

DATA BASES DESIGN & RELATIONAL ALGEBRA

DataBase Foundations

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Outline

1 Basic Concepts

2 Normalization

3 Relational Algebra



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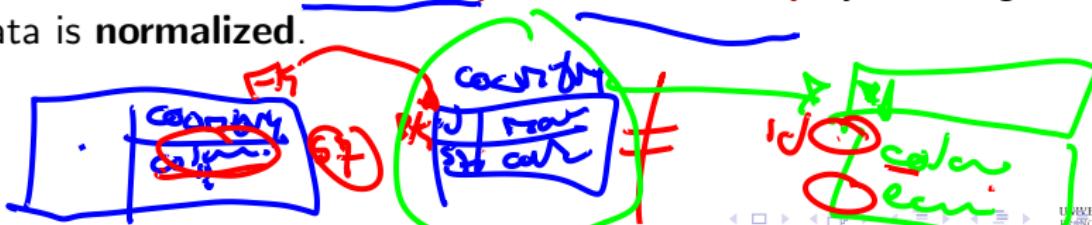


DataBases Design Foundations



- In the context of databases, the **design** of a database is the process of producing a **detailed** data model of a database.
- This **data model** contains all the needed **logical and physical design choices** and **physical storage parameters** needed to generate a design in a ***data definition language***, which can then be used to create a database.
- A **fully attributed data model** contains detailed attributes for **each entity**.
- Data models avoid **redundancy** and **inconsistency** by ensuring that data is **normalized**.

→ NoSQL

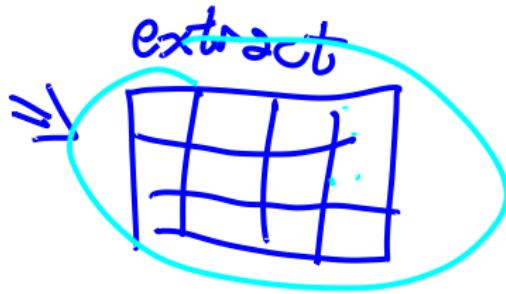
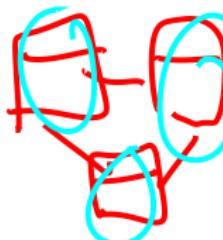


Set Theory in Databases

NoSQL \Rightarrow view SQL



- The **set theory** is a branch of mathematical logic that studies sets, which are **collections** of objects.
- The set theory is applied in databases to define the **relational model** and the **relational algebra**.
- The **relational model** is a **mathematical model** of data for large shared data banks and it has a **solid theoretical foundation**.
- The **relational algebra** is a procedural query language, which takes relations as input and produces relations as output.



*SQL
Deductive*



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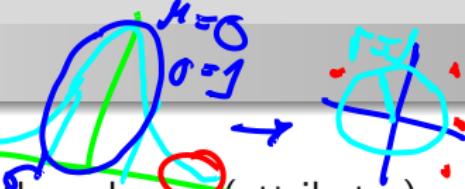
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Normalization in Databases

