



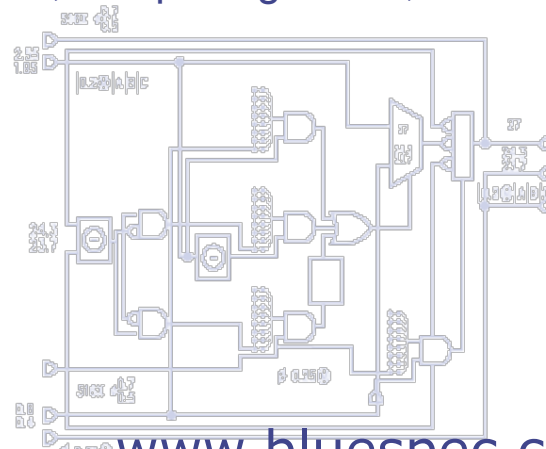
# BSV Training

## Lec\_StmtFSM

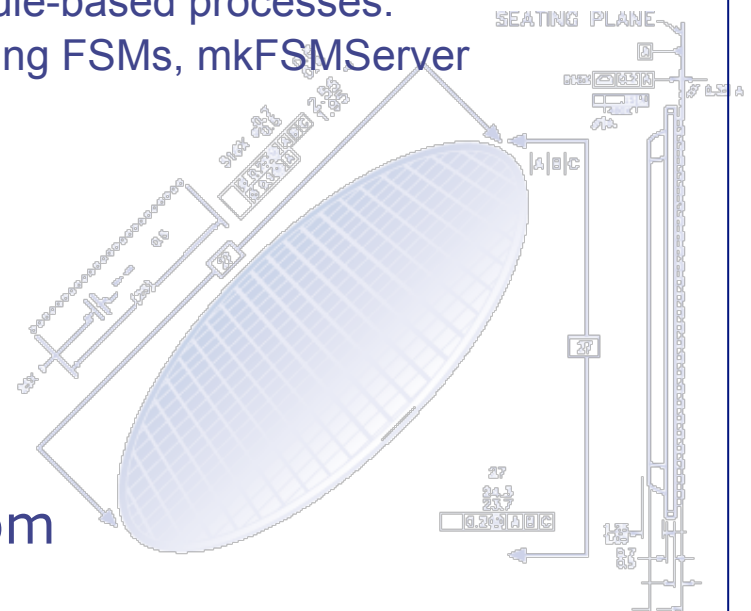
This is a facility for describing structured rule-based processes.  
 StmtFSM, mkFSM, composing FSMs, controlling FSMs, mkFSMServer

```

import PFCtrl;
typedef Bit[24] StateT;
module ex_hdl_ctrl_fsm(empty);
  Integer nfa_depth = 16;
  function Bit[24] determine_pump(StateT s);
    return (s[0]);
  endfunction
  PFCtrl(bstate) inboud0;
  mkFSM(PFCtrl(nfa_depth)) the_inboud0(inboud0);
  PFCtrl(bstate) outboud0;
  mkFSM(PFCtrl(nfa_depth)) the_outboud0(outboud0);
  PFCtrl(bstate) subboud0;
  mkFSM(PFCtrl(nfa_depth)) the_subboud0(subboud0);
  rule exp1 (True);
    bstate[0] = inboud0[0];
    PFCtrl(bstate) out_pump =
      determine_pump(bstate) == 0 ? outboud0 : subboud0;
  inboud0;
  outboud0;
  subboud0;
  control : exp1;
endmodule : ex_hdl_ctrl_fsm
  
```



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# Motivations

Finite State Machines (FSMs) are very common in hardware design.

With BSV rules, you can encode arbitrary FSMs.

