

-CS 315-

Project

Report

Done by:

Dana AL-Duayji	421202227
Emtenan AL-fozan	421215099
Shrooq AL-Qaied	4212202213

Supervised by Ms. Norah Alosily

Table of Contents

Merge Sort	3
Implementation.....	3
Execution Time.....	7
Result Table.....	8
Result Chart.....	8
Analyzation.....	9
 BST.....	10
Implementation.....	10
Execution Time.....	12
Result Table.....	13
Result Chart.....	13
Analyzation.....	14
 Hash Table	15
Implementation.....	15
Result Table.....	18
Result Chart.....	18
Analyzation.....	19

MERGE SORT ALGORITHM

by Dana AL-Duayji

-Implementation

```
✓ [351] import timeit  
0s     import numpy  
         import pandas as pd  
         import plotly.express as px
```

```
✓ ⏴ def merge(arr, l, m, r):  
 0s      n1 = m - l + 1  
      n2 = r - m  
      L = [0] * (n1)  
      R = [0] * (n2)  
      for i in range(0, n1):  
          L[i] = arr[l + i]  
  
      for j in range(0, n2):  
          R[j] = arr[m + 1 + j]  
      i = 0  
      j = 0  
      k = l  
      while i < n1 and j < n2:  
          if L[i] <= R[j]:  
              arr[k] = L[i]  
              i += 1  
          else:  
              arr[k] = R[j]  
              j += 1  
          k += 1  
      while i < n1:  
          arr[k] = L[i]  
          i += 1  
          k += 1  
      while j < n2:  
          arr[k] = R[j]  
          j += 1  
          k += 1  
  def mergeSort(arr, l, r):  
      if l < r:  
          m = l + (r - 1) // 2  
  
          mergeSort(arr, l, m)  
          mergeSort(arr, m + 1, r)  
          merge(arr, l, m, r)
```

MERGE SORT

ALGORITHM

creating an array of random numbers and call merge sort

```
✓ [353] def randomarray(siz):
    numarray = numpy.random.randint(1,100,siz)
    n = len(numarray)
    print(numarray)
    mergeSort(numarray,0,n-1)
    print(numarray)
```

creating an array of ascending numbers and call merge sort

```
✓ [354] def ascendingarray(siz):
    numarray = numpy.random.randint(1,100,siz)
    n = len(numarray)
    sortednumarray=sorted(numarray)
    print(sortednumarray)
    mergeSort(sortednumarray,0,n-1)
    print(sortednumarray)
```

creating an array of ascending numbers and call merge sort

```
✓ [355] def descendingarray(siz):
    numarray = numpy.random.randint(1,100,siz)
    n = len(numarray)
    desnumarray = sorted(numarray , reverse = True)
    print(desnumarray)
    mergeSort(desnumarray,0,n-1)
    print(desnumarray)
```

array of 100

```
✓ [356] st1 = timeit.default_timer()
randomarray(100)
end1= (timeit.default_timer()-st1)*1000
print("time taken for sorting is " , (end1) , "milliseconds " )

st2 = timeit.default_timer()
ascendingarray(100)
end2= (timeit.default_timer()-st2)*1000
print("time taken for sorting is " , (end2) , "milliseconds " )

st3 = timeit.default_timer()
descendingarray(100)
end3= (timeit.default_timer()-st3)*1000
print("time taken for sorting is " , (end3) , "milliseconds " )
```

MERGE SORT

ALGORITHM

array of 1000

```
✓ [357] 0s stt1 = timeit.default_timer()
randomarray(1000)
endd1= (timeit.default_timer()-stt1)*1000
print("time taken for sorting is " , (endd1) , "milliseconds " )

stt2 = timeit.default_timer()
ascendingarray(1000)
endd2= (timeit.default_timer()-stt2)*1000
print("time taken for sorting is " , (endd2) , "milliseconds " )

stt3 = timeit.default_timer()
descendingarray(1000)
endd3= (timeit.default_timer()-stt3)*1000
print("time taken for sorting is " , (endd3) , "milliseconds " )
```

