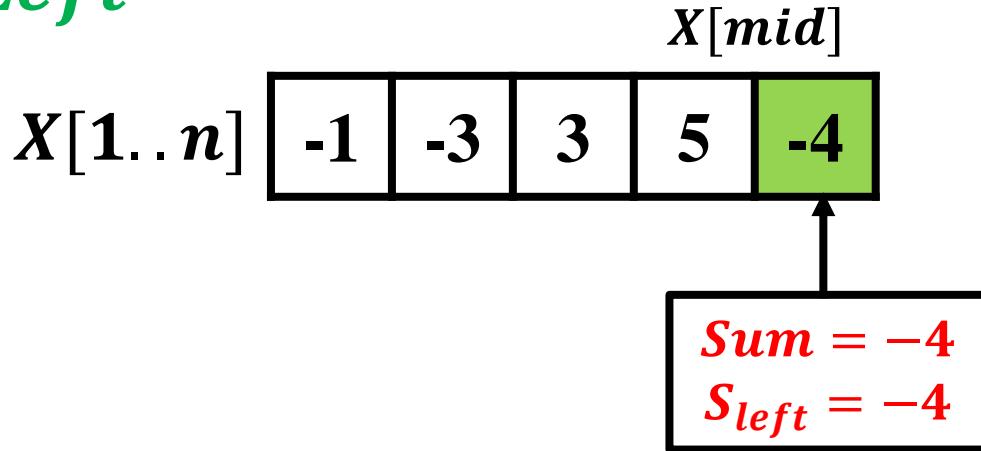
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

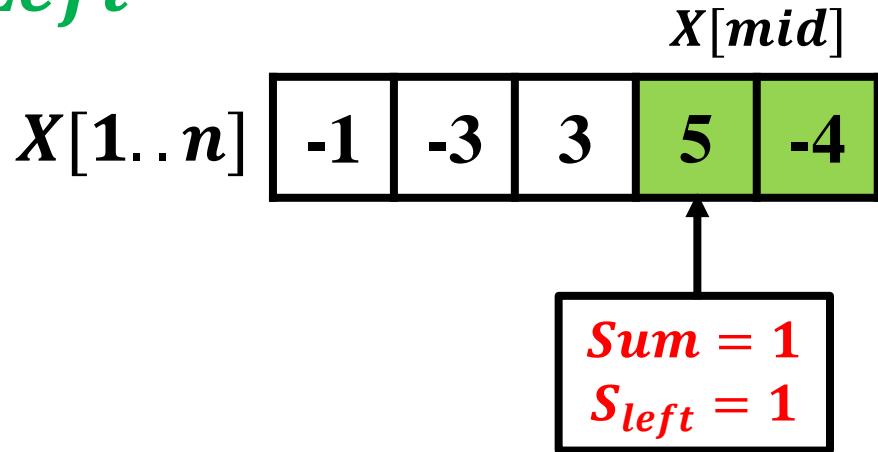
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

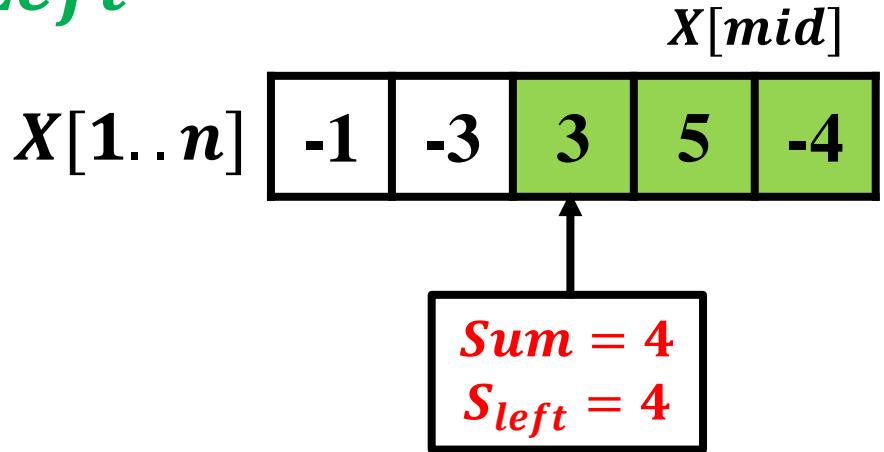
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

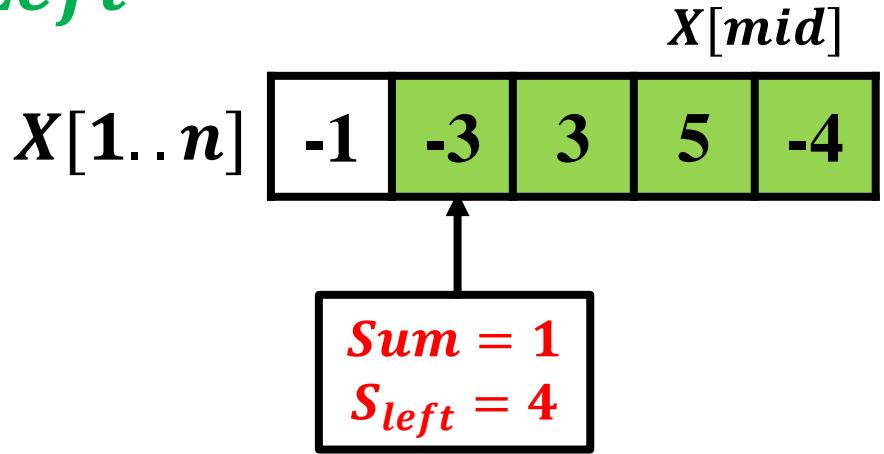
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进历求和，并记录最大值

合并问题解：求解 S_3

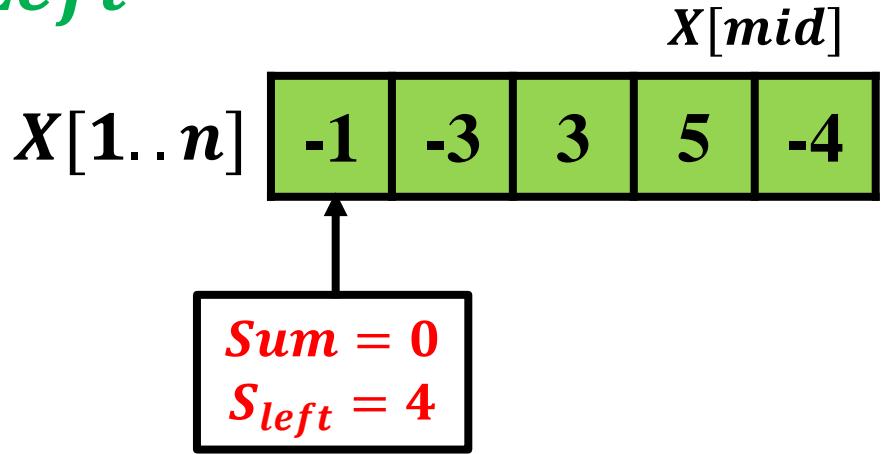
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值

合并问题解：求解 S_3

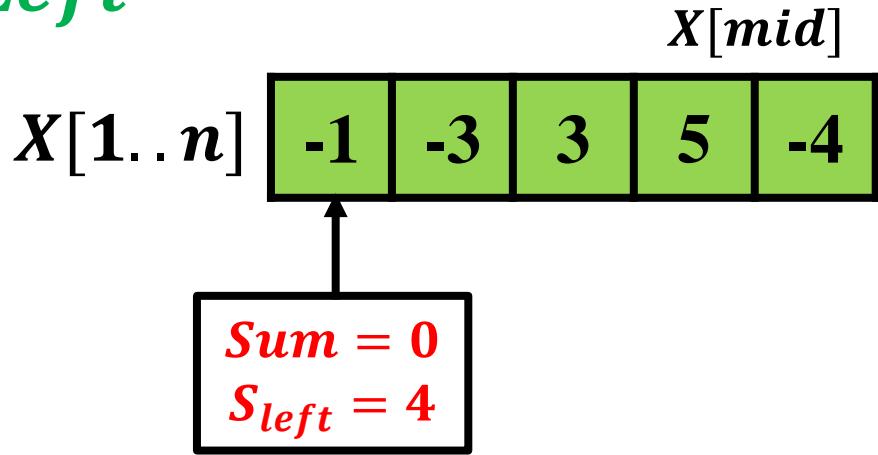
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进历求和，并记录最大值

合并问题解：求解 S_3

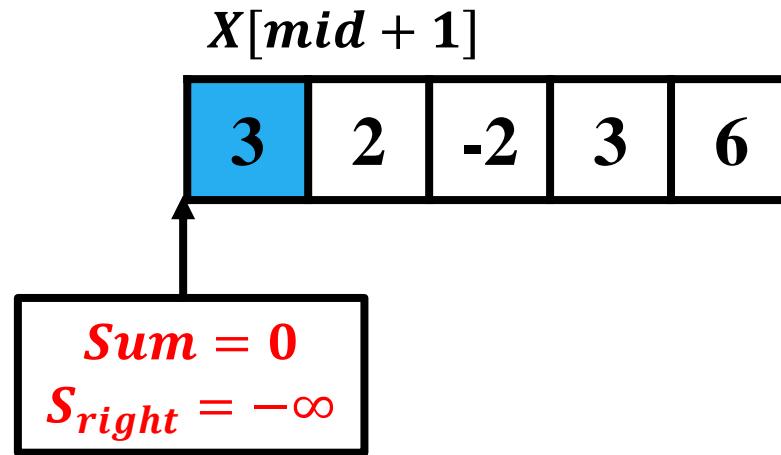
- 求解 $Left$



- 记 $mid = n/2$
- 从 $X[mid]$ 向前进遍历求和，并记录最大值
- 以 $X[mid]$ 为结尾的最大子数组之和 $S_{left} = 4$

合并问题解：求解 S_3

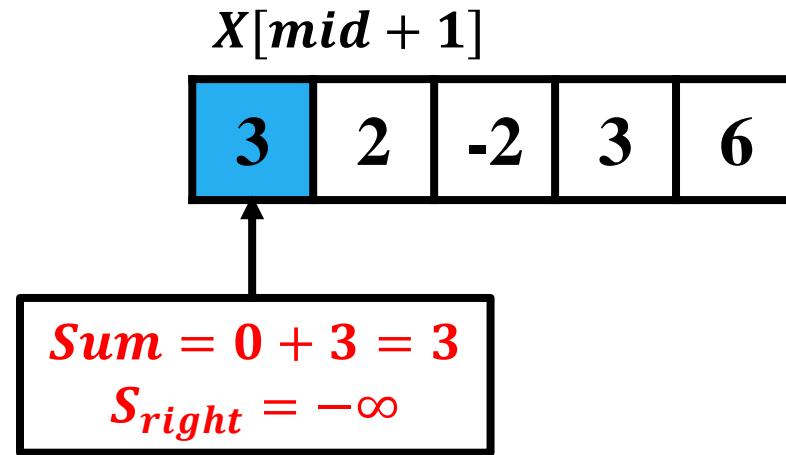
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

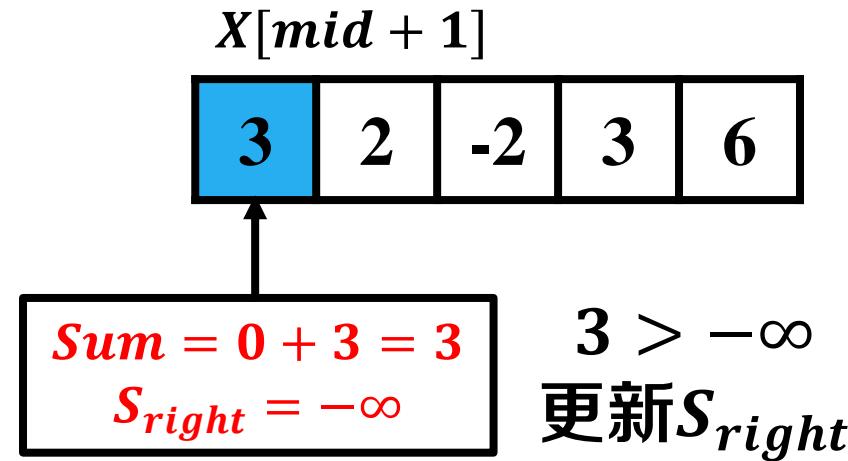
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

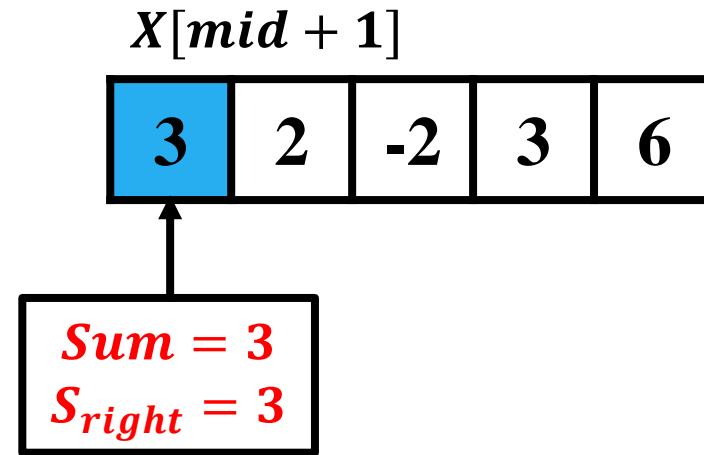
- 求解 $Right$



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

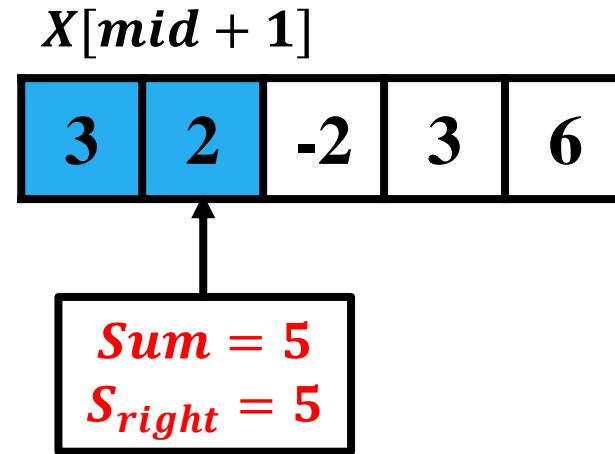
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

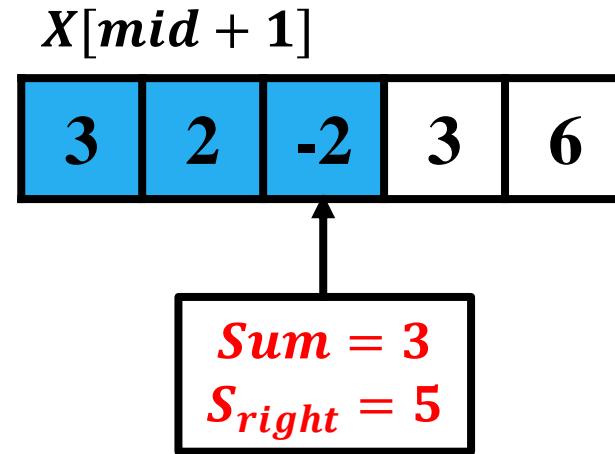
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

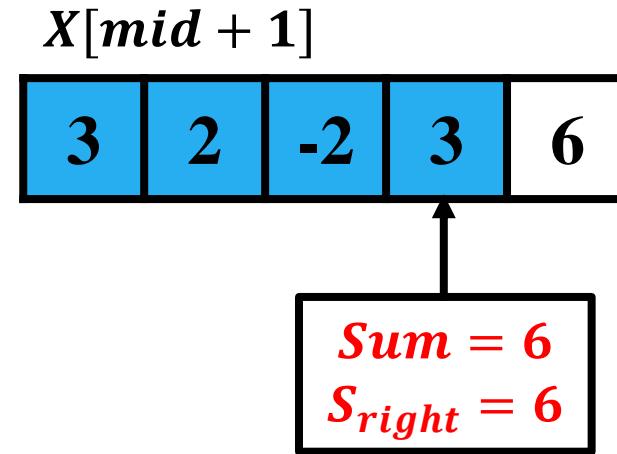
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

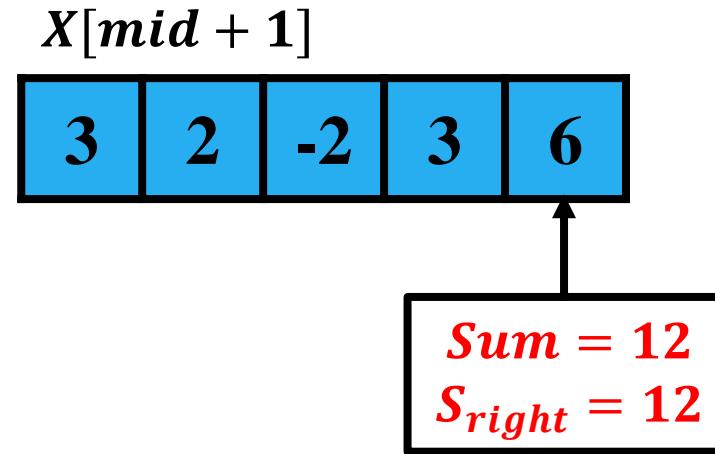
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

合并问题解：求解 S_3

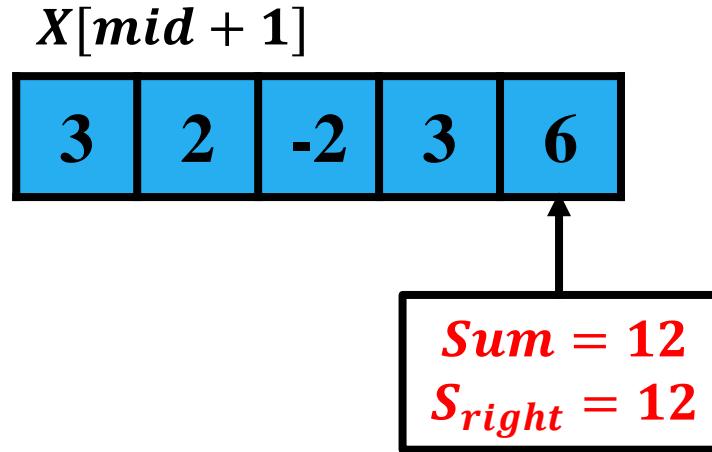
- 求解*Right*



- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值

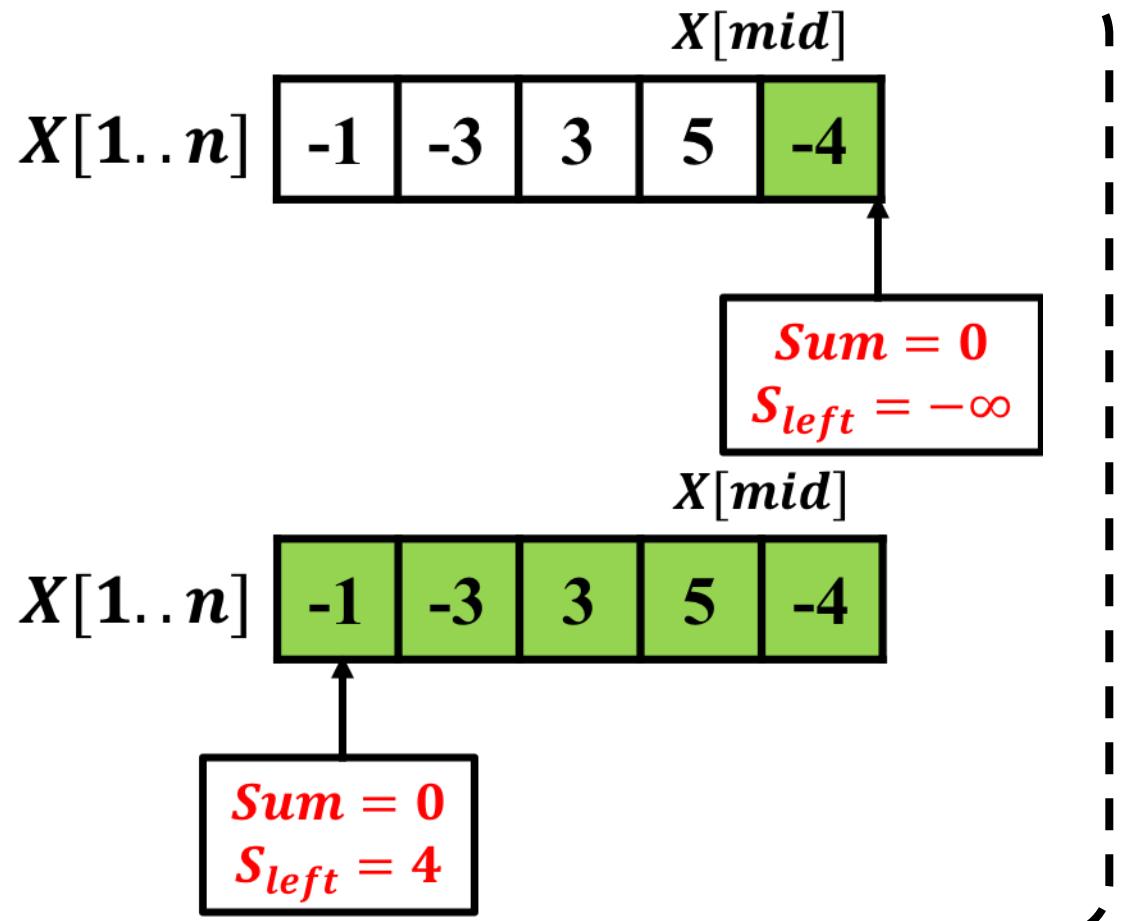
合并问题解：求解 S_3

- 求解*Right*



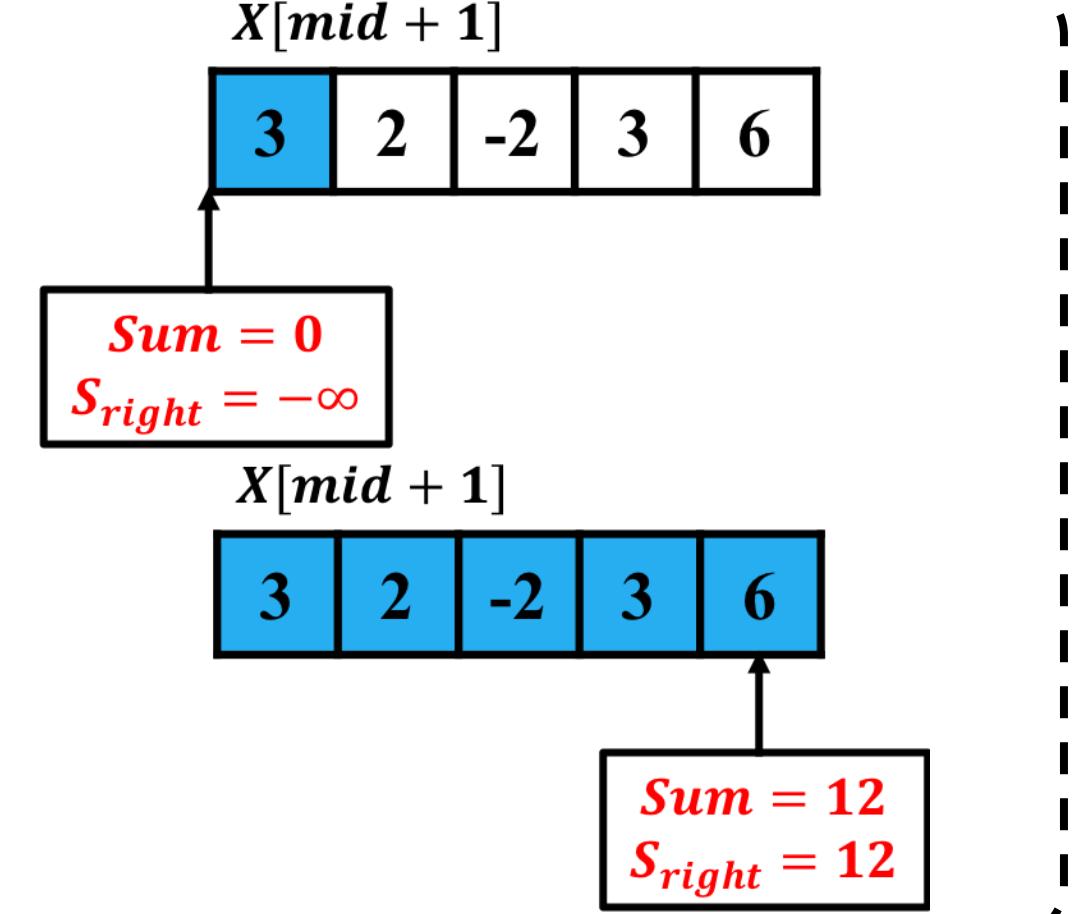
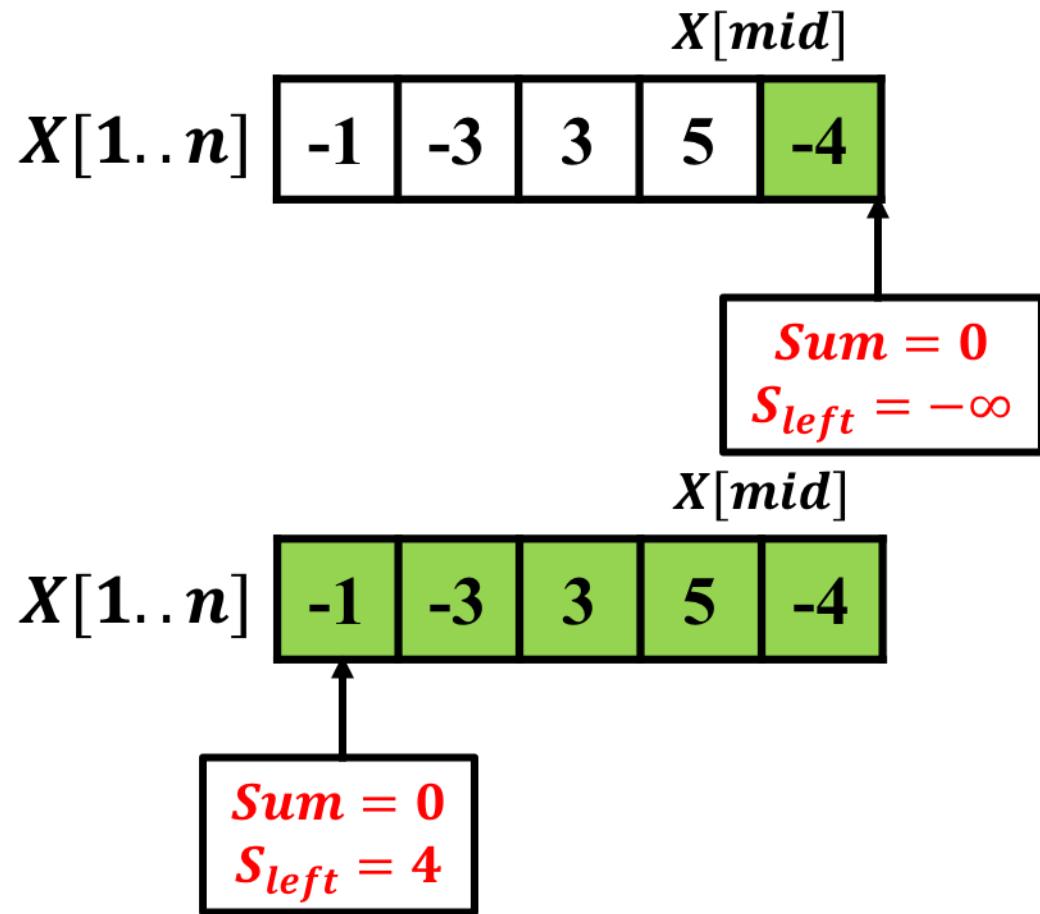
- 记 $mid = n/2$
- 从 $X[mid + 1]$ 向后遍历求和，并记录最大值
- 以 $X[mid + 1]$ 为开头的最大子数组之和 $S_{right} = 12$