outliers

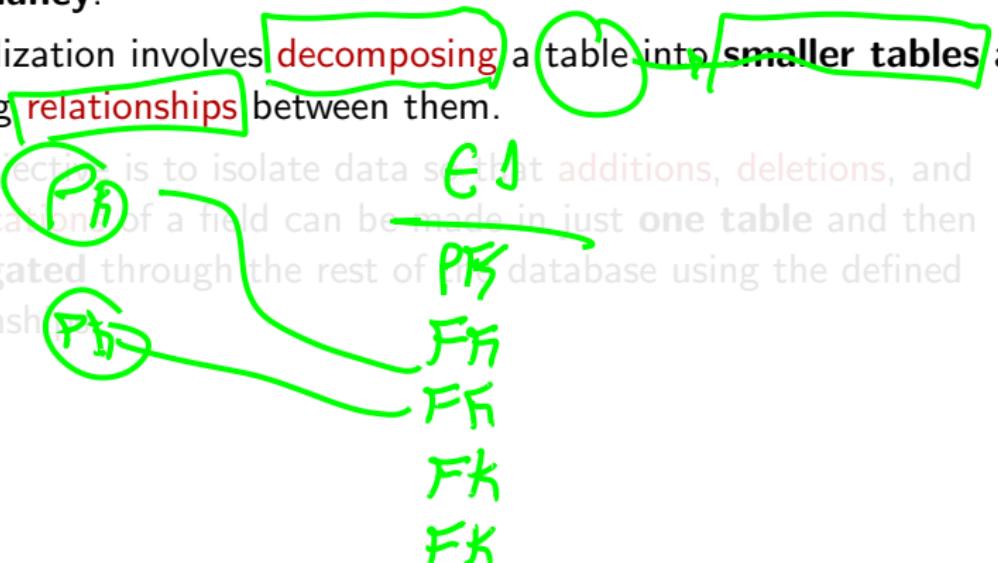


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- The objective is to isolate data so that additions, deletions, and modifications of a field can be made in just **one table** and then propagated through the rest of the database using the defined relationships.



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Ontologies

- An **ontology** is a formal naming and definition of the types, properties, and interrelationships of the **entities** that really or fundamentally exist for a particular domain of discourse.
- Ontologies are used in databases to **define** the **schema** of the database.
- The **schema** of a database is a formal definition of the **structure** of the database: the types of data that are stored, the relationships between the data, and the constraints on the data.



Normal Forms

- ① First normal form (1NF): The table is a two-dimensional table with rows and columns. Each column contains **atomic values**, and there are no repeating groups or arrays.
- ② Second normal form (2NF): The table is in first normal form and all the non-key attributes are fully functionally dependent on the primary key.
- ③ Third normal form (3NF): The table is in second normal form and all the non-key attributes are non-transitively dependent on the primary key.
- ④ Fourth normal form (4NF): The table is in third normal form and there are no multi-valued dependencies.



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What is relational algebra?

SQL

- The **relational algebra** is a **procedural query language**, which takes relations as input and produces relations as output.
- The relational algebra is a **set of operations** that can be performed on a **relation**. Also, it is used to define the **relational model**, which is a mathematical model of data for large shared data banks.
- Let's take a look at the **basic operations** of the relational algebra. First, remember next table called **Students**:

ID	Name	Lastname	Address	Phone	Age
1	John	Doe	123 Fake St	555-1234	25
2	Jane	Smith	456 Elm St	555-5678	30
3	Mike	Johnson	789 Evergreen St	555-9012	35