array of 10000

```
✓ 0s ➜ sttt1 = timeit.default_timer()
randomarray(10000)
enddd1= (timeit.default_timer()-sttt1)*1000
print("time taken for sorting is " , (enddd1) , "milliseconds " )

sttt2 = timeit.default_timer()
ascendingarray(10000)
enddd2= (timeit.default_timer()-sttt2)*1000
print("time taken for sorting is " , (enddd2) , "milliseconds " )

sttt3 = timeit.default_timer()
descendingarray(10000)
enddd3= (timeit.default_timer()-sttt3)*1000
print("time taken for sorting is " , (enddd3) , "milliseconds " )
```

result table

```
✓ 0s  data = {'n': [100, 100, 100, 1000, 1000, 10000, 10000, 10000],  
           'case': ['random', 'ascending', 'descending', 'random', 'ascending', 'descending', 'random', 'ascending', 'descending'],  
           'time': [end1, end2, end3, endd1, endd2, endd3, enddd3, enddd2, enddd3]}  
  
df = pd.DataFrame(data)  
  
df
```

result chart

```
[ ] fig = px.line(data, x='n', y='time', color='case')  
fig.update_layout(xaxis_type='category')  
fig.update_traces(mode='markers+lines')  
fig.show()
```

Source Code: [Click Here](#)

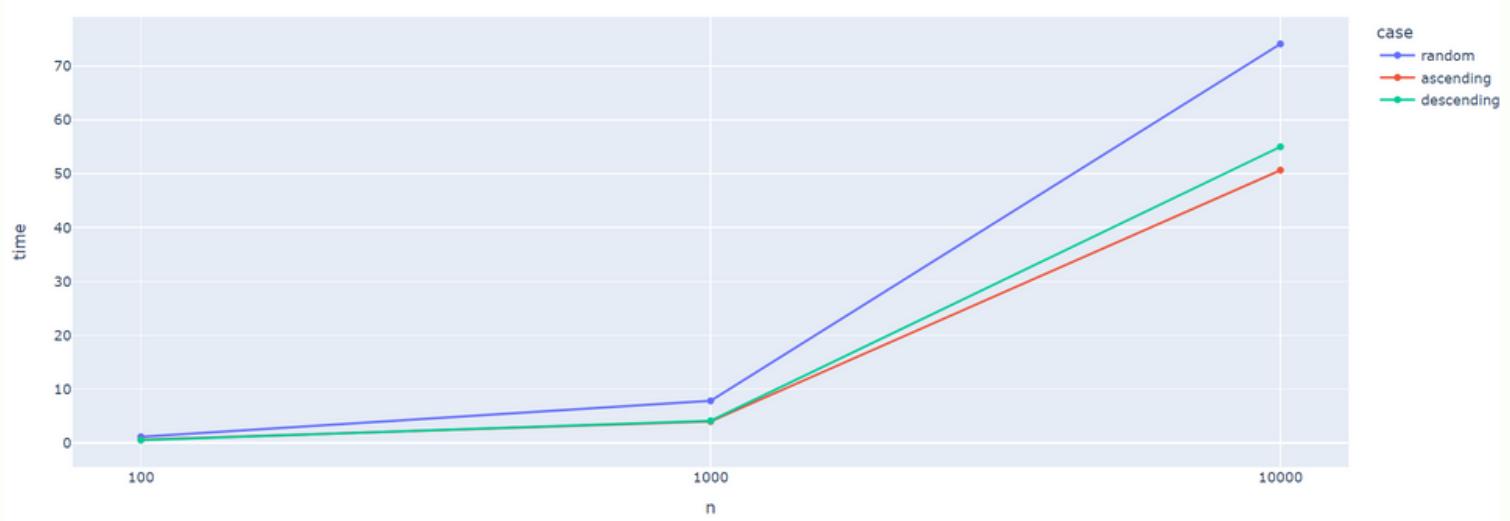
-Execution Time

SIZE=100	time taken for sorting random array is 1.1328909999974712 milliseconds time taken for sorting ascending array is 0.5772740000793419 milliseconds time taken for sorting descending is 0.5054270000073302 milliseconds
SIZE=1000	time taken for sorting random array is 7.800822999797674 milliseconds time taken for sorting ascending array is 3.968220999922778 milliseconds time taken for sorting descarray is 4.098275000160356 milliseconds
SIZE=10000	time taken for sorting random array is 74.12565999993603 milliseconds time taken for sorting ascending array is 50.66710500022964 milliseconds time taken for sorting descending array is 55.013711000356125 milliseconds