求解 S_3 : 时间复杂度



求解 $Left$ 时间复杂度: $O(mid)$

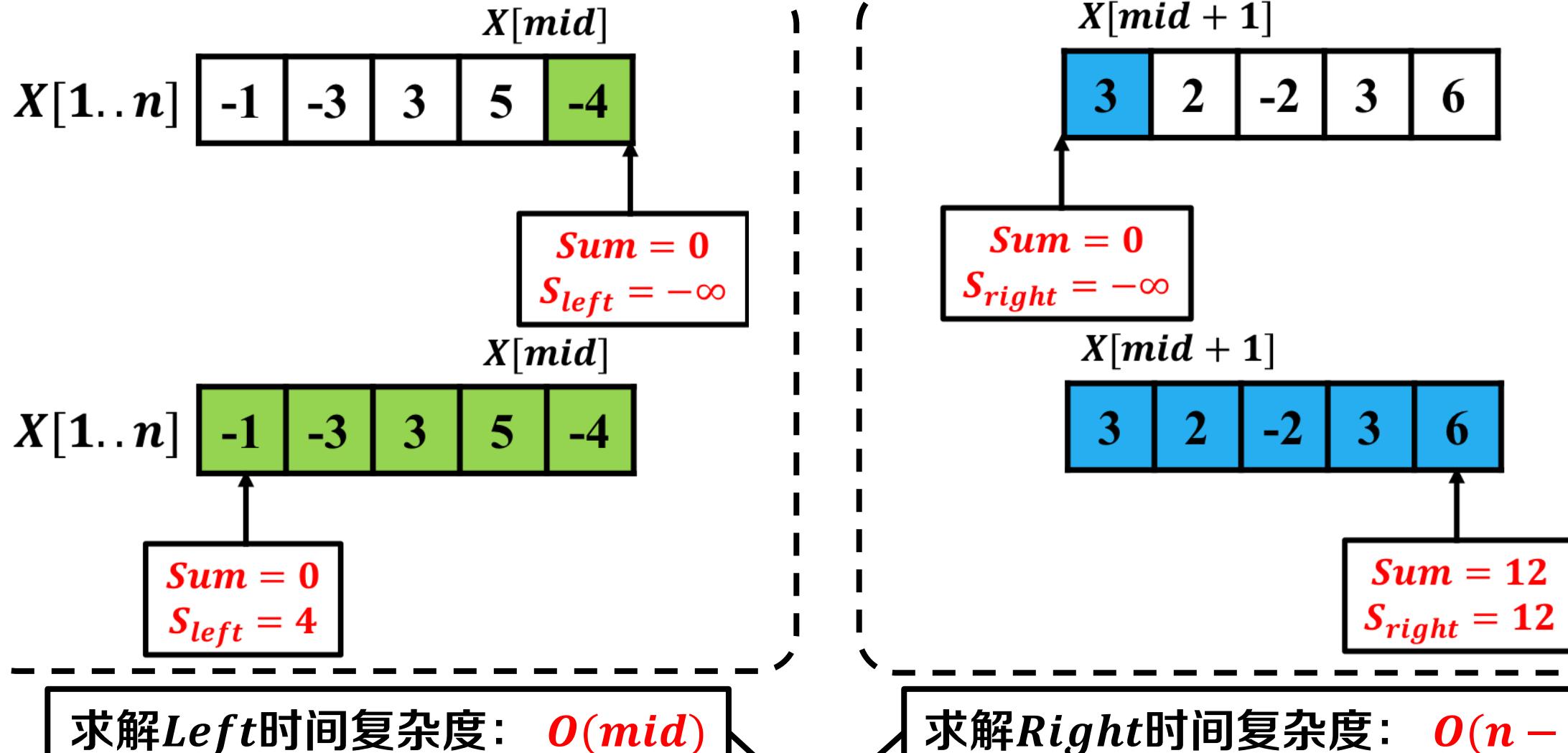
求解 S_3 : 时间复杂度



求解Left时间复杂度: **$O(mid)$**

求解Right时间复杂度: **$O(n - mid)$**

求解 S_3 : 时间复杂度



求解 S_3 时间复杂度: $O(n)$

求解 S_3 : 伪代码

- $\text{CrossingSubArray}(X, low, mid, high)$

求解 S_3 : 伪代码

- CrossingSubArray($X, low, mid, high$)

输入: 数组 X , 数组下标 $low, mid, high$

输出: 跨越中点的最大子数组之和 S_3

```
Sleft ← -∞
Sum ← 0
for l ← mid downto low do
    Sum ← Sum + X[l]
    Sleft ← max{Sleft, Sum}
end
Sright ← -∞
Sum ← 0
for r ← mid + 1 to high do
    Sum ← Sum + X[r]
    Sright ← max{Sright, Sum}
end
S3 ← Sleft + Sright
return S3
```

求解Left

求解 S_3 : 伪代码

- **CrossingSubArray($X, low, mid, high$)**

输入: 数组 X , 数组下标 $low, mid, high$

输出: 跨越中点的最大子数组之和 S_3

```
 $S_{left} \leftarrow -\infty$ 
```

```
 $Sum \leftarrow 0$ 
```

```
for  $l \leftarrow mid$  downto  $low$  do
```

```
    |  $Sum \leftarrow Sum + X[l]$ 
```

```
    |  $S_{left} \leftarrow \max\{S_{left}, Sum\}$ 
```

```
end
```

```
----- |  $S_{right} \leftarrow -\infty$ 
```

```
|  $Sum \leftarrow 0$ 
```

```
for  $r \leftarrow mid + 1$  to  $high$  do
```

```
    |  $Sum \leftarrow Sum + X[r]$ 
```

```
    |  $S_{right} \leftarrow \max\{S_{right}, Sum\}$ 
```

```
end
```

```
----- |  $S_3 \leftarrow S_{left} + S_{right}$ 
```

```
return  $S_3$ 
```

求解 $Right$

求解 S_3 : 伪代码

- CrossingSubArray($X, low, mid, high$)

输入: 数组 X , 数组下标 $low, mid, high$

输出: 跨越中点的最大子数组之和 S_3

```
 $S_{left} \leftarrow -\infty$ 
 $Sum \leftarrow 0$ 
for  $l \leftarrow mid$  downto  $low$  do
    |  $Sum \leftarrow Sum + X[l]$ 
    |  $S_{left} \leftarrow \max\{S_{left}, Sum\}$ 
end
 $S_{right} \leftarrow -\infty$ 
 $Sum \leftarrow 0$ 
for  $r \leftarrow mid + 1$  to  $high$  do
    |  $Sum \leftarrow Sum + X[r]$ 
    |  $S_{right} \leftarrow \max\{S_{right}, Sum\}$ 
end
 $S_3 \leftarrow S_{left} + S_{right}$ 
return  $S_3$ 
```

求解 S_3

求解 S_3 : 伪代码

- **CrossingSubArray($X, low, mid, high$)**

输入: 数组 X , 数组下标 $low, mid, high$

输出: 跨越中点的最大子数组之和 S_3

```
 $S_{left} \leftarrow -\infty$ 
 $Sum \leftarrow 0$ 
for  $l \leftarrow mid$  downto  $low$  do
    |  $Sum \leftarrow Sum + X[l]$ 
    |  $S_{left} \leftarrow \max\{S_{left}, Sum\}$ 
end
 $S_{right} \leftarrow -\infty$ 
 $Sum \leftarrow 0$ 
for  $r \leftarrow mid + 1$  to  $high$  do
    |  $Sum \leftarrow Sum + X[r]$ 
    |  $S_{right} \leftarrow \max\{S_{right}, Sum\}$ 
end
 $S_3 \leftarrow S_{left} + S_{right}$ 
return  $S_3$ 
```

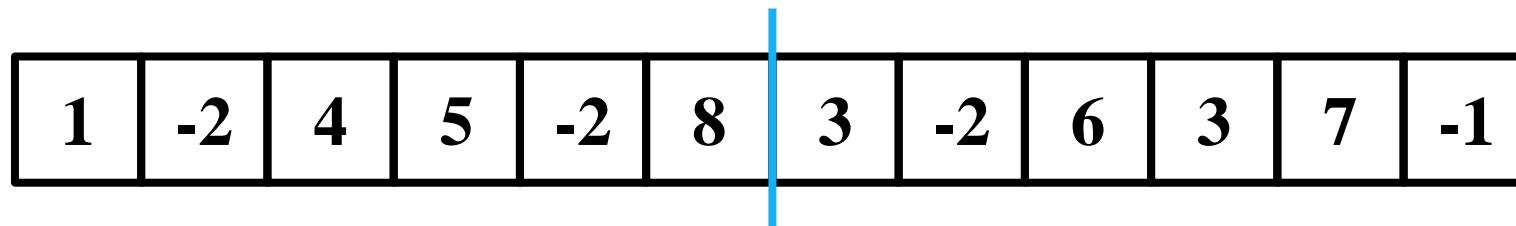
$\left. \begin{matrix} O(mid) \\ O(n - mid) \end{matrix} \right\}$