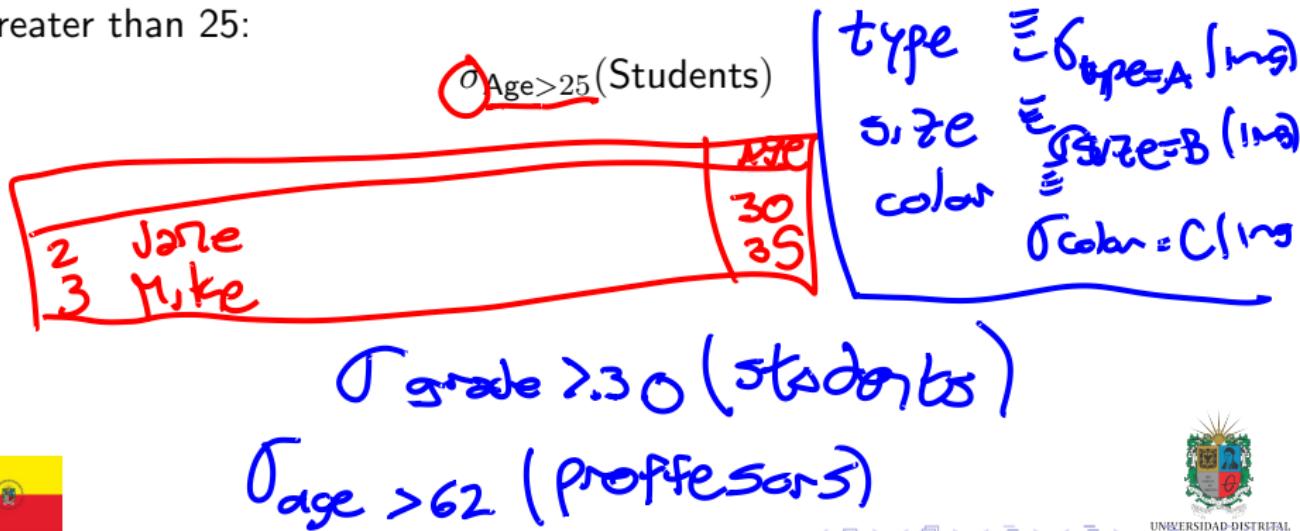
Select Operation

filter

Definition

Select: $\sigma_{\text{condition}}(R)$, is a unary operation that returns the rows (subset) of R that satisfy the condition.

For example, the following expression selects the students whose age is greater than 25:



Project Operation

Definition

Project: $\pi_{\text{column list}}(R)$, is a unary operation that returns the columns (subset) of R that are specified in the column list.

For example, the following expression projects the name and lastname of the students:

$\pi_{\text{Name, Lastname}}(\text{Students})$	
name	lastname
John	Doe
Jane	Smith
Mike	John son

$\pi_{\text{Name}}(\text{Students})$ → $\pi_{\text{Name}}(\text{Students} \rightarrow \text{Project})$

name
John

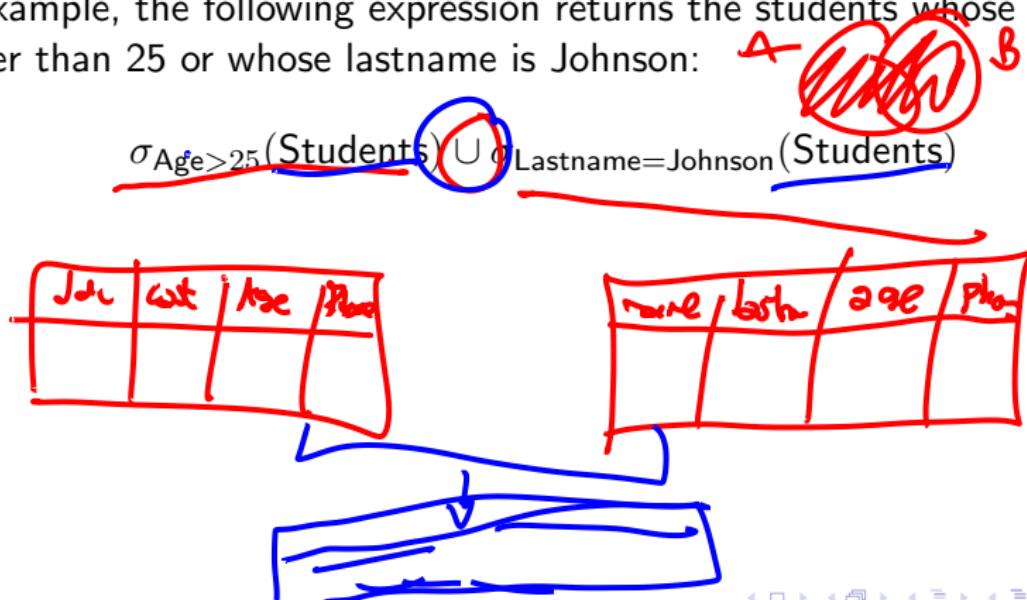


Union Operation

Definition

Union: $R \cup S$, is a binary operation that returns the rows that are in R or in S .

For example, the following expression returns the students whose age is greater than 25 or whose lastname is Johnson:

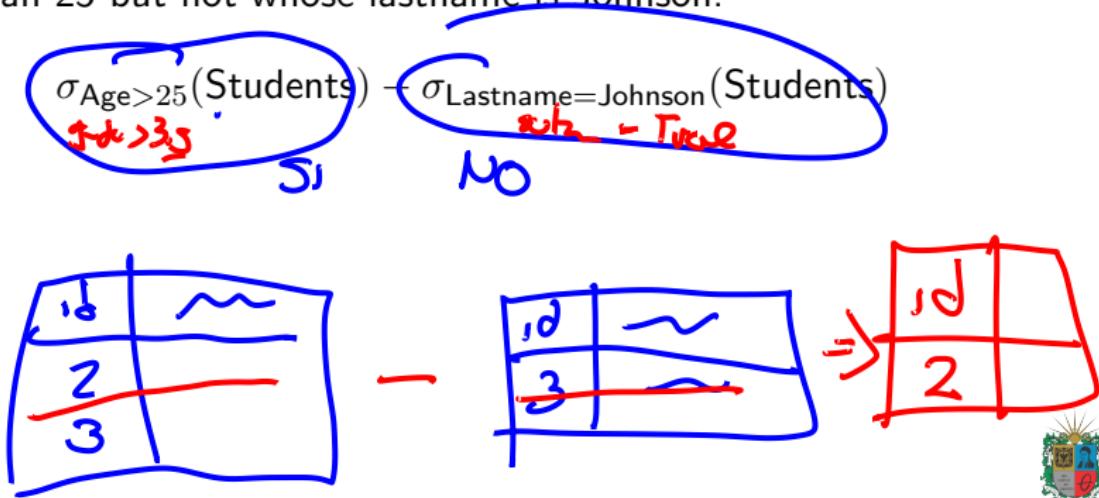


Set Different Operation

Definition

Set Different: $R - S$, is a binary operation that returns the rows that are in R but not in S .

For example, the following expression returns the students whose age is greater than 25 but not whose lastname is Johnson:



Cartesian Product Operation

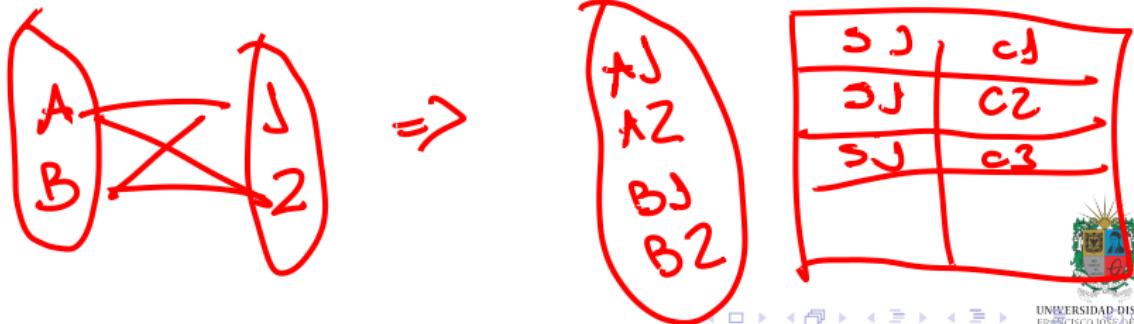
Definition

Cartesian Product: $R \times S$, is a binary operation that returns the Cartesian product of R and S . A formal definition is:

$$R \times S = \{r \cup s \mid r \in R \wedge s \in S\}$$

For example, the following expression returns the Cartesian product of the students and the courses:

Students \times Courses



Rename Operation

Definition

Rename: $\rho_{\text{new_name}}(R)$, is a unary operation that returns the relation R with the name R changed to new_name .

For example, the following expression returns the students relation with the name changed to **People**:

$\rho_{\text{People}}(\text{Students})$

$\rho_{\text{People}}(\sigma_{\text{grades} > 4.8}(\text{Students}))$



Exercises

~~$\sigma_{age > 25}(\text{Students}) \bowtie \sigma_{lastname = "Johnson"}(\text{Students})$~~

$\sigma_{age > 25 \& lastname = "Johnson"}(\text{Students})$

- ① Select the students whose age is greater than 25 and whose lastname is Johnson.
- ② Project the name and lastname of the students whose age is greater than 25. $\pi_{name, lastname}(\sigma_{age > 25}(\text{Students}))$
- ③ Select the students whose age is greater than 25 and whose lastname is Johnson, and project the name and lastname of the students, and rename the relation to People.

People ($\pi_{name, lastname}(\sigma_{age > 25 \& lastname = "Johnson"}(\text{Students}))$)



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Thanks!

Questions?



Repo:

 github.com/EngAndres/ud-public/tree/main/courses/databases-foundations