-Result Table

	n	case	time
0	100	random	1.132891
1	100	ascending	0.577274
2	100	descending	0.505427
3	1000	random	7.800823
4	1000	ascending	3.968221
5	1000	descending	4.098275
6	10000	random	74.125660
7	10000	ascending	50.667105
8	10000	descending	55.013711

-Result Chart



-Analyzation

To conclude, merge sort is a comparison-based sorting algorithm that divides an input array into smaller sub-arrays, compares them and then sorts those sub-arrays, and then merges them back together to produce a sorted array.

The time complexity of merge sort is $O(n \log n)$ in all cases, making it efficient for sorting large datasets. This $O(n \log n)$ complexity is the same for best, average, and worst-case scenarios. Here's why:

- Merge sort follows a divide-and-conquer strategy, where the array is repeatedly divided into two halves until you have sub-arrays of size 1. The time complexity for this division step is $O(\log n)$, where "n" is the size of the initial array. Then, it merges these sub-arrays in a sorted manner. So, regardless of the initial order of the array, this process will continue with the same time complexity.
- It compares and combines two sorted sub-arrays in a way that ensures the resulting merged array is sorted. This merge operation takes $O(n)$ time for any input, which is independent of the order of the elements.

BST ALGORITHM

-Implementation

```
[199] import time
import random
import numpy as np
import matplotlib.pyplot as plt
import pandas as pd
import plotly.express as px

[200] class Node:
    def __init__(self, key):
        self.left = None
        self.right = None
        self.val = key

Definition of the BinarySearchTree class, operation

[201] class BinarySearchTree:
    def __init__(self):
        self.root = None

    def insert(self, key):
        self.root = self._insert(self.root, key)

    def _insert(self, node, key):
        if node is None:
            return Node(key)

        if key < node.val:
            node.left = self._insert(node.left, key)
        else:
            node.right = self._insert(node.right, key)

        return node

    def search(self, key):
        return self._search(self.root, key)

    def _search(self, node, key):
        if node is None or node.val == key:
            return node
        if key < node.val:
            return self._search(node.left, key)
        return self._search(node.right, key)

    def inorder_traversal(self):
        if self.root:
            self._inorder_traversal(self.root)

    def _inorder_traversal(self, root):
        if root:
            self._inorder_traversal(root.left)
            print(root.val, end=' ')
            self._inorder_traversal(root.right)

    def postorder_traversal(self):
        if self.root:
            self._postorder_traversal(self.root)

    def _postorder_traversal(self, root):
        if root:
            self._postorder_traversal(root.left)
            self._postorder_traversal(root.right)
            print(root.val, end=' ')

    def delete(self, key):
        if self.root is None:
            return self.root
        if key < self.root.val:
            self.root.left = self.delete(self.root.left, key)
        elif key > self.root.val:
            self.root.right = self.delete(self.root.right, key)
        else:
            if self.root.left is None:
                return self.root.right
            elif self.root.right is None:
                return self.root.left
            self.root.val = self.find_min(self.root.right)
            self.root.right = self.delete(self.root.right, self.root.val)
        return self.root

    def find_min(self, root):
        while root.left:
            root = root.left
        return root.val

Function to measure the runtime of BST search

[202] def search_runtime(bst, number):
    start = time.time()
    bst.search(number)
    end = time.time()
    return (end - start) * 1000

Function to convert a sorted array to a balanced BST

[203] def sortedarraytobst(arr):
    if len(arr) == 0:
        return None

    mid = len(arr) // 2
    root = Node(arr[mid])
    root.left = sortedarraytobst(arr[:mid])
    root.right = sortedarraytobst(arr[mid+1:])
    return root

Function to insert numbers to a BST

[204] def add_numbers_to_bst(numbers, bst):
    for num in numbers:
        bst.insert(num)

Function to calculate average runtime over multiple runs

[205] def calc_avg_runtime(case_runtime):
    runs = 50
    total_time = 0
    for i in range(runs):
        total_time += case_runtime[i]
    return total_time / runs

Generate random data for each size

[206] size1 = 50
size2 = 150
size3 = 500

data1 = np.random.randint(1, 1000, size=size1)
data2 = np.random.randint(1, 1000, size=size2)
data3 = np.random.randint(1, 1000, size=size3)

Creating a best-case BST where the tree is balanced

[207] print(data1)
#print(data2)
#print(data3)

bst_balanced = BinarySearchTree()
start = time.time()
bst_balanced.root = sortedarraytobst(sorted(data1))
end = time.time()
best_case_50v=end - start)*1000
```