For example, a simple FSM involving some sequencing and a loop:

```
typedef enum { S0, S1, S2, ... } State deriving (Bits, Eq);

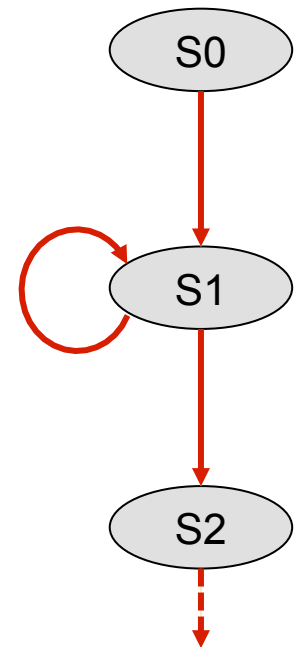
module mkFoo (...);
  Reg #(State) state <- mkReg (S0);

  rule r0 (state == S0);
    ... do state S0 actions ...
    state <= S1;                // next state
  endrule

  rule r1 (state == S1);
    ... do state S1 actions ...
    state <= (some cond ? S1 : S2); // loop back to S1 state, or exit loop
  endrule

  rule r2 (state == S2);
    ... do state S2 actions ...
    ... transition to next state ....
  endrule

endmodule
```



# Structured FSMs can be expressed more succinctly

There are common structured design patterns in FSMs:

- sequences, conditionals, loops, parallel threads, etc.

BSV provides a powerful FSM sub-language to express these more succinctly than having to write out the rules explicitly:

- But note that the semantics are *identical* to rules
- In fact, the compiler expands the FSM spec into the rules that you would have written by hand if you were to express them directly as rules

To use this facility:

- Import the “StmtFSM” package: 

```
import StmtFSM :: *;
```
- Create an FSM specification (an expression of type Stmt)
- Create an FSM module by giving it the specification
  - This returns an “FSM” interface with “start” and “done” methods
- Operate the FSM using the “start” and “done” methods

# The FSM interface

```
interface FSM;  
  method Action start;  
  method Bool   done;  
  method Action waitTillDone;  
  method Action abort;  
endinterface
```

*FSM interface*

```
module mkFSM #( Stmt s ) ( FSM );
```

*mkFSM module*

Note:

- ‘done’ and ‘waitTillDone’ are just alternative ways for knowing when the FSM is done. You can either test the boolean ‘fsm.done’, or you can execute the Action statement ‘fsm.waitTillDone’, whose implicit condition is the same as fsm.done.

In different situations, one of the other may be more convenient.

- ‘abort’ allows an external agent to force the FSM to reset to its initial state, no matter what state it is in, and no matter how deeply nested it is (more about nesting later)

# Example revisited, with FSMs instead of rules

```
import StmtFSM :: *;

module mkFoo (...);
  Stmt stmt = seq
    ... do state S0 actions ...
    ... do state S1 actions ...
    while (some cond) ... do state S1 actions ...
    ... do state S2 actions ...
  endseq;

  FSM fsm <- mkFSM (stmt);

  rule init (...);
    fsm.start; ...
  endrule

  rule done (fsm.done);
    ...
  endrule
endmodule
```

- Import the “StmtFSM” package:
- Create an FSM specification (an expression of type Stmt)
- Create an FSM module by giving it the specification
  - This returns an “FSM” interface with “start” and “done” methods
- Operate the FSM using the “start” and “done” methods

```
while (some cond) ... do state S1 actions ...
```

→  
can also be  
written as

```
while (True) seq
  ... do state S1 actions ...
  if (some cond) break;
endseq
```