时间复杂度: $O(n)$

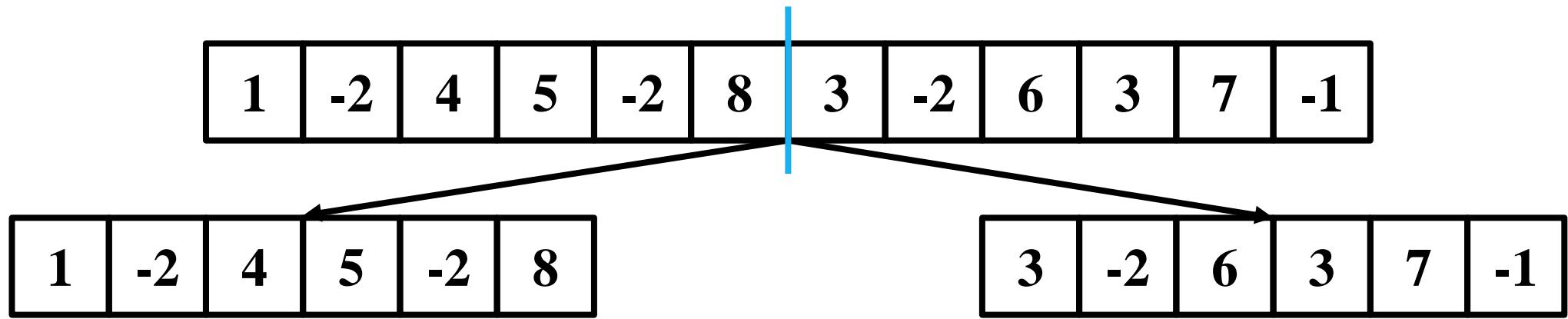
算法实例

1	-2	4	5	-2	8	3	-2	6	3	7	-1
---	----	---	---	----	---	---	----	---	---	---	----

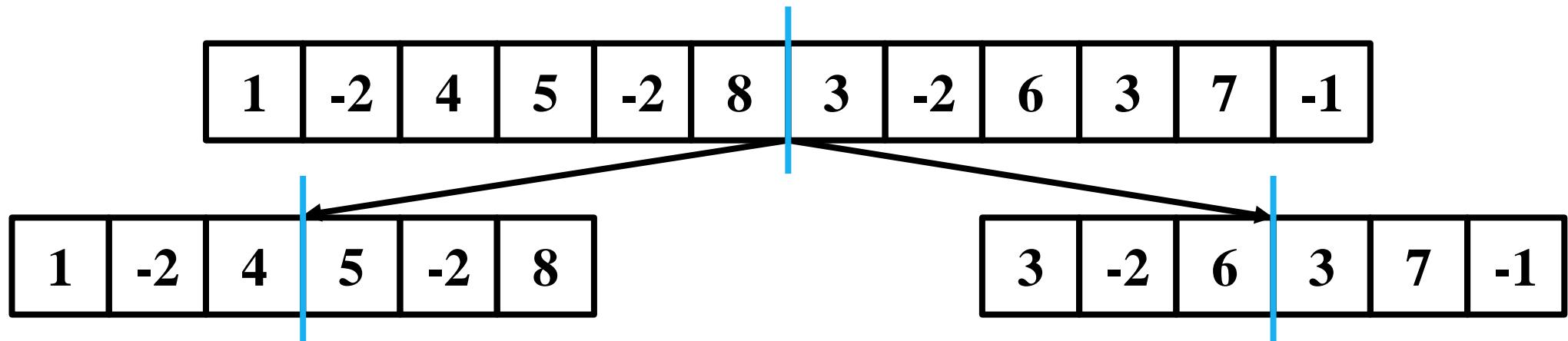
算法实例



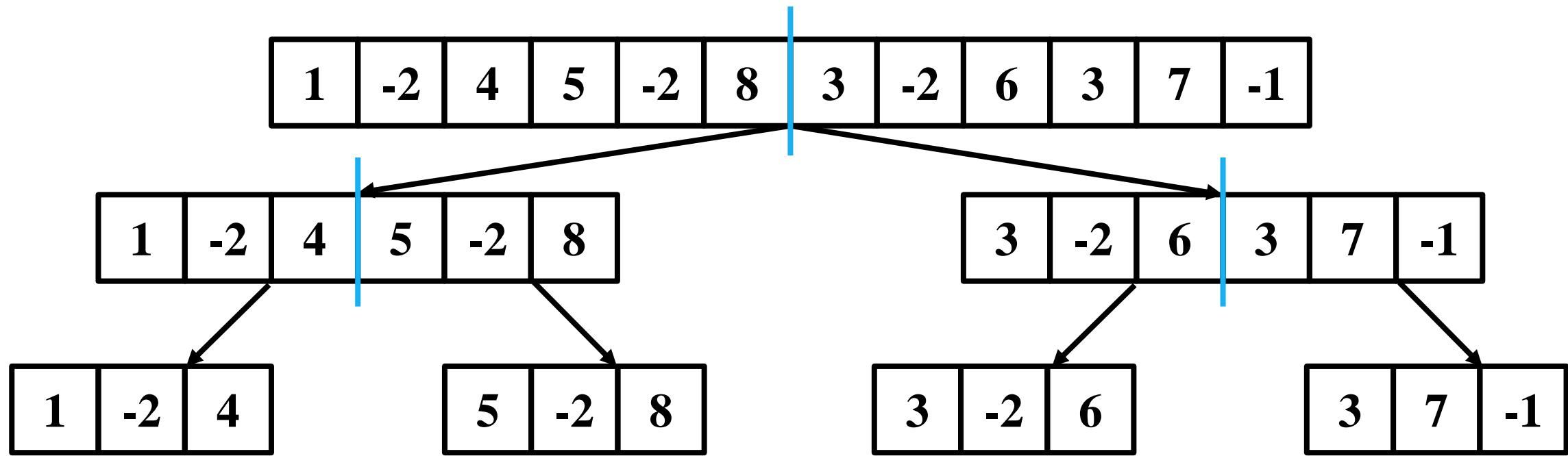
算法实例



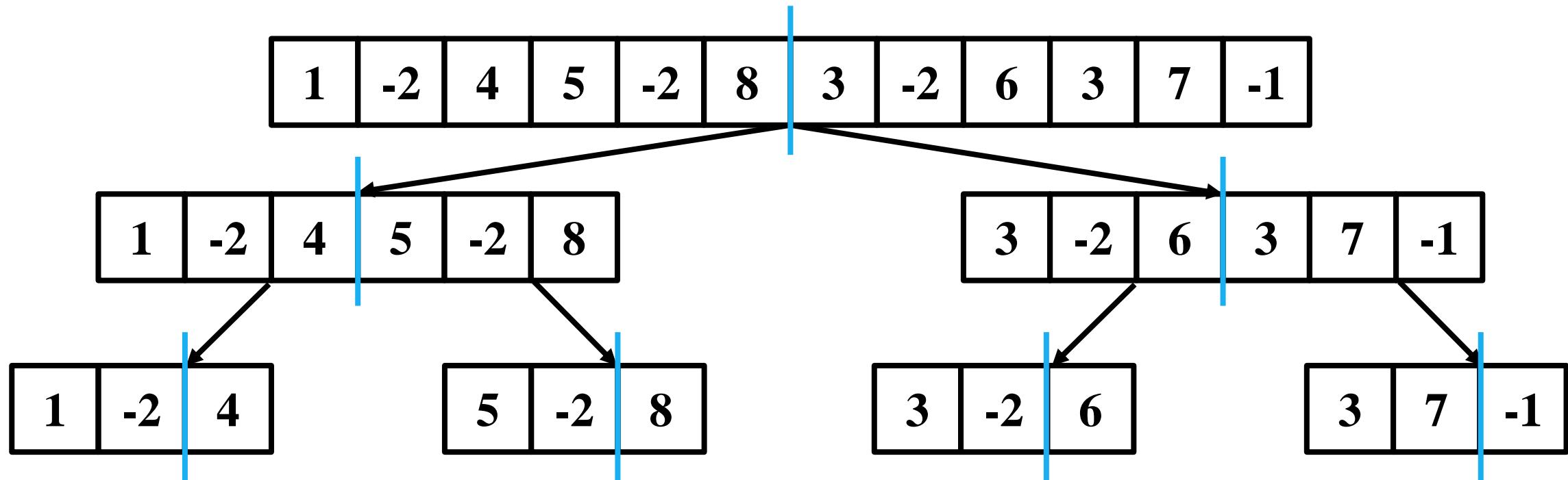
算法实例



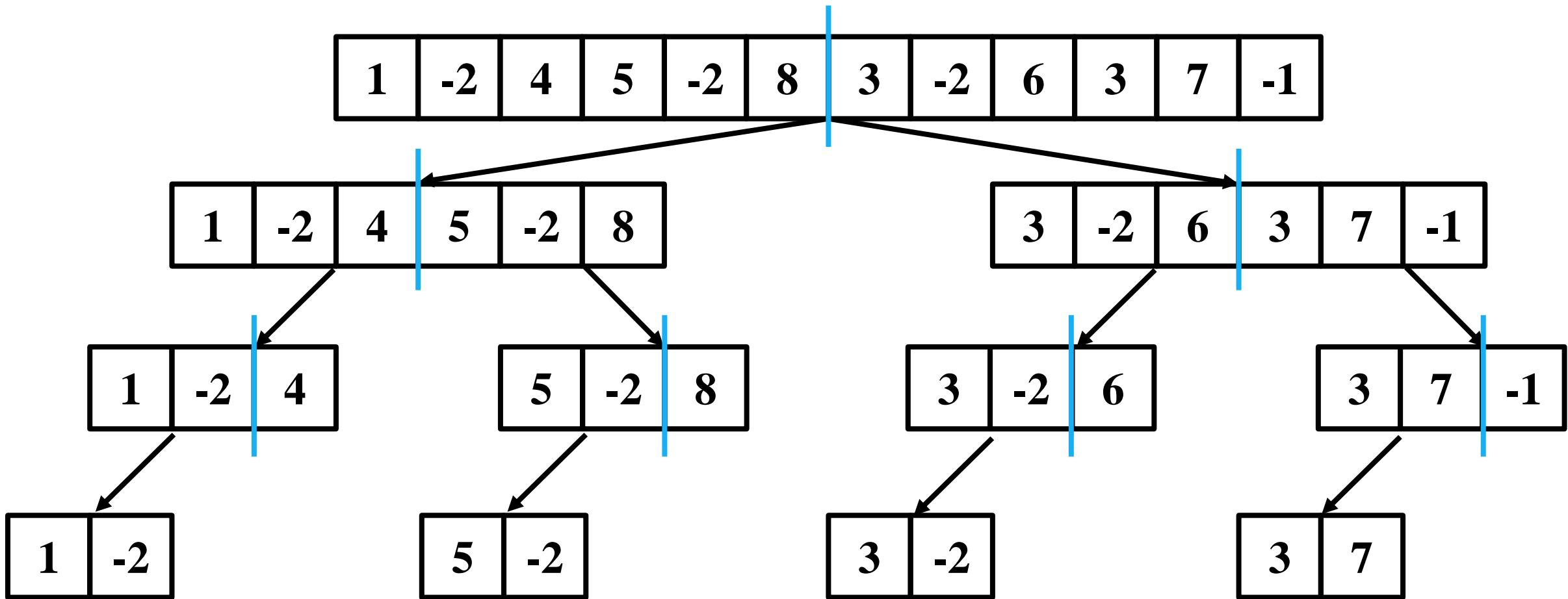
算法实例



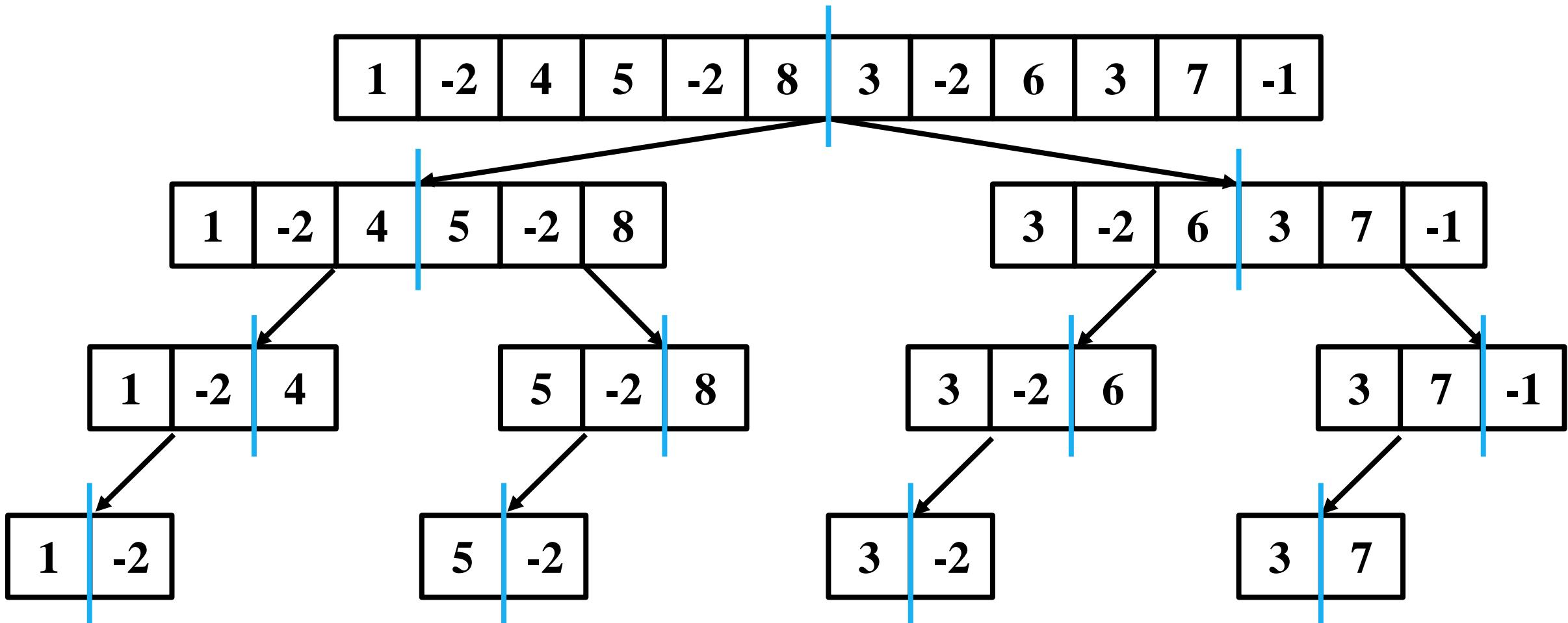
算法实例



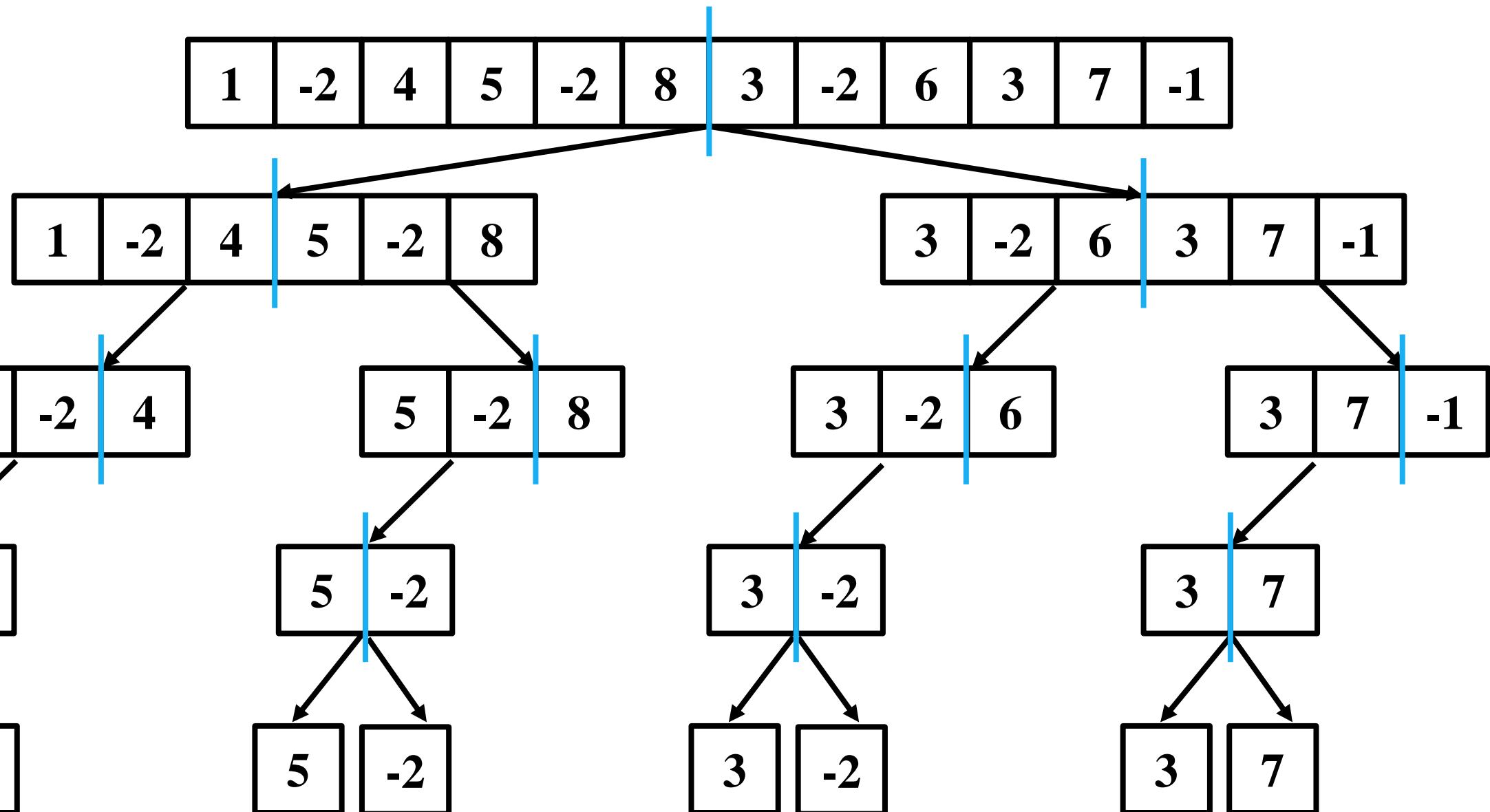
算法实例



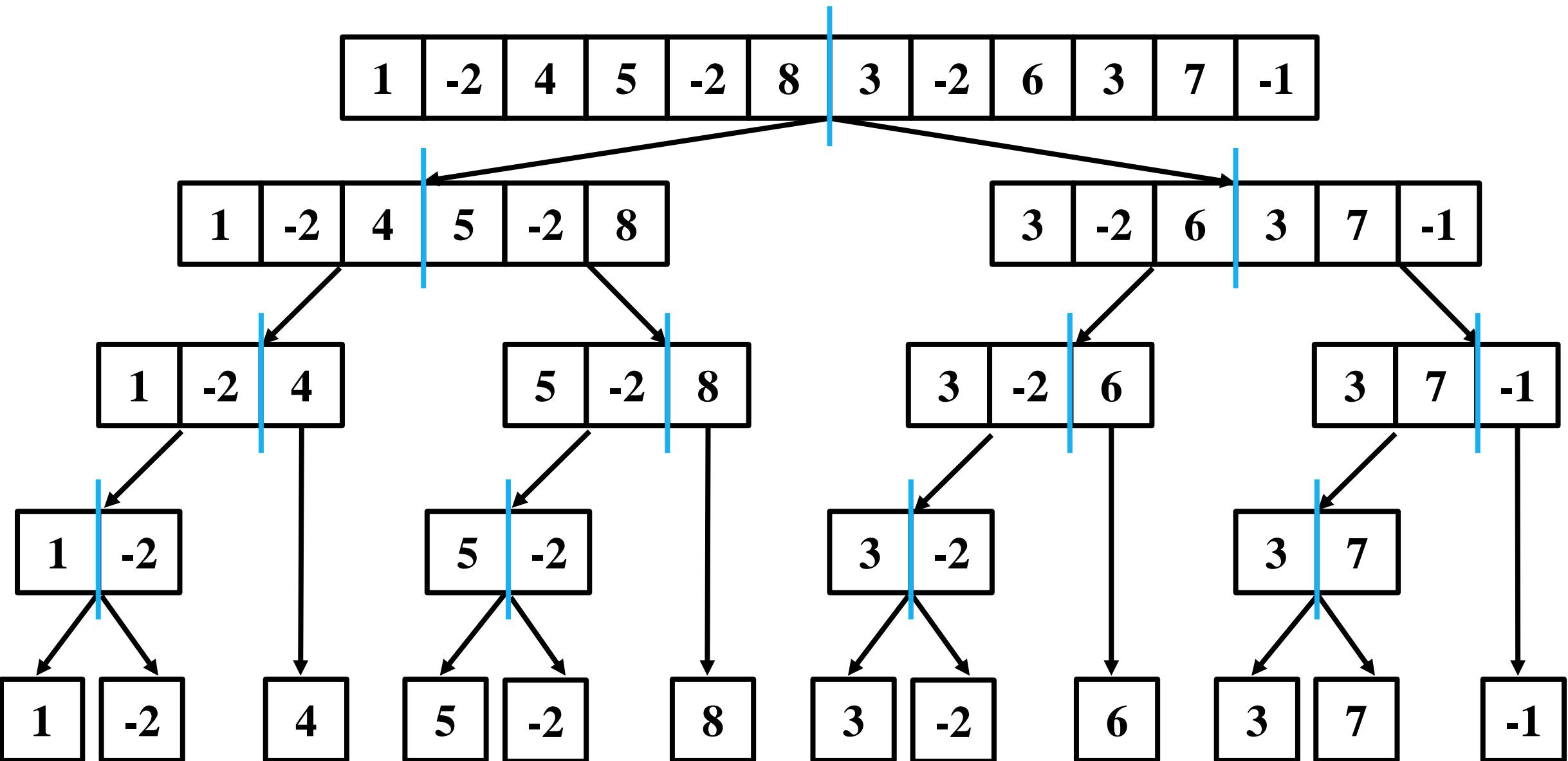
算法实例



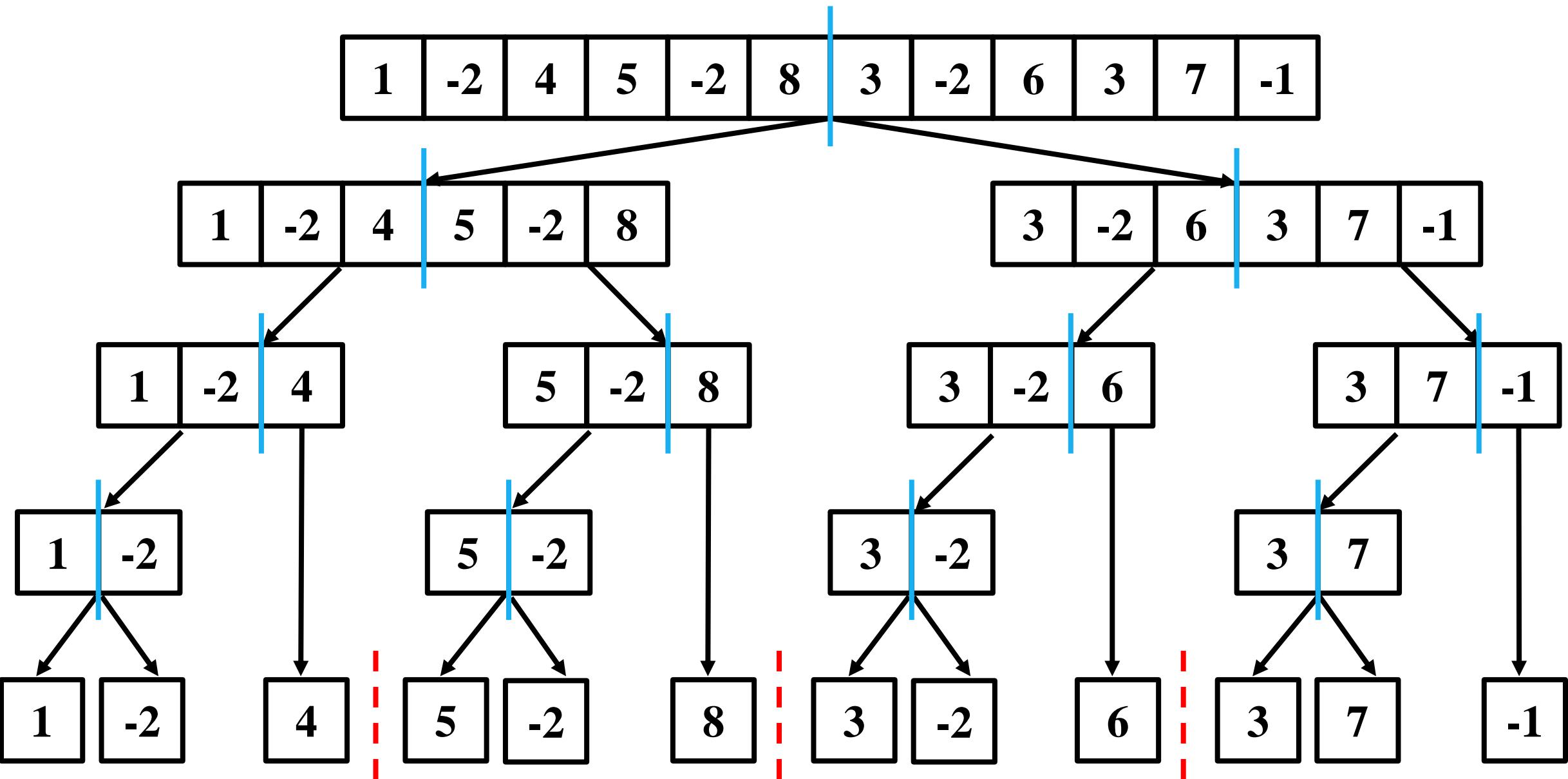
算法实例



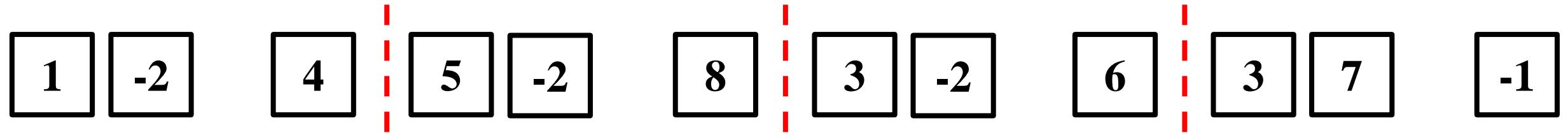
算法实例



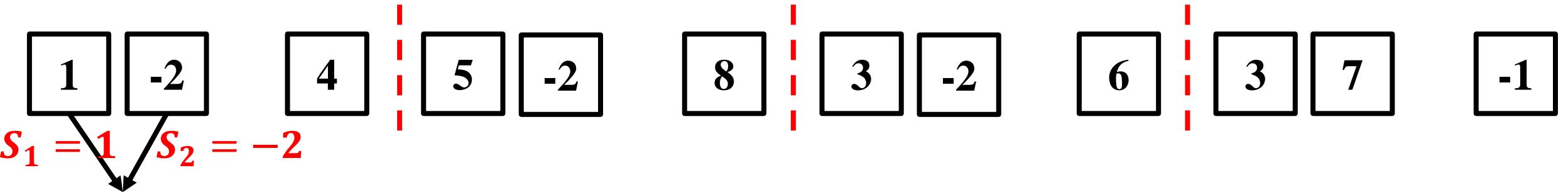
算法实例



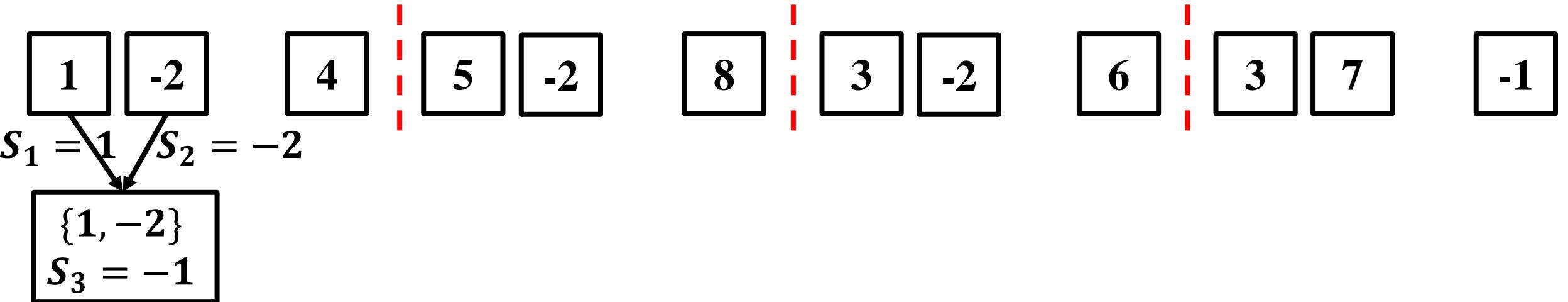
算法实例



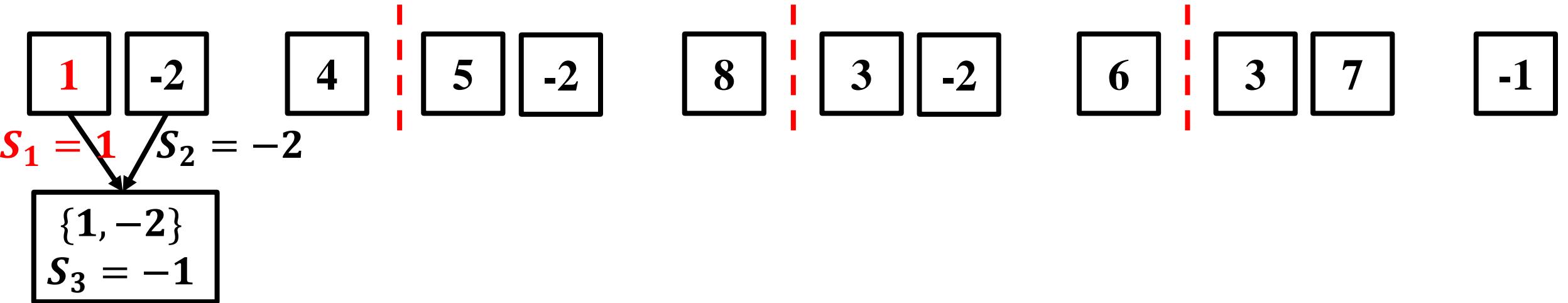
算法实例



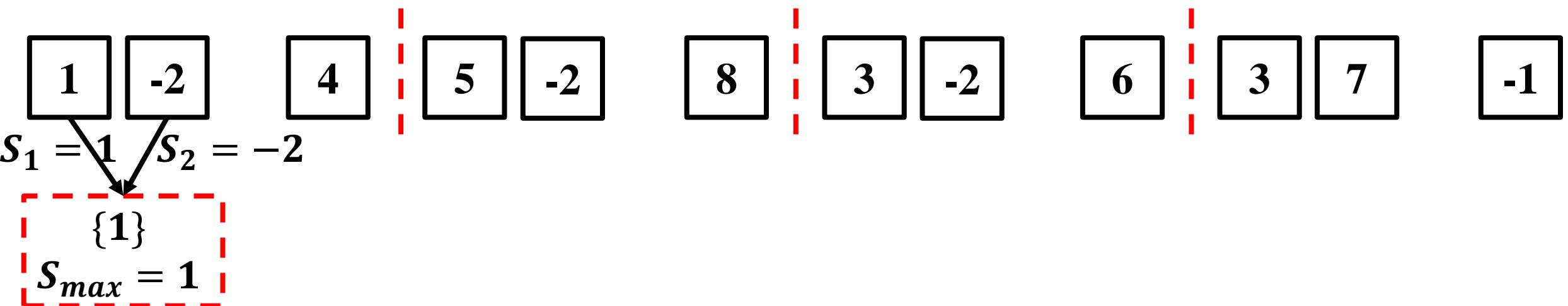
算法实例



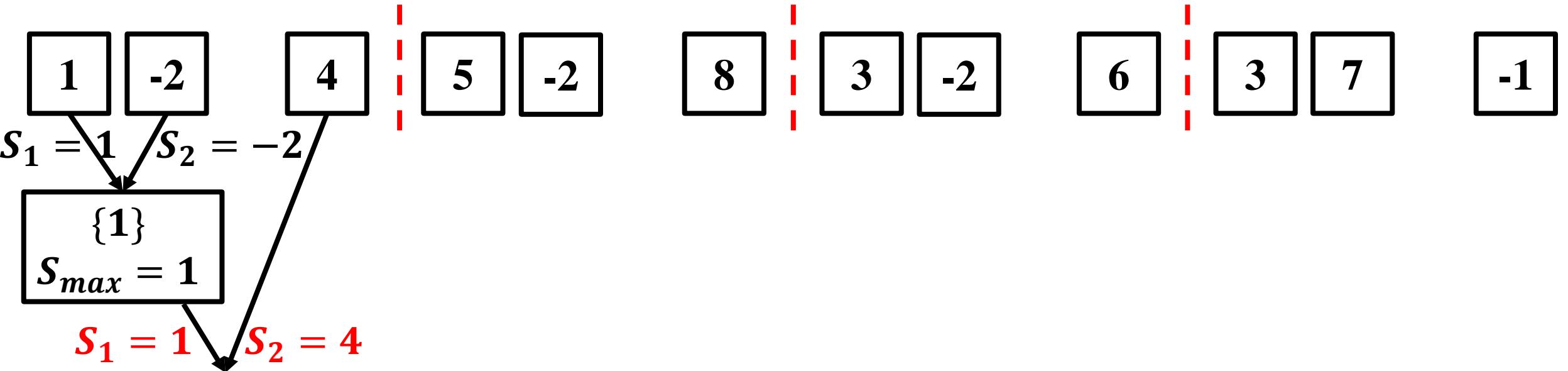
算法实例



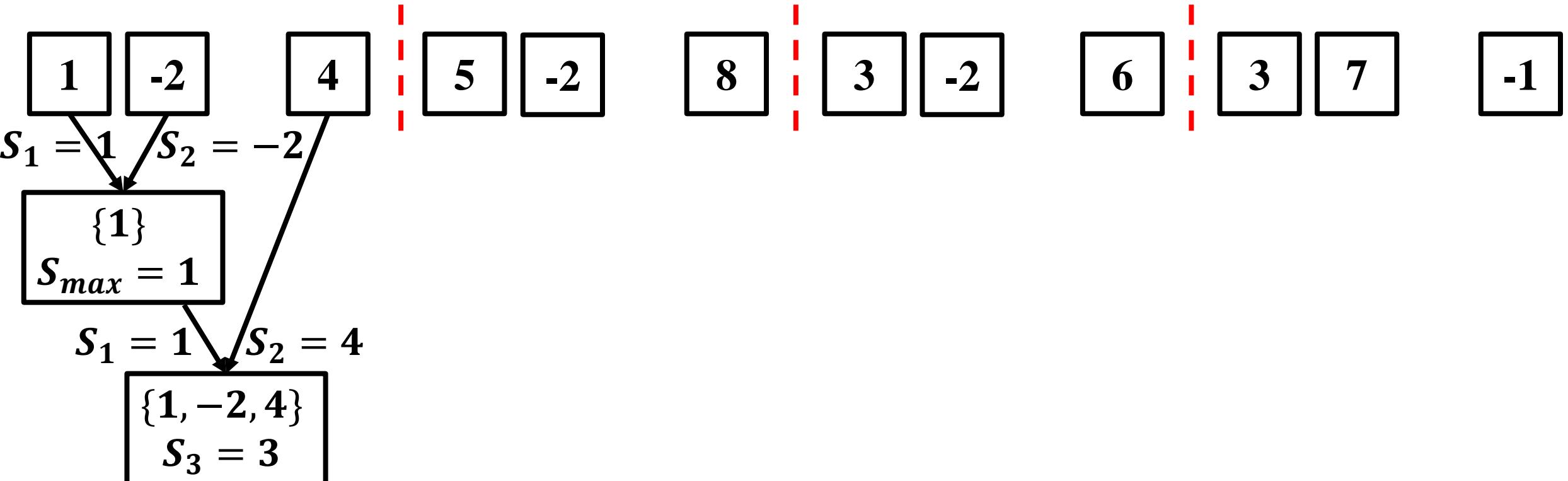
算法实例



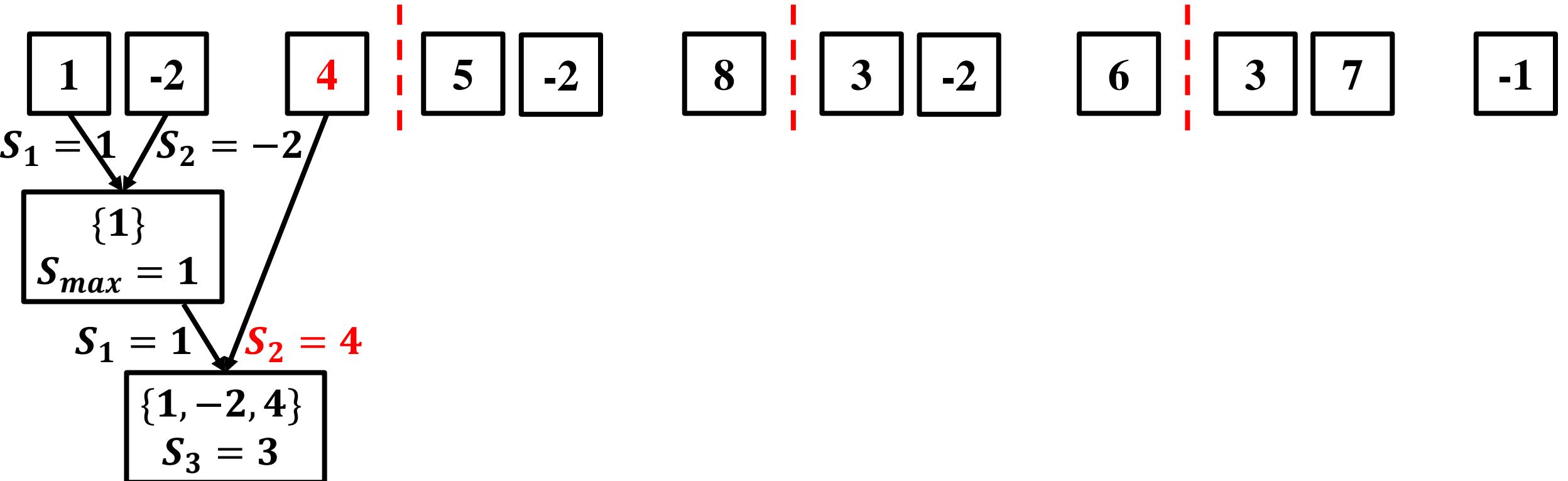
算法实例



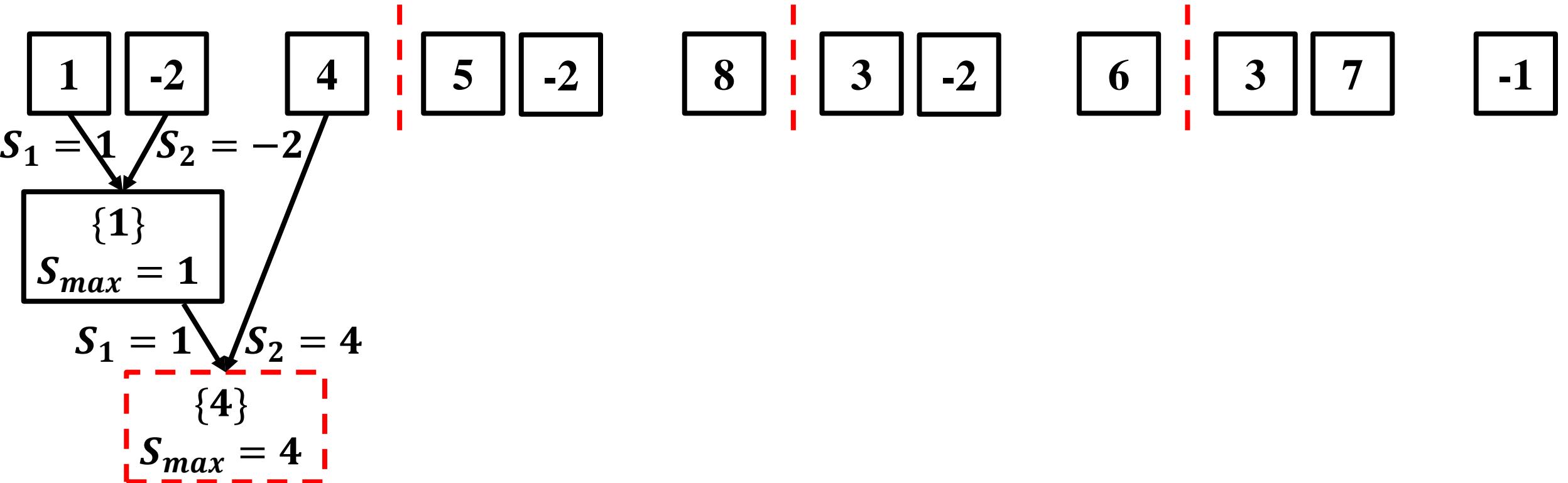
算法实例



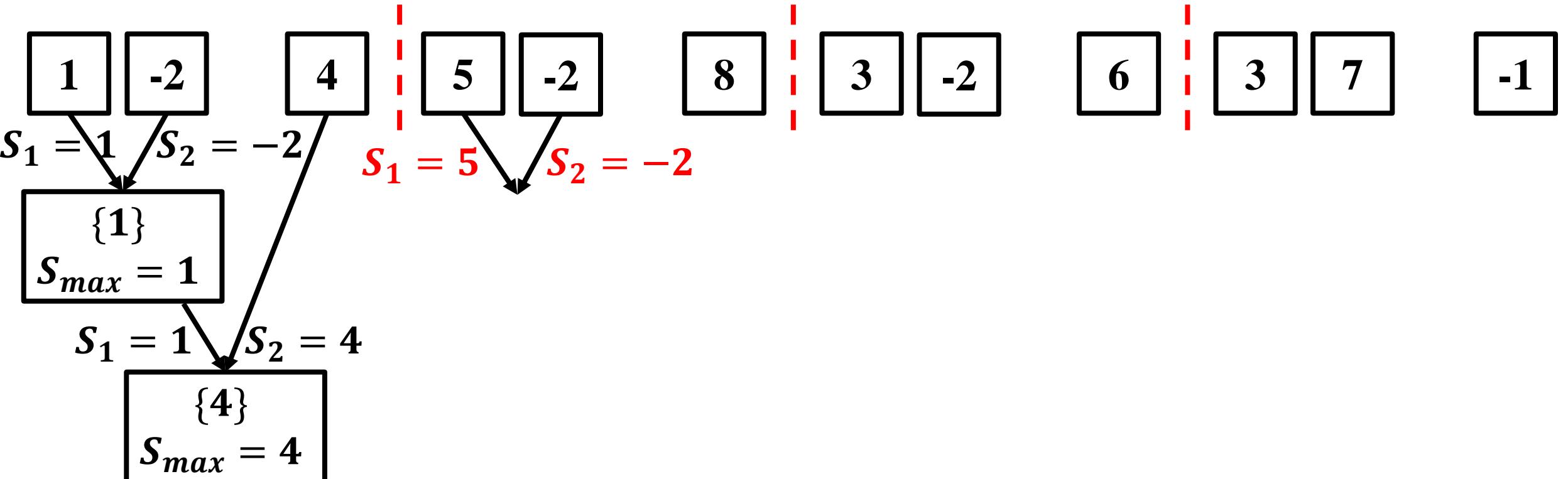
算法实例



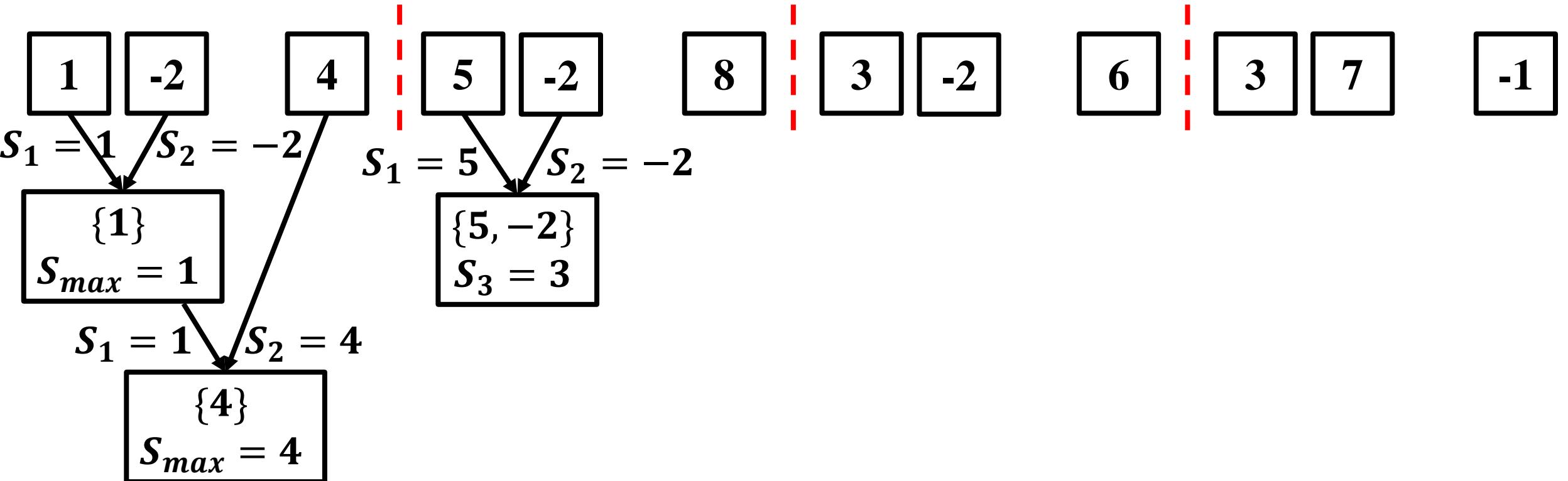