Connected to Python 3 Google Compute Engine backend

BST

ALGORITHM

```
✓ [285] bst_balanced3 = BinarySearchTree()
start = time.time()
bst_balanced3.root = sortedArrayToBST(sorted(data3))
end = time.time()
best_case_500=(end - start) * 1000
print("Elements in-order for bst_balanced:")
bst_balanced.inorder_traversal()
print()
print("Elements post-order for bst_balanced:")
bst_balanced.postorder_traversal()
print()

[785 78 426 501 944 639 15 475 374 460 899 876 458 878 741 112 2 208
961 739 266 296 346 388 227 253 206 90 330 920 677 911 581 169 348 991
188 204 190 348 909 672 378 553 390 348 511 568 899 442]
Elements in-order for bst_balanced:
2 15 78 90 112 169 188 190 204 206 208 237 253 266 296 330 346 348 348 374 378 388 390 426 442 458 460 475 501 511 553 568 581 639 672 677 739 741 785 876 878 899 909 920 931 944 963 991
Elements post-order for bst_balanced:
2 78 15 112 169 96 190 204 208 237 206 188 266 330 296 348 348 364 374 378 390 426 388 348 253 458 475 460 511 553 501 581 639 677 739 672 568 785 876 899 878 920 931 963 991 944 909 741 442
```

Creating an average-case BST where the tree has random input

```
✓ [208] bst_avg1 = BinarySearchTree()
start = time.time()
add_numbers_to_bst(data1, bst_avg1)
end = time.time()
Average_case50 =(end - start) * 1000

bst_avg2 = BinarySearchTree()
start = time.time()
add_numbers_to_bst(data2, bst_avg2)
end = time.time()
Average_case150 =(end - start) * 1000

bst_avg3 = BinarySearchTree()
start = time.time()
add_numbers_to_bst(data3, bst_avg3)
end = time.time()
Average_case500 =(end - start) * 1000
```

Creating a worst-case BST where the tree is unbalanced

```
✓ [209] bst_worst1 = BinarySearchTree()
data1.sort()
start = time.time()
add_numbers_to_bst(data1, bst_worst1)
end = time.time()
```

```
239] Worst_case50 =(end - start) * 1000
bst_worst2 = BinarySearchTree()
data2.sort()
start = time.time()
add_numbers_to_bst(data2, bst_worst2)
end = time.time()
Worst_case150 =(end - start) * 1000

bst_worst3 = BinarySearchTree()
data3.sort()
start = time.time()
add_numbers_to_bst(data3, bst_worst3)
end = time.time()
Worst_case500 =(end - start) * 1000
```

Calculating average runtimes for best, average, and worst cases

```
240] best_case_avg_runtime50 = calc_avg_runtime(lambda: best_case_50)
best_case_avg_runtime150 = calc_avg_runtime(lambda: best_case_150)
best_case_avg_runtime500 = calc_avg_runtime(lambda: best_case_500)

average_case_avg_runtime50 = calc_avg_runtime(lambda: Average_case50)
average_case_avg_runtime150 = calc_avg_runtime(lambda: Average_case150)
average_case_avg_runtime500 = calc_avg_runtime(lambda: Average_case500)

worst_case_avg_runtime50 = calc_avg_runtime(lambda: Worst_case50)
worst_case_avg_runtime150 = calc_avg_runtime(lambda: Worst_case150)
worst_case_avg_runtime500 = calc_avg_runtime(lambda: Worst_case500)
```

240]

