Python → Pandas SQL mean sort df.groupby(column)

Advanced Databases

Databases II

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Queue

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index

date

2025-I

Data Engineer

trade-off

Outline

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What is an OODB?

(2000) → Java → 95%

- Combines object-oriented programming and database principles.
- Stores **objects** with their methods and attributes.
- Supports complex data and inheritance.

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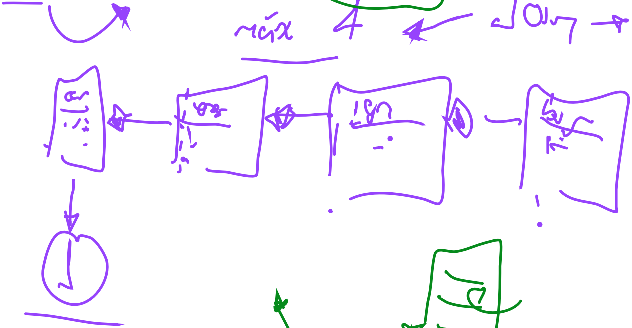
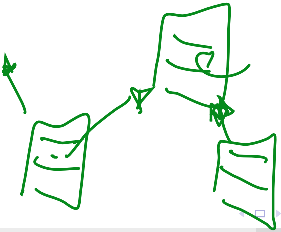
Design

máx 4 ← Java → 6-8-10





JSP/JSF

Java
Serve-
Pages

OODB Features

- Encapsulation, Inheritance, Polymorphism.
- Direct object storage and identity.
- Supports OQL (*Object Query Language*).

common
diff
diff

Serialization

Student
-grade

public isScholarship : bool {
return grade > 4.5 ? true : false;
}

public isApprove : bool {
return grade >= 3.0 ? true : false;

get set
grant

Validation
grade >= 0 && <= 5.0;

Persistence Mechanisms

- **By Reachability**: Root object persistence.
- **By Explicit Marking**: Annotate persistable objects.
- **By Modification**: Modified objects are persisted.

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Persistence Mechanisms

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```
@time
public int getAge(){
  :
}

@persist
public class Student(){
}
```

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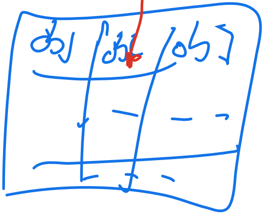

OODB Query Example (OQL)

OQL

```
SELECT s.name FROM Student s WHERE s.average() > 4.0
```

- **Queries** can follow object references.
- Better for nested, recursive, or complex types.

SELECT s.name FROM Student s

WHERE s.isScholarship();

(Legacy)

OODB Architecture

- Integrates:
  - **Object-Oriented Language Runtime**
  - **Persistence Engine**
  - **Query Processor (OQL)**
- Common deployment: embedded or client-server.

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Outline

Types of NoSQL Databases

- **Key-Value**: Redis, DynamoDB
- **Document**: MongoDB, CouchDB
- **Column-Family**: Cassandra, HBase
- **Graph**: Neo4j, ArangoDB

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CAP Theorem Deep Dive

Consistency (C)

Latest write is visible to all reads.

Availability (A)

System always responds.

Partition Tolerance (P)

Tolerates network splits.

Only two can be guaranteed at the same time.

CAP: Practical Examples

Database	CAP Preference	Comment
MongoDB	A + P	Eventual Consistency
Cassandra	A + P	Tunable consistency
Spanner	C + P	Sacrifices availability

NoSQL Architecture Patterns

- **Shared-nothing architecture** enables scalability.
- **Sharding**: Distributes data across partitions.
- **Replication**: Ensures fault-tolerance.
- **Routers** coordinate requests across shards (e.g., `mongos` in MongoDB).

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Case Study: MongoDB vs PostgreSQL

- **MongoDB**: JSON-like schema, scalable, flexible.
- **PostgreSQL**: ACID-compliant, strong relational support.
- **Trade-offs** depend on data model complexity vs transactional needs.

Outline

Parallel DB Concepts

- Uses multiple processors to speed up **query execution**.
- Exploits **parallelism** in data access and computation.
- **Reduces** query response time for big datasets.

Partitioning Strategies

- **Horizontal**: Distributes rows.
- **Vertical**: Splits columns.
- **Hash Partitioning**: Uniform distribution via hash.
- **Range Partitioning**: Based on value intervals.

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Parallel Query Cost Model

- **Startup cost (S)**: Fixed overhead.
- **Communication cost (C)**: Inter-node data transfer.
- **Computation cost (T)**: Local processing time.
- Total: $T_{total} = S + C + T$

Parallel DB Architectures

- **Shared Memory:** Easy communication, poor scalability.
- **Shared Disk:** Easier fault-tolerance, contention risk.
- **Shared Nothing:** Best scalability, harder coordination.

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Case Study: Greenplum

- Open-source **parallel DB** based on `PostgreSQL`.
- MPP architecture (massively parallel processing).
- Good for OLAP and analytics workloads.

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What is a Distributed DB?

- Data stored across multiple physical nodes.
- Appears as a single logical database.
- Must ensure consistency, availability, and partition tolerance.

Transparency Goals

- **Location**: Hide physical location of data.
- **Replication**: Hide duplication.
- **Fragmentation**: Hide data partitioning.

Distributed DBMS Architecture

- **Client-Server Model**: Clients query; servers store data.
- **Federated Model**: Semi-autonomous DBs collaborate.
- **Peer-to-Peer**: Nodes act as both client and server.
- **Layers**: Global schema, transaction manager, local engines.

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Two-Phase Commit Protocol (2PC)

1. **Prepare Phase**: Coordinator asks all participants to prepare.
2. **Commit Phase**: If all vote yes, coordinator commits.

Issue: Blocking if coordinator crashes during commit.

Paxos (Simplified Consensus)

- Needed for agreement in distributed systems.
- **Roles**:
  - *Proposer*: Suggests a value.
  - *Acceptor*: Votes.
  - *Learner*: Learns chosen value.
- Ensures consistency despite node failures.

Case Study: Google Spanner

- Globally distributed RDBMS.
- Uses *TrueTime API* (atomic clocks + GPS).
- Offers external consistency and global transactions.

Conclusion

- Data systems have evolved for **scalability** and **complexity**.
- Choosing the right DB model depends on workload needs.
- Understanding design trade-offs is key for architects.

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Thanks!

Questions?

Repo: `https://github.com/EngAndres/ud-public/tree/main/courses/databases-ii`