算法实例



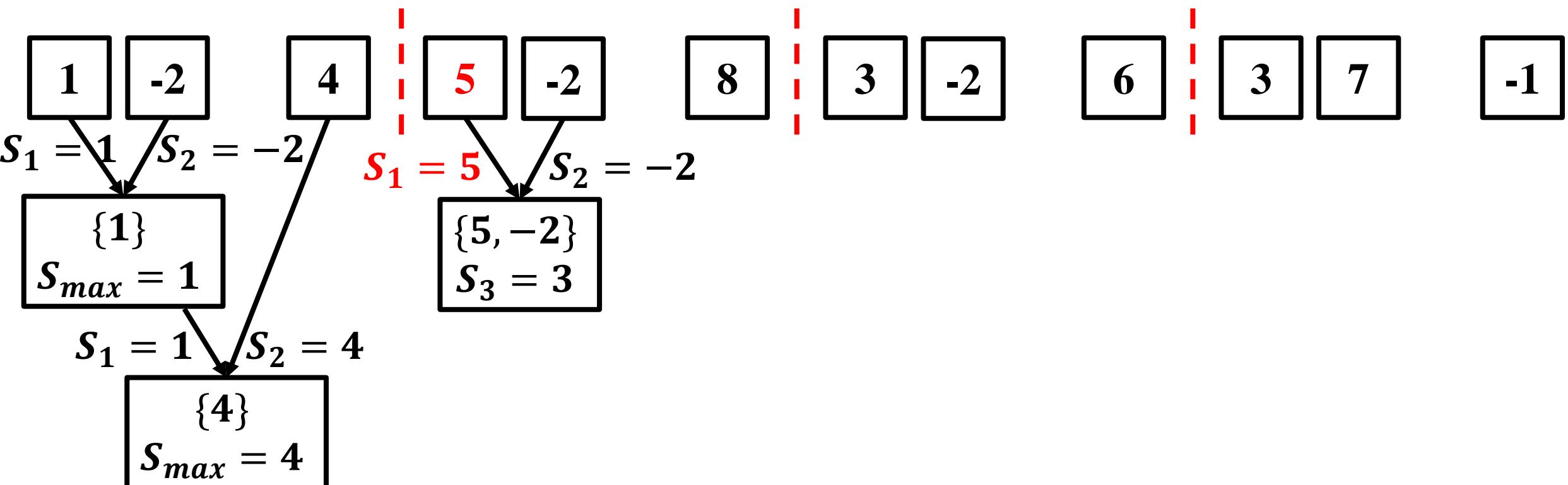
算法实例



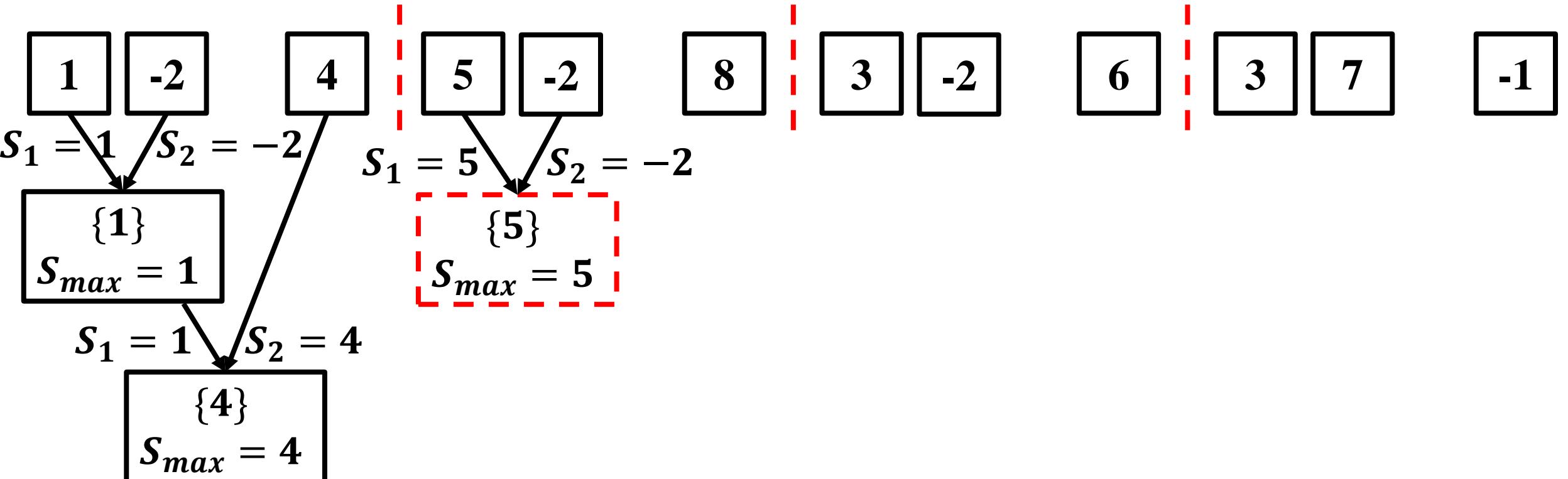
算法实例



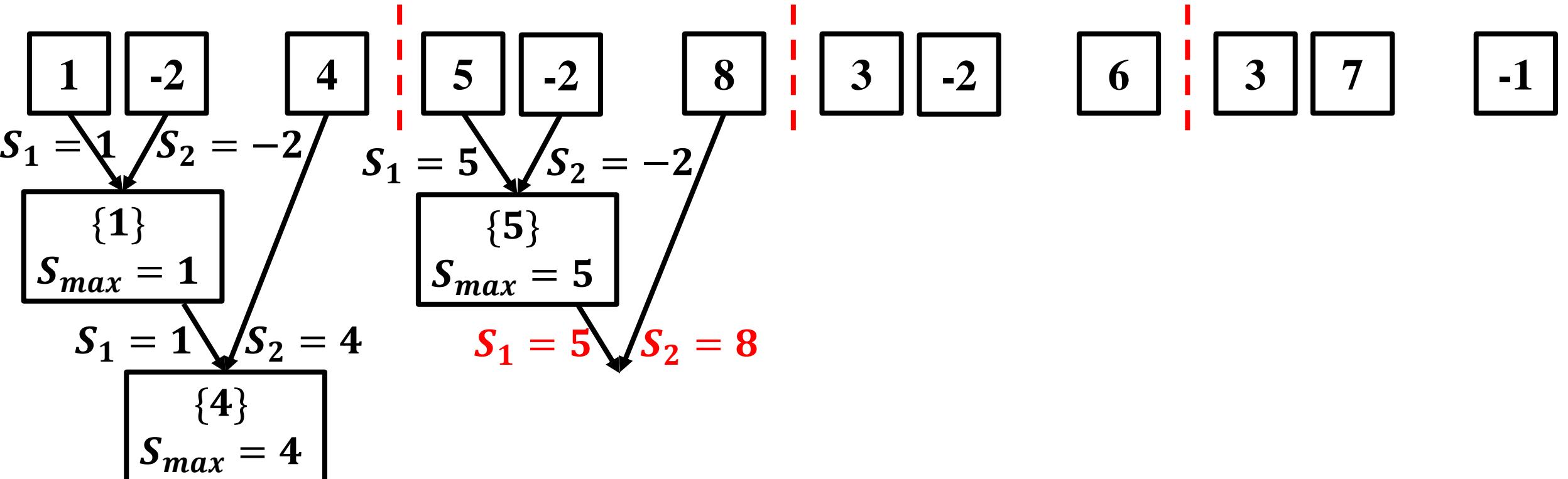
算法实例



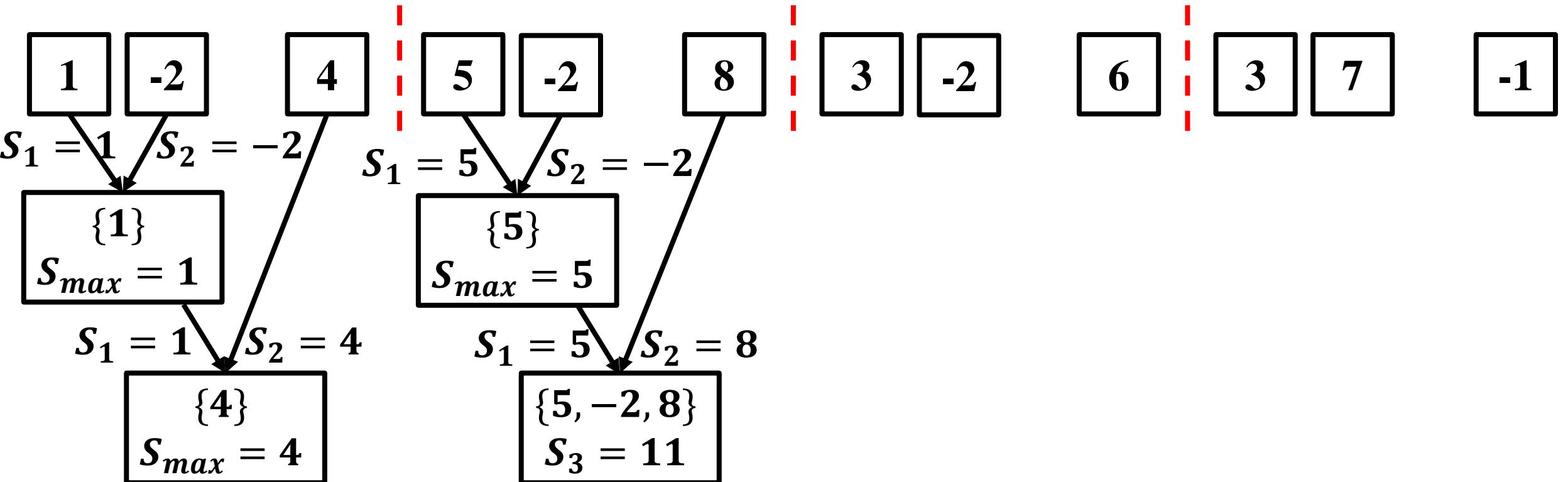
算法实例



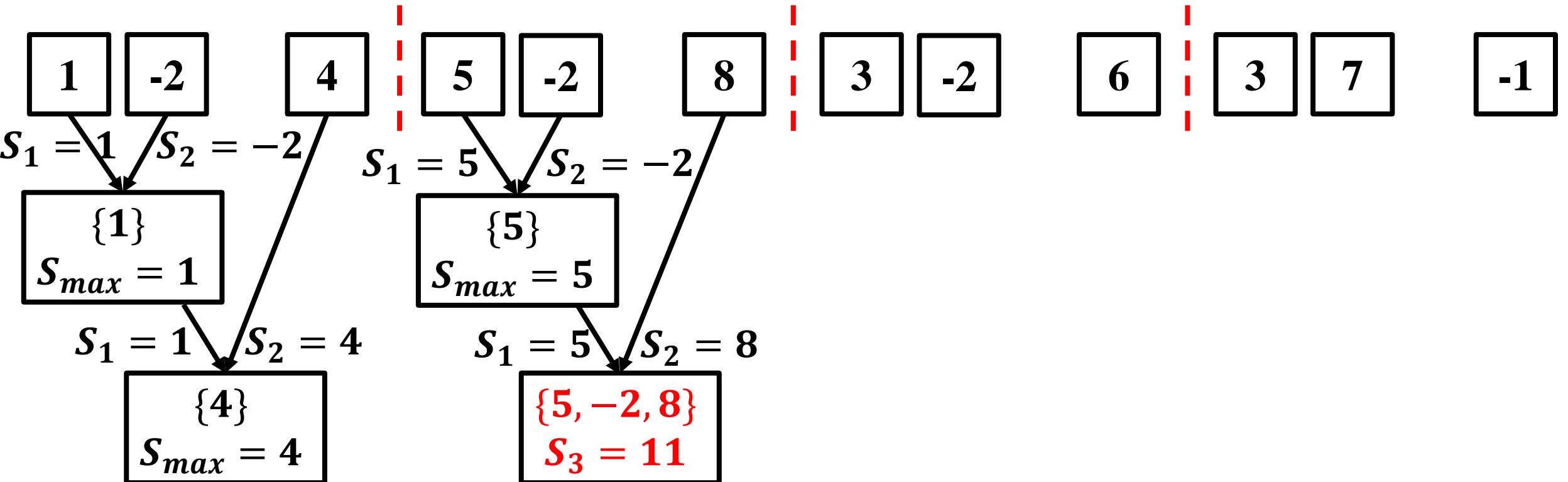
算法实例



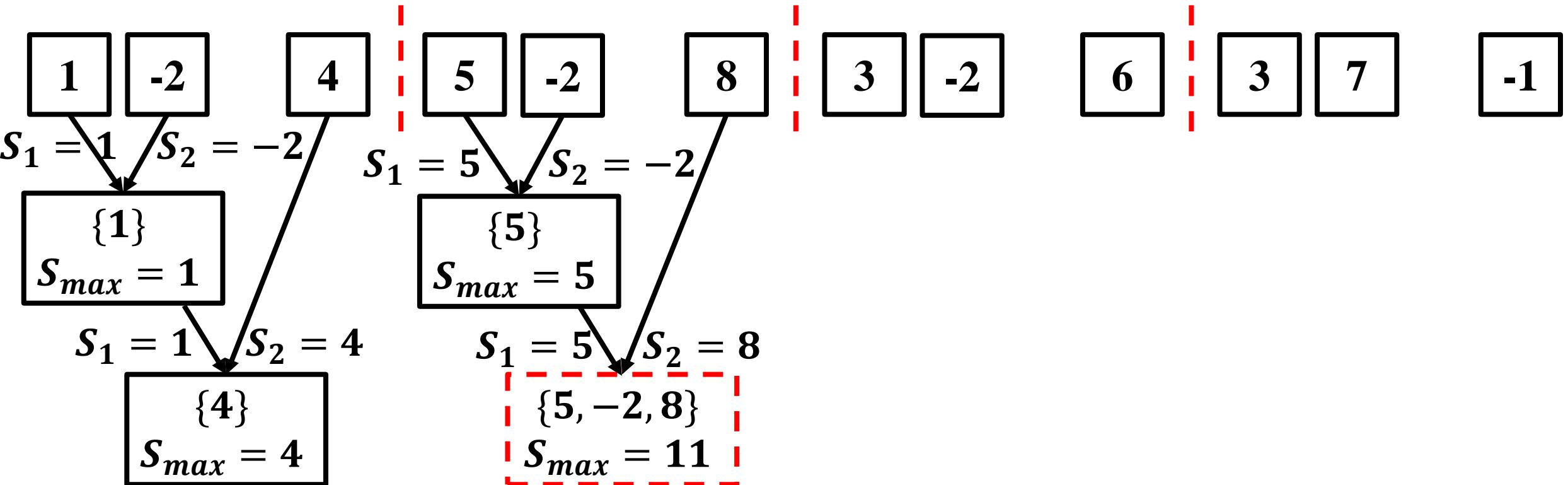
算法实例



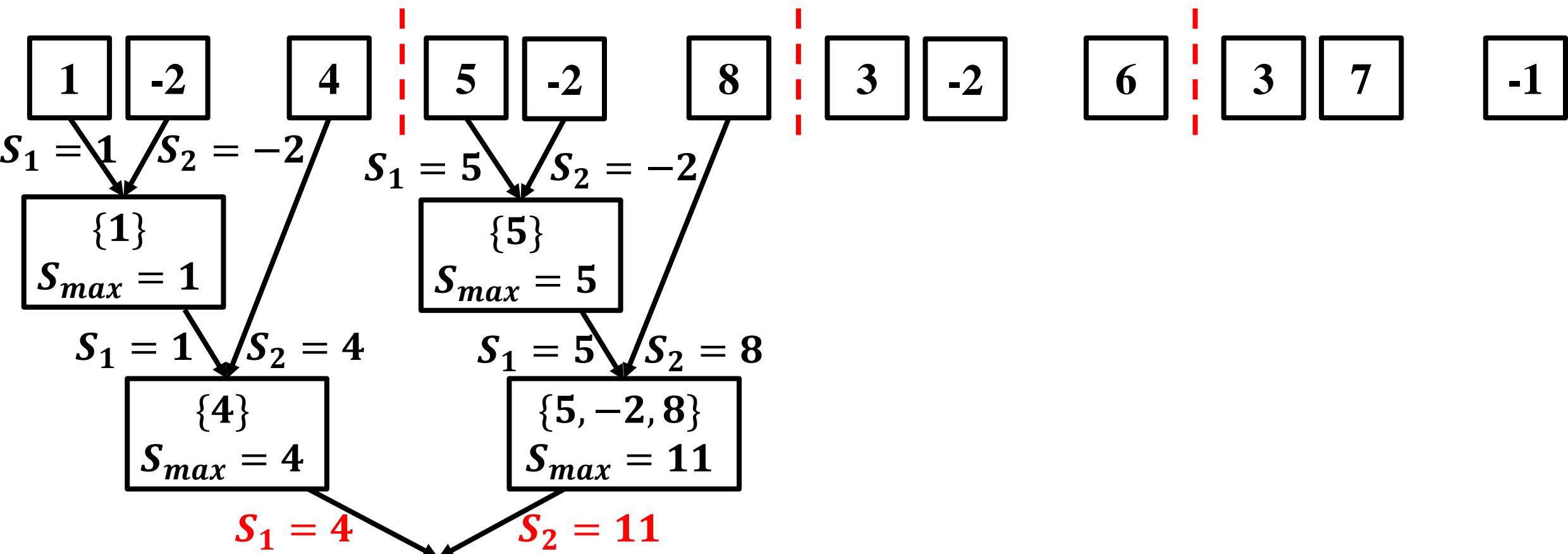
算法实例



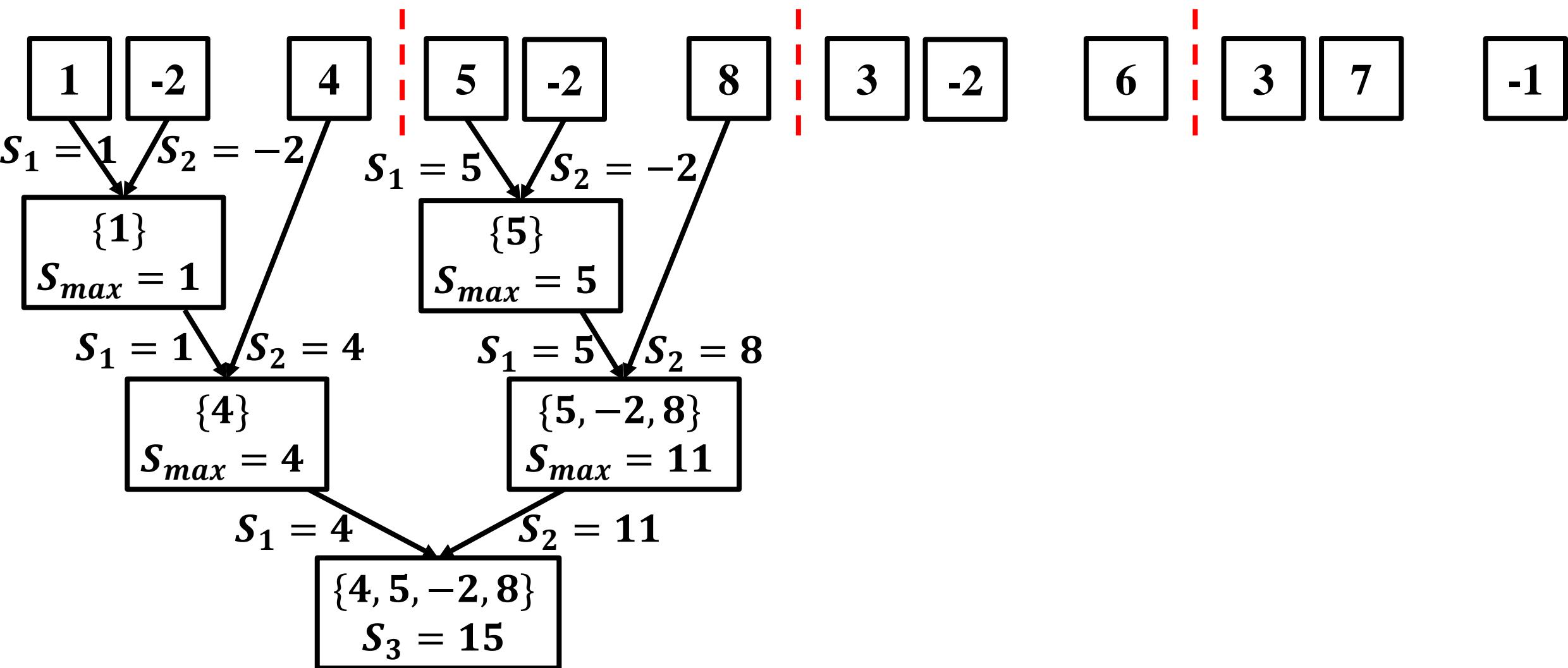
算法实例



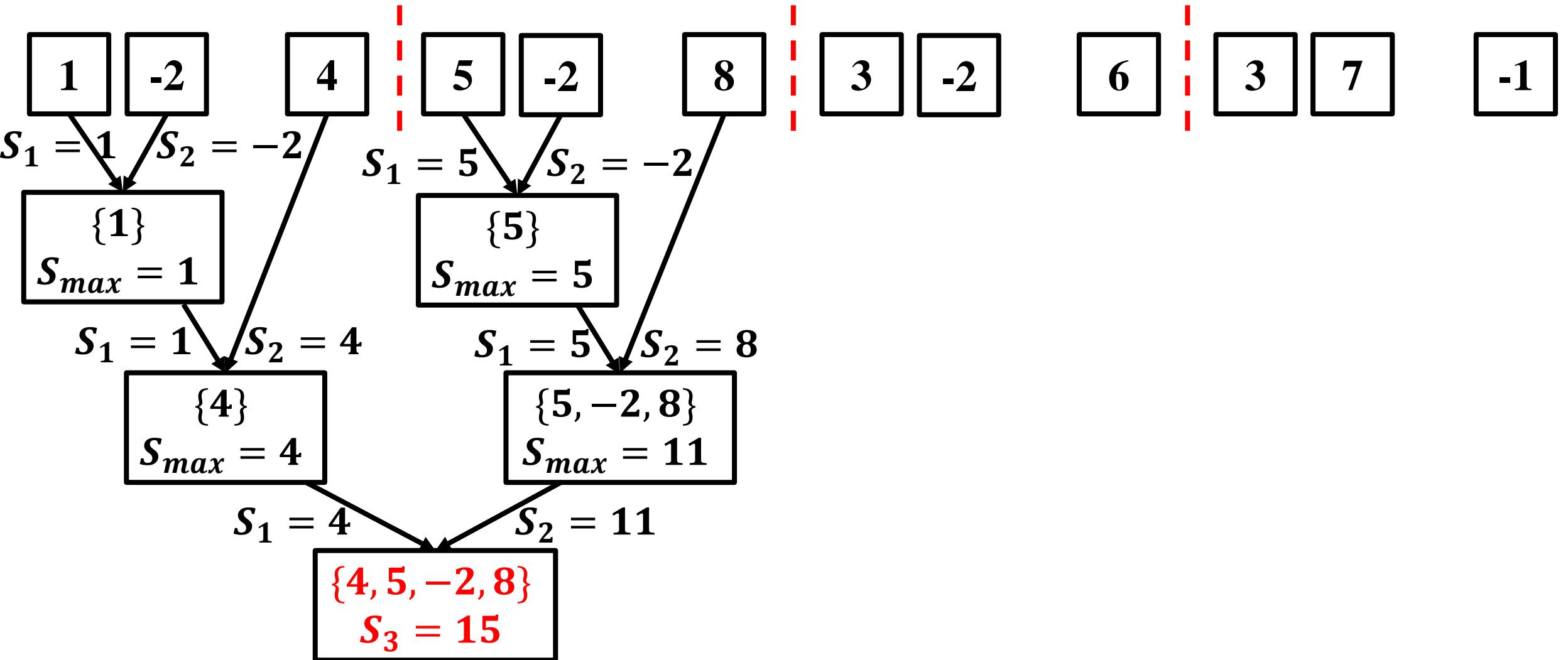
算法实例



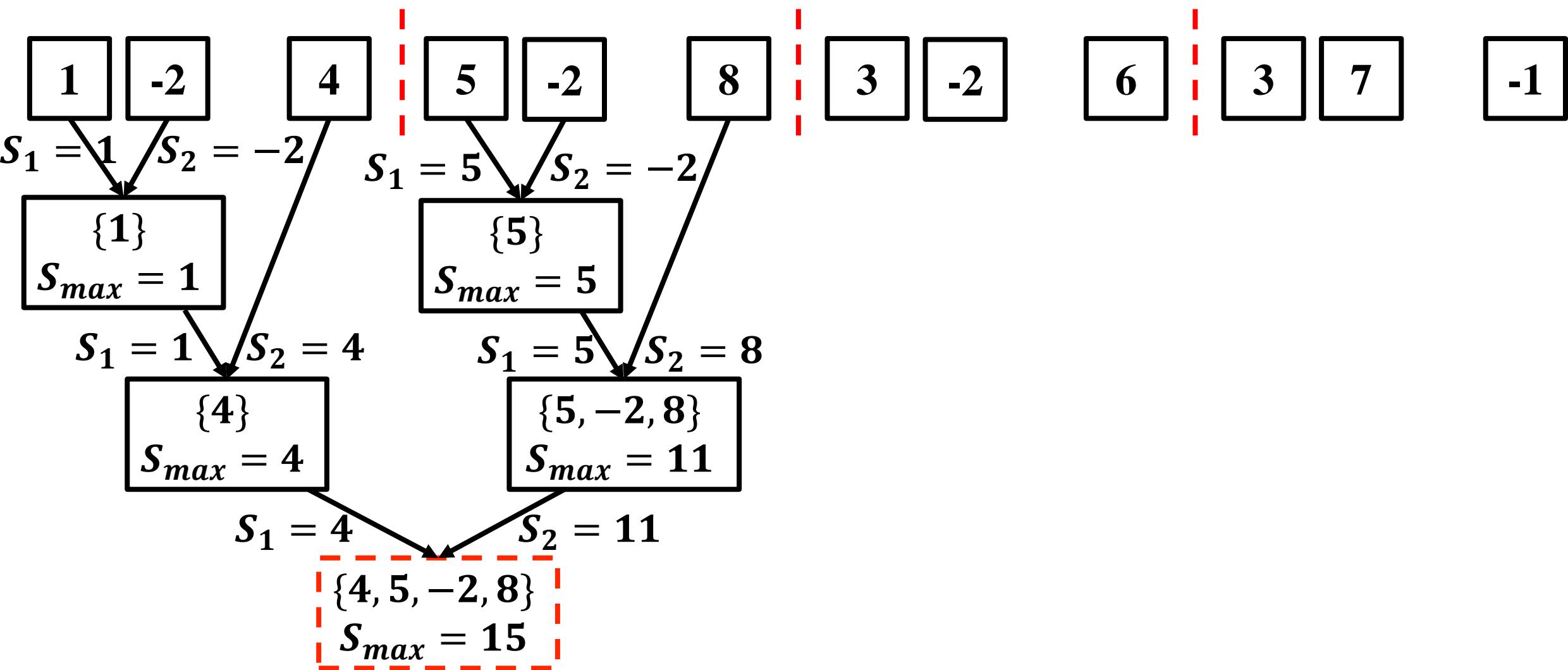
算法实例



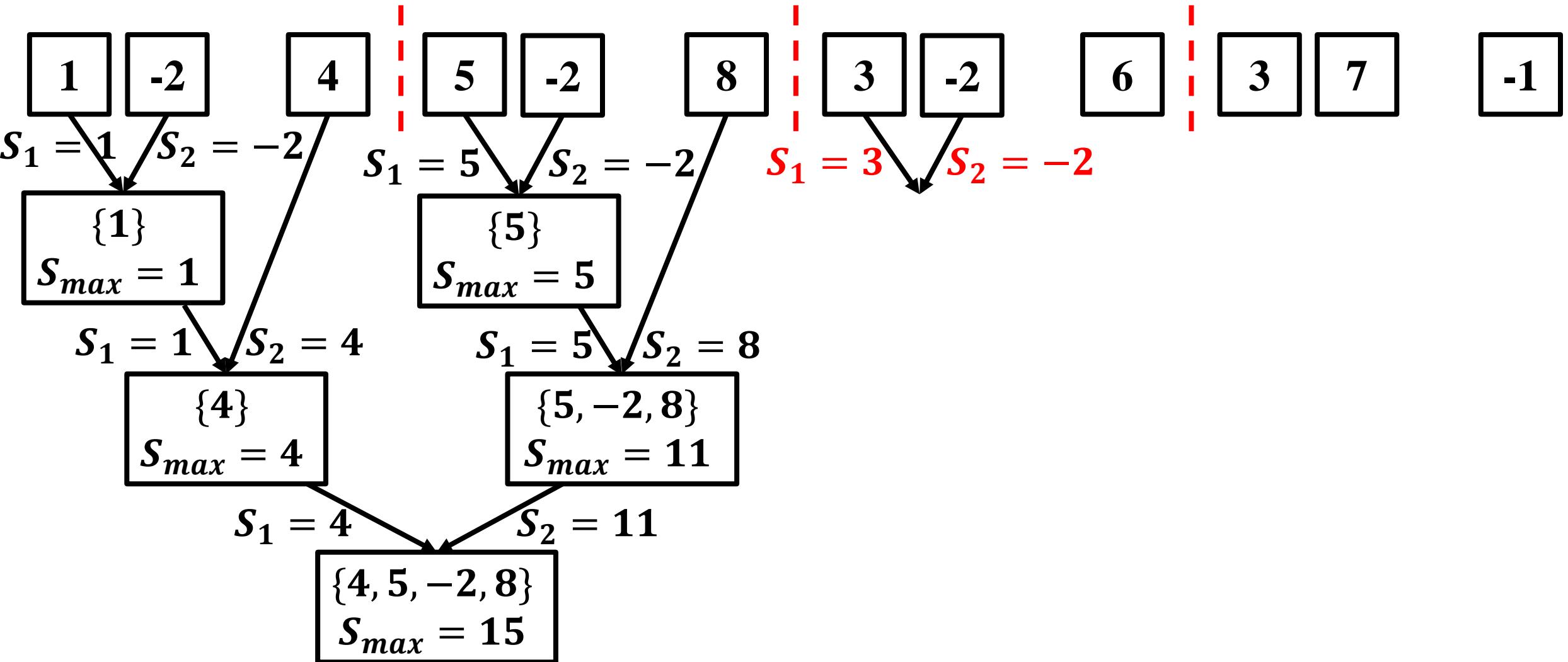
算法实例



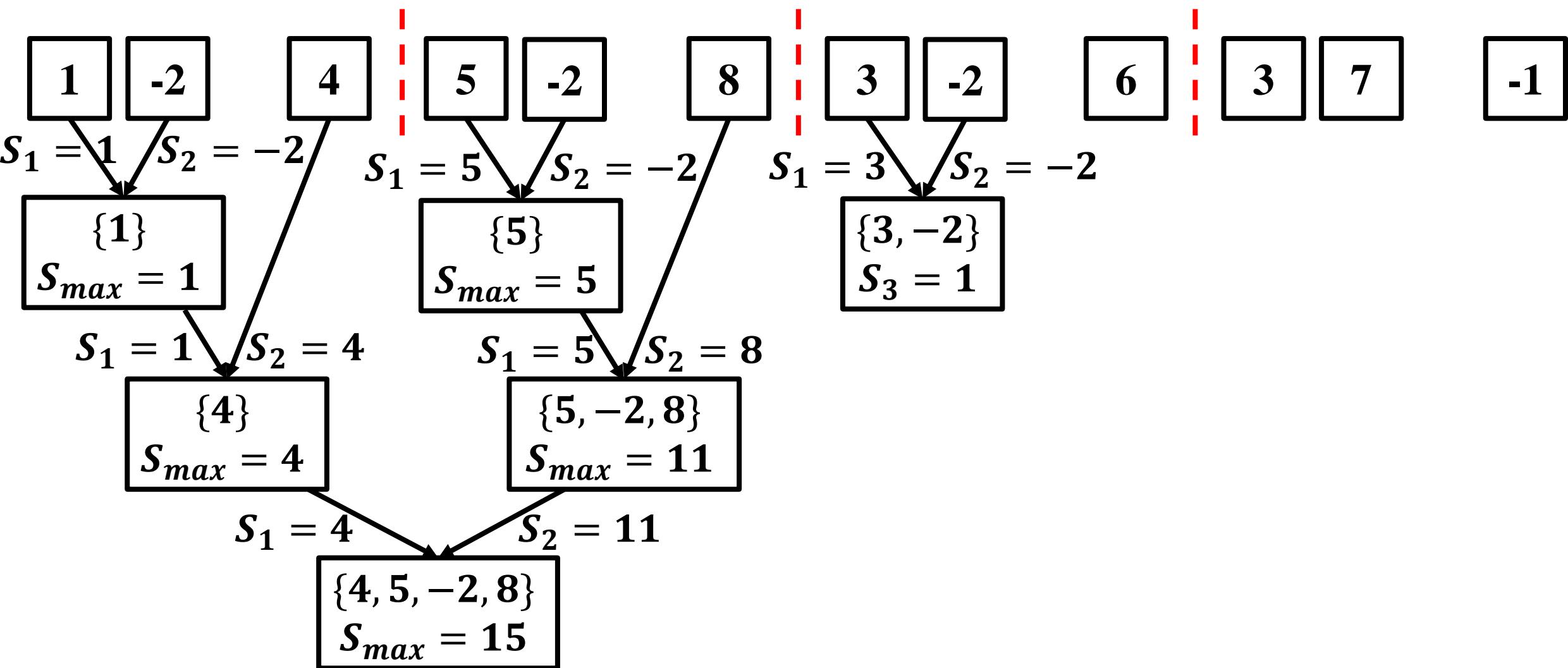
算法实例



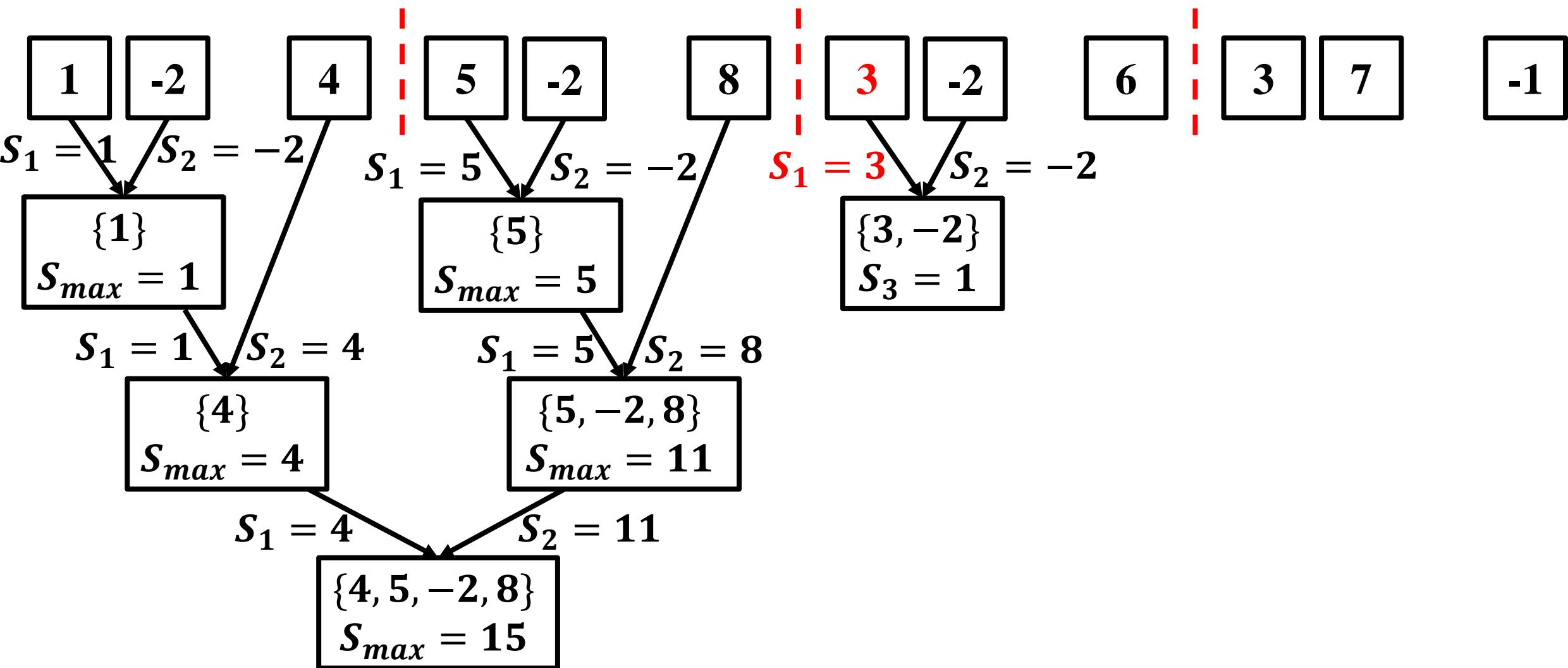
算法实例



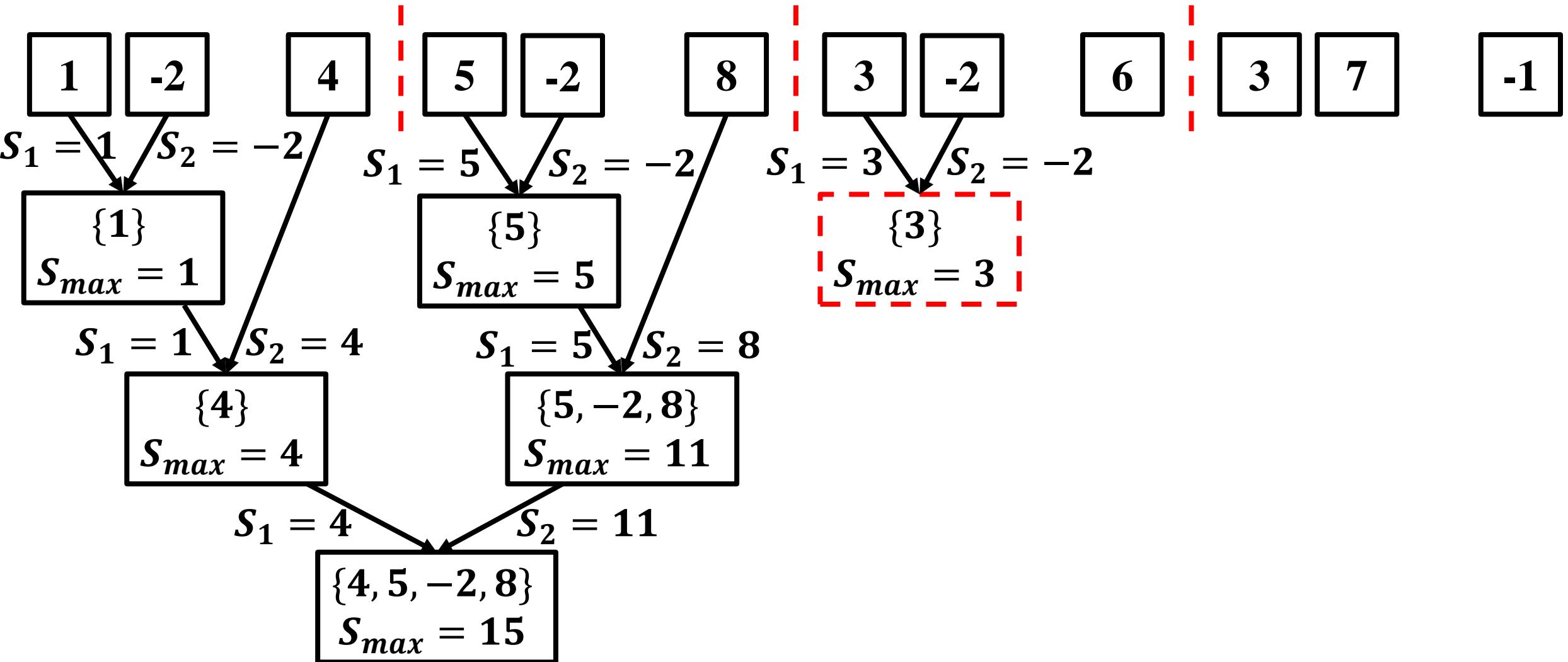
算法实例



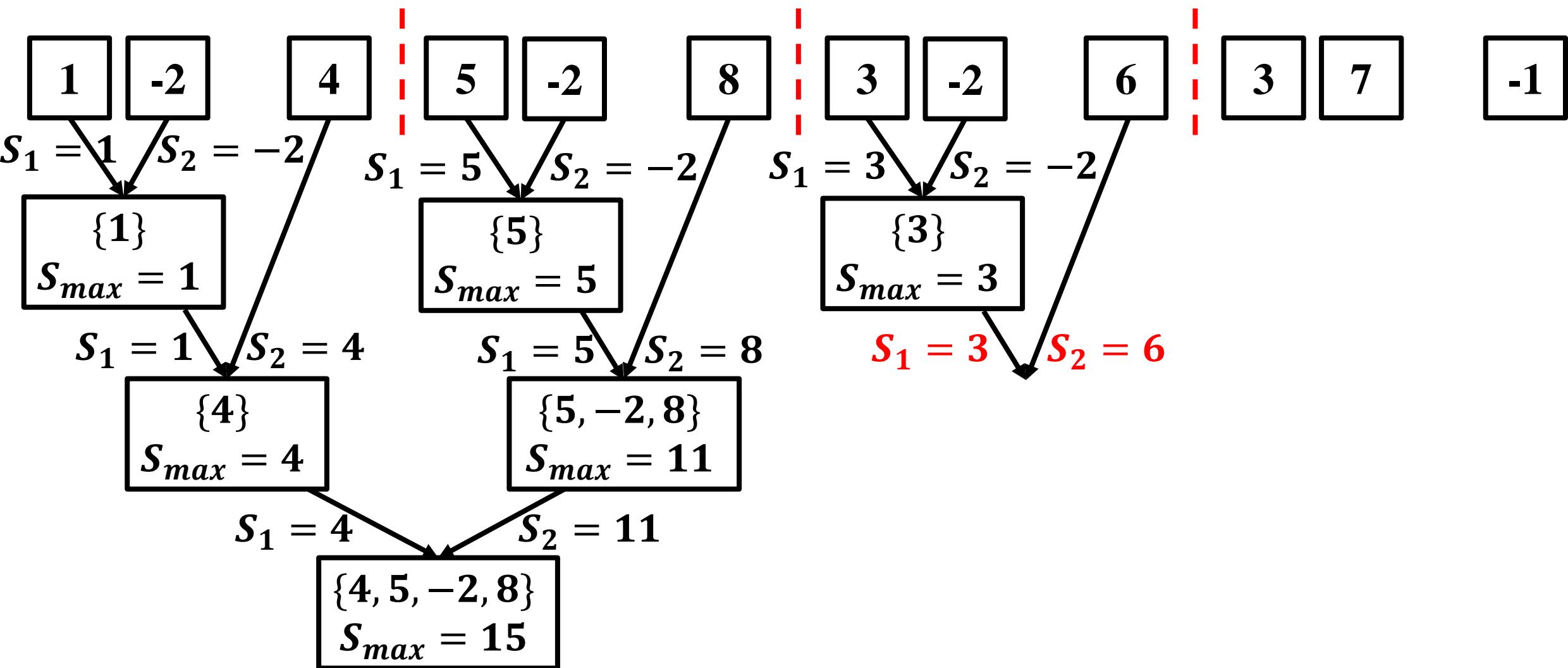
算法实例



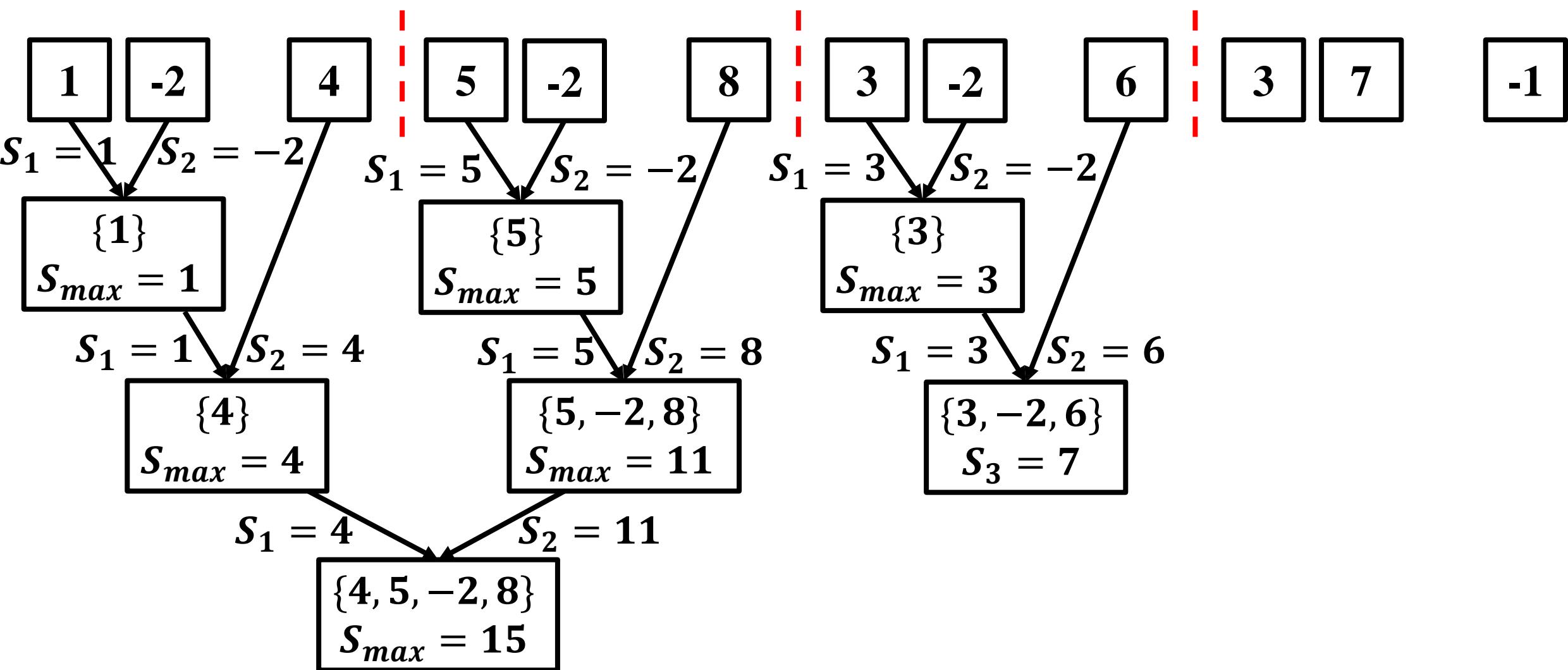
算法实例



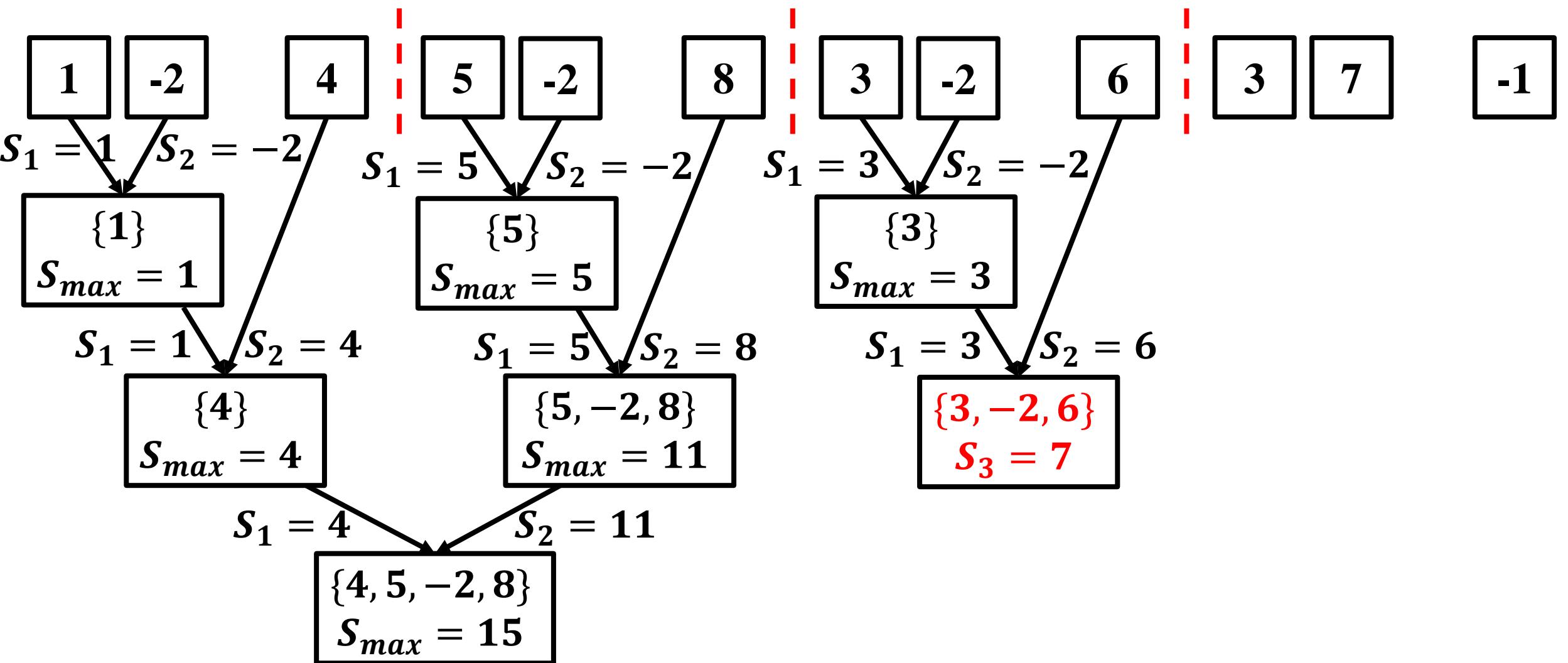
算法实例



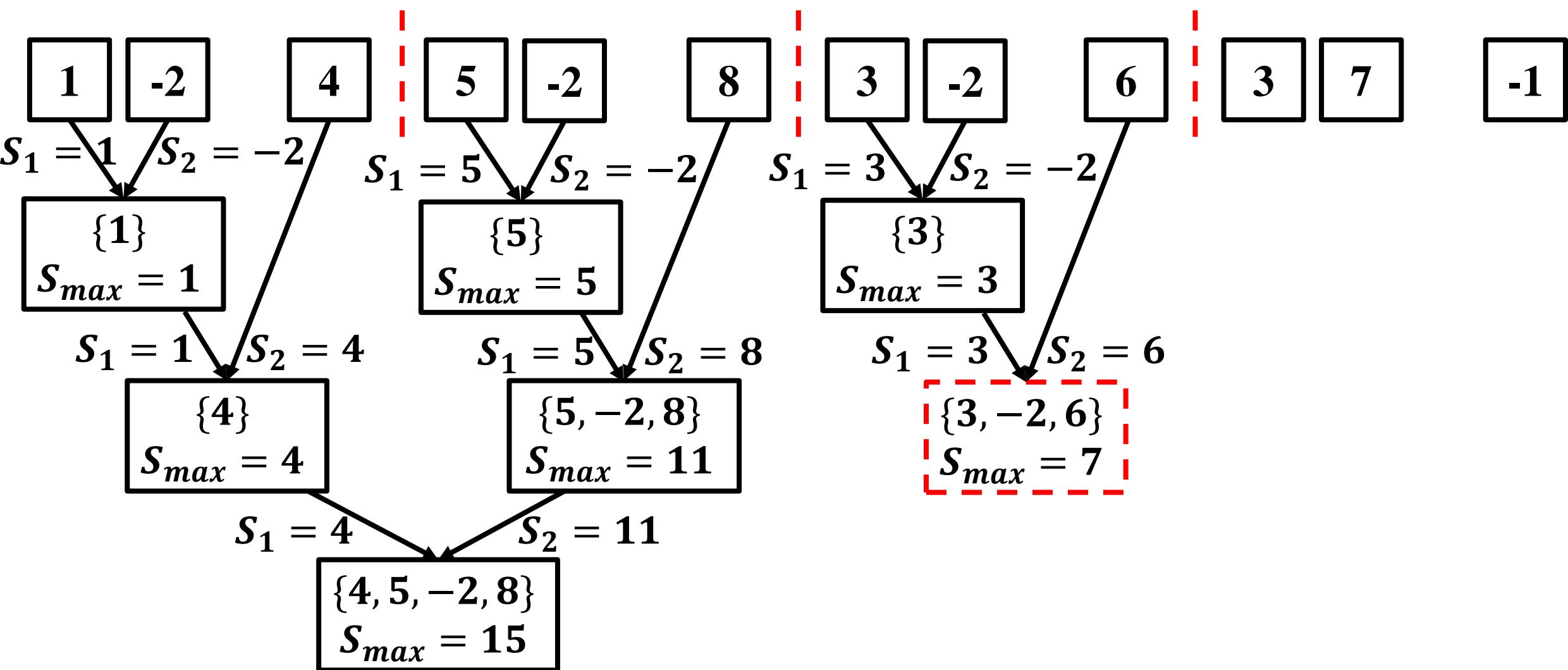
算法实例



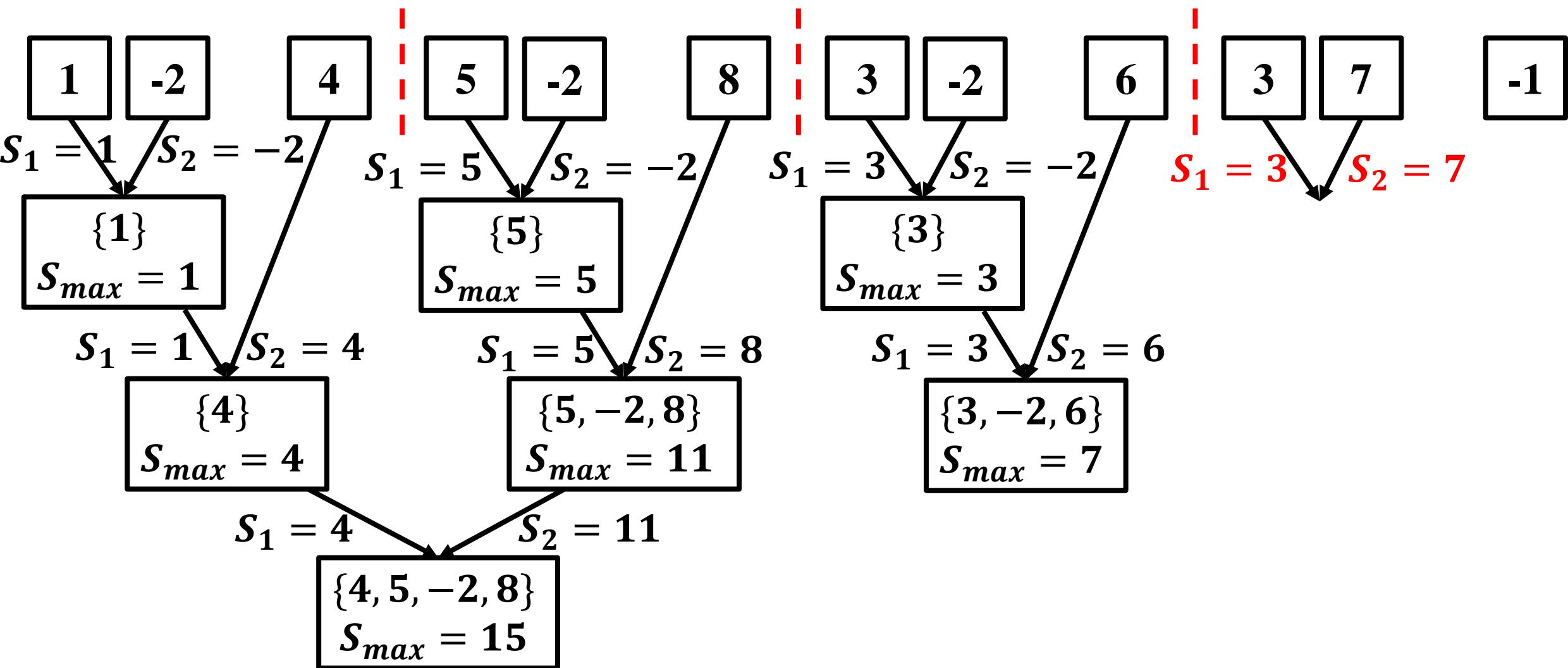
算法实例



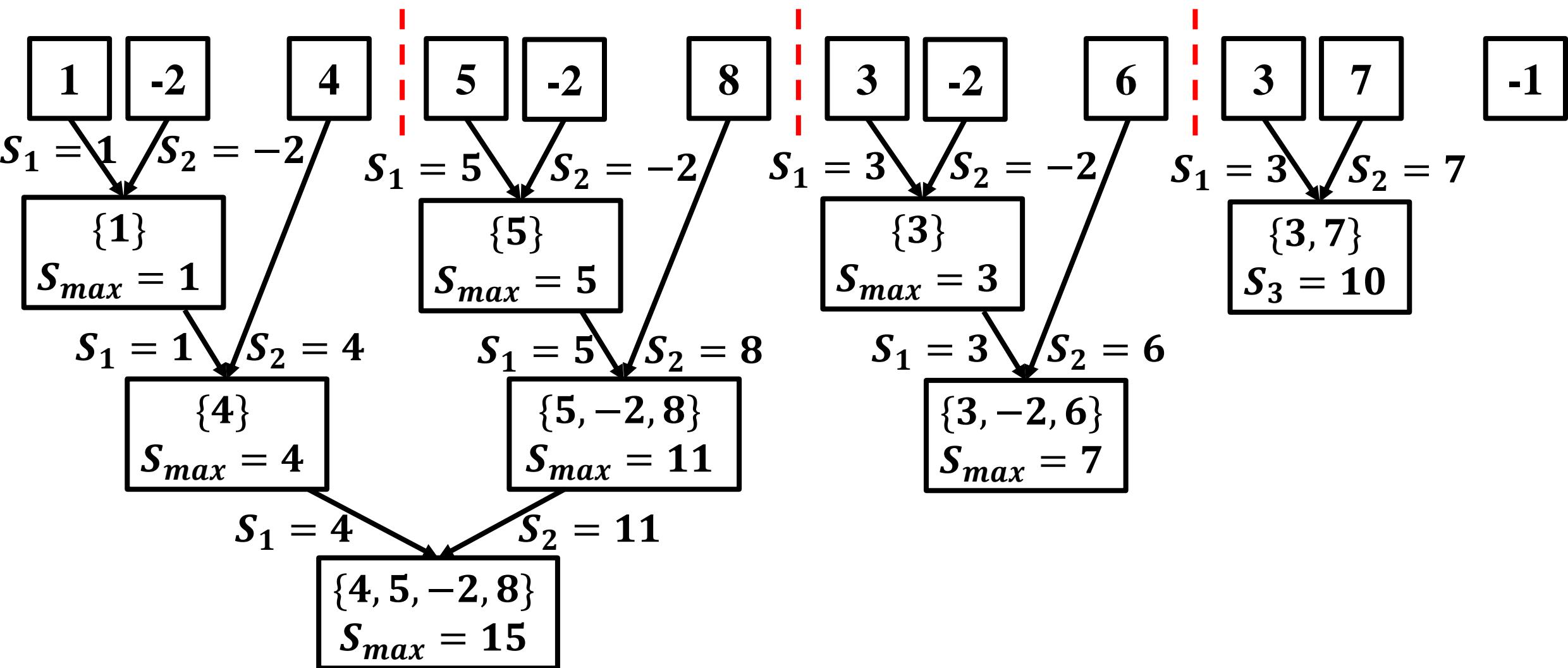
算法实例



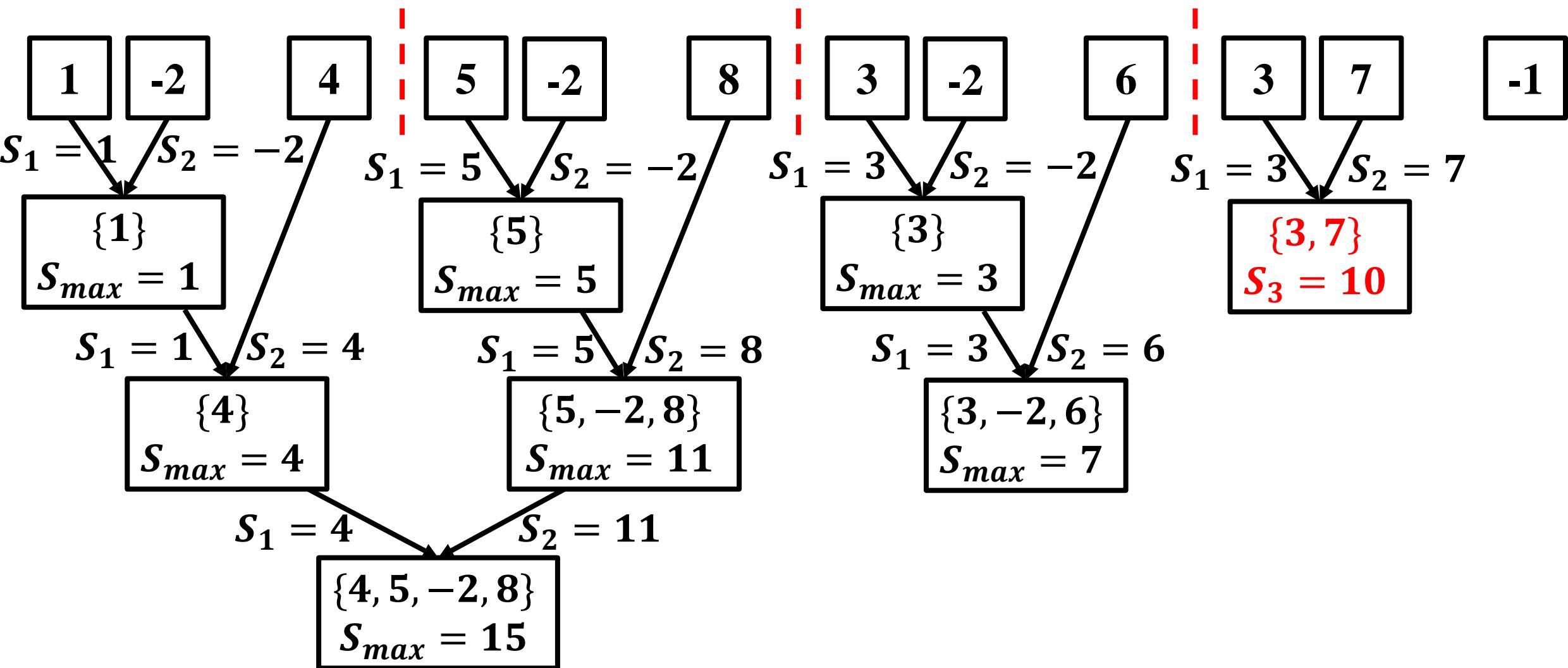
算法实例