Printing average runtimes for different cases

```
0 print('Best Case average Runtime for size=50:', best_case_avg_runtime50)
print('Average Case average Runtime for size=50 : ', average_case_avg_runtime50)
print('Worst Case average Runtime for size=50:', worst_case_avg_runtime50)

print('Best Case average Runtime for size=150:', best_case_avg_runtime150)
print('Average Case average Runtime for size=150:', average_case_avg_runtime150)
print('Worst Case average Runtime for size=150:', worst_case_avg_runtime150)

print('Best Case average Runtime for size=500:', best_case_avg_runtime500)
print('Average Case average Runtime for size=500:', average_case_avg_runtime500)
print('Worst Case average Runtime for size=500:', worst_case_avg_runtime500)

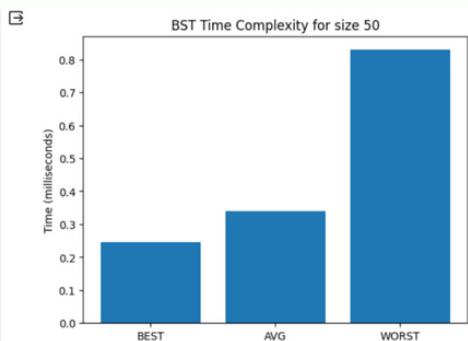
Best Case average Runtime for size=50: 0.141143798828125
Average Case average Runtime for size=50 : 0.211477272964308594
Worst Case average Runtime for size=50: 0.37169456481931594
Best Case average Runtime for size=150: 0.484596328732353156
Average Case average Runtime for size=150: 0.4661083221435547
Worst Case average Runtime for size=150: 4.796508974365234
Best Case average Runtime for size=500: 0.8156299591864453
Average Case average Runtime for size=500: 1.701115618996484
Worst Case average Runtime for size=500: 33.66923322143555
```

Source Code: [Click here](#)

-Execution Time

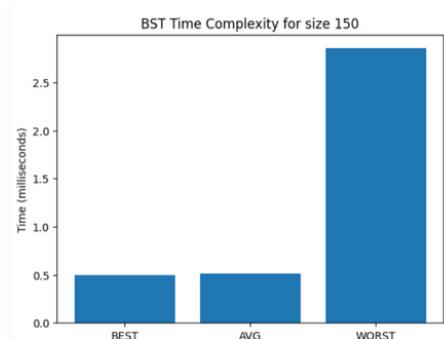
SIZE=50

Best Case average Runtime for size 50: 0.24461746215820312
 Average Case average Runtime for size 50: 0.3409385681152344
 Worst Case average Runtime for size 50: 0.8289813995361328



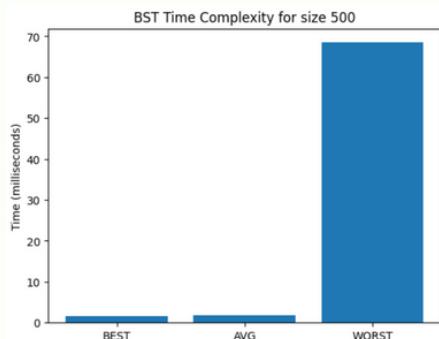
SIZE=150

Best Case average Runtime for size 150: 0.4947185516357422
 Average Case average Runtime for size 150: 0.5099773406982422
 Worst Case average Runtime for size 150: 2.854585647583008



SIZE=500

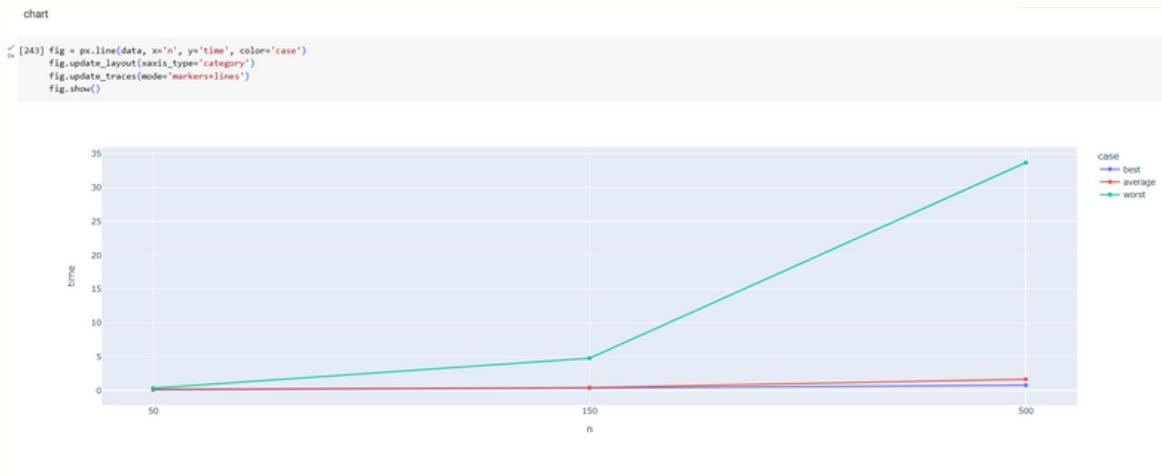
Best Case average Runtime for size 500: 1.6121864318847656
 Average Case average Runtime for size 500: 1.814126968383789
 Worst Case average Runtime for size 500: 68.47238540649414