# Composing FSMs

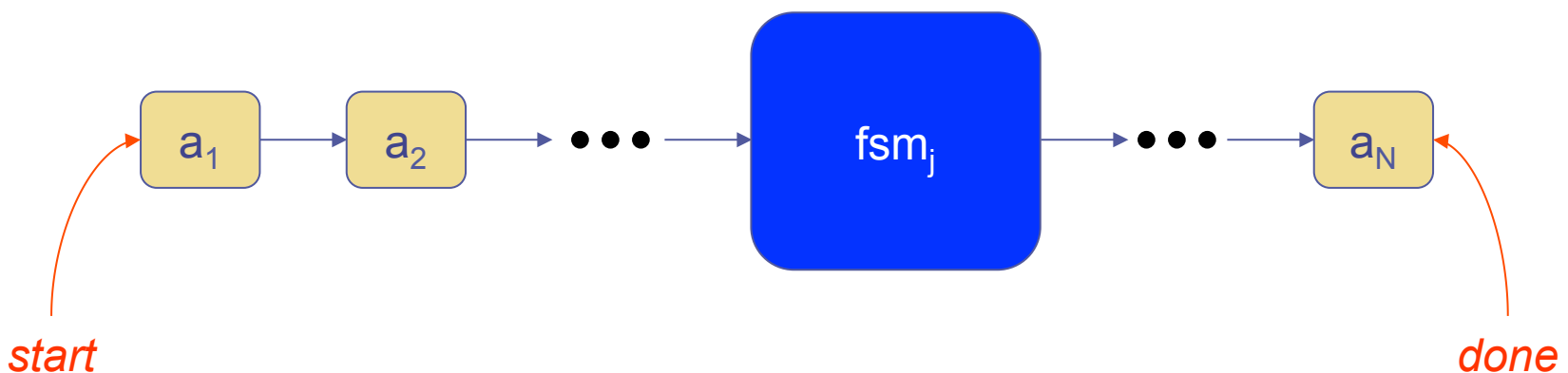
FSMs can be composed—individual FSMs can be combined systematically into larger FSMs.

The compositional principle is:

- Every FSM has a *start* and *done* state, embodied in the FSM interface methods
- When composing a larger FSM, the *start* and *done* for the larger FSM is systematically derived from the *start* and *done* of the component FSMs

Example: Linear sequencing

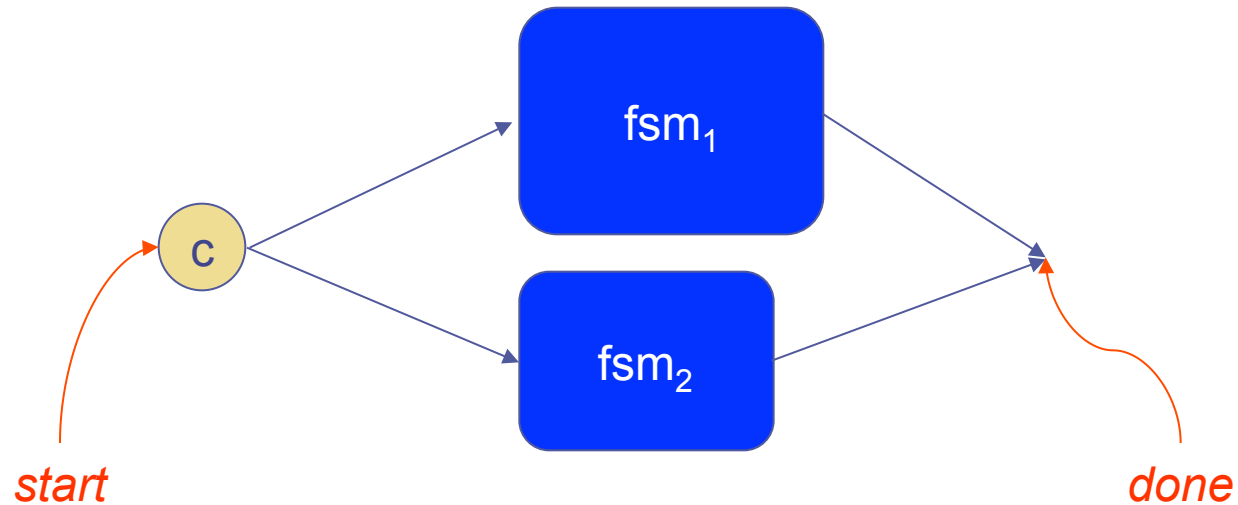
- Syntax: **seq** a<sub>1</sub>; a<sub>2</sub>; ... ; a<sub>N</sub>; **endseq**
- Where each a<sub>j</sub> is either an expression of type Action, or itself a sub-fsm (expression of type Stmt)