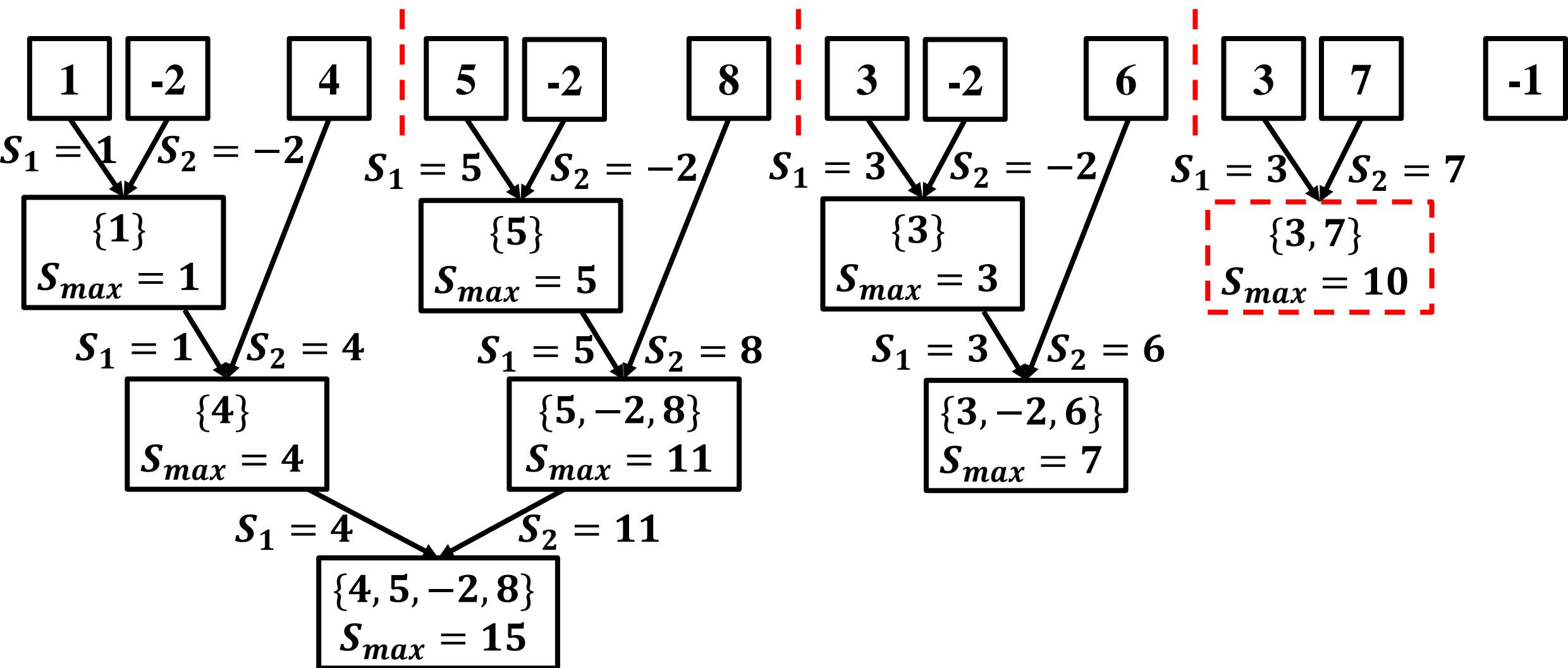
算法实例



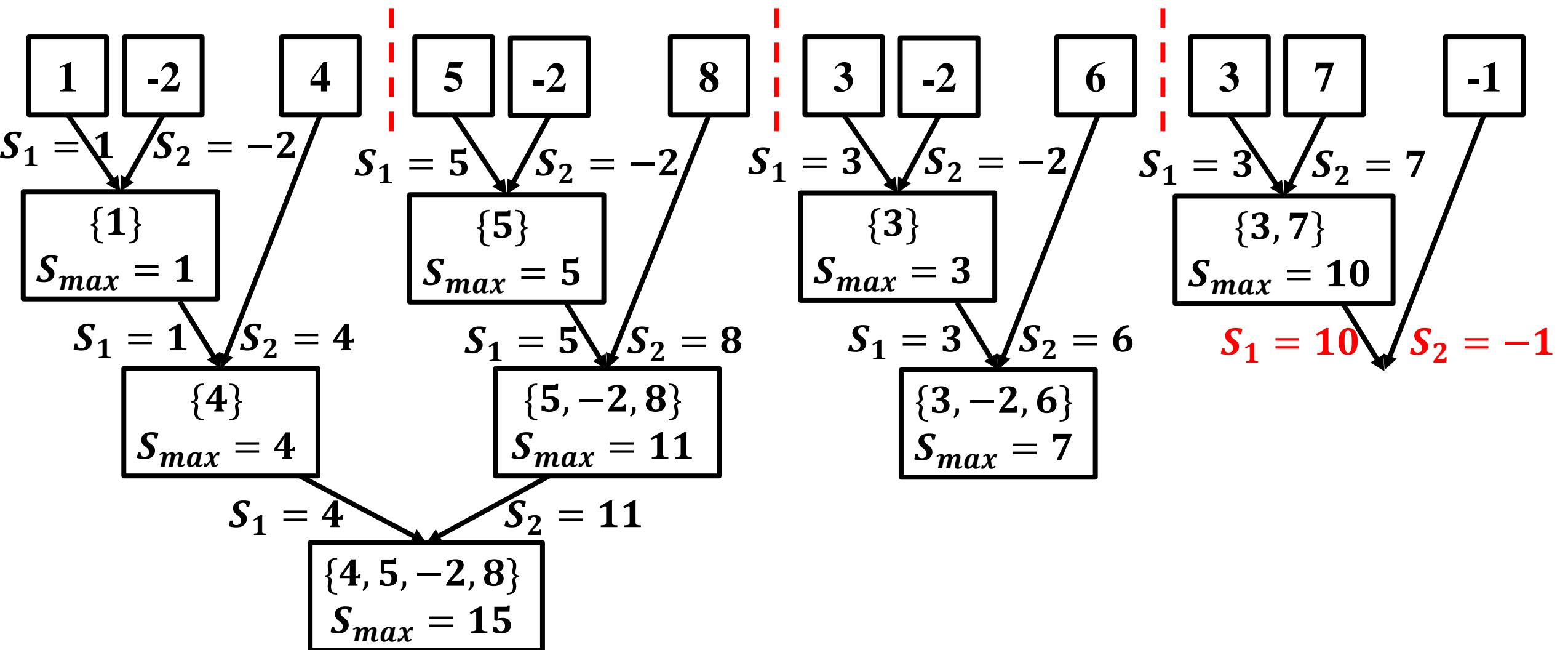
算法实例



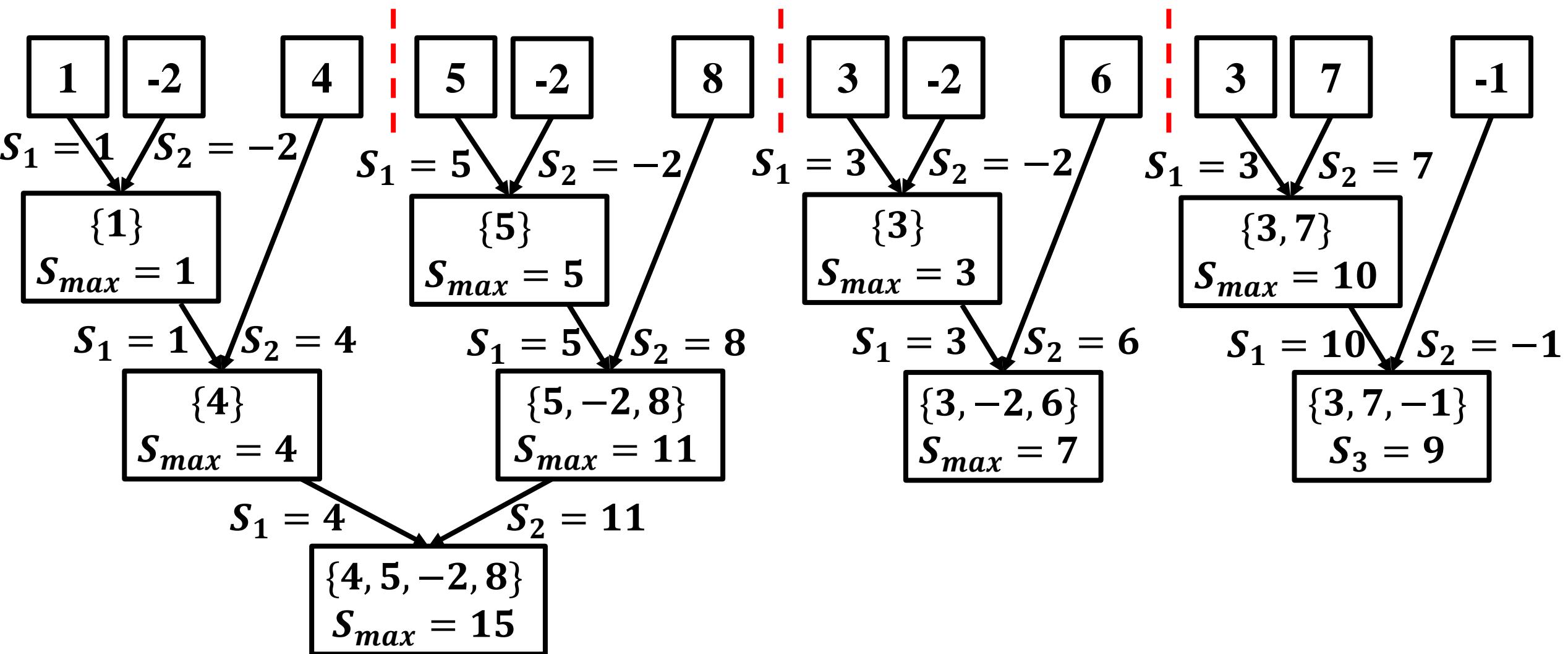
算法实例



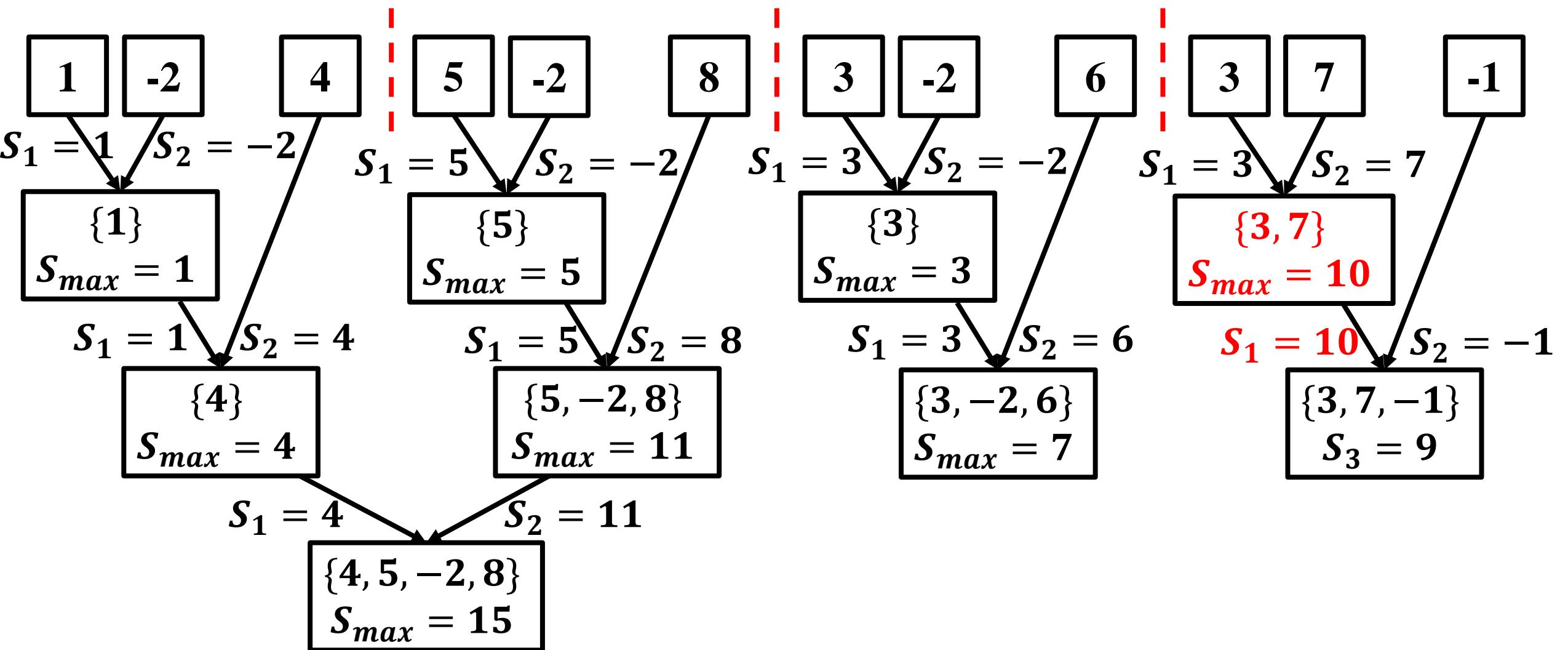
算法实例



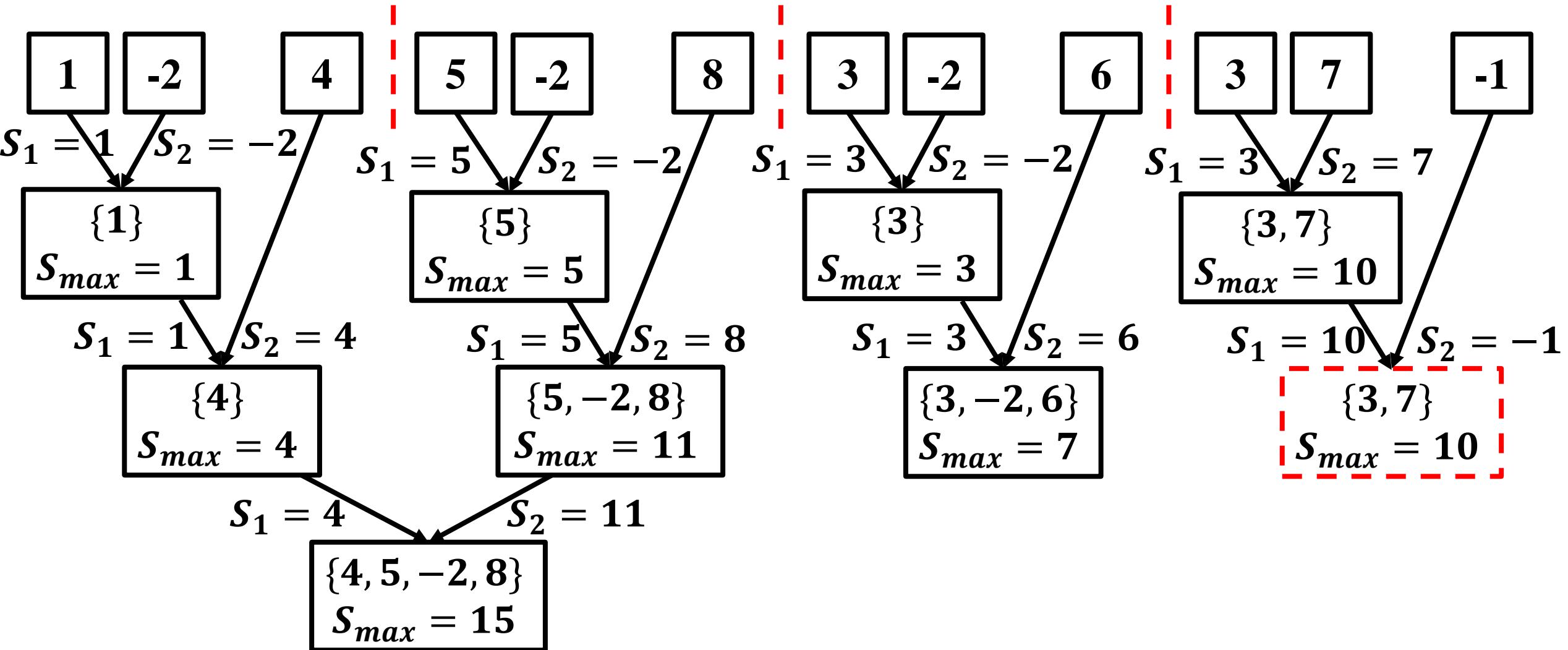
算法实例



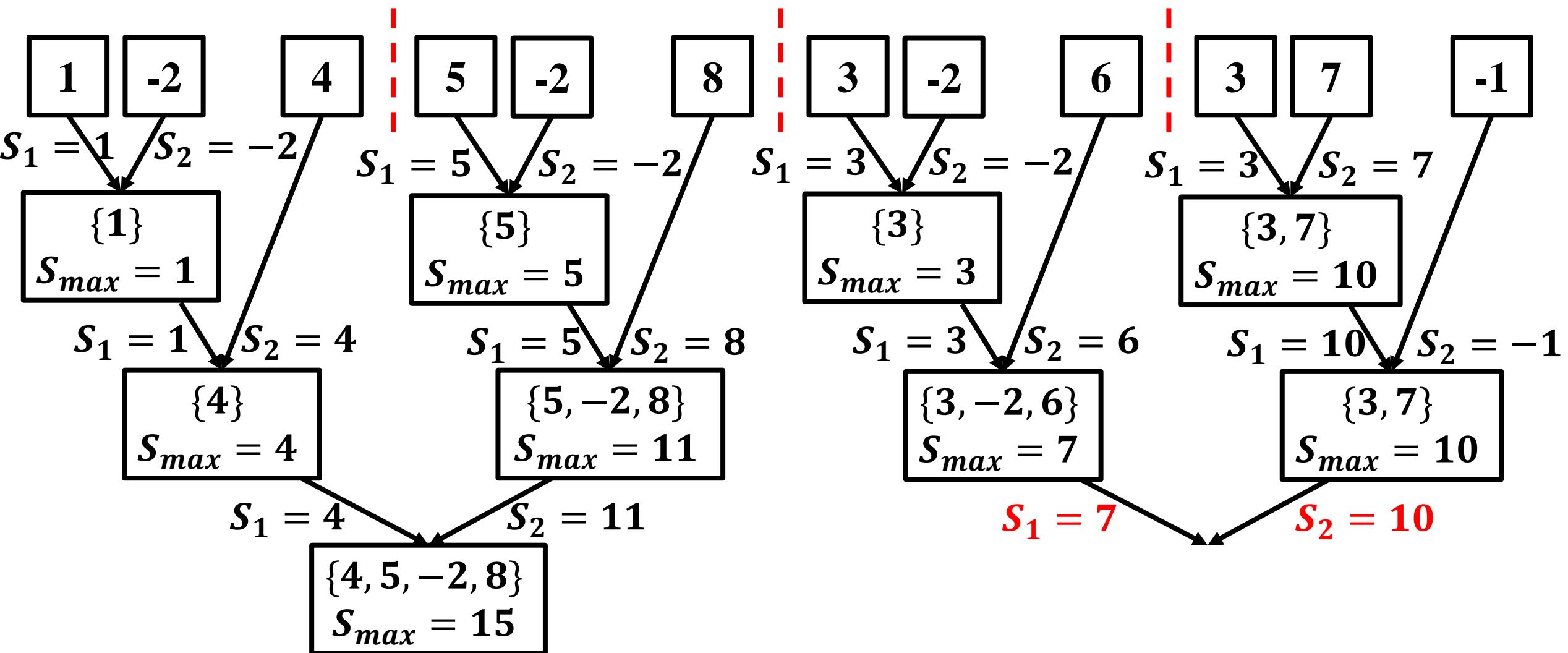
算法实例



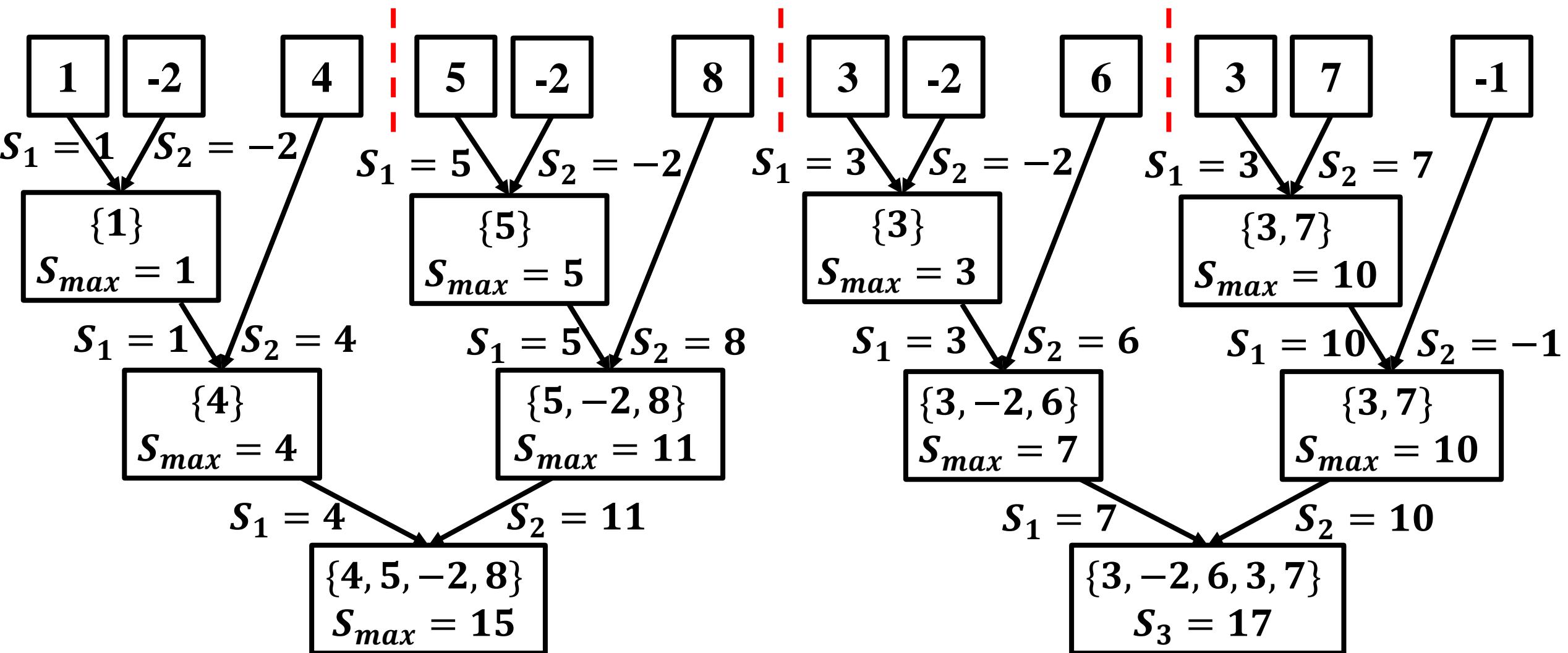
算法实例



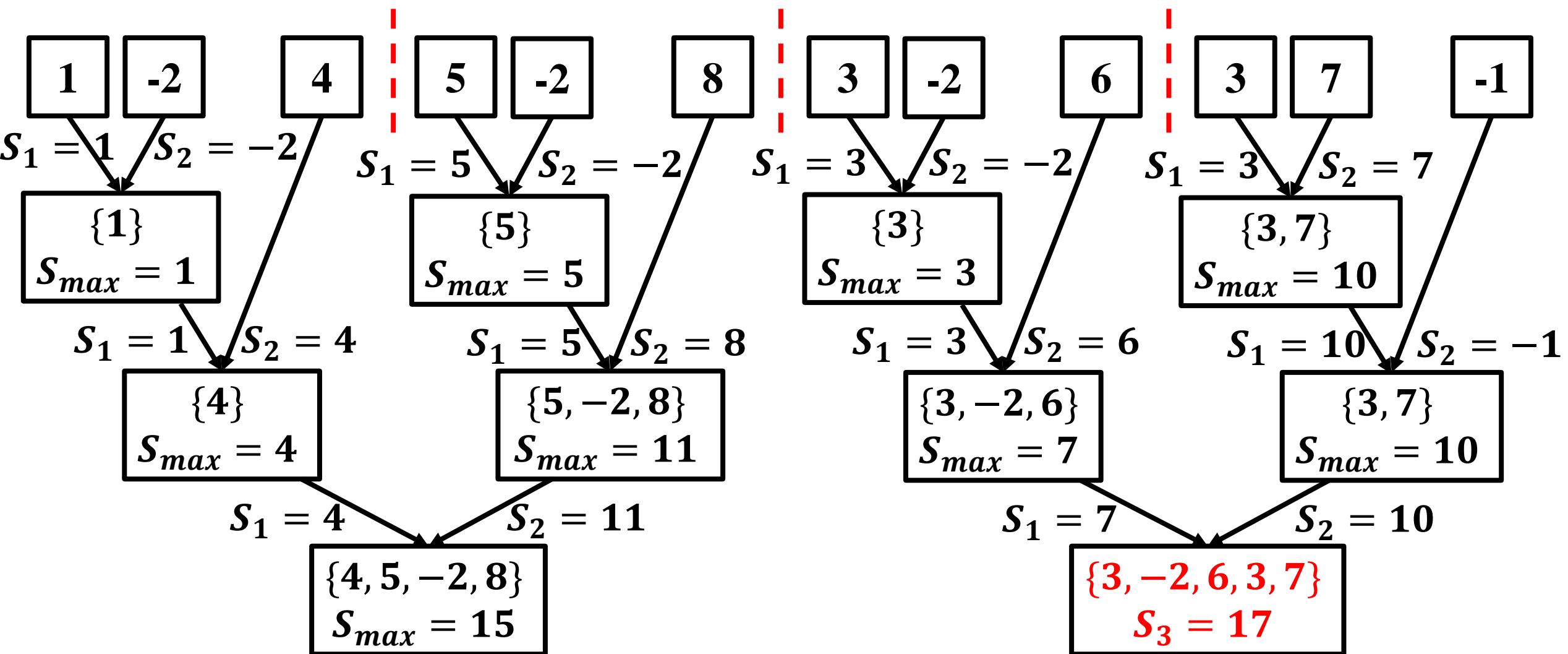
算法实例



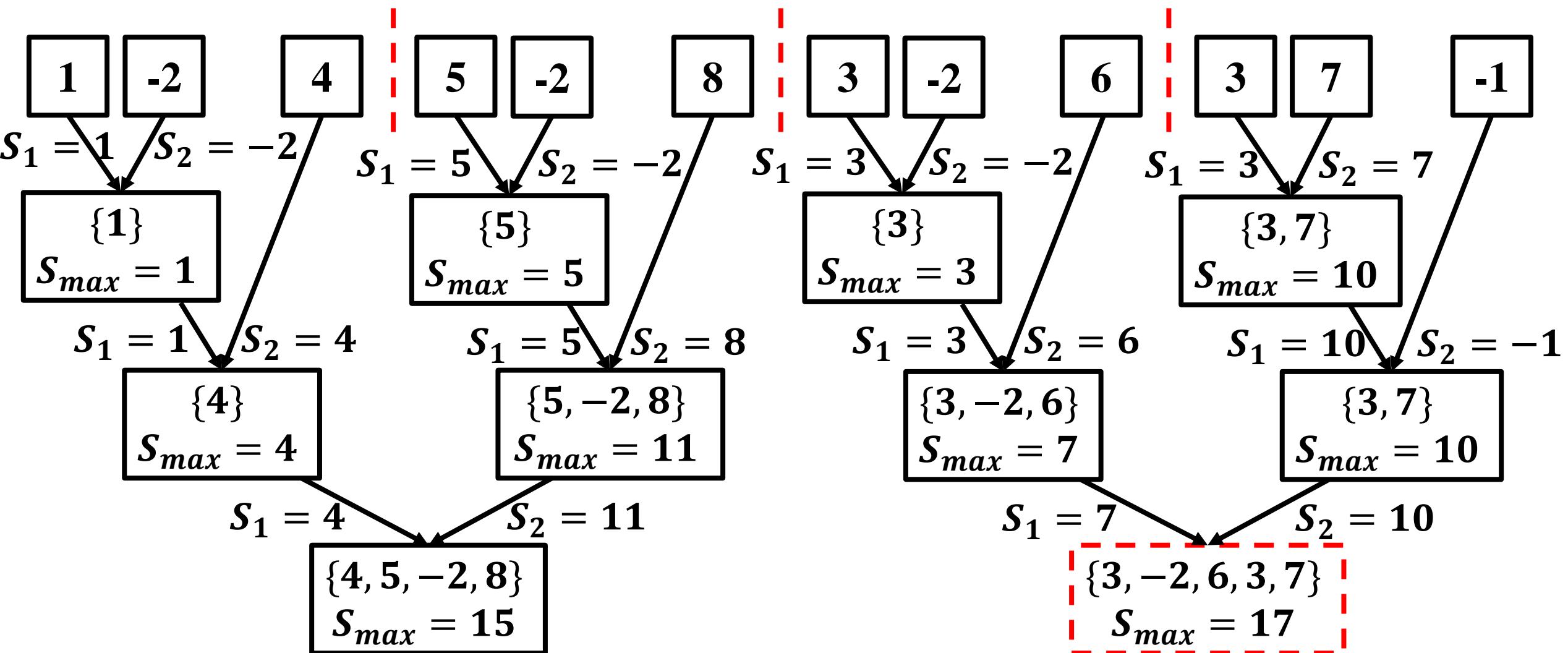
算法实例



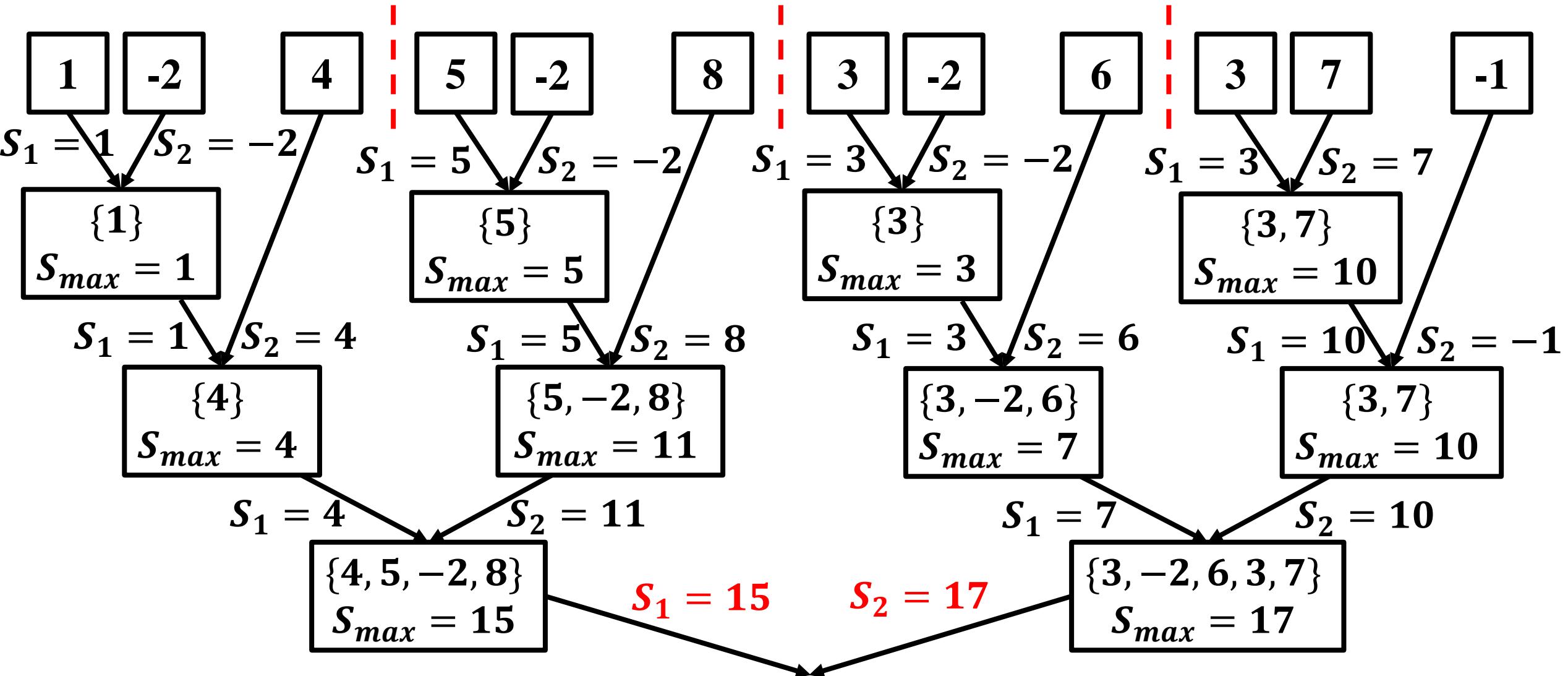
算法实例



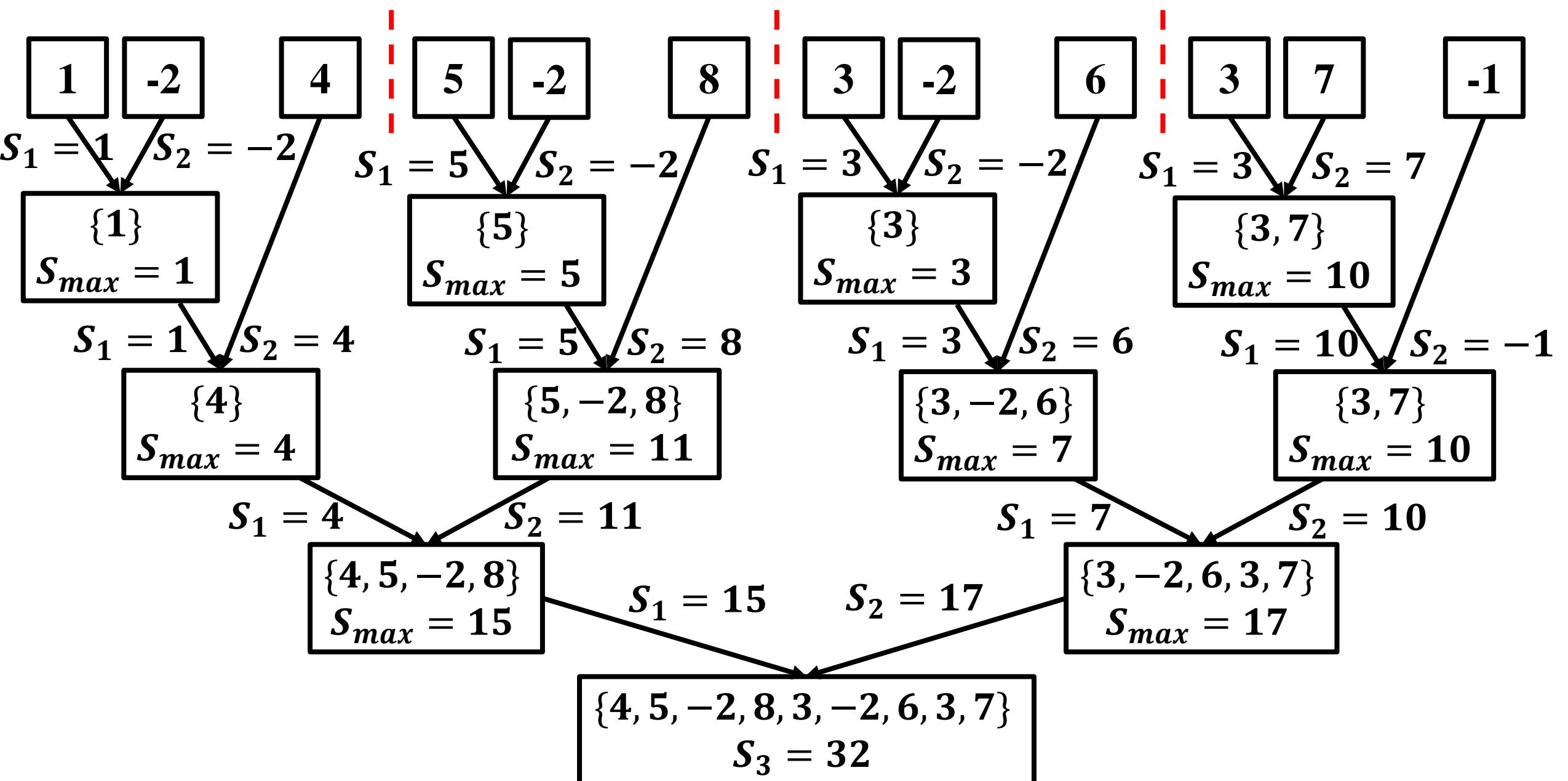
算法实例



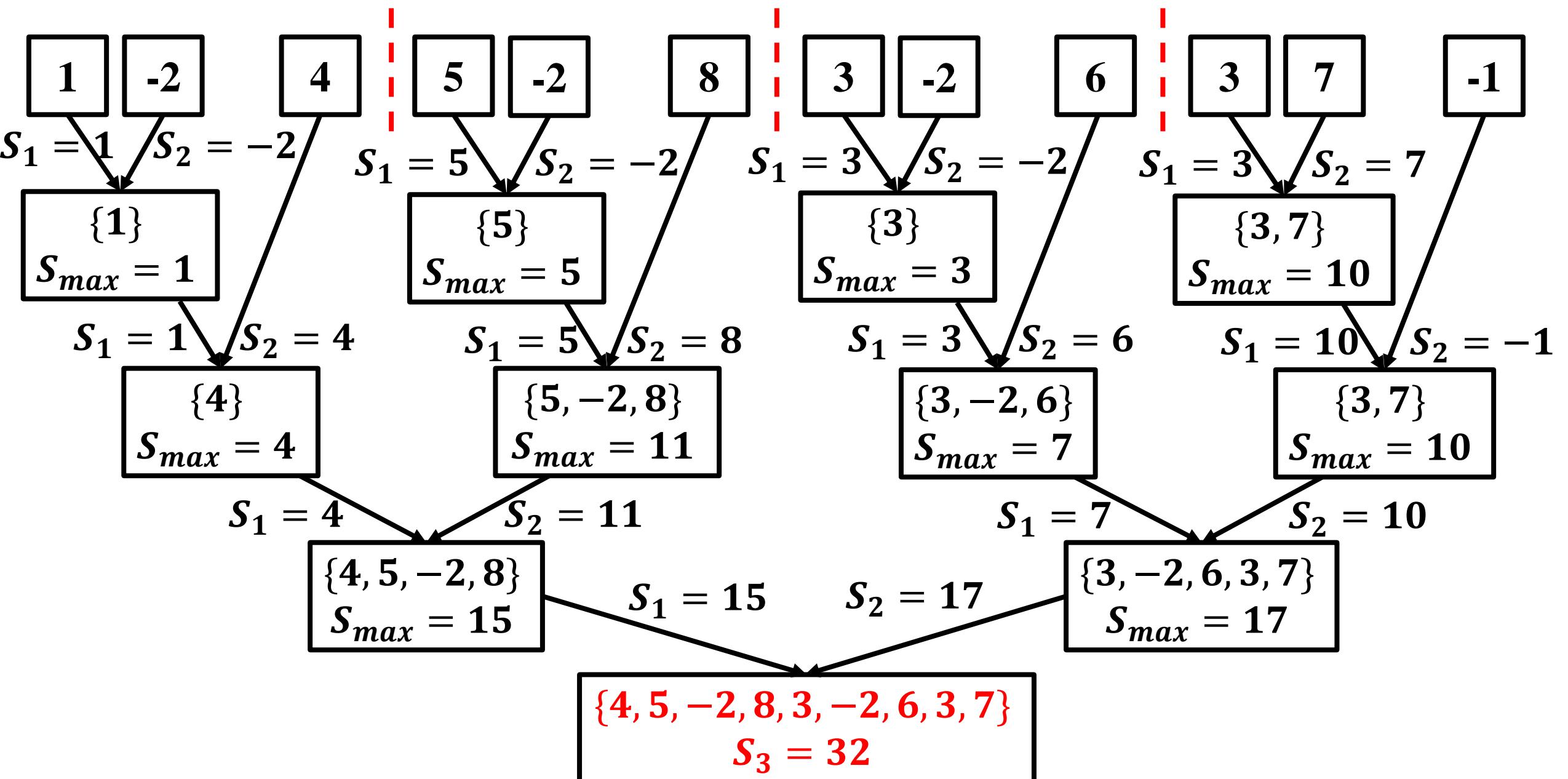
算法实例



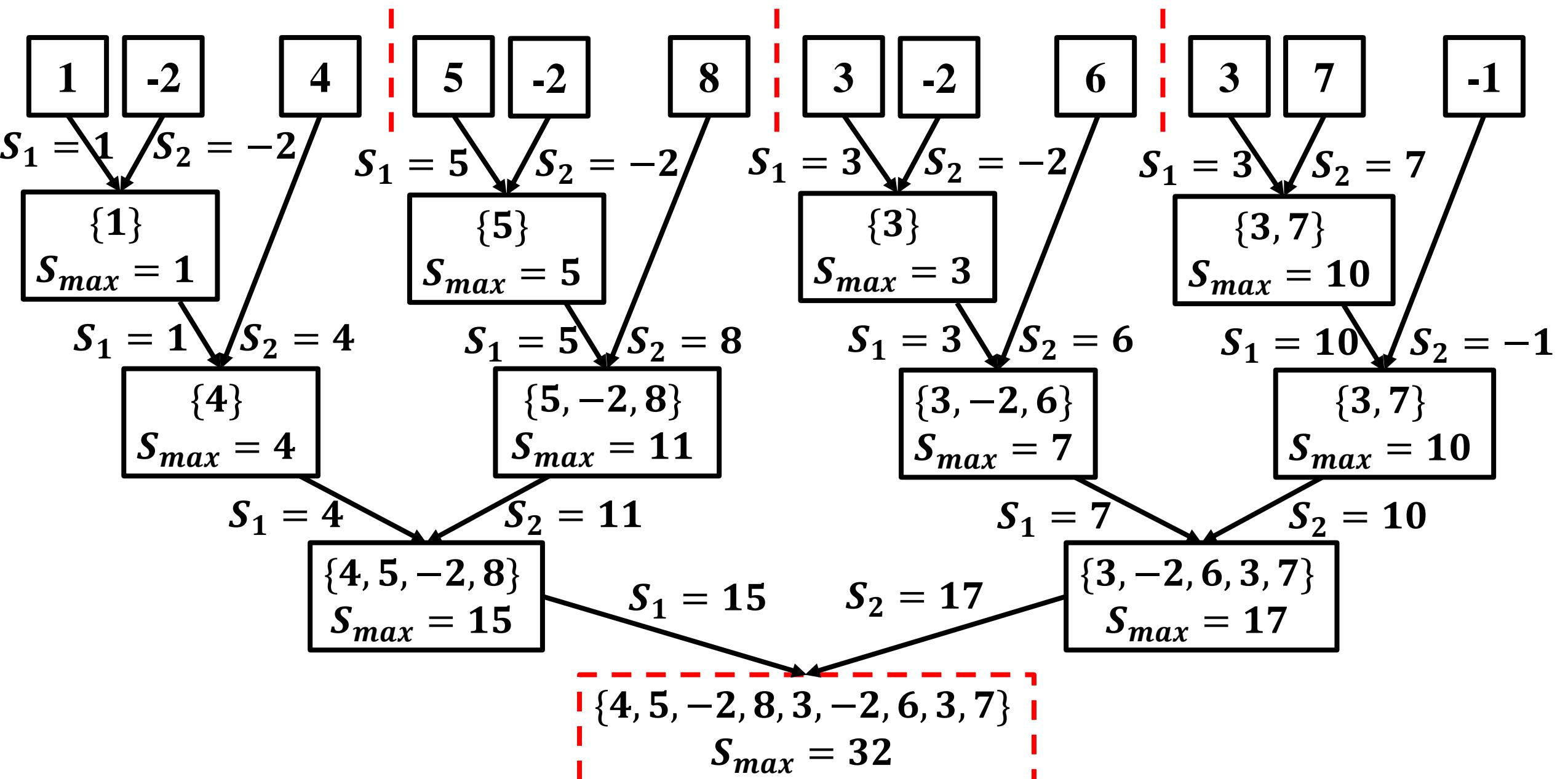
算法实例



算法实例



算法实例



伪代码

- **MaxSubArray($X, low, high$)**

输入: 数组 X , 数组下标 $low, high$

输出: 最大子数组之和 S_{max}

```
(if  $low = high$  then
| | return  $X[low]$ 
| |
| end
else
     $mid \leftarrow \lfloor \frac{low+high}{2} \rfloor$ 
     $S_1 \leftarrow \text{MaxSubArray}(X, low, mid)$ 
     $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$ 
     $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$ 
     $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$ 