-Result Table

n	case	time
0	50 best	0.141144
1	50 average	0.211477
2	50 worst	0.371695
3	150 best	0.404595
4	150 average	0.466108
5	150 worst	4.796505
6	500 best	0.815630
7	500 average	1.701117
8	500 worst	33.669233

-Result Chart



-Analyzation

In conclusion, the analysis of the Binary Search Tree (BST) algorithm provides valuable insights into its performance across different scenarios.

The algorithm illustrates efficient and reliable behavior in **best-case scenarios**, where the tree is balanced, with a complexity of $O(\log n)$.

The **average-case scenario** observed with random input data, also maintains a logarithmic time complexity of $O(\log n)$.

The **worst-case scenario** happens if the tree becomes unbalanced(sorted data) with a complexity $O(n)$.

I measured the average run time of each case by different data sizes [50,150,500], and I concluded that the best, average cases are quietly close to each other in their average runtime.

I have provided many operations of BST and I have tested the inorder, post-order traversal.

Overall, this analysis serves as a valuable foundation for understanding the practical importance and trade-offs associated with the Binary Search Tree algorithm.

HASH TABLE

By shrooq AL-Qaied

-Implementation

```
[ ] import time
import numpy as np
import pandas as pd
import matplotlib.pyplot as plt

[ ] class HashTable:
    def __init__(self, size):
        self.size = size
        self.table = [None] * size

    def _hash(self, key):
        return hash(key) % self.size

    def insert(self, key, value):
        index = self._hash(key)
        if self.table[index] is None:
            self.table[index] = []
        for item in self.table[index]:
            if item[0] == key:
                item[1] = value
                return
        self.table[index].append([key, value])

    def get(self, key):
        index = self._hash(key)
        if self.table[index] is not None:
            for item in self.table[index]:
                if item[0] == key:
                    return item[1]
        raise KeyError(f"Key '{key}' not found")
```

HASH TABLE

ALGORITHM

```
[ ] def insert_best_case(hash_table, size):
    for i in range(size):
        key = i * 2
        value = f"value {i}"
        hash_table.insert(key, value)
```

```
[ ] def insert_average_case(hash_table, size):
    for i in range(size):
        key = np.random.randint(0, size)
        value = f"value {i}"
        hash_table.insert(key, value)
```

```
[ ] def insert_worst_case(hash_table, size):
    for i in range(size):
        key = i * size
        value = f"value {i}"
        hash_table.insert(key, value)
```

```
[ ] def measure_runtime(function, *args, num_iterations=100):

    total_time = 0

    for _ in range(num_iterations):
        start_time = time.time()
        function(*args)
        end_time = time.time()
        elapsed_time = (end_time - start_time) * 1000
        total_time += elapsed_time

    average_time = total_time / num_iterations
    return average_time
```