# Conditionals

## Syntax

- **if** (Bool expr) fsm1
- **if** (Bool expr) fsm1 **else** fsm2

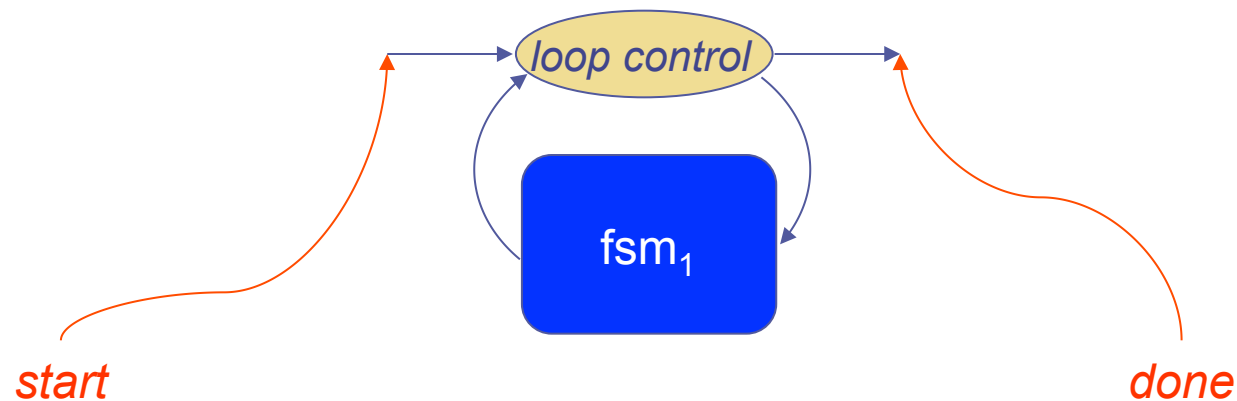


# Iteration (loops)

## Syntax

- **for** (loop control) fsm1
- **while** (Bool expr) fsm1
- **repeat** (Integer expr) fsm1

Loop bodies can contain **break** and **continue** keywords, with the usual meaning





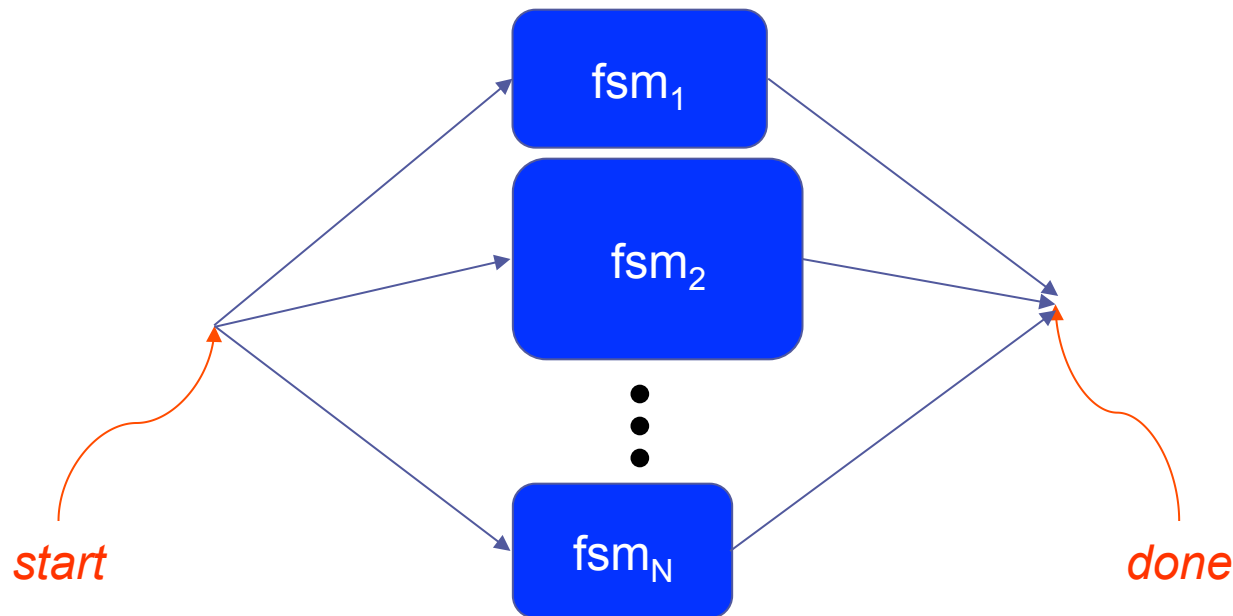
# Parallel composition (fork-join)

