Homework Set 2

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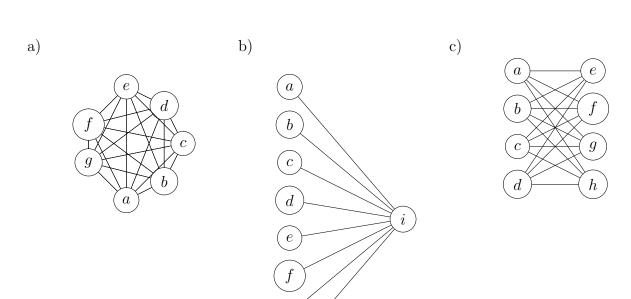
10 Graphs

10.1 Graphs and Graph Models

- 3. The graph has undirected edges and no loops, making it a simple graph.
- 4. The graph has multiple undirected edges and no loops, making it a multigraph.
- 5. The graph has multiple undirected edges and loops, making it a psuedograph.
- 6. The graph has multiple undirected edges and no loops, making it a multigraph.
- 7. The graph has directed edges and loops, making it a digraph.
- 8. The graph has multiple directed edges and loops, making it a directed multigraph.
- 9. The graph has multiple directed edges and loops, making it a directed multigraph.

10.2 Graph Terminology and Special Types of Graphs

- 1. |V|=6, |E|=6, $\deg a=2$, $\deg b=4$, $\deg c=1$ (pendant), $\deg d=0$ (isolated), $\deg e=2$, $\deg f=3$
- 2. |V| = 5, |E| = 13, $\deg a = 6$, $\deg b = 6$, $\deg c = 6$, $\deg d = 5$, $\deg e = 3$
- 3. |V| = 9, |E| = 12, $\deg a = 3$, $\deg b = 2$, $\deg c = 4$, $\deg d = 0$ (isolated), $\deg e = 6$, $\deg f = 0$ (isolated), $\deg g = 4$, $\deg h = 2$, $\deg i = 3$
- 5. A simple graph with 15 vertices each of degree 5 cannot exist, as all graphs must have an even number of vertices of odd degree.
- 7. |V| = 4, E = 7, $\deg^- a = 3$, $\deg^+ a = 1$, $\deg^- b = 1$, $\deg^+ b = 2$, $\deg^- c = 2$, $\deg^+ c = 1$, $\deg^- d = 1$, $\deg^+ d = 3$
- 8. |V| = 4, |E| = 8, $\deg^- a = 1$, $\deg^- b = 3$, $\deg^- c = 2$, $\deg^- d = 1$, $\deg^+ a = 2$, $\deg^+ b = 4$, $\deg^+ c = 1$, $\deg^+ d = 1$
- 9. |V| = 5, E = 13, $\deg^- a = 6$, $\deg^+ a = 1$, $\deg^- b = 1$, $\deg^- b = 5$, $\deg^- c = 2$, $\deg^+ c = 5$, $\deg^- d = 4$, $\deg^- d = 2$, $\deg^- e = 0$, $\deg^- e = 0$



d) e) f) h

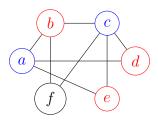


The graph is bipartite.

22.

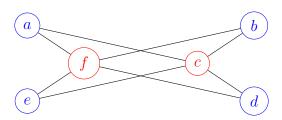
The graph is bipartite.

23.



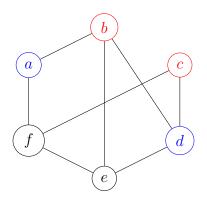
This graph is not bipartite, due to f.

24.



This graph is bipartite.

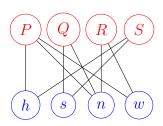
25.



This graph is not bipartite, due to e and f.

- 26. a) K_1 and K_2 are bipartite, but K_n for $n \ge 3$ is not bipartite, as any 3 vertices are connected pairwise, so there is no way to partition them into 2 disjoint sets.
- b) C_n is bipartite whenever n is even, as the vertices can simply alternate.
- c) W_n is never bipartite, as every vertex is connected to the center of the wheel.
- d) Q_n is always bipartite.

27. a)



37.

a)
$$|V| = n, |E| = \binom{n}{2}$$
 b) $|V| = n, |E| = n$ c) $|V| = n + 1, |E| = 2n$

b)
$$|V| = n, |E| = n$$

c)
$$|V| = n + 1, |E| = 2n$$

d)
$$|V| = m + n, |E| = mn$$
 e) $|V| = 2^n, |E| = n2^{n-1}$

53. why

Representing Graphs and Graph Isomorphism 10.3

1.

Vertex	Adjacent Vertices
a	b, c, d
b	a, d
c	a, d
d	a, b, c

3.

Vertex	Terminal Vertices
a	a, b, c, d
b	d
c	a, b
d	a, b b, c, d

5.

7.

$$\begin{bmatrix} a & b & c & d \\ 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ d & 0 & 1 & 1 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \end{bmatrix} \qquad \begin{bmatrix} 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

c)

$$\begin{bmatrix} 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

d) e)

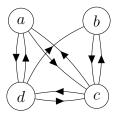
0	1	0	1
1	0	1	$\begin{bmatrix} 1 \\ 0 \\ 1 \\ 0 \end{bmatrix}$
0	1	0	1
1	0	1	0

 $\begin{bmatrix} 0 & 1 & 1 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \end{bmatrix}$

 $\begin{bmatrix} 0 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 0 \end{bmatrix}$

f)

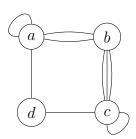
11.



13.

15.