    return  $S_{max}$ 
end
```

递归终止条件

伪代码

- **MaxSubArray($X, low, high$)**

输入: 数组 X , 数组下标 $low, high$

输出: 最大子数组之和 S_{max}

if $low = high$ then

| return $X[low]$

end

else

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| $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$

| $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$

| $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$

| return S_{max}

end

拆分原问题

伪代码

- **MaxSubArray($X, low, high$)**

输入: 数组 X , 数组下标 $low, high$

输出: 最大子数组之和 S_{max}

```
if  $low = high$  then  
    | return  $X[low]$ 
```

```
end
```

```
else
```

```
    |  $mid \leftarrow \lfloor \frac{low+high}{2} \rfloor$ 
```

```
    |  $S_1 \leftarrow \text{MaxSubArray}(X, low, mid)$ 
```

```
    |  $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$ 
```

```
    |  $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$ 
```

```
    |  $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$ 
```

```
    | return  $S_{max}$ 
```

```
end
```

求解子问题

伪代码

- **MaxSubArray($X, low, high$)**

输入: 数组 X , 数组下标 $low, high$

输出: 最大子数组之和 S_{max}

if $low = high$ then

| return $X[low]$

end

else

| $mid \leftarrow \lfloor \frac{low+high}{2} \rfloor$

| $S_1 \leftarrow \text{MaxSubArray}(X, low, mid)$

| $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$

| $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$

| $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$

| return S_{max}

end

合并问题解

伪代码

- **MaxSubArray($X, low, high$)**

初始调用： MaxSubArray($X, 1, n$)

输入：数组 X , 数组下标 $low, high$

输出：最大子数组之和 S_{max}

if $low = high$ then

| return $X[low]$

end

else

| $mid \leftarrow \lfloor \frac{low+high}{2} \rfloor$

| $S_1 \leftarrow \text{MaxSubArray}(X, low, mid)$

| $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$

| $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$

| $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$

| return S_{max}

end

时间复杂度分析

- 输入规模为: n

- $$T(n) = \begin{cases} 1, & n = 1 \\ 2 \cdot T(n/2) + O(n), & n > 1 \end{cases}$$

输入: 数组 X , 数组下标 $low, high$