## Syntax

- **par** fsm1; fsm2; ...; fsmN; **endpar**

The sub-FSMs start at the same time and proceed independently.  
The FSM is done when the all sub-FSMs are done (they need not complete simultaneously; we wait for the last one to be done).

Note: the sub-FSM can contain Actions that conflict with each other—these are resolved automatically using standard rule scheduling.



# Another FSM example

```
Stmt specfsm =  
  seq  
    write( 15, 51 );  
    read( 15 );  
    ack ;  
    ack ;  
    write( 16, 61 );  
    write( 17, 71 );  
    // a memory operation and an  
    // acknowledge can occur  
    // simultaneously  
    action  
      read( 16 );  
      ack ;  
    endaction  
    action  
      read( 17 );  
      ack ;  
    endaction  
    ack ;  
    ack ;  
  endseq ;  
  
FSM testfsm <- mkFSM (specfsm);
```

- action/endaction blocks compose larger entities of type Action. Since an Action is always within a rule, it guarantees that the sub-actions will be simultaneous (atomic).

```
rule run ( True );  
  testfsm.start ;  
endrule  
  
rule done (testfsm.done);  
  $finish(0);  
endrule
```

# FSMs are often used as testbench stimulus generators

Example:

```
// Specify an FSM generating a test sequence
Stmt test_seq =
  seq
    for (i <= 0; i < NI; i <= i + 1)           // each source
      for (j <= 0; j < NJ; j <= j + 1) action   // each destination
        let pkt <- gen_packet ();
        send_packet (i, j, pkt);                // test i-j path in isolation
      endaction
    // then, test arbitration by sending packets simultaneously to same dest
    action
      send_packet (0, 1, pkt0);                 // to dest 1
      send_packet (1, 1, pkt1);                 // to dest 1 (so, collision)
    endaction
  endseq;

mkAutoFSM (test_seq);    // Generate the FSM and code to run it automatically
```

mkAutoFSM is another module provided in the library:

- It has an Empty interface
- Internally, it uses mkFSM to create the FSM, and it creates rules
  - to automatically start the FSM
  - to invoke \$finish when the FSM is done

# Revisiting our testbench from the EHRs lecture ...

```
module mkTest (Empty);
  UpDownSatCounter_ifc ctr <- mkUpDownSatCounter;
  Reg #(int) step <- mkReg (0);
  Reg #(Bool) flag0 <- mkReg (False); Reg #(Bool) flag1 <- mkReg (False);

  function Action count_show (Integer rulenum, Bool a_not_b, Int #(4) delta);
    action
      let x <- (a_not_b ? ctr.countA (delta) : ctr.countB (delta));
      $display ("cycle %0d, r%0d: is %0d, count (%0d)", cur_cycle, rulenum, x, delta);
    endaction
  endfunction

  // Rules 0-9 are sequential, just testing one method at a time
  rule r0 (step == 0); count_show (0, True, 3); step <= 1; endrule
  rule r1 (step == 1); count_show (1, True, 3); step <= 2; endrule
  ... and similarly, sequentially feed deltas of 3,3, -6,-6,-6,-6, 7, 3,
  // Concurrent execution
  rule r10 (step == 10 && !flag0); count_show (10,True, 6); flag0 <= True; endrule
  rule r11 (step == 10 && !flag1); count_show (11,False, -3); flag1 <= True; endrule

  // Show final value
  rule r12 (step == 10 && flag0 && flag1); count_show (12,True, 0); $finish; endrule
endmodule: mkTest
```

*These parts just constitute a structured FSM.  
On the next slide we use StmtFSM instead of explicit rules*

