P → {S} '.' ' S{S} ']' | '[' E '?' S{S} ']' Next Question ? $W \rightarrow '\{' E '?' S\{S\} '\}'$ Tell me all the tormols A → L '=' E ';' $E \rightarrow T \{('+' \mid '-') T\}$ When we write the parser that checks if the program is syntactically right or wrong, we will embed (or hardcode) $T \rightarrow U \{('*' \mid '/' \mid '\%') U\}$ the grammar rules in the parser's parse method. And we will implement a recursive parser that parses the program U → F '^' U | F in a top-down fashion. F → '(' E ')' | L | D For every non-terminal symbol we will have a function. For every terminal symbol in the grammar we will check if it is okay to see that and if so, consume the next token, if not stop with an error. **PSEUDO-CODE** Get the first token of the program. D → '0' | '1' | ... | If the current token is "." then this must be an empty program parse() therefore stop with a success. else S() S() if(current token is [) then currenttoken = scanner.nextToken() E() //function E will check if the next thing is an expression if(currenttoken is ?) //everything is okay currenttoken = scanner.nextToken() S() // we are inside S and we are calling S (recursion) (current les is I) / endoy if him

else if (curentables is:) litter une on in if run-else