HASH TABLE

ALGORITHM

```
[ ] if __name__ == "__main__":
    sizes = [1000, 5000, 10000]
    data = []

    for size in sizes:
        best_case_ht = HashTable(size)
        average_case_ht = HashTable(size)
        worst_case_ht = HashTable(size)

        average_time_best_case = measure_runtime(insert_best_case, best_case_ht, size)
        average_time_average_case = measure_runtime(insert_average_case, average_case_ht, size)
        average_time_worst_case = measure_runtime(insert_worst_case, worst_case_ht, size)

        data.append([size, 'Best Case', average_time_best_case])
        data.append([size, 'Average Case', average_time_average_case])
        data.append([size, 'Worst Case', average_time_worst_case])

df = pd.DataFrame(data, columns=['Hash Table Size', 'Case', 'Average Search Time (seconds)'])
df
```

```
[ ] pivot_df = df.pivot(index='Hash Table Size', columns='Case', values='Average Search Time (seconds)')
pivot_df.plot(kind='bar', figsize=(10, 6))
plt.xlabel('Hash Table Size')
plt.ylabel('Average Search Time (seconds)')
plt.title('Average Search Time in Hash Table for Different Cases and Sizes')
plt.legend(title='Case')

plt.show()
```

Source Code: Click [here](#)

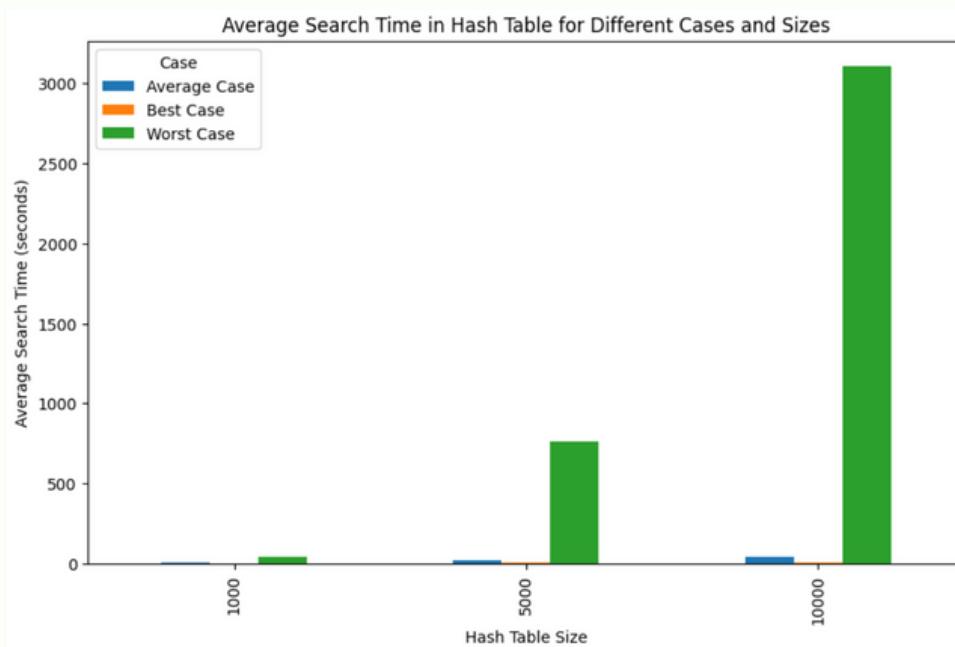
HASH TABLE

ALGORITHM

-RESULT TABLE

Hash Table Size	Case	Average Search Time (seconds)
0	1000 Best Case	0.773826
1	1000 Average Case	3.981521
2	1000 Worst Case	42.209091
3	5000 Best Case	3.878448
4	5000 Average Case	20.675890
5	5000 Worst Case	762.954707
6	10000 Best Case	7.715747
7	10000 Average Case	43.334467
8	10000 Worst Case	3108.881736

-RESULT CHART



-Analyzation

In conclusion,

The hash table is a data structure that allows for efficient key-value pair storage and recall.

It operates on the principle of using a hash function to map keys to specific locations in an array, providing fast access to values associated with those keys.

for the best, average and wort cases i can conclusion that:

BEST CASE:

- In the **best-case scenario**, where keys are well-distributed and there are no collisions OR there are equal number of collisions, which lead to highly efficient performance.
- The search and insertion operations have a time complexity of O(1).

AVERAGE CASE:

- In the **average-case scenario**, where the keys are randomly distributed and there are random numbers of collisions, which lead to a good performance.
- The search and insertion operations also have a time complexity of O(1), with some little differences due to the number of collisions.

WORST CASE:

- In the **worst-case scenario**, where all the keys are inserted into the same index, which lead to poor performance.
- The search and insertion operations have a time complexity of O(n), where "n" is the number of item in the collision chain.