# Revisiting our testbench from the EHRs lecture ...

```
module mkTest (Empty);
  UpDownSatCounter_Ifc ctr <- mkUpDownSatCounter;

  function Action count_show (Integer rulenum, Bool a_not_b, Int #(4) delta);
    action
      let x <- (a_not_b ? ctr.countA (delta) : ctr.countB (delta));
      $display ("cycle %0d, r%0d: is %0d, count (%0d)", cur_cycle, rulenum, x, delta);
    endaction
  endfunction

  mkAutoFSM (
    seq
      count_show (0, True, 3);
      count_show (1, True, 3);
      ... and similarly, sequentially feed deltas of 3,3, -6,-6,-6,-6, 7, 3,
    par
      count_show (10,True, 6);
      count_show (11,False, -3);
    endpar
    count_show (12,True, 0);
  endseq);
endmodule: mkTest
```

*mkAutoFSM is just a BSV library function that*

- *takes a Stmt argument (here, “seq...endseq”),*
- *creates an FSM from the Stmt,*
- *runs it once,*
- *and calls \$finish*



```
module mkAutoFSM #(Stmt s) (Empty);
  FSM fsm <- mkFSM (s);
  rule rA:
    fsm.start;
  endrule
  rule rB (fsm.done);
    $finish;
  endrule
endmodule: mkAutoFSM
```

# Suspendable FSMs

The library provides another FSM constructor:

```
module mkFSMWithPred #(Stmt s, Bool b) (FSM);
```

An external agent can “asynchronously” start/stop the FSM by controlling the boolean predicate b.

*[ Language gurus: With parallel composition, nesting, abort and suspend, you get similar expressive power to FSMs in the language Esterel. The Esterel literature characterizes this power—it allows FSM descriptions to be exponentially smaller than descriptions without these capabilities. ]*

# FSM “servers”

Another useful composition facility is the FSM Server.

Normally, to perform a multi-cycle request to a server, your FSM would have to express it in a “split-phase” fashion:

```
Stmt s = seq
    mem.request.put (Req {op: Load, addr: a});
    let response <- mem.response.get;
endseq
```

With FSM servers, you can express this in a more traditional “procedure call” style.

```
function RStmt #(Data) fn_memServer (Req req);
    seq
        ... do the work for reading mem data ...
        return data;      // RStmt is a generalization of Stmt with ‘return’
                           values
    endseq
endfunction

FSMserver #(Addr, Data) memServer <- mkFSMServer (fn_memServer);

Stmt s = seq
    response <- callServer (memServer, Req {op: Load, addr:a });
endseq
```

## More info on FSMs

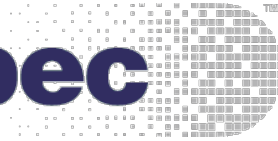
Section C.6.1 in the Reference Guide goes into a lot more detail on FSMs, including describing further functions, modules, etc. and providing many more examples.



# Hands-on

- BSV-by-Example book: Examples in Chapter 14

# bluespec



## End

```

import P2P000;
typedef Bit[24] StateT;
module ex_hdl_csrR_hdl {
  Integer nfa_depth;
  function Bit[24] determine_pump(StateT s);
  return (s[0]);
  endfunction

  P2P000(bstate) inbounds;
  address P2P000(nfa_depth) the_inbounds(inbounds);
  P2P000(bstate) outbounds;
  address P2P000(nfa_depth) the_outbounds(outbounds);
  P2P000(bstate) self_inbounds;
  address P2P000(nfa_depth) the_self_inbounds(self_inbounds);

  rule exp1 (True) {
    bstate in_data = inbounds.first;
    P2P000(bstate) out_data =
      determine_pump(in_data) == 0 ? outbounds : self_inbounds;
    self_inbounds.in_data =
      the_outbounds;
    out_data;
  }
  endrule; exp1;
endmodule; ex_hdl_csrR_hdl

```

Questions?

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or send an e-mail to [support@bluespec.com](mailto:support@bluespec.com)