输出: 最大子数组之和 S_{max}

if $low = high$ then

| return $X[low]$

end

else

| $mid \leftarrow \lfloor \frac{low+high}{2} \rfloor$

| $S_1 \leftarrow \text{MaxSubArray}(X, low, mid)$

| $S_2 \leftarrow \text{MaxSubArray}(X, mid+1, high)$

| $S_3 \leftarrow \text{CrossingSubArray}(X, low, mid, high)$

| $S_{max} \leftarrow \max\{S_1, S_2, S_3\}$

| return S_{max}

end

}

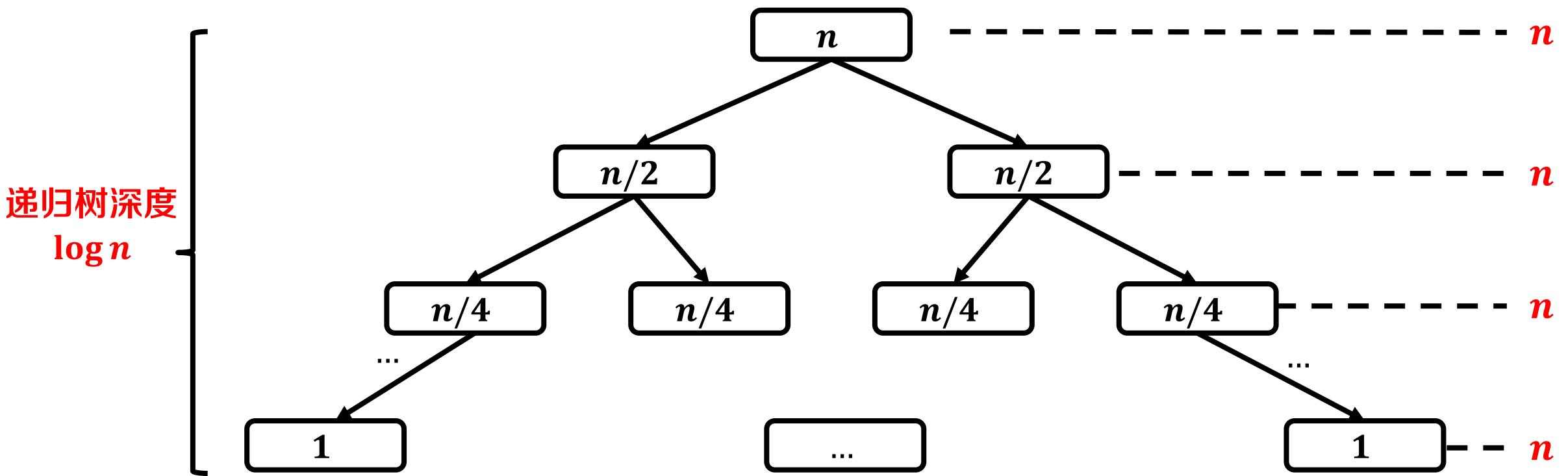
$2 \cdot T(n/2)$

$O(n)$

时间复杂度分析

- 输入规模为： n

- $T(n) = \begin{cases} 1, & n = 1 \\ 2 \cdot T(n/2) + O(n), & n > 1 \end{cases}$



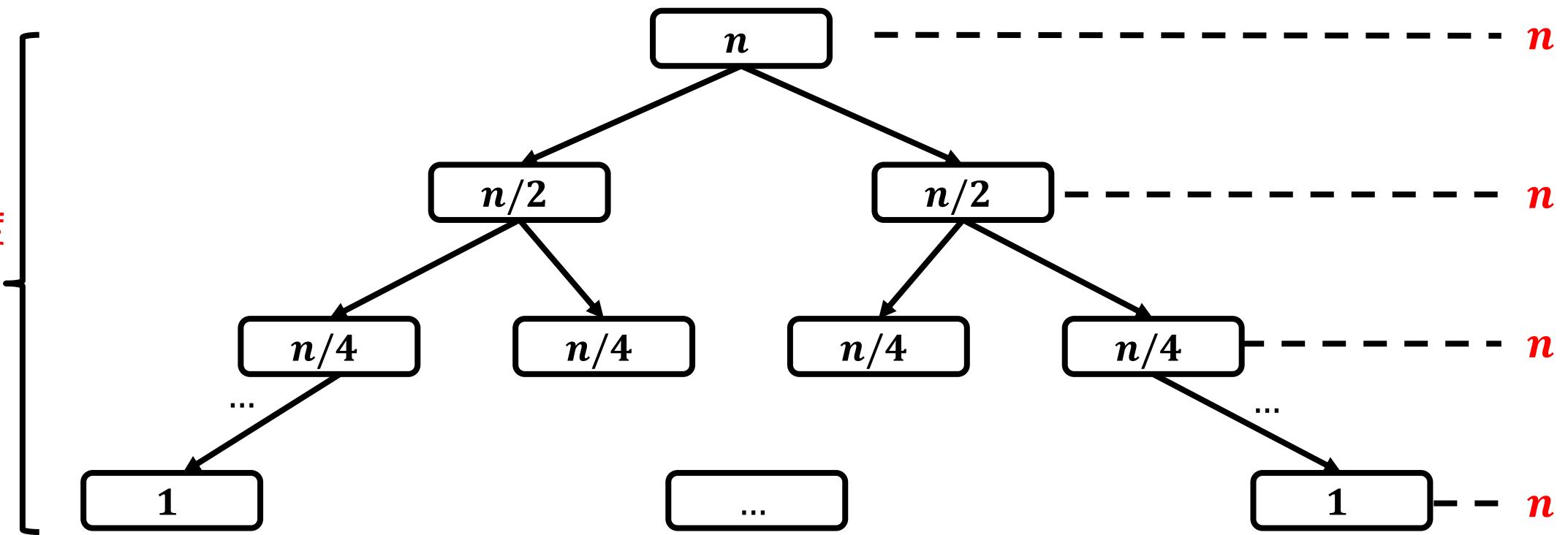
时间复杂度分析

- 输入规模为： n

- $T(n) = \begin{cases} 1, & n = 1 \\ 2 \cdot T(n/2) + O(n), & n > 1 \end{cases}$

$$T(n) = n \log n$$

递归树深度
 $\log n$



小结

- 本节讲述了三种不同的算法，用于解决最大子数组和问题：

算法名称	时间复杂度
蛮力枚举	$O(n^3)$
优化枚举	$O(n^2)$
分而治之	$O(n \log n)$

小结

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?	$O(n)$

问题：是否可以设计一个时间复杂度为 $O(n)$ 的算法？

小结

- 本节讲述了三种不同的算法，用于解决最大子数组和问题：

算法名称	时间复杂度
蛮力枚举	$O(n^3)$
优化枚举	$O(n^2)$
分而治之	$O(n \log n)$
动态规划	$O(n)$

问题：是否可以设计一个时间复杂度为 $O(n)$ 的算法？

答案：动态规划！