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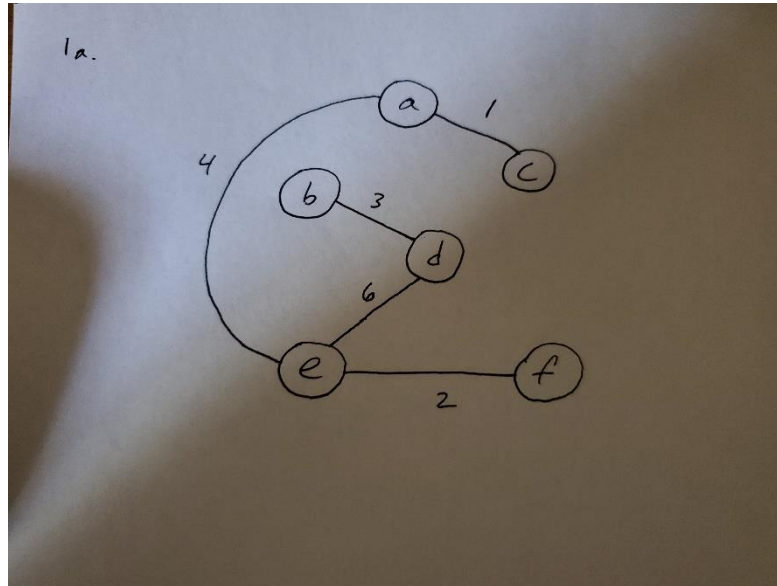
4/21/2021

Com S 311 Exam 2

### Exam 2

1.

a.



b.

	1	2	3	4	5
Kruskal's	(a, c)	(e, f)	(b, d)	(a, e)	(d, e)
Prim's	(a, c)	(a, e)	(e, f)	(d, e)	(b, d)

c. I could show that G has a unique minimum spanning tree by calculating G's minimum cost.

$$\text{Minimum cost} = 1 + 2 + 3 + 4 + 6 = 16$$

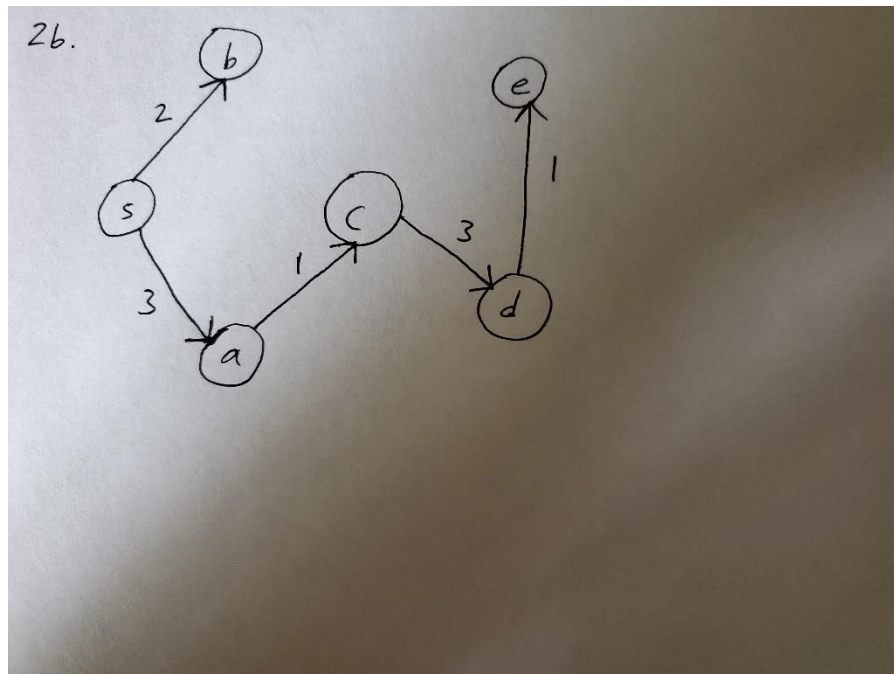
2.

a.

iter.	d[]						Selected Nodes
	s	a	b	c	d	e	
0	0	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	s
1	0	3	2	6	$\infty$	$\infty$	sb
2	0	3	2	5	$\infty$	11	sba

3	0	3	2	4	10	11	sbac
4	0	3	2	4	7	11	sbacd
5	0	3	2	4	7	8	sbacde
6	—	—	—	—	—	—	sbacde

b.



c. I do not know how to solve this question.

3. To prove that  $L$  is in  $P$  we have to build an algorithm that can make decisions in polynomial time that accepts all string in  $L$  but rejects all strings not in  $L$ . This vertex algorithm will work where a user inputs: a binary string  $X_n$  where  $n$  is the length.

Pseudo code:

- i)  $y \leftarrow \text{Reverse } x$   $O(n)$
- ii)  $z \leftarrow \text{compliment each bit of } y$   $O(n)$
- iii) if  $y$  accepts  $z$  in polynomial time, then the algorithm/program outputs “YES”
- iv) else  $y$  rejects  $z$  in polynomial time, output “NO”

Correctness Check:

Repeated complementation and reversal of a string will return the string in  $L$ . Therefore if  $y$  accepts  $z$  or rejects  $z$  each of the steps will run in polynomial time. As such,  $L$  is in  $P$ .

4. This algorithm was written in C!

```
// C program to count string num

#include <stdio.h>
#include <string.h>
#include <stdlib.h>

//k is the total num of characters

int strCount(int k, int b, int c){
//base case

if(b < 0 || c < 0){
return 0;
}

if(k == 0) {
return 1;
}

if(b == 0 && c == 0){
return 1;
}

//in these three cases we will go with the first, second, or third choice, k decreases by 1

int alg = strCount(n - 1, b, c);
alg += strCount(n - 1, b - 1, c - 1);
alg += strCount(n - 1, b, c);

return alg;
}
```

5. I do not know how to solve this problem.