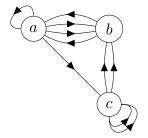
17.



19.

21.

23.



31.

$$\mathbf{M}_{13} = \begin{bmatrix} a \\ b \\ c \\ d \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \qquad \mathbf{M}_{14} = \begin{bmatrix} a \\ b \\ c \\ d \end{bmatrix} \begin{bmatrix} 1 & 1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}$$

- 33. For an undirected graph, the sum of the entries in a column of the adjacency matrix is the number of edges that are connected to that column's vertex (loops only being counted once). For a digraph, it is the in-degree of the vertex.
- 35. For an undirected graph, the sum of the values of a column in the incidence matrix is equal to the number of nodes that the column's edge is incident to. This can only be 1 (if the edge is a loop) or 2.

36. a) b)
$$\begin{bmatrix}
1 & 1 & \cdots & 1 \\
1 & 1 & \cdots & 1 \\
\vdots & \vdots & & \vdots \\
1 & 1 & \cdots & 1
\end{bmatrix}$$

$$\begin{bmatrix}
0 & 1 & 0 & \cdots & 0 & 1 \\
1 & 0 & 1 & \cdots & 0 & 0 \\
0 & 1 & 0 & \cdots & 0 & 0 \\
\vdots & \vdots & \vdots & & \vdots & \vdots \\
1 & 0 & 0 & \cdots & 1 & 0
\end{bmatrix}$$

c)
$$\begin{bmatrix} 1 & 1 & 1 & \cdots & 1 & 1 \\ 1 & 0 & 1 & 0 & \cdots & 0 & 1 \\ 1 & 1 & 0 & 1 & \cdots & 0 & 0 \\ 1 & 0 & 1 & 0 & \cdots & 0 & 0 \\ \vdots & \vdots & \vdots & & \vdots & \vdots & \vdots \\ 1 & 0 & 0 & \cdots & 1 & 0 \end{bmatrix}$$

37. a) b)
$$\begin{bmatrix}
1 & 1 & \cdots & 1 & 0 & \cdots & 0 \\
1 & 0 & \cdots & 0 & 1 & \cdots & 0 \\
0 & 1 & \cdots & 0 & 1 & \cdots & 0 \\
\vdots & \vdots & & \vdots & \vdots & & \vdots \\
0 & 0 & \cdots & 0 & 0 & \cdots & 1 \\
0 & 0 & \cdots & 1 & 0 & \cdots & 1
\end{bmatrix}$$
b)
$$\begin{bmatrix}
1 & 0 & \cdots & 0 & 1 \\
1 & 1 & \cdots & 0 & 0 \\
0 & 1 & \cdots & 0 & 0 \\
\vdots & \vdots & & \vdots & \vdots \\
0 & 0 & \cdots & 1 & 0 \\
0 & 0 & \cdots & 1 & 1
\end{bmatrix}$$

39.						
	v	u_1	u_2	u_3	u_4	u_5
	f(v)	v_1	v_3	v_2	v_5	v_3
-						

41.								
	v	u_1	u_2	u_3	u_4	u_5	u_6	u_7
	f(v)	v_1	v_3	v_5	v_7	v_2	v_4	v_6

- 45. u_5 is connected to exactly 2 other nodes, both of which are of degree 3. There is no node in the second graph that has this property. The graphs are therefore nonisomorphic.
- 47. u_1 u_3 u_4 u_{10} u_2 u_5 u_6 u_7 u_8 u_9 v_6 v_5 v_8 v_{10} v_7 v_3 v_4 v_1 v_9 v_2

63. a)

Yes

c) No, as there is no row in the first matrix with only 1 1.

67.

v	u_1	u_2	u_3	u_4
f(v)	v_3	v_4	v_2	v_1

b) No, as there is no row in the first matrix

with only 1 1.

69.

v	u_1	u_2	u_3	u_4
f(v)	v_3	v_4	v_2	v_1

10.5 Euler and Hamilton Paths

- 1. a, b and c are all of degree 3, which means that 3 nodes are of odd degree, so the graph has neither an Euler path or circuit.
- 3. a and d are of odd degree while every other is of even degree, so an Euler path may exist. Such a path is (a, e, c, e, b, e, d, b, a, c, d).
- 5. Every vertex is of even degree, so the graph may have an Euler circuit. Such a circuit is (a, e, a, e, c, d, c, b, e, d, b, a).
- 7. Every vertex is of even degree, so the graph may have an Euler circuit. Such a circuit is (a, ih, g, d, e, f, g, c, e, h, d, c, a, b, i, c, b, h, a).
- 13. This image can be drawn in a single stroke by treating the intersections of edges as nodes.
- 15. This image cannot be drawn in a single stroke.

19.

v	a	b	c	d
$\deg^- v$	1	3	1	2
$\deg^+ v$	2	1	2	2

 $\deg^- v$ is not always equal to $\deg^+ v$, so an Euler circuit may not exist. $\deg^+ a - \deg^- a = \deg^+ c - \deg^- c = 1$, so an Euler path may not exist either.

21. An Euler path exists: a, d, e, d, b, a, e, c, e, b, c, b, e.

23.

v		a	b	c	d	e	f	g	h	i	j	k	l
\deg^-	v	1	2	1	1	2	1	2	2	2	1	1	1
\deg^+	\dot{v}	1	1	1	2	2	2	1	2	1	1	2	1

 $\deg^- v$ is not always equal to $\deg^+ v$, so an Euler circuit may not exist. $\deg^+ d - \deg^- d = \deg^+ f - \deg^- f = 1$, so an Euler path may not exist either.

26. a)